

Large-Scale Convex Optimization

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June 19, 2025

1 Introduction

Large Scale Problem of dimension n but iterations $\ll n$ desired

Convex One of the only problem classes that are “solvable”

Optimization with decision variable x , objective function f and

feasible set $C = \{x \in \mathbb{R}^n : g(x) \leq 0, h(x) = 0\}$

x^* **local minimum** if $\exists \epsilon > 0$ s.t. $f(x^*) \leq f(x)$,

$\forall x \in C \cap B_\epsilon(x^*), B_\epsilon(x^*) := \{x \in \mathbb{R}^n : |x - x^*| < \epsilon\}$

Proposition 1. f (lower-semi)-continuous, $f(x) \rightarrow \infty$ for

$|x| \rightarrow \infty, C$ closed $\Rightarrow \exists$ of OP with: $\min_{x \in C} f(x)$ and $x^* \in$

$\text{argmin}_{x \in C} f(x)$

Definition 1 (Lipschitz continuity). $q : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is Lipschitz

with constant L if: $|q(x) - q(y)| \leq L|x - y| \forall x, y \in \mathbb{R}^n$

Class of OP P with $C = [0, 1]^n$ and f is L^∞ -Lipschitz with L

Proposition 2. For any algorithm \exists problem in P , s.t. achieving

$|f(x_N) - f(x^*)| < \epsilon$ requires $N \geq (\lfloor \frac{L}{2\epsilon} \rfloor)^n - 1$

Definition 2. OP convex if, f and g_i convex functions, h affine.

Definition 3. $q : \mathbb{R}^n \rightarrow \mathbb{R}$ convex (affine) if $\forall x, y \in \mathbb{R}^n$

$q(\theta x + (1-\theta)y) \leq \theta q(x) + (1-\theta)q(y) \quad \forall \theta \in [0, 1]$

Proposition 3. If OP convex, local minimum == global minimum

2 Convex Optimization Problem

Definition 4 (Convex Set). A set C is convex if and only if

$\theta x + (1 - \theta)y \in C, \forall x, y \in C, \forall \theta \in [0, 1]$

(hyperplane \parallel half-space) $\{x \in \mathbb{R}^n \mid a^\top x (= \parallel \leq) b\}$

polyhedra $\{x \in \mathbb{R}^n \mid A^q \times^n x \preceq b^q \times 1, C^r \times^n x = d^r \times 1\}$

Operations that preserve convexity (sets)

Intersection C_1, C_2 cv $\Rightarrow C_1 \cap C_2$ convex (cv)

Image under affine map $C \subseteq \mathbb{R}^n$ cv $\Rightarrow \{Ax + b \mid x \in C\}$ cv

Inverse loaM $C \subseteq \mathbb{R}^m$ cv $\Rightarrow \{x \in \mathbb{R}^n \mid Ax + b \in C\}$ cv

Separating Hyperplane Theorem

Theorem 1. $C \subseteq \mathbb{R}^n$ non-empty closed (cl) convex set, $y \notin C$

$\Rightarrow \exists a \neq 0, b \in \mathbb{R}$ s.t. $a^\top x + b < a^\top y + b, \forall x \in C$

Corollary 1. $C_{\text{cl,cv}}$: intersection of cl half-spaces that contain C

Support function

Idea represent any cl,cv set by its supporting hyperplanes

$\sigma_C(a) = \sup_{x \in C} a^\top x$ if known, one can construct

$$C = \bigcap_{a \in \mathbb{R}^n} \{x \in \mathbb{R}^n \mid a^\top x - \sigma_C(a) \leq 0\} \\ = \{x \in \mathbb{R}^n \mid \sup_{a \in \mathbb{R}^n} a^\top x - \sigma_C(a) \leq 0\}$$

Definition 5. $f : \mathbb{R}^n \rightarrow \mathbb{R}$ cv \Leftrightarrow epigraph of f is cv set

$\text{epi}(f) := \{(x, t) \in \mathbb{R}^{n+1} \mid f(x) \leq t\}$

\rightarrow this provides a link between convex sets and functions

Operations that preserve convexity (functions)

- the point wise maximum of convex functions is convex

- the sum of convex functions is convex

- $f(Ax + b)$ is convex if f is convex

Check Convexity f is convex if it is composition of simple conv

conv function with convexity preserving operations or if

$f : \mathbb{R}^n \rightarrow \mathbb{R}$ twice differentiable, $\partial^2 f / \partial x^2 \succeq 0 \forall x \in \mathbb{R}^n$

$g : \mathbb{R} \rightarrow \mathbb{R}$ with $g(t) = f(x + tv)$ convex in $t \forall x, v \in \mathbb{R}^n$

Extended real numbers $\mathbb{R} = \mathbb{R} \cup \{+\infty, -\infty\}$

Indicator function $\psi_C(x) := \begin{cases} +\infty & \text{if } x \notin C \leq 0 \\ 0 & \text{if } x \in C \end{cases}$

\rightarrow this provides another link between convex sets and functions

We can write $\min_{x \in C} f(x)$ as $\min_{x \in \mathbb{R}^n} f(x) + \psi_C(x)$

Definition 6 (3). $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is called proper if f is bounded

below and if $\exists x \in \mathbb{R}^n$ s.t. $f(x) < \infty$

Definition 7 (Legendre Transformation). The conjugate function

of $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is defined as $f^*(y) = \sup_{x \in \mathbb{R}^n} y^\top x - f(x)$

3 KKT and Lagrange Duality

Basic 2d Example for derivation: $\min_{x \in \mathbb{R}^2} f(x)$ s.t. $h(x) = 0$

$\rightarrow \nabla f(x^*), \nabla h(x^*)$ co-linear $\Rightarrow \exists \nu^* \in \mathbb{R} : \nabla f(x^*) +$

$\nu^* \nabla h(x^*) = 0 \Leftrightarrow f(x) + \nu^* h(x)$ is stationary at x^*

Generalization for $n \rightarrow \infty$ and with constraints

We consider $f^* = \inf_{x \in \mathbb{R}^n} f(x)$ s.t. $h(x) = 0, g(x) \leq 0$ (1)

Lagrange f. $\mathcal{L}(x, \lambda, \nu) = f(x) + \lambda^\top g(x) + \nu^\top h(x)$ (2)

Proposition 4 (Weak Duality). The dual function $d(\lambda, \nu) =$

$\inf_{x \in \mathbb{R}^n} \mathcal{L}(x, \lambda, \nu)$ satisfies $d(\lambda, \nu) \leq f^*, \forall \lambda \geq 0, \nu \in \mathbb{R}^{n_h}$

Definition 8 (Constraint qualification). C convex, **Slater's Condi-**

tion holds if $\exists \hat{x} \in \mathbb{R}^n$ s.t. $h(\hat{x}) = 0$ and $g(\hat{x}) < 0$

Proposition 5 (Strong Duality). If Slater's condition holds and

(1) is convex $\Rightarrow \exists \lambda \geq 0, \nu \in \mathbb{R}^{n_h}$ s.t. $d(\lambda, \nu) = f^*$

KKT

Theorem 2 (KKT Conditions). Slater's condition holds and (1)

is convex $\rightarrow x^* \in \mathbb{R}^n$ is a minimizer of the primal (1) and

$(\lambda^* \geq 0, \nu^*) \in \mathbb{R}^{n_g} \times \mathbb{R}^{n_h}$ is a maximizer of the dual \Leftrightarrow

$\nabla_x \mathcal{L}(x^*, \lambda^*, \nu^*) = 0$ KKT-1 (Stationary Lagrangian)

$g(x^*) \leq 0, h(x^*) = 0$ KKT-2 (primal feasibility)

$\lambda^* \leq 0, \nu^* \in \mathbb{R}^{n_h}$ KKT-3 (dual feasibility)

$\lambda^{*\top} g(x^*) = 0 = \nu^{*\top} h(x^*)$ KKT-4 (complementary slackness)

In addition we have: $\sup_{\lambda \geq 0, \nu \in \mathbb{R}^{n_h}} q(\lambda, \nu) = \inf_{x \in C} f(x)$

Remark Without Slater, KKT-1-4 still implies x^* minimizes (1)

and λ, ν maximizes dual, but the converse is no longer true.

There can be primal-minimizer/dual-maximizer not satisfy KKT.

Subdifferential

For cv f we have $f(x) \geq f(\bar{x}) + \nabla f(\bar{x})^\top (x - \bar{x}), \forall x, \bar{x} \in \mathbb{R}^n$

Definition 9. $f : \mathbb{R}^n \rightarrow \mathbb{R}$ cv, the subdifferential of f at \bar{x} is:

$\partial f(\bar{x}) := \{\lambda \in \mathbb{R}^n \mid f(x) \geq f(\bar{x}) + \lambda^\top (x - \bar{x}), \forall x \in \mathbb{R}^n\}$

Proposition 6. f (like D9), $x^* \in \text{argmin}_x f(x) \Leftrightarrow 0 \in \partial f(x^*)$

Proposition 7 (Relation to conjugate functions). For convex f

with $\text{epi}(f)$ closed: $y \in \partial f(x) \Leftrightarrow x \in \partial f^*(y)$

4 Convex Optimization Problems

Optimal value $f^* = \inf\{f(x) \mid g_i(x) \leq 0, h_j = 0\}$

$f^* = +\infty$ OP is infeasible, $f^* = -\infty$ OP is unbound below

Feasibility Problem

Special case $f(x) = 0, \forall x \Leftrightarrow \min_s$ s.t. $g_i(x) \leq s, h_j(x) = 0$

Linear Programming minimize $c^\top x$ s.t. $Ax - b \geq 0, x \geq 0$

Step 1: $\mathcal{L}(x, \lambda_1, \lambda_2) = c^\top x - \lambda_1^\top (Ax - b) - \lambda_2^\top x, \lambda_i \geq 0$

Step 2: $\inf_{x \in \mathbb{R}^n} \mathcal{L}(x, \lambda_1, \lambda_2) = \begin{cases} \lambda_1^\top b & \text{if } c - A^\top \lambda_1 - \lambda_2 = 0 \\ -\infty & \text{if } c - A^\top \lambda_1 - \lambda_2 = 0 \end{cases}$

Step 3: Dual, maximize $b^\top \lambda$ s.t. $c - A^\top \lambda \geq 0, \lambda \geq 0$ (again LP)

Proposition 8. The optimal solution of a linear program (if it

exists) lies always on the boundary of the feasible set and there

exists an optimal solution that is a vertex of the feasible set.

Quadratic Programming convex if $P = P^\top$ positiv semi-definite

minimize $\frac{1}{2} x^\top P x + q^\top x$ s.t. $Gx \leq h, Ax = b$

Second-Order Cone Program

minimize $f^\top x$ s.t. $|A_i x + b| \leq c_i^\top x + d_i, Fx = g$

Second-order cone $C_{n+1} = \{(x, t) \mid x \in \mathbb{R}^n, t \in \mathbb{R}, |x| \leq t\}$

Semi-Definite Programming with symmetric F_i, X, A_i

minimize $c^\top x$ s.t. $\sum_{i=1}^n x_i F_i + G \preceq 0, Ax = b$

Standard form minimize $\text{tr}(CX)$ s.t. $X \succeq 0, \text{tr}(A_i X) = b_i$

$\text{tr}(CX) = \sum_{i=1}^n \sum_{j=1}^n C_{ij} X_{ij}, C \in \mathbb{R}^{n \times n}, i = 1, \dots, m$

LP \subset QP \subset QCQP (Quadratically Constrained QP) \subset SOCP \subset SDP

5 Gradient methods - Part I

Definition 10 (smoothness). $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is L -smooth (L -sm)

if $\nabla f(x)$ satisfies $|\nabla f(x) - \nabla f(y)| \leq L|x - y| \forall x, y \in \mathbb{R}^n$

Taylor $\rightarrow f(y) \leq f(x) + \nabla f(x)^\top (y - x) + \frac{L}{2}|x - y|^2$

Definition 11 (strong convexity). $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is μ -strongly

convex (μ -scv) if $f(y) \geq f(x) + \nabla f(x)^\top (y - x) + \frac{\mu}{2}|x - y|^2$

Gradient Descent

$x_{k+1} = x_k - T \nabla f(x_k)$ for $k = (k_0, \dots, k_N)$ given x_0, T

Assume $f(x) = c_0 + b^\top x + \frac{1}{2} x^\top H x, H \succ 0 \Rightarrow H x^* = -b$

$x_{k+1} - x^* = x_k - x^* - T(b + H x_k) = (I - TH)(x_k - x^*)$

Convergence given by eigenvalues of $I - TH$, use $H = U \Lambda U^\top$

$x_N - x^* = U(I - T \Lambda)^N U^\top (x_0 - x^*) \rightarrow \text{conv-rate } 1 - T \lambda_i$

f : L -sm, μ -scv $\rightarrow \mu \leq \min \lambda_i, \max \lambda_i \leq L, \rightarrow \text{conv-rate } \rho(T) =$

f : $\max_{\mu \leq h \leq L} |1 - Th| \rightarrow |x_N - x^*| \leq \rho(T)^N |x_0 - x^*|$

$T^* = \frac{2}{L + \mu}$, with condition number $\kappa := \frac{L}{\mu}$ and $1 - \xi \leq e^{-\xi}$

$\rho(T^*) = \frac{L - \mu}{L + \mu} = \frac{\kappa - 1}{\kappa + 1} = (1 - \frac{2}{\kappa + 1}) \leq e^{-\frac{2}{\kappa + 1}} \rightarrow \text{algebraic}$

complexity $N \geq \frac{\kappa + 1}{2} \ln(\frac{|x_0 - x^*|}{\epsilon})$ to achieve $|x_N - x^*| \leq \epsilon$

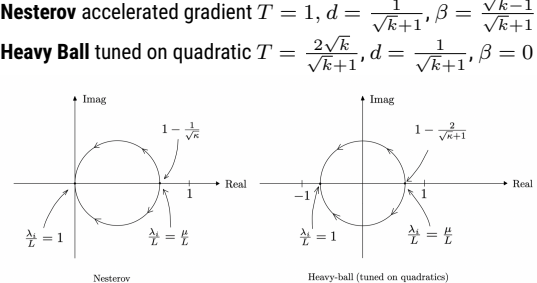
Momentum-based methods

$$q_{k+1} = q_k + T p_{k+1}$$

$$p_{k+1} = (1 - 2dT) p_k - T \nabla f(q_k + \beta p_k) / L$$

Nesterov accelerated gradient $T = 1, d = \frac{1}{\sqrt{k+1}}, \beta = \frac{\sqrt{k}-1}{\sqrt{k+1}}$

Heavy Ball tuned on quadratic $T = \frac{2\sqrt{\kappa}}{\sqrt{k+1}}, d = \frac{1}{\sqrt{k+1}}, \beta = 0$



$C_{\text{Nesterov}} (1 - \frac{1}{\sqrt{\kappa}})^N \approx \frac{|q_N - q^*|}{|q_0 - q^*|} \approx C_{\text{HeavyBall}} (1 - \frac{2}{\sqrt{\kappa+1}})^N$

Theorem 3. f : L -sm, μ -scv \rightarrow Nesterov's method satisfies:

$$|q_N - q^*| \leq \sqrt{\kappa + 1} (1 - 1/\sqrt{\kappa})^{N/2} |q_0 - q^*| \\ f(q_N) - f^* \leq \frac{L + \mu}{2} (1 - 1/\sqrt{\kappa})^N |q_0 - q^*|^2$$

Requires $N \geq 2\sqrt{\kappa} \ln(\frac{|q_0 - q^*|}{\epsilon})$ to achieve $|x_N - x^*| \leq \epsilon$

Theorem 4. For any first-order method $\exists f : \mathbb{R}^\infty \rightarrow \mathbb{R}, \mu$ -scv,

L -sm, s.t. $|x_k - x^*| \geq (1 - \frac{2}{\sqrt{\kappa+1}})^k |x_0 - x^*| \forall k \geq 0$

Line search

Adaptive Methods

6 Gradient Methods - Part II

Definition 12. $\text{prox}_C(x) = \text{argmin}_{y \in C} \frac{1}{2}|x - y|^2$ with $C \subset \mathbb{R}^n$

Lemma 1. $\text{cl}, \text{cv } C \subset \mathbb{R}^n \rightarrow |\text{prox}_C(x) - \text{prox}_C(y)| \leq |x - y|$

$\leftarrow |\text{prox}_C(x) - \text{prox}_C(y)|^2 \leq (\text{prox}_C(x) - \text{prox}_C(y))^\top (x - y)$

Projected Gradient Descent

$x_{k+1} = \text{prox}_C(x_k - T \nabla f(x_k))$, for $x_0, k_{0..N}, T \in (0, 2/L)$

Proposition 9. f : L -sm, μ -scv \rightarrow projected GD with $T = \frac{2}{L + \mu}$

satisfies $|x_N - x^*| \leq |x_0 - x^*| (1 - \frac{2}{\kappa + 1})^N$ (κ still $\frac{L}{\mu}$)

Lemma 2. $f : \mathbb{R}^n \rightarrow \mathbb{R}, L$ -sm, $\text{cv} \rightarrow \tilde{f}$ strongly-cv

$\hat{f}(x) = f(x) + \frac{\mu}{2}|x - x_0|^2$ and $|\tilde{x}^* - x_0| \leq |x^* - x_0|$

and $f(x) - f(x^*) \leq \tilde{f}(x) - \tilde{f}(\tilde{x}^*) + \frac{\mu}{2}|x^* - x_0|^2, \mu > 0$

\rightarrow from here one can apply GD or Nesterov, which results in:

$f(x_N) - f(x_0) \leq \epsilon$ after $N \sim L|x^* - x_0|^2/\epsilon$ iterations

Proposition 10 (Subgradient Method). $\text{cl}, \text{cv } C$ contained in ball

of radius $R, x_{0..N-1}$ satisfy $f(\frac{1}{N} \sum_{k=0}^{N-1} x_k) - f(x^*) \leq \frac{RLf}{\sqrt{N}}$

under $x_{k+1} = \text{prox}_C(x_k - T g_k) \quad g_k \in \partial f(x_k), T = \frac{R}{L_f \sqrt{N}}$

Assumptions on f	Method	$N : f(x_N) - f(x^*) \leq \epsilon$	Optimal
μ -strongly convex	gradient descent	$N \sim \kappa \ln(1/\epsilon)$	No
L -smooth	Nesterov	$N \sim \sqrt{\kappa} \ln(1/\epsilon)$	Yes
	gradient descent	$N \sim 1/\epsilon$	No
	Nesterov (varying stepsize)	$N \sim 1/\sqrt{\epsilon}$	Yes
L_f -Lipschitz, compact set	subgradient method	$N \sim 1/\epsilon^2$	Yes