Impact of memory allocation on the performance of a delegation synchronization algorithm

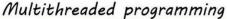
Presented by **SID-LAKHDAR Riyane** (M2 MoSIG: ENSIMAG / UJF)
Supervised by **ROPARS Thomas** (LIG, team ERODS)







MOTIVATION: THE MULTITHREADING LIMITATION





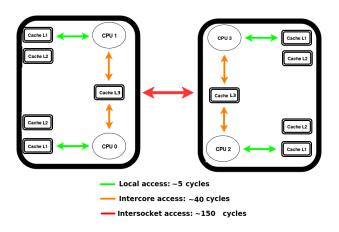


► Theoretical objective of multithreading:

$$Time(multithreaded task) = \frac{Time(Sequential task)}{Number of thread}$$

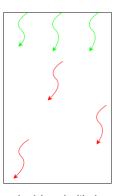
► Objective is still way of the mark

MOTIVATION: COMPLEX PROCESSOR \rightarrow COMPLEX ALLOCATOR

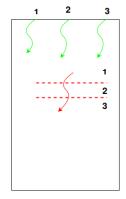


► Memory allocator is an important cause of this limitation

MOTIVATION: DELEGATION APPROACH



Lock based critical section



Delegation based critical section



OUR CONTRIBUTION

Prove the impact of dynamic memory allocation on

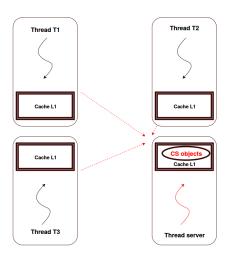
- ► The absolute performances of a custom multithreaded queue
- ► The relative performance of two custom delegation algorithms

THE DELEGATION APPROACH FOR MUTUAL EXCLUSION

MEMORY ALLOCATION: CHALLENGES AND PROPOSED SOLUTIONS

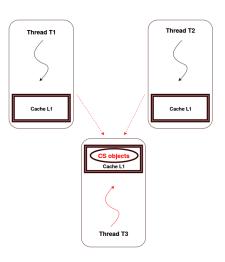
EXPERIMENTAL EVALUATIONS

THE DELEGATION APPROACH

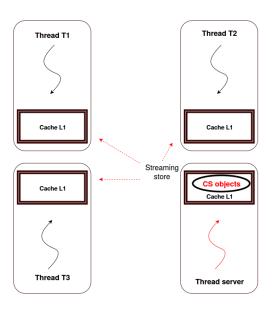


► Improve cache locality: memory accesses of the CS are mostly locals

THE COMBINER ALGORITHM



THE EXPECTED OPTIMIZATION: STREAMING STORE



LIMIT OF THE EVALUATION: THE ALLOCATOR

Used allocator:

- ► Large thread-local buffer used for all the allocations
- ► No returned memory

Unrealistic allocator:

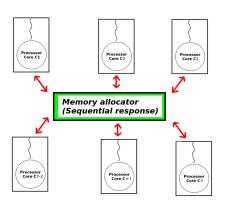
- ► Outperforms the general purpose allocators
- Cannot scale (memory footprint)

THE DELEGATION APPROACH FOR MUTUAL EXCLUSION

MEMORY ALLOCATION: CHALLENGES AND PROPOSED SOLUTIONS

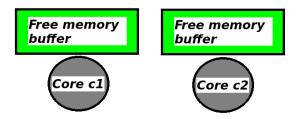
EXPERIMENTAL EVALUATIONS

MEMORY ALLOCATOR: A BOTTLENECK



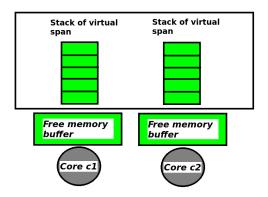
- ► Sequential memory allocation
- ► Inefficient due to contention

CORE LOCAL BUFFER



- ► Improved allocation time
- ▶ Buffer made of contiguous addresses \rightarrow low page-fault probability

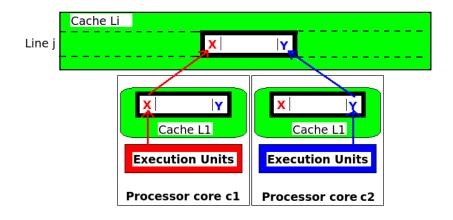
FURTHER CONTENTION LIGHTENING



One stack per core:

- ► No synchronization between cores
- ► Lock free concurrent data structure

SOFTWARE ISSUES: FALSE SHARING



THE SPAN STRUCTURE

- ► **Span size is a multiple of a cache line size:** no false sharing
- ► Simplified allocation algorithm
- Extra memory aligned to page size: swapped out of the RAM

THE DELEGATION APPROACH FOR MUTUAL EXCLUSION

MEMORY ALLOCATION: CHALLENGES AND PROPOSED SOLUTIONS

EXPERIMENTAL EVALUATIONS

DESIGNED TEST ENVIRONMENT

Two main rules: Avoid overhead linked to scheduling

- ► No more threads than cores
- ► Pin threads on independent cores

Environments

- ► A 20 cores **Intel Xeon E5-2620** with a **Debian** kernel 3.16.7
- ► A 24 cores **AMD Opteron 6164** with a **Debian** kernel 3.16.7

TESTBED

Tested algorithms:

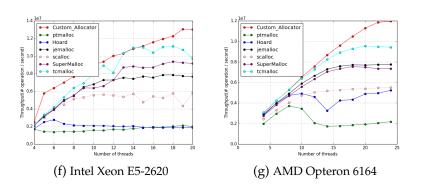
- ► Compare contention on a **Michael-Scott** queue
- ► Compare the **combiner** and the **server based** algorithms.

Considered memory allocators:

- ► **Custom allocator** (used in previous delegation evaluation)
- ► ptmalloc (C standard library)
- ▶ hoard
- ► jemalloc
- ► scalloc
- ▶ supermalloc
- ► tcmalloc

ALLOCATOR COMPARISON (WITH THE COMBINER ALGORITHM)

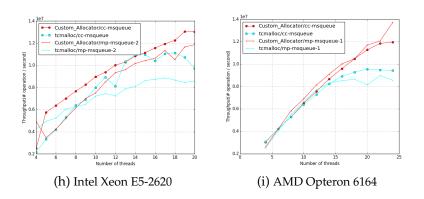
Compare the throughput of the **allocators**: The higher the better



► tcmalloc outperforms the other general purpose allocators

DELEGATION ALGORITHMS COMPARISON

Compare the throughput of the **delegation algorithms**: The higher the better



► With TCmalloc, the relative performance of the algorithms change on AMD

RECAP

During this study, we have:

- ► Highlighted the impact of memory allocation on multithreaded algorithms
- ► Infirmed the performance property of two delegation algorithms

Open issues: understand the reasons of the noticed behavior

- ► The allocation design of tcmalloc
- ► The processor allocation policy

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