

Model Predictive Control

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github.com/silvasta/summary-mpc



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1 Introduction

Requirements for MPC

1. A model of the system
2. A state estimator
3. Define the optimal control problem
4. Set up the optimization problem
5. Get the optimal control sequence (solve the optimization problem)
6. Verify that the closed-loop system performs as desired

1.1 Exact ODE solution of a Linear System

$$x(t) = e^{A^c(t-t_0)}x_0 + \int_{t_0}^t e^{A^c(t-\tau)}B^cu(\tau)d\tau$$

Problem Most physical systems are nonlinear

Idea use First Order Taylor expansion $f(\bar{x}) + \frac{\partial f}{\partial x} \Big|_{\bar{x}} (x - \bar{x})$

1.2 Linearization

$$\begin{aligned} \dot{x}_s &= g(x_s, u_s) = 0 & \Delta \dot{x} &= \dot{x} - \dot{x}_s = A^c \Delta x + B^c \Delta u \\ y_s &= h(x_s, u_s) & \Delta y &= y - y_s = C \Delta x + D \Delta u \end{aligned}$$

$$A^c = \frac{\partial g}{\partial x^T} \Big|_{x_s} B^c = \frac{\partial g}{\partial u^T} \Big|_{u_s} C = \frac{\partial h}{\partial x^T} \Big|_{x_s} D = \frac{\partial h}{\partial u^T} \Big|_{u_s}$$

1.3 Discretization

For general nonlinear systems only approximate discretization methods exist, such as Euler, quality depends on sampling time

Approximation	Notation
$\dot{x}^c \approx \frac{x^c(t+T_s) - x^c(t)}{T_s}$	$x(k) := x^c(t_0 + kT_s)$
	$u(k) := u^c(t_0 + kT_s)$

Exact Discretization of Linear Time-Invariant Models

$$\begin{aligned} x(t_{k+1}) &= \underbrace{e^{A^c T_s}}_{=A} x(t_k) + \underbrace{\int_0^{T_s} e^{A^c(T_s-\tau)} B^c d\tau}_{B=(A^c)^{-1}(A-I)B^c} u(t_k) \\ x(k+N) &= A^N x(k) + \sum_{i=0}^{N-1} A^i B u(k+N-1-i) \end{aligned}$$

1.4 Analysis of LTI Discrete-Time Systems

Controllable if $\text{rank}(C) = n$, $C = [B \quad \dots \quad A^{n-1}B]$

$\forall (x(0), x^*) \exists$ finite time N with inputs \mathcal{U} , s.t. $x(N) = x^*$

Stabilizable iff all uncontrollable modes stable

Observable if $\text{rank}(O) = n$, $[C^T \quad \dots \quad (CA^{n-1})^T]^T$

$\forall x(0) \exists$ finite time N , s.t. the measurements

$y(0), \dots, y(N-1)$ uniquely distinguish initial state $x(0)$

Detectability iff all unobservable modes stable

1.5 Lyapunov

Stability is a property of an **equilibrium point** \bar{x} of a system

Definition 1 (Lyapunov Stability). \bar{x} is **Lyapunov stable** if:

$\forall \epsilon > 0 \exists \delta(\epsilon)$ s.t. $\|x(0) - \bar{x}\| < \delta(\epsilon) \Rightarrow \|x(k) - \bar{x}\| < \epsilon$

Definition 2 (Globally asymptotic stability). If \bar{x} is Lyapunov stable and attractive, i.e., $\lim_{k \rightarrow \infty} \|x(k) - \bar{x}\| = 0$, $\forall x(0)$ then \bar{x} is **globally asymptotic stable**.

Definition 3 (Global Lyapunov function). For $\bar{x} = 0$, function $V : \mathbb{R}^n \rightarrow \mathbb{R}$ is called **Lyapunov function** if it is continuous at the origin, finite $\forall x \in \mathbb{R}^n$,

$$\|x\| \rightarrow \infty \Rightarrow V(x) \rightarrow \infty$$

$$V(x) > 0 \quad \forall x \in \mathbb{R}^n \setminus \{0\} \quad V(0) = 0$$

$$V(g(x)) - V(x) \leq -\alpha(x) \quad \forall x \in \mathbb{R}^n$$

where $\alpha : \mathbb{R}^n \rightarrow \mathbb{R}$ continuous positive definite

Lyapunov Theorem

Theorem 1. If a system admits a Lyapunov function $V(x)$, then $\bar{x} = 0$ is **globally asymptotically stable**.

Theorem 2 (Lyapunov indirect method). For linearization of system around $\bar{x} = 0$ and resulting matrix $A = \frac{\partial g}{\partial x^T} \Big|_{x=0}$ with eigenvalues

$$|\lambda_i| := \begin{cases} \forall i := |\lambda_i| < 1 & x=0 \text{ is asymptotically stable} \\ \exists i := |\lambda_i| > 1 & \text{origin is unstable} \\ \exists i := |\lambda_i| = 1 & \text{no info about stability} \end{cases}$$

Discrete-Time Lyapunov equation

$$A^T P A - P = -Q, \quad Q > 0$$

Theorem 3 (Existence of solution of DT Lyapunov equation). The discrete-time Lyapunov equation (3) has a unique solution $P > 0$ if and only if A has all eigenvalues inside the unit circle, i.e. if and only if the system $x(k+1) = Ax(k)$ is stable.

1.6 Optimal Control

Unconstrained Finite Horizon Control Problem

$$\begin{aligned} J^*(x(0)) &:= \min_U x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i \\ \text{subject to } &x_{i+1} = Ax_i + Bu_i \quad i = 0, \dots, N-1 \\ &x_0 = x(0) \end{aligned}$$

$P \succeq 0$, with $P = P^T$ terminal weight

$Q \succeq 0$, with $Q = Q^T$ state weight

$R \succ 0$, with $R = R^T$ input weight

1.7 Batch Approach

expresses cost function in terms of $x(0)$ and input sequence U

$$\begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} I \\ A \\ \vdots \\ A^N \end{bmatrix} x(0) + \begin{bmatrix} 0 & \dots & 0 \\ B & 0 & 0 \\ AB & B & 0 \\ \vdots & \ddots & 0 \\ A^{N-1}B & \dots & B \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix}$$

$$\bar{Q} := \text{blockdiag}(Q, \dots, Q, P) \quad \bar{R} := \text{blockdiag}(R, \dots, R)$$

Optimal Input (from $\nabla_U J(x(0), U) = 2HU + 2F^T x(0) = 0$)

$$U^*(x(0)) = - \underbrace{(S^u)^T \bar{Q} S^u + \bar{R}}_{H(\text{Hessian})^{-1}} \underbrace{(S^u)^T \bar{Q} S^x}_{F^T} x(0)$$

Optimal Cost

$$J^*(x(0)) = x(0)^T (S_x^T \bar{Q} S_x - S_x^T \bar{Q} S_u (S_u^T \bar{Q} S_u + \bar{R})^{-1} S_u^T \bar{Q} S_x) x(0)$$

1.8 Recursive Approach

uses dynamic programming to solve problem backwards from N

$$J_j^*(x(j)) := \min_{u_j \rightarrow N} x_N^T P x_N + \sum_{i=j}^{N-1} x_i^T Q x_i + u_i^T R u_i$$

Ricatti Equations

RDE - Riccati Difference Equation

$$P_i = A^T P_{i+1} A + Q - A^T P_{i+1} B (B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A$$

RDE - Riccati Difference Equation solved recursively

$$P_i = A^T P_{i+1} A + Q - A^T P_{i+1} B (B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A$$

ARE - Algebraic Riccati Equation solved analytically

$$P_\infty = A^T P_\infty A + Q - A^T P_\infty B (B^T P_\infty B + R)^{-1} B^T P_\infty A$$

From Principle Of Optimality

Optimal Cost-To-Go

$$J_j^*(x_j) = \min_{u_j} I(x_j, u_j) + J_{j+1}^*(x_{j+1}) \quad J_i^*(x_i) = x_i^T P_i x_i$$

Optimal Control Policy

$$u_i^* = F_i x_i = -(B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A \cdot x(i)$$

1.9 Comparison of Batch and Recursive Approaches

Batch optimization returns sequence $U^*(x(0))$ of **numeric values** depending only on $x(0)$, dynamic programming yields **feedback policies** $u_i^* = F_i x_i$ depending on each x_i .

Choice of P

1. Match infinite solution, use ARE

2. Assume no control needed after N , use Lyapunov Equation (makes only sense when asymptotically stable, otherwise P not positive definite)

3. set constraint $x_{i+N} = 0$

1.10 Infinite Horizon LQR

LQR

$$J_\infty^*(x(k)) = \min \sum_{i=0}^{\infty} x_i^T Q x_i + u_i^T R u_i$$

$$\text{subj. to } x_{i+1} = Ax_i + Bu_i, \quad x_0 = x(k)$$

Same u as for finite problem but with ARE Constant Feedback Matrix F_∞ asymptotically stable for. Q,R,stab,detect

1.11 Optimization

A mathematical optimization problem is generally formulated as:

Mathematical Optimization Problem

Decision variable $x \in \mathbb{R}^n$
Objective function $f: \text{dom}(f) \rightarrow \mathbb{R}$
Inequality constraints g_i ($i \in \# \text{constraints}$)
Equality constraints h_i ($i \in \# \text{constraints}$)
Fesabile set $\mathcal{X} := \{x | g(x) \leq 0, h(x) = 0\}$

minimize $f(x)$
subject to:
 $g_i(x) \leq 0$
 $h_i(x) = 0$

Feasible point $x \in \text{dom}(f)$ with $g_i(x) \leq 0, h_i(x) = 0$
Strictly feasible point x with strict inequality $g_i(x) < 0$
Optimal value f^* (or p^*) = $\inf\{f(x) | g_i(x) \leq 0, h_i(x) = 0\}$
 $f^* = +\infty$: OP infeasible, $f^* = -\infty$: OP unbound below
Optimizer set: $\text{argmin}_{x \in \mathcal{X}} f(x) := \{x \in \mathcal{X} | f(x) = f^*\}$

x^* is a **Global Minimum** if $f(x^*) \leq f(x)$
 x^* is a **Local Minimum** if $\exists \epsilon > 0$ s.t. $f(x^*) \leq f(x)$
 $\forall x \in \mathcal{X} \cap B_\epsilon(x^*)$, open ball with center x^* and radius ϵ

Linear Programming minimize $c^T x$ s.t. $Ax - b \geq 0, x \geq 0$
Step 1: $\mathcal{L}(x, \lambda_1, \lambda_2) = c^T x - \lambda_1^T (Ax - b) - \lambda_2^T x, \lambda_i \geq 0$
Step 2: $\inf_{x \in \mathbb{R}^n} \mathcal{L} = \lambda_1^T b$, if $c - A^T \lambda_1 - \lambda_2 = 0$, else $-\infty$

Step 3: Dual, maximize $b^T \lambda$ s.t. $c - A^T \lambda \geq 0, \lambda \geq 0$ (again LP)

Quadratic Programming min ...

1.14 Optimality Conditions

Lagrange Duality

Consider $f^* = \inf_{x \in \mathbb{R}^n} f(x)$ s.t. $g(x) \leq 0, h(x) = 0$

Lagrangian $\mathcal{L}(x, \lambda, \nu) = f(x) + \lambda^T g(x) + \nu^T h(x)$

Dual Function $d(\lambda, \nu) = \inf_{x \in \mathbb{R}^n} \mathcal{L}(x, \lambda, \nu)$

Proposition 1 (Weak Duality). $d(\lambda, \nu) \leq f^*, \forall \lambda \geq 0, \nu \in \mathbb{R}^h$
Definition 14 (Constraint qualification). **Slater's Condition** holds if \exists at least one strictly feasible point \hat{x} ($h(\hat{x}) = 0, g(\hat{x}) < 0$)
Proposition 2 (Strong Duality). If Slater's condition holds and OP is convex $\Rightarrow \exists \lambda \geq 0, \nu \in \mathbb{R}^h$ s.t. $d(\lambda, \nu) = f^*$

KKT Conditions (Karush-Kuhn-Tucker)

Theorem 5 (KKT Conditions). If Slater's condition holds and (1.14) is convex $\rightarrow x^* \in \mathbb{R}^n$ is a minimizer of the primal (1.14) and $(\lambda^* \geq 0, \nu^*) \in \mathbb{R}^{n_g} \times \mathbb{R}^{n_h}$ is a maximizer of the dual \Leftrightarrow is equivalent to the following statements:

KKT-1 (Stationary Lagrangian) $\nabla_x \mathcal{L}(x^*, \lambda^*, \nu^*) = 0$
KKT-2 (primal feasibility) $g(x^*) \leq 0, h(x^*) = 0$
KKT-3 (dual feasibility) $\lambda^*, \nu^* \in \mathbb{R}^{n_h} \geq 0$
KKT-4 (complementary slackness) $\lambda^{*T} g(x^*) = 0, \nu^{*T} h(x^*) = 0$

In addition we have: $\sup_{\lambda \geq 0, \nu \in \mathbb{R}^{n_h}} q(\lambda, \nu) = \inf_{x \in \mathcal{C}} f(x)$

Remark Without Slater, KKT1-4 still implies x^* minimizes (1.14) and λ, ν maximizes dual, but the converse is no longer true. There can be primal-minimizer/dual-maximizer not satisfy KKT.

2 Nominal-MPC

2.1 CFTOC

CFTOC Constrained Finite Time Optimal Control problem

$$J(x(k)) = x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i$$

$$J^*(x(k)) = \min_U I_f(x_N) + \sum_{i=0}^{N-1} I(x_i, u_i)$$

s.t. $x_{i+1} = Ax_i + Bu_i, i = 0, \dots, N-1$
 $x_i \in \mathcal{X}, u_i \in \mathcal{U}, \mathcal{X}_N \in \mathcal{X}_f, x_0 = x(k)$

N is the time horizon and X, U, Xf are polyhedral regions

2.2 Transform Quadratic Cost CFTOC into QP
Goal $\min_{z \in \mathbb{R}^n} \frac{1}{2} z^T H z + q^T z + r$ s.t. $Gz \leq h, Az = b$

2.2.1 Substitute without substitution

Idea Keep state equations as equality constraints

Define variable $z = [x_1^T \dots x_N^T u_0^T \dots u_{N-1}^T]^T$

Equalities from system dynamics $x_{i+1} = Ax_i + Bu_i$

$$G_{eq} = \left[\begin{array}{c|c} \begin{matrix} \mathbb{I} & & & \\ -A & \mathbb{I} & & \\ & \ddots & \ddots & \\ & & -A & \mathbb{I} \end{matrix} & \begin{matrix} -B \\ & & & -B \end{matrix} \end{array} \right] E_{eq} = \begin{bmatrix} A \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

Inequalities $G_{in} z \leq w_{in} + E_{in} x(k)$ from $\mathcal{X} = \{x | A_x x \leq b_x\}, \mathcal{U} = \{u | A_u u \leq b_u\}, \mathcal{X}_f = \{x | A_f x \leq b_f\}$

$$G_{in} = \left[\begin{array}{cc|cc} 0 & 0 & & \\ A_x & 0 & & \\ & \ddots & \ddots & \\ & & A_x & A_f & 0 & 0 \\ 0 & & & A_u & \ddots & \\ & \ddots & 0 & & A_u & A_u \end{array} \right] w_{in} = \begin{bmatrix} b_x \\ b_x \\ \vdots \\ b_x \\ b_f \\ b_u \\ \vdots \\ b_u \end{bmatrix}$$

$$E_{in} = \begin{bmatrix} -A_x \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

Cost Matrix $\bar{H} = \text{diag}(Q, \dots, Q, P, R, \dots, R)$

Finally the resulting quadratic optimization problem

$$J^*(x(k)) = \min_z \begin{bmatrix} z^T & x(k)^T \end{bmatrix} \begin{bmatrix} \bar{H} & 0 \\ 0 & Q \end{bmatrix} \begin{bmatrix} z^T & x(k)^T \end{bmatrix}^T$$

s.t. $G_{in} z \leq w_{in} + E_{in} x(k)$ $G_{eq} z = E_{eq} x(k)$

2.2.2 Substitute with substitution

Idea Substitute the state equations.

Step 1 Rewrite cost as

$$J(x(k)) = U^T X X X X$$

$$= \begin{bmatrix} U^T & x(k)^T \end{bmatrix} \begin{bmatrix} H & F^T \\ F & Y \end{bmatrix} \begin{bmatrix} U^T & x(k)^T \end{bmatrix}^T$$

Step 2 Rewrite constraints compactly as $GU \leq w + Ex(k)$

Step 3 Rewrite constrained problem as

$$J^*(x(k)) = \min_U \begin{bmatrix} U^T & x(k)^T \end{bmatrix} \begin{bmatrix} H & F^T \\ F & Y \end{bmatrix} \begin{bmatrix} U^T & x(k)^T \end{bmatrix}^T$$

subj. to $GU \leq w + Ex(k)$

2.3 Invariance

Definition 15 (Positively Invariant Set \mathcal{O}). For an autonomous or closed-loop system, the set \mathcal{O} is positively invariant if:

$$x(k) \in \mathcal{O} \Rightarrow x(k+1) \in \mathcal{O}, \quad \forall k \in \{0, 1, \dots\}$$

Definition 16 (Maximal Positively Invariant Set \mathcal{O}_∞). A set that contains all \mathcal{O} is the maximal positively invariant set $\mathcal{O}_\infty \subset \mathcal{X}$
Definition 17 (Pre-Sets). The set of states that in the dynamic system $x(k+1) = g(x(k))$ in one time step evolves into the target set \mathcal{S} is the **pre-set** of $\mathcal{S} \Rightarrow \text{pre}(\mathcal{S}) := \{x | g(x) \in \mathcal{S}\}$

Theorem 6 (Geometric condition for invariance). Set \mathcal{O} is positively invariant set iff $\mathcal{O} \subseteq \text{pre}(\mathcal{O}) \Leftrightarrow \text{pre}(\mathcal{O}) \cap \mathcal{O} = \mathcal{O}$

Proof. Necessary if $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$, then $\exists \bar{x} \in \mathcal{O}$ s.t. $\bar{x} \notin \text{pre}(\mathcal{O}) \rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \text{pre}(\mathcal{O})$, thus \mathcal{O} not positively invariant

Sufficient if \mathcal{O} not pos invar set, then $\exists \bar{x} \in \mathcal{O}$ s.t. $g(\bar{x}) \notin \mathcal{O} \rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \text{pre}(\mathcal{O})$ thus $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$ \square

Computing Invariant Sets

Conceptual Algorithm

```
first line
 $\Omega_0 \leftarrow \mathcal{X}$ 
loop
   $\Omega_{i+1} \leftarrow$ 
   $\text{pre}(\Omega_i) \cap \Omega_i$ 
  if  $\Omega_{i+1} = \Omega_i$  then
    return
  end if
 $\mathcal{O}_\infty = \Omega_i$ 
end loop
```

(Same but much harder for control invariant sets)

Conceptual Algorithm

```
first line
 $\Omega_0 \leftarrow \mathcal{X}$ 
loop
   $\Omega_{i+1} \leftarrow \text{pre}(\Omega_i) \cap \Omega_i$ 
  if  $\Omega_{i+1} = \Omega_i$  then
    return  $\mathcal{O}_\infty = \Omega_i$ 
  end if
end loop
```

(Same but much harder for control invariant sets)

Conceptual Algorithm

```
@decorator()
# Example Python code
def hello_world():
  # This is a comment
  print("Hello, World!")
```

(Same but much harder for control invariant sets)

2.4 Control Invariance

Definition 18 (Control Invariant Set). $\mathcal{C} \subseteq \mathcal{X}$ control invariant if

$$x(k) \in \mathcal{C} \Rightarrow \exists u(k) \in \mathcal{U} \text{ s.t. } g(x(k), u(k)) \in \mathcal{C} \quad \forall k$$

Definition 19 (Maximal Control Invariant Set \mathcal{C}_∞). A set that contains all \mathcal{C} is the maximal positively invariant set $\mathcal{C}_\infty \subset \mathcal{X}$
Intuition For all states in \mathcal{C}_∞ exists control law s.t constraints are never violated \rightsquigarrow **The best any controller could ever do**

Pre-set $\text{pre}(\mathcal{S}) := \{x | \exists u \in \mathcal{U} \text{ s.t. } g(x, u) \in \mathcal{S}\}$

Set \mathcal{C} is control invariant iff: $\mathcal{C} \subseteq \text{pre}(\mathcal{C}) \Leftrightarrow \text{pre}(\mathcal{C}) \cap \mathcal{C} = \mathcal{C}$

Control Law from Control Invariant Set

Let \mathcal{C} control invariant set for $x(k+1) = g(x(k), u(k))$
Control law $\kappa(x(k))$ will **guarantee** that system satisfies constraints **for all time** if: $g(x, \kappa(x)) \in \mathcal{C} \forall x \in \mathcal{C}$
We can use this fact to **synthesize** control law κ with f as any function (including $f(x, u) = 0$)

$$\kappa(x) := \operatorname{argmin}\{f(x, u) \mid g(x, u) \in \mathcal{C}\}$$

Does not ensure that system will converge

Difficult because calculating control invariant sets is hard
MPC implicitly describes \mathcal{C} s.t easy to represent/compute

Theorem 7. Minkowski-Weyl

The following statements are equivalent for $\mathcal{P} \subseteq \mathbb{R}^d$

- \mathcal{P} is a polytope and there exists A, b s.t $\mathcal{P} = \{x \mid Ax \leq b\}$
- \mathcal{P} finitely generated, \exists finite set $\{v_i\}$ s.t $\mathcal{P} = \operatorname{co}(\{v_1, \dots, v_s\})$

MOST COMMON Polytopic

1

Lemma 1. Invariant Sets from Lyapunov Functions

If $V: \mathbb{R}^n \rightarrow \mathbb{R}$ is a Lyapunov function for $x(k+1) = g(x(k))$, then $Y := \{x \mid V(x) \leq \alpha\}$ is an invariant set for all $\alpha \geq 0$

Proof. Lyapunov property $V(g(x)) - V(x) < 0$ implies that once $V(x(k)) \leq \alpha$, $V(x(j)) < \alpha, \forall j \geq k \rightarrow$ Invariance \square

Example System for $x(k+1) = Ax(k)$ with $P \succ 0$ that satisfies $A^\top P A - P \prec 0 \rightsquigarrow$ then $V(x(k)) = x(k)^\top P x(k)$ is Lyap. function

Goal - find largest α s.t set $Y_\alpha \in \mathcal{X}$

$Y_\alpha := \{x \mid x^\top P x \leq \alpha\} \subset \mathcal{X} := \{x \mid Fx \leq f\}$
Equivalent to $\max_\alpha \alpha$ subj. to $h_{Y_\alpha}(F_i) \leq f_i \forall i \in \{1 \dots n\}$

2.5 Feasibility and Stability

What can go wrong with "standard" MPC?

- No feasibility guarantee, i.e., the MPC problem may not have a solution
- No stability guarantee, i.e., trajectories may not converge to the origin

MPC Mathematical Formulation

V1

$$J^*(x(k)) = \min \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

V3

$$J^*(x(k)) = \min \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

$$\text{subj. to } x_{i+1} = Ax_i + Bu_i \\ x_0 = x(k) \quad x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}$$

Stability of MPC - Main Result

Theorem 8. The closed-loop system under the MPC control law $u_0^*(x)$ is asymptotically stable and the set \mathcal{X}_f is positive invariant for the system $x(k+1) = Ax(k) + Bu_0^*(x(k))$ under the following assumptions:

- Stage cost is positive definite, i.e. it is strictly positive and only zero at the origin
- Terminal set is **invariant** under the local control law $\kappa_f(x_i)$:

$$x_{i+1} = Ax_i + B\kappa_f(x_i) \in \mathcal{X}_f \forall x_i \in \mathcal{X}_f$$

All state and input **constraints are satisfied** in \mathcal{X}_f :

$$\mathcal{X}_f \in X, \kappa_f(x_i) \in U \forall x_i \in \mathcal{X}_f$$

- Terminal cost is a continuous **Lyapunov function** in the terminal set \mathcal{X}_f and satisfies:

$$I_f(x_{i+1}) - I_f(x_i) \leq -I(x_i, \kappa_f(x_i)) \quad \forall x_i \in \mathcal{X}_f$$

Finite-horizon MPC may not satisfy constraints for all time!
Finite-horizon MPC may not be stable!

- An infinite-horizon provides stability and invariance.
- Infinite-horizon *faked* by forcing final state into an invariant set for which there exists invariance-inducing controller, whose infinite-horizon cost can be expressed in closed-form.
- Extends to non-linear systems, but compute sets is difficult!

3 Practical-MPC

3.1 Steady-state Target Problem

- Reference is achieved by the target state x_s if $z_s = Hx_s = r$
- Target state should be a steady-state, i.e. $x_s = Ax_s + Bu_s$

$$x_s = Ax_s + Bu_s \iff \begin{bmatrix} \mathbb{I} - A & -B \\ H & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r \end{bmatrix}$$

\nexists solution $\rightarrow \min (Hx_s - r)^\top Q_s (Hx_s - r)$ (closest x to r)

If \exists multiple feasible $u_s \rightarrow$ compute $\min u_s^\top R_s u_s$ (cheapest)

$$\min_U |z_N - Hx_s|_{P_z}^2 + \sum_{i=1}^{N-1} |z_i - Hx_s|_{Q_z}^2 + |u_i - u_s|_{R}^2$$

3.2 Reference Tracking

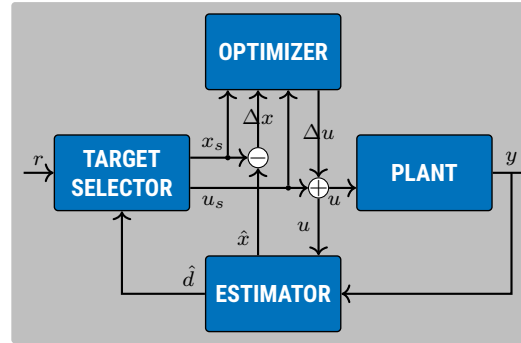
$$\Delta x = x - x_s \Rightarrow \Delta x_{k+1} = x_{k+1} - x_s \\ \Delta u = u - u_s = A\Delta x_k + Bu_k - (Ax_s + Bu_s) \\ = A\Delta x_k + B\Delta u_k$$

$$\begin{matrix} G_x x \leq h_x \\ G_u u \leq h_u \end{matrix} \Rightarrow \begin{matrix} G_x \Delta x \leq h_x - G_x x_s \\ G_u \Delta u \leq h_u - G_u u_s \end{matrix}$$

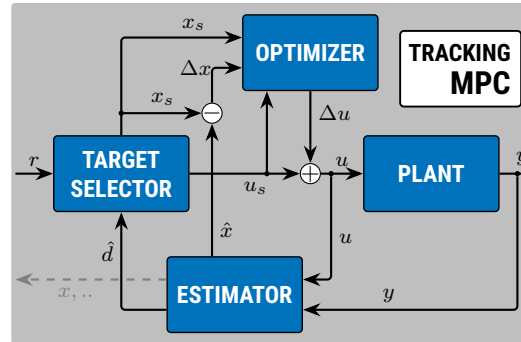
```
def get_next_u(y: Measurement, r: Reference):
    """
    System handler for offset-free tracking
    """
    # approximate state, disturbance
    x, d = estimator(y)
    # find steady state und generate delta
    x_s, u_s = target_selector(x, r, d)
```

```
x_delta = x - x_s
# call solver with new parameter
u_delta = mpc_regulator(x_delta, x_s, u_s)
u = u_delta + u_s

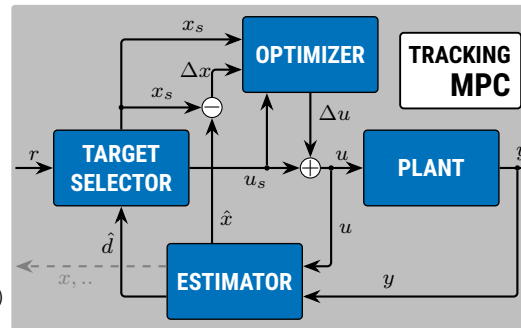
return u
```



Tracking MPC Diagramm V2.1



Tracking MPC Diagramm V2.1.tight



Assume target feasible with $x_s \in \mathcal{X}, u_s \in \mathcal{U}$, choose terminal weight $V_f(x)$ and constraint \mathcal{X}_f as in regulation case satisfying

- $\mathcal{X}_f \subseteq \mathcal{X}, Kx \in \mathcal{U} \quad \forall x \in \mathcal{X}_f$
- $V_f(x(k+1)) - V_f(x(k)) \leq -l(x(k), Kx(k)) \quad \forall x \in \mathcal{X}_f$

If in addition the target reference x_s, u_s is such that

- $x_s \oplus \mathcal{X}_f \subseteq \mathcal{X}, K\Delta x + u_s \in \mathcal{U}, \quad \forall \Delta x \in \mathcal{X}_f$

then CL system converges to target reference

$$x(k) \rightarrow x_s, z(k) = Hx(k) \xrightarrow{k \rightarrow \infty} r$$

Proof. • Invariance under local ctrl law inherited from regulation case

- Constraint satisfaction provided by extra conditions

$$\begin{aligned} - x_s \oplus \mathcal{X}_f &\subseteq \mathcal{X} \rightarrow x \in \mathcal{X} \forall \Delta x \in \mathcal{X}_f \\ - K\Delta x + u_s &\in \mathcal{U} \forall \Delta x \in \mathcal{X}_f \rightarrow u \in \mathcal{U} \end{aligned}$$

- From asympt stability of the regulation problem:

$$\Delta x(k) \xrightarrow{k \rightarrow \infty} 0$$

Terminal set use $\mathcal{X}_f^{\text{scaled}} = \alpha \mathcal{X}_f$ (s.t. constraints satisfied)

3.3 Reference Tracking without Offset

Approach Model the disturbance, use the measurements and model to estimate the state and disturbance and find control inputs that use the disturbance estimate to remove offset.

Augmented Model

$$\begin{aligned} x_{k+1} &= Ax_k + Bu_k + B_d d_k \\ y_k &= Cx_k + C_d d_k \end{aligned}$$

Constant disturbance $d_{k+1} = d_k$

Observable iff $\begin{bmatrix} A - \mathbb{I} & B_d \\ C & C_d \end{bmatrix}$ has full rank ($= n_x + n_d$)

Observer For Augmented Model

$$\begin{bmatrix} \hat{x}_{k+1} \\ \hat{d}_{k+1} \end{bmatrix} = \begin{bmatrix} A & B_d \\ 0 & \mathbb{I} \end{bmatrix} \begin{bmatrix} \hat{x}_k \\ \hat{d}_k \end{bmatrix} + \begin{bmatrix} B \\ 0 \end{bmatrix} u_k + \begin{bmatrix} L_x \\ L_d \end{bmatrix} (Cx_k + C_d \hat{d}_k - y_k)$$

Error Dynamics \Rightarrow choose L s.t error dynamics converge to 0

$$\begin{bmatrix} x_{k+1} - \hat{x}_{k+1} \\ d_{k+1} - \hat{d}_{k+1} \end{bmatrix} = \left(\begin{bmatrix} A & B_d \\ 0 & \mathbb{I} \end{bmatrix} + \begin{bmatrix} L_x \\ L_d \end{bmatrix} \begin{bmatrix} C & C_d \end{bmatrix} \right) \begin{bmatrix} x_k - \hat{x}_k \\ d_k - \hat{d}_k \end{bmatrix}$$

Lemma 2. Steady-state of an asym. stable observer satisfies:

$$\begin{bmatrix} A - \mathbb{I} & B \\ C & 0 \end{bmatrix} \begin{bmatrix} \hat{x}_\infty \\ u_\infty \end{bmatrix} = \begin{bmatrix} -B_d \hat{d}_\infty \\ y_\infty - C_d \hat{d}_\infty \end{bmatrix} \quad (\text{for } n_y = n_d)$$

\Rightarrow Observer output $C\hat{x}_\infty + C_d \hat{d}_\infty$ tracks y_∞ without offset

3.4 Offset-free Tracking

Goal Track constant reference: $Hy(k) = z(k) \rightarrow r, k \rightarrow \infty$

$$\begin{aligned} x_s &= Ax_s + Bu_s + B_d \hat{d}_\infty \\ z_s &= H(Cx_s + C_d \hat{d}_\infty) = r \end{aligned} \quad \begin{bmatrix} A - \mathbb{I} & B \\ HC & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r - HC_d \hat{d} \end{bmatrix}$$

Theorem 9 (Offset-free Tracking: Main Result). Assuming $n_d = n_y$, RHC recursively feasible, unconstrained for $k \geq j$, control law $\kappa(\cdot) = \kappa(\hat{x}(k), \hat{d}(k), r)$ and closed loop system

$$\begin{aligned} x(k+1) &= Ax(k) + B\kappa(\cdot) + B_d d \\ \hat{x}(k+1) &= (A + L_x C)\hat{x}(k) + (B_d + L_x C_d)\hat{d}(k) \\ &\quad + B\kappa(\cdot) - L_x y(k) \\ \hat{d}(k+1) &= L_d C\hat{x}(k) + (\mathbb{I} + L_d C_d)\hat{d}(k) - L_d y(k) \end{aligned}$$

converges, then $z(k) = Hy(k) \rightarrow r$ as $k \rightarrow \infty$

3.5 Soft Constraints

Input constraints are dictated by physical constraints on the actuators and are usually hard

- State/output constraints arise from practical restrictions on the allowed operating range and are **rarely hard**

- Hard state/output constraints always lead to **complications in the controller implementation**

Soft Constrained MPC Problem Setup

$$\min_u \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i + l_\epsilon(\epsilon_i) + x_N^\top P x_i + l_\epsilon(\epsilon_N)$$

$$\begin{aligned} \text{subj. to } & x_{i+1} = A x_i + B u_i \\ & H_x x_i \leq k_x + \epsilon_i \\ & H_u u_i \leq k_u \\ & \text{slack variable } \epsilon_i \geq 0 \end{aligned}$$

Quadratic penalty $l_\epsilon(\epsilon_i) = \epsilon_i^\top S \epsilon_i$ (e.g $S = Q$)

Linear Penalty $v|\epsilon_i|_{1/\infty}$

Requirement on $l_\epsilon(\epsilon)$ If the original problem has a feasible solution z^* , then the softened problem should have the same solution z^* , and $\epsilon = 0$.

Theorem 10 (Exact Penalty Function). $l_\epsilon(\epsilon) = v \cdot \epsilon$ satisfies requirement for any $v > \lambda^* \geq 0$, where λ^* is optimal Lagrange multiplier for original problem

4 Robust-MPC

Uncertain System $x(k+1) = g(x(k), u(k), w(k); \theta)$

4.1 Robust Invariance

Definition 20 (Robust Positive Invariant Set \mathcal{O}^W). For the autonomous system $x(k+1) = g(x(k), w(k))$, the set \mathcal{O}^W is robust positive invariant if:

$$x \in \mathcal{O}^W \Rightarrow g(x, w) \in \mathcal{O}^W, \quad \forall w \in W$$

Given set Ω and dynamic system $x(k+1) = g(x(k), w(k))$,

$$\text{pre}^W(\Omega) := \{x \mid g(x, w) \in \Omega \forall w \in W\}$$

Definition 21 (Robust Pre-Sets). The set of states that in the dynamic system $x(k+1) = g(x(k), w(k))$ for all disturbance $w \in W$ in one time step evolves into the target set Ω is the **pre-set** of $\Omega \Rightarrow \text{pre}^W(\Omega) := \{x \mid g(x, w) \in \Omega \forall w \in W\}$

Computing Robust Pre-Sets for Linear Systems

System $Ax(k) + w(k)$, set $\Omega := \{x \mid Fx \leq f\}$

$$\begin{aligned} \text{pre}^W(\Omega) &= \{x \mid FAx + Fw \leq f\} \\ &= \{x \mid FAx \leq f - \max_{w \in W} Fw\} \\ &= \{x \mid FAx \leq f - h_{W^i}(F)\} \end{aligned}$$

where $h_{W^i}(F)$ is the support function

Theorem 11 (Geometric condition for robust invariance). Set \mathcal{O}^W is robust positive invariant iff $\mathcal{O}^W \subseteq \text{pre}^W(\mathcal{O}^W)$

Definition 22 (Minkowski Sum). For $A, B \subset \mathbb{R}^n$, the Minkowski Sum is $A \oplus B := \{x + y \mid x \in A, y \in B\}$

Definition 23 (Pontryagin Difference). For $A, B \subset \mathbb{R}^n$, the Pontryagin Difference is $A \ominus B := \{x \mid x + e \in A, \forall e \in B\}$

4.2 Impact of Bounded Additive Noise

Defining a Cost to Minimize Expected value, worst case, max W nominal case $w=0$

4.3 Robust Constraint Satisfaction

The idea: Compute a set of tighter constraints such that if the nominal system meets these constraints, then the uncertain system will too. We then do MPC on the nominal system.

Goal: Ensure that constraints are satisfied for the MPC sequence.

Terminal State Constraint

...is called disturbance reachable set,

4.4 Robust open loop MPC

$$\begin{aligned} \min_U & \left[l_f(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i) \right] \\ \text{subj. to } & x_{i+1} = A x_i + B u_i \\ & x_i \in \mathcal{X} \ominus \left(\bigoplus_{j=0}^{i-1} A^j \mathcal{W} \right), \quad u_i \in \mathcal{U} \\ & x_0 = x(k), \quad x_N \in \mathcal{X}_f \ominus \left(\bigoplus_{j=0}^{N-1} A^j \mathcal{W} \right) \end{aligned}$$

4.5 Robust closed loop MPC

Idea Separate the available control authority into two parts:

1. $z(k+1) = Az(k) + Bv(k)$ steers noise-free *nominal* system to origin

2. $u_i = K(x_i - z_i) + v_i$ compensates for deviations, i.e. a *tracking* controller, to keep the real trajectory close to the nominal system.

\Rightarrow We fix the linear feedback controller K offline, and optimize over the nominal inputs $\{v_0, \dots, v_{N-1}\}$ and nominal trajectory $\{z_0, \dots, z_N\}$, which results in a convex problem.

Minimum Robust Invariant Set

$$F_\infty = \bigoplus_{j=0}^\infty A_K^j \mathcal{W}, F_0 := \{0\} \Rightarrow F_n = F_{n+1} = F_\infty$$

4.6 Robust Constraint-Tightening MPC

$$\begin{aligned} \min_{Z, V} & \sum_{i=0}^{N-1} l(z_i, v_i) + l_f(z_N) \\ \text{subj. to } & z_{i+1} = A z_i + B v_i \\ & z_i \in \mathcal{X} \ominus \mathcal{F}_i \\ & u_i \in \mathcal{U} \ominus K(\mathcal{F}_i) \\ & z_N \in \mathcal{X}_f \ominus \mathcal{N}_N \\ & z_0 = x(k) \\ & F_i := \mathcal{W} \oplus A_K \mathcal{W} \oplus \dots \oplus A_K^i \mathcal{W} \\ & \text{Control Law } u(k) = v_0^* + K(x(k) - z_0) = v_0^* \end{aligned}$$

Motivation can robustly ensure constraint satisfactkon at each time step

Note need terminal set \mathcal{X}_f that is robust invariant under tube controller K

4.7 Robust Tube MPC

Idea Ignore noise and plan the nominal trajectory, bound maximum error at any time with RPI set $\mathcal{E} : \epsilon_i \in \mathcal{E} \epsilon_{i+1} \in \mathcal{E}$

Ideally \mathcal{E} is selected as the minimum RPI set F_∞

We know that the real trajectory stays 'nearby' the nominal one because we plan to apply the controller in the future (we won't actually do this, but it's a valid sub-optimal plan)

We must ensure that all possible state trajectories satisfy the constraints This is now equivalent to ensuring that (address input constraints later)

What do we need to make this work?

- Compute the set \mathcal{E} that the error will remain inside

Previously we wanted the **maximum robust invariant set**, or the largest set in which our terminal control law works.

We now want the **minimum robust invariant set**, or the smallest set that the state will remain inside despite the noise.

- Modify constraints on nominal trajectory $\{z_i\}$

$$x_i \in z_i \oplus \mathcal{E} = \{z_i + e \mid e \in \mathcal{E}\}$$

- Formulate as convex optimization problem

BOX

... and then prove that

- Constraints are robustly satisfied

- The closed-loop system is robustly stable

Tube MPC

$$\begin{aligned} \text{Feasible set: } \mathcal{Z}(x_0) &:= \begin{cases} z_{i+1} &= A z_i + B v_i \\ z_i &\in \mathcal{X} \ominus \mathcal{E} \\ v_i &\in \mathcal{U} \ominus K \mathcal{E} \\ z_N &\in \mathcal{X}_f \\ x_0 &\in z_0 \oplus \mathcal{E} \end{cases} \\ \text{Cost function: } J(Z, V) &:= \sum_{i=0}^{N-1} l(z_i, v_i) + l_f(z_N) \\ \text{Optimization: } (V^*(x_0), Z^*(x_0)) &= \underset{\arg \min_{V, Z}}{\text{argmin}} \{ J(Z, V) \mid (Z, V) \in \mathcal{Z}(x_0) \} \\ \text{Control law: } \mu_{\text{tube}}(x) &:= K(x - z_0^*(x)) + v_0^*(x) \end{aligned}$$

$$\begin{aligned} \text{Feasible set: } \mathcal{Z}(x_0) &:= \begin{cases} z_{i+1} &= A z_i + B v_i \\ z_i &\in \mathcal{X} \ominus \mathcal{E} \\ v_i &\in \mathcal{U} \ominus K \mathcal{E} \\ z_N &\in \mathcal{X}_f \\ x_0 &\in z_0 \oplus \mathcal{E} \end{cases} \\ \text{Cost function: } J(Z, V) &:= \sum_{i=0}^{N-1} l(z_i, v_i) + l_f(z_N) \\ \text{Optimization: } (V^*(x_0), Z^*(x_0)) &= \underset{\arg \min_{V, Z}}{\text{argmin}} \{ J(Z, V) \mid (Z, V) \in \mathcal{Z}(x_0) \} \\ \text{Control law: } \mu_{\text{tube}}(x) &:= K(x - z_0^*(x)) + v_0^*(x) \end{aligned}$$

ASSUMPTIONS

Theorem 12 (Robust Invariance of Tube MPC). The set $\mathcal{Z} := \{x \mid \mathcal{Z}(x) \neq \emptyset\}$ is a robust invariant set of the system $x(k+1) = Ax(k) + B\mu_{\text{tube}}(x(k)) + w(k)$ subject to the constraints $x, u \in \mathcal{X} \times \mathcal{U}$.

Theorem 13 (Robust Stability of Tube MPC). The state $x(k)$ of the system $x(k+1) = Ax(k) + B\mu_{\text{tube}}(x(k)) + w(k)$ converges to the limit of the set \mathcal{E} .

Putting it all together: Tube MPC

To implement tube MPC:

– **Offline** –

1. Choose a stabilizing controller K so that $A + BK$ is (Schur) stable

2. Compute the minimal robust invariant set $E = F_\infty$ for the system $x(k+1) = (A + BK)x(k) + w(k), w \in W^1$

3. Compute the tightened constraints $\bar{\mathcal{X}} := \mathcal{X} \ominus \mathcal{E}, \bar{\mathcal{U}} := \mathcal{U} \ominus K \mathcal{E}$

4. Choose terminal weight function l_f and constraint \mathcal{X}_f satisfying assumptions*

– **Online** –

1. Measure / estimate state x

2. Solve the problem $(V^*(x_0), Z^*(x_0)) = \underset{\arg \min_{V, Z}}{\text{argmin}} \{ J(Z, V) \mid (Z, V) \in \mathcal{Z}(x_0) \}$

3. Set the input to $u = K(x - z_0^*(x)) + v_0^*(x)$

5 Implementation

Two options:

- Iterative optimization methods

- Explicit solution

EXPLICIT:

The CFTOC problem is a **multiparametric quadratic program (mp-QP)**

Let $I := 1, \dots, m$ be the set of constraint indices.

Definition 24 (Active Set). $A(x)$ and it's complement $NA(x)$

$$\begin{aligned} A(x) &:= \{j \in I : G_j z^*(x) - S_j x = w_j\} \\ NA(x) &:= \{j \in I : G_j z^*(x) - S_j x < w_j\} \end{aligned}$$

Definition 25 (Critical Region). CR_A is set of parameters x for which set $A \subseteq I$ of constraints i active at the optimum. For given $\bar{x} \in \mathcal{K}^*$ let $(A, NA) := (A(\bar{x}), NA(\bar{x}))$. Then

$$CR_A := \{x \in \mathcal{K}^* : A(x) = A\} \quad (\text{states share active set})$$

Online evaluation: Point location

Sequential search

Logarithmic search

OPTIMIZATION

L-Smooth

(UN-)CONSTRAINED OPTIMIZATION

Projected Gradient Method