## OS 2020 project

## - Project 1 -

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Kernel Version:

Added syscall 334: sys\_my\_printk() prints a string to dmesg

Linux 4.14.25 -> https://cdn.kernel.org/pub/linux/kernel/v4.x/linux-4.14.25.tar.xz

Design

• Read the input (number of process, scheduling policy, process name, ready time, execution

time)

main.c

• Use scheduler() funtion to implement sheduling. process.c

Define the followin function: • TIME\_UNIT(): define a basic unit of execution time.

• assign\_proc\_core(pid, core): decide which core the process "pid" will run on.

• proc\_out(pid): use sched\_setscheduler() to reduce the priority "pid" of process and then

assign the child process back to the core 0 where the parent process is running.

• proc\_wakeup(pid): assign the child process to the core where the core 1 and then use

sched\_setscheduler() to raise the priority "pid" of process.

• proc\_exec(Process):

(1) Scheduling process call fork() to simulate the process which is ready and stops child process by reduecing its priority until parent process wakes it up. Because parent and child

process are in the same core. So due to the priority, child process won't run if it shouldn't run. But in case the child process run in the unavailable time and unfortunately start the timer, we set a while() loop to avoid the above problem and it will break the loop when child process priority is raised by parent process.

(2) When the timer starts, child process will enter a while() loop for execution time of TIME\_UNIT().

(3) When the timer ends, use system call to output the message into dmesg. scheduler.c

• First, we will assign a particular core 0 to scheduling process and raise its priority to the highest level to prevent potential preemptive problem between scheduling process and the child processes which are are generated by fork().

Second, initializes child process by -1 to represent not ready process or already finished

process.

There are four policies:

FIFO():

SJF():

Step4:

**Unit time** 

FIFO\_1.txt

Task

FIFO\_2.txt

FIFO\_3.txt

Task

P7

P2

Р3

P4

P2

P1

P4

P2

P4

P1

PSJF\_5.txt

Task

P1

P3

P2

P4

P5

P1

P2

Р3

P4

P5

RR\_2.txt

Task

P1

P2

RR\_3.txt

P1

RR\_4.txt

Task

P4

P5

P6

Р3

P7

P2

P1

P7

P2

P1

SJF\_1.txt

Task

P2

P3

P4

P1

SJF\_2.txt

Task

P6

P7

P2

P3

P8

P2

P3

P4

Task

P0

P1

P2

P3

P4

P5

P6

P7

P8

P9

500

500

500

TIME\_MEASUREMENT.txt

500

500

500

500

Conclusion

Correctness

than 10%

Work Loading

4000

4000

5000

7000

9000

500

500

7500

8500

**Expected time** 

**Expected time** 

4000

4000

4000

13500

15000

19500

23000

15000

19500

23000

2000

1000

4000

7000

**Expected time** 

**Expected time** 

3000

4000

7000

**Expected time** 

1000

4000

2000

5000

3000

1000

FIFO\_4.txt

4000

0.001377

**Expected time** 

Assume there are some ready processes.

shortest execution time process in the ready queue.

• In while(1) loop, we will kepp doing tje following five steps until all processes are done. Step1:

Check whether there are some process are already done in last UNIT\_TIME.If so,label

process's pid into -1 and finished processes number += 1.If finished processes number is equal to total process number, break the while (1) loop and finish scheduling. Step2:

Step3: Use switch to choose the scheduling policy to find the next process to implement it.

Check whether there are some processes which are ready and if so, implement proc\_exec().

When the implementing process i finish, the next implement process will be the shortest execution time process in the ready queue. RR():In each time slice 500 UNIT\_TIME or the implementing process finish in time slice, the next implementing process will be the next in the ready queue. PSJF():

NO matter the implementing process i is finished, the next implement process will be the

When the implementing process i finish, the next implement process will be i+1.

If the next running process isn't the same as now running process, then scheduling process will reduce the priority of now runnig process and raise the priority of next running process. Step5: run a TIME\_UNIT in parent and now runnig child process simutaneously. Result

**Error rate** 

**Error rate** 

4.2225%

3.2686%

5.1997%

3.0721%

2.2419%

6.1839%

4.8034%

0.2939%

2.3451%

2.1377%

**Error rate** 

3.9726%

4.0296%

**Error rate** 

0.4008%

3.5670%

**Error rate** 

7.0323%

4.7821%

5.1519%

5.1509%

4.1140%

3.3350%

6.2028%

3.7472%

3.8594%

3.0402%

**Error rate** 

3.7596%

5.6285%

3.8875%

4.0647%

**Error rate** 

1.6261%

0.8665%

4.1976%

4.4535%

4.7925%

500 P1 489.9160 2.0168% P2 500 487.8090 2.4382%

**Execution time** 

Р3 500 482.1090 3.5782% P4 500 5.7884% 471.0580 P5 500 3.0100% 484.9500

**Expected time** Task **Execution time Error rate** P1 80000 81020.7230 1.2759% P2 5000 5120.1950 2.4039% P3 1000 1064.8440 6.4844% 1050.3880 1000 P4 5.0388%

**Execution time Expected time** P1 8000 8272.3870 3.4048% 5000 5038.0870 P2 0.7617% 4.5003% 3000 3135.0080 Р3 P4 1000 1002.2240 0.2224% P5 1000 959.5250 4.0475% 3.7264% 962.7360 1000 P6

3831.1020

**Execution time Expected time Error rate** Task 2000 1908.3270 4.5836% P1 P2 500 485.3650 2.9270% P3 200 186.5370 6.7315% P4 500 490.9060 1.8188% FIFO\_5.txt **Expected time Execution time** Task **Error rate** P1 7615.4980 8000 4.8063%

4836.5720

2844.0080

969.2790

P5 1000 965.2920 3.4708% P6 1000 972.8310 2.7169% P7 4000 3875.7150 3.1071% PSJF\_1.txt **Expected time Execution time** Task **Error rate** P4 3000 2819.2620 6.0246% P3 8000 7580.1310 5.2484% 14535.3480 3.0977% P2 15000 P1 25000 24435.5740 2.2577% PSJF\_2.txt **Expected time Execution time Task Error rate** 

1000 5.1960% P5 948.0400 7000 P3 0.1954% 7013.6780 PSJF\_3.txt **Execution time Expected time Error rate** Task P2 500 531.5470 6.3094% 513.2470 Р3 500 2.6494% 500 P4 500.7710 0.1542% P1 3500 3587.3940 2.4970% PSJF\_4.txt **Expected time Execution time** Task **Error rate** Р3 1000 992.0110 0.7989%

3008.8180

4093.8050

7149.6360

**Execution time** 

1022.4190

4247.3570

1903.9320

100 97.8610 2.1390% 200 190.5140 4.7430% 4000 4163.1560 4.0789% 4000 9.2737% 4370.9480 7000 7082.2750 1.1754% RR\_1.txt **Execution time** Task **Expected time Error rate** 500 476.9680 4.6064% 500 486.6720 2.6656% 500 484.8670 3.0266%

480.1370

479.8520

**Execution time** 

7469.9370

8196.8040

**Execution time Task Expected time Error rate** P3 14000 13126.0550 6.2425% 19000 18079.2580 4.8460% P2 18000 5.6129% 16989.6840 21000 P6 19965.1630 4.9278% P5 23500 22642.8740 3.6473% 25000 23560.3680 P4 5.7585%

**Execution time** 

3718.7060

3808.7160

3793.9230

12804.6280

14382.8930

18849.6690

21573.3470

14437.9160

18747.4210

22300.7630

**Execution time** 

1924.8090

943.7150

3844.5010

6715.4740

**Execution time** 

RR\_5.txt **Execution time Expected time Error rate** Task P4 4000 3842.8120 3.9297% P5 4000 3813.7980 4.6551% 3829.5170 P6 4000 4.2621% Р3 13500 13016.0900 3.5845%

P1	100	93.8440	6.1560%
Р3	200	186.3710	6.8145%
P2	4000	3794.1680	5.1458%
P4	4000	3766.8300	5.8293%
P5	7000	6937.6590	0.8906%
SJF_3.txt	t		
SJF_3.txt	Expected time	Execution time	Error rate
_		Execution time 2922.5990	Error rate 2.5800%
Task	Expected time		

3934.9560

3965.3410

4790.1220

6688.2530

8568.6760

453.0360

489.2080

457.9260

483.5790

517.1410

491.5000

495.2670

SJF_4.txt	į		
Task	Expected time	<b>Execution time</b>	Error rate
P1	3000	2956.2080	1.4597%
P2	1000	993.3720	0.6628%
Р3	4000	3733.0290	6.6743%
P5	1000	980.1890	1.9811%
P4	2000	1865.4100	6.7295%
SJF_5.txt	-		
Task	<b>Expected time</b>	<b>Execution time</b>	Error rate
P1	2000	1951.3840	2.4308%

<b>Expected time</b>	<b>Execution time</b>	Error rate
500	502.1140	0.4228%
500	520.3880	4.0776%
500	511.5010	2.3002%
500	512.6380	2.5276%
500	478.3490	4.3302%
500	487.5220	2.4956%

3.2842%

3.4282%

1.7000%

0.9466%

The error rate between theoretical execution time and real execution time aren't bigger

When CPU was busy or has other jobs, the performance would be affected.

9.3928%

2.1584%

8.4148%