

# Generators of $H^1(\Gamma, \partial\Gamma^c)$ with $\partial\Gamma^c \subset \partial\Gamma$ for Triangulated Surfaces $\Gamma$ : Construction and Classification of Global Loops

Silvano Pitassi

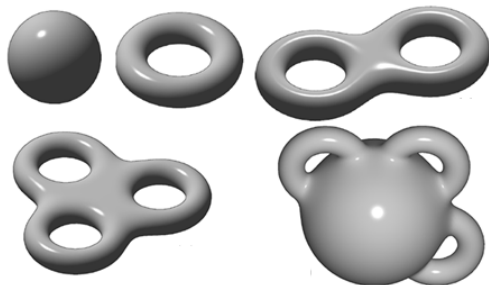
IMAG, University of Montpellier, France

# Introduction

- 1 Introduction
- 2 Key Mathematical Concepts
- 3 Review of past work: Hiptmair-Ostrowski algorithm
- 4 Main algorithm overview
- 5 Conclusion and implications

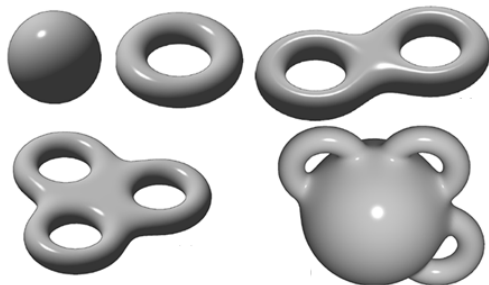
# Motivation: Electromagnetic scattering

Let  $\Gamma$  be a triangulated embedded **surface** (i.e. a closed 2-manifold)



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- **Electrical Field Integral Equation (EFIE):** find the induced current  $J$  on  $\Gamma$  from an *incident* (source) electromagnetic field  $E^i$

$$\gamma_T \circ \mathcal{U}_k(J) = -\gamma_T E^i$$

# Motivation: Low-frequency breakdown

- **Problem:** poor conditioning of BEM matrices and numerical instability as  $\omega \rightarrow 0$

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- **Problem:** poor conditioning of BEM matrices and numerical instability as  $\omega \rightarrow 0$
- **Loop-Star decomposition:** decompose the current  $\mathbf{J}$  as

$$\mathbf{J} = \underbrace{\mathbf{J}_{\text{loop}}}_{\text{Solenoidal}} + \underbrace{\mathbf{J}_{\text{star}}}_{\text{Non-solenoidal}}$$

and introduce the frequency-dependent scalings

$$\mathbf{J}_{\text{loop}} = \tilde{\mathbf{J}}_{\text{loop}} \quad \text{and} \quad \mathbf{J}_{\text{star}} = \omega \tilde{\mathbf{J}}_{\text{star}}$$

# Why Cohomology?

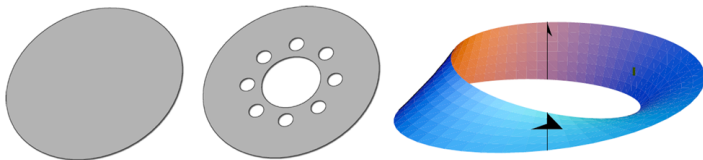
$J_{\text{loop}}$  and  $J_{\text{star}}$  are determined as

$$J = \underbrace{\text{curl}_{\Gamma} \psi + \mathbf{g}}_{=: J_{\text{loop}}} + \underbrace{\text{grad}_{\Gamma} \varphi}_{=: J_{\text{Star}}},$$

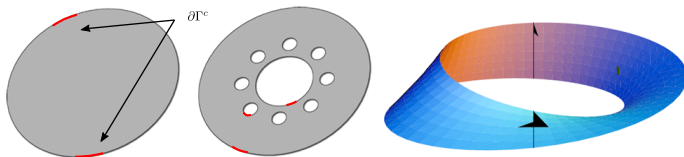
with  $\mathbf{g} \in H^1(\Gamma)$  a **cohomology generator (global loop)**

# The general problem

- Surface with boundary



- Presence of “contacts”  $\partial\Gamma^c$



We need a general algorithm to compute generators of the  
**first relative cohomology group**  $H^1(\Gamma, \partial\Gamma^c)$



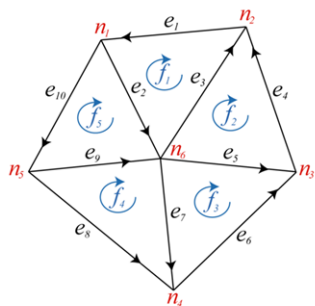
# Key Mathematical Concepts

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- 2 Key Mathematical Concepts**
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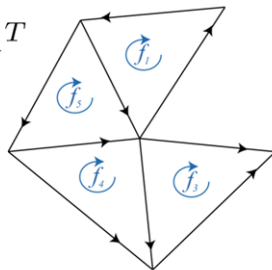
# Chains

Formal **linear combinations** of  $k$ -cells with coefficients in an commutative ring  $G$  ( $\mathbb{Z}$  or  $\mathbb{R}$ )

Example:



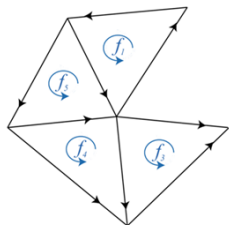
$$w = \{1, 0, 1, 1, 1\}^T$$



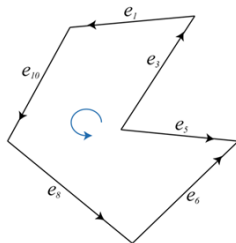
Chains can be added and the  $k$ -cells are a basis for the **chain group**  $C_k(\Gamma, G)$

# Boundary operator

Boundary operator:  $\partial_k : C_k(\Gamma; \mathbb{Z}) \rightarrow C_{k-1}(\Gamma; \mathbb{Z})$



$$W = \{1, 0, 1, 1, 1\}^T$$



$$\partial_2 = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & -1 \\ 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix} \begin{matrix} \text{Number of faces} \\ \text{Number of edges} \end{matrix}$$

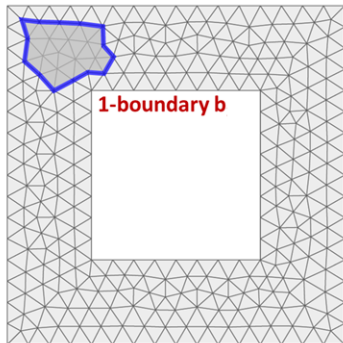
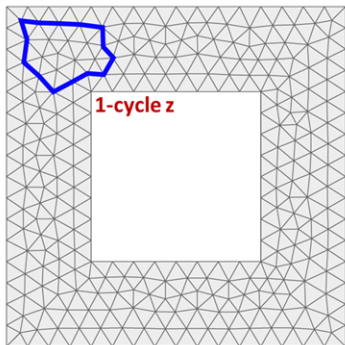
$$Z = \partial_2 W = \{1, 0, 1, 0, -1, 1, 0, 1, 0, 1\}^T$$

# Cycles and boundaries

Boundary operator  $\partial_k$  gives rise to a classification of  $k$ -chains

**Example:**  $k = 1$

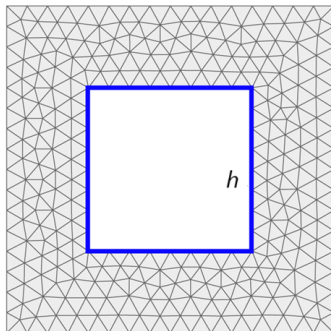
- 1-chains  $z$  whose **boundary is zero**  $\partial_1 z = 0$  are 1-cycles in  $Z_1(\Gamma)$
- 1-chains  $b$  **that are boundary** of a 2-chain  $\partial_2 c = b$  are 1-boundaries in  $B_1(\Gamma)$



# First homology group

All 1-boundaries are 1-cycles, but the converse is not true

**Example:** hollow disk



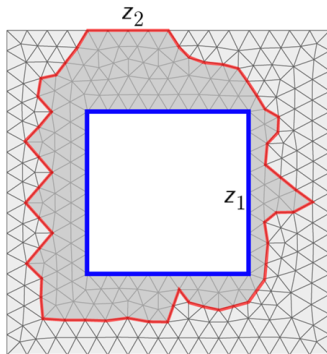
The **first homology group** is defined as the **quotient group**

$$H_1(\Gamma) := Z_1(\Gamma)/B_1(\Gamma)$$

# Homology class

**Equivalence relation:** two 1-cycles  $z_1$  and  $z_2$  are **homologous** if their difference is a 1-boundary:  $z_1 \sim z_2 \iff z_1 - z_2 \in B_1(\Gamma)$

**Example:**



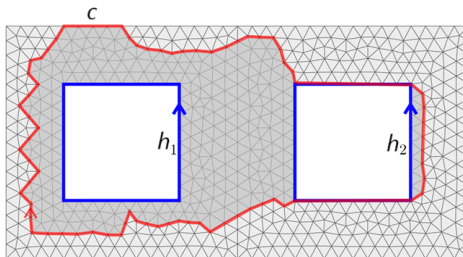
$z_1$  and  $z_2$  are in the same **homology class**  $[z_1]$ . Each homology class is represented by any of its elements (called **representative**)

# Homology generators and homology basis

**Homology generators:** a minimal set of homology classes  $[h_i]$  such that

$$c = \sum_{i=1}^{\beta_1(\Gamma)} a_i h_i + b$$

Example:



A  $k$ -cochain  $\mathbf{c}$  is a **linear functional**  $\mathbf{c} : C_k(\Gamma; G) \rightarrow G$

Cochains can be added in the natural way and form the **cochain group**  $C^k(\Gamma; G)$  (**dual space** of  $C_k(\Gamma; G)$ )



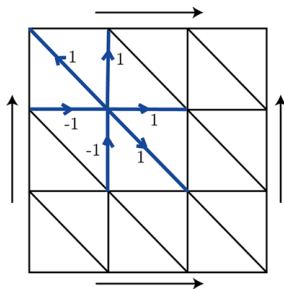
# Coboundary, cocycles and coboundaries

The coboundary operator  $\delta^k : C^k(\Gamma; G) \rightarrow C^{k+1}(\Gamma; G)$  is designed in such a way that the **Generalized Stokes theorem** holds:

$$\langle \delta^k \mathbf{c}, w \rangle = \langle \mathbf{c}, \partial_{k+1} w \rangle$$

$\delta^k$  gives rise to a classification of 1-cochains:

- 1-cochains  $\mathbf{c}$  whose **coboundary is zero**  $\delta^1 \mathbf{c} = \mathbf{0}$  are 1-cocycles in  $Z^1(\Gamma)$



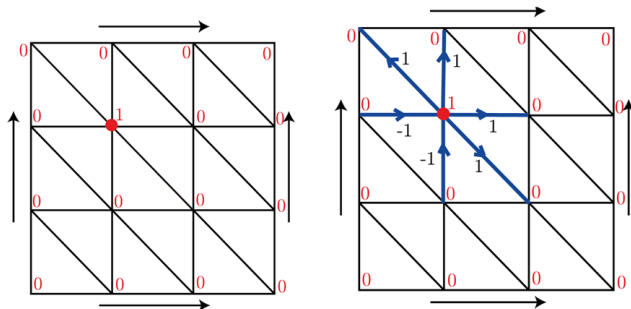
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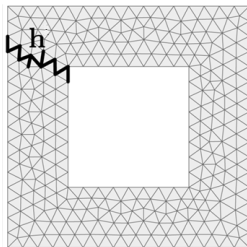
- 2 1-cochains **that are coboundary** of a 0-cochain are 1-coboundaries in  $B^1(\Gamma)$



# Cohomology group

All  $k$ -coboundaries are  $k$ -cocycles, but the converse is not true (as not all irrotational fields are gradients when the domain is not simply connected)

Example:



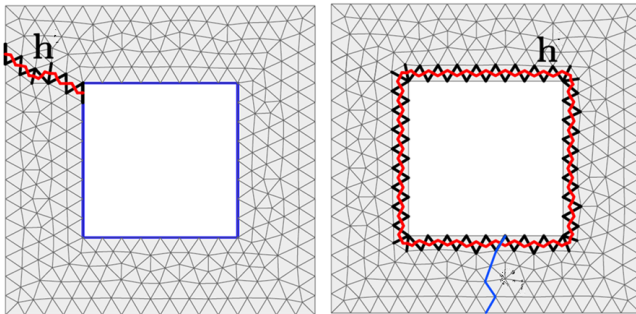
The **first cohomology group** is defined as the following **quotient group**

$$H^1(\Gamma) := Z^1(\Gamma)/B^1(\Gamma)$$

# Relative vs absolute cohomology

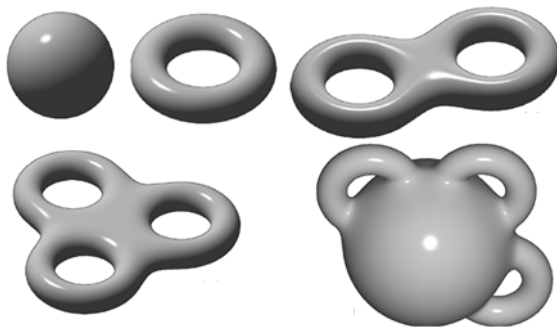
Relative cohomology takes into account **boundary conditions**: the current  $\mathbf{J}$  cannot flow outside  $\Gamma$

**Example:**  $\mathbf{h} \in H^1(\Gamma, \partial\Gamma)$



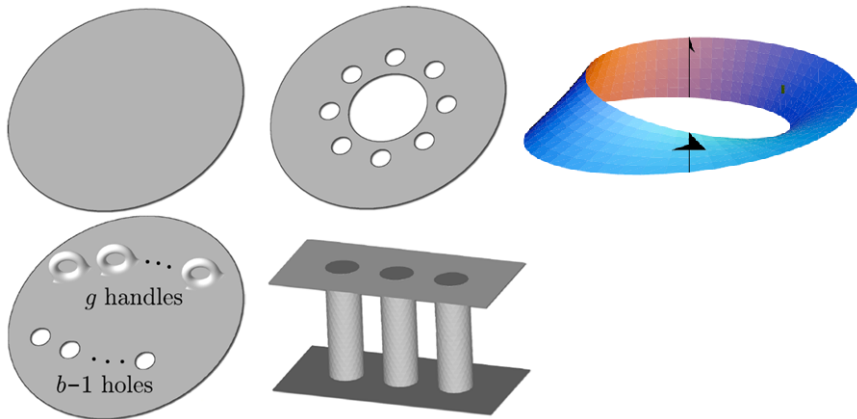
# How to compute relative generators?

For closed surfaces they are absolute generators !



# How to compute relative generators?

**Harder !** No combinatorial algorithm has been devised to compute relative cohomology generators



# Review of past work: Hiptmair-Ostrowski algorithm

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# Hiptmair-Ostrowski algorithm

- **Input:** closed and connected triangulated surface  $\Gamma$
- **Output:** generators of  $H_1(\Gamma)$

Introduced in:

- *R. Hiptmair, J. Ostrowski, Generators of  $H_1(\Gamma_h, \mathbb{Z})$  for Triangulated Surfaces: Construction and Classification, SIAM J. Comput., Vol. 31, No. 5, pp. 1405-1423 (2002)*

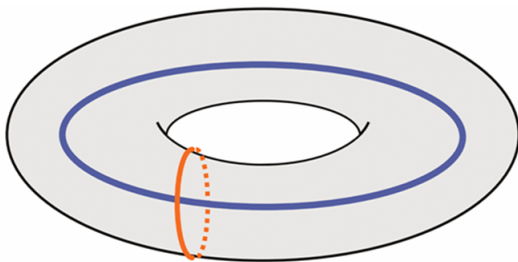
Reinvented several times, for example:s

- *D. Eppstein, Dynamic generators of topologically embedded graphs. Proc. ACM-SIAM Symp. on Discrete Algorithms, pp. 599–608 (2003)*
- *G. Rubinacci, A. Tamburrino, Automatic Treatment of Multiply Connected Regions in Integral Formulations, IEEE Trans. Magn., vol. 46, no. 8, pp. 2791-2794 (2010)*



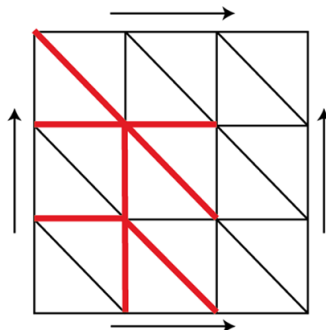
# Hiptmair-Ostrowski algorithm

**Example:** find generators for  $H_1(\Gamma)$ , where  $\Gamma$  is a **torus**



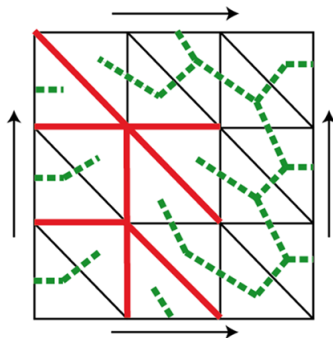
# Hiptmair-Ostrowski algorithm

- 1 Produce a **spanning tree**  $T$  of 1-skeleton of  $\Gamma$  (usually by Breadth First Search)



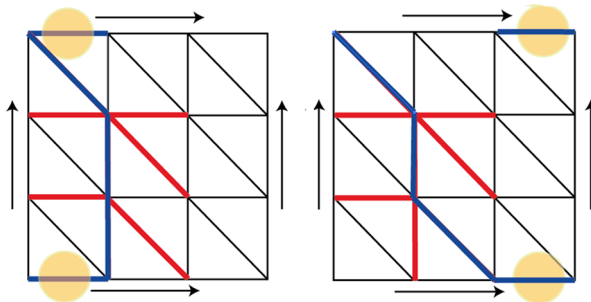
# Hiptmair-Ostrowski algorithm

- 2 Produce a spanning tree  $\tilde{T}$  of the **dual** 1-skeleton



# Hiptmair-Ostrowski algorithm

- ③ Add each edge  $e$  neither in  $T$  nor in  $\tilde{T}$  to the tree



The (unique) cycle in  $T$  is the support of the homology generator

# Hiptmair-Ostrowski algorithm for cohomology?

Usually in numerical analysis and engineering applications the **cohomology generators** are needed in place of homology generators

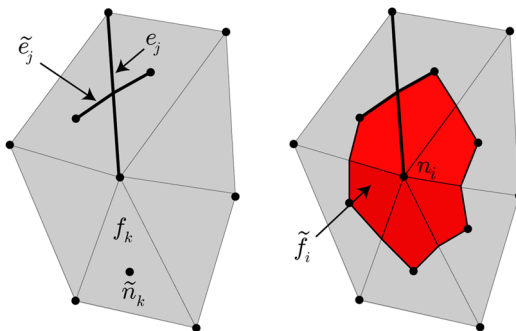
**Natural question:** is there an algorithm such that

- **Input:** triangulated surface with boundary  $\Gamma$  and
- **Output:** representatives of  $H^1(\Gamma, \partial\Gamma^c)$

# Hiptmair-Ostrowski algorithm for cohomology

If  $\Gamma$  is closed, leverage **Poincaré-Lefschetz duality**

$$H^1(\Gamma) \cong H_1(\tilde{\Gamma})$$

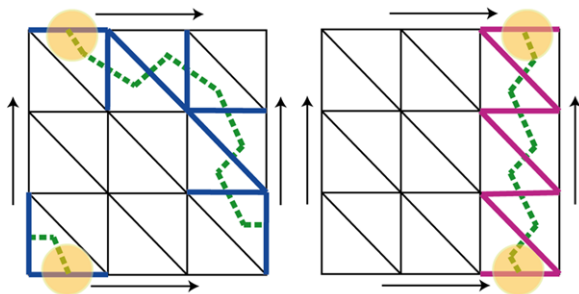


# Hiptmair-Ostrowski algorithm for cohomology

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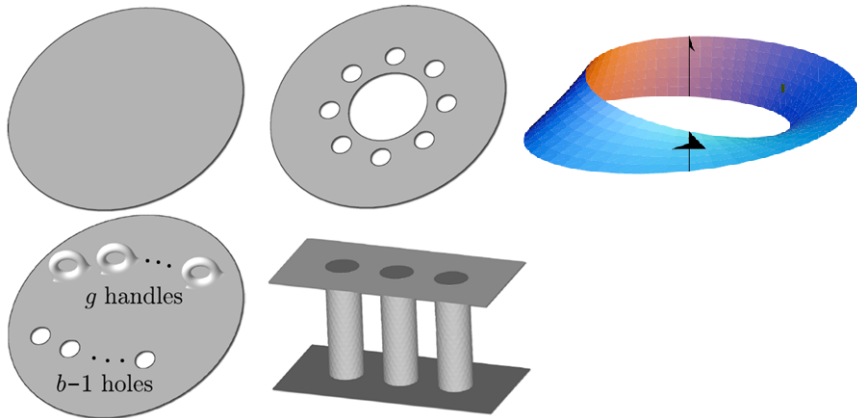
**Example:** algorithm on the torus, spanning trees  $T$  and  $\tilde{T}$  as before



Find **dual cycle** on  $\tilde{T}$  and use  $H^1(\Gamma) \cong H_1(\tilde{\Gamma})$

# Hiptmair-Ostrowski algorithm for open surfaces

**Harder !** No combinatorial algorithm has been devised for this case





# Main algorithm overview

- 1 Introduction
- 2 Key Mathematical Concepts
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# Main algorithm overview

**Goal:** Compute generators for  $H^1(\Gamma, \partial\Gamma^c)$

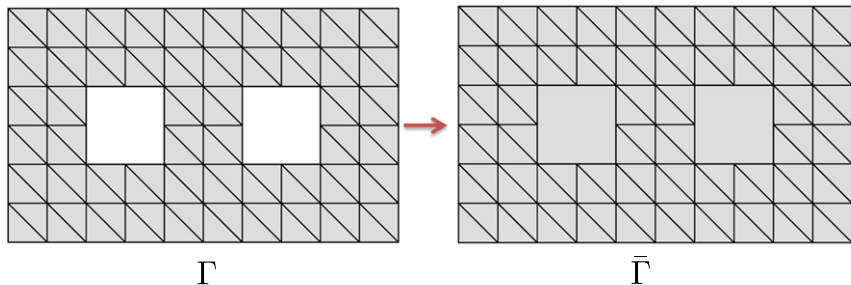
A **three-step** approach:

- 1 Compute  $N^{\text{ha}}$  generators  $\mathbf{h}_j^{\text{ha}}$  due to “handles”
- 2 Compute  $N^{\text{h}}$  generators  $\mathbf{h}_k^{\text{h}}$  due to “holes”
- 3 Compute  $N^{\text{c}}$  generators  $\mathbf{h}_l^{\text{c}}$  due to “contacts”

# Compute generators due to “handles”

Less easy task: proceed in two further steps

- 1 “Glue” **topological disks** (i.e. polygons) along the connected components of  $\partial\Gamma$  to obtain a closed surface  $\bar{\Gamma}$

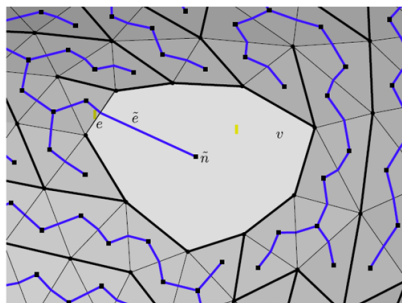


# Compute generators due to “handles”

Less easy task: proceed in two further steps

- 2 Apply the same ideas of the standard HO-algorithm on  $\bar{\Gamma}$

**Warning:** assure that generators do not pass through the attached disks



# Characterization of cohomology generators

## Universal Coefficient Theorem for cohomology

*If a chain complex  $C$  of free Abelian groups has homology groups  $H_n(C; G)$ , then the cohomology groups  $H^n(C; G)$  of the cochain complex  $\text{Hom}(C; G)$  are determined by the exact sequence*

$$0 \rightarrow \text{Ext}(H_{n-1}(C); G) \rightarrow H^n(C; G) \xrightarrow{\varphi} \text{Hom}(H_n(C); G) \rightarrow 0.$$

*with  $\varphi(\mathbf{h}) \mapsto \langle \mathbf{h}, \cdot \rangle$*

**Strategy:** find a set  $\{\mathbf{h}_1^{\text{ha}}, \dots, \mathbf{h}_{N^{\text{ha}}}^{\text{ha}}\}$  such that

$$\langle \mathbf{h}_i^{\text{ha}}, c_j \rangle = \delta_{i,j}$$

with  $\{c_j\}_{j=1}^{N^{\text{ha}}}$  a basis of **(the free part)** of  $H_1(\bar{\Gamma}; \mathbb{Z})$

# Challenge: torsion subgroup

## Structure of finitely generated abelian groups

*Let  $G$  be a finitely generated abelian group. There exists some natural number  $m \geq 0$ , and some positive natural numbers  $n_1, \dots, n_q$ , such that  $G$  is isomorphic to the direct sum*

$$\mathbb{Z}^m \oplus \mathbb{Z}/n_1\mathbb{Z} \oplus \dots \oplus \mathbb{Z}/n_q\mathbb{Z},$$

*where  $n_i$  divides  $n_{i+1}$  for all  $1 \leq i \leq q-1$ .*

**Corollary:** the homology group  $H_1(\bar{\Gamma}; \mathbb{Z})$  is isomorphic to a **direct sum**

$$H_1(\bar{\Gamma}; \mathbb{Z}) \cong \mathbb{Z}^m \oplus T$$

where  $T$  is the **torsion subgroup** of  $H_1(\bar{\Gamma}; \mathbb{Z})$  and  $\mathbb{Z}^m$  is the **free part** (a **free abelian group**) of dimension  $\beta_1(\bar{\Gamma})$

# A glimpse of Discrete Morse Theory

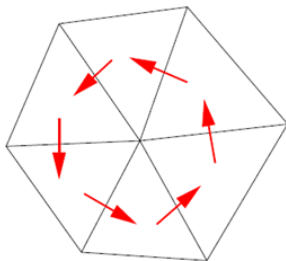
## Acyclic matching

A 1-matching  $M$  on  $\bar{\Gamma}$  is a set of **matched** pairs  $(f, e)$  such that:

- 1 If  $(f, e) \in M$ , then  $e \prec f$ .
- 2 If  $(f_1, e), (f_2, e) \in M$ , then  $f_1 = f_2$ .

A matching  $M$  is **acyclic** if there does not exist a cycle of of the form

$$f_1 \succ e_1 \prec f_2 \succ e_2 \prec \cdots \prec f_l \succ e_l \prec f_1$$



# Acyclic matching from spanning trees

## Lemma

*Let  $T$  be a spanning tree of  $\bar{\Gamma}$ .*

*There exist an acyclic matching  $M_T$  such that every vertex of  $T$ , except for one arbitrary vertex, is matched in  $M_T$ .*



# Morse boundary formula

## Theorem

Let  $M$  be an acyclic matching on  $\bar{\Gamma}$ , and denote by  $d_m$  the number of **non-matched**  $m$ -cells of  $\bar{\Gamma}$ . For every critical 2-cell  $f$  and 1-cell  $e$ , let  $S(f, e)$  be the set of all sequences  $s$  of matched cells of the form

$$s := (f \succ e_1 \prec f_1 \succ e_2 \prec f_2 \succ \cdots \succ e_l \prec f_l \succ e).$$

Then,  $\bar{\Gamma}$  is **homotopy equivalent** to a CW complex  $\Sigma$  with exactly  $d_m$  cells of dimension  $m$ , and the corresponding **Morse incidence number**  $\iota^M(f, e)$  in  $\Sigma$  is given by

$$\iota^M(f, e) := \sum_{s \in S(f, e)} \langle \partial_2 w(s), e \rangle,$$

where  $w(s)$  is the **weight** 2-chain computed from each sequence  $s$ .

# Applying the Morse framework

$\Sigma$  consists of:

- **One 0-cell**, the unique non-matched vertex of  $\bar{\Gamma}$  in  $M_T$ .
- $(|E| - |V| + 1) - (|F| - 1) = 2 - \chi$  **1-cells**, where  $\chi$  is the **Euler characteristic** of  $\bar{\Gamma}$ .
- **One 2-cell**, the unique non-matched face (a vertex, by duality) of  $\bar{\Gamma}$  in  $M_{\tilde{T}}$ .

**Observation:**  $\Sigma$  is homotopy equivalent to  $\bar{\Gamma}$  and the boundary of  $\bar{\Gamma}$  is **empty**  $\Rightarrow$  each Morse incidence number  $\iota^M(f, e)$  in  $\Sigma$  **can only take one of two possible values**:

- ①  $\iota^M(f, e) = 0$ .
- ②  $\iota^M(f, e) = 2\eta_e$ , with  $\eta_e \in \{-1, +1\}$ .

# Homology characterization of $\Sigma$ (and hence of $\bar{\Gamma}$ )

## Lemma

Let  $E_c^{\text{II}}$  be the set of critical 1-cells such that  $\iota^M(f, e) \neq 0$  in  $\Sigma$ . Denote by  $\mathbf{g} := 2 - \chi$  the **genus** of  $\Sigma$ . We distinguish between the following two cases:

- 1 If  $|E_c^{\text{II}}| = 0$ , then  $H_1(\Sigma; \mathbb{Z}) = \mathbb{Z}^{\mathbf{g}}$ .
- 2 If  $|E_c^{\text{II}}| > 0$ , then  $H_1(\Sigma; \mathbb{Z}) = \mathbb{Z}^{\mathbf{g}-1} \oplus \mathbb{Z}/2\mathbb{Z}$ . In this case, the **torsion generator**  $\mathbf{c}^*$  is given by

$$\mathbf{c}^* := \sum_{e \in E_c^{\text{II}}} \eta_e e,$$

where each  $\eta_e \in \{-1, +1\}$  is the sign of the Morse coefficient  $\iota^M(f, e)$

# Missing piece: identify the 1-cells of $\Sigma$ in $E_c^{\text{II}}$

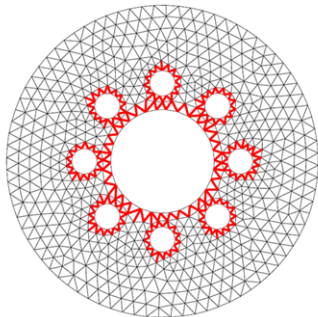
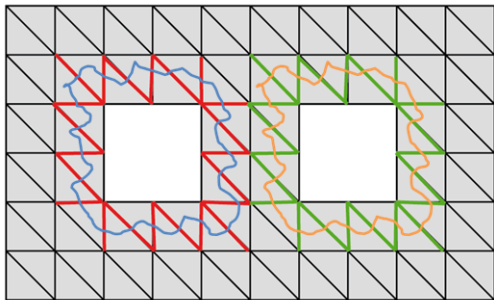
Easy ! Simply algorithm:

- 1 pick any triangle, assign any orientation to it
- 2 then go to neighbouring triangles (via some algorithm, it does not matter), and decide the orientation for those, which is forced to be one or another due to already oriented neighbors
- 3 if at some point you find inconsistency, that is, several neighbors force a single triangle to have different orientations, then the surface is non-orientable

Otherwise, when there are no more neighbors, you have found an orientable connected component of the mesh !

# Compute generators due to “holes”

**Very easy** ! Simply select the edges that go outside all connected components except one



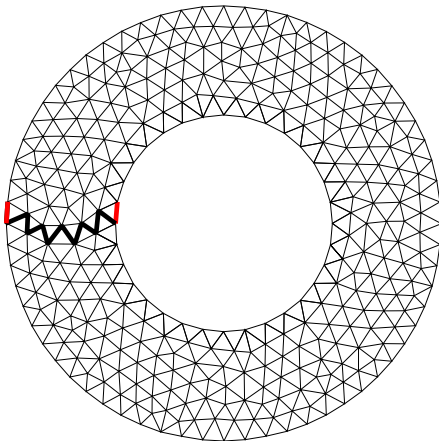
## First decomposition Lemma

$$H^1(\Gamma, \partial\Gamma) \cong H^1(\bar{\Gamma}) \oplus \langle [h_1^h], \dots, [h_{N^h}^h] \rangle$$

# Compute generators due to “contacts”

Also easy, modulo a technical detail

- 1 If  $|E_c^{\text{II}}| = 0$ , then simply select one generator for all contact regions except one (called **ground contact**)



# Compute generators due to “contacts”

This is also easy, modulo a technical detail involving the presence of a **torsion generator**

- 1 If  $|E_c^{\text{II}}| = 0$ , then simply select one generator for all contact regions except one (called **ground contact**)
- 2 If  $|E_c^{\text{II}}| > 0$ , there an additional generator associated with the chosen ground contact and the **torsion generator**

## Second decomposition Lemma

$$H^1(\Gamma, \partial\Gamma^c) \cong H^1(\Gamma, \partial\Gamma) \oplus \langle [\mathbf{h}_1^c], \dots, [\mathbf{h}_{N^c}^h] \rangle$$

We have the following isomorphism

## Cohomology characterization

$$\begin{aligned} H^1(\Gamma, \Gamma^c) &\cong \langle [\mathbf{h}_1^{\text{ha}}], \dots, [\mathbf{h}_{N^{\text{ha}}}^{\text{ha}}] \rangle && \text{(generators due to “handles”)} \\ &\oplus \langle [\mathbf{h}_1^{\text{h}}], \dots, [\mathbf{h}_{N^{\text{h}}}^{\text{h}}] \rangle && \text{(generators due to “holes”)} \\ &\oplus \langle [\mathbf{h}_1^{\text{c}}], \dots, [\mathbf{h}_{N^{\text{c}}}^{\text{h}}] \rangle && \text{(generators due to “contacts”)} \end{aligned}$$



# Conclusion and implications

- 1 Introduction
- 2 Key Mathematical Concepts
- 3 Review of past work: Hiptmair-Ostrowski algorithm
- 4 Main algorithm overview
- 5 Conclusion and implications**

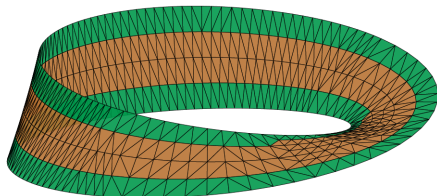
# Summary

- Novel algorithm to compute  $H^1(\Gamma, \partial\Gamma^c)$  with linear complexity
- Applications to EFIE and topological Data Analysis
- Handles non-orientable surfaces and torsion

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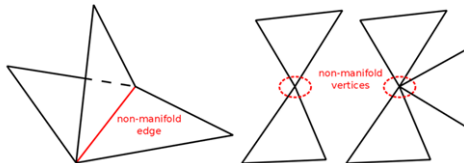
*B. Hofmann, T. F. Eibert, F. P. Andriulli, S. B. Adrian, A Low-Frequency Stable, Excitation Agnostic Discretization of the Right-Hand Side for the Electric Field Integral Equation on Multiply-Connected Geometries, IEEE Transactions on Antennas and Propagation, Vol. 71, No. 2, (2023)*

they say that **green triangles** are the support of a global loop

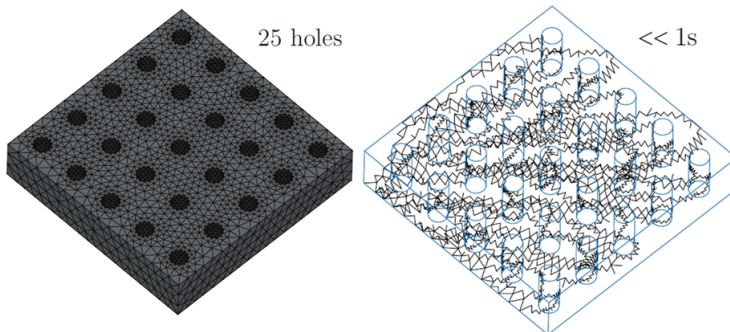


# Glimpses on future work

- Extension to non-manifold surfaces



- Minimal cohomology basis generation



# References



S. Pitassi, *Generators of  $H^1(\Gamma, \partial\Gamma^c)$  with  $\partial\Gamma^c \subset \partial\Gamma$  for Triangulated Surfaces  $\Gamma$ : Construction and Classification of Global Loops*, arXiv preprint (2025).