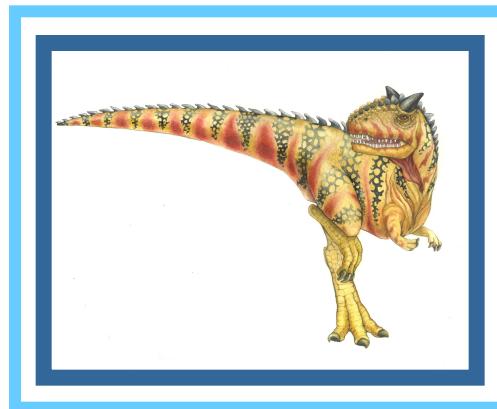


Chapter 9: Main Memory





Chapter 9: Memory Management

- Background
- Contiguous Memory Allocation
- Paging
- Structure of the Page Table
- Swapping
- Example: The Intel 32 and 64-bit Architectures
- Example: ARMv8 Architecture





Objectives

- To provide a detailed description of various ways of organizing memory hardware
- To discuss various memory-management techniques,
- To provide a detailed description of the Intel Pentium, which supports both pure segmentation and segmentation with paging





Background

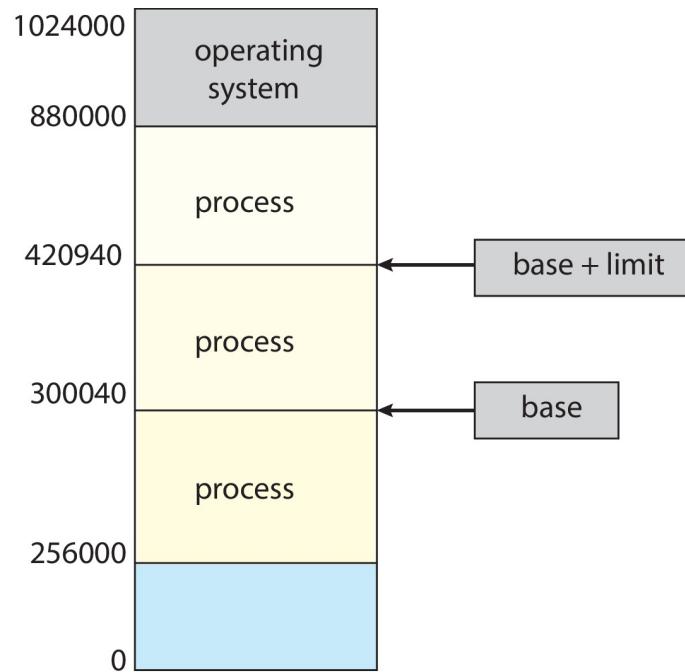
- Program must be brought (from disk) into memory and placed within a process for it to be run
- Main memory and registers are only storage CPU can access directly
- Memory unit only sees a stream of:
 - addresses + read requests, or
 - address + data and write requests
- Register access is done in one CPU clock (or less)
- Main memory can take many cycles, causing a **stall**
- **Cache** sits between main memory and CPU registers
- Protection of memory required to ensure correct operation





Protection

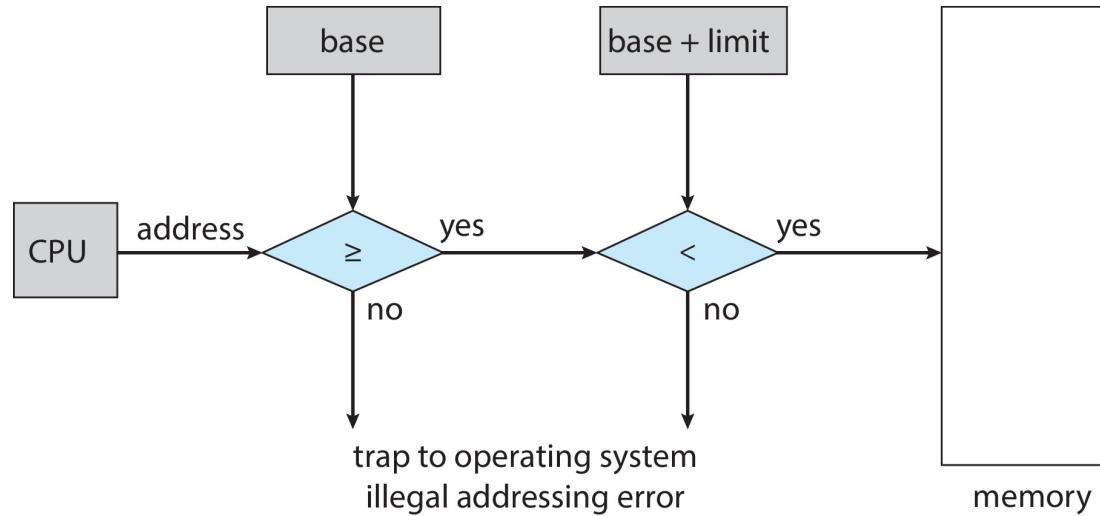
- Need to ensure that a process can access only those addresses in its address space.
- We can provide this protection by using a pair of **base** and **limit registers** define the logical address space of a process





Hardware Address Protection

- CPU must check every memory access generated in user mode to be sure it is between base and limit for that user



- the instructions to loading the base and limit registers are privileged





Address Binding

- Programs on disk, ready to be brought into memory to execute from an **input queue**
 - Without support, must be loaded into address 0000
- Inconvenient to have first user process physical address always at 0000
 - How can it not be?
- Addresses represented in different ways at different stages of a program's life
 - Source code addresses usually symbolic
 - Compiled code addresses **bind** to relocatable addresses
 - ▶ i.e., "14 bytes from beginning of this module"
 - Linker or loader will bind relocatable addresses to absolute addresses
 - ▶ i.e., 74014
 - Each binding maps one address space to another

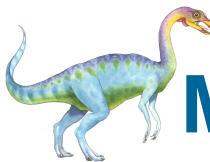




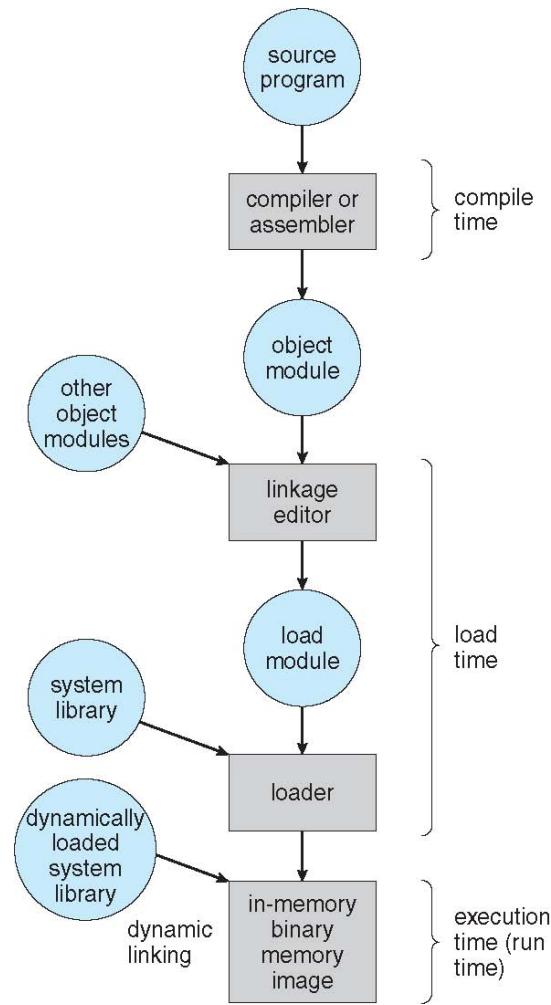
Binding of Instructions and Data to Memory

- Address binding of instructions and data to memory addresses can happen at three different stages
 - **Compile time:** If memory location known a priori, **absolute code** can be generated; must recompile code if starting location changes
 - **Load time:** Must generate **relocatable code** if memory location is not known at compile time
 - **Execution time:** Binding delayed until run time if the process can be moved during its execution from one memory segment to another
 - ▶ Need hardware support for address maps (e.g., base and limit registers)





Multistep Processing of a User Program





Logical vs. Physical Address Space

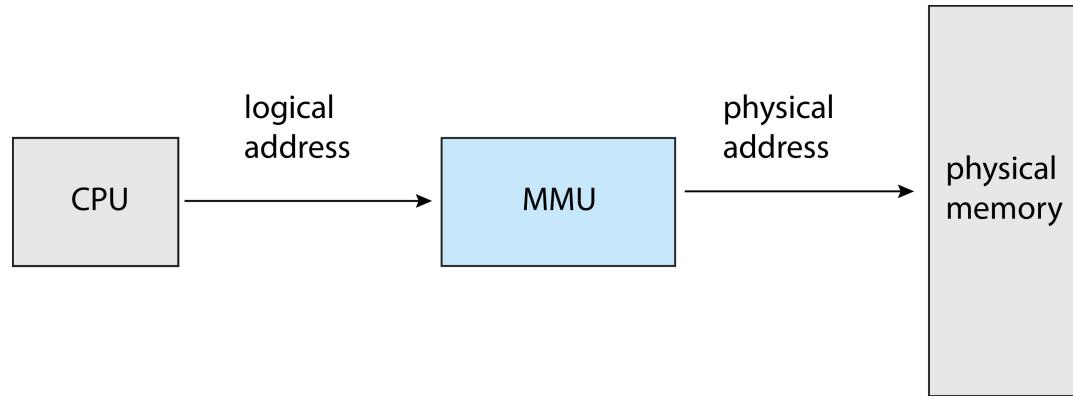
- The concept of a logical address space that is bound to a separate **physical address space** is central to proper memory management
 - **Logical address** – generated by the CPU; also referred to as **virtual address**
 - **Physical address** – address seen by the memory unit
- Logical and physical addresses are the same in compile-time and load-time address-binding schemes; logical (virtual) and physical addresses differ in execution-time address-binding scheme
- **Logical address space** is the set of all logical addresses generated by a program
- **Physical address space** is the set of all physical addresses generated by a program





Memory-Management Unit (MMU)

- Hardware device that at run time maps virtual to physical address



- Many methods possible, covered in the rest of this chapter





Memory-Management Unit (Cont.)

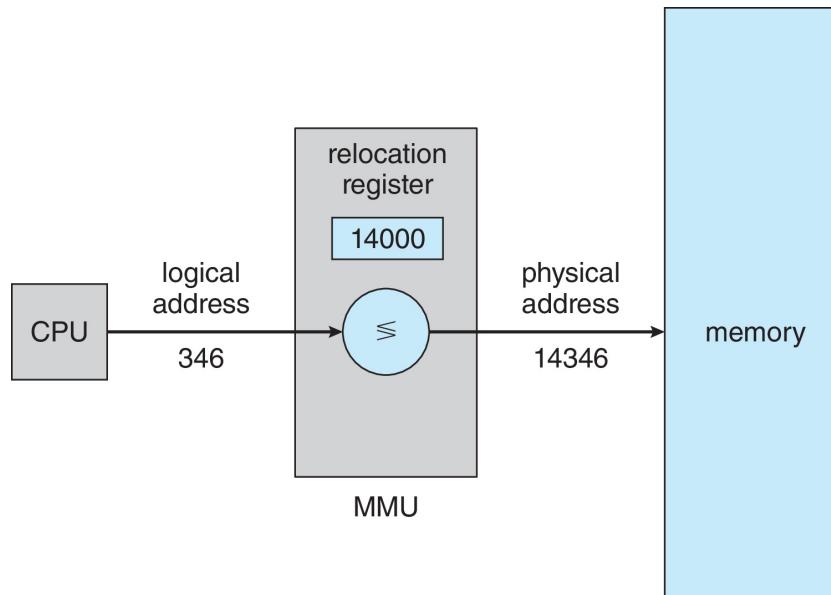
- Consider simple scheme, which is a generalization of the base-register scheme.
- The base register now called **relocation register**
- The value in the relocation register is added to every address generated by a user process at the time it is sent to memory
- The user program deals with *logical* addresses; it never sees the *real* physical addresses
 - Execution-time binding occurs when reference is made to location in memory
 - Logical address bound to physical addresses





Memory-Management Unit (Cont.)

- Consider simple scheme, which is a generalization of the base-register scheme.
- The base register now called **relocation register**
- The value in the relocation register is added to every address generated by a user process at the time it is sent to memory





Dynamic Loading

- The entire program does need to be in memory to execute
- Routine is not loaded until it is called
- Better memory-space utilization; unused routine is never loaded
- All routines kept on disk in relocatable load format
- Useful when large amounts of code are needed to handle infrequently occurring cases
- No special support from the operating system is required
 - Implemented through program design
 - OS can help by providing libraries to implement dynamic loading





Dynamic Linking

- **Static linking** – system libraries and program code combined by the loader into the binary program image
- Dynamic linking –linking postponed until execution time
- Small piece of code, **stub**, used to locate the appropriate memory-resident library routine
- Stub replaces itself with the address of the routine, and executes the routine
- Operating system checks if routine is in processes' memory address
 - If not in address space, add to address space
- Dynamic linking is particularly useful for libraries
- System also known as **shared libraries**
- Consider applicability to patching system libraries
 - Versioning may be needed





Contiguous Allocation

- Main memory must support both OS and user processes
- Limited resource, must allocate efficiently
- Contiguous allocation is one early method
- Main memory usually into two **partitions**:
 - Resident operating system, usually held in low memory with interrupt vector
 - User processes then held in high memory
 - Each process contained in single contiguous section of memory

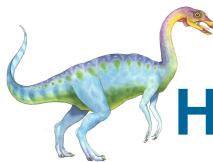




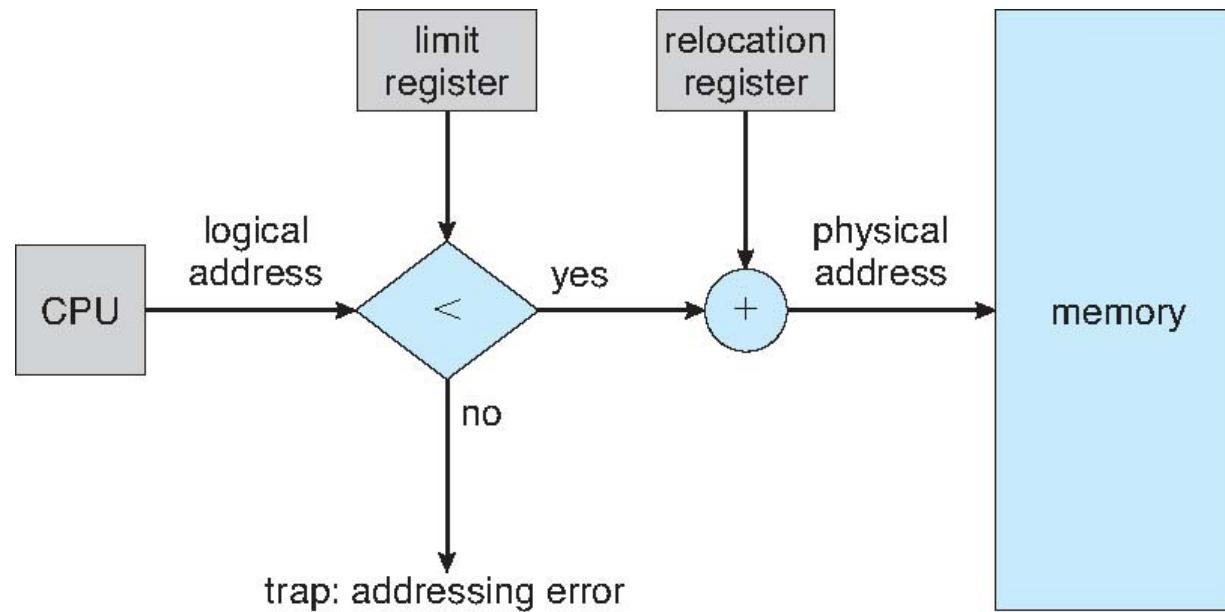
Contiguous Allocation (Cont.)

- Relocation registers used to protect user processes from each other, and from changing operating-system code and data
 - Base register contains value of smallest physical address
 - Limit register contains range of logical addresses – each logical address must be less than the limit register
 - MMU maps logical address *dynamically*
 - Can then allow actions such as kernel code being **transient** and kernel changing size





Hardware Support for Relocation and Limit Registers

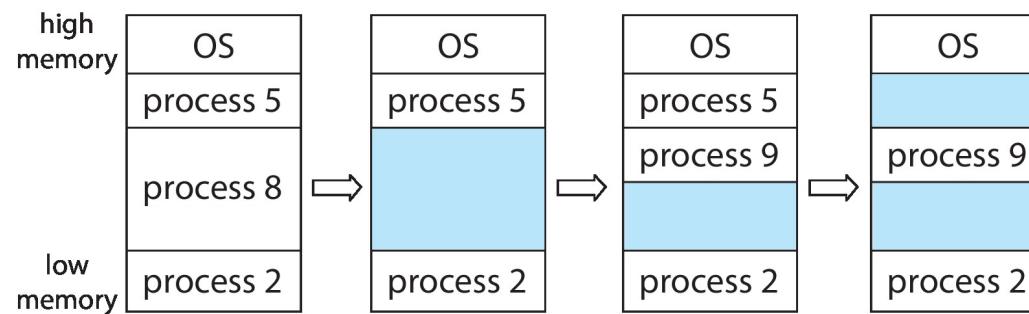




Variable Partition

- Multiple-partition allocation

- Degree of multiprogramming limited by number of partitions
- **Variable-partition** sizes for efficiency (sized to a given process' needs)
- **Hole** – block of available memory; holes of various size are scattered throughout memory
- When a process arrives, it is allocated memory from a hole large enough to accommodate it
- Process exiting frees its partition, adjacent free partitions combined
- Operating system maintains information about:
 - a) allocated partitions
 - b) free partitions (hole)





Dynamic Storage-Allocation Problem

How to satisfy a request of size n from a list of free holes?

- **First-fit:** Allocate the *first* hole that is big enough
- **Best-fit:** Allocate the *smallest* hole that is big enough; must search entire list, unless ordered by size
 - Produces the smallest leftover hole
- **Worst-fit:** Allocate the *largest* hole; must also search entire list
 - Produces the largest leftover hole

First-fit and best-fit better than worst-fit in terms of speed and storage utilization





Fragmentation

- **External Fragmentation** – total memory space exists to satisfy a request, but it is not contiguous
- **Internal Fragmentation** – allocated memory may be slightly larger than requested memory; this size difference is memory internal to a partition, but not being used
- First fit analysis reveals that given N blocks allocated, $0.5 N$ blocks lost to fragmentation
 - 1/3 may be unusable -> **50-percent rule**





Fragmentation (Cont.)

- Reduce external fragmentation by **compaction**
 - Shuffle memory contents to place all free memory together in one large block
 - Compaction is possible *only* if relocation is dynamic, and is done at execution time
 - I/O problem
 - Latch job in memory while it is involved in I/O
 - Do I/O only into OS buffers
- Now consider that backing store has same fragmentation problems





Paging

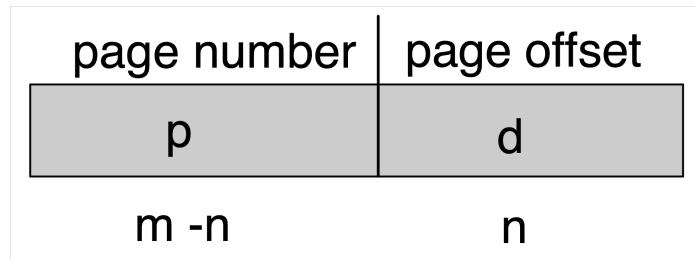
- Physical address space of a process can be noncontiguous; process is allocated physical memory whenever the latter is available
 - Avoids external fragmentation
 - Avoids problem of varying sized memory chunks
- Divide physical memory into fixed-sized blocks called **frames**
 - Size is power of 2, between 512 bytes and 16 Mbytes
- Divide logical memory into blocks of same size called **pages**
- Keep track of all free frames
- To run a program of size N pages, need to find N free frames and load program
- Set up a **page table** to translate logical to physical addresses
- Backing store likewise split into pages
- Still have Internal fragmentation





Address Translation Scheme

- Address generated by CPU is divided into:
 - **Page number (p)** – used as an index into a **page table** which contains base address of each page in physical memory
 - **Page offset (d)** – combined with base address to define the physical memory address that is sent to the memory unit

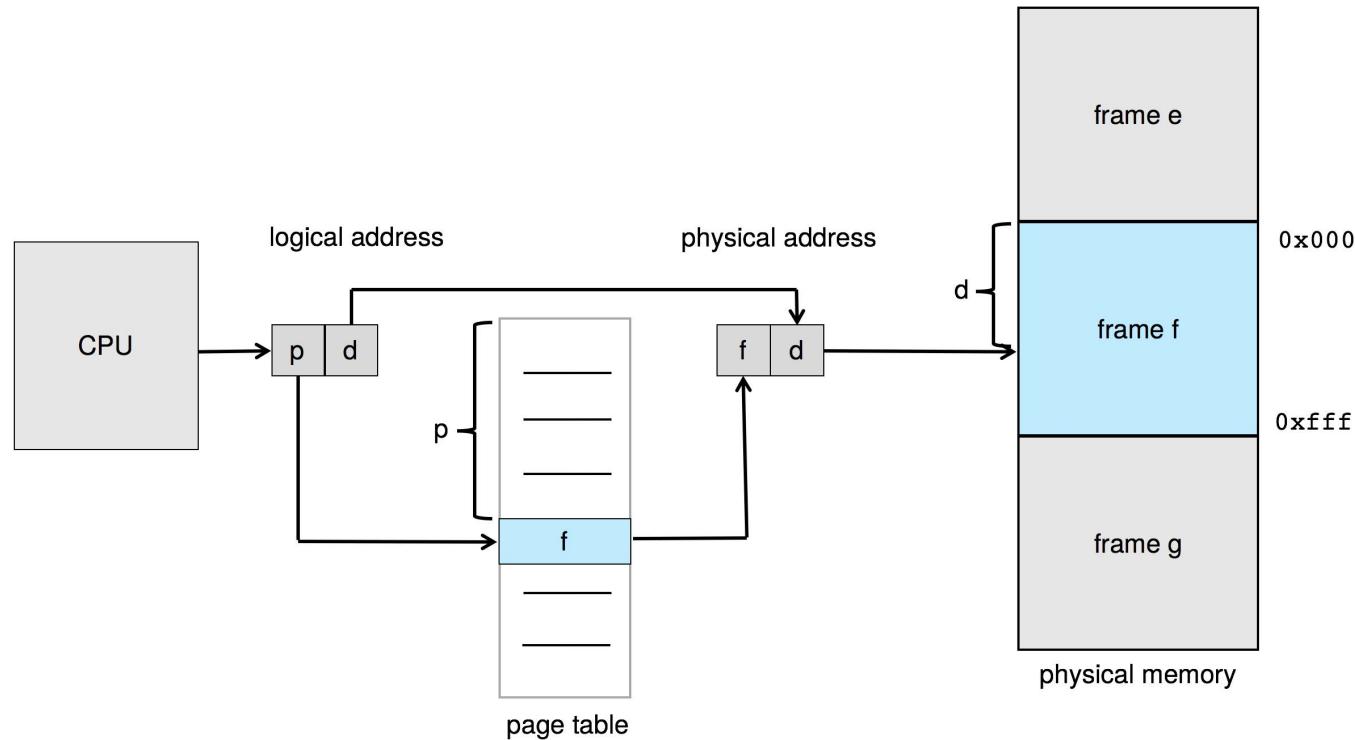


- For given logical address space 2^m and page size 2^n



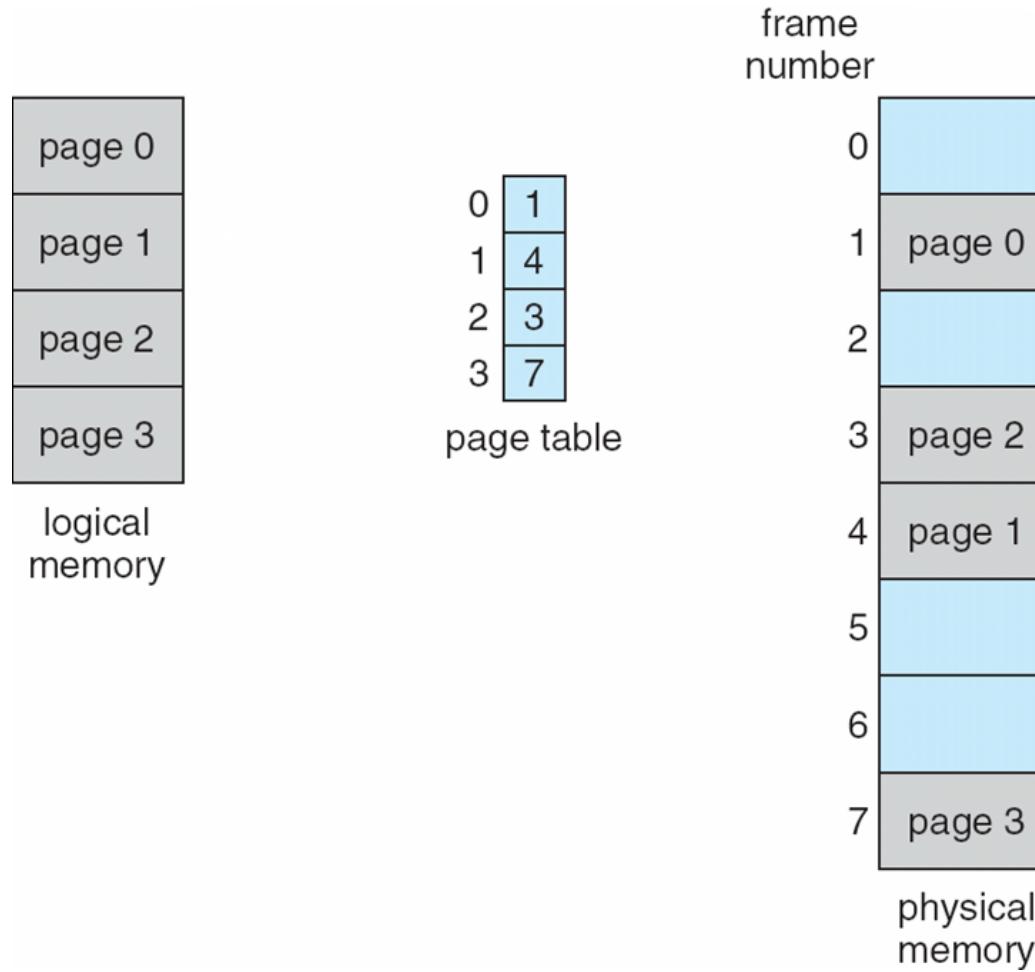


Paging Hardware





Paging Model of Logical and Physical Memory





Paging Example

- Logical address: $n = 2$ and $m = 4$. Using a page size of 4 bytes and a physical memory of 32 bytes (8 pages)

0	a
1	b
2	c
3	d
4	e
5	f
6	g
7	h
8	i
9	j
10	k
11	l
12	m
13	n
14	o
15	p

logical memory

0	5
1	6
2	1
3	2

page table

0	
4	
8	
12	
16	
20	a
21	b
22	c
23	d
24	e
25	f
26	g
27	h

physical memory





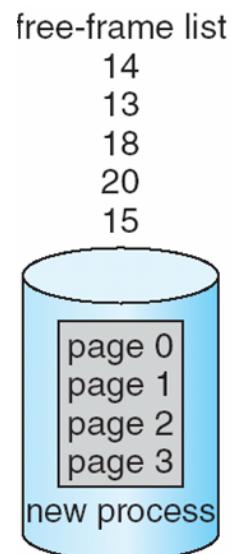
Paging -- Calculating internal fragmentation

- Page size = 2,048 bytes
- Process size = 72,766 bytes
- 35 pages + 1,086 bytes
- Internal fragmentation of $2,048 - 1,086 = 962$ bytes
- Worst case fragmentation = 1 frame – 1 byte
- On average fragmentation = 1 / 2 frame size
- So small frame sizes desirable?
- But each page table entry takes memory to track
- Page sizes growing over time
 - Solaris supports two page sizes – 8 KB and 4 MB



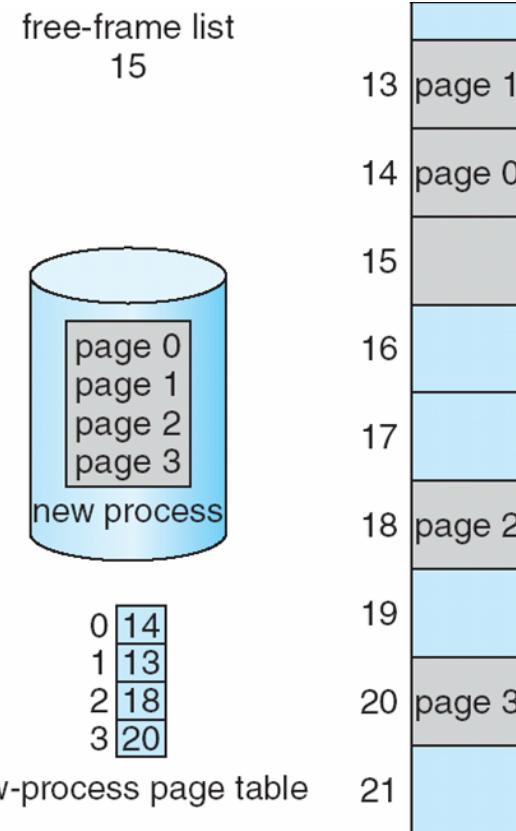


Free Frames



(a)

Before allocation



(b)

After allocation





Implementation of Page Table

- Page table is kept in main memory
 - **Page-table base register (PTBR)** points to the page table
 - **Page-table length register (PTLR)** indicates size of the page table
- In this scheme every data/instruction access requires two memory accesses
 - One for the page table and one for the data / instruction
- The two-memory access problem can be solved by the use of a special fast-lookup hardware cache called **translation look-aside buffers (TLBs)** (also called **associative memory**).





Translation Look-Aside Buffer

- Some TLBs store **address-space identifiers (ASIDs)** in each TLB entry – uniquely identifies each process to provide address-space protection for that process
 - Otherwise need to flush at every context switch
- TLBs typically small (64 to 1,024 entries)
- On a TLB miss, value is loaded into the TLB for faster access next time
 - Replacement policies must be considered
 - Some entries can be **wired down** for permanent fast access





Hardware

- Associative memory – parallel search

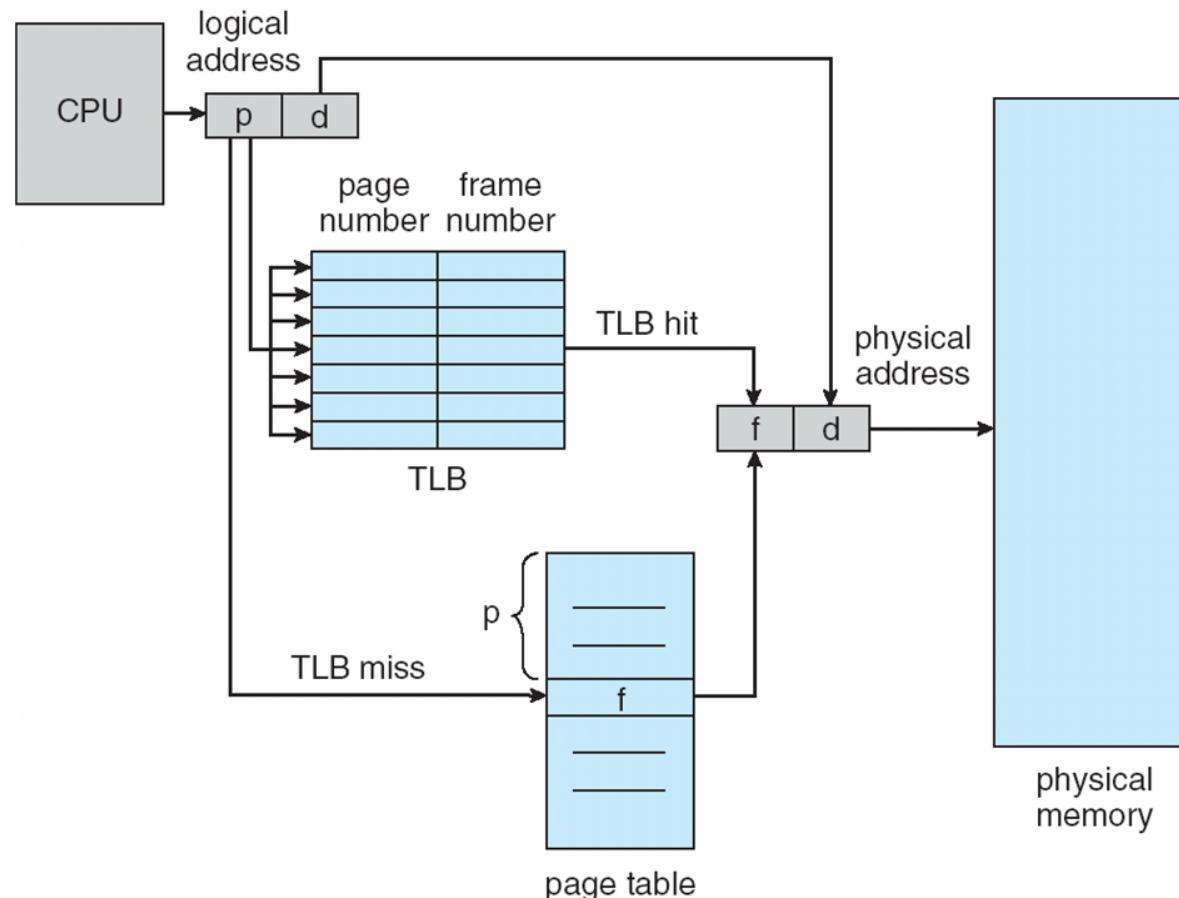
Page #	Frame #

- Address translation (p, d)
 - If p is in associative register, get frame # out
 - Otherwise get frame # from page table in memory





Paging Hardware With TLB





Effective Access Time

- Hit ratio – percentage of times that a page number is found in the TLB
- An 80% hit ratio means that we find the desired page number in the TLB 80% of the time.
- Suppose that 10 nanoseconds to access memory.
 - If we find the desired page in TLB then a mapped-memory access take 10 ns
 - Otherwise we need two memory access so it is 20 ns
- **Effective Access Time (EAT)**

$$EAT = 0.80 \times 10 + 0.20 \times 20 = 12 \text{ nanoseconds}$$

implying 20% slowdown in access time

- Consider a more realistic hit ratio of 99%,

$$EAT = 0.99 \times 10 + 0.01 \times 20 = 10.1 \text{ ns}$$

implying only 1% slowdown in access time.





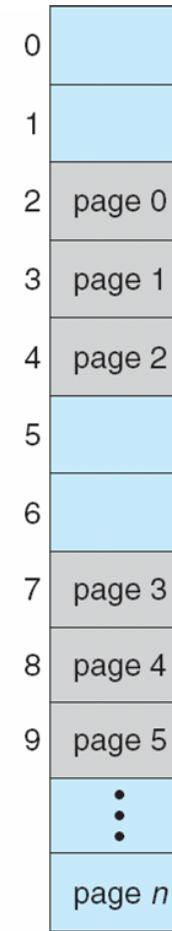
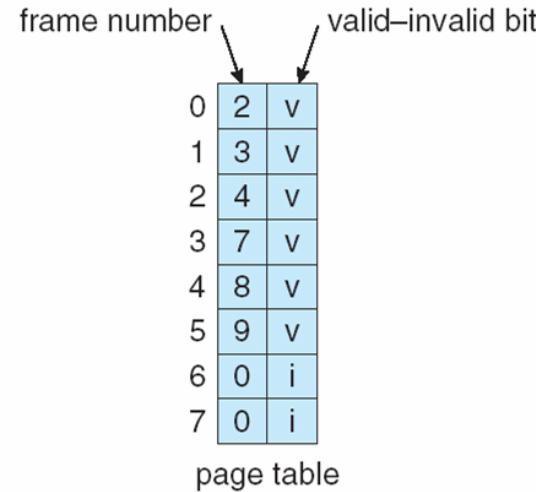
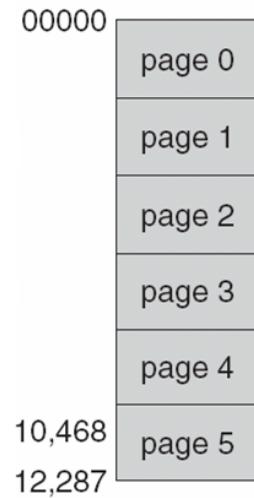
Memory Protection

- Memory protection implemented by associating protection bit with each frame to indicate if read-only or read-write access is allowed
 - Can also add more bits to indicate page execute-only, and so on
- **Valid-invalid** bit attached to each entry in the page table:
 - “valid” indicates that the associated page is in the process’ logical address space, and is thus a legal page
 - “invalid” indicates that the page is not in the process’ logical address space
 - Or use **page-table length register (PTLR)**
- Any violations result in a trap to the kernel





Valid (v) or Invalid (i) Bit In A Page Table





Shared Pages

■ Shared code

- One copy of read-only (**reentrant**) code shared among processes (i.e., text editors, compilers, window systems)
- Similar to multiple threads sharing the same process space
- Also useful for interprocess communication if sharing of read-write pages is allowed

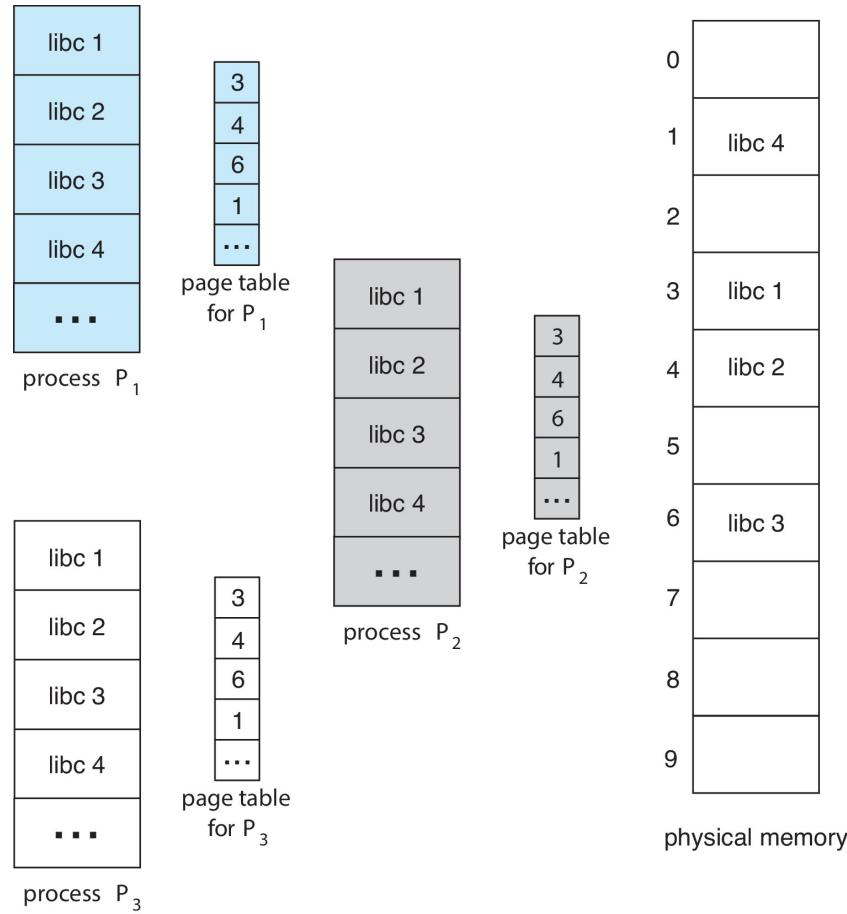
■ Private code and data

- Each process keeps a separate copy of the code and data
- The pages for the private code and data can appear anywhere in the logical address space





Shared Pages Example





Structure of the Page Table

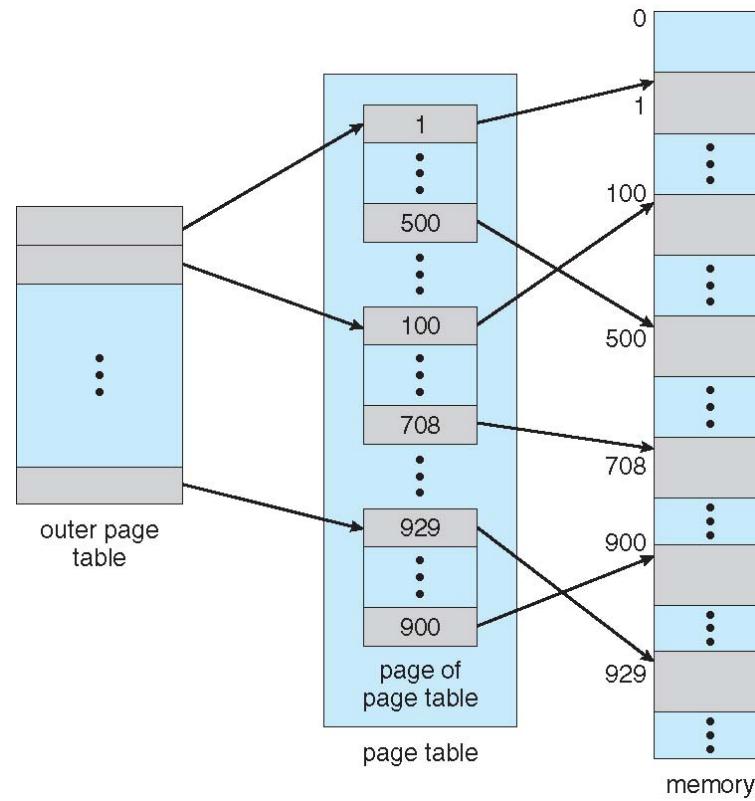
- Memory structures for paging can get huge using straight-forward methods
 - Consider a 32-bit logical address space as on modern computers
 - Page size of 4 KB (2^{12})
 - Page table would have 1 million entries ($2^{32} / 2^{12}$)
 - If each entry is 4 bytes → each process 4 MB of physical address space for the page table alone
 - Don't want to allocate that contiguously in main memory
 - One simple solution is to divide the page table into smaller units
 - Hierarchical Paging
 - Hashed Page Tables
 - Inverted Page Tables





Hierarchical Page Tables

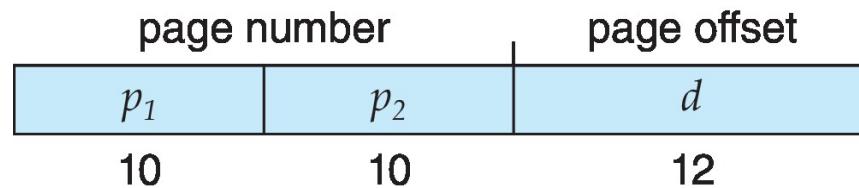
- Break up the logical address space into multiple page tables
- A simple technique is a two-level page table
- We then page the page table





Two-Level Paging Example

- A logical address (on 32-bit machine with 4K page size) is divided into:
 - a page number consisting of 20 bits
 - a page offset consisting of 12 bits
- Since the page table is paged, the page number is further divided into:
 - a 10-bit page number
 - a 10-bit page offset
- Thus, a logical address is as follows:

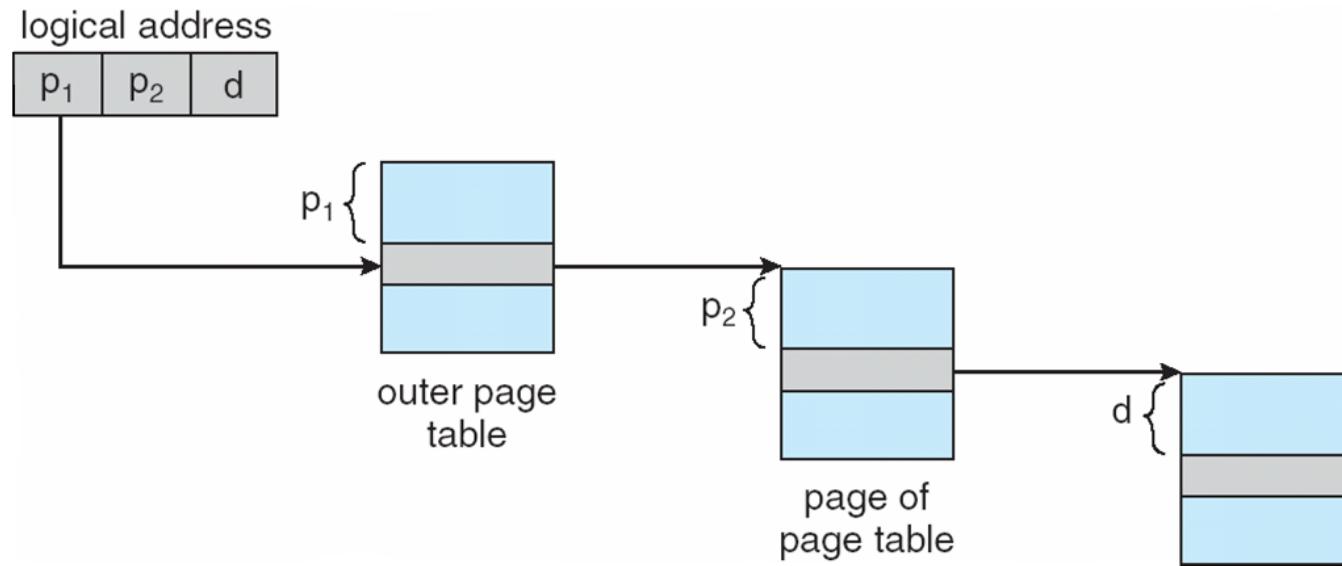


- where p_1 is an index into the outer page table, and p_2 is the displacement within the page of the inner page table
- Known as **forward-mapped page table**





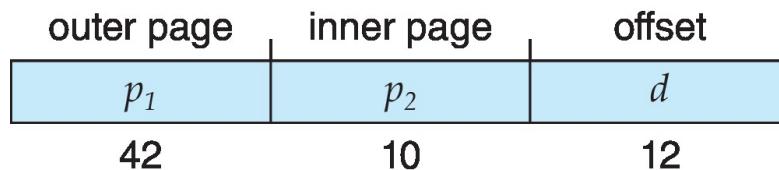
Address-Translation Scheme





64-bit Logical Address Space

- Even two-level paging scheme not sufficient
- If page size is 4 KB (2^{12})
 - Then page table has 2^{52} entries
 - If two level scheme, inner page tables could be 2^{10} 4-byte entries
 - Address would look like

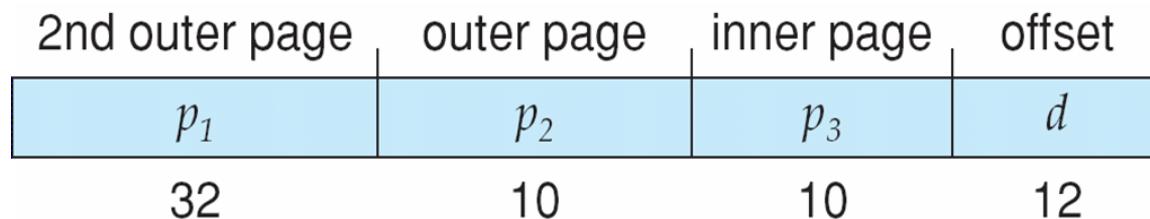
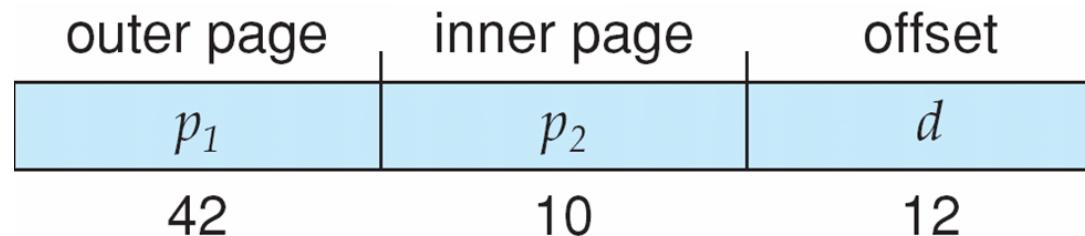


- Outer page table has 2^{42} entries or 2^{44} bytes
- One solution is to add a 2nd outer page table
- But in the following example the 2nd outer page table is still 2^{34} bytes in size
 - ▶ And possibly 4 memory access to get to one physical memory location





Three-level Paging Scheme





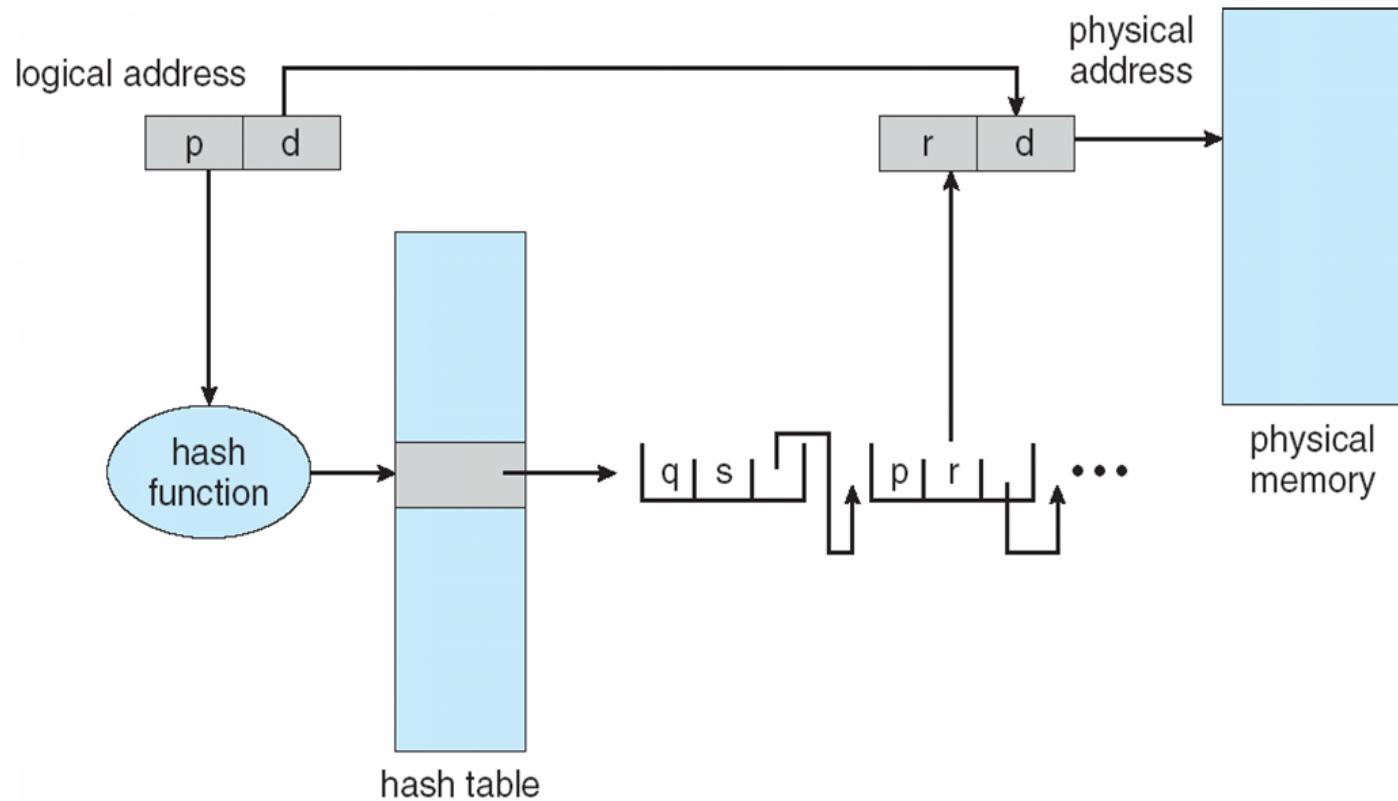
Hashed Page Tables

- Common in address spaces > 32 bits
- The virtual page number is hashed into a page table
 - This page table contains a chain of elements hashing to the same location
- Each element contains (1) the virtual page number (2) the value of the mapped page frame (3) a pointer to the next element
- Virtual page numbers are compared in this chain searching for a match
 - If a match is found, the corresponding physical frame is extracted
- Variation for 64-bit addresses is **clustered page tables**
 - Similar to hashed but each entry refers to several pages (such as 16) rather than 1
 - Especially useful for **sparse** address spaces (where memory references are non-contiguous and scattered)





Hashed Page Table





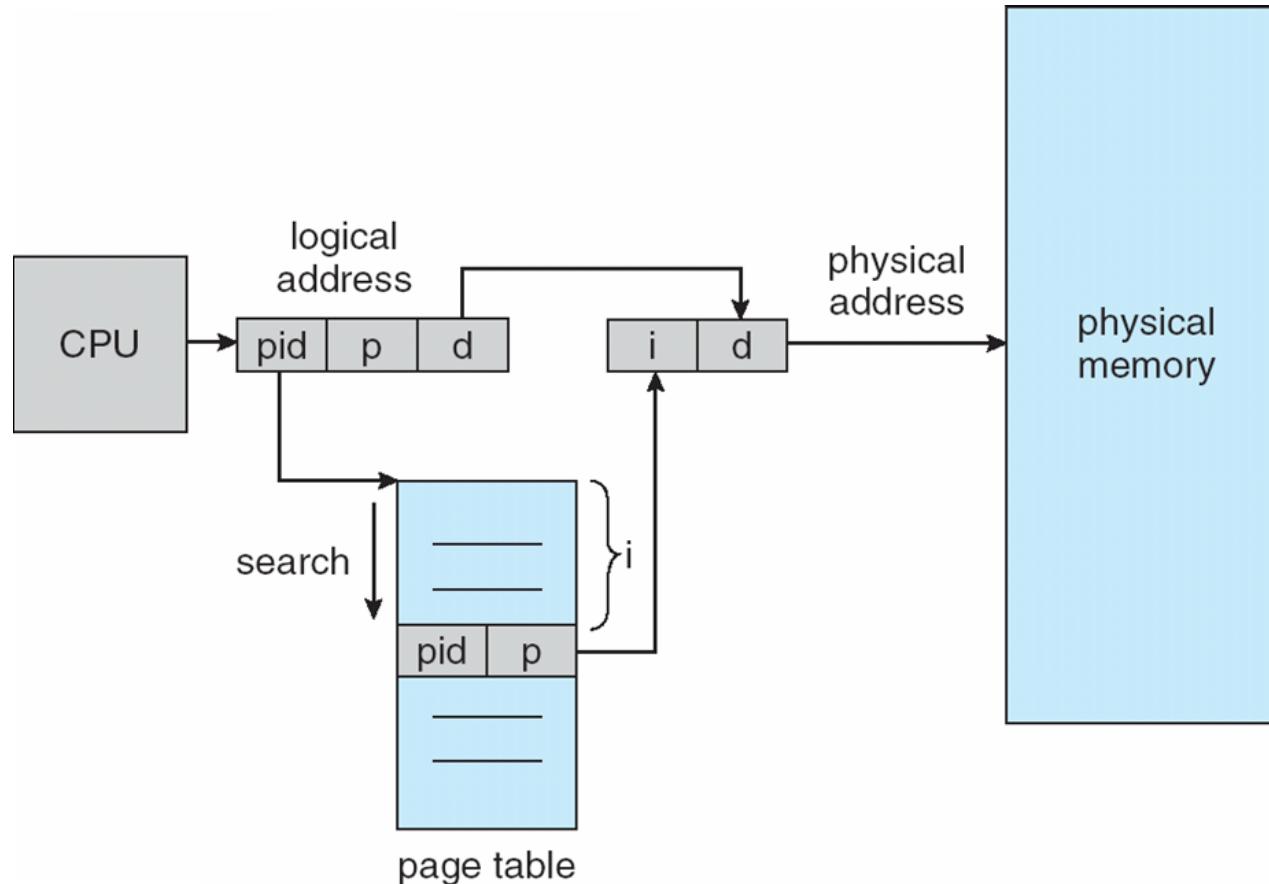
Inverted Page Table

- Rather than each process having a page table and keeping track of all possible logical pages, track all physical pages
- One entry for each real page of memory
- Entry consists of the virtual address of the page stored in that real memory location, with information about the process that owns that page
- Decreases memory needed to store each page table, but increases time needed to search the table when a page reference occurs
- Use hash table to limit the search to one — or at most a few — page-table entries
 - TLB can accelerate access
- But how to implement shared memory?
 - One mapping of a virtual address to the shared physical address





Inverted Page Table Architecture





Oracle SPARC Solaris

- Consider modern, 64-bit operating system example with tightly integrated HW
 - Goals are efficiency, low overhead
- Based on hashing, but more complex
- Two hash tables
 - One kernel and one for all user processes
 - Each maps memory addresses from virtual to physical memory
 - Each entry represents a contiguous area of mapped virtual memory,
 - ▶ More efficient than having a separate hash-table entry for each page
 - Each entry has base address and span (indicating the number of pages the entry represents)





Oracle SPARC Solaris (Cont.)

- TLB holds translation table entries (TTEs) for fast hardware lookups
 - A cache of TTEs reside in a translation storage buffer (TSB)
 - ▶ Includes an entry per recently accessed page
- Virtual address reference causes TLB search
 - If miss, hardware walks the in-memory TSB looking for the TTE corresponding to the address
 - ▶ If match found, the CPU copies the TSB entry into the TLB and translation completes
 - ▶ If no match found, kernel interrupted to search the hash table
 - The kernel then creates a TTE from the appropriate hash table and stores it in the TSB, Interrupt handler returns control to the MMU, which completes the address translation.





Swapping

- A process can be **swapped** temporarily out of memory to a backing store, and then brought **back** into memory for continued execution
 - Total physical memory space of processes can exceed physical memory
- **Backing store** – fast disk large enough to accommodate copies of all memory images for all users; must provide direct access to these memory images
- **Roll out, roll in** – swapping variant used for priority-based scheduling algorithms; lower-priority process is swapped out so higher-priority process can be loaded and executed
- Major part of swap time is transfer time; total transfer time is directly proportional to the amount of memory swapped
- System maintains a **ready queue** of ready-to-run processes which have memory images on disk





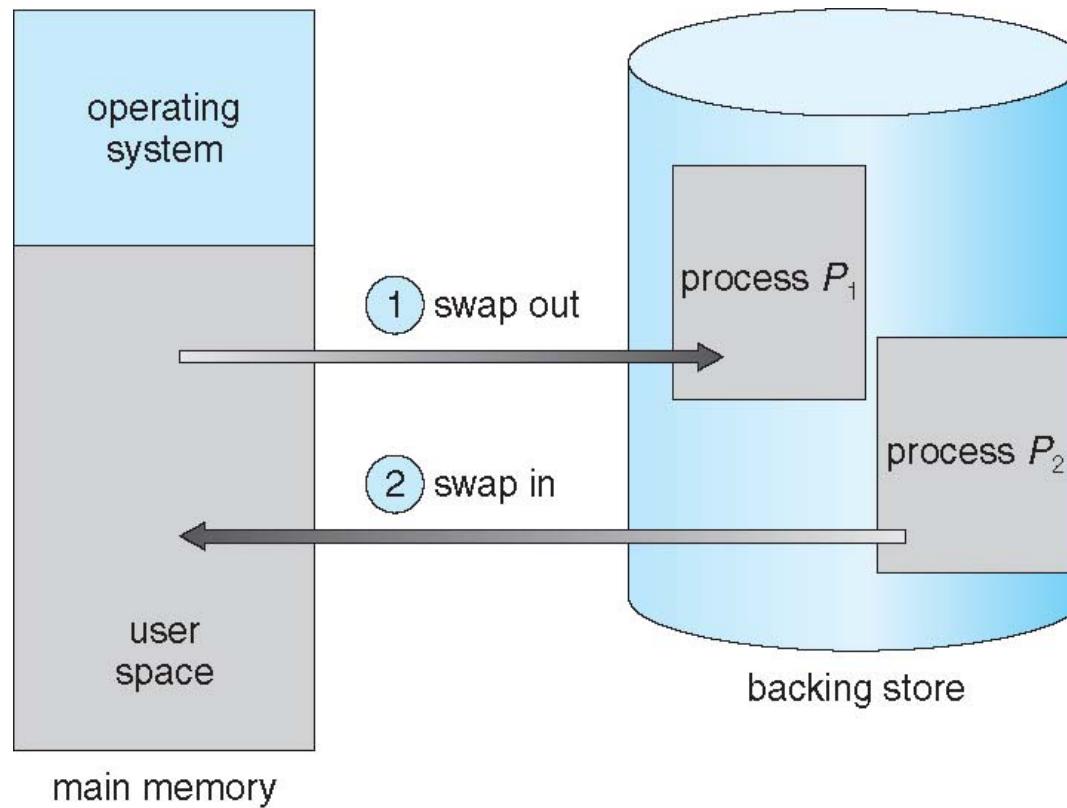
Swapping (Cont.)

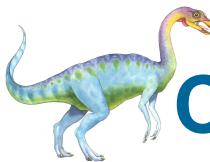
- Does the swapped out process need to swap back in to same physical addresses?
- Depends on address binding method
 - Plus consider pending I/O to / from process memory space
- Modified versions of swapping are found on many systems (i.e., UNIX, Linux, and Windows)
 - Swapping normally disabled
 - Started if more than threshold amount of memory allocated
 - Disabled again once memory demand reduced below threshold





Schematic View of Swapping





Context Switch Time including Swapping

- If next processes to be put on CPU is not in memory, need to swap out a process and swap in target process
- Context switch time can then be very high
- 100MB process swapping to hard disk with transfer rate of 50MB/sec
 - Swap out time of 2000 ms
 - Plus swap in of same sized process
 - Total context switch swapping component time of 4000ms (4 seconds)
- Can reduce if reduce size of memory swapped – by knowing how much memory really being used
 - System calls to inform OS of memory use via `request_memory()` and `release_memory()`





Context Switch Time and Swapping (Cont.)

- Other constraints as well on swapping
 - Pending I/O – can't swap out as I/O would occur to wrong process
 - Or always transfer I/O to kernel space, then to I/O device
 - ▶ Known as **double buffering**, adds overhead
- Standard swapping not used in modern operating systems
 - But modified version common
 - ▶ Swap only when free memory extremely low





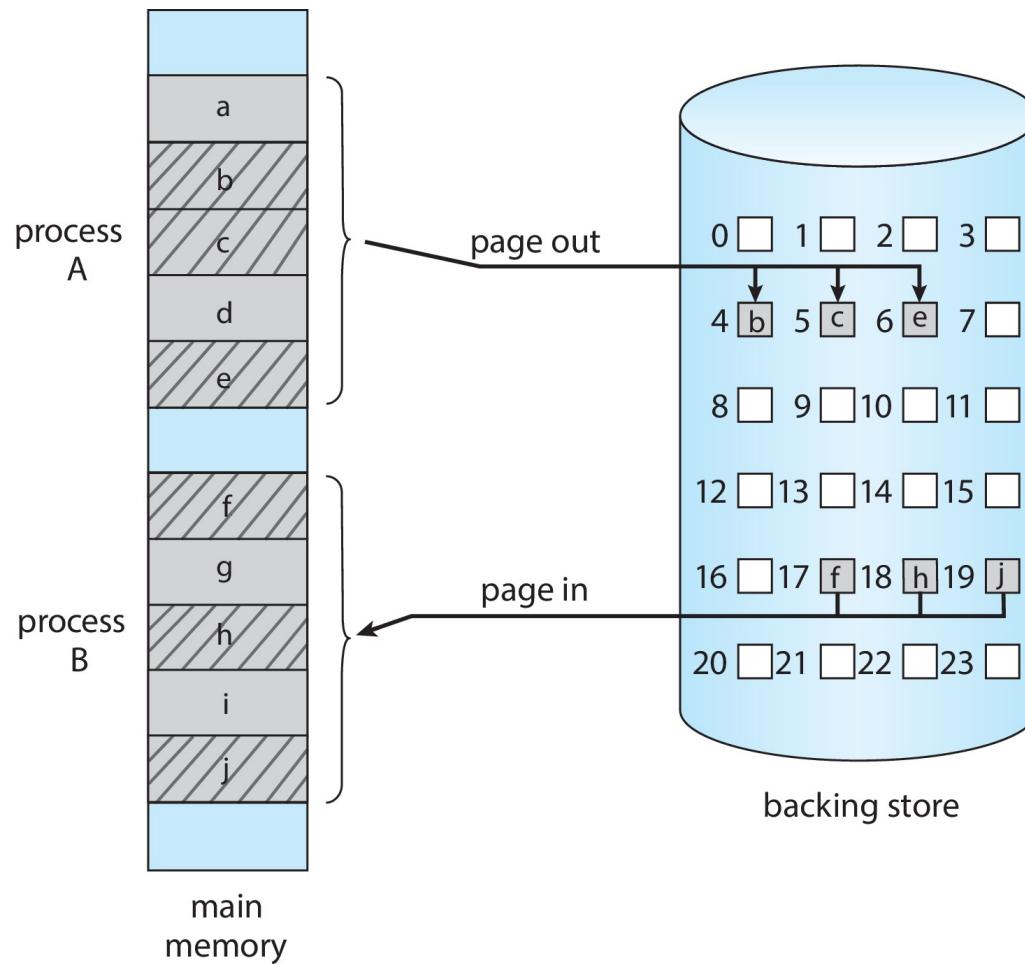
Swapping on Mobile Systems

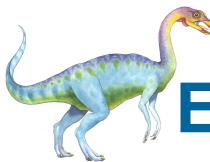
- Not typically supported
 - Flash memory based
 - ▶ Small amount of space
 - ▶ Limited number of write cycles
 - ▶ Poor throughput between flash memory and CPU on mobile platform
- Instead use other methods to free memory if low
 - iOS **asks** apps to voluntarily relinquish allocated memory
 - ▶ Read-only data thrown out and reloaded from flash if needed
 - ▶ Failure to free can result in termination
 - Android terminates apps if low free memory, but first writes **application state** to flash for fast restart
 - Both OSes support paging as discussed below





Swapping with Paging





Example: The Intel 32 and 64-bit Architectures

- Dominant industry chips
- Pentium CPUs are 32-bit and called IA-32 architecture
- Current Intel CPUs are 64-bit and called IA-64 architecture
- Many variations in the chips, cover the main ideas here

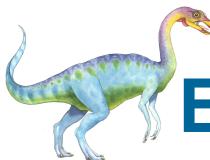




Example: The Intel IA-32 Architecture

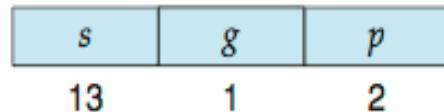
- Supports both segmentation and segmentation with paging
 - Each segment can be 4 GB
 - Up to 16 K segments per process
 - Divided into two partitions
 - ▶ First partition of up to 8 K segments are private to process (kept in **local descriptor table (LDT)**)
 - ▶ Second partition of up to 8K segments shared among all processes (kept in **global descriptor table (GDT)**)





Example: The Intel IA-32 Architecture (Cont.)

- CPU generates logical address
 - Selector given to segmentation unit
 - ▶ Which produces linear addresses

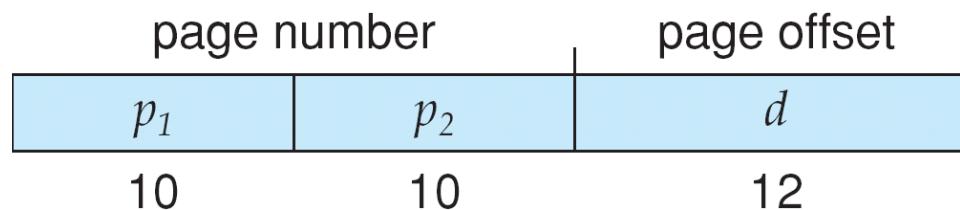
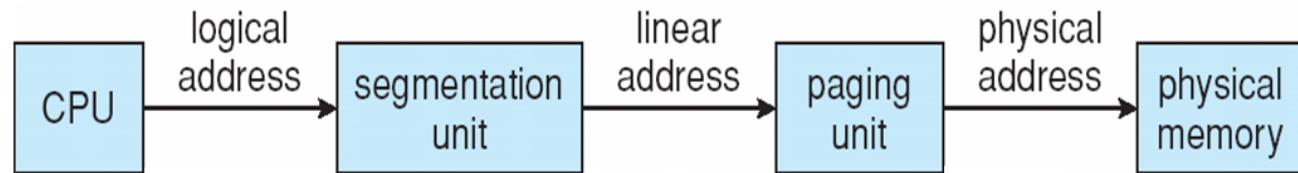


- Linear address given to paging unit
 - ▶ Which generates physical address in main memory
 - ▶ Paging units form equivalent of MMU
 - ▶ Pages sizes can be 4 KB or 4 MB



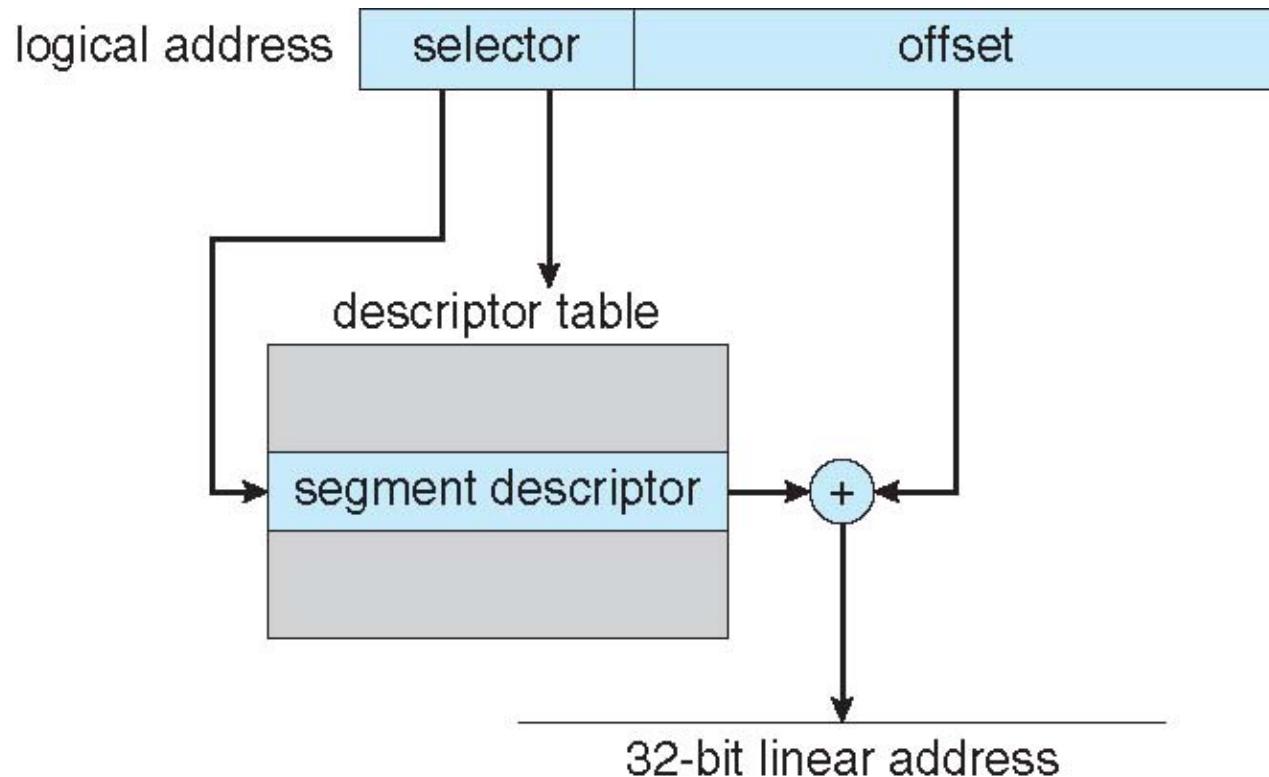


Logical to Physical Address Translation in IA-32



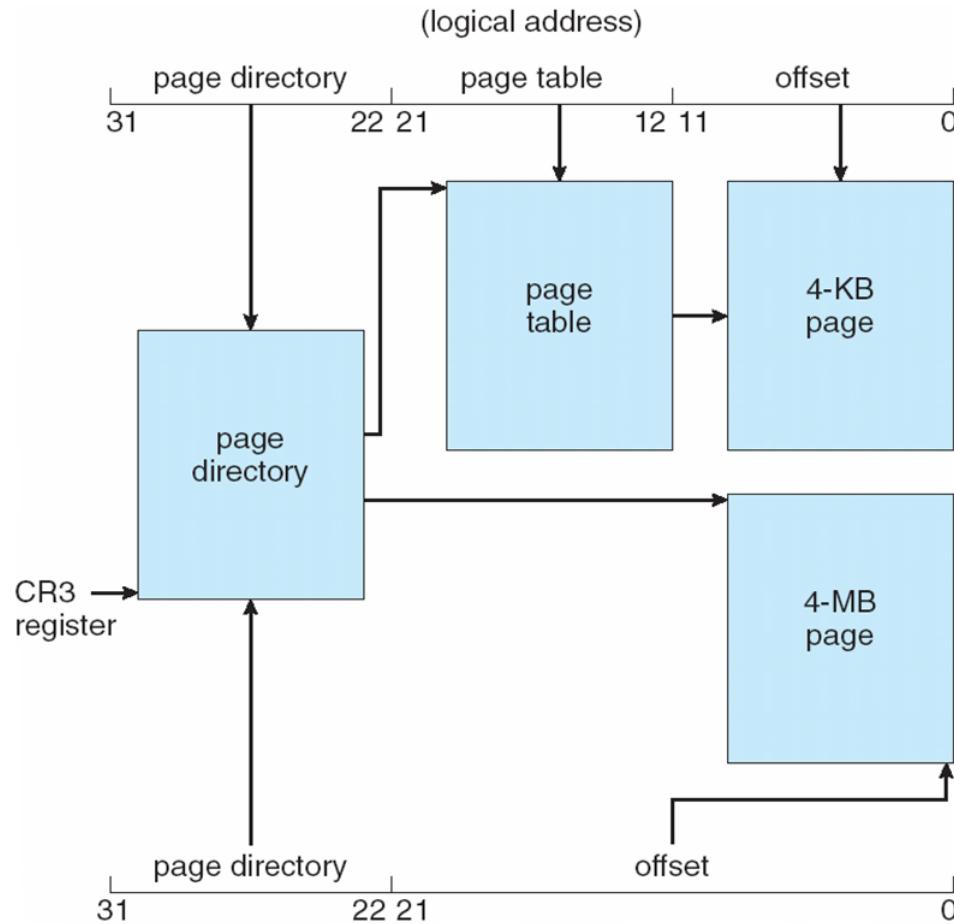


Intel IA-32 Segmentation





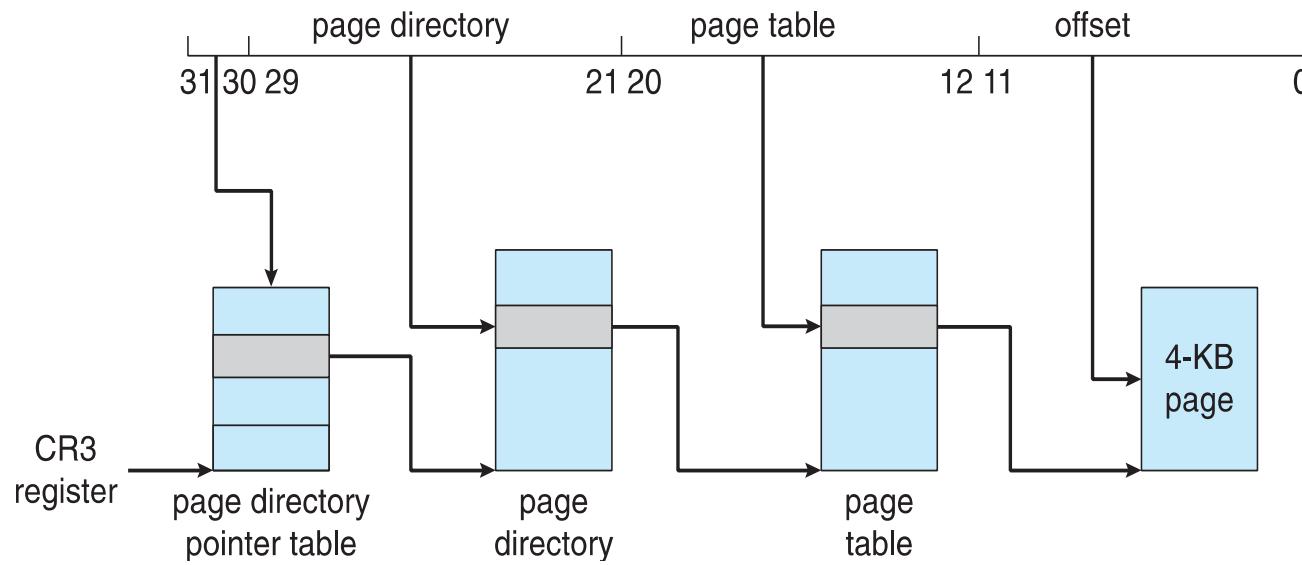
Intel IA-32 Paging Architecture





Intel IA-32 Page Address Extensions

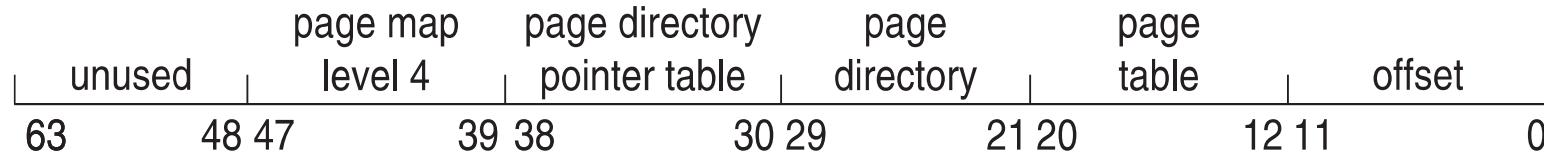
- 32-bit address limits led Intel to create **page address extension (PAE)**, allowing 32-bit apps access to more than 4GB of memory space
 - Paging went to a 3-level scheme
 - Top two bits refer to a **page directory pointer table**
 - Page-directory and page-table entries moved to 64-bits in size
 - Net effect is increasing address space to 36 bits – 64GB of physical memory





Intel x86-64

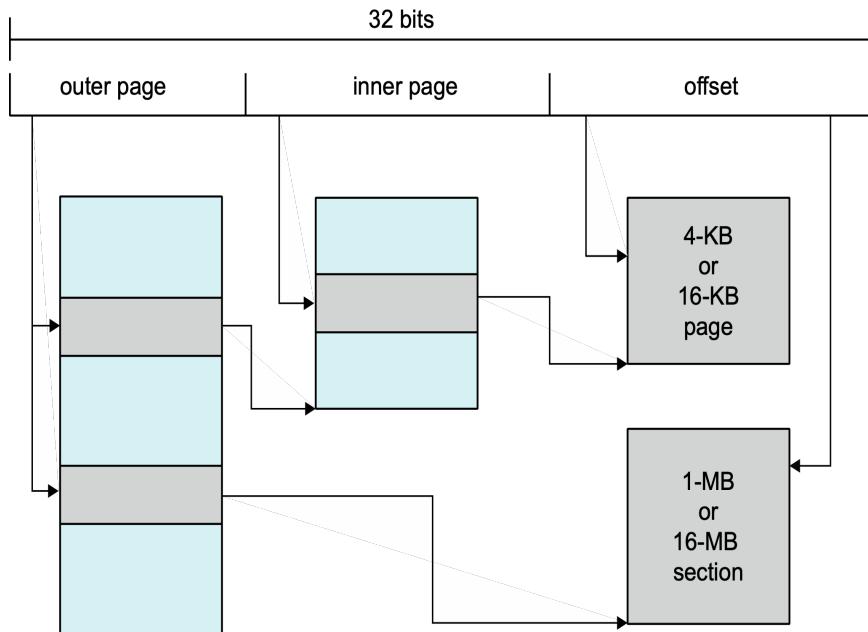
- Current generation Intel x86 architecture
- 64 bits is ginormous (> 16 exabytes)
- In practice only implement 48 bit addressing
 - Page sizes of 4 KB, 2 MB, 1 GB
 - Four levels of paging hierarchy
- Can also use PAE so virtual addresses are 48 bits and physical addresses are 52 bits





Example: ARM Architecture

- Dominant mobile platform chip (Apple iOS and Google Android devices for example)
- Modern, energy efficient, 32-bit CPU
- 4 KB and 16 KB pages
- 1 MB and 16 MB pages (termed **sections**)
- One-level paging for sections, two-level for smaller pages
- Two levels of TLBs
 - Outer level has two micro TLBs (one data, one instruction)
 - Inner is single main TLB
 - First inner is checked, on miss outers are checked, and on miss page table walk performed by CPU



End of Chapter 9

