

# Model checking with edge-valued decision diagrams

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**1 Decision Diagrams**

**2 EVMDDs**

**3 Implementation**

# The State of Symbolic Model Checking

## Evolution and Impact of Decision Diagrams

- Early 90s : the wow factor, BDDs are (re)discovered
- Late 90s - early 2000s : real progress
  - Extensions, generalizations (MTBDDs, BMDs, EVMDDs, etc)
  - New algorithms (saturation, bounded MC, CEGAR)
- Since then ...
  - Interest has shifted to other areas of verification
  - There are even rumors out there that symbolic MC has entered a Brezhnev era ( $\sim$  stagnation)
  - Fact or fiction ?

# Purpose of this work

## Stagnation: fact or fiction?

- A little bit of both
- New ideas exist, but are disparate
- Example of untapped resources:
  - Edge-valued decision diagrams (EVMDD)
  - Identity-reduced decision diagrams
  - Hashing, caching, garbage collection
  - Heuristics for SAT/SMT solving

## Our goal

Represent in one formalism (some of) the best techniques available at the moment across a spectrum of existing tools

# Encoding of functions

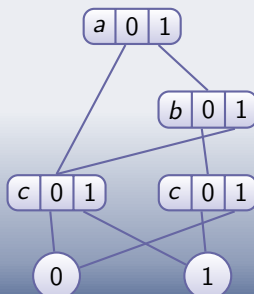
The advent of symbolic MC: **compact** representation of

- boolean functions  $f : \{0,1\}^n \rightarrow \{0,1\}$
- sets  $\{x \in \{0,1\}^n \mid f(x) = 1\}$

Evolution:

- Truth table:  $2^n$  **entries**
- Binary Decision Diagram (BDD): merge common subtrees  
**still exponential size in worst case, often better in practice**

<i>a</i>	<i>b</i>	<i>c</i>	$f(a, b, c)$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1



# Integer/arithmetic functions

- $f : \{0, 1\}^n \rightarrow \mathbb{Z}$
- Extend BDD to **Multi-Terminal BDD (MTBDD)**

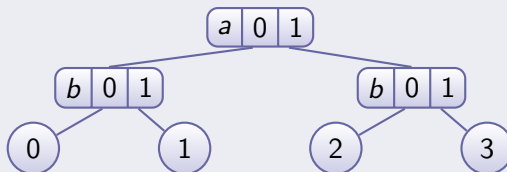


Figure:  $f : (a, b) \mapsto 2a + b$

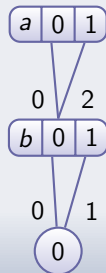
- Inefficient if  $\text{Img}(f)$  is large: less chances to share subtrees

Other forms of DDs:

- Multiway DDs (MDD):  $f : \{0, \dots, k_1\} \times \dots \times \{0, \dots, k_n\} \rightarrow \{0, 1\}$
- Binary Moment Diagrams (BMD):
  - work well for multipliers, but not much else

# Edge Valued MDDs (EVMDDs)

- EVBDDs introduced in 1992, but not sufficiently exploited  
⇒ (*Reed-Müller spectrum* !?!)
- From MTBDDs to EVMDDs:  
merge all terminals (0) and assign (integer) values to edges



- Value of  $f$ : composition of edge-values (e.g. addition,  $+$ ) along the path from root to terminal node

# EVMDD characteristics

- EVMDD encoding cannot have more nodes than MTBDDs  
     $\Rightarrow$  proved in this paper
- Size can be linear instead of exponential (e.g. linear functions)
- Composition  $\Rightarrow$  a generic algorithm for all binary operators:  
    for  $f, g$  encoded by EVMDDs of size  $|f|$  and  $|g|$   
     $f * g$  computed in  $O(|f| |g| |\text{Img}(f)| |\text{Img}(g)|)$
- This algorithm has **exactly the same complexity**  
    as its equivalent for MTBDDs, hence  
    **no gain** in (worst-case) time complexity
- Is there room for improvement ?



# EVMDD algorithms

Yes, for following operations:

- Addition:  
 $f + g$  computed in  $O(|f| \cdot |g|)$
- Multiplication by constant:  
 $f \times c$  computed in  $O(|f|)$
- Multiplication:  
 $f \times g$  computed in  $O(|f|^2 \cdot |g|^2 \cdot |f \times g|)$ 
  - exponential in worst case
  - much better in many “practical” cases
- Remainder and Euclidean division by constant:  
 $f/c$  and  $f \% c$  computed in  $O(c \cdot |f|)$

# An EVMDD-based Model Checker

We have developed an EVMDD library featuring:

- EVMDDs for arithmetic expressions
- (Regular) MDDs for boolean expressions
- Identity-reduced encoding of transition relations
- Saturation-based state space construction
- Unsophisticated (i.e. fast) garbage collector (mark & sweep)

Some stats:

- 7 kLOC of ANSI C : library
- 4 kLOC : model checking front-end

Available at <http://research.nianet.org/~radu/evmdd/>

# Results

Building state space vs CUDD (BFS) and SMART (saturation)

Model	Model size	Reachable states	CUDD (sec)	SMART (sec)	EVMDD (sec)
Dining philosophers	100	$4 \times 10^{62}$	11.42	1.49	0.03
	200	$2 \times 10^{125}$	3054.69	3.03	0.07
	15000	$2 \times 10^{9404}$	—	—	195.29
Round robin mutual exclusion protocol	40	$9 \times 10^{13}$	4.44	0.44	0.08
	100	$2 \times 10^{32}$	—	2.84	1.17
	200	$7 \times 10^{62}$	—	20.02	9.14
Slotted ring protocol	10	$8 \times 10^9$	1.16	0.19	0.01
	20	$2 \times 10^{20}$	—	0.71	0.04
	200	$8 \times 10^{211}$	—	412.27	25.97

On Intel Core 2, 1.2GHz, 1.5GB mem (“—” means “> 1h”).

# Results

Building state space vs CUDD (BFS) and SMART (saturation)

Model	Model size	Reachable states	CUDD (sec)	SMART (sec)	EVMDD (sec)
Kanban assembly line	15	$4 \times 10^{10}$	80.43	3.41	0.01
	20	$8 \times 10^{11}$	2071.58	8.23	0.02
	400	$6 \times 10^{25}$	—	—	74.89
Knights problem	5	$6 \times 10^7$	1024.42	5.29	0.27
	7	$1 \times 10^{15}$	—	167.41	3.46
	9	$8 \times 10^{24}$	—	—	32.20
Randomized leader election protocol	6	$2 \times 10^6$	4.22	8.42	0.86
	9	$5 \times 10^9$	—	954.81	18.89
	11	$9 \times 10^{11}$	—	—	109.25

On Intel Core 2, 1.2GHz, 1.5GB mem (“—” means “> 1h”).

# Questions

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