

Sum of Squares programming

TK18 2020 - Fifth colloquium Morten Hovd

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Aims

- Introducing Sum of Squares (SoS)
 - Introducing SoS decomposition as a means of proving non-negativity of polynomials
 - Illustrating SoS programming as a tool for control analysis and design for polynomial non-linear systems.
- Illustrations based on recent work on discrete-time bilinear systems

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Polynomials and monomials

• A multivariate polynomial f(x) has the form

$$f(x) = \sum_{k} c_k x_1^{a_{k1}} x_2^{a_{k2}} \cdots x_n^{a_{kn}}$$

where a_{ki} are non-negative integers

• A monomial is a term in a polynomial (without the coefficient c_k):

$$m_k(x) = x_1^{a_{k1}} x_2^{a_{k2}} \dots x_n^{a_{kn}}$$

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Sum of Squares decomposition

• A polynomial f(x) is a sum of squares if it can be expressed as a sum of squared terms

$$f(x) = \sum_{i=1}^{N} h_i^2(x) = \sum_{i=1}^{N} (q_i^T v(x))^2 = v^T(x) Q v(x)$$

for some vector of monomials v(x) and PSD matrix Q

• Similarly, a symmetric polynomial matrix M(x) is an SoS matrix if it can be decomposed as

$$M(x) = H^{T}(x)H(x)$$

for some polynomial matrix H(x)

• The primary use of an SoS decomposition is to find a guarantee for the non-negativity (or positivity) of a polynomial expression

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Positivity and SoS



 The existence of a SoS decomposition is only a sufficient condition for non-negativity of a polynomial. For instance, the following polynomial is non-negative, but no SoS decomposition exists:

$$M(x, y, z) = x^4y^2 + x^2y^4 + z^6 - 3x^2y^2z^2$$

- The existence of an SoS decomposition is equivalent to the nonnegativity of a polynomial <u>only</u> in the following cases:
 - Polynomials in two variables
 - Quadratic forms (polynomials of maximum order 2)
 - Fourth order polynomials in three variables
- Trivial conditions for the existence of an SoS decomposition:
 - Any constant term must be non-negative
 - The highest order of the polynomial must be even

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 From definition on previous slide, we have an SoS polynomial if it can be expressed as

$$f(x) = \sum_{i=1}^{N} h_i^2(x) = \sum_{i=1}^{N} (q_i^T v(x))^2 = v^T(x) Q v(x)$$

for some vector of monomials v(x) and some $Q \ge 0$

• However, since the elements of v(x) are not independent, the matrix Q is not unique.

Example



Consider the following example (Parillo, thesis, p. 41):

$$f(x_1, x_2) = 2x_1^4 + 2x_1^3x_2 - x_1^2x_2^2 + 5x_2^4 = \begin{bmatrix} x_1^2 \\ x_2^2 \\ x_1x_2 \end{bmatrix}^T \begin{bmatrix} 2 & 0 & 1 \\ 0 & 5 & 0 \\ 1 & 0 & -1 \end{bmatrix} \begin{bmatrix} x_1^2 \\ x_2^2 \\ x_1x_2 \end{bmatrix}$$

• The matrix Q is trivially not positive (semi-)definite. However, we can use the fact that the elements of v(x) are not independent to introduce a new variable in the matrix Q:

$$f(x_1, x_2) = 2x_1^4 + 2x_1^3 x_2 - x_1^2 x_2^2 + 5x_2^4 = \begin{bmatrix} x_1^2 \\ x_2^2 \\ x_1 x_2 \end{bmatrix}^T \begin{bmatrix} 2 & -\lambda & 1 \\ -\lambda & 5 & 0 \\ 1 & 0 & -1 + 2\lambda \end{bmatrix} \begin{bmatrix} x_1^2 \\ x_2^2 \\ x_1 x_2 \end{bmatrix}$$

- Easy to choose λ such that Q is positive definite
 - E.g., $\lambda = 1$
 - The set of parameters λ which makes $Q \geq 0$ is convex
 - Holds in general for symmetric matrices that depend linearly on one or more parameters

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SoS programming



- SoS constraints are convex
 - Assuming linear dependence on decision variables in control design or analysis formulations
 - Linear in free variables resulting from algebraic dependence of elements of monomial vector v(x)
- Can be combined with a convex optimization criterion to define a convex semi-definite programming (SDP) problem.
- Available software
 - YALMIP
 - SOSTOOLS http://www.cds.caltech.edu/sostools/
 - Both are free for non-commercial use, and are toolboxes for Matlab
 - Personal experience only with Yalmip
 - In addition, an appropriate SDP solver is required

Numerical issues



- The optimal matrix Q will often be severely ill-conditioned, with some very small eigenvalues
- Numerical conditioning may be improved by preprocessing:
 - Removing from the vector v(x) elements (monomials) that will not be required
 - Imposing a block-diagonal structure on Q, depending on symmetries in the polynomial
- Analysis of an initial solution may give hints about additional structure in the optimal Q that can further improve conditioning
 - Re-run after implementing such additional structure
- Good software (e.g. Yalmip) will perform the pre-processing automatically, and allow for automatic post-processing
- For the user, there is therefore little need to know the details of these manipulations
 - Interested readers may consult Löfberg (IEEE TAC, 2009)

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Proof of positivity of a polynomial



- SoS optimization will always yield a positive definite Q (if a 'primal' optimization algorithm is chosen). Does this mean that the original polynomial is an SOS???
- Important to consider the accuracy of the solution
- The *residual* is defined as the largest coefficient in the polynomial $f(x) v^T(x)Qv(x)$
 - Due to floating point arithmetic in solvers, it is rare to find a zero residual
- Define
 - $\lambda_{min}(Q) = \min eig(Q)$
 - M: number of monomials in v(x)
- The polynomial f(x) is guaranteed to be positive provided
 - $-\lambda_{min}(Q) \geq 0$
 - $\lambda_{min}(Q) \ge M \times abs(residual)$
 - See Löfberg (2009)

Proof of positivity



- Q is often ill-conditioned, can be important to check
- A 'primal' solver will always give $Q \ge 0$
 - But still important to check the eigenvalues compared to the size of the residual
- A 'dual'/infeasible path solver may give an indefinite Q
- Occasionally, chosing a dual solver may (for numerical reasons) give a valid solution that is not found using a primal formulation.

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LMIs and SoS programming



- SoS problems are solved as a 'special case' of LMI problems: does that mean that the same tricks can be used for SoS's as for LMI's?
 - Answer: only some of them!
- Linearizing change of variables:
 - Let a problem depend (linearly) on P>0 and PK, but without K appearing independently. May then introduce L=PK, to formulate a problem linear in L and P.
 - Note that
 - K can always be calculated from L and P.
 - A constraint such as L = PK is bilinear in the variables P and K, and is hence NOT convex.

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Congruence transform



• For a matrix valued SoS problem described by $M(x) \ge 0$, and a full rank W(x) (not necessarily symmetric or positive definite)

$$M(x) \ge 0 \Leftrightarrow W^{T}(x)M(x)W(x) \ge 0$$

 However, this congruence transform can NOT be used together with the scalarized Schur complement (to be shown below)

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Scalarization



- For LMIs, $L > 0 \Leftrightarrow x^T Lx > 0 \forall x$
- For SoS, M(x) > 0 and $x^T M(x) x > 0$ are very different
 - M(x) > 0 implies $z^T M(x) z > 0 \forall x$, $\forall z$ (with no relationship between x and z).
 - It is therefore an advantage to formulate scalar-valued SoS problems (i.e., to scalarize the problems), see below for illustrations.

A scalarized Schur complement



Consider:

$$M(x) = \begin{bmatrix} E(x) & F^T(x) \\ F(x) & P(x) \end{bmatrix}$$
 with $P(x)$ symmetric and invertible.

Then:

$$\begin{bmatrix} x \\ z \end{bmatrix}^T M(x) \begin{bmatrix} x \\ z \end{bmatrix} > 0, \forall \{x, z\} \neq \{0, 0\}$$

is equivalent to

$$x^T \Big(E(x) - F^T(x) P^{-1}(x) F(x) \Big) x > 0 \forall x \neq 0$$
and $z^T P(x) z > 0 \forall z \neq 0, \forall x$
(Note that $z^T P(x) z > 0 \forall z \neq 0, \forall x \Leftrightarrow z^T P^{-1}(x) z > 0 \forall z \neq 0, \forall x$)

This follows from

$$M(x) = \begin{bmatrix} I_E & F^T(x)P^{-1}(x) \\ 0 & I_P \end{bmatrix} \begin{bmatrix} E(x) - F^T(x)P^{-1}(x)F(x) & 0 \\ 0 & P(x) \end{bmatrix} \begin{bmatrix} I_E & 0 \\ P^{-1}(x)F(x) & I_P \end{bmatrix}$$

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A scalarized Schur product



Consider

$$\begin{bmatrix} x \\ w \end{bmatrix} = \begin{bmatrix} I_E & 0 \\ P^{-1}(x)F(x) & I_P \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix}$$

- Clearly, w can take any value whatever the value of x, and for any w there is a corresponding value of z whatever the value of x.
- Note that using the congruence transform pre- and postmultiplying the matrix M(x) with non-singular matrices will ruin the equality between the values of x in M(x) and the values in the vectors pre-and postmultiplying M(x).
 - Hence the congruence transform should not be used with the scalarized Scur complement.

The S-procedure



• This applies in much the same way as for LMIs:

We want to prove that f(x)>0 wherever g(x)<0. This may be formulated as

$$f(x) + s(x)g(x) > 0$$

for some SoS polynomial s(x).

Note that

- s(x) is not restricted to be a positive constant, it can be any SoS polynomial
- Optimizing over parameters in both s(x) and g(x) (at the same time) will lead to a bilinear (non-convex) problem.

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SoS programming for bilinear discrete-time systems



Bilinear discrete-time systems



• Systems of the form

$$x_{k+1} = Ax_k + \sum_{i=1}^{m} (B_i x_k + b_i) u_{i,k} = Ax_k + (B_x + B) u_k$$

 Assuming (possibly non-linear) state feedback, stability is ensured by fulfilling the Lyapunov difference inequality

$$x_{k}^{T} P x_{k} - x_{k+1}^{T} P x_{k+1} > 0$$

$$x_{k}^{T} P x_{k} - \left(A x_{k} + (B_{x} + B) u_{k}(x_{k})\right)^{T} P \left(A x_{k} + (B_{x} + B) u_{k}(x_{k})\right) > 0$$

- If system is unstable with $u_k=0$ (for a state affected by bilinearity):
 - No $u_k(x_k)$ polynomial in the state can give global quadratic stability
 - $u_k(x_k)$ must be *a ratio* of two polynomials of the same order to achieve global quadratic stability
 - Note that continuous-time bilinear systems may allow polynomial u(x) and still achieve global quadratic stability

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Proposed controller

• The proposed controller takes the form

$$u_k(\mathbf{x}_k) = \frac{C(x_k)x_k}{c_0(x_k) + 1}$$

Where $\mathcal{C}(x_k)$ is a polynomial matrix and $c_0(x_k)$ is an SoS polynomial

 Adding 1 to the numerator prevents the numerator from becoming very small, which could lead to an 'explosion' in the input



Region of convergence

• Given a quadratic Lyapunov function $V = x_k^T P x_k$, a polynomial matrix $C(x_k)$ and SoS polynomials $c_0(x_k)$ and $s_1(x_k, z)$, the closed loop systems is stable $\forall x_k \mid x_k^T P x_k < \gamma$ provided

$$\begin{bmatrix} x_k \\ z \end{bmatrix}^T M(x) \begin{bmatrix} x_k \\ z \end{bmatrix} - s_1(x_k, z)(\gamma - x_k^T P x_k) > 0$$

where

$$M(x_k) = \begin{bmatrix} (c_0(x_k) + 1)P & ((c_0(x_k) + 1)A + (B_x + B)C(x_k))^T P \\ P((c_0(x_k) + 1)A + (B_x + B)C(x_k)) & (c_0(x_k) + 1)P \end{bmatrix}$$

• Follows from $(c_0(x_k)+1)$ being strictly positive, the scalarized Schur complement and the S-procedure

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Input saturation

- Consider input constraints of the form $-u_{i,max} \le u_i \le u_{i,max}$
- Let $c_i(x_k)$ denote an element of the polynomial vector $(C(x_k)x_k)$, $c_o(x_k)$ be an SoS polynomial (as before), and $q_i(x_k)$ be an SoS polynomial. Then the input constraint is satisfied provided $\forall x_k \mid x_k^T P x_k < \gamma$

$$\begin{bmatrix} (c_0(x_k) + 1) u_{i,\max}^2 - q_i(x_k) (\gamma - x_k^T P x_k) & c_i(x_k) \\ c_i(x_k) & (c_0(x_k) + 1) \end{bmatrix} > 0$$

• Follows from Schur complement and S-procedure, and $(c_0(x_k)+1)$ being strictly positive

Rate of convergence



• An exponential rate of convergence specified by the scalar α is easily included by changing the matrix $M(x_k)$ above:

$$M(x_k) = \begin{bmatrix} (1-\alpha)(c_0(x_k)+1)P & \left((c_0(x_k)+1)A+(B_x+B)C(x_k)\right)^T P \\ P\left((c_0(x_k)+1)A+(B_x+B)C(x_k)\right) & \left(c_0(x_k)+1)P \end{bmatrix}$$

 Maximizing the region of convergence can easily lead to slow control, specifying a certain rate of convergence can avoid this (at the expense of having a smaller region with this specified rate of convergence)

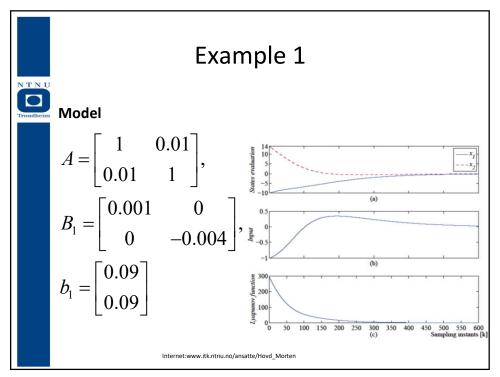
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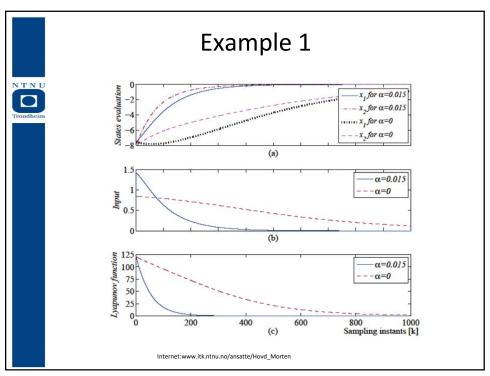
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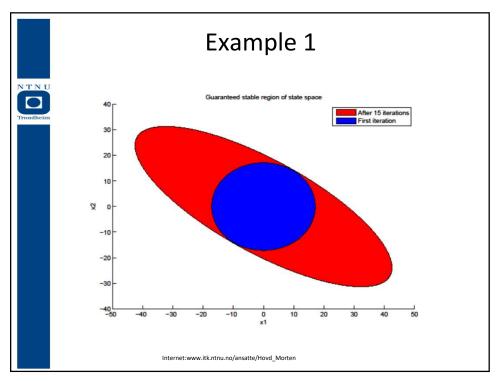
Optimization formulation

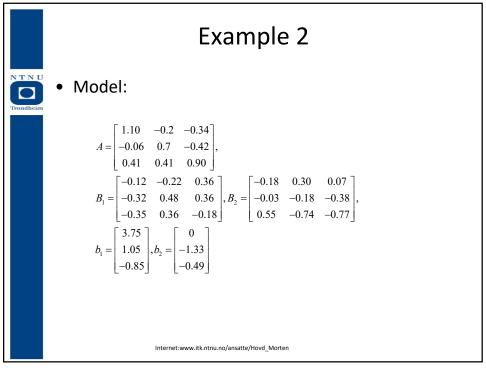


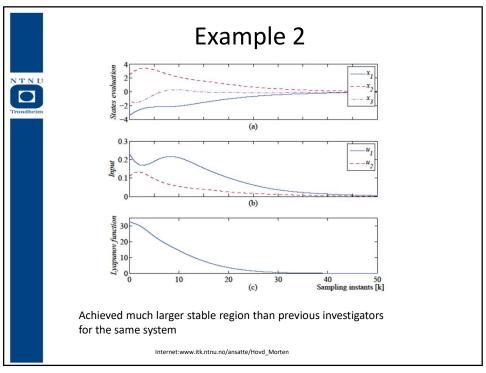
- If α , γ , and P are specified, the coefficients of $C(x_k)$, $c_0(x_k)$, $q_i(x_k)$, and $s_1(x_k,z)$ all enter linearly in the inequalities above (and vice versa).
- For given γ , and P we can find $C(x_k)$, $c_0(x_k)$, $q_i(x_k)$, and $s_1(x_k,z)$
 - Formulated as a feasibility problem, i.e., an optimization with an empty objective in Yalmip
- For given $C(x_k)$, $c_0(x_k)$, $q_i(x_k)$, and $s_1(x_k, z)$, we can maximize γ with P as a free variable.
 - To avoid γ and P growing without bound (without describing a larger region), we will need to add a normalizing constraint such as trace(P) = constant.
- Iterating between the two optimization problems, we gan design controllers stabilizing gradually increasing regions.
- Will typically initialize with LQ solution of linearized system















x_{k+1} in an augmented 'state space'

- NTNU
- Previously: use model equations to express x_{k+1} using x_k and u_k .
- May instead express the problem using both x_k and x_{k+1} .
- Given a polynomial LF V(x) and model $x_{k+1} = f(x_k, u_k)$ where $f(x_k, u_k)$ is a polynomial vector:

$$V(x_k) - V(x_{k+1}) + m^T(x_k, x_{k+1}) (x_{k+1} - f(x_k, u_k)) - s(x_k, x_{k+1}) (\gamma - V(x_k)) > 0$$

- Here: $s(x_k, x_{k+1})$ is SoS, $m^T(x_k, x_{k+1})$ free (not necessarily SoS)
- If expression holds globally (in $[x_k^T \ x_{k+1}^T]^T$), it also has to hold along the trajectories of the system, where $(x_{k+1} f(x_k, u_k))$ =0.

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x_{k+1} in an augmented 'state space'

- Allows for simultaneously optimizing the parameters of LF and controller.
 - May be an advantage in some cases
 - The vector $m(x_k, x_{k+1})$ may introduce many new parameters if there are many states
 - Higher order LFs more easily expressed than with previous formulation
- Straight forward if u_k enters linearly in $f(x_k, u_k)$.
 - For controller parametrization as used previously, or purely polynomial controllers
- Need additional SoS constraints to ensure
 - _ IF>0
 - Stable region increasing from one iteration to the next
 - Details omitted for brevity
- Both formulations can relatively easily be adapted to continuoustime systems
 - Using the gradient of the LF instead of the LF difference



u_k non-linear in $f(x_k, u_k)$

- Augmented state space including also u_k
- Expand $f(x_k, u_k)$ to include controller, e.g., for $u_k = c(x_k)$:

$$V(x_k) - V(x_{k+1}) + m^T(x_k, x_{k+1}, u_k) \left(\begin{bmatrix} x_{k+1} \\ u_k \end{bmatrix} - \begin{bmatrix} f(x_k, u_k) \\ c(x_k) \end{bmatrix} \right) - s(x_k, x_{k+1}, u_k) (\gamma - V(x_k)) > 0$$

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APPROXIMATING OPTIMIZATION-BASED CONTROLLERS



Optimization problem P_{...}

- Original problem
- $\min_{u_k, x_{k+1}} \frac{1}{2} x_{k+1}^T Q x_{k+1} + u_k^T R u_k \quad \bullet \quad \min_{u_k} \frac{1}{2} u_k^T H u_k + x_k^T F u_k$ $x_{k+1} = A x_k + B u_k \quad \bullet \quad H_u u_k \le 1$
- $H_u u_k \leq 1$
- $V(x_{k+1}) \le \gamma V(x_k)$
- 'One-step' MPC
 - Can be generalized to more common MPC formulations
- Imposing decrease of Lyapunov function directly $(\gamma < 1)$, a.k.a. a contraction constraint
- Instead of using the value

• Equivalent problem

- $V(x_{k+1}) \le \gamma V(x_k)$

Convex problem if $V(x_{k+1})$ is convex

Unconstrained solution can be found. LMI techniques can be used to find a corresponding ellipsoid inside which this is valid. Subsequent developments are applied in the constrained region, inside the set where $V(x_k) < 1$

function as an LF Internet:www.itk.ntnu.no/ansatte/Hovd_Morten

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KKT conditions of problem P₁₁

- $Hu_k + F^T x_k + H_u^T \lambda_u + V_{k+1}^1(u_k) \lambda_q = 0$
- $H_u u_k \leq 1$
- $V(x_{k+1}) \gamma V(x_k) \le 0$
- $\lambda_u \ge 0$
- $\lambda_u^T(H_u u_k 1) + \lambda_q(V(x_{k+1}) \gamma V(x_k)) = 0$



Optimization problem P_c

- $\min_{c,u_k,\lambda_u,\lambda_q} c$
- $Hu_k + F^T x_k + H_u^T \lambda_u + V_{k+1}^1(u_k) \lambda_q = 0$
- $H_u u_k \leq 1$
- $V(x_{k+1}) \gamma V(x_k) \le 0$
- $\lambda_u \ge 0$
- $\lambda_q \geq 0$
- $\lambda_u^T(H_u u_k 1) + \lambda_q(V(x_{k+1}) \gamma V(x_k)) \le c$
- A feasible solution to P_c is also a feasible solution to P_u
- If, at optimum, c=0, the optimal solution to ${\rm P_c}$ is also an optimal solution to ${\rm P_u}$

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Approximate solution to Pu

- Let $J(u_k^*(c))$ be the value of the objective function of P_u evaluated using the input u_k^* from the solution to problem P_c .
- Let $J(u_k^*(0))$ be the (optimal) value of the objective function of ${\sf P_c}$.
- Then $J(u_k^*(c)) J(u_k^*(0)) < c$
 - The degree of suboptimality of P_c (relative to P_u) is bounded by the degree of violation of the complementarity constraints
 - Also holds for other convex quadratic objective functions with convex constraints



Optimization formulation

- Basic idea: use polynomial expressions for u_k , λ_e , λ_q , V
- However, when optimizing also over V, must include additional constraint $V_{xx}>0$
 - Where V_{xx} is the matrix of second derivatives, i.e., $[V_{xx}]_{i,j}=rac{\partial^2 V}{\partial i\partial j}$
 - Otherwise the initial formulation is not convex, and we cannot replace the optimization problem by the KKT conditions
- However, using the model equation to eliminate x_{k+1} will lead to a non-linear dependence of V(k+1) on the controller parameters
 - And hence a severely non-convex problem
- Instead, the problem is lifted to a higher-dimensional space described by (x_k, x_{k+1})
 - Starting from the original problem formulation on slide 35

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Optimization formulation

- $\min_{c,u_k,\lambda_u,\lambda_q,\lambda_e,\mu_i,V} c$
- $H_u u_k \le 1$, $\lambda_u \ge 0$, $\lambda_q \ge 0$
- $Ru_k B^T \lambda_e + H_u^T \lambda_u + \mu_1^T (x_{k+1} Ax_k Bu_k) = 0$
- $Qx_{k+1} + \lambda_e + \lambda_q \nabla V(x_{k+1}) + \mu_2^T (x_{k+1} Ax_k Bu_k) = 0$
- $V(x_{k+1}) \gamma V(x_k) + \mu_3^T(x_{k+1} Ax_k Bu_k) \le 0$
- $\lambda_u^T(H_u u_k 1) + \lambda_q(V(x_{k+1}) \gamma V(x_k)) + \mu_4^T(x_{k+1} Ax_k Bu_k) \le c$
- $V_{xx} > 0$
- V SOS, no sign constraint for μ 's
- λ_u, λ_q only need to be positive in the constrained region, inside the set $V(x_k) < 1$. The same holds for all other inequality constraints. Implemented through Sprocedure.
- c scalar, u_k function of x_k (only), V function of x_k or x_{k+1} as appropriate (depending on which time instant it refers to)
- Other variables functions of x_k , x_{k+1}
- Optimizing over polynomial coefficients! These enter linearly or bi-linearly above.
- Where $(x_{k+1}-Ax_k-Bu_k)=0$ and c=0 the original KKT conditions are fulfilled (complementarity constraint relaxed)

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Optimization formulation



- · Additional features to
 - Account for state constraints by ensuring the level set of the LF is inside the state constraint
 - Systematically 'pushing out' the level set to make the guaranteed stable region as large as possible
 - Otherwise, a similar iterative solution procedure as described before
- Note the additional constraint specifying that the Hessian (second derivative matrix) of the LF is an SOS matrix
 - To guarantee the convexity of the contraction constraint.
 - The second derivative matrix of a polynomial is found by a standard function in Yalmip

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Non-polynomial systems

- SoS design assumes polynomial or rational polynomial dynamics and LFs
- Some non-polynomial systems can be handled with the appropriate reformulation (see Papachhristodoulou and Pranja)

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Furuta pendulum

- Pendulum on a rotary base
 - Rotate base to stabilize pendulum in the upright vertical position

$$\left(m_p L_r^2 + \frac{1}{4} m_p L_p^2 - \frac{1}{4} m_p L_p^2 \cos(\alpha)^2 + J_r\right) \ddot{\theta} - \left(\frac{1}{2} m_p L_p L_r \cos(\alpha)\right) \ddot{\alpha} \qquad (1)$$

$$+ \left(\frac{1}{2} m_p L_p^2 \sin \alpha \cos \alpha\right) \dot{\theta} \dot{\alpha} + \left(\frac{1}{2} m_p L_p L_r \sin(\alpha)\right) \dot{\alpha}^2 = \tau - B_r \dot{\theta}$$

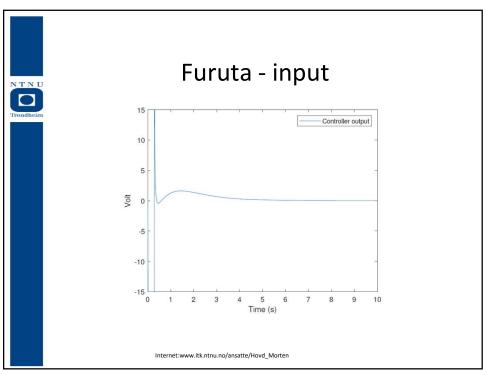
$$-\frac{1}{2} m_p L_p L_r \cos(\alpha) \ddot{\theta} + \left(J_p + \frac{1}{4} m_p L_p^2\right) \ddot{\alpha} - \frac{1}{4} m_p L_p^2 \cos(\alpha) \sin(\alpha) \dot{\theta}^2 \qquad (2)$$

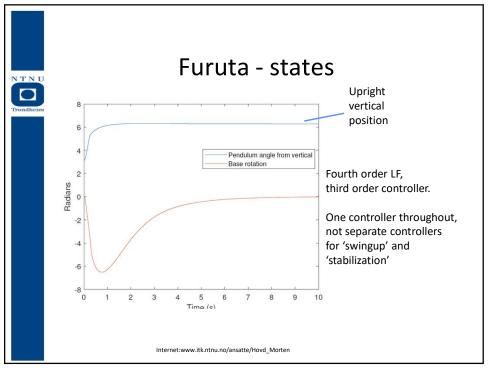
$$-\frac{1}{2} m_p L_p g \sin(\alpha) = -B_p \dot{\alpha}$$
where

where

$$\tau = \frac{\eta_g K_g \eta_m k_t (V_m - K_g k_m \dot{\theta})}{R_m}$$

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Some literature



- P. Parillo, PhD thesis (mainly Chapter 4)
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