

# Final Exam

● Graded

Student

Sihat Afnan

Total Points

96 / 100 pts

Question 1

Q1

4 / 4 pts

- 0 pts Correct

- 4 pts wrong

Question 2

Q2

4 / 4 pts

- 0 pts Correct

- 4 pts wrong

Question 3

Q3

4 / 4 pts

- 0 pts Correct

Question 4

Q4

4 / 4 pts

- 0 pts Correct

- 4 pts wrong

Question 5

Q5

4 / 4 pts

- 0 pts Correct

- 4 pts WRONG

Question 6

Q6

4 / 4 pts

- 0 pts Correct

- 4 pts WRONG

### Question 7

Q7

0 / 4 pts

– 0 pts Correct

✓ – 4 pts WRONG

### Question 8

Q8

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

### Question 9

Q9

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

### Question 10

Q10

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

### Question 11

Q11

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

### Question 12

Q12

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

### Question 13

Q13

4 / 4 pts

✓ – 0 pts Correct

– 4 pts Incorrect

#### Question 14

Q14

4 / 4 pts

✓ - 0 pts Correct

- 4 pts Incorrect

#### Question 15

Q15

11 / 11 pts

✓ - 0 pts Correct

- 11 pts wrong

- 5 pts adding a delay at the start of the video playtime at the client side

#### Question 16

Q16

11 / 11 pts

✓ - 0 pts Correct

- 3 pts slow start- incomplete

- 3 pts congestion avoidance - incomplete

- 3 pts time out,  $w = 1$ , slow start threshold =  $w/2$

- 11 pts wrong/missing

#### Question 17

Q17

11 / 11 pts

✓ - 0 pts Correct

- 3 pts Partially Correct

- 5 pts Partially Correct

- 10 pts Incorrect

#### Question 18

Q18

11 / 11 pts

✓ - 0 pts Correct

- 3 pts Difference between unicast and broadcast

- 3 pts Dangers of broadcast

- 3 pts Mechanisms to avoid dangers

- 11 pts Incorrect

# CS232 Final - 2025

Name	Sihat Afnan	UCInetID - First part of your UCI email before @
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- Use of books/notes is forbidden.
- Write the answers using only the predefined space using a black pen
- Put your answer into the box below for Q1-Q14
- **Duration:** 2 hrs.

Q1	Q2	Q3	Q4	Q5	Q6	Q7
a	b	c	c	c	c	c

  

Q8	Q9	Q10	Q11	Q12	Q13	Q14
a	a	b	c	b	b	b

**Q1** Consider a  $M/M/1/K$  queueing system with arrival and service rate equal to  $\lambda$  and  $\mu$  pkt/s, respectively. If we decrease  $K$ , the probability that a packet is dropped in going to

- a) Increase.
- b) Decrease.
- c) Stay the same.

**Q2** Consider a  $M/M/1$  queueing system with arrival and service rate equal to  $\lambda$  and  $\mu$  pkt/s, respectively. Define  $\rho = \lambda/\mu$ , the expected total time spent by packets in the system ( $E[T]$ ) is equal to

- a)  $1/\mu$ .
- b)  $\frac{1/\mu}{(1-\rho)}$ .
- c)  $\frac{\rho}{(1-\rho)}$ .

**Q3** The application layer protocol used by web browsers (such as Chrome and Firefox) is:

- a) SMTP
- b) FTP
- c) HTTP
- d) IMAP

**Q4** TCP is implemented at:

- a) Switches
- b) Routers
- c) host devices
- d) all the above

**Q5** Little's Formula (no blocking) is

- a)  $E[N] = \sum_{n=0}^N p^n n.$
- b)  $P[N(t) = 0] = 1 - \rho.$
- c)  $E[N] = \lambda E[T].$

**Q6** IMAP (Internet Mail Access Protocol) is an e-mail protocol used between:

- a) the sender and the sender's mail server.
- b) the sender's mail server and the receiver's mail server.
- c) the recipient's mail server and the recipient.

**Q7** Consider a setting where only one node has traffic. Compared to TDMA, polling-based access control achieves

- a) a larger data rate.
- b) a smaller data rate.
- c) the same data rate.

**Q8** in M/M/1/ $\infty$  the probability of having an empty buffer is

- a)  $p_0 = 1 - \rho.$
- b)  $p_0 = \rho.$
- c)  $p_0 = 1/(1 - \rho).$

**Q9** in M/M/c/c the probability that a connection is rejected is

- a)  $\frac{\rho^c / c!}{\sum_{k=0}^c \rho^k / k!}$
- b)  $\sum_{k=0}^c \rho^k / k!$
- c)  $\rho^c.$

**Q10** A socket is implemented between:

- a) transport layer and network layer.
- b) application layer and transport layer.
- c) network layer and link layer.

**Q11** Mark the (one) correct statement

- a) UDP implements packet retransmissions and packet reordering
- b) UDP implements packet retransmissions, but does not implement packet reordering
- c) UDP does not implement packet retransmissions and packet reordering

**Q12** In TCP, if we increase the timeout time with respect to the RTT

- a) The number of retransmitted packets will increase
- b) The number of retransmitted packets will decrease
- c) The number of retransmitted packets will be the same

**Q13** In Distance Vector-based routing algorithms

- a) the routers send broadcast messages to all the routers in the autonomous system
- b) the routers send broadcast messages to all their directly connected routers
- c) the routers do not need to exchange messages.

**Q14** In broadcast routing based on reverse path forwarding, the router forwards the packet

- a) in any case
- b) if the packet comes from their shortest path link to the source
- c) If the packet comes from a link part of the spanning tree

**Q15** Briefly motivate and describe the delayed playout mechanism in multimedia streaming.

The delayed playout mechanism employs a initial delay to each packets before streaming it.

Motivation: This is done so that we can avoid delay or buffering in streaming. The application starts with enough data to be needed in future.

Mechanism: Each packet is marked by its timestamp. A delay time  $q$  is chosen. For packet  $i$ , if its  $t_i + q$  is the delay duration. If the packet arrives before  $t_i + q$ , it's accepted; otherwise discarded. For fixed mode, the value  $q$  is constant. For adaptive mode, the  $q$  value is computed based on network status.

**Q16** Describe the update rules of the transmission window  $W$  in slow start and congestion control states of TCP (Tahoe).

Let  $cwnd$  be the congestion window in TCP Tahoe.

Slow Start State:

initially,  $cwnd = 1$

on each ACK,  $cwnd = cwnd + 1$

So, in 1 RTT,  $cwnd$  gets doubled  
Upon timeout / 3 dup ACK, it makes  $cwnd = 1$  &  $ssthresh = \frac{cwnd}{2}$

Congestion Avoidance State:

enters when  $cwnd \geq ssthreshold$

on each ACK,  $cwnd = cwnd + mss \left( \frac{mss}{cwnd} \right)$

So, in 1 RTT,  $cwnd$  increases by 1

segment size ( $mss = \text{Max. segment size}$ )

**Q17** Describe the leaky bucket mechanism and discuss how the bucket parameters influence the characteristics of the output traffic.

The leaky bucket mechanism is a traffic policing algorithm used to regulate the rate at which packets enter a network. It consists of bucket capacity b and leak rate r.

The parameter  $r$  controls the long term average transmission rate while  $b$  controls the maximum burst size allowed. A larger  $b$  allows larger bursts without packet loss and a larger  $r$  allows higher sustained bandwidth. The output traffic is therefore smoothed and constrained and congestion is reduced.

**Q18** Describe the difference between unicast and broadcast. What are the dangers of broadcast and what mechanisms are implemented to avoid them?

Unicast: One to One communication where sender transmits data to a specific receiver. It's bandwidth efficient and scalable.

Broadcast: In broadcast, one sender transmits data to all hosts in the network. This is not bandwidth efficient neither scalable.

Dangers of Broadcast:

- 1) Broadcast Storm: When too many broadcast packets generated, there's CPU overload or network shutdown.
- 2) Bandwidth Wastage: Packets delivered to all hosts who don't even need them.
- 3) Network Congestion: Queue overflow or Packet drops; increased delay.
- 4) Security Risks: ARP spoofing, MAC flooding, Smurf attacks etc.