

Topic B

RELATIONAL DATA

MODEL

#### RELATIONAL DATA MODEL

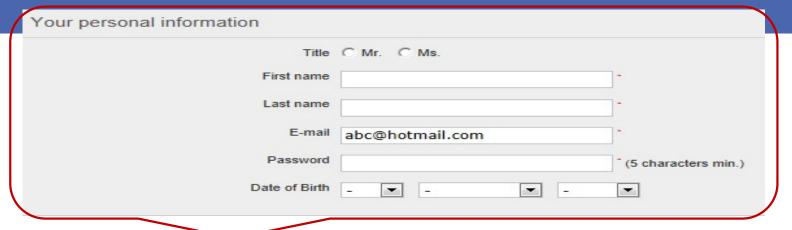
#### CONTENT

- The basic concepts of a database
- ■The relational terminologies
- ■The properties of a relation
- Entity Relationship Diagram (ERD)
- Transform ERD to Database
- Normalization

#### Basic Concepts of Database

- What is a database?
  - A collection of inter-related data
- Inter-related data
  - This means data are linked together by some mechanism
- What is a database management system
  - A collection of programs to manage(update, insert, delete, select), protect and control access to the database

#### The Relational Terminologies – A relation



#### Customer

Title	First_name	Last_name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

A relation is a 2 dimensional table. The intersection between a row and column is called a cell.

#### The Relational Terminologies – Record/Tuple

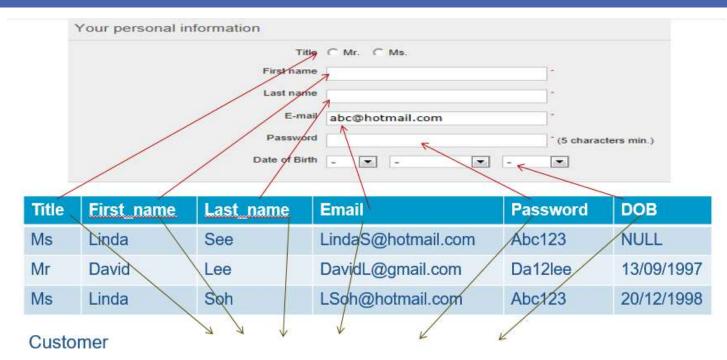
Title	First_name	Last_name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

Customer

Each row in the table is a record which represents an individual object in the relation.

Record of a table is also called tuple.

#### The Relational Terminologies – Attributes



Attributes are column names of a table to define the characteristics of the object.

#### The Relational Terminologies – Attribute Values

Title	First_name	Last_name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998
Customer  Attribute Value		e Value			

#### Concept of NULL Attribute Value -

- NULL is a special value allowed in a relational database
- NULL means the value is UNKNOWN or NOT APPLICABLE

#### The Relational Terminologies – Attribute Domain



Title	First_name	Last name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	01/02/2000
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

Customer

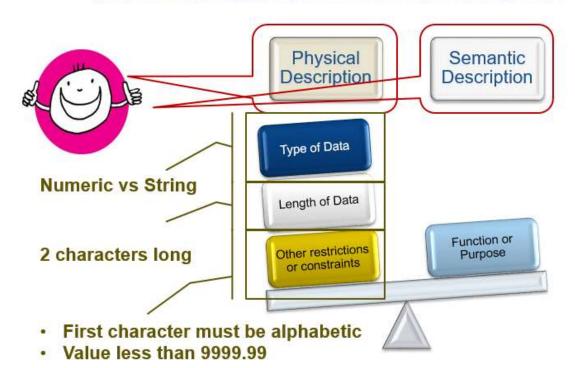
Attribute Domain -

The set of allowable values of that attribute (Example Ms, Mr) A description of the possible values for that attribute

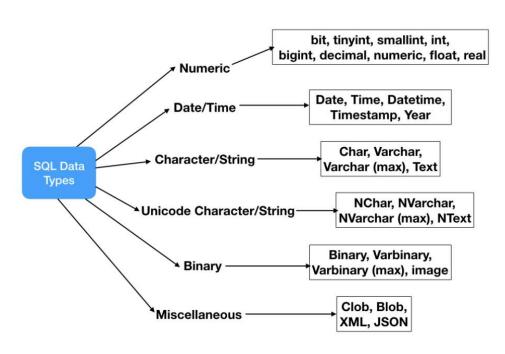
- Physical Description
- Semantic Description

### The Relational Terminologies – Attribute Domain

Domain of Title attribute is set of letters consists of {'Ms', 'Mr'}
 where Ms represents Miss while Mr represents Mister.



#### The Relational Terminologies – Attribute DataType / SQL Date Type



Source: https://www.journaldev.com/16774/sql-data-types, 29 Feb 2020

What is the difference between Char and Varchar?

Not all data types are supported by every relational database management systems.

#### The Relational Terminologies – Relational Database



Cart Data A collection of interrelated proper form relations

#### Customer

Cusic	Customer						
Title	First_ name	Last_na me	Email	Password	DOB		
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL		
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997		
Ms	Linda	Soh	Lsoh@hotmail.com	Abc123	20/12/1998		

#### Cart

Email	Product_code	Cart_ID	Qty
DavidL@gmail.com	HG4872	C10001	1
DavidL@gmail.com	HG5879	C10002	1

#### Properties Of A Relation

#### A table must satisfy the **SIX** properties of a relation.

#I: Name of Relation is unique in a database
You cannot create another relation/table in the same database by the same
name

#2: Every cell must be single-valued

You cannot have multi-valued cell

#3: Attribute name in a relation must be unique

You cannot have two attributes by the same name

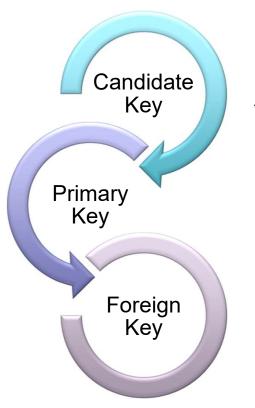
#4: Values of an attribute are from the same domain

#5: Order of tuples or attributes in a relation does not matter

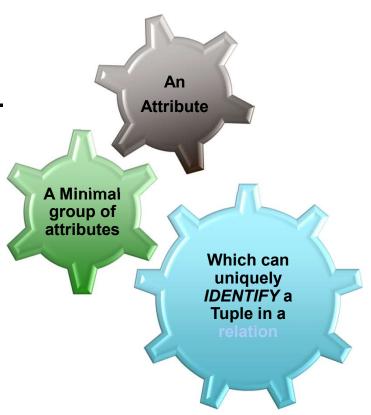
#6: Each tuple in a relation is unique

You cannot have <u>duplicate tuple/row</u>

## The Relationship Integrity Rules - Keys



A candidate key is -



#### The Relationship Integrity Rules – Candidate Keys



Email (an attribute) can uniquely identify a customer tuple (record) in the Customer relation

Customer

## The Relationship Integrity Rules – Candidate Keys

Title	First_name	Last_name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

Customer

#### A minimal group of attributes

First\_name, Last\_name, Password

Identify the group of attributes which can uniquely identify a customer based on the tuples

## The Relationship Integrity Rules – Primary Key

There are TWO candidate keys of customer relation:

- I. Single attribute email
- 2. Minimal group of attributes First name, last name, Password

Do we need so many identifiers for a relation ?

NO! We do not need so many identifiers for a relation.

 Choose one of the most <u>suitable candidate key</u> be the official identifier of a relation - Primary key

#### The Relationship Integrity Rules – Foreign Keys

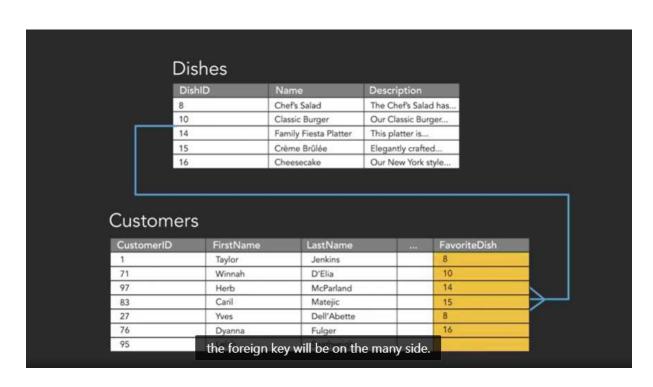
#### What is the Foreign key?

An attribute (column)

OR a group of attributes (2 or more columns) with a relation that matches

the value of the primary key OR the candidate key

of the home OR other relation



#### The Relationship Integrity Rules – Foreign Keys

Title	First_name	Last_name	Email	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

Customer

Foreign key

	Email	Product_code	Cart_ID	Qty
_	DavidL@gmail.com	HG4872	C1001	1
	DavidL@gmail.com	HG5879	C1001	1

Cart

Email of cart relation is the foreign key referencing the primary key in the customer relation

## The Relationship Integrity Rules – Entity Integrity Rule

## Applies to primary key

Title	First_name	Last_name	Email (Primary Key)	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998
Mr	Jonanthan	Sim	NULL	JonS098	NULL

**Email is the attribute of the Primary key, CANNOT BE NULL** 

Primary key = attribute or combination of attributes cannot be NULL

### The Relationship Integrity Rules – Entity Integrity Rule

Composite Primary Key = Email + Product\_code + Cart\_ID

Email (PK)		Product_code (PK)	Cart_ID (PK)	Qty	
DavidL@gmail.co	om	HG4872	C1001	1	
DavidL@gmail.co	om	HG4872	C1001	1	
NULL		HG4872	C1002	2	



No part of a primary key can be duplicated

No part of the primary key can be Null

#### The Relationship Integrity Rules – Referential Integrity Rule

## Applies to foreign key

Title	First_name	Last_name	Email (Primary Key)	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	<b>√DavidL@gmail.com</b>	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

Customer

Email (Foreign Key)	Product_code	Cart_ID	Qty
DavidL@gmail.com	HG4872	C1001	1
DavidL@gmail.com	HG5879	C1001	1
JuneL@gmail.com	HG4872	C1002	2
Cart			

Davidl@gmail.com of cart relation is the foreign key referencing the primary key in the customer relation

#### The Relationship Integrity Rules – Referential Integrity Rule

Title	First_name	Last_name	Email (Primary Key)	Password	DOB
Ms	Linda	See	LindaS@hotmail.com	Abc123	NULL
Mr	David	Lee	DavidL@gmail.com	Da12lee	13/09/1997
Ms	Linda	Soh	LSoh@hotmail.com	Abc123	20/12/1998

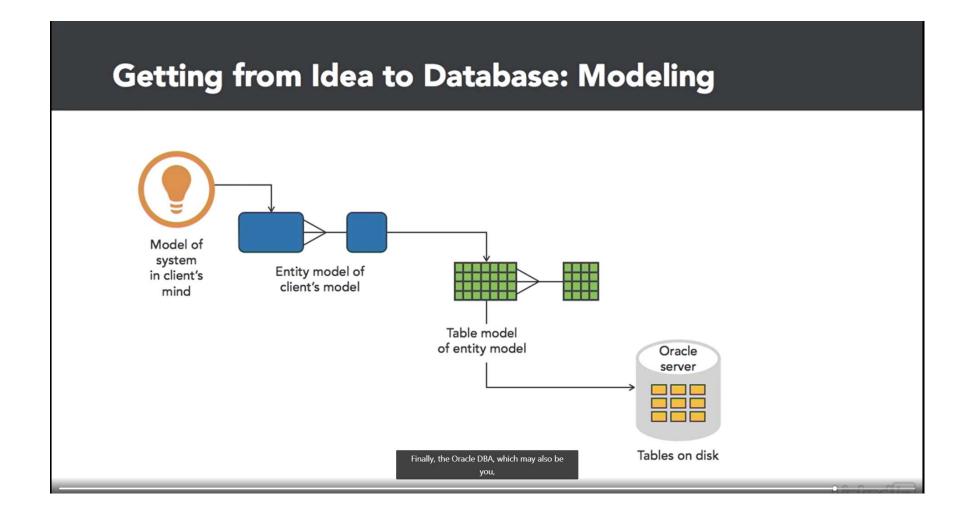
#### Customer

Email (Foreign Key)	Product_code	Cart_ID	Qty
DavidL@gmail.com	HG4872	C1001	1
DavidL@gmail.com	HG5879	C1001	1
JuneL@gmail.com	HG4872	C1002	2

Cart

- > Junel@gmail.com of cart relation is not found in the customer relation
- > Can we set the email in 3<sup>rd</sup> row of cart relation to Null? Is this logical?

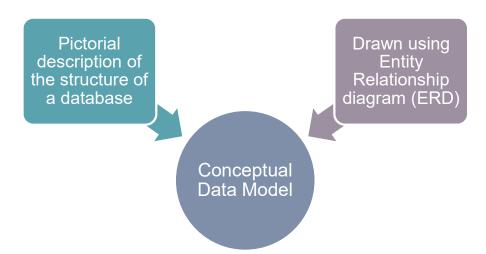
Yes we can set part or whole of the foreign key to Null



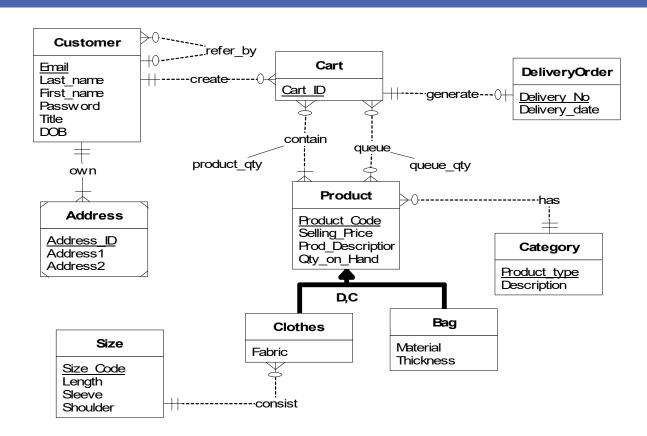
#### Entity Relationship Diagram - Conceptual Data Model

A set of concepts that describe the structure of a database and the associated retrieval and update transactions.

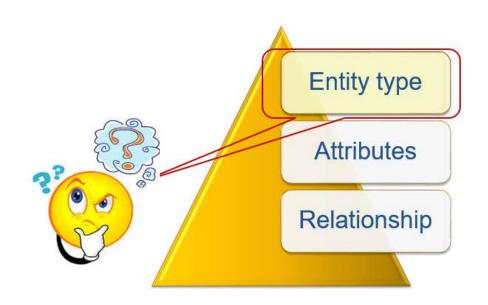
Independent of the particular type of DBMS software, such as Oracle, Microsoft SQL and hardware.



### Entity Relationship Diagram -Conceptual Data Model (ERD)



#### Entity Relationship Diagram – Entity Type

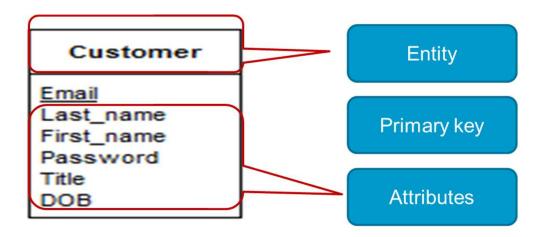


Defined as a <u>collection of entities</u> that <u>has common properties</u> and and are <u>of interest to an organization</u>

#### **Examples:**

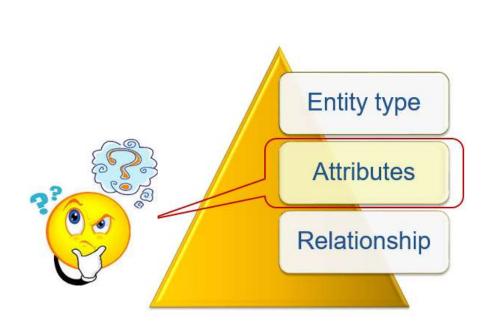
Customer, Address and Cart Product, Clothes and Bag

### Entity Relationship Diagram - Entity Type



Each entity type is identified by <u>a name (singular noun)</u> and a list of <u>attributes</u>.

## Entity Relationship Diagram - Attributes



# <u>Composite</u> <u>Attribute</u> – an attribute composes of multiple components. E.g. Name and Address

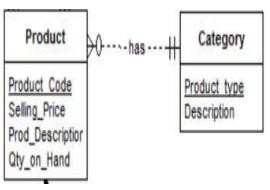
#### **Simple Attribute**

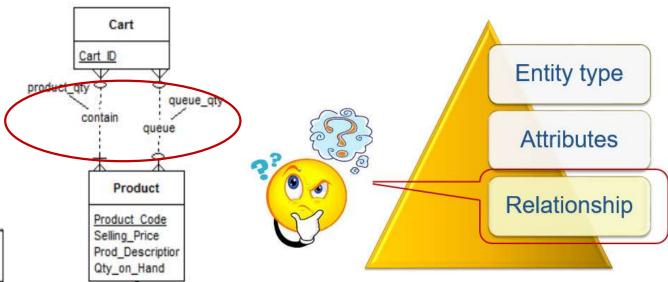
- an attribute that composes of single component.
E.g. Gender, mobile\_no, student\_no

<u>Derived Attribute</u> – an attribute that derives from a set of attributes. E.g. Age

#### Entity Relationship Diagram – Relationships (single and multiple)

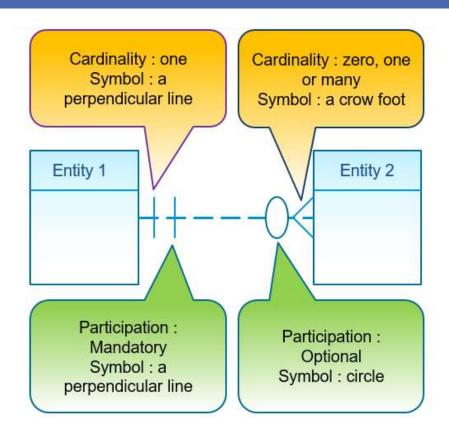
- Relationship is the association between 2 entities.
- Each relationship has a verb-based name.
- E.g. The 'has' relationship between product and product\_type enables the user to know that this product is a blouse



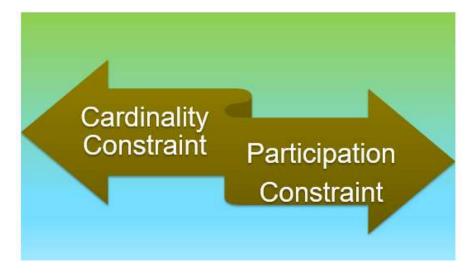


- Multiple relationships can be created when the two entities are associated through two distinct relationships.
- (a) The cart must contain one or many products
- (b) The cart may be queuing for one or more products that are out-ofstock.

#### Entity Relationship Diagram – Structural Constraints



## Reflect business rules set by the organization



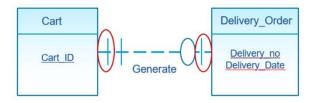
### Entity Relationship Diagram – Cardinality Constraints

Describe the number of entity instances that participate in a relationship

Cardinality Ratios	Read as
1:1	One-to-one
1 : M	One-to-many
M : N	Many-to-many

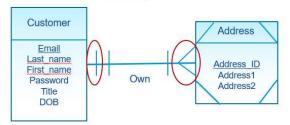
### Entity Relationship Diagram – Cardinality Constraints

#### \* 1:1 Relationship

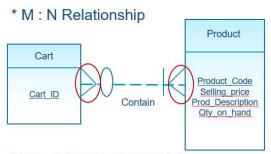


A cart generates one Delivery\_order once the customer checkout the bag. And every Delivery order belongs to one cart.

#### \* 1 : M Relationship



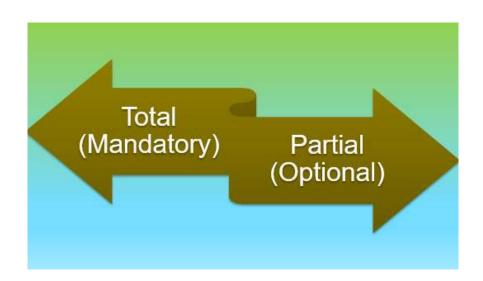
A customer has one or many addresses and each address belongs to only one customer.



A cart contains one or many products.

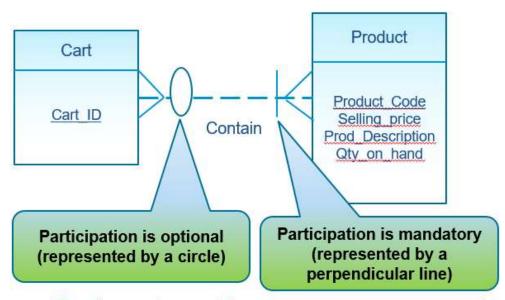
A product can be added in zero, one or many carts.

### Entity Relationship Diagram – Participation Constraints



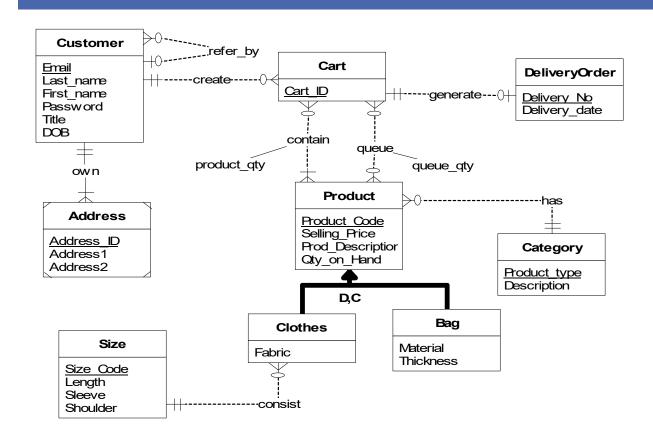
Determines if the existence of an entity type depends on another entity type

#### Entity Relationship Diagram – Participation Constraints



- Each cart must have one or more product.
- A product may not be added to a cart.

### How to transform an Entity Relation Diagram to Database?



- 1. Database tables does not have M:N Relation?
- 2. How to systematically transform an ERD to Database?

## How to transform an Entity Relation Diagram to Database?

Cart (CartID, Email)

Product (ProductCode, SellingPrice, ProductDescription, Qty, ProductType)

Category (ProductType, Description)

Cart Product(CardID, ProductCode, Product Qty)

Customer(Email, Last Name, First Name, Password, Title, DOB)

DeliveryOrder (<u>Delivery No</u>, DeliveryDate, <u>CartID</u>)

- 1. Database tables does not have M:N Relation?
- 2. How to systematically transform an ERD to Database?

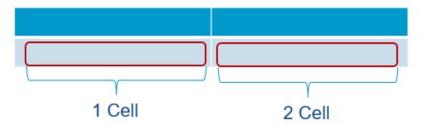


# Entity Relationship Diagram and Database

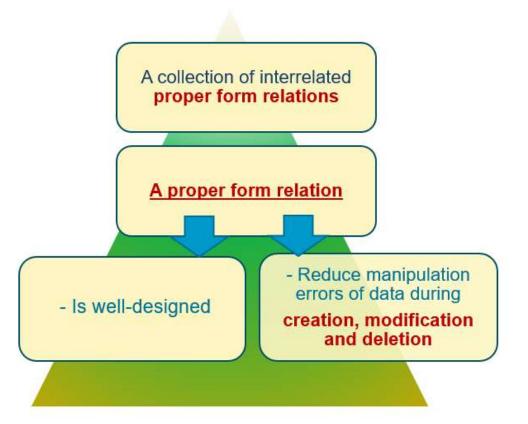
Definition		Example				
ERD	Database	ERD	Database			
Entity type	Relation / Table	Customer Entity type	Customer Relation / Table			
Entity	Tuple / Row	Customer Entity	Customer Tuple / Row			
Attribute	Attribute / Column	Email Attribute	Email Attribute / Column			
Relationship	Foreign key referencing Primary key	Relationship Own	Email_addr (Foreign key) referencing Email (Primary key)			
Primary key	Primary key	Primary key (Email)	Primary key (Email)			

## Normalization – Proper Form Relation

- □No more than one value in
  - each cell of a relation
  - □Each cell must be single value
  - □Each cell must not be multi-value



Intersection of the row and column in a table



# Normalization – Un-normalized Form (0NF/UNF)

The table is called an Un-normalized (UNF) table



- A table with multi-valued cells
- A table with repeating group

#### Customer Product Supplier

First_name	Email	Prod_Co de	Prod Desc	Unit px	Supplier _ID	Supplier Name	Qty	Ttl_Amt
David DavidL@gmail.	AG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70	
		HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
		RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte ltd Best Clothing Pte	1	15.90
		HG6159	Sale Dress Pink	15.40	S1003	Ltd	4	61.60

Single-valued cell

Multi-valued cell

# Normalization – How To Represent a 0NF/UNF?

First_name	Email	Prod_Co de	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
		HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
		RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	1	15.90
		HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

Relational Heading Format	Syntax
Customer_Product_Supplier (First_name, Email, { Prod_Code, Prod_Desc, Unit_px, Supplier_ID, Supplier_Name, Qty, Ttl_Amt } )	<relation> ( Single-valued columns, { Muti-valued columns } )</relation>

#### Normalization – Transform 0NF to 1NF

Normal Form	Description	Problem	Solution
UNF/0NF	Table with multi-valued cells	Violate the properties of relation	Form new row for each multi- valued cell to remove repeating group.

We therefore cannot make use of this improper form relation. We need to transform it to first normal form.

# First Normal Form (INF)

## Customer\_Product\_Supplier

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
David	DavidL@gmail. com	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	1	15.90
Linda	LSoh@hotmail. com	HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

Step	Description	

# Normalization – First Normal Form (INF)

First_name	Email	Prod_Code	Prod Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39,60
David	DavidL@gmail. com	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail.	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	1	15.90
Linda	LSoh@hotmail. com	HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

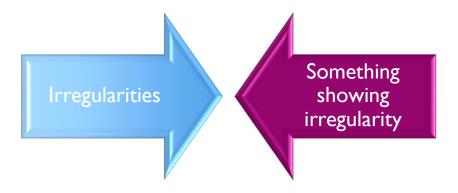
There is **NO REPEATING GROUP** in a First Normal
Form table

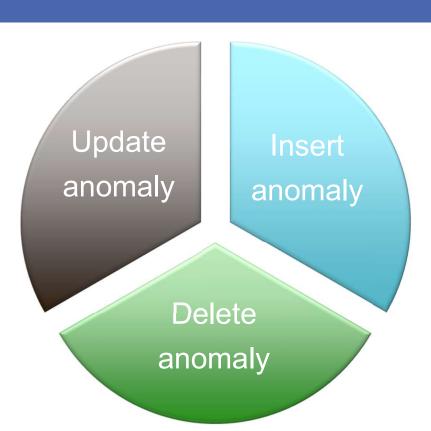


- No Repeating groups of data
- Identify the primary key

#### Normalization - Anomalies

# What are anomalies?





# Normalization-INF: Insert anomaly

Customer\_Product\_Supplier

Primary Key

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
David	DavidL@gmail.	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail.	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	1	15.90
Linda	LSoh@hotmail. com	HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

First_name	<u>Email</u>	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
	Rachel@gmail. com	?	?	?	?	?	?	?

Cannot insert a row when a particular customer (email) does not buy any product

# Normalization-INF: Insert anomaly

Customer Product Supplier

Primary Key

First_name	<u>Email</u>	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
David	DavidL@gmail. com	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail.	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	1	15.90
Linda	LSoh@hotmail. com	HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
?	?	HX1021	Red Dress	10.10	ST1001	King Dress Pte Itd	?	?

Cannot insert a row when the new product does not have any new customer

## Normalization – INF : Insert anomaly

- Which rule does Insert anomaly violates?
- ☐ Entity Integrity Rule
- ☐ Referential Integrity Rule
- Both
- None of the above

# Normalization - INF: Update anomaly

#### Customer\_Product\_Supplier

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Dress Pure White	15.90	S1001	King Kong Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Kong Dress Pte Ltd	2	39.60
David	DavidL@gmail.	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40
Linda	LSoh@hotmail.	HG7160	Dress Pure White	15.90	S1001	King Kong Dress Pte Ltd	1	15.90
Linda	LSoh@hotmail. com	HG6159	Sale Dress Pink	15.40	S1003	Best Clothing Pte Ltd	4	61.60

- a. Change the supplier name from King Dress Pte Ltd to King Kong Dress Pte Ltd
- b. Change the product description of HG7160 to Dress Pure White

Change of supplier name and product description requires updates in several rows

# Normalization – INF: Update anomaly

- What happen during Update anomaly?
- Multi-value cells created
- □ Data inconsistency due to data redundancy
- Violate Entity Integrity Rule
- □ All the above

# Normalization - INF : Delete anomaly

#### Customer\_Product\_Supplier

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60
David	DavidL@gmail. com	RQ0207	Dress White	18.60	S1002	Ladies Green Pte Ltd	4	74.40

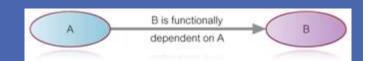
Delete Linda will result in removing product such as HG6159 that is only bought by Linda

Loss of information when rows are deleted from the relation

# Normalization – Anomaly Summary

Anomaly	Description / Examples	Violation / Problem
Insert Anomaly	Primary key (email, product_code)	Entity Integrity Rule
	Cannot insert a row when a particular customer (email) does not buy any product	
	Cannot insert a row when the new product does not have any new customer	
Update Anomaly	Changes of supplier name and product description requires updates in several rows	Data inconsistencies due to data redundancy
	Resulting in <u>data inconsistencies</u> <u>if care is not taken to update all</u> <u>relevant data.</u>	
Delete Anomaly	Delete Linda will result in removing product such as HG6159 that is only bought by Linda	Lost of information when rows are deleted from the relation

### Normalization – Functional Dependency



- A functional dependency is a constraint about two or more columns of a table
- Occurs when the value of an attribute is fully dependent upon another attribute's value
- A relation is said to be in Second Normal Form if all nonkey attributes are fully dependent on whole of its primary key

Attribute on the <u>left hand side</u> of a functional dependency (FD) is known as LHS or determinant.

Attribute B's value is functionally dependent upon attribute A's value. i.e. if there exists at most one value of B for every value of A.

# Normalization – Functional Dependency

Examples of functional dependency

Supplier\_ID -> Supplier\_Name

Supplier\_ID -> Supplier\_Phone

Supplier\_ID -> Supplier\_Address

Supplier\_ID -> Supplier\_Name, Supplier\_Phone, Supplier\_Address

#### Normalization – Transform INF to 2NF

Normal Form	Description	Problem	Solution
1NF	Table with anomalies	Violate the properties of relation	Transform further to second normal form.

As there are anomalies in INF, we therefore need to transform it second normal form.

# Normalization – Second Normal Form (2NF)

#### Customer\_Product\_Supplier

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60

Partial keys	Whole key
Email	Email, Prod_Code
Prod_Code	

Step I – Identify keys

## Normalization – Second Normal Form

Customer\_Product\_Supplier

First_name	Email	Prod_Code	Prod_Desc	Unit_px	Supplier _ID	Supplier_Name	Qty	Ttl_Amt
David	DavidL@gmail. com	HG7160	Sale Dress White	15.90	S1001	King Dress Pte Ltd	3	47.70
David	DavidL@gmail. com	HG9298	Sale Top + Skirt Red	19.80	S1001	King Dress Pte Ltd	2	39.60

Partial Functional Dependencies (One or more non-key attributes functionally dependent on part (but not all) of the primary key)	Functional Dependencies with Whole key (Non-key attributes functionally dependent upon whole of primary key)
Email -> First_Name	Email, Prod_Code -> Qty, Ttl_Amt

Always identify the FD with all the partial keys first, follow by whole key

Step 2 – List functional dependencies with keys

#### Normalization – Second Normal Form

Functional Dependencies	Relations (Relational Header Format)
Email -> First_Name	Customer ( <u>Email</u> , First_Name)
Prod_Code -> Prod_Desc, Unit_Px, Supplier_ID, Supplier_Name	Product ( <u>Prod_Code</u> , Prod_Desc, Unit_Px, Supplier_ID, Supplier_Name)
Email, Prod_Code -> Qty, Ttl_Amt	Customer_Product_Supplier ( <u>Email</u> , <u>Prod_Code</u> , Qty, Ttl_Amt

For each of the partial dependencies identified, create a relation and name them appropriately. Check there are no anomalies in each of the relations.

Step 2 – Derive relations from functional dependencies

## Normalization – Transitive Dependency

- Transitive Dependency is a functional dependency between two (or more) non-key attributes in a relation
  - A relation is said to be in Third Normal Form if there is no functional dependencies among non-key attributes

# Normalization - Transform 2NF to 3NF

Relations (Relational Header Format)	3NF?	Transformation
Customer ( <u>Email</u> , First_Name)	Yes	
Product ( <u>Prod_Code</u> , Prod_Desc, Unit_Px, Supplier_ID, Supplier_Name)	No	Supplier (Supplier_ID, Supplier_Name)  NB: Supplier_ID becomes foreign key ref Supplier  Product (Prod_Code, Prod_Desc, Unix_Px, Supplier_ID)
Customer_Product_Supplier ( <u>Email</u> , <u>Prod_Code</u> , Qty, Ttl_Amt	Yes	

#### Normalization – From 0NF To 3NF

Customer (Email, First\_Name)

Supplier (Supplier\_ID, Supplier\_Name)

Product (Prod\_Code, Prod\_Desc, Unix\_Px, Supplier\_ID)

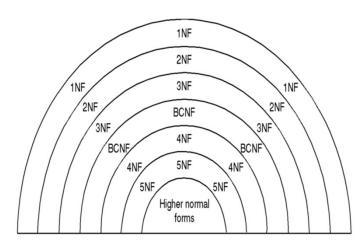
Customer\_Product\_Supplier (Email, Prod\_Code, Qty, Ttl\_Amt)

We started with one table and after the 3<sup>rd</sup> normalization, we have four tables.

# Normalization – Summary

Normal Form	Definition	Problem	Solution
UNF/0NF	A table with multi-valued cells or repeating groups	Violate the properties of relation	Form new row for each multi-valued cell to remove repeating group.
1NF	A table without any repeating groups of data	Contain 3 types of anomalies:  a. Insert anomaly  b. Update anomaly  c. Delete anomaly	Remove partial dependencies.  Determination of primary key
2NF	A relation that is in the 1NF, and, every non-key attribute is functionally dependent on the whole of its primary key	Functional dependency among the non-key attributes	Create new relation(s). The primary key of the new relation will be served as the foreign key in the existing relation.
3NF	A relation that is in 2NF, and, there is no functional dependency among non-key attributes		

## Normalization – Relationships Between Normal Forms



- Normalisation process splits the information across several relations
- Aims to establish relations that are more efficient when we perform insert, update and delete records operations on the relation
- ❖ A process of grouping attributes into wellstructured relations that allow users to insert, delete and modify rows in these relations without errors or inconsistencies resulting from these operations.

# THE END

# Optional Discussion

## Entity Relationship Diagram – Weak Entity and Identification Dependency

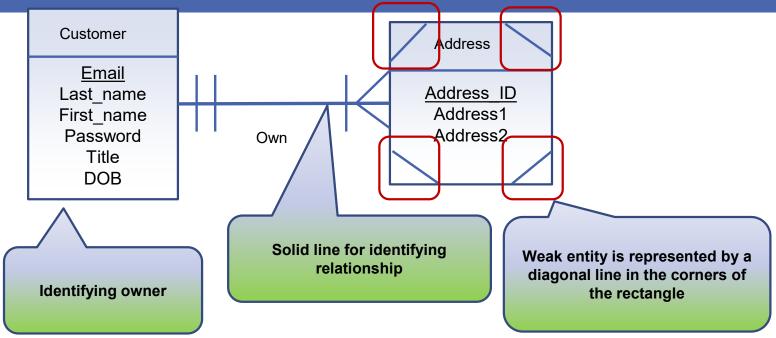
Address ID
Address1
Address2

# Weak entity or Child entity

Dependent on parent/owner entity for existence

Cannot be identified using its own attributes only

Entity Relationship Diagram – Weak Entity and Identification Dependency

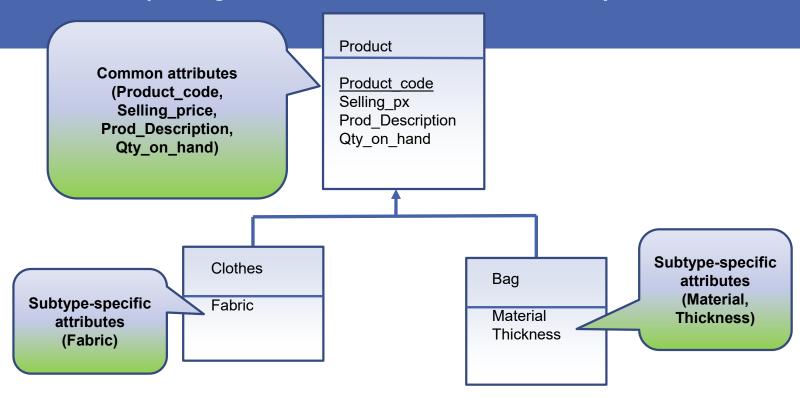


- The **Address** entity is dependent on customer entity for existence.
- Address\_ID needs to combine with Email to form the full-identifier.
  - Address\_ID + Email = Primary Key.

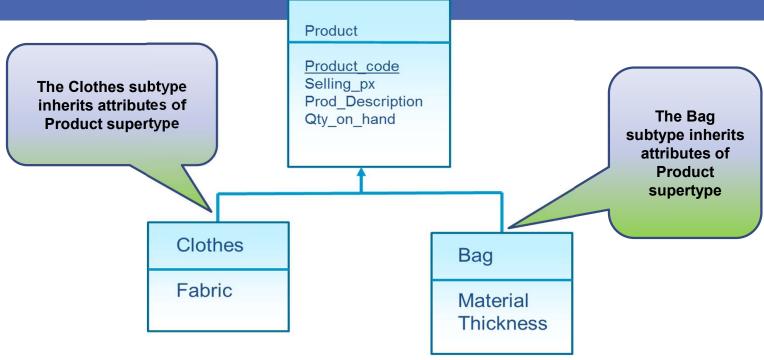
STI501 DENG

# Entity Relationship Diagram - Generalization Hierarchy **Supertype** Supertype (General entity type) General hierarchy symbol (A line connecting the subtypes, an arrow connecting the supertype) Subtype1 Subtype2 **Subtype** (Specialised entity type)

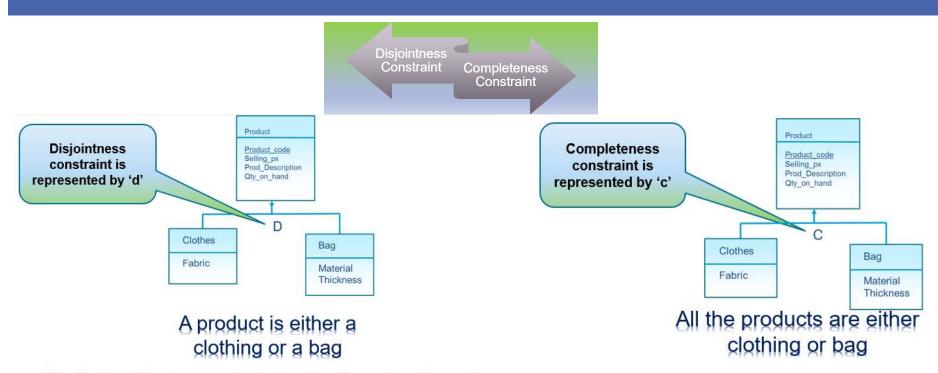
- Also known as IS-A relationship.
- Collection of entity types arrange in hierarchy structure.



- Attributes that are common to all entities are associated with the Supertype.
- Attributes that are unique to a particular subtype are associated with that subtype.

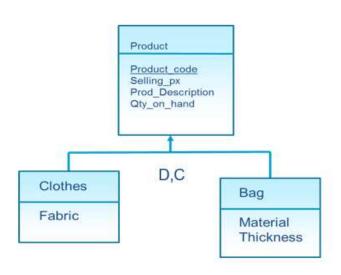


- Subtype entities inherits attributes of the supertype.
- Subtype possesses subtype-specific attributes as well as those associated with supertype attributes.

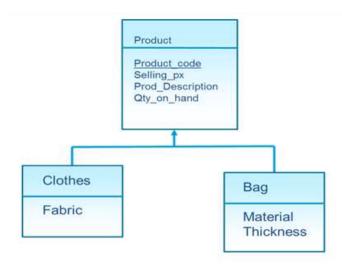


 Specifies that a <u>Supertype</u> can only be a member of one subtype; it cannot be simultaneously be a member of both.

Specifies all the subtypes that make up a supertype entity.



The Disjointness constraint shows that the product is either a clothing or a bag. The Completeness constraint shows that there are only 2 types of products: Clothes or Bag.



The omission of Disjoint constraint implies that a product can be both an article of clothing and a bag at the same time.

The omission of the Completeness constraint implies there are other types of products besides clothes and bags eg. shoes.