#### Plan

Modeling techniques

15 Relaxation linéaire

16 Branch & Bound

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Integer variables are used

- to describe discrete structures subset of  $S \subseteq \{1,\ldots,n\} \Rightarrow \text{vector } (x_1,\ldots,x_n) \in \{0,1\}^n$
- to linearize non linear expressions

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$$\begin{cases} \sum_{i=1}^{k} y_i = 1 \\ x = \sum_{i=1}^{k} p_i y_i \\ y_i \in \{0, 1\} \quad \text{for } i = 1, 2 \dots k \end{cases}$$

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**OR:** x or y must be TRUE x and y being two boolean variables  $\{0,1\}$ 

$$x + y \ge 1$$

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#### **Disjunctive constraints:**

• Two tasks of duration  $d_i$  and  $d_j$  must be processed on the same machine  $\begin{cases} t_i + d_i \leq t_j & \text{if } i \text{ is scheduled before } j \\ t_i + d_i \leq t_i & \text{if } j \text{ is scheduled before } i \end{cases}$ 

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$$\begin{cases} t_i + d_i \leq t_j + M(1 - y_{ij}) \\ t_j + d_j \leq t_i + My_{ij} \\ y_{ij} \in \{0, 1\} \end{cases}$$

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\begin{cases} y \leq x \\ y \leq x' \\ x + x' - 1 \leq y \\ y \in \{0, 1\} \end{cases}
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