



Efficient and Exact Multi-Marginal Optimal Transport with Pairwise Costs



DARTMOUTH

Bohan Zhou,¹ Matthew Parno.¹

¹Dartmouth College

Abstract

We provide an exact and efficient method to solve Multimarginal Optimal Transport (MMOT) under a family of cost functions:

$$\inf_{P \in \Gamma(\mu_1, \dots, \mu_m)} \int c(x_1, \dots, x_m) dP(x_1, \dots, x_m), \quad (1)$$

for the space $\mathbf{X} = X_1 \times \dots \times X_m$ and prescribed marginal probability measures $(\mu_i)_{i=1}^m$. The set of transport plans $\Gamma(\mu_1, \dots, \mu_m)$ is defined by

$$\Gamma(\mu_1, \dots, \mu_m) \stackrel{\text{def}}{=} \{P \in \mathbb{P}(\mathbf{X}) \mid (\pi_i)_\# P = \mu_i, 1 \leq i \leq m\}.$$

We assume the cost function $c(x_1, \dots, x_m)$ satisfies:

- The cost function is a summed pairwise cost functions
- $c(x_1, \dots, x_m) = \sum_{1 \leq i < j \leq m} c_{ij}(x_i, x_j);$
- $c_{ij}(x_i, x_j) = h_{ij}(x_i - x_j)$ for some strictly convex function $h_{ij}.$

Preliminary

Background

① c -transform and Duality theory:

The c -transform of a function $f : X_1 \mapsto \mathbb{R}$ is given by

$$f^c(x_2) = \inf_{x_1} f(x_1, x_2) - f(x_1).$$

It is natural to have $f(x_1) + f^c(x_2) \leq c(x_1, x_2)$. The c -transform is a generalization of the Legendre transform $f^*(y) = \sup_x x \cdot y - f(x)$.

We say (f_1, f_2) are c -conjugate if $f_1 = f_2^c$ and $f_2 = f_1^c$.

The dual problem corresponding to (1) is given by

$$\sup_{(f_1, \dots, f_m)} \sum_{i=1}^m \int_{X_i} f_i(x_i) d\mu_i, \quad (2)$$

where $f_i \in L^1(\mu_i)$ and $\sum_{i=1}^m f_i(x_i) \leq c(x_1, \dots, x_m)$.

[Kel84] provided a general duality theorem: there exists a c -conjugate solution to (2). We have the relationship between the primal solution and dual solution:

$$\sum_{i=1}^m f_i(x_i) = c(x_1, \dots, x_m) \quad P\text{-a.e..} \quad (3)$$

② Gradient in Hilbert space and back-and-forth method by [JL20] to solve 2-marginal OT under cost $c(x_1, x_2) = h(x_1 - x_2)$ for some strictly convex function h :

- For functional $I(f) = \int f d\mu_1 + \int f^c d\mu_2$, first find the first variance δI by a perturbation lemma [GM96];
- Define the gradient in a Hilbert space $(\mathcal{H}, \langle \cdot, \cdot \rangle)$:

$$\langle \nabla_{\mathcal{H}} I(f), g \rangle = \delta I(g; f).$$

[JL20] picked the space \dot{H}^1 with the inner product $\langle f_1, f_2 \rangle_{\dot{H}^1} = \int \nabla f_1 \cdot \nabla f_2 dx$ for the dual variables. The gradient is given by

$$I(f) = (-\Delta)^{-1} (\mu_1 - (S_{f^c})_\# \mu_2), \quad (4)$$

where the Brenier map $S_f(x_1) \stackrel{\text{def}}{=} x_1 - \nabla h^*(\nabla f(x_1))$.

To solve (2) for $m = 2$, [JL20] used a gradient-ascent scheme to update two functionals of type $I(f)$, depending on each dual variables $(f_i)_{i=1}^2$, in a back-and-forth fashion.

Current Computational Methods

Here, we list several MMOT solvers to our best knowledge. In general, entropy-regularized based methods may suffer from numerical instability and blurring issues. LP based methods may not be practical in solving large-scale problems.

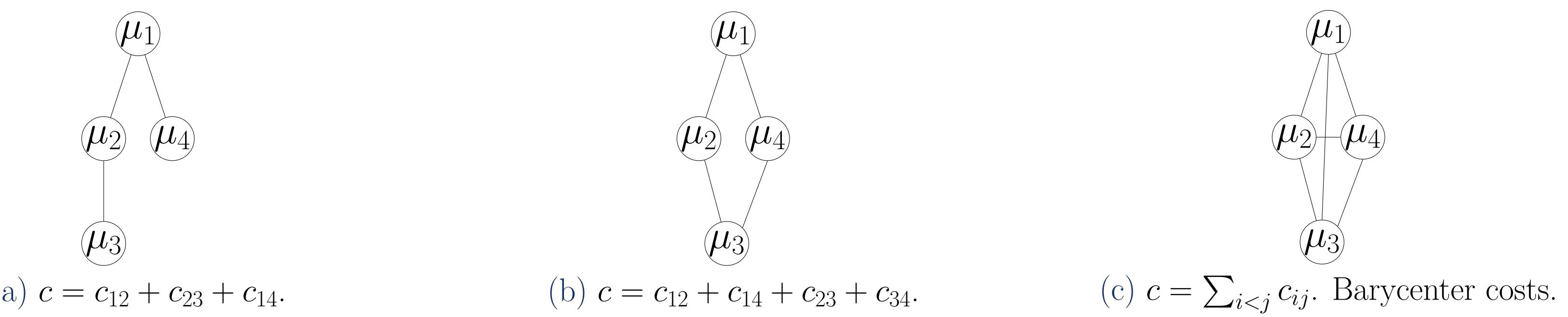
- [BCC+15]: Entropy-regularized MMOT on primal variables.
- [HRCK21]: Entropy-regularized MMOT with structure on dual variables.
- [ABA22]: Solving exact MMOT with structure via ellipsoid algorithm with oracle.
- [NX22]: LP-based method to approximate MMOT with controllable level of sub-optimality.

Main results

Our Strategy

- Graphical Representation of MMOT** Given a summed pairwise cost function, we can represent the relationship between their marginals on a graph, each node stores a marginal μ_i with its dual variable f_i , each edge stores pairwise cost c_{ij} .
- Unrolling MMOT into a tree representation** We prove an equivalent theorem that any MMOT that has a graphical representation with possible cycle is equivalent to another MMOT of a tree representation. The proof is via duplicating nodes and generalized gluing lemma. We also show that the cost of duplicating is limited by the number of edges in original graph.
- Solving MMOT on the rooted tree via gradient-ascent** By leveraging c -transform to get rid of the constraint, we will use gradient ascent on the remaining $(m-1)$ dual variables in the space \dot{H}^1 . The key in this step is that we introduce a “net potential”, which help us to define the gradient and to compute the c -transform.

Step 1: Graphical Representation



Step 2: Tree Representation [Equivalence theorem]

Given a cost function $c(x_1, \dots, x_m)$ that corresponds to $G = (V, E)$ with possible cycles, there exists a cost function $\bar{c}(x_1, \dots, x_n)$ that corresponds to tree $\bar{G} = (\bar{V}, \bar{E})$ with $n = |\bar{V}| = |E| + 1$ nodes, such that

$$\inf_{P^{(m)} \in \Gamma(\mu_1, \dots, \mu_m)} \int c(x_1, \dots, x_m) dP^{(m)} = \inf_{P^{(n)} \in \Gamma(\mu_1, \dots, \mu_n)} \int \bar{c}(x_1, \dots, x_n) dP^{(n)},$$

where $(\mu_k)_{k=m+1}^n$ are duplicated from $(\mu_i)_{i=1}^m$ in the “unrolling” process.

Furthermore, let $P^{(m)}$ and $(f_i)_{i=1}^m$ be the optimal primal and dual solutions to the original MMOT. And $P^{(n)}$ and $(\tilde{f}_i)_{i=1}^n$ be the optimal primal and dual solutions to the new MMOT.

Then for any i , the original dual variable f_i is the sum of all \tilde{f}_j whose nodes are duplicated from μ_i .



Step 3: Gradient-ascent on Rooted Tree

Define $I_r(f_1, \dots, f_{r-1}, f_{r+1}, \dots, f_m) \stackrel{\text{def}}{=} I(f_1, \dots, f_{r-1}, (\sum_{i \neq r} f_i)^c, f_{r+1}, \dots, f_m)$. The updates are:

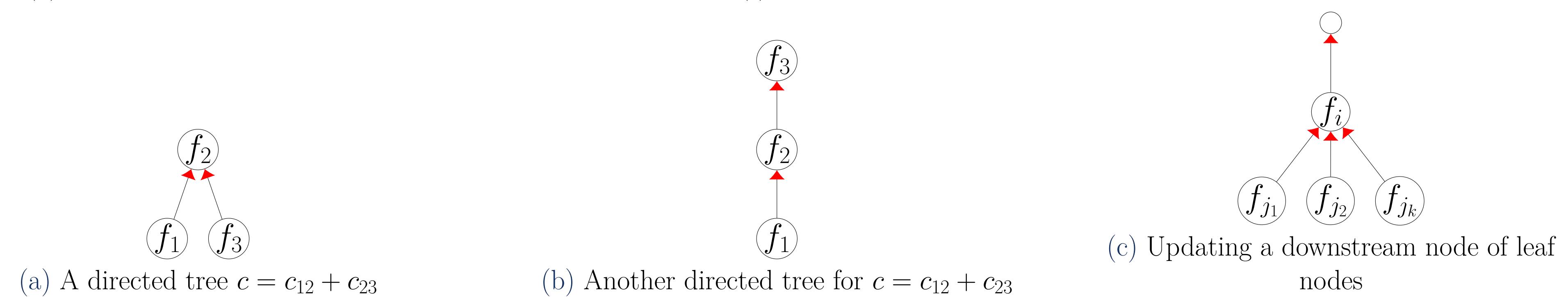
$$\begin{cases} f_i \leftarrow f_i - \sigma \nabla_{H^1} I_r(f_i); \\ f_r \leftarrow (\sum_{i \neq r} f_i)^c. \end{cases} \quad (5)$$

$$\begin{cases} \nabla_{H^1} I_r(f_i) = (-\Delta)^{-1} (\mu_i - (S_{f_i^c})_\# \mu_{N^+(i)}) \\ f_r(x_r) = \sum_{i \in N^-(r)} f'_i(x_r). \end{cases} \quad (6a)$$

$$(6b)$$

where the net potential f'_i at edge $(i, N^+(i))$ we introduced, are recursively defined by $f'_i = \left(f_i - \sum_{j \in N^-(i)} f'_j \right)^{c_{iN^+(i)}}$.

Here $N^-(r)$ are the collections of upstream nodes of root node r and $N^+(i)$ are the one downstream node of node i .



Numerical Results

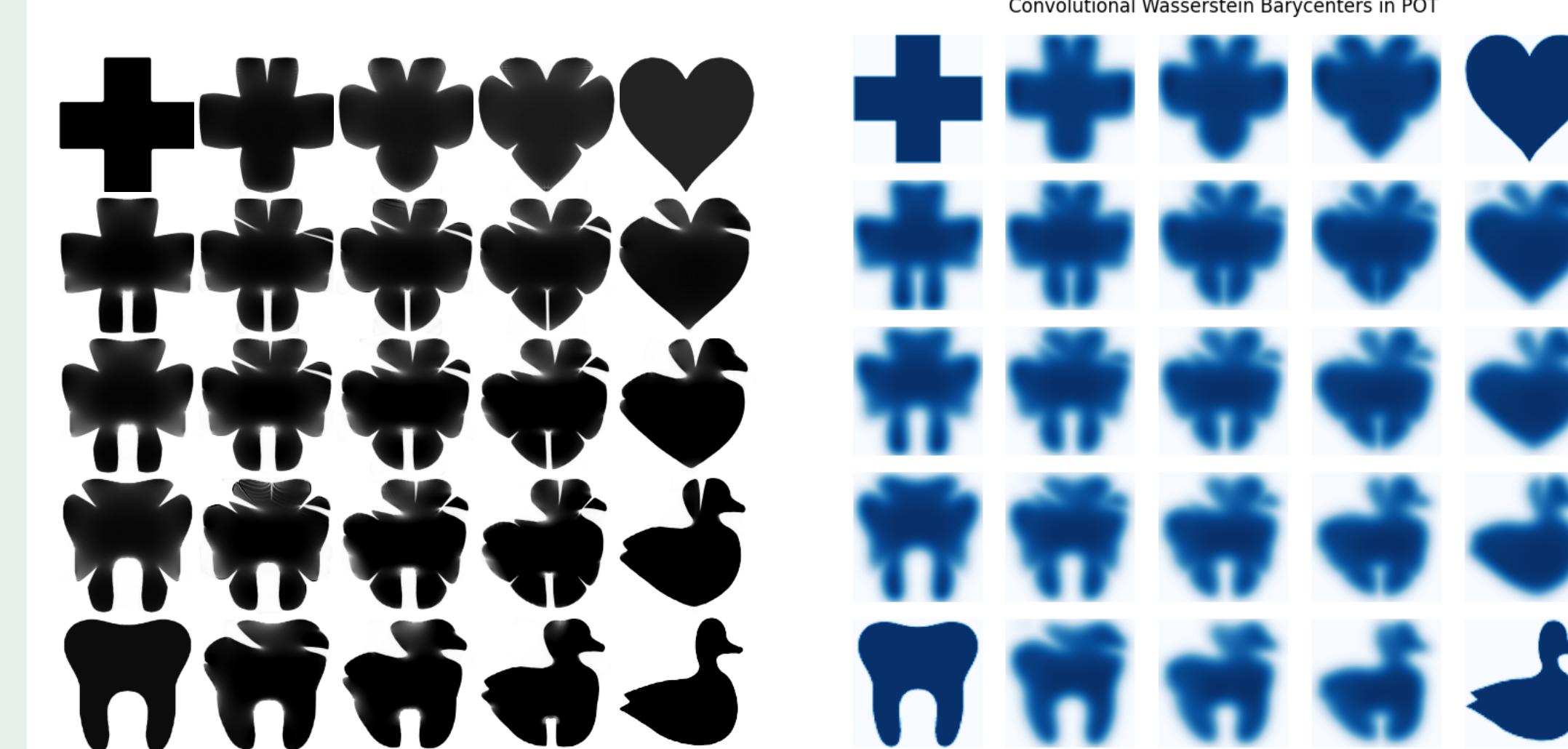
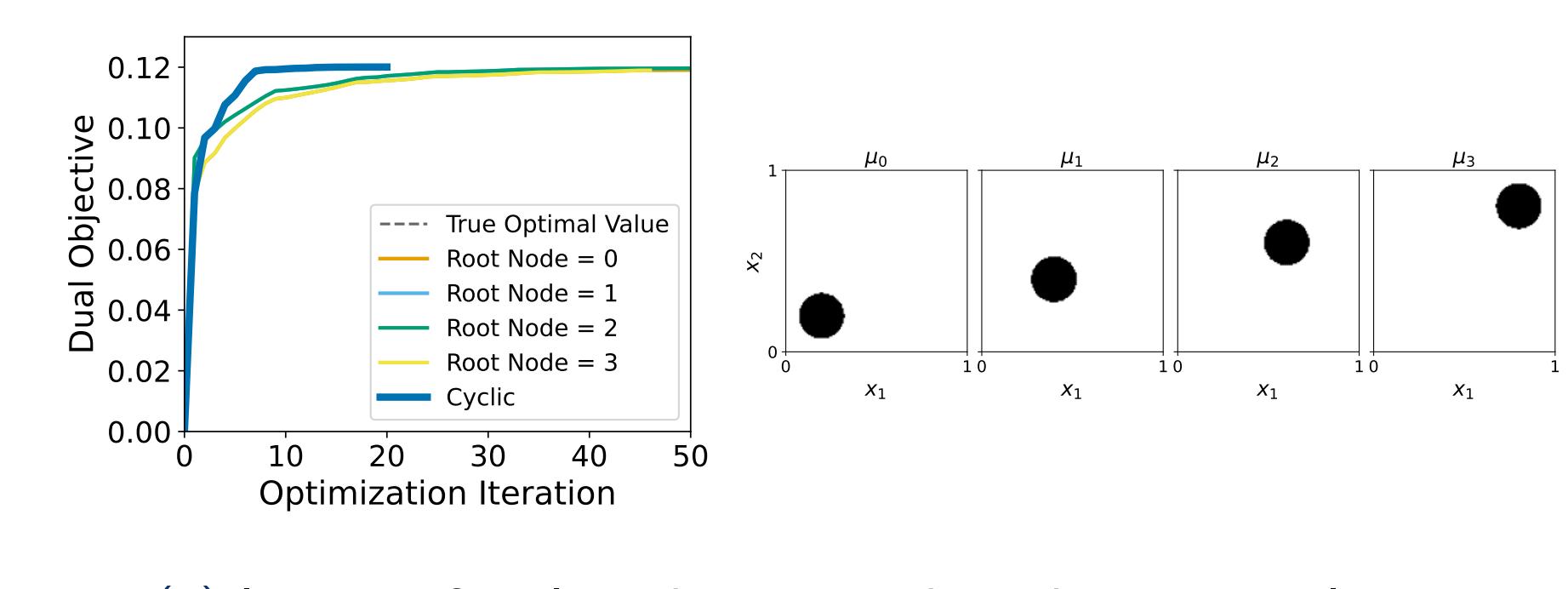
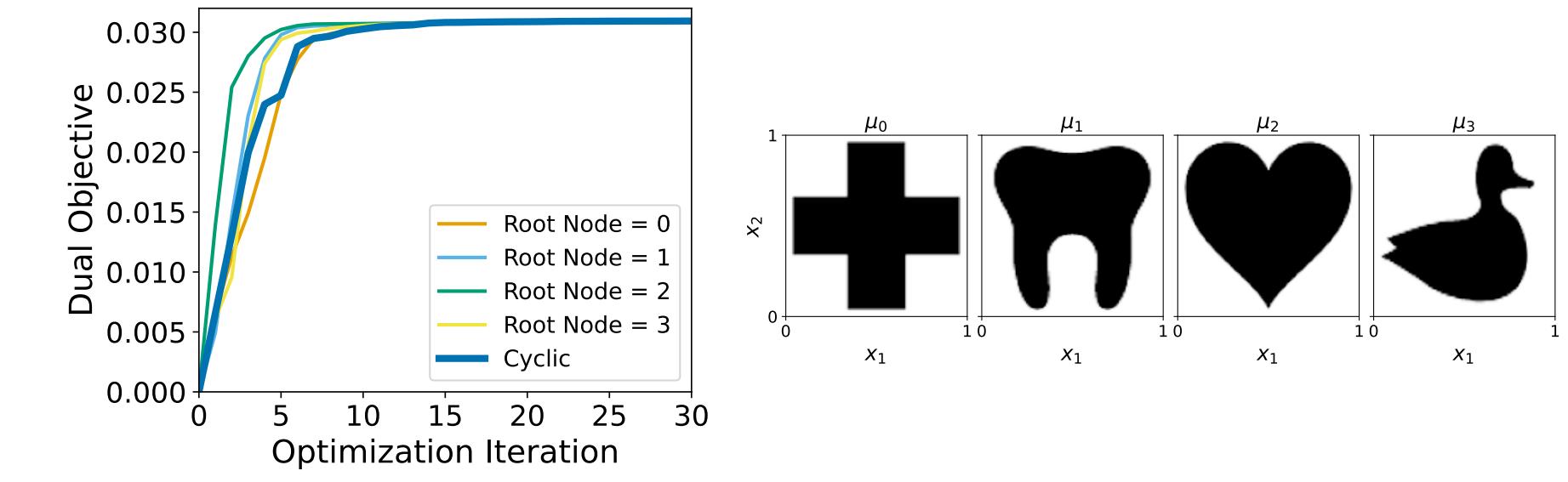


Figure: Left plot: sharp Wasserstein barycenter via our method. Right plot: blurred Wasserstein barycenter via entropy-regularized based method in POT package, regularization parameter is 0.004. Both 4-marginals are given at four corners.



(a) Impact of cycling the root node with pure translation.



(b) Impact of cycling the root node with shape deformation.

Figure: Performance between fixed root node and cyclic root node. The cyclic root node is a variant of the main algorithm which forces c -conjugate condition.

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