

Homework 8

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Exercise 1. The Selection Sort algorithm sorts the input sequence a_1, \dots, a_n of real numbers in ascending order

Algorithm 1 Selection Sort

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1: for  $i = 1$  to  $n - 1$  do
2:    $\min \leftarrow i + 1$ 
3:   for  $j = i + 1$  to  $n$  do
4:     if  $a_{\min} > a_j$  then
5:        $\min \leftarrow j$ 
6:     end if
7:   end for
8:   if  $a_i > a_{\min}$  then
9:     swap  $a_i$  and  $a_{\min}$ 
10:  end if
11: end for
```

Which sequence will result in the most operations when using Selection Sort (**swaps** and **comparisons**)?

☐ 9, 3, 8, 6, 5, 1, 7

☒ 4, 7, 2, 9, 5, 6, 3

☐ 0, 1, 2, 3, 4, 5, 6 \rightarrow really fast

☐ 7, 6, 5, 4, 3, 2, 1 \rightarrow really fast

Exercise 2. Which function below grows fastest when n goes to infinity?

☐ $(\log_3(33))^{n-3}$

☐ 3^n

☐ $n^{3 \log_3(n)}$

☐ $n^3 \log_3(n)$

①
$$= 3^{\log_3((\log_3(33))^{n-3})}$$

$$= 3^{(n-3) \underbrace{\log_3(\log_3(33))}_{> 1}}$$

(b) 3^n

(c) $3^{\log(n^3) \cdot \log(n)} = 3^{3 \log(n)^2}$

(d) $3^{\log(n^3 \cdot \log(n))} = 3^{3 \log(n \cdot \log(n))}$

Exercise 7. Consider the two statements below, where k and ℓ are constants with $k > \ell \geq 2$ and $m \rightarrow \infty$:

$\log_m(k)$ is $\Theta(\log_m(\ell))$ $k^{\log_\ell(m)}$ is $O(\ell^{\log_k(m)})$.

- ☐ They are both false.
- ☒ Only the first is true.
- ☐ Only the second is true.
- ☐ They are both true.

$$\log_m(k) = \frac{\log_\ell(k)}{\log_\ell(m)}$$

$$\log_m(\ell) = \frac{\log_\ell(\ell)}{\log_\ell(m)}$$

Same order

$$\text{Let } c = \log_e(k).$$

$$k^{\log_e(m)} = e^{\log_e(k) \cdot \log_e(m)} = e^{\log_e(m) \cdot c}$$

$$e^{\log_e(m)} = e^{\frac{\log_e(m)}{\log_e(k)}} = e^{\log_e(m) \cdot (-c)}$$

as $c > 1$, when $m \rightarrow \infty$, m^c is not $O(m^{1/c})$

Exercise 9. Given the two statements below, where $d > 0$ is an integer constant and a_i for all $i \in \mathbf{Z}$ are positive integers with $\max_{i \in \mathbf{Z}}(a_i) = D$ for a constant $D > 0$,

$$\sum_{i=0}^n a_i i^d \text{ is } \Theta(n^{d+1})$$

$$\sum_{i=0}^d a_i n^i \text{ is } \Theta(n^d)$$

- ☐ They are both true.
- ☐ Only the first is true.
- ☐ Only the second is true.
- ☐ They are both false.

$$\sum_{i=0}^n a_i i^d \text{ is } \Theta(n^{d+1})$$

$$= a_1 1^d + a_2 2^d + \dots + a_n n^d$$

$$\geq (n \cdot D) n^d = D n^{d+1}$$

$$\sum_{i=0}^n a_i i^d \text{ is } O(n^{d+1})$$

$$\sum_{i=0}^n a_i i^d \leq C n^{d+1} ?$$

$$\sum_{i=0}^n a_i i^d$$

$$= \sum_{i=0}^{n/2-1} a_i i^d + \sum_{i=n/2}^n a_i i^d$$

$$\geq \sum_{i=n/2}^n a_i i^d$$

$$\geq \sum_{i=n/2}^n a_i \left(\frac{n}{2}\right)^d$$

$$\geq m \sum_{i=n/2}^n \left(\frac{n}{2}\right)^d = m \left(\frac{n}{2} + 1\right) \left(\frac{n}{2}\right)^d$$

$$\sum_{i=0}^n a_i i^d \text{ is } \Omega(n^{d+1}) = \frac{m}{2^{d+1}} n^{d+1}$$

try with $n=4$
 $\frac{n}{2}=2$, we will do 2,3,4
 which is $\frac{n}{2}+1$.

$$\sum_{i=0}^d a_i n^i$$

$$= a_1 n + a_2 n^2 + a_3 n^3 + \dots + a_d n^d$$

