

17~18

$\varphi_k = \mathbb{N} \times \mathbb{N} \times \mathbb{N} \times \dots \times \mathbb{N} \mapsto \mathbb{N}$, and $k \in \mathbb{N}$, $k \geq 2$

$\varphi_k(n_1, n_2, n_3, n_4, \dots, n_k) = \max_k \{n_1, n_2, \dots, n_k\}$ 是原始递归的.

$$\varphi_k(n_1, n_2, \dots, n_k) = \begin{cases} \max_2 \{n_1, n_2\} & , k=2 \\ \max_2 [\max_{k-1} \{n_1, n_2, \dots, n_{k-1}\}, n_k] & , k \geq 3 \end{cases}$$

$\max_2 \{n_1, n_2\} = n_1 \cdot (n_1 \geq n_2) + n_2 \cdot (1 \sim (n_1 \geq n_2))$ and $n_1 \geq n_2$ 是原始递归函数, 并且原函数被递归定义, 则证毕.

11~12

$$\varphi(x, y) = \begin{cases} (x+1)^y, & x, y \text{ are composite number} \\ 0, & \text{otherwise} \end{cases}$$

$$\varphi(x, y) = (1 - \text{prime}(x)) (1 - \text{prime}(y)) \exp(x+1, y)$$

$$\text{prime}(x) = (1 - \text{equals}(x, 2)) \prod_{i=2}^{x-1} (1 - \text{rem}(x, i) \neq 0) + \text{equals}(x, 2)$$

$$\exp(x, y) = (1 \sim \text{equals}(y, 0)) (\text{multi}(\exp(x, y-1), x)) + \text{equals}(y, 0)$$

故原函数是原始递归的.

08~09

$g(x, y)$ 是原始递归函数

$$e(x, y) = \begin{cases} 1, & \text{if } \exists t, 0 \leq t < y, (g(x, t) = 0) \\ 0, & \text{otherwise} \end{cases}$$

$$e(x, y) = \sum_{t=0}^{y-1} (g(x, t) = 0)$$

$g(x, y)$ 原始递归, is zero 判定也是, 析取也是, 故证毕

06~07

$$f(x) = \begin{cases} x+1, & x \text{ is odd} \\ 4x, & x \text{ is even} \end{cases}$$

$$f(x) = (\text{rem}(x, 2)) (x+1) + (1 \sim \text{rem}(x, 2)) \cdot 4x$$

04~05

$P(x, y)$ 原始递归, 证明 $\forall y \in \mathbb{N} P(x, y)$

$$\forall y \in \mathbb{N} P(x, y) \Leftrightarrow \prod_{y=0}^{\infty} P(x, y) \neq 0$$

$$\text{类比: } \exists t (0 \leq t < y) (q(x, y) \neq 0) \Leftrightarrow \sum_{t=0}^{y-1} (q(x, y) \neq 0)$$

$$\text{fac}(n) = n : \\ (1 - \text{equals}(n, 0)) \text{fac}(n-1) \cdot n + \text{equals}(n, 0)$$

$$\text{gcd}(m, n) = \begin{cases} n, & \text{if } m \% n = 0 \\ \text{gcd}(n, m \% n), & \text{if } m \% n \neq 0 \end{cases}$$

$$\text{gcd} = (1 \sim \text{rem}(m, n)) n + \text{rem}(m, n) \text{gcd}(n, m \% n)$$

$$\text{prime}(n) = (1 - \text{equals}(n, 2)) \prod_{i=2}^{n-1} (\text{rem}(n, i) \neq 0) + \text{equals}(n, 2)$$

总结: 类似于 muu 编程, 需要细心.

$$\varphi(x) = \begin{cases} |x-y|, & (x \equiv y) \bmod 3 \\ x+y, & \text{otherwise} \end{cases}$$

$$\text{mod } 3(x, y) = ((x-y) \geq 0) (1 \sim \text{rem}(x-y, 3)) + ((x-y) < 0) (1 \sim \text{rem}(y-x, 3))$$

$$\varphi(x) = \text{mod } 3(x, y) ((x-y) \geq 0) (x-y) + ((x-y) < 0) (y-x) + (1 \sim \text{mod } 3(x, y)) (x+y)$$