I-code Working-Notes 2v0
Peter S. Robertson
1 October 1984
Updated by

John D. McMullin 30th September 2023

Note: This document is not intended to be a complete formal description of I-code

Copyright (c) 1984 Lattice Logic Limited 9, Wemyss Place Edinburgh EH3 6DH Scotland

Contents

1.	Philosophy	7
2.	Definitions	8
3.	Conventions	10
4.	ICODE INSTRUCTIONS as ASCII	11
5.	ICODE GROUPS	13
5.	1. Stack Operations	13
	ADD	13
	ADDA	13
	AND	13
	CONCAT	14
	DIVIDE	14
	IEXP	14
	LSH	15
	MOD	15
	MUL	15
	NEGATE	16
	NOT	16
	OR	16
	QUOT	16
	REXP	17
	RSH	17
	SUB	17
	SUBA	17
	XOR	18
5.	2. Jumps and Labels	19
	FOR	19
	GOTO	19
	JE	21
	JG	21
	JGE	22
	JL	22
	JLE	
	JNE	23
	JNZ	24
	IIIMP	24

JZ	Z	. 25
L	ABEL	. 25
L	OCATE	. 25
R	REPEAT	. 26
S.	JUMP	. 26
S	SLABEL	. 26
5.3.	. Assignment	. 27
Α	ASSPAR	. 27
Α	ASSREF	. 27
Α	ASSVAL	. 27
F.	ALSE	. 28
JA	AM	. 28
Ν	ЛАР	. 28
R	RESULT	. 28
R	RETURN	. 29
Т	RUE	. 29
5.4.	. Embedding Machine Code	.30
Ν	MCODE	.30
Р	PLANT	.30
5.5.	. Event/Signal	.31
Ε	VENT	.31
С	DN	.31
S	TOP	.31
5.6.	. Compare/Test	. 32
С	COMPARE	.32
С	COMPAREA	. 32
С	COMPARED	.32
5.7.	. Parameter/Variable Declaration	. 33
Α	ALIAS	. 33
Α	ALT	. 33
В	3OUNDS	.33
D	DEF	. 34
D	DIM	.36
F	INISH	.37
11	NIT	.37
S	SETFORMAT	. 38

	START	38
5.	8. Access Variables	39
	ACCESS	39
	INDEX	39
	PUSH	39
	PUSHI	39
	PUSHR	39
	PUSHS	40
	RESOLVE	40
	SELECT	40
5.	9. Compiler Directives	41
	CONTROL	41
	DIAG	41
	EOF	41
	LANG	41
	LINE	41
	MONITOR	41
5.	10. Block Structure/Call	42
	BEGIN	42
	CALL	42
	END	42
6.	Unused ICODE Instructions	43
	Absolute	43
	Address	43
	Adjust	43
	Byte	43
	Compare-Repeated-Values	44
	Compare-Unsigned-Values	44
	Complement	44
	Define-Range	44
	Duplicate	44
	Eval	45
	Eval-Addr	45
	Float	45
	Include	45
	Init-Type	45

Int		45
Intpt		46
Localise.		46
Null-Set.		46
Pop		46
Reference	ce	46
Remaind	der	47
Round		47
Size-Of		47
Stack-Co	ondition	47
Stack-In.		48
Stack-Ur	nsigned-Condition	48
Swop		48
Test-Boo	blean	48
Test-In		49
Test-Nil .		49
Test-Ran	nge	49
Trunc		49
Variable-	-Call	50
Appendix 1.	Instructions which set the condition code	51
Appendix 2.	Instructions which test the condition code	51

1. Philosophy

I-code is an intermediate code used to provide an interface between the machine-independent and machine-dependent sections of a compiler.

Most intermediate codes in use today describe the execution of an abstract machine which performs the desired computation. For example, the intermediate code for a statement of the form: X = Y+Z would describe the operations of the abstract machine which would compute the value of Y+Z and assign it to X. Using this sort of code, the machine-dependent section of the compiler then maps the abstract machine onto the real target machine.

I-code uses a fundamentally different model. It describes the execution of an abstract compiler which generates target code to perform the desired computation; it does not describe an abstract machine which will perform that computation. It is vital to understand that it does not describe the function of the program directly but describes it indirectly via the abstract compiler. It is this indirection which gives I-code the power to be machine-independent without sacrificing efficiency in the executable programs which it can be used to generate.

There are two important corollaries of this. Firstly, the structures and operations associated with the abstract compiler need have no counterparts in the object program. For example, the target machine need have no hardware or software stack and neither need it have a true condition code. Secondly the code assumes that the operations it describes will be performed by the abstract compiler in the order specified with no omissions. In particular, the control transfer instructions do not transfer control in the abstract compiler but indicate changes of control flow in the program which is being compiled. This also does not mean that any of the operations need have counterparts in the object program, nor that the order of the generated code corresponds to the order of the I-code.

For example, the Algol-60 statement: A := if B then C else D; could not be encoded in the seemingly obvious way:

```
Stack A
```

Stack B; Test-Boolean; BF 1

Stack C; Forward 2

Label 1

Stack D

Label 2

Assign-Value

as this would assign D to C, the last two objects stacked prior to the Assign-Value instruction and leave A on the stack.

Encode instead as:

Stack B; Test-Boolean; BF 1

Stack A; Stack C; Assign-Value; Forward 2

Label 1

Stack A; Stack D; Assign-Value

Label 2

This is similar to that generated from the code fragment: if B then A := C else A := D;

N.B. Stack has the same usage as PUSH.

2. Definitions

The ICODE instructions and parameters are stored as encoded data in a binary file using the following definitions.

byte-order All multi-byte values are specified with the least-significant byte first.					
Unsigned	A natural binary number.				
Signed	A 2's complement binary number.				
	A one-byte unsigned number.				
<integer></integer>	A four-byte signed integer number.				
<label></label>	An unsigned 16-bit value used to identify a simple label.				
<n></n>	An unsigned 16-bit number.				
<string></string>	A byte-counted string constant.				
<real></real>	A real constant in a textual encoding.				
<tag></tag>	An unsigned 16-bit value used to identify a tag (descriptor).				
condition-code	A conceptual flag which is set at run-time by the instructions listed in Appendix 2.				
	This flag need not exist in the target machine but is defined in order to simplify the definitions of certain instructions.				
	The values which this flag may take are: equal, less than, greater than, true, false				
	This setting only remains valid for the duration of the next instruction which must				
	be one of the instructions listed in Appendix 3.				
Integer	A general integer value, including subranges of integers.				
Labels	I-code distinguishes two sorts of label.				
1 – Simple Labels: have the property that they are only jumped to in one direction, that is, all references to an instance of a simple label are either forward or all backward. This means that the same denotation may be use many simple labels. For example, the following is valid:					
	Goto 1 >+				
	Goto 1 >+				
	Goto 1 >+				
	Label 1				
	Label 1 <+				
	All uses of a particular simple label must be in the same block as the definition of				
that label.					
	Simple labels are encoded as two-byte unsigned integers although code generators may assume that their values are within a fairly small range. (1 50 is common.)				
	2. General Labels - have none of the restrictions of simple labels. They are identified by tags and will be defined automatically if necessary when they are first used. General labels are referenced by the instructions:				

	Jump <tag></tag>		
	Locate <tag></tag>		
	Def <tag></tag>		
Real	A floating-point value, either single or double precision.		
SOS	The second-top item on the stack.		
Stack	A first-in, last-out structure used to imply the operands required by various I-code instructions.		
	The first item which can be removed from the stack is called TOS and the item which can be removed after TOS is called SOS.		
Tags	Tags are definitions of objects which are to be manipulated by the compiler. These definitions are created in a nested fashion; all tags defined in a block are deleted when the end of that block is reached.		
	On definition the machine-independent description of the object is converted into the appropriate machine-dependent description of the actual object to be used. Within this document the term 'tag' is used to describe both this descriptor and the unsigned sixteen-bit integer used as an index value to select it from the collection of all tags.		
	Apart from the resolution of forward references to procedures, tags are never altered. When a copy of a tag value is pushed onto the stack the value becomes known as		
	a descriptor. Descriptors may be modified.		
tag list	A tag list is an ordered sequence of tag definitions used to describe either the parameters required by a procedure or the fields of a record. Components of a tag list are referred to either explicitly by their position in the list (see SELECT) or implicitly by sequence starting with the first to be specified (see Assign-Parameter).		
TOS	The top item on the stack.		
list flag	An internal flag which is set during the processing of explicit lists of tag definitions. Its only purpose is to prevent certain nested list structures.		

3. Conventions

- 1. It is assumed that the reader is familiar with the IMP language and its terminology.
- 2. Whenever a pointer-variable is used in the context of a value the value of the data item pointed to will be used.
- 3. In general, diagnostic checks are implied rather than explicitly specified.
- 4. Items on the stack are intended to be 'rules for generating' values or references rather than the values or references themselves. For simplicity the descriptions of instructions will often refer to items as if they contain values or references.
- 5. By convention tag index values will often be replaced in examples by the identifier which is assumed to have been associated with the tag in question.
 - For example, given 'DEFINE 57,Fred......' then 'Stack 57' could be written 'Stack Fred' PUSH is equivalent to Stack
- 6. The term 'error' is used to indicate a condition discovered by the code-generator which should terminate the compilation with a suitable error message.

4. ICODE INSTRUCTIONS as ASCII

ASCII Value	ASCII Character	ICODE INSTRUCTION	ICODE PARAMETERS
09	<nul><ht></ht></nul>	unused	
10	<lf></lf>	EOF	
1131	<vt><us></us></vt>	unused	
32	<space></space>	unused	
33	!	OR	
34	"	COMPARED	
35	#	JNE	Tag
36	\$	DEF	Tag String ',' Tag ',' Tag ',' Tag
37	%	XOR	0 07 07 07
38	&	AND	
39	,	PUSHS	String
40	(JLE	Tag
41)	JGE	Tag
42	*	MUL	
43	+	ADD	
44	,	parameter separator	
45	-	SUB	
46		CONCAT	
47	/	QUOT	
4857	09	unused	
58	:	LOCATE	Tag
59	;	END	
60	<	JL	Tag
61	=	JE	Tag
62	>	JG	Tag
63	?	COMPARE	
64	@	PUSH	Tag
65	Α	INIT	Tag
66	В	REPEAT	Tag
67	С	COMPAREA	
68	D	PUSHR	Real
69	E	CALL	
70	F	GOTO	Tag
71	G	ALIAS	String
72	Н	BEGIN	
73	1	unused	
74	J	JUMP	Tag
75	К	FALSE	
76	L	LABEL	Tag
77	М	MAP	
78	N	PUSHI	Integer
79	0	LINE	Tag
80	Р	PLANT	Byte
81	Q	DIVIDE	
82	R	RETURN	
83	S	ASSVAL	
84	Т	TRUE	

85	U	NEGATE	
86	V	RESULT	
87	W	SJUMP	Tag
88	X	IEXP	108
89	Υ	unused	
90	Z	ASSREF	
91	[LSH	
92	\	NOT	
93	1	RSH	
94	V .	SETFORMAT	Tag
95		SLABEL	Tag
96	`	unused	l
97	а	ACCESS	
98	b	BOUNDS	
99		unused	
100	c d	DIM	Tag ' ' Tag
			Tag ',' Tag
101	e f	EVENT FOR	Tag
102			Tag
103	g	unused	
104	h	unused	
105	i	INDEX	
106	J	JAM	
107	k	JZ	Tag
108	1	LANG	Tag
109	m	MONITOR	
110	n	SELECT	Tag
111	0	ON	Tag ',' Tag
112	р	ASSPAR	
113	q	SUBA	
114	r	RESOLVE	Tag
115	S	STOP	
116	t	JNZ	Tag
117	u	ADDA	
118	V	MOD	
119	W	MCODE	String ';'
120	Х	REXP	
121	У	DIAG	Tag
122	Z	CONTROL	Tag
123	{	START	
124	1	unused	
125	}	FINISH	
126	~	ALT	Byte
127255		unused	

All other byte values (in the range 0..255) are currently ILLEGAL/UNUSED.

Extra ICODE instructions can be added by using unused ASCII characters with parameters as needed.

5. ICODE GROUPS

5.1. Stack Operations

These ICODE instructions represent operations placed on an operation stack for code generation on the IMP variables.

Instruction:	ADD	'+'	Load Operation(ADDx)	
Effect:	TOS is replaced by the sum of values of TOS, SOS. Integer values will be converted into floating-point if one operand is integer and the other is real.			
Notes:				
Error:	 The stack is empty. TOS,SOS are not integer or real. 			
Example:	X = Y+Z	PUSH X PUSH Y PUSH Z ADD ASSVAL		

Instruction:	ADDA	ʻu'	Load Double(ADDx)
Effect:	TOS is replaced by the sum of values of TOS, SOS.		
Notes:	Used to add two address values/two double real values		
Error:	1. The stack is empty.		
	TOS,SOS are not int	eger or real.	
Example:	X = Y + Z	PUSH X	
		PUSH Y	
		PUSH Z	
		ADDA	
		ASSVAL	

Instruction:	AND	' &'	Load Operation(ANDx)	
Effect:	TOS,SOS are replaced by the logical and of the values of TOS, SOS.			
Notes:	Used to AND two logical values			
Error:	1. The stack has less than 2 items.			
	TOS,SOS are not int	2. TOS,SOS are not integer values.		
Example:	X = Y & Z	PUSH X		
		PUSH Y		
		PUSH Z		
		AND		
		ASSVAL		

Instruction:	CONCAT		Load Operation(CONCx)	
Effect:	TOS and SOS are removed from the stack and the string concatenation of their			
	values, SOS.TOS, is stacked.			
Notes:				
Error:	 The stack contains 	less than two values.		
	2. TOS is not a string value.			
	SOS is not a string v	SOS is not a string value.		
Example: S = T.U.V PUSH S				
		PUSH T		
		PUSH U		
		CONCAT		
		PUSH V		
		CONCAT		
		ASSVAL		

Instruction:	DIVIDE	'Q'	Load Operation(RDIVx)	
Effect:	TOS and SOS are removed from the stack and the real quotient, SOS / TOS, is stacked. Integer values will be converted into floating-point before the division is			
	attempted.			
Notes:	·			
Error:	The stack contains less than two items.			
	TOS is neither integ	er nor real type.		
	SOS is neither integ	neither integer nor real type.		
Example:	A = B / C	PUSH A		
		PUSH B		
	PUSH C			
		DIV		
		ASSVAL		

Instruction:	IEXP	'X'	Load Operation(EXPx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS raised to the integer			
	power of the value of TOS,	SOS^^TOS, is stacked.		
Notes:				
Error:	The stack contains less than two items.			
	TOS is not an integer	er value.		
	3. SOS is not an integer value.			
Example:	J = K^^3	PUSH J		
		PUSH K		
	PUSHI 3			
		IEXP		
		ASSVAL		

Instruction:	LSH	<u>([</u> '	Load Operation(LSHx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS logically shifted left			
	by the value of TOS, SOS <<	TOS, is stacked.		
Notes:				
Error:	1. The stack contains	less than two items.		
	TOS is not an integer	er value.		
	SOS is not an integer	er value.		
	4. The value of TOS is negative or greater than or equal to the number of bits			
	in an integer.			
Example:	A = B << C	PUSH A		
		PUSH B		
	PUSH C			
	LSH			
		ASSVAL		

Instruction:	MOD	'v'	Load Operation(ABSx)		
Effect:	TOS and SOS are removed from the stack and are replaced by the value 'SOS MOD				
	TOS' where MOD is as defin	ed in section 6.7.2.2 of the P	ascal standard BS		
	6192:1982.				
Notes:					
Error:	 The stack contains I 	less than two items.			
	2. TOS is not an integer value.				
	SOS is not an integer	not an integer value.			
Example:	m := p MOD q;	PUSH m			
		PUSH p			
		PUSH q			
		MOD			
		ASSVAL			

Instruction:	MUL	(*)	Load Operation(MULx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS multiplied by the			
	value of TOS, SOS * TOS, is s	stacked. Integer values will be	e converted into floating-	
	point if one operand is integ	ger and the other is real.		
Notes:				
Error:	The stack contains less than two items.			
	2. TOS is neither integer nor real type.			
	3. SOS is neither integer nor real type.			
Example:	A = B * C	PUSH A		
	PUSH B			
	PUSH C			
		MUL		
		ASSVAL		

Instruction:	NEGATE	'U'	Load Operation(NEGx)	
Effect:	The value in TOS is negated	The value in TOS is negated and left as TOS.		
Notes:				
Error:	1. The stack is empty.			
	2. TOS is not an integer or real value.			
Example:	A = -B PUSH A			
	PUSH B			
	NEGATE			
		ASSVAL		

Instruction:	NOT	\\'	Load Operation(NOTx)
Effect:	The value in TOS is logically inverted and left on TOS		
Notes:			
Error:	1. The stack is empty.		
	2. TOS is not an integer or real value.		
Example:			

Instruction:	OR	'!'	Load Operation(ORx)
Effect:	TOS and SOS are removed from the stack and the value of SOS logically ORed with		
	the value of TOS, SOS! TOS	, is stacked.	
Notes:			
Error:	The stack contains less than two items.		
	TOS is not an integer	er value.	
	SOS is not an integer	er value.	
Example:	A = B ! C	PUSH A	
	1	PUSH B	
	PUSH C		
	OR		
		ASSVAL	

Instruction:	QUOT	<i>'/'</i>	Load Operation(DIVx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS integer-divided by			
	the value of TOS, SOS // TO	S, is stacked.		
Notes:				
Error:	The stack contains less than two items.			
	TOS is not an integer	er.		
	3. SOS is not an integer.			
Example:	A = B // C	PUSH A		
		PUSH B		
		PUSH C		
		QUOT		
		ASSVAL		

Instruction:	REXP	ʻx'	Load Operation(REXPx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS raised to the integer			
	power TOS, SOS^TOS, is stacked. The type of this value is real.			
Notes:				
Error:	The stack contains less than two items.			
	2. SOS is neither an integer nor a real value.			
	3. TOS is not an integer value.			
Example:	R = 12 ^X PUSH R			
		PUSHI 12; PUSH X; REXP;		
		ASSVAL		

Instruction:	RSH	(']'	Load Operation(RSHx)
Effect:	TOS and SOS are removed from the stack and the value of SOS logically left shifted by the value of TOS, SOS >> TOS, is stacked.		
Notes:			
Error:	 The stack contains less than two items. TOS is not an integer value. SOS is not an integer value. The value of TOS is negative or greater than or equal to the number of bits in an integer. 		
Example:	A = B >> C	PUSH A PUSH B; PUSH C; RSH ASSVAL	

Instruction:	SUB	1_1	Load Operation(SUBx)	
Effect:	TOS and SOS are removed from the stack and the value of SOS minus the value of			
	TOS, SOS - TOS, is stacked.			
	Integer values will be conve	erted into floating-point if one	e operand is integer and the	
	other is real.			
Notes:				
Error:	1. The stack contains less than two items.			
	2. TOS is neither integer nor real type.			
	3. SOS is neither integer nor real type.			
Example:	A = B - C PUSH A			
		PUSH B; PUSH C; SUB		
		ASSVAL		

Instruction:	SUBA	ʻqʻ	Load DoubleOp(SUBx)	
Effect:	Similar to Sub except general	ated code is for addresses		
Notes:	Equivalent to			
Error:	The stack contains less than two items.			
	2. TOS is neither integer nor real type.			
	3. SOS is neither integer nor real type.			
Example:				

Instruction:	XOR	' %'		Load Operation(XORx)
Effect:	TOS and SOS are removed fr			of SOS exclusively ORed
	with the value of TOS, SOS!	! TOS, is stacke	ed.	
Notes:				
Error:	The stack contains less than two items.			
	2. TOS is not an integer value.			
	3. SOS is not an integer value.			
Example:	A = B !! C PUSH A			
			PUSH B; PUSH	HC; XOR
			ASSVAL	

5.2. Jumps and Labels

These ICODE instructions are used to control the flow of executable code by the program.

There are two types of labels:

- User specified labels
- Compiler generated labels

Instruction:	FOR <label></label>	'f'		Compile For(Tag)
Effect:	This instruction marks the s <			+1> is the label to be
	jumped to on an <u>exit</u> .			
Notes:	The corresponding <u>repeat</u> w Backward < label>.	vill be the next i	instruction of	the form:
	<label+1> should only be de</label+1>	efined explicitly	if the loop co	ntains an explicit <u>exit</u>
	instruction.			
Error:	The stack contains less than four items.			
	2. The top three items on the stack are not integer values.			
	3. The fourth item on	the stack is not	a reference to	o an integer variable.
	4. No Backward <labe< td=""><td>l> instruction o</td><td>ccurs before t</td><td>he end of the current block.</td></labe<>	l> instruction o	ccurs before t	he end of the current block.
Example:	A(j) = 0 for J = 1, 1, N		PUSH J	
			PUSHI 1	
			PUSHI 1	
			PUSH N	
			FOR 40	
		PUSH A		
		PUSH J		
	ACCESS			
			PUSHI 0	
			ASSVAL	
			REPEAT 40	

Instruction:	GOTO <label></label>	'F'		Jump Forward(Tag, Always)	
Effect:	Control is transferred for	ward unconditio	nally to <labe< td=""><td>el></td></labe<>	el>	
Notes:					
Error:	 <label> does not get defined by a Label instruction before the end of the current block.</label> 				
Example:	<u>if</u> X=0 <u>then</u> Y=1 <u>else</u> Y=2		PUSH X		
			PUSHI 0		
			COMPARE		
			JNE 20		
			PUSH Y		
			PUSHI 1		
			ASSVAL		
			GOTO 21		
			LABEL 20		
			PUSH Y		
			PUSHI 2		
			ASSVAL		
			LABEL 21		

Instruction:	JE <label></label>	'='	Jump Forward(Tag, EQ)	
Effect:	When execution of the obje	ect program rea	aches this point, control is to be	
	transferred to the given sim	ple label if the	condition code is set 'equal', otherwise	
	control is to pass onto the n	next instruction	1.	
Notes:				
Error:	The previous instruction did not set the condition-code.			
	2. < label> does not get defined by a LABEL instruction before the end of the			
	current block.			
Example:	IF x <> y THEN p = q; PUSH x; PUSH y; COMPARE; JE 12			
		PUSH p; PUSH q; ASSVAL		
	LABEL 12			

Instruction:	JG <label></label>	'>'		Jump Forward(Tag, GT)	
Effect:	When execution of the obje	ect program rea	ches this poin	it, control is to be	
	transferred to the given sim	ple label if the	condition cod	le is set 'greater than',	
	otherwise control is to pass	onto the next	instruction.		
Notes:	<label> must refer to a simple.</label>	ole label which	must be defin	ed somewhere after the JG	
	instruction, that is JG can or	nly specify a for	ward jump (al	Ithough the object program	
	may use a backward jump).				
Error:	 The previous instru- 	ction did not se	t the conditio	n-code.	
	<label> does not ge</label>	t defined by a l	ABEL instruct	ion before the end of the	
	current block.				
Example:	<u>if</u> X <= Y <u>then</u> P = Q		PUSH X		
			PUSH Y		
			COMPARE		
	JG 12				
	PUSH P				
	PUSH Q				
			ASSVAL		
			LABEL 12		

Instruction:	JGE <label></label>	')'		Jump Forward(Tag, GE)		
Effect:	When execution of the obje	ct program rea	ches this poin	nt control is to be		
	transferred to the given sim	ple label if the	condition cod	le is set 'greater than' or		
	'equal' , otherwise control is	to pass onto t	he next instru	iction.		
Notes:	<label> must refer to a simp</label>	le label which	must be defin	ed somewhere after the		
	JGE instruction, that is JGE of	an only specify	a forward jur	mp (although the object		
	program may use a backwar	rd jump).				
Error:	 The previous instruction 	ction did not se	t the conditio	n-code.		
	2. <label> does not get</label>	t defined by a l	ABEL instruct	ion before the end of the		
	current block.	_				
Example:	<u>if</u> X < Y <u>then</u> P = Q		PUSH X			
			PUSH Y			
			COMPARE			
	JGE 12					
	PUSH P					
	PUSH Q					
		ASSVAL				
			LABEL 12			

Instruction:	JL <label></label>	' <'	Jump Forward(Tag, LT)		
Effect:	When execution of the obje	ect program reaches this poin	nt control is to be		
	transferred to the given sim	ple label if the condition cod	le is not set 'equal',		
	otherwise control is to pass	onto the next instruction.			
Notes:	<label> must refer to a simple</label>	ole label which must be defin	ed somewhere after the		
	BLT instruction, that is BLT	can only specify a forward jur	mp (although the object		
	program may use a backwa	rd jump).			
Error:	 The previous instru 	ction did not set the conditio	n-code.		
	<label> does not ge</label>	t defined by a LABEL instruct	ion before the end of the		
	current block.				
Example:	<u>if</u> X >= Y <u>then</u> P = Q	PUSH X			
		PUSH Y			
		COMPARE			
		JL 12			
		PUSH P			
		PUSH Q			
		ASSVAL			
		LABEL 12			

Instruction:	JLE <label></label>	'('	Jump Forward(Tag, LE)		
Effect:	When execution of the object program reaches this point, control is to be				
	transferred to the given sim	ple label if the condition coc	le is set 'less than' or 'equal',		
	otherwise control is to pass	onto the next instruction.			
Notes:	<label> must refer to a simp</label>	le label which must be defin	ed somewhere after the JLE		
	instruction, that is JLE can or	nly specify a forward jump (a	although the object		
	program may use a backwar	d jump).			
Error:	 The previous instruct 	tion did not set the condition	n-code.		
	<label> does not get</label>	defined by a LABEL instruct	ion before the end of the		
	current block.				
Example:	<u>if</u> X > Y <u>then</u> P = Q	PUSH X			
		PUSH Y			
		COMPARE			
	JLE 12				
		PUSH P			
		PUSH Q			
		ASSVAL			
		LABEL 12			

Instruction:	JNE <label></label>	'#'	Jump Forward(Tag, NE)		
Effect:	When execution of the obje	ect program reaches this poir	it control is to be		
	transferred to the given sim	ple label if the condition cod	le is not set 'equal',		
	otherwise control is to pass	onto the next instruction.			
Notes:	<label> must refer to a simple</label>	ole label which must be defin	ed somewhere after the		
	JNE instruction, that is JNE	can only specify a forward jur	np (although the object		
	program may use a backwa	rd jump).			
Error:	 The previous instru 	ction did not set the conditio	n-code.		
	<label> does not ge</label>	et defined by a LABEL instruct	ion before the end of the		
	current block.				
Example:	<u>if</u> X = Y <u>then</u> P = Q	PUSH X			
		PUSH Y			
		COMPARE			
		JNE 12			
		PUSH P			
		PUSH Q			
		ASSVAL			
		LABEL 12			

Instruction:	JNZ <label></label>	't'	JNZ			
Effect:	When execution of the obje	ect program reaches this poir	nt, control is to be			
	transferred to the given sim	ple label if the condition coc	le is set 'true', otherwise			
	control is to pass onto the r	next instruction.				
Notes:	Branch on TRUE (# 0)					
	<label> must refer to a simple.</label>	ole label which must be defin	ed somewhere after the			
	JNZ instruction, that is JNZ	can only specify a forward jui	mp (although the object			
	program may use a backwa	rd jump).				
Error:	 The previous instru 	ction did not set the condition	n-code.			
	<label> does not ge</label>	et defined by a LABEL instruct	ion before the end of the			
	current block.					
Example:	if not B then P := Q	PUSH B				
		Test-Boolean				
	JNZ 12					
	PUSH P					
		PUSH Q				
		ASSVAL				
		LABEL 12				

Instruction:	JUMP <tag></tag>	'J'	User Jump(Tag)		
Effect:	Control is transferred uncor	nditionally to the general labe	el <tag>.</tag>		
Notes:	Jump to user label				
	This control transfer may pa	ass over block boundaries.			
Error:	 <tag> is defined but</tag> 	is not a general label.			
	<tag> is undefined </tag>	out does not become defined	l in a suitable block before		
	the end of the prog	ram. A suitable block is one v	vhich is either the same		
	block as the Jump instruction or is a block which properly contains that				
	instruction.				
Example:	X = 1	PUSH X; PUSHI 1; ASSVAL			
	Pos: Y = Y+1	LOCATE Pos			
	->Pos <u>if</u> Y < 0	PUSH Y; PUSH Y; PUSHI 1;	ADD; ASSVAL		
	PUSH Y; PUSHI 0; COMPARE				
	JGE 18				
		JUMP Pos			
		LABEL 18			

Instruction:	JZ <label></label>	'k'	Jump Forward(Tag, FF)		
Effect:	When execution of the object program reaches this point, control is to be				
	transferred to the given sim	ple label if the condition cod	e is set 'false', otherwise		
	control is to pass onto the r	ext instruction.			
Notes:	Branch on FALSE (= 0)				
	<label> must refer to a simp</label>	ole label which must be define	ed somewhere after the BF		
	instruction, that is BF can only specify a forward jump (although the object program				
	may use a backward jump).				
Error:	The previous instruction did not set the condition-code.				
	2. <label> does not get defined by a LABEL instruction before the end of the</label>				
	current block.				
Example:	if B then P := Q; PUSH B; Test-Boolean; JZ 12				
		PUSH P; PUSH Q; ASSVAL			
		LABEL 12			

Instruction:	LABEL <label></label>	'L'		Define User Label(Tag)	
Effect:	The simple label <label> is defined to be here. If outstanding references to the label exist they are satisfied and the label ceases to be defined. This ensures that all</label>				
	references to this label are			a. This ensures that all	
	references to this laber are	iii tile saille uii	ection.		
Notes:	Define user label				
Error:					
Example:	S = "**" <u>if</u> S = ""		PUSH S; PUSH	HS "**"; COMPARE	
	JNE 26				
	PUSH S; PUSHS ""; ASSVAL				
			LABEL 26		

Instruction:	LOCATE <tag></tag>	<i>'.'</i>	Define Compiler Label(Tag)	
Effect:	The tag is defined as a gene	eral label if necessary and m	nade to reference the current	
	position in the program.			
Notes:	Define compiler label			
Error:	1. <tag> is already defined.</tag>			
Example:	X = 1	PUSH X; PUSHI 1; ASSVAL		
	Pos: Y = Y+1	LOCATE Pos		
	->Pos <u>if</u> Y < 0	PUSH Y; PUSH Y; PUSHI 1; ADD; ASSVAL		
		PUSH Y; PUSHI 0; COMPARE		
		JGE 18		
		JUMP Pos		
		LABEL 18		

Instruction:	REPEAT <label></label>	'B'	Jump Backward(Tag)	
Effect:	Control is to be transferred	unconditionally to <label> at</label>	run-time.	
Notes:	Backward Jump			
Error:	1. <label> is currently undefined.</label>			
Example:	X=X+1 while $A(X) = 0$	LABEL 16		
		PUSH A; PUSH X; ACCESS		
		PUSHI 0; COMPARE; BNE 17		
		PUSH X; PUSH X; PUSHI 1; ADD; ASSVAL		
		REPEAT 16		
		LABEL 17		

Instruction:	SJUMP <tag></tag>	'W'	Switch Jump(Tag)	
Effect:	TOS is used to index into the switch vector and control is then transferred to the			
	selected label. TOS is removed from the stack.			
Notes:	Jump to switch			
Error:	1. The stack is empty.			
	2. <tag> is not a switch.</tag>			
Example:	->Sw(J)	Stack J; Switch-Jump Sw		

Instruction:	SLABEL <tag></tag>		Switch Label(Tag)		
Effect:	The label selected from the	The label selected from the switch vector is defined to be here. TOS is removed			
	from the stack.				
Notes:	Define switch label				
Error:	1. The stack is empty.				
	2. <tag> is not a switch in the current block.</tag>				
	3. TOS is not an integer value within the bounds of <tag>.</tag>				
Example:	Sw(12):	PUSHI 12; SLABEL Sw			

5.3. Assignment

These ICODE instructions are used to access and store data within the program.

Instruction:	ASSPAR	ʻp'	Load Assign(-1)	
Effect:	TOS is passed as the next pa	arameter to SOS. TOS is remo	ved from the stack leaving	
	SOS as the new TOS.			
Notes:	Parameters must be specific	ed in the order of the definiti	on of the parameter list. If	
	the parameter list is empty	the occurrence of this instruc	ction implies that the	
	procedure has a variable nu	ımber of parameters and so t	he parameters are to be	
	passed in a C-like manner; t	his also requires that the pro	cedure be called using the	
	Variable-Call instruction.			
Error:	The stack contains less than two items.			
	2. SOS is not a procedure descriptor.			
	3. TOS is unsuitable for this parameter.			
Example:	J = Calc(1, K) PUSH J			
	PUSH Calc			
	PUSHI 1; ASSPAR			
	PUSH K; ASSPAR			
		CALL		

Instruction:	ASSREF	ʻZ'	Load Assign(0)	
Effect:	The pointer variable refere	nced by SOS is made to point	at the variable referenced	
	by TOS. Both TOS and SOS a	are removed from the stack.		
Notes:	Assign address '=='			
Error:	The stack contains less than two items.			
	2. SOS is not a reference to a pointer variable.			
	3. TOS is not a reference to a variable.			
	4. The types of TOS ar	The types of TOS and SOS are different.		
Example:	P == Q	PUSH P; PUSH Q; ASSREF		

The value of TOS is assigned			
The value of TOS is assigned to the data item referenced by SOS. Both TOS and SOS are removed from the stack.			
Integer values will be converted to real if necessary. Normal value assignment			
The stack contains less than two items.			
A = B+C	PUSH A PUSH B; PUSH C; ADD		
N	teger values will be conve ormal value assignment 1. The stack contains l	teger values will be converted to real if necessary. ormal value assignment 1. The stack contains less than two items. = B+C PUSH A	

Instruction:	FALSE	'K'		Load Return(False)
Effect:	The current block returns fa	ılse.		
Notes:	%false			
	This is usually accomplished by setting the true condition-code appropriately.			
Error:	The current block is not a predicate.			
Example:	<u>false if</u> Flag = 0 PUSH Flag; PUSHI 0; COMPARE			
	JNE 42			
	FALSE			
			LABEL 42	

Instruction:	JAM	ʻj'	Load Assign(2)	
Effect:	The value of TOS is assigned to the data item referenced by SOS. Both TOS and SOS are removed from the stack.			
Notes:	JAM transfer Truncate/Jam the source to fit into the destination Excess byte data will be omitted.			
Error:	The stack contains less than two items.			
Example:	! A is a string(15) ! B is a string(31) A = B	PUSH A PUSH B JAM		

Instruction:	MAP	'M'		Load Return(Map)	
Effect:	The address of the object re	The address of the object referenced by TOS is returned as the result of the map.			
Notes:	MAP result				
Error:	1. The current block is not a map.				
	2. The stack is empty.				
	3. TOS is not a reference to a data object.				
Example:	<u>result</u> == X		PUSH X; MAP		

Instruction:	RESULT	'V'		Load Return (Fn)
Effect:	TOS is removed from the sta	ack and return	ed as the resul	t of the function defined by
	the current block.			
Notes:	FN result			
	Integer values will be conve	rted to real if r	necessary.	
Error:	1. The stack is empty.			
	2. The current block is not a function.			
Example:	<u>result</u> = "Hello"	ult = "Hello" PUSHS "Hello"		
			RESULT	

Instruction:	RETURN	'R'		Load Return(Routine)
Effect:	A return sequence is genera	ated to return o	control from th	e current block.
Notes:	RETURN			
Error:				
Example:	return if X # 0 PUSH X; PUSHI 0; COMPARE			HI 0; COMPARE
	JE 19			
	RETURN			
			LABEL 19	

Instruction:	TRUE	'T'		Load Return(True)
Effect:	The current block returns tr	ue.		
Notes:	%true			
	This is usually accomplished by setting the true condition-code appropriately.			-code appropriately.
Error:	The current block is not a predicate.			
Example:	<u>true</u> if Flag = 0 PUSH Flag; PUSHI 0; COMPARE			JSHI 0; COMPARE
			JNE 42	
		TRUE		
			LABEL 42	

5.4. Embedding Machine Code

These ICODE instructions allow a programmer to insert assembler code or binary values into the program's instruction sequence. Multiple PLANT instructions can be used to insert legal machine code instructions that the embedded assembler function (MCODE) does not yet recognise.

Instruction:	MCODE <string></string>	'w'		Machine Code(Get Ascii(';'))
Effect:	The string is decoded as assembler text and added to the instruction stream			
Notes:	escape to assembler code			
	Assembler code is encoded by the IMP Compiler Pass1			
Error:				
Example:	*MOV_ %eax,#59		MCODE	

Instruction:	PLANT <byte></byte>	'P'		Load Plant
Effect:	Stores the byte in the instruction stream			
Notes:	Machine code literal			
Error:				
Example:	*45		PLANT 45	

5.5. Event/Signal

The IMP language has a simple mechanism to signal events/exceptions and insert handler code which the following ICODE instructions direct the code-generator to cater for.

Instruction:	EVENT <n></n>	'e'	Signal Event(Tag)	
Effect:	The event <n>,SOS,TOS is si</n>	The event <n>,SOS,TOS is signalled.</n>		
Notes:	%signal event			
Error:	The stack contains less than two items.			
	2. TOS is not an integer value.			
	3. SOS is not an integer value.			
Example:	<u>signal</u> 1,2,3	PUSHI 3; PUSHI 2; EVENT 1		

Instruction:	ON <n> <label></label></n>	ʻoʻ		EventTrap(TagC, Tag)
Effect:	This instruction marks the start of an on event block. <n> is a sixteen-bit set of flags where each trapped event is represented by a 1-bit, with the least-significant bit corresponding to event 0 and the most-significant bit event 15.</n>			
Notes:	%on %event block <a href="</td">			
Error:	 <n> does not have any bits set.</n> <label> is not defined before the end of the current block</label> 			
Example:	on 9 start; return; finish		ON 512 17 RETURN LABEL 17	

Instruction:	STOP	's'	Load Perm(stop, 0)
Effect:	Generates code to stop the target program (not the compiler).		
Notes:	%stop		
	This is equivalent to %signal 0,0,0		
Error:			
Example:			

5.6. Compare/Test

These ICODE instructions are used to indicate code required to compare values.

Instruction:	COMPARE	'?'	Load Compare Values	
Effect:	SOS is compared to TOS and	the condition-code is set ap	propriately. TOS and SOS	
	are then removed from the	stack.		
Notes:	The types of both TOS and S	SOS must be one of the follow	ving:	
	Integer, Real, String, Recor	d, Set		
	Integer values will be converted to real if necessary.			
	The comparison is signed w	d where integer values are concerned.		
Error:	1. The stack contains less than two items.			
	The types of TOS ar	2. The types of TOS and SOS are incompatible.		
Example:	<u>if</u> S < "123" <u>then</u> X = 0	PUSH S; PUSHS "123"; COM	PARE; BGE 13	
		PUSH X; PUSHI 0; ASSVAL		
		LABEL 13		

Instruction:	COMPAREA	'C'	Lo	oad Compare Addresses
Effect:	The address of SOS is comp	ared to the address	of TOS and t	he condition-code is set
	appropriately. TOS and SOS	are then removed	rom the stac	ck.
Notes:				
Error:	The stack contains less than two items.			
	2. TOS is not a reference for a data object.			
	SOS is not a referen	ce for a data object		
Example:	integername M; integer N			
	<u>if</u> N == M <u>then</u> Newline PUSH N; PUSH M; COMPAREA; BNE 14			
		PUSH Newline; CALL		
		LAB	EL 14	

Instruction:	COMPARED	(III)	Load Compare Double	
Effect:	SOS is compared to TOS and	d the condition-code is set ap	propriately. TOS and SOS	
	are then removed from the	stack.		
Notes:	The types of both TOS and S	SOS must be one of the follow	ving:	
	Integer, Real, String, Record	d, Set		
	Integer values will be conve	erted to real if necessary.		
	The comparison is signed where integer values are concerned.			
Error:	1. The stack contains less than two items.			
	2. The types of TOS and SOS are incompatible.			
Example:	<u>if</u> S < "123" <u>then</u> X = 0	PUSH S; PUSHS "123"; COM	PARED; BGE 13	
		PUSH X; PUSHI 0; ASSVAL		
		LABEL 13		

5.7. Parameter/Variable Declaration

These ICODE instructions help specify the type of variables present inside the program source.

Although IMP allows spaces within a variable name, internally the spaces are removed from variable names by the first stage of the compiler when generating the ICODE sequence.

The ICODE instruction, DEF, indicates the data-type, size and structure of a variable. When referencing the variable, the tag value is used.

Instruction:	ALIAS <string></string>	'G'		Get Alias Value(ReadString)
Effect:	<string> is noted as the curr</string>	ent alias.		
Notes:	See 'Begin' and 'Define'			
Error:	None			
Example:	external integer Thing alias	"SS\$THING"	ALIAS "SS	\$THING"
			DEF THIN	G

Instruction:	ALT 	'~' AlternateFormat(ReadByte)								
Effect:	This instruction marks an alternative sequence of tag definitions.									
Notes:	The instructions ALT-start and ALF-end are brackets and must be properly nested.									
	 = 'A' =>ALT-start									
	= 'C' =>ALT-next									
	 = 'B' =>ALT-end									
Error:	1. List flag is not set.									
Example:	recordformat F(integer X, (integer Y or real Z))									
	DEF F									
	START									
	DEF X									
	ALT-start									
	DEF Y									
	ALT-next									
	DEF Z									
	ALT-end									
	FINISH									

	•	•						
Instruction:	BOUNDS	ʻb'	Load Constant Bounds					
Effect:	The value of TOS is noted as 'upper-bound' and the value of SOS is noted as 'lower-							
	bound'. TOS and SOS are removed from the stack.							
Notes:	This is used to define a constant bounded Dope Vector							
	This instruction is used as a preliminary to defining switch vectors and own, const							
	and <u>external</u> arrays.							
Error:	1. The stack contains less than two items.							
	2. TOS is not an integer value.							
	3. SOS is not an integer value.							
	4. The value of TOS is less than the value of SOS.							
Example:	switch Sw(-3:3)	PUSHI 3; Negate; PUSHI 3; BOUNDS						
		DEF Sw						

Instruction:	DEF <tag> [id] <a> <c> '\$' Define Var(Tag, AsciiC, TagC, TagC, Tag)</c></tag>						
Effect:	A new tag value is created.						
	<tag> defines the tag index which will be used to select the value.</tag>						
	[id] specifies the actual identifier associated with the described object. It is a						
	sequence of zero or more characters terminated by a comma. This identifier will be used for external linkage if necessary unless overridden by an Alias. [id] will also be						
	used for run-time diagnostic information.						
	asea for fair time diagnostic information.						
	<a> A two-byte value: <a> = T<<4+F where:						
	T = 0 : void						
	T = 1 : integer {qualified by }						
	T = 2 : real {qualified by }						
	T = 3 : string {maximum length }						
	T = 4 : record {format }						
	T = 5 : boolean T = 6 : set						
	T = 7 : 8-bit-enumerated {format }						
	T = 8 : 16-bit-enumerated {format }						
	T = 9 : pointer						
	T = 10 : char						
	= 11 : unsigned {qualified by }						
	T = 12-15 : undefined {error}						
	F = 0 : void						
	F = 1 : simple {byte}						
	F = 2 : indirect {bytename}						
	F = 3 : general label						
	F = 4 : recordformat						
	F = 5 : undefined {error}						
	F = 6: switch						
	F = 7 : routine						
	F = 8 : function						
	F = 9 : map F = 10 : predicate						
	F = 11 : array {array}						
	F = 12 : array indirect {arrayname}						
	F = 13 : indirect array {namearray}						
	F = 14 : indirect array indirect {namearrayname}						
	F = 15 : undefined {error}						
	 If T is INTEGER takes the following meanings b=:						
	1, full range (depends on machine i.e. 32-bit or 64-bit)						
	2, range 0 255 (=02^8-1)						
	3, range -32768 32767 (=-2^152^15-1)						
	4, range -2^312^31-1						
	5, range -2^632^63-1						

If T is UNSIGNED takes the following meanings b=:

- 1, full range (depends on machine i.e. 32-bit or 64-bit)
- 2, range 0..2^8-1
- 3, range 0..2¹⁶⁻¹
- 4, range 0..2³²⁻¹
- 5, range 0..2^64-1

If T is REAL takes the following meanings b=:

- 1, normal precision
- 4, double precision

If T is STRING gives the maximum length of the string.

If T is RECORD gives the tag of the corresponding recordformat.

If T is enumerated gives the tag of the dummy format used to identify the enumerated value identifiers.

<c> is a two-byte value: U<<5+I<<4+S<<3+X where:

U = 1 check the object for unassigned, U = 0 otherwise

I = 1 if the object is an indirect object, I = 0 otherwise

S = 1 if this is a <u>spec</u>, S = 0 otherwise

X = 0 :: automatic (stack) allocation

X = 1 :: own

X = 2 :: <u>constant</u>

X = 3 :: external

X = 4 :: system

X = 5 :: dynamic

X = 6 :: primitive

X = 7 :: permanent

An indirect object (I=1) differs from F=2 in that F=2 implies that the actual object created will be a pointer and will be dereferenced whenever used unless explicit action is taken (e.g. use of Assign-Reference).

If I=1 a pointer will be created (usually as an integer) and will be treated as an integer (or address) with no automatic dereferencing taking place.

Notes:

The tag values within a block should be dense and preferably consecutive.

All tag values within a block must have values greater than the maximum tag value yet defined within the enclosing block.

Tag definitions remain valid until the end of the enclosing block.

The tag values used within a recordformat definition must all be zero; the fields of a record are selected by their position in the format, numbered starting from one.

	x = i	x = illegal combination of T,F values															
		F															
			0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
		0		Х			Х	Х			Х			Χ		Х	
		1	Х			Х	Х	Х	Χ	Χ			Χ				
		2	Х			Х	Х	Х	Χ	Χ			Χ				
		3	Х			Х	Х	Х	Χ	Χ			Χ				
		4	Х			Х		Х	Χ	Х			Х				
	T	5	Х			Х	Х	Х	Χ	Х			Х				
		6	Х			Х	Х	Х	Х	Х			Х				
		7	Х			Х	Х	Х	Χ	Χ			Χ				
		8	Х			Х	Х	Х	Χ	Χ			Χ				
		9	Χ			Х	Χ	Х	Χ	Х			Χ				
		10	Х			Х	Х	Х	Х	Х			Х				
Error:																	
Example:																	

Instruction:	DIM <n><d></d></n>	'd'	Dimension(TagC, Tag)						
Effect:	<d>pairs of integer values on the stack are used to define the bounds of the last</d>								
	<n> arrays to have been defined. Code is generated, if necessary, to allocate the</n>								
	arrays and the definitions are adjusted to reference the appropriate storage.								
Notes:	The pairs of values are stacked in order of the declaration, that is, first dimension								
	first.								
	In each pair of values the lower bound is stacked before the upper bound.								
	The last <n> tags must have had consecutive index values.</n>								
	NB in params: d =0 -> simple array, # 0 -> array-in-record								
Error:	1. The stack contains less than 2* <d> items.</d>								
	2. The last <n> definitions were not all arrays.</n>								
Example:	integerarray A, B, C(1:2, Low:	4)	DEF A						
			DEF B						
	DEF C								
	PUSHI 1; PUSHI 2								
	PUSH Low; PUSHI 4								
	DIM 3 2								

Instruction:	FINISH	<i>'</i> }'		Finish Params	
Effect:	This instruction marks the e		_		
	parameter list or a recordfo		n. List flag is cle	eared and the tag list is	
	processed in any ways nece	•			
	If the tag list is associated w	•	•	cordformat this instruction	
	also marks the 'end' of the a		ck'.		
Notes:	The list of definitions may b	e empty.			
Error:	1. List flag is clear.				
	2. There has been an unmatched Alt-Start.				
Example:	routine Test(integer j,k)	DEF Test			
		START			
			DEF j		
			DEF k		
			FINISH		
	recordformat F(integer P or	r real R) DEF F			
		START			
		DEF P			
		ALT-next			
			DEF R		
			FINISH		

Instruction:	INIT <n></n>	'A'		Init(Tag)
Effect:	<n> copies of the init-value are added to the list of values associated with the init-variable. The init-value is either the default value (unassigned) if the stack is empty, the value of TOS (possibly converted to real) if TOS is a constant, or the address of TOS if TOS is a variable. The init-variable is the last static object to have been defined using Define.</n>			
Notes:	Used to initialise an OWN v	ariable		
Error:	1. TOS, if it exists, is not of the same type as the init-variable.			
Example:	ownintegerarray A(1:5) = 1(3), 4, 99	PUSHI 1; PUS	HI 5; BOUNDS
			DEF A	
			PUSHI 1; PUS	HI 3
	PUSHI 4; INIT 1			
	PUSHI 99; INIT 1			
	owninteger P = -1 PUSHI 1; NEGATE			
			DEF P	
			INIT 1	

Instruction:	SETFORMAT <tag></tag>	'A'	Set Record Format(Tag)
Effect:	TOS is converted to be a red	cord of format <tag>. This nev</tag>	ver involves any instructions
	being executed in the objec	t program.	
Notes:			
Error:	1. <tag> is not a record</tag>	dformat.	
	2. The stack is empty.		
	TOS is not a referen	ice to a variable.	
Example:			

Instruction:	START	' {'	StartParams	
Effect:	This instruction marks the s procedure or the componer		tags defining the parameters to a . List flag is set.	
Notes:	This instruction must always follow the definition of a procedure or recordformat tag. There must be a matching 'Finish' before the end of the block.			
Error:	 List flag is set. The last instruction was not a 'Define' which introduced a procedure or recordformat. 			
Example:	routine Newline;; end		DEF Newline START FINISH END	

5.8. Access Variables

These ICODE instructions are used to access the variables defined in the previous set of ICODE instructions.

Instruction:	ACCESS	'a'	Load Array Ref(0)	
Effect:	TOS is used as the final inde	x into SOS. TOS is removed fr	rom the stack leaving SOS	
	as the new TOS.			
Notes:	This instruction is used to p	rocess the final dimension of	an N dimensional array.	
	See Index for the previous dimensions.			
Error:	1. The stack contains less than two items.			
	TOS is not an intege	er value.		
	3. SOS does not describe an array.			
Example:	A(j) = 0 PUSH A			
	PUSH j; ACCESS			
	PUSHI 0; ASSVAL			

Instruction:	INDEX	ʻi'	Load Array Ref(1)	
Effect:	TOS is used as the next index into SOS. TOS is removed from the stack leaving SOS as the new TOS.			
Notes:	This instruction is used to process the first N-1 dimensions of an N dimensional array. See Access for the final dimension.			
Error:	 The stack contains less than two items. TOS is not an integer value. SOS is not an array descriptor. 			
Example:				

Instruction:	PUSH <tag></tag>	'@'	Load Var(Tag)
Effect:	The tag with index value <tag> is pushed onto the stack.</tag>		
Notes:	Load variable descriptor (calls Stack Var)		
Error:	1. <tag> has not been defined.</tag>		
Example:	A = B	PUSH A; PUSH B; ASSVAL	

Instruction:	PUSHI <integer></integer>	'N'	Load Const(ReadInteger)	
Effect:	The integer value <integer></integer>	is pushed into the stack and	becomes the new TOS.	
Notes:	Load integer constant (calls Push Const) The instruction 'Byte' is an abbreviation for this instruction when the value is in the range 0255.			
Error:				
Example:	X = 2500	PUSH X; PUSHI 2500; ASSVA	L	

Instruction:	PUSHR <real></real>	'D'		Input Real Value(ReadReal)	
Effect:	The real constant <real> is pushed onto the stack.</real>				
Notes:	Stack real constant				
Error:					
Example:	Z = -1.23 PUSH Z; PUSHR <1.23>; NEGATE				
			ASSVAL		

Instruction:	PUSHS <string></string>	m	Input String Value(ReadString)		
Effect:	The string constant <string></string>	is pushed onto the sta	ck.		
Notes:	Stack string constant	Stack string constant			
Error:					
Example:	S = "Hello"	PUSH S			
		PUSHS "Hello"			
		ASSVAL			

Instruction:	RESOLVE <flag></flag>	'r'	Resolve(Tag)
Effect:			
Notes:	Split the string S ->A.(B).C		
Error:			
Example:			

Instruction:	SELECT <n></n>	'n'		SelectField(Tag)	
Effect:	TOS is replaced by the <n>'t</n>	h item in the format	of TOS. Fi	ields within records are	
	numbered from 1; alternativ	ve markers have no	effect on t	the numbering.	
Notes:	Select member from record	format			
Error:	 The stack is empty. 				
	TOS is not a record.				
	The format of TOS of	loes not contain at le	east <n> it</n>	ems.	
Example:	recordformat F(integer P or	record (F)name Q)			
	record (F) R				
	$R_QP = 0$ PUSH R				
	SELECT 2				
	SELECT 1				
	PUSHI 0				
			ASSVAL		

5.9. Compiler Directives

These ICODE instructions are used to control the operation of the various compiler stages.

Instruction:	CONTROL <n></n>	ʻz'	Set Control(ReadTag)
Effect:	The value <n> is of the form p<<14+q. The value q is to be used by the p'th pass of</n>		
	the compiler in an implementation-specific manner.		
Notes:			
Error:			
Example:			

Instruction:	DIAG <n></n>	'y'	Set Diagnose(ReadTag)
Effect:	The value <n> is of the form p<<14+q. The value q is to be used by the p'th pass of</n>		
	the compiler in an implementation-specific manner.		
Notes:			
Error:			
Example:			

Instruction:	EOF	10	
Effect:	The compilation is to be abandoned with an error message.		
Notes:	This instruction is required at the end of an I-code file even though in correct programs it will never be executed. It is there to provide a check on the operation of the code-generator and to permit the code-generator to use a one-character look-ahead.		
Error:			
Example:			

Instruction:	LANG <flags></flags>	η,	Load Language Flags(Tag)
Effect:	The 16-bit parameter <flag< td=""><td></td></flag<>		
Notes:	Currently only support standard IMP		
Error:			
Example:			

Instruction:	LINE <n></n>	'O'	Update Line(Tag)
Effect:	The current position is asso	ciated with the start of the co	ode for source line <n>.</n>
Notes:			
Error:	 The stack is not em 	pty.	
Example:	X = 1		
	Y = 3; Z = 4		
	LINE 1; PUSH X; PUSHI 1; AS	SSVAL	
	LINE 2; PUSH Y; PUSHI 3; AS	SSVAL	
	LINE 2; PUSH Z; PUSHI 4; AS	SSVAL	

Instruction:	MONITOR	'm'	Load Monitor
Effect:	Plant code to execute the specific monitor action		
Notes:			
Error:			
Example:			

5.10. Block Structure/Call

These ICODE instructions indicate the block structure of the program, and calling sequence between the different blocks of code.

Instruction:	BEGIN	'H'		Load Compile Begin
Effect:	An anonymous procedure is defined here and called once. A new block is entered.			
	If an alias has been noted (s	see Alias) that s	tring will be us	sed for identifying the block
	in diagnostic information le	aving no alias r	noted.	
Notes:	The sequence "Begin End" may always be replaced by a sequence of the form:			y a sequence of the form:
	Define X; Start; Finish; End; Stack X; Call			
	where X is a suitable unique tag.			
Error:	None			
Example:	begin; Newline; end BEGIN			
			PUSH Newlin	e; CALL
			END	

Instruction:	CALL	'E'	Load Compile Call(top)	
Effect:	The procedure described by TOS is called. If TOS is a procedure which returns a			
	result TOS is replaced by the	at result, otherwise TOS is rei	moved from the stack.	
Notes:	Predicates do not return a r	esult but set the condition-co	ode.	
Error:	1. The stack is empty.			
	2. TOS does not describe a procedure.			
	3. There has not been the same number of parameters assigned using Assign-			
	Parameter as is specified by the parameter list.			
Example:	Newlines(4)	ewlines(4) PUSH Newlines		
		PUSHI 4; ASSPAR		
		CALL		

Instruction:	END	(,))	Load End of Block
Effect:	This instruction marks the end of a block. All tags defined within the block are deleted (made undefined) and become available for re-use. If the block is a routine this instruction also implies a 'return' instruction.		
Notes:			
Error:	The stack is not empty.		
Example:			

6. Unused ICODE Instructions

The current compiler implementation uses a limited set of ICODE instructions which are described in the previous section.

The following ICODE instructions are currently not implemented in this version of the IMP compiler suite.

Some of these unused ICODE instructions could be used to represent instances of calls to IMP runtime library routines or be replaced by embedded implicit code to be inserted by the code generation phase of the compiler.

Other unused ICODE instructions could be used to define constructs in the non-IMP compilers (i.e. C/Pascal) implemented by P.S. Robertson.

Instruction:	Absolute		
Effect:	TOS is replaced by the absolute value of TOS.		
Notes:			
Error:	3. The stack is empty.		
	4. TOS is not integer or real.		
Example:	X = Y+Z	Stack X	
		Stack Y; Stack Z; Add; Absolute	
		Assign-Value	

Instruction:	Address			
Effect:	TOS is replaced by the address of the object it describes.			
Notes:	Commonly the type of an address will be indistinguishable from integer.			
Error:	1. The stack is empty.			
	2. TOS does not have an address.			
Example:	P = Addr(Q)	Stack P		
		Stack Q; Address		
		Assign-Value		

Instruction:	Adjust			
Effect:	The address of SOS is adjusted forwards or backwards by TOS items of the same size as SOS. This may be thought of as an array accessing instruction where SOS defines the zero'th element.			
Notes:				
Error:	The stack contains less than two items.			
	2. TOS is not an integer.			
	3. SOS does not reference a data object.			
Example:	N == N[X]	Stack N		
		Stack N; Stack X; Adjust		
		Assign-Reference		

Instruction:	Byte 		
Effect:	The unsigned byte value is stacked.		
Notes:	This is a compact form for the Integer instruction when small values are to be		
	stacked.		
Error:	None		
Example:	X = 200	Stack X; Byte 200; Assign-Va	lue

Instruction:	Compare-Repeated-Values		
Effect:	SOS is compared to TOS and the condition-code is set appropriately. SOS is then removed from the stack leaving TOS		
Notes:	The types of both TOS and SOS must be one of the following:		
	Integer, Real, String, Reco	rd, Set	
	Integer values will be converted to real if one operand is real.		
Error:	 The stack contains 	less than two items.	
		and SOS are incompatible.	
Example:	$if 1 \le X \le 12 then X = 0$	Byte 1; Stack X; Compare-Repeated-Values; BGT 15	
		Byte 12; Compare-Values; BGT 15	
		Stack X; Byte 0; Assign-Values	
		Label 15	
Instruction:			
	Compare-Unsigned-Values		
Effect:		OS with the values being interpreted as unsigned values	
Nietes	The condition-code is set accordingly and TOS and SOS are removed from the stack.		
Notes: Error:	1. The steel contains less than two theres.		
EITOI:	1. The stack contains less than two items.		
	 TOS is not an integer value. SOS is not an integer value. 		
Example:	if U1 < U2 then U2 = 0	Stack U1; Stack U2; Compare-Unsigned-Values; Bge 31	
Example.	<u> </u>	Stack U2; Byte 0; Assign-Value	
		Label 31	
		2000101	
Instruction:	Complement		
Effect:	TOS is replaced by the ones-complement of TOS.		
Notes:			
Error:	1. The stack is empty	,	
	2. TOS is not an integer value.		

Instruction:	Define-Range <tag></tag>			
Effect:	The given tag is defined to be the integer range defined by lower-bound and upper- bound.			
Notes:				
Error:	The two bounds have not been defined.			
Example:	Var x:110; x := i; Byte 1; Byte 10; Bounds			
		Define-Range 99; Define x		
		Stack x; Stack i; Test-Range 99; Assign-Value		

Instruction:	Duplicate		
Effect:	A copy of TOS is pushed onto the stack.		
Notes:	After this instruction TOS and SOS are identical.		
Error:	1. The stack is empty.		
Example:	int A[10],x; A[x]++; Stack A; Stack x; Adjust Duplicate		
		Byte 1; Add; Assign-Value	

Instruction:	Eval			
Effect:	The value described by TOS is protected against being altered as the side-effect of			
	alterations to any variables	which make up that value.		
Notes:	Commonly this instruction	loads the value of TOS into a machine register.		
Error:	 The stack is empty. 			
Example:	a = b + c++	Stack a; Stack b; Stack c; Add; Eval		
		Stack c; Stack c; Byte 1; Add; Assign-Value		
	Assign-Value			
	T			
Instruction:	Eval-Addr			
Effect:	•	escribed by TOS is protected against alteration.		
Notes:	-	loads the address of the object referred to be TOS into a		
	machine register.			
Error:	1. The stack is empty.			
	2. TOS is not a referer			
Example:	int *p; *p++ = 10;	Stack p; Eval-Addr; Stack p; Duplicate		
		Byte 1; Adjust; Assign-Value		
		Byte 10; Assign-Value		
la star sti sa		T T		
Instruction:	Float			
Effect:	The value described by TOS is converted into a real value if necessary.			
Notes:	If TOS is already a real this operation is a no-op.			
Error:	1. The stack is empty.			
F I .	2. TOS is neither integ	ger nor real.		
Example:				
Instruction:	In alcodo catain as	T T		
Effect:	Include <string></string>	start (or and) of the code generated from source		
Ellect:	This instruction marks the start (or end) of the code generated from source contained in an include file.			
	If <string> is null it marks t</string>			
Notes:	ii sting is nun temans e	ne end of an include me.		
Error:				
Example:				
zxampiei				
Instruction:	Init-Type <n></n>			
Effect:		e (see Init) is <n> for the purposes of subsequent 'Init'</n>		
	1	he type as in the type field of 'Define'.		
Notes:		71		
Error:	1. <n> does not corre</n>	spond to a valid type.		
Example:				
	1			
Instruction:	Int			
Effect:		, that is, the nearest integer to the value of TOS. The		
	type of the new TOS is integer.			
Notes:		ion for a discussion of the details of the INT function.		
Error:	The stack is empty.			
	2. TOS is neither integ			
Example:	I = Int(R+0.3)	Stack I; Stack R; Real <0.3>; Add		
		Int		
		Assign-Value		

	T		
Instruction:	Intpt		
Effect:	TOS is replaced by Intpt(TOS), that is, the integer part of TOS. The type of the new TOS is integer.		
Notes:	See the IMP Library Definit	ion for a discussion of the details of the INTPT function.	
Error:	 The stack is empty. 		
	2. TOS is neither integ	ger nor real value.	
Example:	I = Intpt(R-S)	Stack I; Stack R; Stack S; Sub Intpt Assign-Value	
Instruction:	Localise		
Effect:	The area pointed at by <a> point at the new area.	is copied into the local stack frame and <a> is updated to	
Notes:	If the first byte of the new	area is at X, <a> is updated to the address X-<c>.</c>	
Error:	1. The stack contains		
	2. <a> is not a referer	nce.	
	3. and <c> are no</c>	t integer values.	
Example:			
	T		
Instruction:	Null-Set		
Effect:	A descriptor of a null set is	stacked.	
Notes:			
Error:			
Example:	SetA := [];	Stack SetA; Null-Set; Assign-Value	
	T		
Instruction:	Pop		
Effect:	TOS is removed from the st	ack.	
Notes:			
Error:	 The stack is empty. 		
Example:	X := Y := 0	Stack X; Duplicate Stack Y; Duplicate Byte 0; Assign-Value; Assign-Value Pop	
	I ~ .	,	
Instruction:	Reference <n></n>		
Effect:	TOS is replaced by a reference to an object of type <n> at the address given by the original TOS. The value of <n> is encoded in the same way as the type information, <a>, in the Define instruction.</n></n>		
Notes:			
Error:	 The stack is empty. TOS is not an integ. 		
Example:	P = Integer(Q)	Stack P; Stack Q; Reference 1	
		Assign-Value	

Instruction:	Remainder			
Effect:	TOS and SOS are removed from the stack and are replaced by the value REM(SOS,			
	TOS). The exact definition of	of REM is given in the IMP Library Definition.		
Notes:	·			
Error:	The stack contains less than two items.			
	2. TOS is not an integer value.			
	3. SOS is not an integer value.			
Example:	Digit = Rem(N, 10)	Stack Digit; Stack N; Byte 10; Remainder		
		Assign-Value		

Instruction:	Round		
Effect:	TOS is replaced by the value ROUND(TOS) where ROUND is as defined in section		
	6.6.6.3 of the Pascal standa	rd BS 6192:1982, with the ext	tension that ROUND
	returns the value of its para	ameter if that value is already	an integer.
Notes:			
Error:	1. The stack is empty.		
	2. TOS is neither integer nor real.		
Example:	i := Round(r+0.1);	Stack i; Stack r; Real <0.1>; A	Add
		Round; Assign-Value	

Instruction:	Size-Of		
Effect:	TOS is replaced by the integer value giving the number of bytes in the object referenced by the original TOS.		
Notes:	Tereford at the original root		
Error:	 The stack is empty. TOS is not a reference to a data object. 		
Example:	P = Sizeof(R)	Stack P; Stack R; Size-Of; Assign-Value	

Instruction:	Stack-Condition 			
Effect:	The values of SOS and TOS are compared as in Compare-Values but instead of the			
	condition-code being set, T	OS and SOS are replaced by t	he constant 1 (true) or 0	
	(false) depending on wheth	er the condition specified by	 is true or false. The	
	values of are the encod	ings of the instructions:		
	BEQ, BNE, BLT, BLE, BGT, B	GE, BT, BF		
Notes:	The comparison is signed when integers are concerned.			
Error:	The stack contains less than two items.			
	2. TOS and SOS cannot be compared.			
	3. is not a valid condition.			
Example:	B := (X=Y);	B := (X=Y); Stack B		
	Stack X; Stack Y; Stack-Condition BEQ			
	Assign-Value			

Instruction:	Stack-In			
Effect:	TOS and SOS are removed from the stack and are replaced by an integer value			
	which is 1 (true) if SOS is IN	the set TOS or 0 (false) if SOS	S is not IN the set TOS.	
Notes:				
Error:	1. The stack contains less than two items.			
	2. TOS is not a set.			
	3. SOS is not an integer.			
Example:	B := (x IN s); Stack B			
	Stack x; Stack s; Stack-In			
	Assign-Value			

Instruction:	Stack-Unsigned-Condition 			
Effect:	The values of SOS and TOS are compared as in Compare-Unsigned-Values but			
	instead of the condition-code b	eing set TOS and SOS are	replaced by the constant 1	
	(true) or 0 (false) depending on whether the condition specified by is true or			
	false. The values of are the	encodings of the instruct	ions:	
	BEQ, BNE, BLT, BLE, BGT, BGE, BT, BF			
Notes:	The comparison is unsigned when integers are concerned.			
Error:	The stack contains less than two items.			
	2. TOS and SOS cannot be compared.			
	3. is not a valid condition.			
Example:	B := (Ux < Uy);	B := (Ux < Uy); Stack B		
		Stack Ux; Stack Uy; Stack-Unsigned-Condition BLT		
	Assign-Value			

Instruction:	Swop		
Effect:	The top two items on the stack are reversed. That is, TOS becomes the new SOS, and SOS becomes the new TOS.		
Notes:			
Error:	1. The stack contains less than two items.		
Example:	X = Y	Stack Y; Stack X; Swop; Assign-Value	

Instruction:	Test-Boolean			
Effect:	The condition-code is set to 'false' or 'true' depending on whether TOS is 0 (false) or			
	1 (true). TOS is removed from the stack.			
Notes:				
Error:	1. The stack is empty.			
	2. TOS is not a boolean value.			
Example:	If B Then DoltNow;	Stack B; Test-Boolean; BF 43		
		Stack DoltNow; Call		
		Label 43		

Instruction:	Test-In		
Effect:	The value of SOS is tested for inclusion within the set TOS. The condition-code is set		
	'true' or 'false' accordingly. Both TOS and SOS are removed from the stack.		
Notes:			
Error:	The stack contains less than two items.		
	2. SOS is not an integer value.		
	3. TOS is not a set value.		
Example:	If NOT x IN s Then x := 0;	Stack x; Stack s; Test-In	
		Bt 31	
		Stack x; Byte 0; Assign-Value	ue
		Label 31	

Instruction:	Test-Nil		
Effect:	A check is performed to ensure that TOS is not NIL. An event is signalled if it is, or if		
	TOS points to a heap item which has been returned to the heap (disposed).		
Notes:	This test can also perform an unassigned variable check.		
Error:	1. The stack is empty.		
	2. TOS is not a pointer variable.		
Example:	P^ := 0;	Stack P; Test-Nil	
		Reference <1>	
		Byte 0; Assign-Value	

Instruction:	Test-Range <tag></tag>		
Effect:	The value of TOS is checked to be within the range defined by the tag. If the value is		
	not in the range an event is signalled.		
Notes:	The event is signalled at run-time.		
Error:	1. The stack is empty.		
	2. TOS is not an integer value.		
	3. <tag> does not define a range.</tag>		
Example:	Byteval := Bigval;	Stack Byteval	
	Stack Bigval; Test-Range Byterange		yterange
		Assign-Value	

Instruction:	Trunc		
Effect:	TOS is replaced by the value TRUNC(TOS) where TRUNC is as defined in section		
	6.6.6.3 of the Pascal standard BS 6192:1982, with the extension that Trunc returns		
	the value of its parameter if that value is already an integer.		
Notes:			
Error:	1. The stack is empty.		
	2. TOS is neither integer nor real.		
Example:	i := Trunc(r+0.1);	Stack i; Stack r; Real <0.1>; A	Add
		Trunc; Assign-Value	

Instruction:	Variable-Call			
Effect:	The procedure describe	procedure described by TOS is called. This differs from 'Call' in that the		
	procedure may have a v	procedure may have a variable number of parameters.		
Notes:	This instruction is used t	ed to call 'C' procedures and its definition is as woolly as the		
	definition of that langua	age.		
Error:	 The stack is emp 	1. The stack is empty.		
	TOS is not a pro	rocedure descriptor.		
Example:	try(1); try(1,2);	Stack try		
		Byte 1; Assign-Parameter		
		Variable-Call		
Stack try;				
		Byte 1; Assign-Parameter		
	Byte 2; Assign-Parameter			
		Variable-Call		

Appendix 1. Instructions which set the condition code Call {predicate} Compare-Values Compare-Unsigned-Values Compare-References Compare-Repeated-Values Test-Boolean Test-In Appendix 2. Instructions which test the condition code JEQ JΖ JGE JGT JLE JLT JNE JNZ