

Exercise 1

$$n = 7$$

$$min = 2 \quad a_{min} = 3$$

$$j = 2 \quad a_2 = 3 \quad j =$$

① comp

$$j = 3 \quad a_3 = 8$$

① comp

$$j = 4 \quad a_4 = 6$$

① comp

$$j = 5 \quad a_5 = 5$$

① comp

$$j = 6 \quad a_6 = 1$$

Exercice 1

même nb de comparaisons.

$$n + (n-1) + \dots + 1$$

①

9	38	65	17			
1	38	65	97			
1	3	5	68	97		
1	3	5	6	7	98	
1	3	5	6	7	8	9

④

②

4	7	2	9	5	6	3
2	7	4	9	5	6	3
2	3	4	9	5	6	7
2	3	4	5	9	6	7
2	3	4	5	6	9	7
2	3	4	5	6	7	9

⑤

③ 0

④

7	6	5	4	3	2	1
1	6	5	4	3	2	7
1	2	5	4	3	6	7
1	2	3	4	5	6	7

③

2 worst swap
eq. for comp.

✓

$$c^{n-3}$$

$$, c > 1$$

$$3^n$$

comme $c > 3^3 \geq 27$

on a ① > ②

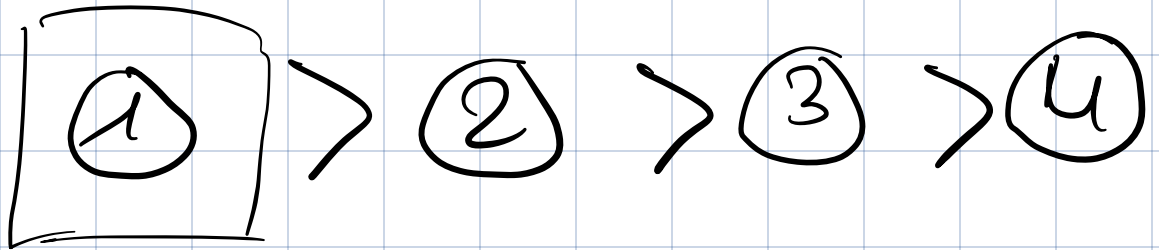
$$n^{3 \log_3(n)} \geq n^3 \log_3(n)$$

la puissance \nearrow de multipl:

$$n^{3x} \geq n^3 x,$$

or $x^n \rightarrow \oplus$ like n^x

done



Ex 3

if $f_1(x)$ is $O(g_1(x))$

$$|f_1(x)| \leq C_1 |g_1(x)|$$

$$|f_2(x)| \leq C_2 |g_2(x)|$$

$$|f_1(x)| + |f_2(x)| \leq C_1 |g_1(x)| + C_2 |g_2(x)|$$

$$\Leftrightarrow |f_1(x) + f_2(x)| \leq$$
$$|f_1(x)| + |f_2(x)|$$

$$\leq C_1 |g_1(x) + g_2(x)|$$

$$\leq C_1 |g_1(x)| + C_2 |g_2(x)|$$

$$\leq (C_1 + C_2) \left(2 |\max(g_1(x), g_2(x))| \right)$$

Ex 4

① f is $o(g(x))$

$$\text{if } \lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$$

$$\lim_{x \rightarrow \infty} \frac{5x}{x^2} = \frac{x(5)}{x(x)} = 0$$

$$\textcircled{2} \lim_{x \rightarrow \infty} \frac{2x^2}{x^2} = 2 \neq 0$$

$$\textcircled{3} \lim_{x \rightarrow \infty} \frac{1}{x} - \frac{1}{x} = 0.$$

$E \times S$

$$f(0) = 0$$

$$f(n) = f(n-1) - 1$$

$$f(n-1) = f(n) - 1$$

0, -1, -2, -3

$$f(n) = -n.$$

Base's step

$$f(0) = 0 \quad \boxed{\text{ok}}$$

Inductive step

$$f(k) = -k$$

$$f(k+1) = -k-1$$

$$= f(k) - 1 \quad \text{Ok.}$$

\Rightarrow therefore $f(n)$ is true $\forall n \in \mathbb{N}$

$$f(0) = 0$$

$$f(1) = 1$$

$$f(n) \stackrel{?}{=} 2f(n-2)$$

$$0, 1, 0, 2, 0, 4, 0$$

$$f(n) = \frac{2^{n-2} \cdot (2^{n-2}) \cdot (-1)^{n+1}}{2}$$

0, 1, 0, 2, 0, 4

0 1 2 3 4 5

2^0

2^1

2^2

$$2^{\frac{n-1}{2}} - \left(2^{\frac{n-1}{2}}\right) \cdot (-1)^n$$

2

$$2^{\frac{n-1}{2}} \left(1 - (-1)^n\right)$$

2

=

$$2^{\frac{n-3}{2}} \left(1 + (-1)^{n-1}\right)$$

Base Step

$$f(0) = \frac{2^{1/2} (1-1)}{2} = 0$$

Inductive Step

We assume $P(k)$ is true.

- $f(n+1) = 2 f(n)$.

We assume $f(n) = \frac{2^{\frac{n-1}{2}} (1 - (-1)^n)}{2}$

- $f(n+1) = \frac{2^{\frac{n}{2}} (1 - (-1)^{n+1})}{2}$

$$= 2^{\frac{n-2}{2}} (1 + (-1)^n)$$

$$f(n) = 2^{\frac{n-3}{2}} \left(1 + (-1)^{n-1} \right)$$

Ex 6

g_1 is $O(f)$
 \Leftrightarrow

③

$$|g_1(x)| \leq C_1 |f(x)|$$

$$|g_2(x)| \leq C_2 |f(x)|$$

f is $\Omega(g_1)$

$$|g_1(x)| \geq C_3 |f(x)| \quad \Leftrightarrow$$

$$|g_2(x)| \geq C_4 |f(x)|$$

$$|g_1(x) + g_2(x)|$$

$$\leq |g_1(x)| + |g_2(x)|$$

$$\leq (C_1 + C_2) |f(x)| \quad \Leftrightarrow (g_1 + g_2)(x) \text{ is } O(f)$$

I
 N
 V
 A
 L
 I
 D

$$\begin{aligned}
 & |g_1(x)| + |g_2(x)| \geq \\
 & \quad (C_3 + C_4) |f(x)| \\
 & |g_1(x)| + |g_2(x)| \\
 & \geq |g_1(x) + g_2(x)| \\
 & \geq (C_3 + C_4) |f(x)|
 \end{aligned}$$

$$\begin{aligned}
 g_1(x) &= -5x \\
 g_2(x) &= 5x
 \end{aligned}$$

$$\begin{aligned}
 g_1 &\text{ is } \Theta(x) \\
 g_2 &\text{ is } \Theta(x)
 \end{aligned}$$

$$\begin{aligned}
 (g_1 + g_2) &\text{ is } \Theta(1) \\
 &\neq \Theta(x)
 \end{aligned}$$

because it is
 not $\Omega(x)$

② + ④

g_1 is $O(f)$

$$|g_1(x)| \leq C_1 |f(x)| \quad \Leftrightarrow$$

$$|g_2(x)| \leq C_2 |f(x)|$$

f is $\Omega(g_1)$

$$|g_1(x)| \geq C_3 |f(x)| \quad \Leftrightarrow$$

$$|g_2(x)| \geq C_4 |f(x)|$$

$$|g_1 g_2| \leq C_1 C_2 |f^2|$$

ok parce
qu'on
multiplie
par la valeur
absolue

$$|g_1 g_2| \geq C_3 C_4 |f^2|$$

①

g_1 is $O(f)$

\Leftrightarrow

$$|g_1(x)| \leq C_1 |f(x)|$$

$$|g_2(x)| \leq C_2 |f(x)|$$

f is $\Omega(g_1)$

\Leftrightarrow

$$|g_1(x)| \geq C_3 |f(x)|$$

$$|g_2(x)| \geq C_4 |f(x)|$$

$$|g_1(x) + g_2(x)|$$

$$\leq |g_1(x)| + |g_2(x)|$$

$$\leq (C_1 + C_2) |f(x)| \quad \Leftrightarrow (g_1 + g_2)(x) \text{ is } O(f)$$

$$|g_1(x)| + |g_2(x)| \geq$$

$$(C_3 + C_4) |f(x)|$$

$$|g_1(x)| + |g_2(x)|$$

$$= |g_1(x) + g_2(x)|$$

$$\geq (C_3 + C_4) |f(x)|$$

supérieur à g_1

donc aussi
supérieur à

g_1

car quel que

> 0

⑤

$$|g(x)| \leq C |f(x)|$$
$$C > 0$$
$$\forall x > k.$$

$$\Leftrightarrow 2^{|g(x)|} \leq 2^{C|f(x)|}$$

arg g ev $f > 0$

$$\Leftrightarrow 2^g \leq (2^C)^f$$
$$2 \leq C_2 2^f$$

2 1

$$2x \leq 4 \times 1x$$

$$2^{2x} < 2^{1x}$$

$$4^x < 2^x$$

Exercice 7

• $\log_m(k) = x \Leftrightarrow m^x = k$

• $\log_m(l) = y \Leftrightarrow m^y = l$

or $k > l \Rightarrow m^x > m^y$.

$\Rightarrow x > y$.

donc $\log_m(k)$ is $O(\log_m(l))$

$$\log_m(k) = \frac{\log_2(k)}{\log_2(m)} = aX$$

$$\log_m(l) = \frac{\log_2(l)}{\log_2(m)} = bX$$

\Rightarrow même ordre, **VRAT**

$$\log_a(b) = \frac{\log_c(b)}{\log_c(a)}$$

$$\log_a(b) = x$$

$$\Leftrightarrow b = a^x$$

$$\Leftrightarrow \log_c(b) = \log_c(a^x)$$

$$\Leftrightarrow \log_c(b) = x \log_c(a)$$

$$\Leftrightarrow x = \frac{\log_c(b)}{\log_c(a)}$$

$$\log_k(m) = \frac{\log_e(m)}{\log_e(k)}$$

$$k = e^{\log_e(k)}$$

$$\frac{\log_k(m)}{\log_k(e)}$$

$$\left(e^{\log_e(k)} \right)^{\frac{\log_k(m)}{\log_k(e)}} = e^{\log_e(k) \cdot \log_e(m)}$$

$$\left(e^{\log_k(m)} \right) = e^{\log_e(m) / \log_e(k)}$$

As $k > e \geq 2$,

$$k^{\log_e(m)} = m^c$$

with $c = \log_e(k) > 1$

$$e^{\log_k(m)} = m^{1/c}$$

$$\Rightarrow \text{ndr } O(e^{\log_k(m)})$$

$$c < 1.$$

Exercice 8

f is $o(f)$

$$\Leftrightarrow \lim_{x \rightarrow \infty} \frac{f(x)}{f(x)} = 0$$

faux.

f is $o(g)$

$$\Leftrightarrow \lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$$

$$\exists x_0, \forall x \geq x_0, g(x) \geq f(x).$$

$$\Rightarrow f \text{ is } o(g).$$

Exercise 9

$$\sum_{i=0}^n 1+2+3+4+\dots+n \quad O(n^2)$$

$$\frac{n(n+1)}{2}$$

$$\leq n+n+\dots+n$$

$$= n \cdot n$$

$$= n^2$$

AK

$$= a_0 n^0 + a_1 n^1 + \dots + a_n n^d$$

$$\leq D n^d + D n^d + \dots + D n^d$$

$$= D \cdot n^d \leq \underline{D \cdot n^{d+1}}$$

$$\geq n^0 + n^1 + \dots + n^d \quad \text{Vrai?}$$
$$= \frac{1 - n^{d+1}}{1 - n^d} =$$

$$\left[\begin{aligned} & a_0 n^0 + a_1 n^1 + \dots + a_n n^d \\ & \ll D n^d + D n^d + \dots + D n^d \\ & = D \cdot n^d \cdot n = D n^{d+1} \end{aligned} \right.$$

$$O(n^{d+1})$$

$$a_0 n^0 + a_1 n^1 + \dots + a_n n^d$$

$$\begin{aligned} & \gg n^0 + n^1 + \dots + n^d \\ & = \sum_{k=0}^d n^k = \frac{1 - n^{(d+1)}}{1 - n} \end{aligned}$$

$$\gg n^{d+1}$$

$$\sum a_0 n^0 + a_1 n^1 + \dots + a_r n^d$$

$$\begin{aligned} a_1 n^1 + a_2 n^2 + a_3 n^3 + \dots + a_d n^d \\ \leq a_d n^d + a_d n^d + \dots + a_d n^d \\ = a_d n^d \cdot d \end{aligned}$$

Vrai

Exercise 10

Base Step

$$4! = 24 > 2^4 = 16$$

Inductive Step

We assume that $P(k): k! > 2^k$ is T.
Does it imply $P(k+1)$ T?

$$P(k+1): (k+1)! > 2^{k+1}$$

$$\Rightarrow k! (k+1) > 2^k \cdot 2.$$

$$\Rightarrow k! > 2^k \cdot \frac{2}{k+1}$$

2
$k+1$

$\nwarrow < 1$

Exercise 11

$$AB = BA. \quad P(k): AB^k = B^k A$$

Base Step

$$AB^1 = B^1 A \quad \boxed{\text{OK}}$$

Inductive step

$$P(k+1): AB^{k+1} = B^{k+1} A.$$

$$\Leftrightarrow AB^k B = B^k B A.$$

$$\Leftrightarrow B^k A B = B^k B A.$$

$$\Leftrightarrow B^k B A = B^k B A. \text{ OK}$$

$$\frac{f(x+h) - f(x)}{h}$$

$$-(-1) = \underline{\underline{1}}$$