

# Projecting chains of constrained points

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## 1 Formulation

We're going to solve

$$\begin{aligned} \min_{\mathbf{x}} \quad & \frac{1}{2} (\mathbf{x} - \mathbf{p})^\top \mathbf{M} (\mathbf{x} - \mathbf{p}) \\ \text{s.t.} \quad & \mathbf{z} = \mathbf{R}\mathbf{x} \quad \text{and} \quad \mathbf{z}_i^\top \mathbf{z}_i = 1 \end{aligned}$$

where

$\mathbf{x} \in \mathbb{R}^{3n}$  is the solution vector,

$\mathbf{p} \in \mathbb{R}^{3n}$  is the vector of target/free fall positions,

$\mathbf{z} \in \mathbb{R}^{3(n-1)}$  is the vector of difference vectors in the chain,

$\mathbf{z}_i \in \mathbb{R}^3$  is the  $i$ th 3-vector in  $\mathbf{z}$ ,

$\mathbf{R} \in \mathbb{R}^{3(n-1) \times 3n}$  is the per vector finite-difference matrix  
so  $(\mathbf{R}\mathbf{x})_i = \mathbf{x}_i - \mathbf{x}_{i+1}$

$\mathbf{M} \in \mathbb{R}^{3n \times 3n}$  is the mass matrix, simply having  
the mass (or weight) of each particle  
in each 3x3 block diagonal

We want to be working only on the difference vectors in  $\mathbf{z}$ . To do that, we skip the details and jump directly to the minimizer for  $\mathbf{x}$ , treating  $\mathbf{z}$  as a constant for now:

$$\mathbf{x} = \mathbf{p} - \mathbf{M}^{-1}\mathbf{R}^T\mathbf{S}^{-1}(\mathbf{R}\mathbf{p} - \mathbf{z})$$

where  $\mathbf{S} = \mathbf{R}\mathbf{M}^{-1}\mathbf{R}^T$  is a block tridiagonal matrix. Inserting this into the main problem eliminates  $\mathbf{z} = \mathbf{R}\mathbf{x}$  and we are left with

$$\begin{aligned} \min_{\mathbf{x}} \quad & \frac{1}{2} (\mathbf{z} - \mathbf{R}\mathbf{p})^T \mathbf{S}^{-1} (\mathbf{z} - \mathbf{R}\mathbf{p}) \\ \text{s.t.} \quad & \mathbf{z}_i^T \mathbf{z}_i = 1 \quad i = 1 \dots n \end{aligned}$$

In this reformulated problem we must find a  $\mathbf{z}$  that consist of unit length 3-vectors that minimize the distance to  $\mathbf{R}\mathbf{p}$  under the  $\mathbf{S}^{-1}$  norm. It worth noting that it is trivial to project  $\mathbf{z}$  to satisfy the unit length constraints, and likewise it is trivial to obtain  $\mathbf{x}$  given any  $\mathbf{z}$ . This means that whenever we take a newton step on  $\mathbf{z}$  we can trivially reign it in to its feasible solution, and in effect only changing the directions of each  $\mathbf{z}_i$  vector, never their length.

## 2 Algorithm

Here we'll outline the algorithm in full before explaining each step in detail:

```

L = cholesky(S)                                1
z = Rx                                          2
while true do                                  3
    Q,z = normalize(z)                          4
    y = Rp - z                                  5

```

$\lambda = \text{solve}(\mathbf{Q}\mathbf{S}\mathbf{Q}^T, \mathbf{Q}\mathbf{y})$	6
$\mathbf{b}_z = \text{solve}(\mathbf{L}, \mathbf{y})$	7
$\mathbf{b}_\lambda = \mathbf{L}^T \mathbf{Q}^T \lambda$	8
if $\ \mathbf{b}_z - \mathbf{b}_\lambda\ ^2 < \epsilon$ then	9
$\mathbf{s} = \text{solve}(\mathbf{L}^T, \mathbf{b}_z)$	10
$\mathbf{x} = \mathbf{p} - \mathbf{M}^{-1} \mathbf{R}^T \mathbf{s}$	11
return $\mathbf{x}$	12
$\mathbf{D} = \text{diagonal}(\max(0, \lambda))$	13
$\Delta \mathbf{z} = \mathbf{L} \text{ solve}(\mathbf{I} + \mathbf{L}^T \mathbf{D} \mathbf{L}, \mathbf{b}_z - \mathbf{b}_\lambda)$	14
$\mathbf{z} = \mathbf{z} + \Delta \mathbf{z}$	15

Each iteration the algorithm does these main steps:

1. Normalize the 3-vectors in  $\mathbf{z}$
2. Compute the  $n - 1$  Lagrange multipliers of the system at  $\mathbf{z}$
3. Check the solution residual, if below threshold compute and return  $\mathbf{x}$
4. Compute the Newton step on  $\mathbf{z}$  and go to step 1

Note the following:

- we compute the cholesky decomposition  $\mathbf{L}\mathbf{L}^T = \mathbf{S}$  to keep the systems we need to solve sparse and symmetric. See more details below. Computing cholesky for a block-tridiagonal SPD matrix is simple and fast, and if  $\mathbf{M}$  is fixed it can be done off-line/once.
- every `solve()` is at least a block tri-diagonal system (or simpler), which are fast and easy to solve
- to compute the Lagrange multipliers we build the Jacobian matrix  $\mathbf{Q} \in \mathbb{R}^{3(n-1) \times (n-1)}$ , which is simply the  $\mathbf{z}_i$  vectors vertically stacked

- the matrix  $\mathbf{D}$  is the contribution to the Hessian stemming from the non-linear unit length constraints. We clamp it to be positive to keep the Hessian PD. This is a hack but seems to play out well in practice

## 2.1 Lagrange multipliers

## 2.2 Hessian system