## Physical & Link layer

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(Wilson Bentley via wikipedia...)

#### Serial line



- One pin for transmission, one for reception
- Byte by byte transmission (or groups of 5, 6 or 7 bits)
  - √ When idle, the line is in state I
  - √ start bit: 0
  - √ stop bit: I (during I, I,5 or 2 periods)
- Optional : parity bit (error detection)
- ⇒ This is suitable for a console, this is not enough for a data link (PPP or SLIP are mandatory)

The signal is often modulated by a modem for long distance transmission  $\rightarrow$  other pins carry utility signals.

# Serial line data link PPP (Point to Point Protocol) and SLIP (Serial Line IP)

- Framing (framing) and transparency (sole purpose of SLIP)
- error detection / correction
- PDU data description (eg. IP packet or netbios ?)
- Authentication before link establishment
- Configuration/negociation of layers 2 & 3
- Headers or packet compression

## PPP in HDLC framing



- FCS : frame check sequence
- after FCS computation : for byte-stuffing lines:
  - √ control bytes escaping in the data:

$$0x7e \rightarrow 0x7d$$
,  $0x5e$  (Flag);  $0x03 \rightarrow 0x7d$ ,  $0x23$ . (Control);  $0x7d \rightarrow 0x7d$ ,  $0x5d$ . (Escape)

#### For bit-stuffing lines

- √ between the flags
- √ insertion of a zero after any sequence of 5 "1"
- √ Removed upon reception

## **PPP:** associated protocols

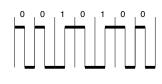
- Négociation LCP (Link Control Protocol)
  - √ Link establishment & control
  - √ Connection parameters negociation (for instance: PPP header compression — eg. address field is useless! FCS suppression)
  - ✓ Transparency: avoiding use of bit sequences significant to the PHY layer
- CCP (Compression...)
  - √ layer 3 headers compression (Similar packets often follow each other)
  - √ Whole packet compression (zip-style)
- Authentification, to choose from :
  - ✓ Unix-style login: authentication before PPP starts (see lab session)
  - ✓ PAP (username and password transmission in LCP packets)
  - √ CHAP (Cryptography)
- Network layer configuration: NCP protocol family
  - ✓ IPCP Protocol (config. IP)
  - √ ATCP (Appletalk config.)

## Physical layer

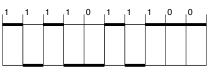
- Transmit symbols
- · Maintaining synchronization

## **Transmission symbols**

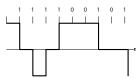
 Manchester coding: one transition per bit



• NRZI : Non-Return to Zero, Invert on one



• MLT-3 : Multi-Level Transitions



## **Maintaining synchronization**

- Manchester: I transition per bit
  - √ Ethernet I0Mb/s
- NRZI
  - √ At most I transition per bit
  - ightarrow Bandwidth usage gain
  - → Clock drift :

A transcoding step introduces transitions

100base-TX : 4B5B 
$$\rightarrow$$
 MLT-3

100base-FX:  $4B5B \rightarrow NRZI$ 

Scrambling

## Why is the preamble useful?

- The sampling point matters!
  - √ With manchester coding, a sequence of Is is very similar to a sequence of 0s
- → A preamble precedes the frame Sequence of 0, 1, 0, 1...

#### 100base-X

- Only uses 2 twisted pairs (-TX) or 2 fibers (-FX)
- · After negociation, the cable is never idle
- 4B5B transcoding creates transitions
   For each group of 4 bits, there is a dictionary lookup
   ex: 0x0 =: 00001; 0xE =: 11100
- Start of frame delimiter (1100010001)[replaces the first preamble transcoded byte]
   End of frame delimiter 0110100111 [overlaps start of IFS]
- MLT-3 (NRZI-3) (MultiLevel Transitions)
  - √ Three logical levels: +1;0;-1

#### What is a hub?

- Repeats from one wire to all others
  - $\rightarrow$  offset attenuation
    - \* Inter-connects fiber to copper

(physical heterogeneity, same speed)

- Detects collisions, generates jam
- The Ethernet hub is a repeater
  - √ The network topology is a star, with possibly sub-stars...

It's a snowflake!

✓ One uplink and one downlink per interface

## CSMA/CD

```
i = 1
while (i <= maxAttempts) do
  listen until (channel is idle)
  transmit and listen at the same time
  wait until (end of transmission) or (collision detected)
  if (collision detected) then
    stop transmitting /* after 32 bits ("jam")*/
    wait random_time ## picked in [0, CW-1]
    increment i
  else
   wait for interframe_gap
   exit
end do
```

## CSMA/<u>CD</u> (cont.)

- Initially, CW=1 (no random wait), doubles after each collision until CW=2<sup>10</sup>
- Max number of transmissions attempts: 16 (that's 16 collisions...)
   then the frame is dropped, start from the beginning with fresh frame (and CW=I!)

### CSMA/CD, the numbers

- I0Mb/s:
  - √ slotTime 512 bit times
  - ✓ interFrameGap 9.6  $\mu$ s
  - √ jamSize 32 bits
  - √ maxUntaggedFrameSize 1518 octets; minFrameSize 512 bits
- 100Mb/s:
  - ✓ slotTime 512 bit times
  - ✓ interFrameGap 0.96  $\mu$ s
  - √ jamSize 32 bits
  - √ maxUntaggedFrameSize 1518 octets; minFrameSize 512 bits
- I000Mb/s: slotTime 4096 bit times
- propagation speed.

```
Thin Coax. and fiber : \simeq 200.10^6~\rm m.s^{-1} ; twisted pair : 177.10^6~\rm m.s^{-1}
```

## **Bridge**

- Collision-free interconnection
- Allows to extend the Ethernet network further than the largest collision domain
- More than one frame may be crossing the switch at a given moment
- · A switch is a multi-interface bridge
  - ✓ Interconnect heterogeneous networks (100Mb/s to 10Mb/s, to wireless...)
  - √ Simple routing (...)
- VLANs

#### **Switched Ethernet**

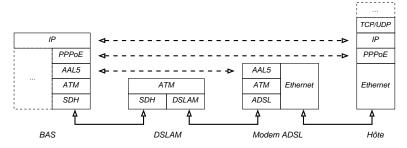
- Only the interface towards the receivers gets the frame
- ⇒ Ethernet «Full Duplex»
  - √ If no collision occurs, it is useless to listen while transmitting (How do we detect that the wire was unplugged?)
    - ⇒ Independent use of uplink and downlink
  - ⇒ Full Duplex Ethernet: No CSMA/CD

#### **ADSL: architecture**

• DSLAM: DSL Access Module

• BAS : Broadband Access Server





Typical architecture...

#### **PPPoE**

 PPPoE : No HDLC framing! Ethernet payload:

				1											2											3					
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
+-+	-+-	-+-	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	-+	+	4	4	4	+	+-+	+	+	+	+	+	+	+-+
1	VE	ER   TYPE						CODE						9					SESSION_ID									- 1			
+-															+-+																
1		LENGTH													- [	payload									~						
+-	+-+-+-+-+-+-+																														

- PADI, PADO, PADR, PADS, PADT: PPPoE Active Discovery Initiation, Offer, Request, Session-confirmation, Terminate
  - ✓ PADI : broadcast, expects a PADO
  - √ PADR: The client picks a server
  - ✓ PADS: The server attributes a session #
  - ✓ PADT : Sent anytime

## PPPoE (cont.)

- Then LCP packets exchanges, like other PPP variants
  - √ The session number lasts for subsequent packets. It identifies the client (useful on shared medium).
  - √ The PPPoE payload is a PPP frame:

 Other architectures: IP/AAL5/ATM (BAS and DSLAM is the same entity)