

Code Generation for Arithmetical Expressions

$\text{GCodeAExp}(x)$	=	(ε, x)
$\text{GCodeAExp}(n)$	=	(ε, n)
$\text{GCodeAExp}(\text{tab}[i])$	=	Let $t1 = \text{newTemp}, t2 = \text{newTemp}$ in $(t1 := T*i \parallel$ $t2 := \text{tab}[t1], t2)$
$\text{GCodeAExp}(\text{tab}[i, j])$	=	Let $t1 = \text{newTemp}(), t2 = \text{newTemp}()$ $t3 = \text{newTemp}(), t4 = \text{newTemp}()$ $t5 = \text{newTemp}()$ in $(t1 := T*i \parallel$ $t2 := N \times T \parallel$ $t3 := t2 \times j \parallel$ $t4 := t1 + t3 \parallel$ $t5 := \text{tab}[t4], t5)$
$\text{GCodeAExp}(a_1 + a_2)$	=	Let $(C_1, t_1) = \text{GCodeAExp}(a_1),$ $(C_2, t_2) = \text{GCodeAExp}(a_2),$ $t = \text{newTemp}$ in $(C_1 \parallel C_2 \parallel t := t_1 + t_2, t)$

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Code Generation for Boolean Expressions

$\text{GCodeBExp}(a_1 < a_2, \text{true}, \text{false})$	=	Let $(C_1, t_1) = \text{GCodeAExp}(a_1),$ $(C_2, t_2) = \text{GCodeAExp}(a_2),$ in $C_1 \parallel C_2 \parallel$ if $t_1 < t_2$ goto ltrue goto lfalse
$\text{GCodeBExp}(b_1 \wedge b_2, \text{true}, \text{false})$	=	Let $l = \text{newLabel}()$ in $\text{GCodeBExp}(b_1, l, \text{false}) \parallel$ $l: \parallel$ $\text{GCodeBExp}(b_2, \text{true}, \text{false})$
$\text{GCodeBExp}(\neg b, \text{true}, \text{false})$	=	$\text{GCodeBExp}(b, \text{false}, \text{true})$

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Code Generation for Statements

Assignment and sequential composition

To each node of the abstract syntax tree, we associate code.

$\text{GCodeStm}(x := a)$	=	Let $(C, t) = \text{GCodeAExp}(a)$ in $C \parallel x := t$
$\text{GCodeStm}(\text{tab}[i] := a)$	=	Let $t1 = \text{newTemp},$ $(C, t) = \text{GCodeAExp}(a)$ in $(t1 := T*i \parallel$ $C \parallel \text{tab}[t1] := t)$
$\text{GCodeStm}(S_1 ; S_2)$	=	Let $C_1 = \text{GCodeStm}(S_1),$ $C_2 = \text{GCodeStm}(S_2)$ in $C_1 \parallel C_2$

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Code Generation for Statements

Iterative statement

$\text{GCodeStm}(\text{while } b \text{ do } S \text{ od})$	=	Let $l\text{begin} = \text{newLabel}(),$ $l\text{true} = \text{newLabel}(),$ $l\text{false} = \text{newLabel}()$ in $l\text{begin}: \parallel$ $\text{GCodeBExp}(b, l\text{true}, l\text{false}) \parallel$ $l\text{true}: \parallel$ $\text{GCodeStm}(S) \parallel$ goto lbegin lfalse:
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Code Generation for Statements

Conditional statement

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GCodeStm (if b then S1 else S2) = Let lnext=newLabel(),  
                                     ltrue=newLabel(),  
                                     lfalse=newLabel()  
                                     in GCodeBExp(b,ltrue,lfalse)||  
                                     ltrue:  
                                     GCodeStm(S1)||  
                                     goto lnext ||  
                                     lfalse:||  
                                     GCodeStm(S2)||  
                                     lnext:
```

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Summary - Intermediate-Code Generation

Intermediate-Code Generation

- ▶ From While to 3-address code.
- ▶ 3-address code = general-purpose representation of code:
 - ▶ easy to generate,
 - ▶ suitable for optimization,
 - ▶ easy to generate to assembly code.
- ▶ Ready for optimization!

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