

# Assignment 4

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## 1 Question 1

The system is an M/M/1/K queue with limited capacity  $K = 3$ .

- Arrival Rate:  $\lambda = 4 \text{ pkt/s}$
- Service Rate:  $\mu = 12 \text{ pkt/s}$
- System Capacity (Queue + Server):  $K = 3$

$$\rho = \frac{\lambda}{\mu} = \frac{4}{12} = \frac{1}{3}$$

### 1.1 1a

For an M/M/1/ $K$  queue where  $\rho \neq 1$ , the probability of zero packets is:

$$P_0 = \frac{1 - \rho}{1 - \rho^{K+1}}$$

Substituting  $\rho = 1/3$  and  $K = 3$ :

$$P_0 = \frac{1 - \frac{1}{3}}{1 - \left(\frac{1}{3}\right)^{3+1}} = \frac{2/3}{1 - 1/81} = \frac{2/3}{80/81} = \frac{2}{3} \cdot \frac{81}{80} = \frac{27}{40}$$

The probability of  $N$  packets is given by  $P_N = P_0\rho^N$ . The system has 1 packet being served and 1 packet waiting, which means the total number of packets in the system is  $N = 2$ . We compute  $P_2$ :

$$P_2 = P_0\rho^2 = \left(\frac{27}{40}\right) \left(\frac{1}{3}\right)^2$$

$$P_2 = \frac{27}{40} \cdot \frac{1}{9} = \frac{3}{40}$$

So, the probability that the system has 2 packets is  $3/40 \approx 0.075$ .

## 1.2 1b

We find out the expected time a packet stays in the system  $W$  (queueing + service). We proceed via Little's Law.

Mean number in system  $L$  : Using  $p_n = P\{N = n\} = P_0 \rho^n$  with  $\rho = \frac{1}{3}$  and  $K = 3$ ,

$$p_1 = \frac{27}{40} \cdot \frac{1}{3} = \frac{9}{40}, \quad p_2 = \frac{27}{40} \cdot \frac{1}{9} = \frac{3}{40}, \quad p_3 = \frac{27}{40} \cdot \frac{1}{27} = \frac{1}{40}.$$

Hence

$$L = \sum_{n=0}^3 n p_n = 0 \cdot p_0 + 1 \cdot p_1 + 2 \cdot p_2 + 3 \cdot p_3 = \frac{9 + 6 + 3}{40} = \frac{18}{40}.$$

Effective arrival rate  $\lambda_{\text{eff}}$ : Arrivals are blocked when the system is full, with blocking probability

$$P_{\text{block}} = p_3 = \frac{1}{40}.$$

Thus the admitted (throughput) rate is

$$\lambda_{\text{eff}} = \lambda (1 - P_{\text{block}}) = 4 \left(1 - \frac{1}{40}\right) = \frac{156}{40} = 3.9 \text{ pkt/s.}$$

Applying  $L = \lambda_{\text{eff}} W$ ,

$$W = \frac{L}{\lambda_{\text{eff}}} = \frac{\frac{18}{40}}{\frac{156}{40}} = \frac{18}{156} = \boxed{\frac{3}{26} \text{ s} \approx 0.11538 \text{ s}}.$$

## 2 Question 2

The system is a standard M/M/1 queue with infinite capacity.

- Arrival Rate:  $\lambda = 6 \text{ pkt/s}$
- Service Rate:  $\mu = 12 \text{ pkt/s}$
- Traffic Intensity:  $\rho = \frac{\lambda}{\mu} = \frac{6}{12} = \frac{1}{2}$

### 2.1 2a

The expected time a packet spends in the system ( $\mathbf{W}$ ) includes both queueing time and service time. The formula for the average time in the system for an M/M/1 queue is:

$$W = \frac{1}{\mu - \lambda}$$

Substituting the rates:

$$W = \frac{1}{12 - 6} = \frac{1}{6} \text{ seconds}$$

The expected time a packet spends in the system is **1/6** seconds.

## 2.2 2b

The expected time a packet waits in the buffer before getting served is the average queueing time ( $\mathbf{W}_q$ ). The formula for the average queueing time for an M/M/1 queue is:

$$W_q = \frac{\lambda}{\mu(\mu - \lambda)}$$

Substituting the rates:

$$W_q = \frac{6}{12(12 - 6)} = \frac{6}{12(6)} = \frac{6}{72} = \frac{1}{12} \text{ seconds}$$

The expected waiting time in the buffer is **1/12** seconds.

### 3 Question 3

Given Parameters

- Number of applicationsstreams:  $c = 5$
- Individual Arrival Rate:  $\lambda_i = 3 \text{ pkt/s}$
- Individual Service Rate:  $\mu_i = 6 \text{ pkt/s}$

Packet delay is defined as the average time a packet spends in the system ( $W$ ).

#### 3.1 3a

For a single dedicated server:

- $\lambda = 3 \text{ pkt/s}$
- $\mu = 6 \text{ pkt/s}$

The average time in the system ( $W_{\text{ded}}$ ) is:

$$W_{\text{ded}} = \frac{1}{\mu - \lambda} = \frac{1}{6 - 3} = \frac{1}{3} \text{ seconds}$$

For a single Powerful Server (1 Combined M/M/1 Queue), the total rates for the single combined server are:

- Total Arrival Rate:  $\lambda_{\text{total}} = 5 \cdot 3 = 15 \text{ pkt/s}$ . (The reason the total arrival rate must be  $5\lambda$  for the single powerful server is due to the principle of Superposition of Poisson Processes)
- Total Service Rate:  $\mu_{\text{total}} = 5 \cdot 6 = 30 \text{ pkt/s}$

The average time in the system ( $W_{\text{single}}$ ) is:

$$W_{\text{single}} = \frac{1}{\mu_{\text{total}} - \lambda_{\text{total}}} = \frac{1}{30 - 15} = \frac{1}{15} \text{ seconds}$$

Comparing the delays:

$$W_{\text{single}} = \frac{1}{15} \quad \text{and} \quad W_{\text{ded}} = \frac{1}{3} = \frac{5}{15}$$

Since  $W_{\text{single}} < W_{\text{ded}}$ , the single powerful server provides a lower packet delay. The single powerful server is better in terms of packet delay.

### 3.2 3b

This is based on the single powerful server from Case 2, which is an M/M/1 queue with  $\rho = 15/30 = 1/2$ . We want to compute  $P_N$  for  $N = 6$ . The steady-state probability of having  $N$  packets in an M/M/1 system is  $P_N = (1 - \rho)\rho^N$ .

$$P_6 = (1 - \rho)\rho^6$$

Substituting  $\rho = 1/2$ :

$$P_6 = \left(1 - \frac{1}{2}\right) \left(\frac{1}{2}\right)^6 = \frac{1}{2} \cdot \frac{1}{64} = \frac{1}{128}$$

The probability that there are 6 packets in the system is **1/128**.

## 4 Question 4

Given Parameters

- System Type: M/M/1/ $K$
- Fixed Capacity:  $K = 2$

The average number of customers in the system ( $L$ ) is  $L = \sum_{N=1}^K N \cdot P_N$ , where  $P_N = P_0 \rho^N$  and  $P_0 = \frac{1-\rho}{1-\rho^{K+1}}$ .

### 4.1 4a

In the startup Phase, the parameters are:

- $\lambda = 2$  pkt/s
- $\mu = 10$  pkt/s
- $\rho = 2/10 = 0.2$

Let's find the probabilities of having 0,1,2 customers in the system.

1.  $P_0 = \frac{1-0.2}{1-(0.2)^3} = \frac{0.8}{0.992} \approx 0.80645$
2.  $P_1 = P_0 \cdot (0.2) \approx 0.16129$
3.  $P_2 = P_0 \cdot (0.2)^2 \approx 0.03226$

So the average number of customers is

$$L = 1 \cdot P_1 + 2 \cdot P_2 \approx 0.16129 + 0.06452 \approx 0.2258$$

## 4.2 4b

In the Adolescent Phase, the parameters are:

- $\lambda = 8 \text{ pkt/s}$
- $\mu = 10 \text{ pkt/s}$
- $\rho = 8/10 = 0.8$

Let's find the probabilities of having 0, 1, 2 customers in the system.

$$1. P_0 = \frac{1-0.8}{1-(0.8)^3} = \frac{0.2}{0.488} \approx 0.40984$$

$$2. P_1 = P_0 \cdot (0.8) \approx 0.32787$$

$$3. P_2 = P_0 \cdot (0.8)^2 \approx 0.26229$$

So the average number of customers is

$$L = 1 \cdot P_1 + 2 \cdot P_2 \approx 0.32787 + 0.52458 \approx 0.8525$$

### 4.3 4c

In the Viral Spike Phase, the parameters are:

- $\lambda = 15 \text{ pkt/s}$
- $\mu = 10 \text{ pkt/s}$
- $\rho = 15/10 = 1.5$

Let's find the probabilities of having 0, 1, 2 customers in the system.

1.  $P_0 = \frac{1-1.5}{1-(1.5)^3} = \frac{-0.5}{-2.375} \approx 0.21053$
2.  $P_1 = P_0 \cdot (1.5) \approx 0.31579$
3.  $P_2 = P_0 \cdot (1.5)^2 \approx 0.47368$

So the average number of customers is

$$L = 1 \cdot P_1 + 2 \cdot P_2 \approx 0.31579 + 0.94736 \approx 1.2632$$

Note that since  $\rho > 1$ , the server is a major bottleneck, and the probability of the system being full ( $P_2$ ) is the highest.

#### 4.4 4d

We consider the viral-spike parameters  $\lambda = 15$  pkt/s,  $\mu = 10$  pkt/s so that  $\rho = \lambda/\mu = 1.5 > 1$ . If we double the buffer from  $K = 2$  to  $K = 4$  (M/M/1/ $K$ ), the steady-state empty probability is

$$p_0 = \frac{1 - \rho}{1 - \rho^{K+1}} = \frac{1 - 1.5}{1 - 1.5^5} = \frac{-0.5}{-6.59375} \approx 0.07586.$$

Hence the blocking probability is

$$p_K = p_4 = p_0 \rho^4 = 0.07586 \times 1.5^4 = 0.07586 \times 5.0625 \approx 0.384.$$

The admitted (throughput) rate is therefore

$$\lambda_{\text{eff}} = \lambda(1 - p_K) = 15(1 - 0.384) = 15 \times 0.616 = 9.24 \text{ pkt/s.}$$

(For comparison, with the original  $K = 2$ ,  $p_2 = 0.4737$  and  $\lambda_{\text{eff}} = 15(1 - 0.4737) \approx 7.89$  pkt/s.)

Increasing the buffer reduces blocking and raises throughput from  $\approx 7.89$  to  $9.24$  pkt/s, but it does not fix the overload: the offered load  $\lambda = 15$  still exceeds the service capacity  $\mu = 10$ , so the system remains congested and spends much of the time near full. To handle the spike one must increase service capacity (larger  $\mu$  or more servers).

**General Rule of Networking:** The response implies the following general rule:

**Buffer Space is Not a Substitute for Service Capacity (Bandwidth).**

A system that is unstable ( $\rho > 1$ ) must address the service bottleneck ( $\mu$ ), not the buffer size ( $K$ ). The fundamental condition for stability is  $\lambda < \mu$ .

## 5 Question 5

Given Parameters (M/M/c/c System)

- Number of Servers:  $c = 3$
- Arrival Rate:  $\lambda = 2 \text{ pkt/s}$
- Service Rate (per server):  $\mu = 4 \text{ pkt/s}$

The traffic intensity ( $\rho$ ) is  $\rho = \frac{\lambda}{c\mu} = \frac{2}{3 \cdot 4} = \frac{2}{12} = \frac{1}{6}$ . Since  $\rho < 1$ , the system is stable. The term  $\lambda/\mu$  is  $\alpha = \frac{2}{4} = \frac{1}{2}$ .

### 5.1 5a

We assume the system has reached steady-state. We first find  $P_0$ . The steady-state probability of zero packets is:

$$P_0 = \left[ \sum_{n=0}^{c-1} \frac{(\lambda/\mu)^n}{n!} + \frac{(\lambda/\mu)^c}{c!} \frac{1}{1-\rho} \right]^{-1}$$

Substituting  $\alpha = 1/2$ ,  $c = 3$ , and  $\rho = 1/6$ :

$$\begin{aligned} P_0 &= \left[ \frac{(1/2)^0}{0!} + \frac{(1/2)^1}{1!} + \frac{(1/2)^2}{2!} + \frac{(1/2)^3}{3!} \frac{1}{1-1/6} \right]^{-1} \\ P_0 &= \left[ 1 + \frac{1}{2} + \frac{1}{8} + \frac{1/8}{6} \cdot \frac{6}{5} \right]^{-1} \\ P_0 &= \left[ 1 + \frac{1}{2} + \frac{1}{8} + \frac{1}{40} \right]^{-1} = \left[ \frac{40 + 20 + 5 + 1}{40} \right]^{-1} = \left[ \frac{66}{40} \right]^{-1} = \frac{40}{66} = \frac{20}{33} \end{aligned}$$

Since  $N = 3$  is less than or equal to  $c = 3$ , we use the formula:

$$\begin{aligned} P_3 &= \frac{(\lambda/\mu)^3}{3!} P_0 \\ P_3 &= \frac{(1/2)^3}{6} \cdot \frac{20}{33} = \frac{1/8}{6} \cdot \frac{20}{33} = \frac{1}{48} \cdot \frac{20}{33} = \frac{5}{12 \cdot 33} = \frac{5}{396} \end{aligned}$$

So the probability is **5/396**  $\approx 0.0126$ .

## 5.2 5b

The probability  $P_3$  represents the steady-state probability that there are exactly 3 customers in the system. Since the number of customers ( $N = 3$ ) equals the number of servers ( $c = 3$ ), this probability specifically represents the state where All 3 Servers are Busy, and the Queue is Empty. The system is operating at maximum utilization without any customer needing to wait in the buffer.