# **Model Predictive Control**

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github.com/silvasta/summary-mpc



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# Requirements for MPC

# A model of the system

#### 2. A state estimator

- 3. Define the optimal control problem
- 4. Set up the optimization problem
- 5. Get the optimal control sequence (solve the optimization problem)
- 6. Verify that the closed-loop system performs as desired

## 1.1 Exact ODE solution of a Linear System

$$x(t) = e^{A^{c}(t-t_{0})}x_{0} + \int_{t_{0}}^{t} e^{A^{c}(t-\tau)}B^{c}u(\tau)d\tau$$

**Problem** Most physical systems are nonlinear

**Idea** use First Order Taylor expansion  $f(\bar{x}) + \frac{\partial f}{\partial x^{\top}}\Big|_{\bar{x}} (x - \bar{x})$ 

# 1.2 Linearization

$$\dot{x_s} = g(x_s, u_s) = 0$$
  $\Delta \dot{x} = \dot{x} - \dot{x_s} = A^c \Delta x + B^c \Delta u$   
 $y_s = h(x_s, u_s)$   $\Delta y = y - y_s = C \Delta x + D \Delta u$ 

$$A^c = \frac{\partial g}{\partial x^T}\Big|_{\substack{x_s \\ u_s}} B^c = \frac{\partial g}{\partial u^T}\Big|_{\substack{x_s \\ u_s}} C = \left.\frac{\partial h}{\partial x^T}\right|_{\substack{x_s \\ u_s}} D = \left.\frac{\partial h}{\partial u^T}\right|_{\substack{x_s \\ u_s}}$$

For general nonlinear systems only approximate discretization methods exist, such as Euler, quality depends on sampling time

# $\dot{x}^c \approx \frac{x^c(t+T_s) - x^c(t)}{T_s}$ $x(k) := x^c(t_0 + kT_s)$ $u(k) := u^c(t_0 + kT_s)$

#### **Exact Discretization of Linear Time-Invariant Models**

$$x(t_{k+1}) = \underbrace{e^{A^c T_s}}_{=A} x(t_k) + \underbrace{\int_0^{T_s} e^{A^c (T_s - \tau)} B^c d\tau}_{B = (A^c)^{-1} (A - \mathbb{I}) B^c} u(t_k)$$
$$x(k+N) = A^N x(k) + \sum_{i=0}^{N-1} A^i B u(k+N-1-i)$$

# 1.4 Analysis of LTI Discrete-Time Systems

Controllabe if rank(C) = n,  $C = \begin{bmatrix} B & \cdots & A^{n-1}B \end{bmatrix}$  $\forall (x(0),x^*) \exists$  finite time N with inputs  $\mathcal{U}$ , s.t.  $x(N)=x^*$ 

Stabilizable iff all uncontrollable modes stable **Observable** if rank $(\mathcal{O}) = n$ ,  $\begin{bmatrix} C^{\top} & \cdots & (CA^{n-1})^{\top} \end{bmatrix}^{\top}$ 

 $\forall x(0) \exists$  finite time N, s.t. the measurements

 $y(0), \dots, y(N-1)$  uniquely distinguish initial state x(0)

**Detectablitiy** iff all unobservable modes stable

# 1.5 Lyapunov

Stability is a property of an equilibrium point  $\bar{x}$  of a system **Definition 1** (Lyapunov Stability).  $\bar{\mathbf{x}}$  is Lyapunov stable if:

 $\forall \, \epsilon > 0 \, \exists \, \delta(\epsilon) \, \text{s.t.} \, ||x(0) - \bar{x}|| < \delta(\epsilon) \rightarrow ||x(k) - \bar{x}|| < \epsilon$ **Definition 2** (Globally asymptotic stability). If  $\bar{\mathbf{x}}$  is Lyapunov stable and attractive, i.e.,  $\lim_{k\to\infty} ||x(k)-\bar{x}||=0, \ \forall x(0)$ then  $\bar{x}$  is globally asymptotic stable.

### **Definition 3** (Global Lyapunov function). For $\bar{\mathbf{x}} = 0$ , function $V:\mathbb{R}^n o \mathbb{R}$ is called **Lyapunov function** if it is continuous at the origin, finite $\forall x \in \mathbb{R}^n$ .

$$||x|| \to \infty \Rightarrow V(x) \to \infty$$

$$V(x) > 0 \,\forall \, x \in \mathbb{R}^n \setminus \{0\} \quad V(0) = 0$$

$$V(g(x)) - V(x) \le -\alpha(x) \quad \forall x \in \mathbb{R}^n$$

where  $\alpha:\mathbb{R}^n\to\mathbb{R}$  continuous positive definite

# **Lyapunov Theorem**

**Theorem 1.** If a system admits a Lyapunov function V(x), then  $\bar{\mathbf{x}} = 0$  is globally asymptotically stable.

Theorem 2 (Lyapunov indirect method). For linearization of system around  $\bar{\mathbf{x}}=0$  and resulting matrix  $A=\left.\frac{\partial g}{\partial x^T}\right|_{x=0}$ with eigenvalues

$$|\lambda_i| := \begin{cases} \forall i := |\lambda_i| < 1 & \text{x=0 is asymptotically stable} \\ \exists i := |\lambda_i| > 1 & \text{origin is unstable} \\ \exists i := |\lambda_i| = 1 & \text{no info about stability} \end{cases}$$
 Discrete-Time Lyapunov equation

$$A^T P A - P = -Q, \quad Q > 0$$

Theorem 3 (Existence of solution of DT Lyapunov equation). The discrete-time Lyapunov equation (3) has a unique solution P > 0 if and only if A has all eigenvalues inside the unit circle, i.e. if and only if the system x(k + 1) = Ax(k) is stable.

# 1.6 Optimal Control

Unconstrained Finite Horizon Control Problem

$$\begin{split} J^{\star}(x(0)) &:= \min_{U} x_N^\top P x_N + \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i \\ \text{subject to } x_{i+1} &= A x_i + B u_i \quad i = 0, \dots, N-1 \\ x_0 &= x(0) \end{split}$$

 $P \succ 0$ , with  $P = P^T$  terminal weight

 $Q \succ 0$ , with  $Q = Q^T$  state weight  $R \succ 0$ , with  $R = R^T$  input weight

#### 1.7 Batch Approach

expresses cost function in terms of x(0) and input sequence U

$$\begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} \mathbb{I} \\ A \\ \vdots \\ A^N \end{bmatrix} x(0) + \begin{bmatrix} 0 & \cdots & 0 \\ B & 0 & 0 \\ AB & B & 0 \\ \vdots & \ddots & 0 \\ A^{N-1}B & \cdots & B \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix}$$

 $\overline{Q} := \mathsf{blockdiag}(Q, \dots, Q, P) \quad \overline{R} := \mathsf{blockdiag}(R, \dots, R)$ Optimal Input (from  $\nabla_U J(x(0), U) = 2HU + 2F^{\top} x(0) = 0$ )

$$U^{\star}(x(0)) = -\underbrace{\left( \underbrace{\mathcal{S}^{u}}\right)^{\top} \overline{Q} \mathcal{S}^{u} + \overline{R} \right)}_{H(\text{Hessian})^{-1}} \underbrace{\left( \mathcal{S}^{u}\right)^{\top} \overline{Q} \mathcal{S}^{x}}_{F^{\top}} x(0)$$

# **Optimal Cost**

 $J^{\star}(x(0)) {=} x(0)^{\top} (\mathcal{S}_{r}^{\top} \overline{Q} \mathcal{S}_{x} {-} \mathcal{S}_{x}^{\top} \overline{Q} \mathcal{S}_{u} (\mathcal{S}_{u}^{\top} \overline{Q} \mathcal{S}_{u} {+} \overline{R})^{-1} \mathcal{S}_{u}^{\top} \overline{Q} \mathcal{S}_{x}) x(0)$ 

# 1.8 Recursive Approach

uses dynamic programming to solve problem backwards from N

$$J_j^\star(x(j)) := \min_{U_j \to N} x_N^\top P x_N + \sum_{i=j}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

# **Ricatti Equations**

RDE - Riccati Difference Equation

$$P_i = A^{\top} P_{i+1} A + Q - A^{\top} P_{i+1} B (B^{\top} P_{i+1} B + R)^{-1} B^{\top} P_{i+1} A$$

RDE - Riccati Difference Equation solved recursively

ARE - Algebraic Riccati Equation solved analytically

From Principle Of Optimality Optimal Cost-To-Go

$$J_j^{\star}(x_j) = \min_{u_j} I(x_i, u_i) + J_{j+1}^{\star}(x_{j+1}) \ J_i^{\star}(x_i) = x_i^{\top} P_i x_i$$

**Optimal Control Policy** 

$$u_i^{\star} = F_i x_i = -(B^{\top} P_{i+1} B + R)^{-1} B^{\top} P_{i+1} A \cdot x(i)$$

# 1.9 Comparison of Batch and Recursive Approaches

Batch optimization returns sequence  $U^{\star}(x(0))$  of **numeric** values depending only on x(0), dynamic programming yields **feedback policies**  $u_i^{\star} = F_i x_i$  depending on each  $x_i$ .

#### Choice of P

- 1. Match infinite solution, use ARE
- 2. Assume no control needed after N. use Lyapunov Equation (makes only sense when asymptotically stable, otherwise P not positive definite)
- 3. set constraint  $x_{i+N} = 0$
- 1.10 Infinite Horizon LOR

# LOR

$$\begin{split} J_{\infty}^{\star}(x(k)) &= \min \sum_{i=0}^{\infty} x_i^\top Q x_i + u_i^\top R u_i \\ \text{subj. to } x_{i+1} &= A x_i + B u_i, \quad x_0 = x(k) \end{split}$$

Same u as for finite problem but with ARE Constant Feedback Matrix  $F\infty$  asymptotically stable for.. Q,R,stabi,detect

### 1.11 Optimization

A mathematical optimization problem is generally formulated

# **Mathematical Optimization Problem**

minimize f(x)

subject to:

 $g_i(x) \leq 0$ 

 $h_i(x) = 0$ 

Decision variable  $x \in \mathbb{R}^n$ **Objective function**  $f: dom(f) \to \mathbb{R}$ **Inequality constraints**  $q_i$  ( $i \in \#$ constraints) Equality constraints  $h_i$  ( $i \in \#$ constraints) Fesabile set  $\mathcal{X} := \{x | g(x) \le 0, h(x) = 0\}$ 

 $x^*$  is a Global Minimum if  $f(x^*) < f(x)$ 

Step 1:  $\mathcal{L}(x, \lambda_1, \lambda_2) = c^\mathsf{T} x - \lambda_1^\mathsf{T} (Ax - b) - \lambda_2^\mathsf{T} x, \ \lambda_i \ge 0$ Step 2:  $\inf_{\alpha \in \mathbb{R}^n} \mathcal{L} = \lambda_1^\mathsf{T} b$  , if  $c - A^\mathsf{T} \lambda_1 - \lambda_2 = 0$ , else  $-\infty$ 

Step 3: Dual, maximize  $b^{\mathsf{T}}\lambda$  s.t.  $c-A^{\mathsf{T}}\lambda > 0, \lambda > 0$  (again LP) Quadratic Programming min ...

# 1.14 Optimality Conditions

Consider 
$$f^\star = \inf_{x \in \mathbb{R}^n} f(x)$$
 s.t.  $g(x) \le 0, \ h(x) = 0$ 

#### 1.12 Convex Sets

**Definition 4** (Convex Set). Set C is convex if and only if

$$\theta x + (1 - \theta)y \in \mathcal{C}, \ \forall \ x, y \in \mathcal{C}, \ \forall \ \theta \in [0, 1]$$

**Feasible point**  $x \in dom(f)$  with  $q_i(x) < 0$ ,  $h_i(x) = 0$ 

**Strictly feasible point** x with strict inequality  $q_i(x) < 0$ 

Optimal value  $f^*(\text{or }p^*) = \inf\{f(x)|g_i(x) < 0, h_i = 0\}$ 

 $f^* = +\infty$ : OP infeasible,  $f^* = -\infty$ : OP unbound below

**Optimizer** set:  $\operatorname{argmin}_{x \in \mathcal{X}} f(x) := \{x \in \mathcal{X} | f(x) = f^*\}$ 

 $x^*$  is a Local Minimum if  $\exists \epsilon > 0$  s.t.  $f(x^*) < f(x)$ 

 $\forall x \in \mathcal{X} \cap B_{\epsilon}(x^{\star})$ , open ball with center  $x^{\star}$  and radius  $\epsilon$ 

**Definition 5** (Hyperplanes).  $\{x \in \mathbb{R}^n \mid a^\mathsf{T} x = b\}$ **Definition 6** (Halfspaces).  $\{x \in \mathbb{R}^n \mid a^\mathsf{T} x < b\}$ 

can be open (strict inequality) or closed (non-strict inequality) Definition 7 (Polyhedra). intersection of finite number of closed halfspaces: polyhedra  $\{x \in \mathbb{R}^n \mid A^{q \times n} x \prec b^{q \times 1}, \}$ 

**Definition 8** (Polytope). is a **bounded** polyhedron. **Definition 9** (Convex hull). for  $\{v_1, ..., v_k\} \in \mathbb{R}^d$  is:

 $\begin{array}{l} \operatorname{co}(\{v_1,...,v_k\}) := \{x | x = \sum_i \lambda_i v_i, \lambda \geq 0, \sum_i \lambda_i = 1\} \\ \text{Definition 10 (Ellipsoid)}. \ \ \, \operatorname{set:} \ \{x | (x - x_c)^\top A^{-1}(x - x_c) \leq 1\} \end{array}$ where  $x_c$  is center of ellipsoid,  $A \succ 0$  (i.e. positive definite) (Semi-axis lengths are square roots of eigenvalues of A)

**Definition 11** (Norm Ball).  $B_r(x) := \{ \xi \in \mathbb{R}^n : |\xi - x|_p < r \}$ where p defines the  $l_p$  norm,  $p = \{1|2|..|\infty\}$ 

Intersection  $C_1, C_2$  cv  $\Rightarrow C_1 \cap C_2$  convex (cv)

Image under affine map  $\mathcal{C} \subseteq \mathbb{R}^n$   $cv \Rightarrow \{Ax + b \mid x \in \mathcal{C}\}$  cv

Inverse loam  $\mathcal{C} \subseteq \mathbb{R}^m$  cv  $\Rightarrow \{x \in \mathbb{R}^n \mid Ax + b \in \mathcal{C}\}$  cv

# 1.13 Convex Functions

**Definition 12** (Convex Function).  $f: \mathcal{C}_{cv} \to \mathbb{R}$  is convex iff

$$f(\theta x + (1 - \theta)y) \le \theta f(x) + (1 - \theta)f(y), \ \forall \ x, y \in \mathcal{C}, \ \forall \ \theta \in [0, 1]$$

f is strictly convex if this inequality is strict.

**Definition 13** (Epigraph).  $f: \mathbb{R}^n \to \mathbb{R}$  cv  $\Leftrightarrow$  epi(f) is cv set

$$\operatorname{epi}(f) := \{(x,t) \in \mathbb{R}^{n+1} | f(x) \le t\}$$

**Check Convexity** f is convex if it is composition of simple convex function with convexity preserving operations or if  $f: \mathbb{R}^n \to \mathbb{R}$  twice differentiable,  $\partial^2 f/\partial x^2 \succeq 0 \ \forall \ x \in \mathbb{R}^n$  $q:\mathbb{R}\to\mathbb{R}$  with q(t)=f(x+tv) convex in  $t\ \forall\ x,v\in\mathbb{R}^n$  $\rightarrow f$  convex (restriction to a line)

- the point wise maximum of convex functions is convex
- the sum of convex functions is convex
- f(Ax + b) is convex if f is convex

Theorem 4. For a convex optimization problem, any locally optimal solution is globally optimal (local optima are global optima).

# **Lagrange Duality**

**Linear Programming** minimize  $c^{\mathsf{T}}x$  s.t. Ax - b > 0, x > 0

Consider 
$$f^\star = \inf_{x \in \mathbb{R}^n} f(x)$$
 s.t.  $g(x) \le 0, \ h(x) = 0$ 

**Proposition 1** (Weak Duality).  $d(\lambda, \nu) < f^*, \forall \lambda > 0, \nu \in \mathbb{R}^h$ **Definition 14** (Constraint qualification). **Slaters Condition** holds if  $\exists$  at least one strictly feasible point  $\hat{x}$  ( $h(\hat{x}) = 0$ ,  $q(\hat{x}) < 0$ ) Proposition 2 (Strong Duality). If Slater's condition holds and OP is convex  $\Rightarrow \exists \lambda \geq 0, \nu \in \mathbb{R}^{n_h}$  s.t.  $d(\lambda, \nu) = f^*$ 

# **KKT Conditions** (Karush-Kuhn-Tucker)

Theorem 5 (KKT Conditions). If Slater's condition holds and (1.14) is convex  $\to x^* \in \mathbb{R}^n$  is a minimizer of the primal (1.14) and  $(\lambda^{\star} \geq 0, \nu^{\star}) \in \mathbb{R}^{n_g} \times \mathbb{R}^{n_h}$  is a maximizer of the dual  $\Leftrightarrow$  is equivalent to the following statements:

**KKT-1** (Stationary Lagrangian)  $\nabla_x \mathcal{L}(x^*, \lambda^*, \nu^*) = 0$ **KKT-2** (primal feasibility)  $q(x^*) < 0, h(x^*) = 0$ KKT-3 (dual feasibility)  $\lambda^{\star}, \nu^{\star} \in \mathbb{R}^{n_h} > 0$ **KKT-4** (compenentary  $\lambda^{\star T} q(x^{\star}) = 0$  $\nu^{\star T} h(x^{\star}) = 0$ slackness)

In addition we have:  $\sup_{\lambda>0,\nu\in\mathbb{R}^{n_h}}q(\lambda,\nu)=\inf_{x\in\mathcal{C}}f(x)$ 

**Remark** Without Slater, KKT1-4 still implies  $x^*$  minimizes (1.14) and  $\lambda$ ,  $\nu$  maximizes dual, but the converse is no longer true. There can be primal-minimizer/dual-maximizer not satisfy KKT.

# 2 Nominal-MPC

#### 2.1 CFTOC

# **CFTOC** Constrained Finite Time Optimal Control problem

$$J(x(k)) = x_N^{\top} P x_N + \sum_{i=0}^{N-1} x_i^{\top} Q x_i + u_i R u_i$$

$$J^{\star}(x(k)) = \min_{U} I_f(x_N) + \sum_{i=0}^{N-1} I(x_i, u_i)$$

s.t 
$$x_{i+1} = Ax_i + Bu_i, i = 0, ..., N-1$$
  
 $x_i \in \mathcal{X}, u_i \in \mathcal{U}, \mathcal{X}_N \in \mathcal{X}_f, x_0 = x(k)$ 

N is the time horizon and X, U, Xf are polyhedral regions

# 2.2 Transform Quadratic Cost CFTOC into QP

Goal  $\min_{z \in \mathbb{R}^n} \frac{1}{2} z^\top H z + q^\top z + r$  s.t.  $Gz \leq h$ , Az = b2.2.1 Substitute without substitution

Idea Keep state equations as equality constraints

**Define variable**  $z = \begin{bmatrix} x_1^\top & \dots & x_N^\top & u_0^\top & \dots & u_{N-1}^\top \end{bmatrix}^\top$ **Equalities** from system dynamics  $x_{i+1} = Ax_i + Bu_i$ 

$$G_{eq} = \begin{bmatrix} \mathbb{I} & & & -B & \\ -A & \mathbb{I} & & -B & \\ & \ddots & & & -A & \mathbb{I} & \\ & & -A & \mathbb{I} & & -B \end{bmatrix} E_{eq} = \begin{bmatrix} A & \\ 0 & \\ \vdots \\ 0 \end{bmatrix}$$

Inequalities  $G_{in}z \leq w_{in} + E_{in}x(k)$  from  $\mathcal{X} = \{x \mid A_xx \leq$  $b_x$ ,  $\mathcal{U} = \{u \mid A_u u < b_u\}, \mathcal{X}_f = \{x \mid A_f x < b_f\}$ 

Cost Matrix  $\bar{H} = \text{diag}(Q, ..., Q, P, R, ..., R)$ 

Finally the resulting quadratic optimization problem

$$\begin{split} J^{\star}(x(k)) &= \min_{z} \left[ z^{\top} \; x(k)^{\top} \right] \left[ \begin{smallmatrix} \bar{H} & 0 \\ 0 & Q \end{smallmatrix} \right] \left[ z^{\top} \; x(k)^{\top} \right]^{\top} \\ \text{s.t.} \quad G_{in}z &\leq w_{in} + E_{in}x(k) \quad G_{eq}z = E_{eq}x(k) \end{split}$$

#### 2.2.2 Substitute with substitution Idea Substitute the state equations.

Step 1 Rewrite cost as

$$\begin{split} J(x(k)) = & UXXXX \\ = & \begin{bmatrix} U^\top & x(k)^\top \end{bmatrix} \begin{bmatrix} H & F^\top \\ F & Y \end{bmatrix} \begin{bmatrix} U^\top & x(k)^\top \end{bmatrix}^\top \end{split}$$

**Step 2** Rewrite constraints compactly as  $GU \le w + Ex(k)$ Step 3 Rewrite constrained problem as

$$\begin{split} J^{\star}(x(k)) &= \min_{U} \begin{bmatrix} U^{\top} & x(k)^{\top} \end{bmatrix} \begin{bmatrix} H & F^{\top} \\ F & Y \end{bmatrix} \begin{bmatrix} U^{\top} & x(k)^{\top} \end{bmatrix}^{\top} \\ \text{subj. to } GU &< w + Ex(k) \end{split}$$

#### 2.3 Invariance

**Definition 15** (Positively Invariant Set  $\mathcal{O}$ ). For an autonomous or closed-loop system, the set  $\mathcal O$  is positively invariant if:

$$x(k) \in \mathcal{O} \Rightarrow x(k+1) \in \mathcal{O}, \quad \forall k \in \{0, 1, \dots\}$$

**Definition 16** (Maximal Positively Invariant Set  $\mathcal{O}_{\infty}$ ). A set that contains all  $\mathcal{O}$  is the maximal positively invariant set  $\mathcal{O}_{\infty} \subset \mathcal{X}$ **Definition 17** (Pre-Sets). The set of states that in the dynamic system x(k+1) = q(x(k)) in one time step evolves into the target set S is the **pre-set** of  $S \Rightarrow \operatorname{pre}(S) := \{x \mid g(x) \in S\}$ 

**Theorem 6** (Geometric condition for invariance). Set  $\mathcal{O}$  is positively invariant set iff  $\mathcal{O} \subseteq \operatorname{pre}(\mathcal{O}) \Leftrightarrow \operatorname{pre}(\mathcal{O}) \cap \mathcal{O} = \mathcal{O}$ 

 $\bar{x} \notin \operatorname{pre}(\mathcal{O}) \leadsto \bar{x} \in \mathcal{O}, \bar{x} \notin \operatorname{pre}(\mathcal{O}), \text{ thus } \mathcal{O} \text{ not positively}$ invariant **Sufficient** if  $\mathcal{O}$  not pos invar set, then  $\exists \bar{x} \in \mathcal{O}$  s.t  $g(\bar{x}) \notin \mathcal{O}$ 

*Proof.* Necessary if  $\mathcal{O} \not\subset \operatorname{pre}(\mathcal{O})$ , then  $\exists \bar{x} \in \mathcal{O}$  s.t

 $\rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \mathsf{pre}(\mathcal{O}) \text{ thus } \mathcal{O} \notin \mathsf{pre}(\mathcal{O})$ 

**Computing Invariant Sets** 

# **P**re-Set Computation

System with constraints x(k+1) = Ax(k) + Bu(k) $u(k) \in \mathcal{U} := \{u | Gu \le g\}$ and set  $S := \{x | Fx < f\}$  $pre(S) := \{x \mid Ax \in S\}$  $= \{x \mid FAx \leq f\}$ 

loop  $\Omega_{i+1} \leftarrow$  $\operatorname{pre}(\Omega_i) \cap \Omega_i$ if  $\Omega_{i+1} = \Omega_i$  then return  $\mathcal{O}_{\infty} = \Omega_i$ end if

Conceptual Algorithm

first line

 $\Omega_0 \leftarrow \mathcal{X}$ 

(Same but much harder for control invariat sets)

end loop

# **C**onceptual Algorithm

first line  $\Omega_0 \leftarrow \mathcal{X}$  $\Omega_{i+1} \leftarrow \mathsf{pre}(\Omega_i) \cap \Omega_i$ if  $\Omega_{i+1} = \Omega_i$  then return  $\mathcal{O}_{\infty} = \Omega_i$ end if end loop

(Same but much harder for control invariat sets)

# **C**onceptual Algorithm

@decorator() # Example Python code def hello\_world(): # This is a comment

(Same but much harder for control invariat sets)

#### 2.4 Control Invariance

**Definition 18** (Control Invariant Set).  $C \subseteq \mathcal{X}$  control invariant if

$$x(k) \in \mathcal{C} \Rightarrow \exists u(k) \in \mathcal{U} \text{ s.t } q(x(k), u(k)) \in \mathcal{C} \ \forall k$$

**Definition 19** (Maximal Control Invariant Set  $\mathcal{C}_{\infty}$ ). A set that contains all  $\mathcal{C}$  is the maximal positively invariant set  $\mathcal{C}_{\infty} \subset \mathcal{X}$ 

**Intuition** For all states in  $\mathcal{C}_{\infty}$  exists control law s.t constraints are never violated \infty The best any controller could ever do

**Pre-set**  $pre(S) := \{x \mid \exists u \in \mathcal{U} \text{ s.t } g(x, u) \in S\}$ 

Set  $\mathcal{C}$  is control invariant iff:  $\mathcal{C} \subseteq \operatorname{pre}(\mathcal{C}) \Leftrightarrow \operatorname{pre}(\mathcal{C}) \cap \mathcal{C} = \mathcal{C}$ 

#### Control Law from Control Invariant Set

Let C control invariant set for x(k+1) = q(x(k), u(k))Control law  $\kappa(x(k))$  will **guarantee** that system satisfies constraints for all time if:  $q(x, \kappa(x)) \in \mathcal{C} \ \forall x \in \mathcal{C}$ We can use this fact to **synthesize** control law  $\kappa$ with f as any function (including f(x, u) = 0)

$$\kappa(x) := \mathrm{argmin}\{f(x,u) \mid g(x,u) \in \mathcal{C}\}$$

Does not ensure that system will converge Difficult because calculating control invariant sets is hard **MPC** implicitly describes C s.t easy to represent/compute

#### Theorem 7. Minkowski-Weyl

The following statements are equivalent for  $\mathcal{P} \subseteq \mathbb{R}^d$ 

- $\mathcal{P}$  is a polytope and there exists A, b s.t  $\mathcal{P} = \{x \mid Ax < b\}$
- $\mathcal{P}$  finitely generated,  $\exists$  finite set  $\{v_i\}$  s.t  $\mathcal{P} =$  $co(\{v_1,...,v_s\})$

# MOST COMMON Polytopic

#### Lemma 1. Invariant Sets from Lyapunov Functions

If  $V: \mathbb{R}^n \to \mathbb{R}$  is a Lyapunov function for x(k+1) = g(x(k)), then  $Y := \{x \mid V(x) < \alpha\}$  is an invariant set for all  $\alpha > 0$ 

*Proof.* Lyapunov property V(g(x)) - V(x) < 0 implies that once  $V(x(k)) < \alpha$ ,  $V(x(j)) < \alpha$ ,  $\forall j > k \rightarrow \text{Invariance} \quad \Box$ 

**Example System** for x(k+1) = Ax(k) with  $P \succ 0$  that satisfies  $A^{\top}PA - P \prec 0 \rightsquigarrow \text{then } V(x(k)) = x(k)^{\top}Px(k)$ is Lyap, function

Goal – find largest  $\alpha$  s.t set  $Y_{\alpha} \in \mathcal{X}$ 

 $Y_{\alpha} := \{x \mid x^{\top} Px \leq \alpha\} \subset \mathcal{X} := \{x \mid Fx \leq f\}$ Equivalent to  $\max_{\alpha} \alpha$  subj. to  $h_{Y_{\alpha}}(F_i) \leq f_i \ \forall i \in$  $\{1 \dots n\}$ 

# 2.5 Feasibility and Stability

What can go wrong with "standard" MPC?

- No feasibility guarantee, i.e., the MPC problem may not have a solution
- No stability guarantee, i.e., trajectories may not converge to the origin

# **MPC Mathematical Formulation**

$$J^{\star}(x(k)) = \min \sum_{i=0}^{N-1} x_i^{\top} Q x_i + u_i^{\top} R u_i$$

٧3

$$\boldsymbol{J}^{\star}(\boldsymbol{x}(k)) = \min \sum_{i=0}^{N-1} \boldsymbol{x}_i^{\top} \boldsymbol{Q} \boldsymbol{x}_i + \boldsymbol{u}_i^{\top} \boldsymbol{R} \boldsymbol{u}_i$$

subj. to 
$$x_{i+1} = Ax_i + Bu_i$$
 
$$x_0 = x(k) \quad x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}$$

#### Stability of MPC - Main Result

**Theorem 8.** The closed-loop system under the MPC control law  $u_0^{\star}(x)$  is asymptotically stable and the set  $\mathcal{X}_f$  is positive invariant for the system  $x(k+1) = Ax(k) + Bu_0^*(x(k))$ under the following assumptions:

- 1. Stage cost is positive definite, i.e. it is strictly positive and only zero at the origin
- 2. Terminal set is invariant under the local control law  $\kappa_f(x_i)$ :

$$x_{i+1} = Ax_i + B\kappa_f(x_i) \in \mathcal{X}_f \forall x_i \in \mathcal{X}_f$$

All state and input constraints are satisfied in  $\mathcal{X}_f$ :

$$\mathcal{X}_f \in X, \kappa_f(x_i) \in U \forall x_i \in \mathcal{X}_f$$

3. Terminal cost is a continuous Lyapunov function in the terminal set  $\mathcal{X}_f$  and satisfies:

$$I_f(x_{i+1}) - I_f(x_i) \le -I(x_i, \kappa_f(x_i)) \quad \forall x_i \in \mathcal{X}_f$$

#### Finite-horizon MPC may not satisfy constraints for all time! Finite-horizon MPC may not be stable!

- An infinite-horizon provides stability and invariance.
- Infinite-horizon faked by forcing final state into an invariant set for which there exists invariance-inducing controller, whose infinite-horizon cost can be expressed in closed-form.
- Extends to non-linear systems, but compute sets is difficult!

#### 3 Practical-MPC

#### 3.1 Steady-state Target Problem

- Reference is achieved by the target state  $x_s$  if  $z_s = Hx_s = r$
- Target state should be a steady-state, i.e.  $x_s = Ax_s + Bu_s$

$$\begin{array}{c} x_s = Ax_s + Bu_s \\ z_s = Hx_s = r \end{array} \iff \begin{bmatrix} \mathbb{I} - A & -B \\ H & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r \end{bmatrix}$$

 $\nexists$  solution  $\rightarrow \min(Hx_s - r)^{\top} Q_s(Hx_s - r)$  (closest x to r) If  $\exists$  multiple feasible  $u_s \to \text{compute min } u_s^\top R_s u_s$  (cheapest)

$$\min_{U}|z_N - Hx_s|_{P_z}^2 + \sum_{i=1}^{N-1}|z_i - Hx_s|_{Q_z}^2 + |u_i - u_s|_R^2$$

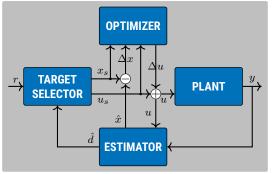
# 3.2 Reference Tracking

$$\Delta x = x - x_s \Rightarrow \Delta x_{k+1} = x_{k+1} - x_s$$

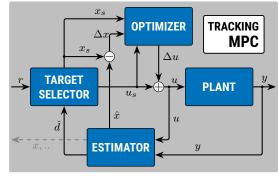
$$\Delta u = u - u_s \Rightarrow = A\Delta x_k + Bu_k - (Ax_s + Bu_s)$$

$$= A\Delta x_k + B\Delta u_k$$

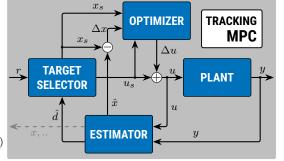
$$\begin{array}{l} G_x x \leq h_x \\ G_u u \leq h_u \end{array} \Rightarrow \begin{array}{l} G_x \Delta x \leq h_x - G_x x_s \\ G_u \Delta u \leq h_u - G_u u_s \end{array}$$



Tracking MPC Diagramm V2.1



Tracking MPC Diagramm V2.1.tight



Assume target feasible with  $x_s \in \mathcal{X}, u_s \in \mathcal{U}$ , choose terminal weight  $V_f(x)$  and constraint  $\mathcal{X}_f$  as in regulation case satisfying

- $\mathcal{X}_f \subseteq \mathcal{X}, Kx \in \mathcal{U} \quad \forall x \in \mathcal{X}_f$
- $V_f(x(k+1)) V_f(x(k)) \le -l(x(k), Kx(k)) \quad \forall x \in \mathcal{X}_f$ If in addition the target reference  $x_s$ ,  $u_s$  is such that
- $x_s \oplus \mathcal{X}_f \subseteq \mathcal{X}, K\Delta x + u_s \in \mathcal{U}, \forall \Delta x \in \mathcal{X}_f$ then CL system converges to target reference

$$x(k) \to x_s, z(k) = Hx(k) \xrightarrow{k \to \infty} r$$

Proof. • Invariance under local ctrol law inherited from regulation case

- Constraint satisfaction provided by extra conditions
  - $x_s \oplus \mathcal{X}_f \subseteq \mathcal{X} \to x \in \mathcal{X} \forall \Delta \in \mathcal{X}_f$
- $K\Delta x + u_s \in \mathcal{U} \forall \Delta x \in \mathcal{X}_f \rightarrow u \in \mathcal{U}$
- Fron asympt stability of the regulation problem:  $\Delta x(k) \xrightarrow{k \to \infty} 0$

$$\Delta x(k) \xrightarrow{k \to \infty}$$

**Terminal set** use  $\mathcal{X}_f^{\text{scaled}} = \alpha \mathcal{X}_f$  (s.t. constraints satisfied)

### 3.3 Reference Tracking without Offset

Approach Model the disturbance, use the measurements and model to estimate the state and disturbance and find control inputs that use the disturbance estimate to remove offset.

#### **Auamented Model**

$$x_{k+1} = Ax_k + Bu_k + B_d d_k$$
$$y_k = Cx_k + C_d d_k$$

Constant disturbance  $d_{k+1} = d_k$ 

Observable iff  $\left[egin{array}{c} A^{-\mathbb{I}} & B_d \\ C & C_d \end{array}
ight]$  has full rank  $(=n_x+n_d)$ 

#### Observer For Augmented Model

$$\begin{bmatrix} \hat{x}_{k+1} \\ \hat{d}_{k+1} \end{bmatrix} = \begin{bmatrix} A & B_d \\ 0 & \mathbb{I} \end{bmatrix} \begin{bmatrix} \hat{x}_k \\ \hat{d}_k \end{bmatrix} + \begin{bmatrix} B \\ 0 \end{bmatrix} u_k + \begin{bmatrix} L_x \\ L_d \end{bmatrix} (C\hat{x}_k + C_d\hat{d}_k - y_k)$$

**Error Dynamics**  $\Rightarrow$  choose L s.t error dynamics converge to 0

$$\left[ \begin{smallmatrix} x_{k+1} - \hat{x}_{k+1} \\ d_{k+1} - \hat{d}_{k+1} \end{smallmatrix} \right] = \left( \left[ \begin{smallmatrix} A & B_d \\ 0 & \mathbb{I} \end{smallmatrix} \right] + \left[ \begin{smallmatrix} L_x \\ L_d \end{smallmatrix} \right] \left[ \begin{smallmatrix} C & C_d \end{smallmatrix} \right] \right) \left[ \begin{smallmatrix} x_k - \hat{x}_k \\ d_k - \hat{d}_k \end{smallmatrix} \right]$$

Lemma 2. Steady-state of an asym. stable observer satisfies:

$$\begin{bmatrix} A - \mathbb{I} & B \\ C & 0 \end{bmatrix} \begin{bmatrix} \hat{x}_{\infty} \\ u_{\infty} \end{bmatrix} = \begin{bmatrix} -B_d \hat{d}_{\infty} \\ y_{\infty} - C_d \hat{d}_{\infty} \end{bmatrix} \text{ (for } n_y = n_d \text{)}$$

 $\Rightarrow$  Observer output  $C\hat{x}_{\infty}+C_{d}\hat{d}_{\infty}$  tracks  $y_{\infty}$  without offset

# 3.4 Offset-free Tracking

**Goal** Track constant reference:  $Hy(k) = z(k) \rightarrow r, k \rightarrow \infty$ 

$$\begin{aligned} x_s &= Ax_s + Bu_s + B_d \hat{d}_{\infty} \\ z_s &= H(Cx_s + C_d \hat{d}_{\infty}) = r \begin{bmatrix} A - I & B \\ HC & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r - HC_d \hat{d} \end{bmatrix} \end{aligned}$$

Theorem 9 (Offset-free Tracking: Main Result). Assuming  $n_d = n_u$ , RHC recursively feasible, unconstrained for  $k \geq j$ , control law  $\kappa(\cdot) = \kappa(\hat{x}(k), \hat{d}(k), r)$  and closed loop system

$$\begin{split} x(k+1) = & Ax(k) + B\kappa(\cdot) + B_d d \\ \hat{x}(k+1) = & (A + L_x C)\hat{x}(k) + (B_d + L_x C_d)\hat{d}(k) \\ & + B\kappa(\cdot) - L_x y(k) \\ \hat{d}(k+1) = & L_d C\hat{x}(k) + (\mathbb{I} + L_d C_d)\hat{d}(k) - L_d y(k) \end{split}$$

converges, then 
$$z(k) = Hy(k) o r$$
 as  $k o \infty$ 

#### 3.5 Soft Constraints

Input constraints are dictated by physical constraints on the actuators and are usually hard

- State/output constraints arise from practical restrictions on the allowed operating range and are rarely hard
- Hard state/output constraints always lead to complications in the controller implementation

# **Soft Constrained MPC Problem Setup**

$$\begin{aligned} \min_{u} \sum_{i=0}^{N-1} x_{i}^{\top}Qx_{i} + u_{i}^{\top}Ru_{i} + l_{\epsilon}(\epsilon_{i}) + x_{N}^{\top}Px_{i} + l_{\epsilon}(\epsilon_{N}) \\ \text{subj. to } x_{i+1} &= Ax_{i} + Bu_{i} \\ H_{x}x_{i} &\leq k_{x} + \epsilon_{i} \\ H_{u}u_{i} &\leq k_{u} \end{aligned}$$

Quadratic penalty 
$$l_{\epsilon}(\epsilon_i) = \epsilon_i^{\top} S \epsilon_i$$
 (e.g  $S = Q$ )

slack variable  $\epsilon_i > 0$ 

Linear Penalty  $v|\epsilon_i|_{1/\infty}$ 

**Requirement on**  $l_{\epsilon}(\epsilon)$  If the original problem has a feasible solution  $z^{\star}$ , then the softened problem should have the same solution  $z^{\star}$ , and  $\epsilon=0$ .

**Theorem 10** (Exact Penalty Funtcion).  $l_{\epsilon}(\epsilon) = v \cdot \epsilon$  satisfies requirement for any  $v > \lambda^{\star} \geq 0$ , where  $\lambda^{\star}$  is optimal Lagrange multiplier for original problem

#### 4 Robust-MPC

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# 5 Implementation

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