

Model Predictive Control

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github.com/silvasta/summary-mpc



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1 Introduction

Requirements for MPC

1. A model of the system
2. A state estimator

3. Define the optimal control problem
4. Set up the optimization problem
5. Get the optimal control sequence (solve the optimization problem)
6. Verify that the closed-loop system performs as desired

1.1 Exact ODE solution of a Linear System

$$x(t) = e^{A^c(t-t_0)}x_0 + \int_{t_0}^t e^{A^c(t-\tau)}B^c u(\tau)d\tau$$

Problem Most physical systems are nonlinear

Idea use First Order Taylor expansion $f(\bar{x}) + \frac{\partial f}{\partial x^T} \Big|_{\bar{x}} (x - \bar{x})$

1.2 Linearization

$$\begin{aligned} \dot{x}_s &= g(x_s, u_s) = 0 & \Delta \dot{x} &= \dot{x} - \dot{x}_s = A^c \Delta x + B^c \Delta u \\ y_s &= h(x_s, u_s) & \Delta y &= y - y_s = C \Delta x + D \Delta u \end{aligned}$$

$$A^c = \frac{\partial g}{\partial x^T} \Big|_{x_s} \quad B^c = \frac{\partial g}{\partial u^T} \Big|_{u_s} \quad C = \frac{\partial h}{\partial x^T} \Big|_{x_s} \quad D = \frac{\partial h}{\partial u^T} \Big|_{u_s}$$

1.3 Discretization

For general nonlinear systems only approximate discretization methods exist, such as Euler, quality depends on sampling time

Approximation

$$\dot{x}^c \approx \frac{x^c(t+T_s) - x^c(t)}{T_s} \quad \begin{array}{l} x(k) := x^c(t_0 + kT_s) \\ u(k) := u^c(t_0 + kT_s) \end{array}$$

Notation

Exact Discretization of Linear Time-Invariant Models

$$x(t_{k+1}) = \underbrace{e^{A^c T_s}}_{=A} x(t_k) + \underbrace{\int_0^{T_s} e^{A^c(T_s-\tau)} B^c d\tau}_{B=(A^c)^{-1}(A-I)B^c} u(t_k)$$

$$x(k+N) = A^N x(k) + \sum_{i=0}^{N-1} A^i B u(k+N-1-i)$$

1.4 Analysis of LTI Discrete-Time Systems

Controllable if $\text{rank}(C) = n$, $C = [B \ \dots \ A^{n-1}B]$

$\forall (x(0), x^*) \exists$ finite time N with inputs \mathcal{U} , s.t. $x(N) = x^*$

Stabilizable iff all uncontrollable modes stable

Observable if $\text{rank}(O) = n$, $[C^T \ \dots \ (CA^{n-1})^T]^T$

$\forall x(0) \exists$ finite time N , s.t. the measurements

$y(0), \dots, y(N-1)$ uniquely distinguish initial state $x(0)$

Detectability iff all unobservable modes stable

1.5 Lyapunov

Stability is a property of an **equilibrium point** \bar{x} of a system

Definition 1 (Lyapunov Stability). \bar{x} is **Lyapunov stable** if:

$\forall \epsilon > 0 \exists \delta(\epsilon)$ s.t. $\|x(0) - \bar{x}\| < \delta(\epsilon) \rightarrow \|x(k) - \bar{x}\| < \epsilon$

Definition 2 (Globally asymptotic stability). If \bar{x} is Lyapunov stable and attractive, i.e., $\lim_{k \rightarrow \infty} \|x(k) - \bar{x}\| = 0$, $\forall x(0)$ then \bar{x} is **globally asymptotic stable**.

Definition 3 (Global Lyapunov function). For $\bar{x} = 0$, function $V : \mathbb{R}^n \rightarrow \mathbb{R}$ is called **Lyapunov function** if it is continuous at the origin, finite $\forall x \in \mathbb{R}^n$,

$$\|x\| \rightarrow \infty \Rightarrow V(x) \rightarrow \infty$$

$$V(x) > 0 \quad \forall x \in \mathbb{R}^n \setminus \{0\} \quad V(0) = 0$$

$$V(g(x)) - V(x) \leq -\alpha(x) \quad \forall x \in \mathbb{R}^n$$

where $\alpha : \mathbb{R}^n \rightarrow \mathbb{R}$ continuous positive definite

Lyapunov Theorem

Theorem 1. If a system admits a Lyapunov function $V(x)$, then $\bar{x} = 0$ is **globally asymptotically stable**.

Theorem 2 (Lyapunov indirect method). For linearization of system around $\bar{x} = 0$ and resulting matrix

$$A = \left. \frac{\partial g}{\partial x^T} \right|_{x=0} \text{ with eigenvalues}$$

$$|\lambda_i| := \begin{cases} \forall i : |\lambda_i| < 1 & x=0 \text{ is asymptotically stable} \\ \exists i : |\lambda_i| > 1 & \text{origin is unstable} \\ \exists i : |\lambda_i| = 1 & \text{no info about stability} \end{cases}$$

Discrete-Time Lyapunov equation

$$A^T P A - P = -Q, \quad Q > 0$$

Theorem 3 (Existence of solution of DT Lyapunov equation). The discrete-time Lyapunov equation (3) has a unique solution $P > 0$ if and only if A has all eigenvalues inside the unit circle, i.e. if and only if the system $x(k+1) = Ax(k)$ is stable.

1.6 Optimal Control

Unconstrained Finite Horizon Control Problem

$$J^*(x(0)) := \min_U x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i$$

subject to $x_{i+1} = Ax_i + Bu_i \quad i = 0, \dots, N-1$
 $x_0 = x(0)$

$P \succeq 0$, with $P = P^T$ terminal weight

$Q \succeq 0$, with $Q = Q^T$ state weight

$R \succ 0$, with $R = R^T$ input weight

1.7 Batch Approach

expresses cost function in terms of $x(0)$ and input sequence U

$$\begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} \mathbb{I} \\ A \\ \vdots \\ A^N \end{bmatrix} x(0) + \begin{bmatrix} 0 & \cdots & 0 \\ B & 0 & 0 \\ AB & B & 0 \\ \vdots & \ddots & 0 \\ A^{N-1}B & \cdots & B \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix}$$

$$\bar{Q} := \text{blockdiag}(Q, \dots, Q, P) \quad \bar{R} := \text{blockdiag}(R, \dots, R)$$

Optimal Input (from $\nabla_U J(x(0), U) = 2HU + 2F^T x(0) = 0$)

$$U^*(x(0)) = - \underbrace{((S^u)^T \bar{Q} S^u + \bar{R})}_{H(\text{Hessian})^{-1}} \underbrace{(S^u)^T \bar{Q} S^x}_{F^T} x(0)$$

Optimal Cost

$$J^*(x(0)) = x(0)^T (S_x^T \bar{Q} S_x - S_x^T \bar{Q} S_u (S_u^T \bar{Q} S_u + \bar{R})^{-1} S_u^T \bar{Q} S_x) x(0)$$

1.8 Recursive Approach

uses dynamic programming to solve problem backwards from N

$$J_j^*(x(j)) := \min_{U_{j \rightarrow N}} x_N^T P x_N + \sum_{i=j}^{N-1} x_i^T Q x_i + u_i^T R u_i$$

Ricatti Equations

RDE - Riccati Difference Equation

$$P_i = A^T P_{i+1} A + Q - A^T P_{i+1} B (B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A$$

RDE - Riccati Difference Equation solved recursively

$$P_i = A^T P_{i+1} A + Q - A^T P_{i+1} B (B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A$$

ARE - Algebraic Riccati Equation solved analytically

$$P_\infty = A^T P_\infty A + Q - A^T P_\infty B (B^T P_\infty B + R)^{-1} B^T P_\infty A$$

From Principle Of Optimality

$$J_j^*(x_j) = \min_{u_j} I(x_j, u_j) + J_{j+1}^*(x_{j+1})$$

Optimal Cost-To-Go

$$J_i^*(x_i) = x_i^T P_i x_i$$

Optimal Control Policy

$$u_i^* = F_i x_i = -(B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A \cdot x(i)$$

1.9 Comparison of Batch and Recursive Approaches

Batch optimization returns sequence $U^*(x(0))$ of **numeric values** depending only on $x(0)$, dynamic programming yields **feedback policies** $u_i^* = F_i x_i$ depending on each x_i .

Choice of P

1. Match infinite solution, use ARE

2. Assume no control needed after N , use Lyapunov Equation (makes only sense when asymptotically stable, otherwise P not positive definite)

3. set constraint $x_{i+N} = 0$

1.10 Infinite Horizon LQR

LQR

$$J_{\infty}^*(x(k)) = \min \sum_{i=0}^{\infty} x_i^{\top} Q x_i + u_i^{\top} R u_i$$

subj. to $x_{i+1} = A x_i + B u_i, \quad x_0 = x(k)$

Same u as for finite problem but with ARE Constant
Feedback Matrix F_{∞} asymptotically stable for..
 $Q, R, \text{stabi, detect}$

1.11 Optimization

A mathematical optimization problem is generally formulated as:

Mathematical Optimization Problem

Decision variable $x \in \mathbb{R}^n$

Objective function $f : \text{dom}(f) \rightarrow \mathbb{R}$

Inequality constraints $g_i \ (i \in \# \text{constraints})$

Equality constraints $h_i \ (i \in \# \text{constraints})$

Feasible set $\mathcal{X} := \{x | g(x) \leq 0, h(x) = 0\}$

minimize $f(x)$

subject to:

$$g_i(x) \leq 0$$

$$h_i(x) = 0$$

Feasible point $x \in \text{dom}(f)$ with $g_i(x) \leq 0, h_i(x) = 0$

Strictly feasible point x with strict inequality $g_i(x) < 0$

Optimal value f^* (or p^*) $= \inf\{f(x) | g_i(x) \leq 0, h_j = 0\}$ $f^* = +\infty$: OP infeasible, $f^* = -\infty$: OP unbound below

Optimizer set: $\text{argmin}_{x \in \mathcal{X}} f(x) := \{x \in \mathcal{X} | f(x) = f^*\}$

x^* is a **Global Minimum** if $f(x^*) \leq f(x)$

x^* is a **Local Minimum** if $\exists \epsilon > 0$ s.t. $f(x^*) \leq f(x)$

$\forall x \in \mathcal{X} \cap B_{\epsilon}(x^*)$, open ball with center x^* and radius ϵ

1.12 Convex Sets

Definition 4 (Convex Set). Set \mathcal{C} is convex if and only if

$$\theta x + (1 - \theta)y \in \mathcal{C}, \quad \forall x, y \in \mathcal{C}, \quad \forall \theta \in [0, 1]$$

Definition 5 (Hyperplanes). $\{x \in \mathbb{R}^n | a^{\top} x = b\}$

Definition 6 (Halfspaces). $\{x \in \mathbb{R}^n | a^{\top} x \leq b\}$

can be **open** (strict inequality) or **closed** (non-strict inequality)

Definition 7 (Polyhedra). intersection of **finite** number of closed halfspaces: polyhedra $\{x \in \mathbb{R}^n | A^{\top} x \leq b^{q \times 1}\}$

Definition 8 (Polytope). is a **bounded** polyhedron.

Definition 9 (Convex hull). for $\{v_1, \dots, v_k\} \in \mathbb{R}^d$ is:

$\text{co}(\{v_1, \dots, v_k\}) := \{x | x = \sum_i \lambda_i v_i, \lambda \geq 0, \sum_i \lambda_i = 1\}$

Definition 10 (Ellipsoid). set: $\{x | (x - x_c)^{\top} A^{-1} (x - x_c) \leq 1\}$ where x_c is center of ellipsoid, $A \succ 0$ (i.e. positive definite) (Semi-axis lengths are square roots of eigenvalues of A)

Definition 11 (Norm Ball). $B_r(x) := \{\xi \in \mathbb{R}^n : |\xi - x|_p < r\}$ where p defines the l_p norm, $p = \{1, 2, \dots, \infty\}$

Intersection $\mathcal{C}_1, \mathcal{C}_2 \text{ cv} \Rightarrow \mathcal{C}_1 \cap \mathcal{C}_2 \text{ convex (cv)}$

Image under affine map $\mathcal{C} \subseteq \mathbb{R}^n \text{ cv} \Rightarrow \{Ax + b | x \in \mathcal{C}\} \text{ cv}$

Inverse loaM $\mathcal{C} \subseteq \mathbb{R}^m \text{ cv} \Rightarrow \{x \in \mathbb{R}^n | Ax + b \in \mathcal{C}\} \text{ cv}$

1.13 Convex Functions

Definition 12 (Convex Function). $f : \mathcal{C}_{cv} \rightarrow \mathbb{R}$ is convex iff

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y), \quad \forall x, y \in \mathcal{C}, \quad \forall \theta \in [0, 1]$$

f is strictly convex if this inequality is strict.

Definition 13 (Epigraph). $f : \mathbb{R}^n \rightarrow \mathbb{R} \text{ cv} \Leftrightarrow \text{epi}(f)$ is cv set

$$\text{epi}(f) := \{(x, t) \in \mathbb{R}^{n+1} | f(x) \leq t\}$$

Check Convexity f is convex if it is composition of simple convex function with convexity preserving operations or if

$f : \mathbb{R}^n \rightarrow \mathbb{R}$ twice differentiable, $\partial^2 f / \partial x^2 \succeq 0 \quad \forall x \in \mathbb{R}^n$

$g : \mathbb{R} \rightarrow \mathbb{R}$ with $g(t) = f(x + tv)$ convex in $t \quad \forall x, v \in \mathbb{R}^n \rightarrow f$ convex (restriction to a line)

- the point wise maximum of convex functions is convex

- the sum of convex functions is convex

- $f(Ax + b)$ is convex if f is convex

Theorem 4. For a convex optimization problem, **any** locally optimal solution is globally optimal (local optima are global optima).

Linear Programming minimize $c^{\top} x$ s.t. $Ax - b \geq 0, x \geq 0$

Step 1: $\mathcal{L}(x, \lambda_1, \lambda_2) = c^{\top} x - \lambda_1^{\top} (Ax - b) - \lambda_2^{\top} x, \lambda_i \geq 0$

Step 2: $\inf_{x \in \mathbb{R}^n} \mathcal{L} = \lambda_1^{\top} b$, if $c - A^{\top} \lambda_1 - \lambda_2 = 0$, else $-\infty$

Step 3: Dual, maximize $b^{\top} \lambda$ s.t. $c - A^{\top} \lambda \geq 0, \lambda \geq 0$ (again LP)

Quadratic Programming min ...

1.14 Optimality Conditions

Lagrange Duality

Consider $f^* = \inf_{x \in \mathbb{R}^n} f(x)$ s.t. $g(x) \leq 0, h(x) = 0$

Lagrangian $\mathcal{L}(x, \lambda, \nu) = f(x) + \lambda^{\top} g(x) + \nu^{\top} h(x)$

Dual Function $d(\lambda, \nu) = \inf_{x \in \mathbb{R}^n} \mathcal{L}(x, \lambda, \nu)$

Proposition 1 (Weak Duality). $d(\lambda, \nu) \leq f^*, \forall \lambda \geq 0, \nu \in \mathbb{R}^h$

Definition 14 (Constraint qualification). **Slater's Condition** holds if \exists at least one strictly feasible point $\hat{x} (h(\hat{x}) = 0, g(\hat{x}) < 0)$

Proposition 2 (Strong Duality). If Slater's condition holds and OP is convex $\Rightarrow \exists \lambda \geq 0, \nu \in \mathbb{R}^{n_h}$ s.t. $d(\lambda, \nu) = f^*$

KKT Conditions (Karush-Kuhn-Tucker)

Theorem 5 (KKT Conditions). If Slater's condition holds and (1.14) is convex $\rightarrow x^* \in \mathbb{R}^n$ is a minimizer of the primal (1.14) and $(\lambda^* \geq 0, \nu^*) \in \mathbb{R}^{n_g} \times \mathbb{R}^{n_h}$ is a maximizer of the dual \Leftrightarrow is equivalent to the following statements:

- KKT-1** (Stationary Lagrangian) $\nabla_x \mathcal{L}(x^*, \lambda^*, \nu^*) = 0$
- KKT-2** (primal feasibility) $g(x^*) \leq 0, h(x^*) = 0$
- KKT-3** (dual feasibility) $\lambda^*, \nu^* \in \mathbb{R}^{n_h} \geq 0$
- KKT-4** (complementary slackness) $\lambda^{*T} g(x^*) = 0$
 $\nu^{*T} h(x^*) = 0$

In addition we have: $\sup_{\lambda \geq 0, \nu \in \mathbb{R}^{n_h}} q(\lambda, \nu) = \inf_{x \in \mathcal{C}} f(x)$

Remark Without Slater, KKT1-4 still implies x^* minimizes (1.14) and λ, ν maximizes dual, but the converse is no longer true. There can be primal-minimizer/dual-maximizer not satisfy KKT.

2 Nominal-MPC

2.1 CFTOC

CFTOC Constrained Finite Time Optimal Control problem

$$J(x(k)) = x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i$$

$$J^*(x(k)) = \min_U I_f(x_N) + \sum_{i=0}^{N-1} I(x_i, u_i)$$

$$\text{s.t. } x_{i+1} = A x_i + B u_i, i = 0, \dots, N-1$$

$$x_i \in \mathcal{X}, u_i \in \mathcal{U}, \mathcal{X}_N \in \mathcal{X}_f, x_0 = x(k)$$

N is the time horizon and X, U, Xf are polyhedral regions

2.2 Transform Quadratic Cost CFTOC into QP

Goal $\min_{z \in \mathbb{R}^n} \frac{1}{2} z^T H z + q^T z + r \quad \text{s.t. } G z \leq h, A z = b$

Substitute without substitution

Idea Keep state equations as equality constraints

Define variable $z = [x_1^T \dots x_N^T u_0^T \dots u_{N-1}^T]^T$

Equalities from system dynamics $x_{i+1} = A x_i + B u_i$

$$G_{eq} = \left[\begin{array}{c|c} \mathbb{I} & -B \\ -A & \mathbb{I} \end{array} \right] E_{eq} = \begin{bmatrix} A \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

Inequalities $G_{in} z \leq w_{in} + E_{in} x(k)$ from $\mathcal{X} = \{x \mid A_x x \leq b_x\}, \mathcal{U} = \{u \mid A_u u \leq b_u\}, \mathcal{X}_f = \{x \mid A_f x \leq b_f\}$

$$G_{in} = \left[\begin{array}{c|c} \begin{matrix} 0 & 0 \\ A_x & 0 \\ & \ddots \\ & A_x & A_f \\ 0 & & & 0 \end{matrix} & \begin{matrix} 0 \\ & \ddots & 0 \\ A_u & & \\ & \ddots & A_u \\ & & A_u \end{matrix} \end{array} \right] w_{in} = \begin{bmatrix} b_x \\ b_x \\ \vdots \\ b_x \\ b_f \\ b_u \\ \vdots \\ b_u \end{bmatrix}$$

$$E_{in} = \begin{bmatrix} -A_x \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

Cost Matrix $\bar{H} = \text{diag}(Q, \dots, Q, P, R, \dots, R)$

Finally the resulting quadratic optimization problem

$$J^*(x(k)) = \min_z [z^T x(k)^T] \begin{bmatrix} \bar{H} & 0 \\ 0 & Q \end{bmatrix} [z^T x(k)^T]^T$$

$$\text{s.t. } G_{in} z \leq w_{in} + E_{in} x(k) \quad G_{eq} z = E_{eq} x(k)$$

Substitute with substitution

Idea Substitute the state equations.

Step 1 Rewrite cost as

$$J(x(k)) = U X X X X$$

$$= [U^T \quad x(k)^T] \begin{bmatrix} H & F^T \\ F & Y \end{bmatrix} [U^T \quad x(k)^T]^T$$

Step 2 Rewrite constraints compactly as $GU \leq w + Ex(k)$

Step 3 Rewrite constrained problem as

$$J^*(x(k)) = \min_U [U^T \quad x(k)^T] \begin{bmatrix} H & F^T \\ F & Y \end{bmatrix} [U^T \quad x(k)^T]^T$$

$$\text{subj. to } GU \leq w + Ex(k)$$

2.3 Invariance

Definition 15 (Positively Invariant Set \mathcal{O}). For an autonomous or closed-loop system, the set \mathcal{O} is positively invariant if:

$$x(k) \in \mathcal{O} \Rightarrow x(k+1) \in \mathcal{O}, \quad \forall k \in \{0, 1, \dots\}$$

Definition 16 (Maximal Positively Invariant Set \mathcal{O}_∞). A set that contains all \mathcal{O} is the maximal positively invariant set $\mathcal{O}_\infty \subset \mathcal{X}$

Definition 17 (Pre-Sets). The set of states that in the dynamic system $x(k+1) = g(x(k))$ in one time step evolves into the target set \mathcal{S} is the **pre-set** of $\mathcal{S} \Rightarrow \text{pre}(\mathcal{S}) := \{x \mid g(x) \in \mathcal{S}\}$

Theorem 6 (Geometric condition for invariance). Set \mathcal{O} is positively invariant set iff $\mathcal{O} \subseteq \text{pre}(\mathcal{O}) \Leftrightarrow \text{pre}(\mathcal{O}) \cap \mathcal{O} = \mathcal{O}$

Proof. Necessary if $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$, then $\exists \bar{x} \in \mathcal{O}$ s.t. $\bar{x} \notin \text{pre}(\mathcal{O}) \rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \text{pre}(\mathcal{O})$, thus \mathcal{O} not positively invariant

Sufficient if \mathcal{O} not pos invar set, then $\exists \bar{x} \in \mathcal{O}$ s.t. $g(\bar{x}) \notin \mathcal{O} \rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \text{pre}(\mathcal{O})$ thus $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$ \square

Computing Invariant Sets

Pre-Set Computation

Conceptual Algorithm

System with constraints

$$x(k+1) = Ax(k) + Bu(k)$$

$$u(k) \in \mathcal{U} := \{u | Gu \leq g\}$$

and set $\mathcal{S} := \{x | Fx \leq f\}$

$$\text{pre}(\mathcal{S}) := \{x | Ax \in \mathcal{S}\}$$

$$= \{x | FAx \leq f\}$$

```

first line
 $\Omega_0 \leftarrow \mathcal{X}$ 
loop
   $\Omega_{i+1} \leftarrow \text{pre}(\Omega_i) \cap \Omega_i$ 
  if  $\Omega_{i+1} = \Omega_i$  then
    return  $\mathcal{O}_\infty = \Omega_i$ 
  end if
end loop

```

(Same but much harder for control invariant sets)

Conceptual Algorithm

```

first line
 $\Omega_0 \leftarrow \mathcal{X}$ 
loop
   $\Omega_{i+1} \leftarrow \text{pre}(\Omega_i) \cap \Omega_i$ 
  if  $\Omega_{i+1} = \Omega_i$  then
    return  $\mathcal{O}_\infty = \Omega_i$ 
  end if
end loop

```

(Same but much harder for control invariant sets)

Conceptual Algorithm

```

@decorator()
# Example Python code
def hello_world():
  # This is a comment
  print("Hello, World!")

```

(Same but much harder for control invariant sets)

2.4 Control Invariance

Definition 18 (Control Invariant Set). $\mathcal{C} \subseteq \mathcal{X}$ control invariant if

$$x(k) \in \mathcal{C} \Rightarrow \exists u(k) \in \mathcal{U} \text{ s.t. } g(x(k), u(k)) \in \mathcal{C} \forall k$$

Definition 19 (Maximal Control Invariant Set \mathcal{C}_∞). A set that contains all \mathcal{C} is the maximal positively invariant set $\mathcal{C}_\infty \subset \mathcal{X}$

Intuition For all states in \mathcal{C}_∞ exists control law s.t constraints are never violated \rightsquigarrow **The best any controller could ever do**

Pre-set $\text{pre}(\mathcal{S}) := \{x | \exists u \in \mathcal{U} \text{ s.t. } g(x, u) \in \mathcal{S}\}$

Set \mathcal{C} is control invariant iff: $\mathcal{C} \subseteq \text{pre}(\mathcal{C}) \Leftrightarrow \text{pre}(\mathcal{C}) \cap \mathcal{C} = \mathcal{C}$

Control Law from Control Invariant Set

Let \mathcal{C} control invariant set for
 $x(k+1) = g(x(k), u(k))$
 Control law $\kappa(x(k))$ will **guarantee** that system satisfies constraints **for all time** if:

$$g(x, \kappa(x)) \in \mathcal{C} \forall x \in \mathcal{C}$$

We can use this fact to **synthesize** control law κ with f as any function (including $f(x, u) = 0$)

$$\kappa(x) := \text{argmin}\{f(x, u) \mid g(x, u) \in \mathcal{C}\}$$

Does not ensure that system will converge
 Difficult because calculating control invariant sets is hard
MPC implicitly describes \mathcal{C} s.t easy to represent/compute

Theorem 7. Minkowski-Weyl

The following statements are equivalent for $\mathcal{P} \subseteq \mathbb{R}^d$

- \mathcal{P} is a polytope and there exists A, b s.t $\mathcal{P} = \{x \mid Ax \leq b\}$
- \mathcal{P} finitely generated, \exists finite set $\{v_i\}$ s.t $\mathcal{P} = \text{co}(\{v_1, \dots, v_s\})$

MOST COMMON Polytopic

1

Lemma 1. Invariant Sets from Lyapunov Functions

If $V : \mathbb{R}^n \rightarrow \mathbb{R}$ is a Lyapunov function for $x(k+1) = g(x(k))$, then $Y := \{x \mid V(x) \leq \alpha\}$ is an invariant set for all $\alpha \geq 0$

Proof. Lyapunov property $V(g(x)) - V(x) < 0$ implies that once $V(x(k)) \leq \alpha, V(x(j)) < \alpha, \forall j \geq k \rightarrow$ Invariance \square

Example System for $x(k+1) = Ax(k)$ with $P \succ 0$ that satisfies $A^T P A - P \prec 0 \rightsquigarrow$ then $V(x(k)) = x(k)^T P x(k)$ is Lyap. function

Goal – find largest α s.t set $Y_\alpha \in \mathcal{X}$

$$Y_\alpha := \{x \mid x^T P x \leq \alpha\} \subset \mathcal{X} := \{x \mid Fx \leq f\}$$

Equivalent to $\max_\alpha \alpha$ subj. to $h_{Y_\alpha}(F_i) \leq f_i \forall i \in \{1 \dots n\}$

2.5 Feasibility and Stability

What can go wrong with “standard” MPC?

- No feasibility guarantee, i.e., the MPC problem may not have a solution

- No stability guarantee, i.e., trajectories may not converge to the origin

MPC Mathematical Formulation

V1

$$J^*(x(k)) = \min \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

V3

$$J^*(x(k)) = \min \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

$$\text{subj. to } x_{i+1} = Ax_i + Bu_i \\ x_0 = x(k) \quad x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}$$

Stability of MPC - Main Result

Theorem 8. The closed-loop system under the MPC control law $u_0^*(x)$ is asymptotically stable and the set \mathcal{X}_f is positive invariant for the system $x(k+1) = Ax(k) + Bu_0^*(x(k))$ under the following assumptions:

1. Stage cost is positive definite, i.e. it is strictly positive and only zero at the origin
2. Terminal set is **invariant** under the local control law $\kappa_f(x_i)$:

$$x_{i+1} = Ax_i + B\kappa_f(x_i) \in \mathcal{X}_f \forall x_i \in \mathcal{X}_f$$

All state and input **constraints are satisfied** in \mathcal{X}_f :

$$\mathcal{X}_f \in X, \kappa_f(x_i) \in U \forall x_i \in \mathcal{X}_f$$

3. Terminal cost is a continuous **Lyapunov function** in the terminal set \mathcal{X}_f and satisfies:

$$I_f(x_{i+1}) - I_f(x_i) \leq -I(x_i, \kappa_f(x_i)) \quad \forall x_i \in \mathcal{X}_f$$

Finite-horizon MPC may not satisfy constraints for all time!

Finite-horizon MPC may not be stable!

- An infinite-horizon provides stability and invariance.
- Infinite-horizon *faked* by forcing final state into an invariant set for which there exists invariance-inducing controller, whose infinite-horizon cost can be expressed in closed-form.
- Extends to non-linear systems, but compute sets is difficult!

3 Practical-MPC

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4 Robust-MPC

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5 Implementation

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