

# r-2012-11-01

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## 1 Processed real data

Table 1 summarizes the so far processed timetables. The regional bus timetables provided by *cp.sk* are the examples of more local timetables with underlying network being the road network. The time range was so far set to 1 day, assuming that each bus operates the same way every day. This is not a correct assumption (some buses may operate e.g. only during weekends), but this way we save space and the resulting timetable should not differ much in properties from the real one. We may later however try to extend the time range to one week and see the changes in properties.

Timetable from *zsr.sk* covers a full year, as connections had a reference to a calendar, stating if the train operates on a given date. That is also the reason for the large number of elementary connections.

Timetable of *United airlines* requires more processing, but should be later added as it provides yet another type of timetable.

Name	Type of network	Provided by	# of stations	# of elementary connections	Time range	Note
tt_cpza	Regional bus	cp.sk	1128	68096	1 day	Area of Žilina
tt_cpru	Regional bus	cp.sk	877	43069	1 day	Area of Ružomberok
tt_zsr	Country-wide rails	zsr.sk	233	932052	1 year	

Table 1: Timetable files

Timetable files have a simple form:

1	7	//number of elementary connections
2	A B 0 10:00 0 10:45	//FROM TO DAY(depart) TIME(depart) DAY(arrive) TIME(arrive)
3	A B 0 11:00 0 11:45	
4	A B 0 12:00 0 12:45	
5	A C 0 9:30 0 10:00	
6	A C 0 10:15 0 10:45	

7	C D 0 11:00 0 11:30
8	C D 0 13:00 0 13:30

Listing 1: Timetable files form

Table 2 summarizes the so far processed maps (or underlying graphs) of networks, on top of which we can later build random timetables with given properties. The maps were created by extracting stations from the timetable and making an edge between the pair where a connection is operating, thus getting the underlying structure of the timetable. Note that we will not really get the real map of the (e.g.) rails - with our approach, there will be a direct link between a pair of large (but distant) cities even though many smaller stations exist on the way. This is due to express lines that do not stop at these smaller stations. If the real underlying map should be desirable, we can later remove these “express links”, although we cannot be sure that the express line did not really follow a new, faster road/railway without any smaller stops on the way (that should not happen in case of railways).

Map of the *road network of Slovakia* should be added, as well as the underlying map of *United airlines*. Also, currently the edges of the underlying graphs have null lengths - the travel time of the fastest connection on the edge can be used instead.

Name	Type of network	Provided by	# of stations	# of connections	Notion of lines	Note
ug_cpza	Regional bus	cp.sk	1128	2034	✓	Area of Žilina
ug_cpzu	Regional bus	cp.sk	877	1784	✓	Area of Ružomberok
ug_zsr	Country-wide rails	zsr.sk	233	588	✓	
ug_london	Underground rails	queensu.ca	321	732	✓	

Table 2: Underlying graph files

Underlying graph files have the following form:

1	4	//number of stations
2	5	//number of edges
3	A [45 32]	//name of the station, optional coordinates
4	B	
5	C [56 34]	
6	D	
7	A B 57 Northern	//FROM TO edge <b>length</b> , list of lines operating
	on the edge	
8	A C <b>null</b> Picadilly Victoria	//edge <b>length</b> may be <b>null</b> (will be e.g. random,
	or calculated from coordinates)	
9	C B 45 Circle Jubilee Picadilly	
10	C D 32 <b>null</b>	//list of lines may be also <b>null</b>
11	D A 90 Metropolitan Victoria	

Listing 2: Underlying graph files form

## 2 Overtaking and express lines

In this section, we will define and discuss two notions: *overtaking* and *express lines*. For clarity, we repeat the definition of the timetable:

### Definition 1. Timetable on a graph $G$

A timetable on a given graph  $G$  is a set  $T_G = \{(x, y, p, q) | (x, y) \in E_G, p, q \in N, p < q\}$ .

- An element of  $T$  is called an **elementary connection** and  $x$  [ $y$ ] and  $p$  [ $q$ ] are the **departure** [**arrival**] **node** and **time**.
- Graph  $G$  is the **underlying graph** of timetable  $T$ .

### Definition 2. Connection from $a$ to $b$ in a given timetable $T_G$

A connection from  $a$  to  $b$  in a given timetable  $T_G$  is a sequence of elementary connections  $(e_1, e_2, \dots, e_k)$ ,  $k \geq 1$ ,  $e_i = (e_x^i, e_y^i, e_p^i, e_q^i)$ , such that  $e_x^1 = a$ ,  $e_q^k = b$  and  $\forall i \in \{2, \dots, k\} : (e_x^i = e_y^{i-1}, e_p^i \geq e_q^{i-1})$

- Connection **starts** at the departure time  $e_p^1$ .

- **Size (length)** of the connection is  $e_q^k - e_p^1$ .

Overtaking, simply stated, means that there are connections in the timetable that overtake one another - that is by choosing a later connection, we will actually arrive sooner to the destination.

**Definition 3. Overtaking**

A timetable  $T_G$  has overtaking property  $\iff$  there exist two connections  $e$  and  $f$  from  $a$  to  $b$  in  $T_G$  such that  $e = (e_1, e_2, \dots, e_k)$ ,  $f = (f_1, f_2, \dots, f_k)$ ,  $\forall i \ e_x^i = f_x^i$ ,  $e_y^i = f_y^i$  and  $e_p^1 < f_p^1$  but  $e_q^k > f_q^k$ .

Informally, express lines are connections that bypass several stops on the way. In our definition of timetable graph, we do not allow such connections, but we can extend the definition.

**Definition 4. Timetable with express lines on a graph  $G$**

A timetable with express lines on a given graph  $G$  is a set  $T_G = \{(x, y, p, q) | \text{there is a path from } x \text{ to } y \text{ in } G, p, q \in N, p < q\}$ .

- Tuples  $(x, y, p, q)$  where  $(x, y) \in E_G$  are **elementary connections**, the remaining are **express connections**.
- Such  $(x, y)$  that  $(x, y) \notin E_G$  but there is a path from  $x$  to  $y$  in  $G$  will be called **express lines**.

Now we can formulate the following observation:

**Observation 1.**

Let  $G$  be a graph with the vertex set  $V$ ,  $T_G$  timetable on  $G$  and  $H$  a connected graph with the vertex set  $V$ . There exists a timetable with express lines  $T_H$  such that  $T_H = T_G$ .

*Proof.* The proof is trivial: By the definition of timetable with express lines, there may be a connection between any pair of vertices in the underlying graph  $H$ , as it is connected. Thus for every elementary connection  $(x, y, p, q) \in T_G$  we will simply have the same (*either elementary or express*) connection  $(x, y, p, q)$  in  $T_H$  □

Informally, this observation just states, that with express lines we can forget about underlying structure of the timetable and consider it as a complete graph. That would however imply that no property of the underlying graph  $G$  can be generally propagated to the graph representing the timetable on  $G$  (not even to certain extent). Unfortunately, express lines are not uncommon in real-world scenarios. On the other hand, usage of express lines is limited in practice and generally express lines act as a shortcut layer on top of the real underlying network. One can imagine this as a hierarchy starting from the road network with nodes being the intersections. Next layer would be the local bus lines, then express inter-city lines etc...

### 3 Adjustment of Gavaille's algorithm for graphs with $r(n)$ separator [?]

### 4 Open points

- What does overtaking really affect?
- Hierarchy of express lines  $\rightarrow$  what properties can be propagated in time-expansion?

### 5 To do

- Adjustment of Gavaille's algorithm - formally
- United airlines and road network of Slovakia - extract data
- Start work on the diagnostic program
- Replace null lengths in underlying graphs of existing timetables