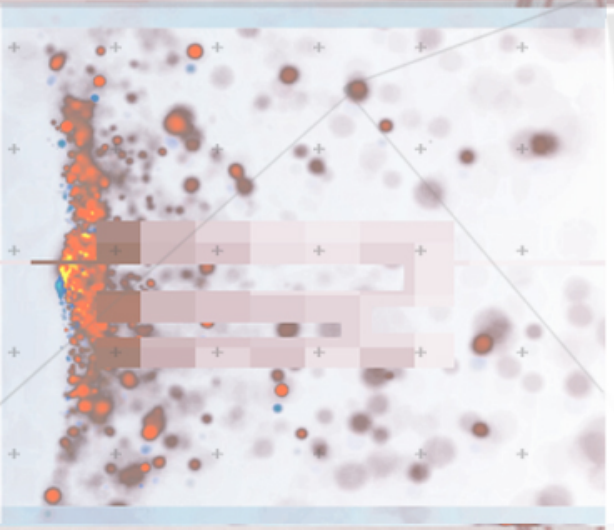
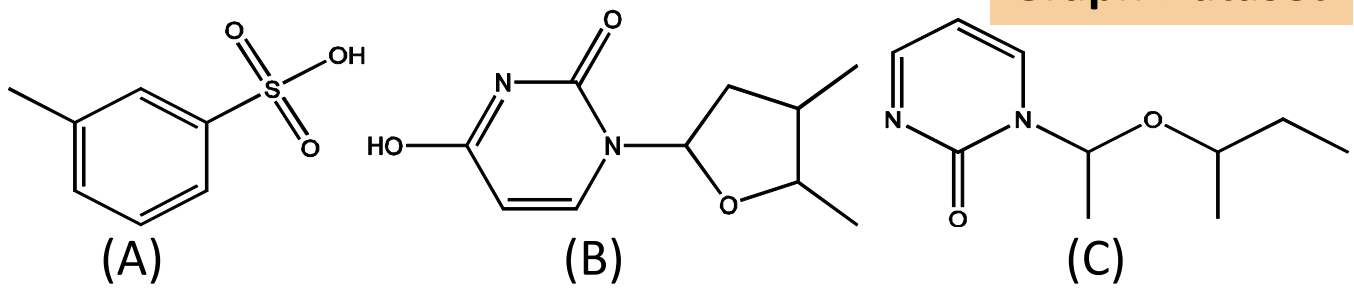


Graph Pattern and Graph Pattern Mining

Frequent (Sub)Graph Patterns

- Given a labeled graph dataset D = {G_1, G_2, …, G_n), the supporting graph set of a subgraph *g* is D_g = {G_i | $g \subseteq G_i$, $G_i \in D$}
  - support(*g*) = $|D_g|/|D|$
- A (sub)graph *g* is ***frequent*** if *support*(*g*) ≥ *min_sup*
- Ex.: Chemical structures
- Alternative:
  - Mining frequent subgraph patterns from a single large graph or network

Graph Dataset

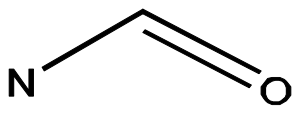


(A) (B) (C)

min_sup = 2

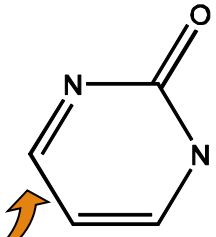
Frequent Graph Patterns

(1)



support = 67%

(2)

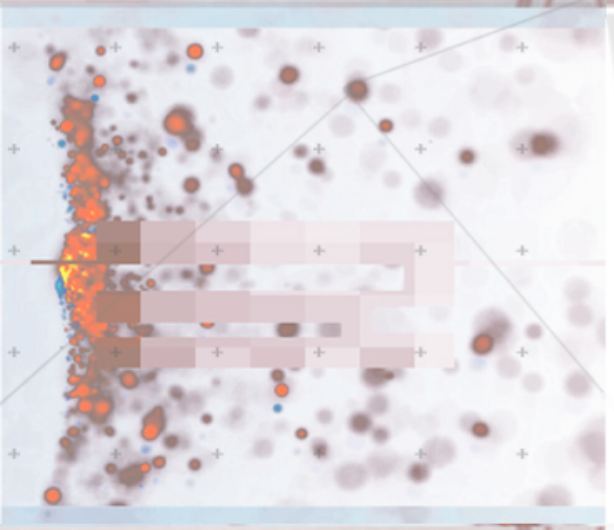


Applications of Graph Pattern Mining

- Bioinformatics
  - Gene networks, protein interactions, metabolic pathways
- Chem-informatics: Mining chemical compound structures
- Social networks, web communities, tweets, …
- Cell phone networks, computer networks, …
- Web graphs, XML structures, Semantic Web, information networks
- Software engineering: Program execution flow analysis
- Building blocks for graph classification, clustering, compression, comparison, and correlation analysis
- Graph indexing and graph similarity search

Graph Pattern Mining Algorithms: Different Methodologies

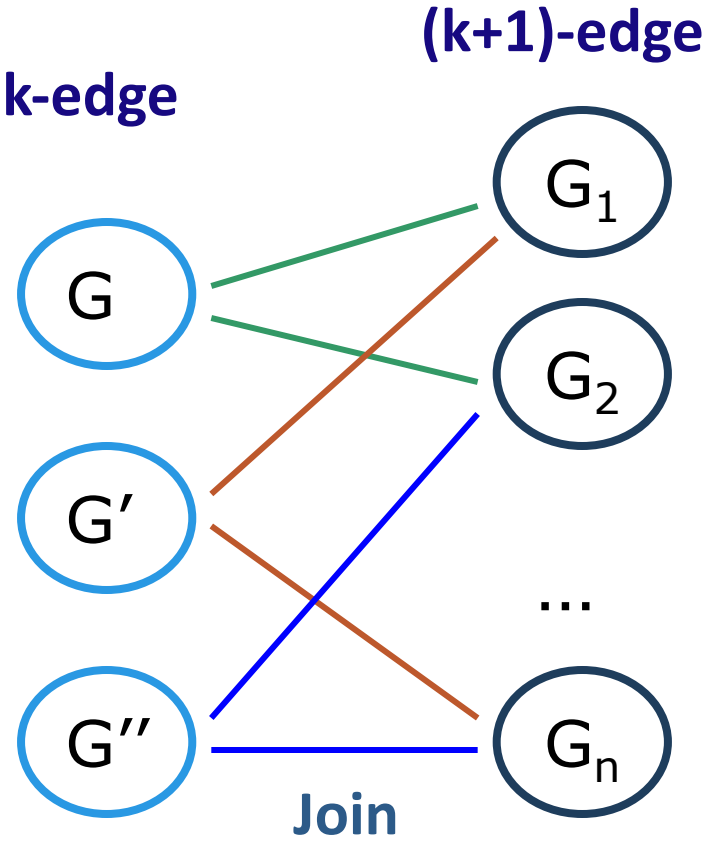
- Generation of candidate subgraphs
  - Apriori vs. pattern growth (e.g., FSG vs. gSpan)
- Search order
  - Breadth vs. depth
- Elimination of duplicate subgraphs
  - Passive vs. active (e.g., gSpan [Yan & Han, 2002])
- Support calculation
  - Store embeddings (e.g., GASTON [Nijssen & Kok, 2004], FFSM [Huan, Wang, & Prins, 2003], MoFa [Borgelt & Berthold, ICDM'02])
- Order of pattern discovery
  - Path → tree → graph (e.g., GASTON [Nijssen & Kok, 2004])



Graph Pattern Mining: Apriori-Based Approach

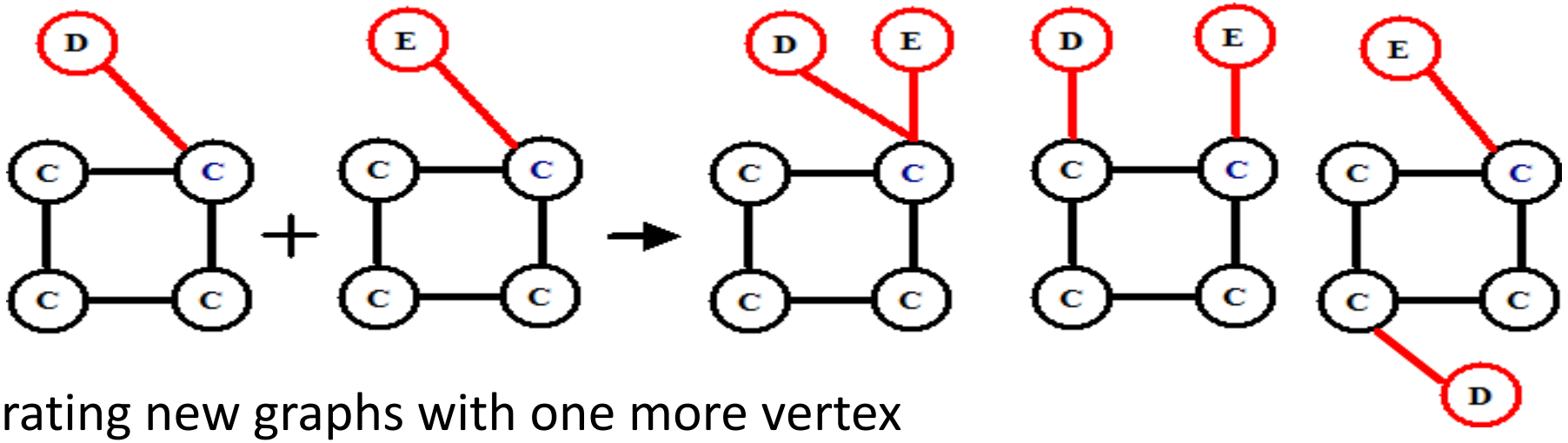
Apriori-Based Approach

- The Apriori property (anti-monotonicity): A size-*k* subgraph is frequent if and only if all of its subgraphs are frequent
- A candidate size-(*k*+1) edge/vertex subgraph is generated if its corresponding two *k*-edge/vertex subgraphs are frequent
- Iterative mining process:
  - Candidate-generation → candidate pruning → support counting → candidate elimination

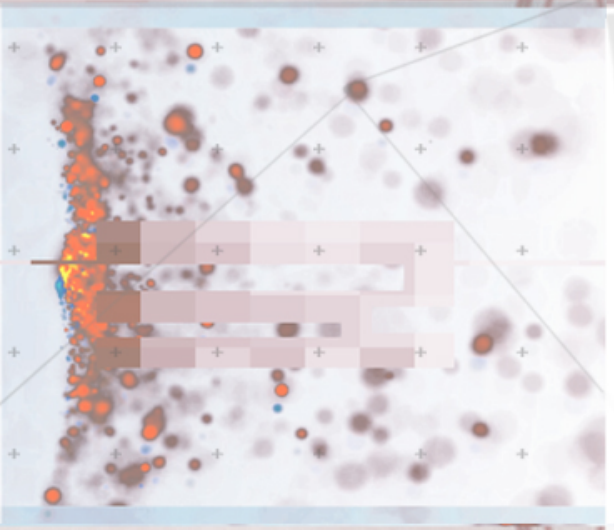


Candidate Generation: Vertex Growing vs. Edge Growing

- Methodology: Breadth-search, Apriori joining two size-*k* graphs
  - Many possibilities at generating size-(*k*+1) candidate graphs



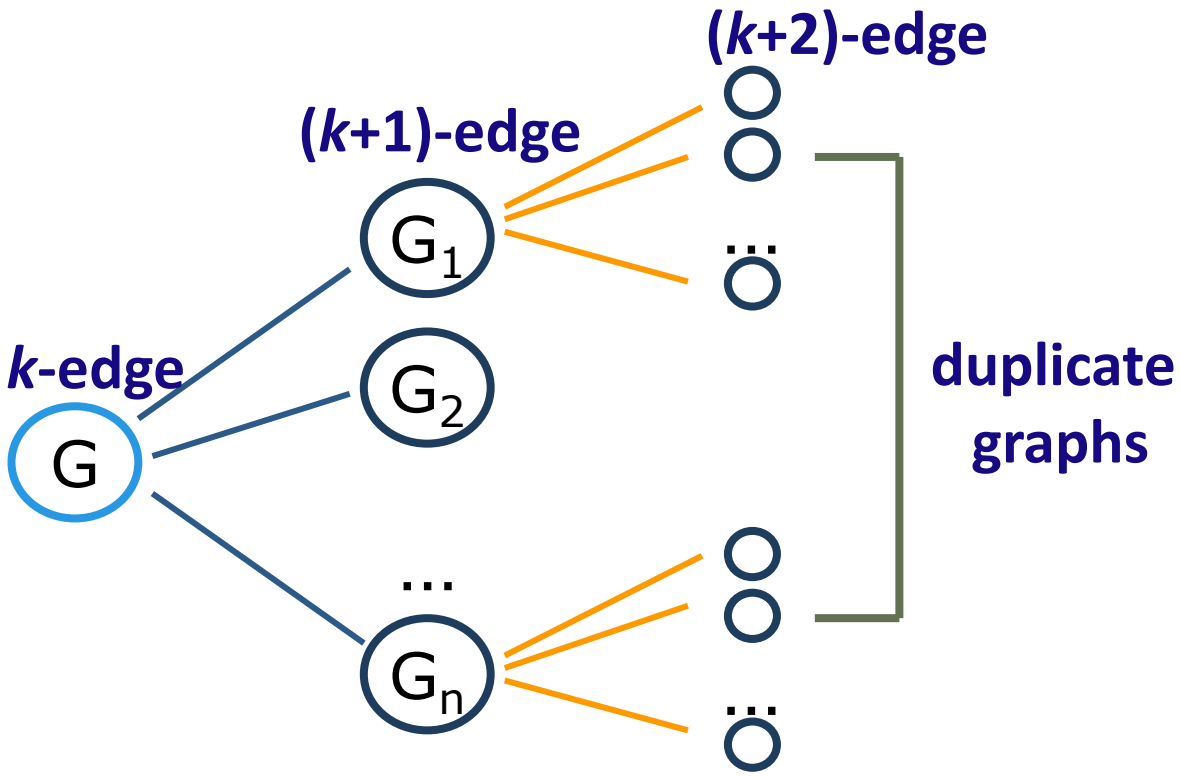
- Generating new graphs with one more vertex
  - AGM (Inokuchi, Washio, & Motoda, PKDD'00)
- Generating new graphs with one more edge
  - FSG (Kuramochi & Karypis, ICDM'01)
- Performance shows *via edge growing* is more efficient



gSpan: A Pattern Growth Approach

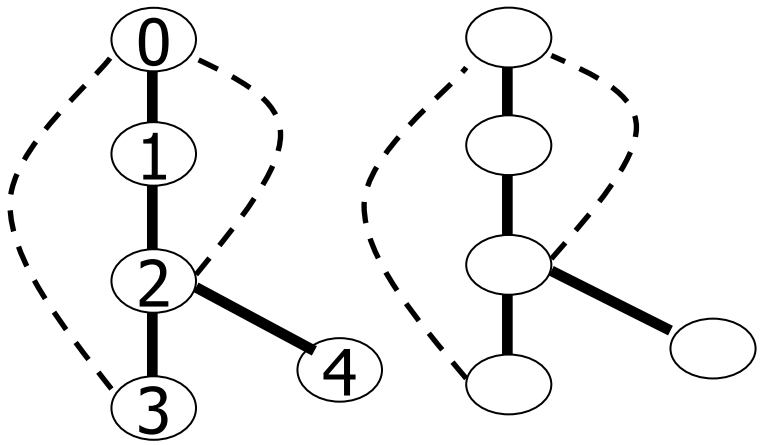
Pattern-Growth Approach

- Depth-first growth of subgraphs from *k*-edge to (*k*+1)-edge, then (*k*+2)-edge subgraphs
- Major challenge
  - Generating many duplicate subgraphs
- Major idea to solve the problem
  - Define an order to generate subgraphs
  - DFS spanning tree: Flatten a graph into a sequence using depth-first search
  - gSpan (Yan & Han, ICDM'02)

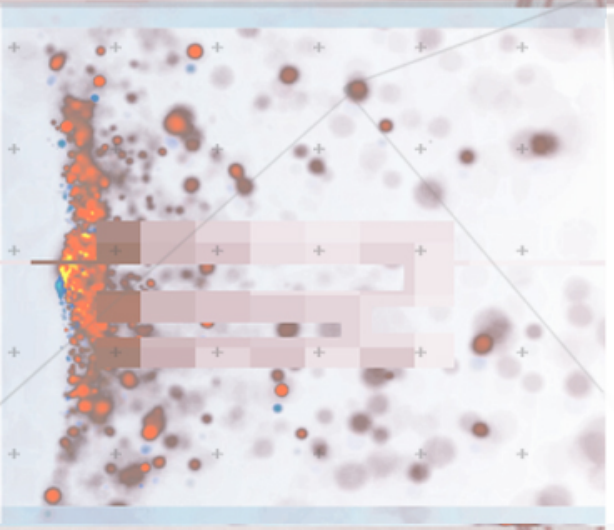


gSPAN: Graph Pattern Growth in Order

- **Right-most path extension** in subgraph pattern growth
  - Right-most path: The path from root to the right-most leaf (choose the vertex with the smallest index at each step)
  - Reduce generation of duplicate subgraphs
- **Completeness:** The enumeration of graphs using right-most path extension is complete
- DFS code: Flatten a graph into a sequence using depth-first search



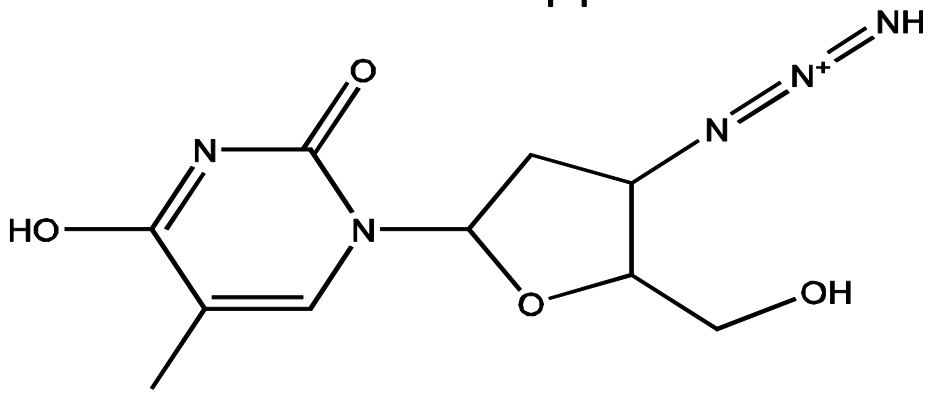
e_0: (0,1)
e_1: (1,2)
e_2: (2,0)
e_3: (2,3)
e_4: (3,0)
e_5: (2,4)



CloseGraph: Mining Closed Graph Patterns

Why Mine Closed Graph Patterns?

- Challenge: An **n**-edge frequent graph may have 2^n subgraphs
- Motivation: Explore *closed frequent subgraphs* to handle graph pattern explosion problem
- A frequent graph G is *closed* if there exists no supergraph of G that carries the same support as G

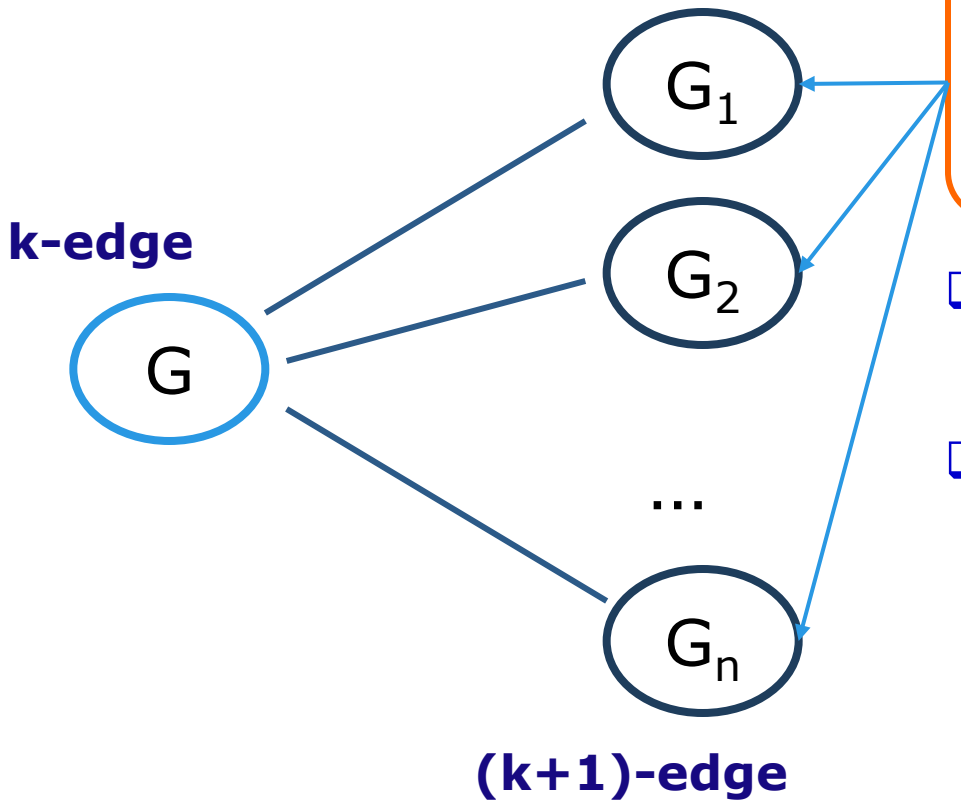


If this subgraph is *closed* in the graph dataset, it implies that none of its frequent super-graphs carries the same support

- *Lossless compression:* Does not contain non-closed graphs, but still ensures that the mining result is complete
- Algorithm CloseGraph: Mines closed graph patterns directly

CloseGraph: Directly Mining Closed Graph Patterns

- CloseGraph: Mining closed graph patterns by extending gSpan (Yan & Han, KDD'03)



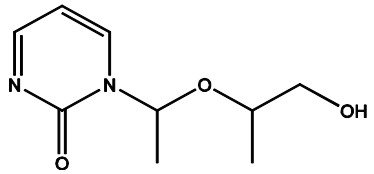
At what condition can we stop searching their children, i.e., early termination?

- Suppose G and G_1 are frequent, and G is a subgraph of G_1
- If **in any part of the graph in the dataset where G occurs, G_1 also occurs**, then we need not grow G (except some special, subtle cases), since none of G's children will be closed except those of G_1

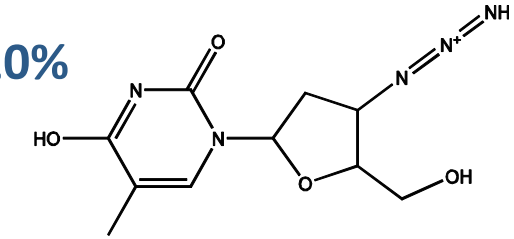
Experiment and Performance Comparison

- The AIDS antiviral screen compound dataset from NCI/NIH
- The dataset contains 43,905 chemical compounds
- Discovered patterns: The smaller minimum support, the bigger and more interesting subgraph patterns discovered

20%



10%



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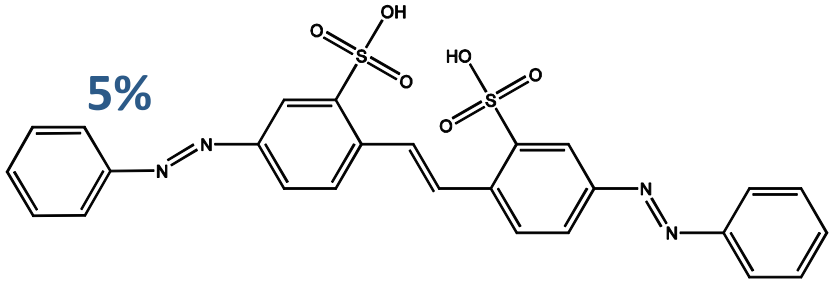
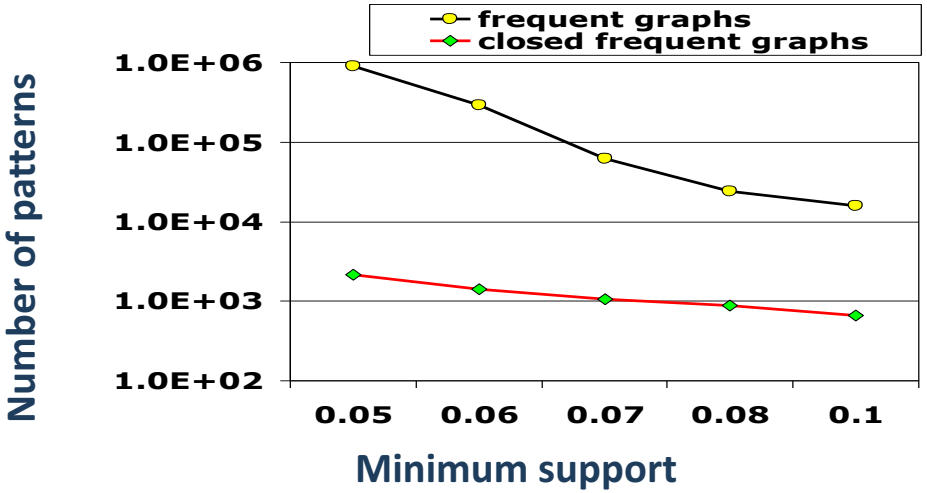
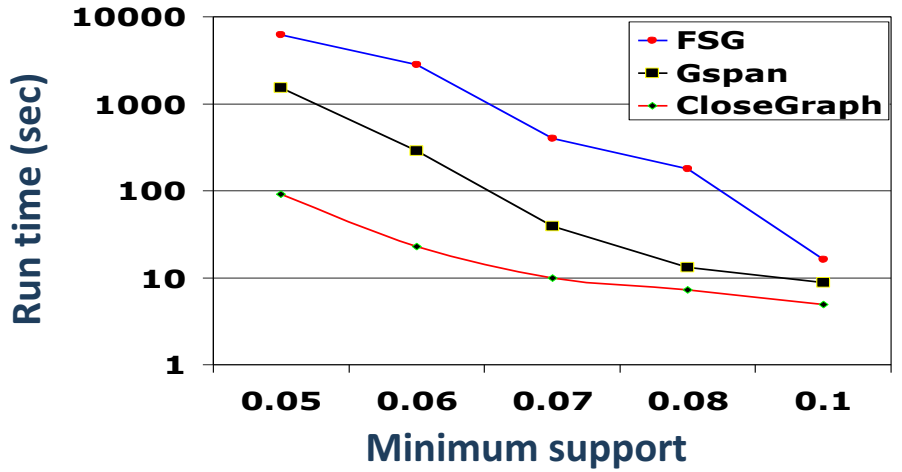
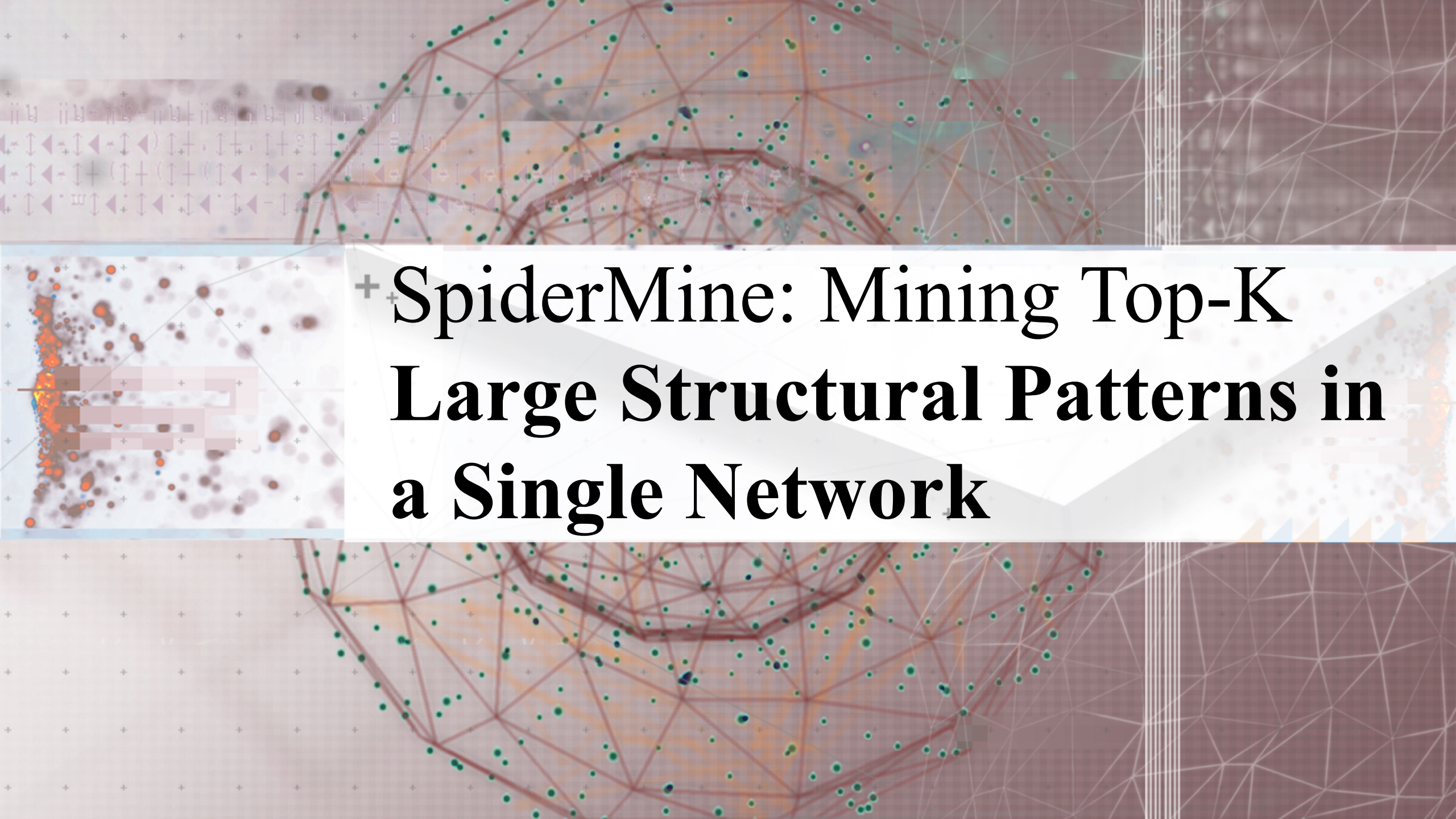


# of Patterns: Frequent vs. Closed



Runtime: Frequent vs. Closed

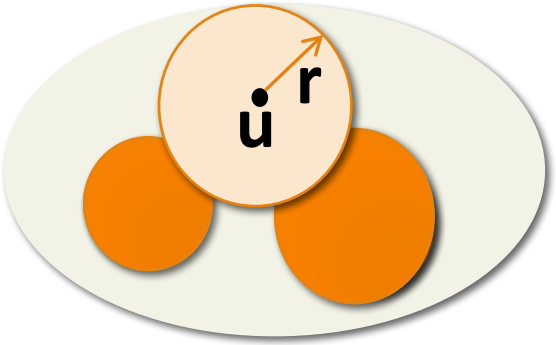




SpiderMine: Mining Top-K **Large Structural Patterns in a Single Network**

SpiderMine: **Mining Top-K Large Structural Patterns in a Massive Network**

- Large patterns are informative to characterize a large network (e.g., social network, web, or bio-network)
- Similar to pattern fusion, mining large patterns should not aim for completeness but for representativeness of the target results
- SpiderMine (Zhu et al., VLDB'11): Mine top-*K* largest frequent substructure patterns whose diameter is bounded by D_{max} with a probability at least $1-\epsilon$
- General idea: Large patterns are composed of a number of small components ("spiders"), which will eventually connect together after some rounds of pattern growth
- **r-Spider:** An r-spider is a frequent graph pattern P such that there exists a vertex u of P, and all other vertices of P are within distance r from u

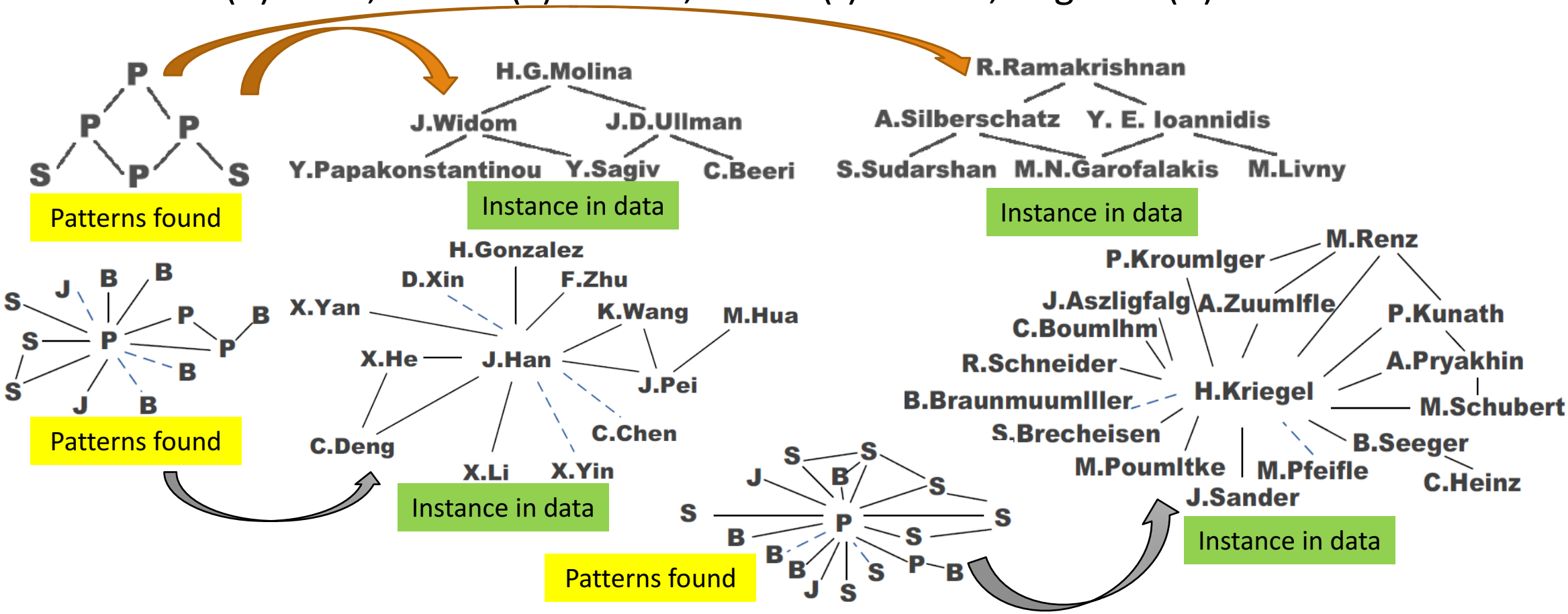


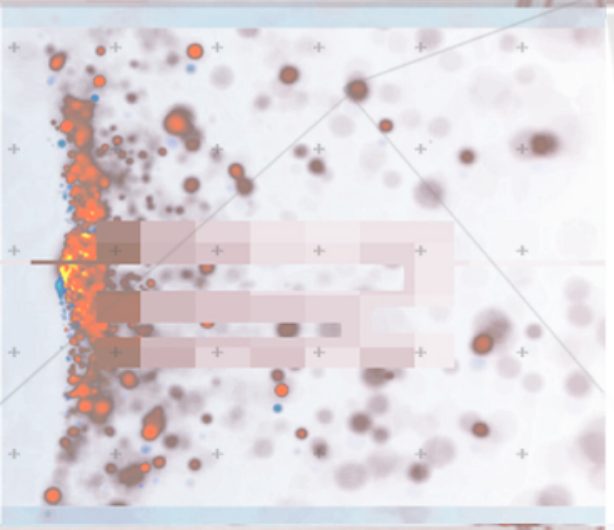
Why Is SpiderMine Good for Mining Large Patterns?

- The SpiderMine algorithm
  - Mine the set S of all the r-spiders
  - Randomly draw M r-spiders
  - Grow these M r-spiders for t = $D_{max}/2$ iterations, and merge two patterns whenever possible
  - Discard unmerged patterns
  - Continue to grow the remaining ones to maximum size
  - Return the top-K largest ones in the result
- Why is SpiderMine likely to retain large patterns and prune small ones?
  - Small patterns are much less likely to be hit in the random draw
  - Even if a small pattern is hit, it is even less likely to be hit multiple times
  - The larger the pattern, the greater the chance it is hit and saved

Mining Collaboration Patterns in DBLP Networks

- Data description: 600 conferences, 9 major CS areas, 15,071 authors in DB/DM
- Author labeled by # of papers published in DB/DM
  - Prolific (P): ≥ 50, Senior (S): 20~49, Junior (J): 10~19, Beginner(B): 5~9

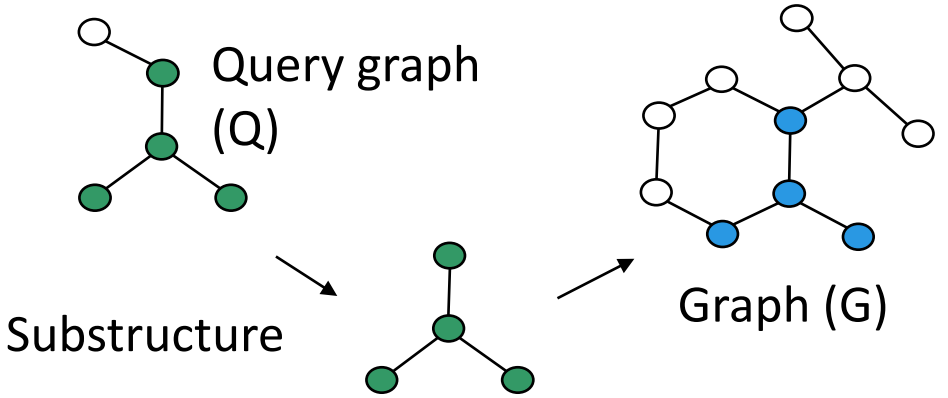




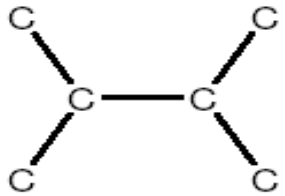
Graph Pattern Mining
Application I: Graph Indexing

Application of Pattern Mining I: Graph Indexing

- Graph query: Find all the graphs in a graph DB containing a given query graph
- Index should be a powerful tool
- Path-index may not work well
- Solution: Index directly on substructures (i.e., graphs)



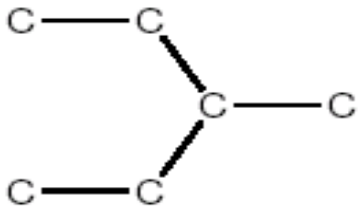
Query Q:



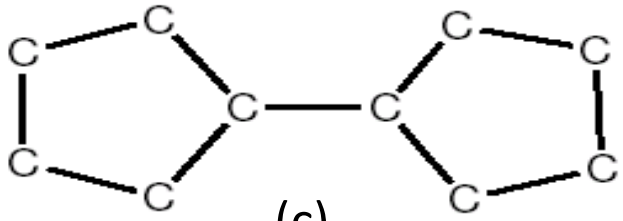
Only graph (c) contains Q

Graph DB:

(a)



(b)



(c)

Path-indices: C, C-C, C-C-C, C-C-C-C cannot prune (a) & (b)

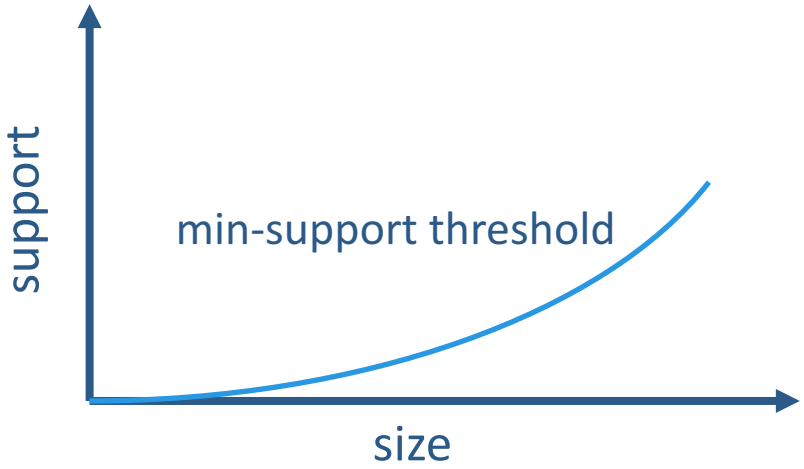
gIndex: Indexing Frequent and Discriminative Substructures

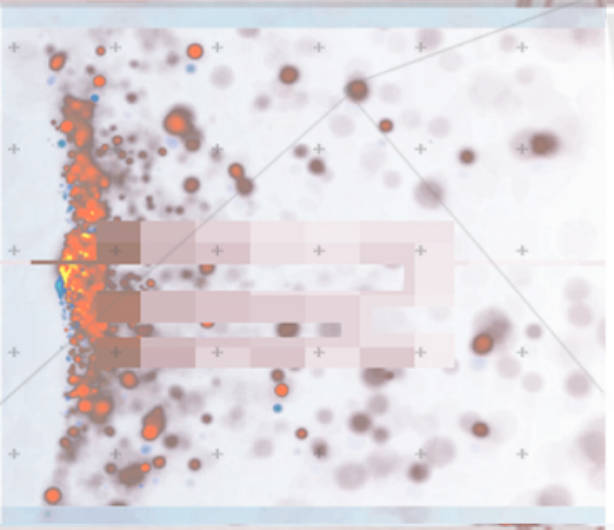
- Why index frequent substructures?
  - Too many substructures to index
  - Size-increasing support threshold
  - Large structures will likely be indexed well by their substructures
- Why discriminative substructures?
  - Reduce the index size by an order of magnitude
- Selection: Given a set of selected structures $f_1, f_2, .. f_n$, and a new structure x, the extra indexing power is measured by

$$\Pr(x | f_1, f_2, \ldots f_n), f_i \subset x$$

when Pr(x|f_1, f_2, …, f_n) is small enough, x is a discriminative structure and should be included in the index

- Experiments show that gIndex is small, effective, and stable

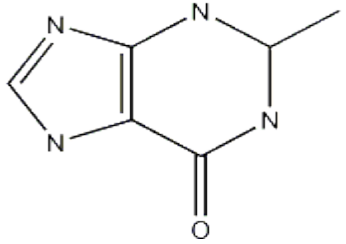


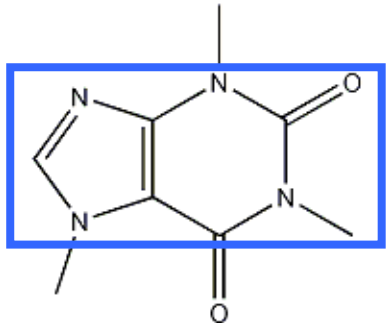


Graph Pattern Mining Application II: Graph Similarity Search

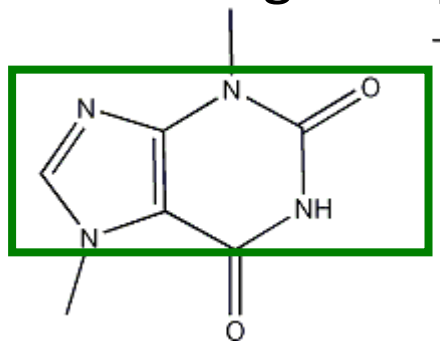
Application II: Support Substructure Similarity Search

- Find graphs in a graph DB containing substructures similar to a given query graph
- Ex. Data: A chemical compound DB
  - A query graph q:

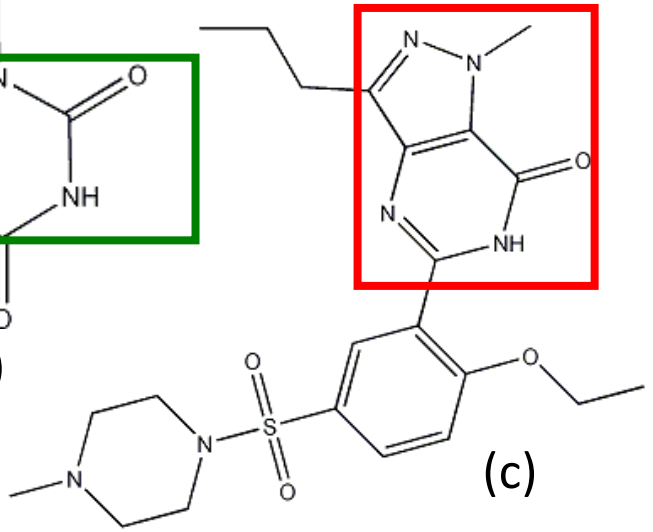




(a)



(b)

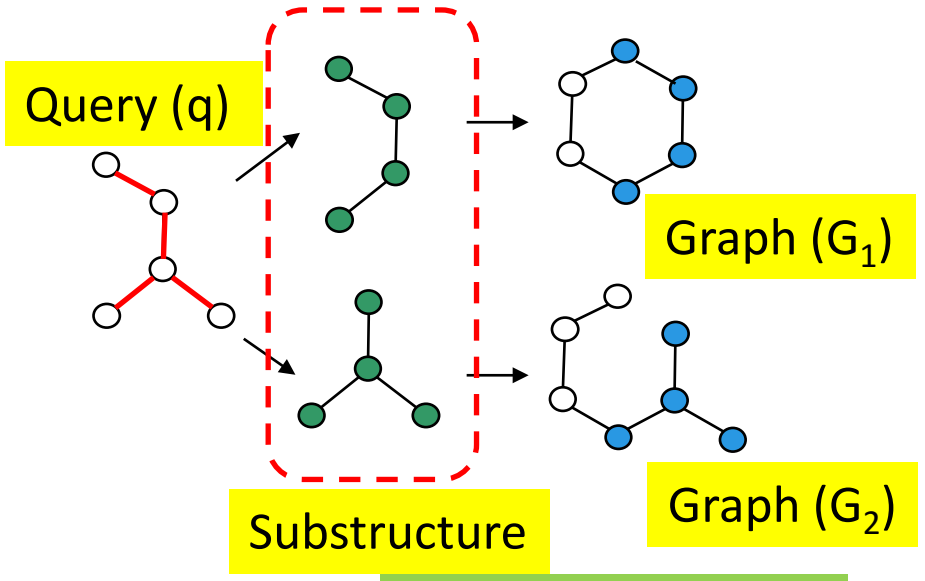


(c)

- How to do similarity search efficiently?
  - No indexing? – Sequential scan + computing subgraph similarity – too costly!
  - Build graph indices to support approximate search?
    - Need an explosive number of subgraphs to cover all the *similar* subgraphs!
- An elegant solution (Yan, Yu, & Han, SIGMOD'05):
  - Keep the graph index structure, but select features in the query space

Feature-Based Similarity Search

- Decompose a query graph into a set of features
- Feature-based similarity measure
  - Each graph is represented as a feature vector $X = \{x_1, x_2, \ldots, x_n\}$
  - Similarity is defined by the distance of their corresponding vectors
- If graph G contains the major part of a query graph q, G should share a number of common features with q
  - Given a relaxation ratio, one can calculate the maximal number of features that can be missed!



Assume: Query graph has 5 features

Relaxation threshold: Can miss at most 2 features

Then: G_1, G_2, G_3 are pruned

Graphs in database

features	G_1	G_2	G_3	G_4	G_5
f_1	0	1	0	1	1
f_2	0	1	0	0	1
f_3	1	0	1	1	1
f_4	1	0	0	0	1
f_5	0	0	1	1	0

A feature-graph matrix