

n = 1 : equivalent to part 1) of Def

n = 2 equivalent to surjectivity in part 2) of
Def g.

n = 3 equivalent to injectivity in part 2) of
Def g. □

- . The fibers of left/right fibrations are Kan complexes. This has some further consequences.

Recall that $\widetilde{s\text{Set}}$ is a simplicial category.

Let C be an ∞ -category. The slice category $\widetilde{s\text{Set}}_C$ has itself a structure of simplicial category.

Prop 12: Let $\begin{cases} p: D \rightarrow C \\ q: D' \rightarrow C \end{cases}$ be left fibrations.

between ∞ -categories.

The simplicial set $\widetilde{s\text{Set}}_C(p, q)$ of $\widetilde{s\text{Set}}_C$ is

a Kan complex.

Cor 13: Let C be an ∞ -category. The simplicial subcategory $\widetilde{L\text{fib}}(C)$ of $\widetilde{s\text{Set}}_C$ spanned by left fibrations is locally Kan.

Def 14: The ∞ -category of left fibrations

$L\text{Fib}(C)$ is defined as $N_\Delta(\widetilde{L\text{Fib}}(C))$.

- Let $\text{Spc}_* := \text{Spc}_{*,/}$ be the ∞ -category of pointed spaces. The canonical functor $\text{Spc}_* \longrightarrow \text{Spc}$ is a left fibration.

Rmk: I mistakenly claimed earlier that

Spc_* is isomorphic to $N_\Delta(\widetilde{\text{Kan}}_{\Delta^0, /})$.

The situation is slightly more complicated:

there is a natural functor

$$N_{\Delta}(\tilde{\text{Kan}}_{\Delta^{\circ}/}) \longrightarrow N_{\Delta}(\tilde{\text{Kan}})_{\Delta^{\circ}/}$$

which is an equivalence of ∞ -categories.

See Kerodon, Prop. 5.5.3.5.

- We can now state the main theorem of this section:

Thm 15 : (Joyal, Lurie)

There is an equivalence of ∞ -categories

$$\mathcal{S}: \text{Fun}(C, \text{Spc}) \xrightarrow{\sim} \text{LFib}(C)$$

which on objects sends a functor

$$F: C \longrightarrow \text{Spc} \text{ to the pullback:}$$

$$\begin{array}{ccc} \mathcal{F} & \longrightarrow & \mathbf{Spc}_* \\ \downarrow & \lrcorner & \downarrow \\ C & \longrightarrow & \mathbf{Spc} \end{array}$$

Dually, there is an equivalence of categories

$$\int : \mathbf{Fun}(C^{\text{op}}, \mathbf{Spc}) \xrightarrow{\sim} \mathbf{RFib}(C)$$

which on objects sends $G : C^{\text{op}} \rightarrow \mathbf{Spc}$
to the pullback:

$$\begin{array}{ccc} \int G & \longrightarrow & \mathbf{Spc}_*^{\text{op}} \\ \downarrow & \lrcorner & \downarrow \\ C & \xrightarrow{G^{\text{op}}} & \mathbf{Spc}^{\text{op}} \end{array}$$

We can thus say that $\mathbf{Spc}_* \rightarrow \mathbf{Spc}$ (resp.
 $\mathbf{Spc}_*^{\text{op}} \rightarrow \mathbf{Spc}^*$) is the universal left (resp.
right) fibration.

Proofs:

- This theorem has a number of different proofs in the literature; among them:

I) [HTT, Thm 2.2.1.2]

II) [Cisinski, Thm 5.4.5 and 7.8.9] (\sim)

III) Heuts and Moerdijk, Left fibrations
and Homotopy colimits II

} the
simplest?

(See also Beardsley-Wong, Operadic nerve..., App B)

- All of them require presenting the equivalence using model structures. The proofs I) and III) also rely on the comparison with simplicial categories as model for ∞ -categories.

• Let me just mention the necessary model structures (in Lurie and Heuts-Moerdijk).

+ Let \mathcal{C} be a simplicial category.

There is a simplicial category of simplicial functors $\text{Fun}(\mathcal{C}, \widetilde{s\text{Set}})$, and we can equip it with a **projective model structure** (with respect to $s\text{Set}_{\text{Kan-Quillen}}$):

* weak equivalences are objectwise weak homotopy equivalences.

* fibrations are objectwise Kan fibrations.

This is a simplicial model structure, and if \mathcal{C} is locally Kan the associated ∞ -category is equivalent to $\text{Fun}(N_{\Delta}(\mathcal{C}), \text{Spc})$.

+ Let X be a simplicial set. Then

\widetilde{sSet}_X has a (simplicial) model structure,

the covariant model structure, where

- * the cofibrations are the monomorphisms
- * the fibrant objects are the left fibrations

(Recall that by a result of Joyal this determines the model structure uniquely if it exists, see e.g. [Riehl, Categorical Homotopy theory, Thm 15.3.1].)

One can describe the weak equivalences quite precisely, see [HTT, Definition 2.1.4.5].

If $X = \mathcal{C}$ is an ∞ -category, the associated ∞ -category is equivalent to $L\text{Fib}(\mathcal{C})$.

- The equivalence of the theorem is

obtained by "deriving" a Quillen equivalence between simplicial model

categories. But what does this mean?

Rmk: Let $F: M \rightleftarrows N: G$ be a Quillen equivalence. It is natural to expect it to give rise to an equivalence of ∞ -categories; this is true (Mazel-Gee, "Quillen adjunctions induce adjunctions of quasicategories") , but surprisingly subtle:

- In this course, we only discussed how to associate ∞ -categories to simplicial model categories. It is possible more generally, but this requires other constructions.
- Even if M, N are simplicial, neither F nor G preserves bifibrant objects in general. (they each preserve half!)

Nevertheless, Lurie proves a result in HTT which is sufficient for the theorem above:

Thm : [HTT, Corollary A.3.1.12]

Let M, N be simplicial model categories.

Let $F: M \rightleftarrows N: G$ be a Quillen equivalence.

Assume that:

- Every object of M is cofibrant.
- G is a simplicial functor.

Then G induces a functor of simplicial categories

$N_{cf} \xrightarrow{G_{cf}} M_{cf}$, and the R.c. nerve

$$N_\Delta(N_{cf}) \xrightarrow{N_\Delta(G_{cf})} N_\Delta(M_{cf})$$

is an equivalence of ∞ -categories. □

This is proven by checking that G_{cf} is a Dwyer-Kan equivalence of simplicial cats.

Rmk 16: This formulation of the theorem

is unsatisfying because it only describes what

happens on objects. It would be better to

have a purely ∞ -categorical characterisation

of the functor, but I don't know of

one!

Terminology:

Lurie calls this equivalence the “straightening / unstraightening” equivalence:

- unstraightening : $\text{Fun}(C, \text{Spc}) \rightarrow \text{Lfib}(C)$
- straightening : $\text{Lfib}(C) \rightarrow \text{Fun}(C, \text{Spc})$.

- Let us explore a little the “straightening” process. We start with a left fibration $p: D \rightarrow C$. We won't quite produce a functor $C \rightarrow \text{Spc}$, but a functor $R_C: \text{R}\text{Spc} \simeq \text{R}\text{Kan}$.

- For $c \in C$, the fiber D_c is a Kan complex.
- Let $f: c \rightarrow c'$ be a morphism in C .

The inclusion $D_c = D_c \times \{0\} \hookrightarrow D_c \times \Delta^1$

is left anodyne because $\{0\} \hookrightarrow \Delta^1$ is.
(Special case of Prop III.19)

So we can solve the lifting problem:

$$\begin{array}{ccc} D_c & \xrightarrow{\quad} & D \\ \downarrow & \nearrow & \downarrow P \\ D_c \times \Delta^1 & \xrightarrow{\quad} & \Delta^1 \xrightarrow{g} C \end{array}$$

Restricting \nearrow to $D_c \times \{1\}$, we get

$$g_! : D_c \longrightarrow D_{c'}$$

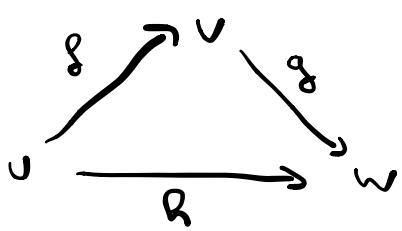
$g_!$ is not uniquely determined. However it is determined up to homotopy, and:

Lemma 17: The assignment

$$c \mapsto D_c, g \mapsto [g_!]$$

determines a functor $R\mathcal{C} \longrightarrow R\text{Kan}$.

proof: Let $\sigma : \Delta^2 \longrightarrow C$ be a 2-simplex



Choose $g_!, g_!, h_!$

Using the fact that $D_c \times \{0\} \hookrightarrow D_c \times \Delta^2$

is also left anodyne, one shows that

$[h_!] = [g_!] \circ [g_!]$ in hKan .

(see [HTT, Lemma 2.1.1.4])

□

Rmk : The functor classified by a

given left/right fibration is only defined

up to isomorphism in $\text{Fun}(C^{(\text{op})}, \text{Spc})$, so

“up to a contractible choice.” This will

apply to all the examples in the next section.

4) Applications

Terminology: As in the 1-categorical case, we call a functor $C^{\text{op}} \rightarrow \text{Spc}$ a **presheaf** on the ∞ -category C and sometimes write $\mathbf{PSh}(C) := \text{Fun}(C^{\text{op}}, \text{Spc})$.

• Representable and corepresentable functors

Let C be an ∞ -category and $x \in C$.

We have :

- a left fibration $C_{x/} \longrightarrow C$
- a right fibration $C_{/x} \longrightarrow C$

The functors classified by them are:

- the corepresentable functor

$$R^x : C \longrightarrow \text{Spc}$$

- the representable functor (or presheaf)

$$R_x : \mathcal{C}^{\text{op}} \rightarrow \text{Spc}$$

Let us compute their values on objects:

R_x^y is by definition isomorphic in Spc

to $(\mathcal{C}_{x/})_y$ and $R_x(y)$ is iso. to $(\mathcal{C}_{/x})_y$.

These simplicial sets are variants of the mapping spaces $\text{map}_{\mathcal{C}}(x, y)$.

Def 18: Let \mathcal{C} be an ∞ -category and

$x, y \in \mathcal{C}_0$. The left(-pinched) mapping space $\text{map}_{\mathcal{C}}^L(x, y)$ is defined as

$$\text{map}_{\mathcal{C}}^L(x, y) := (\mathcal{C}_{x/})_y.$$

More concretely, for any $n \geq 0$, we have

$$\text{map}_C^L(x, y)_n = \left\{ \epsilon \in C_{n+1} \mid \begin{array}{l} \epsilon(0) = x, d_0(\epsilon) \\ \text{constant map } \tilde{\Delta} \rightarrow \{y\} \cap C \end{array} \right\}$$

Dually, the right(-pinched) mapping space

$$\text{map}_C^R(x, y) := (C_{/x})_y$$

More concretely, for any $n \geq 0$, we have

$$\text{map}_C^R(x, y)_n = \left\{ \epsilon \in C_{n+1} \mid \begin{array}{l} d_{n+1}(\epsilon) \text{ constant map at } x \\ \epsilon(n+1) = y \end{array} \right\}$$



As fibers of left/right fibrations, those are Kan complexes. Moreover:

Prop 19: There are natural monomorphisms

$$\text{map}_C^L(x, y) \hookrightarrow \text{map}_C(x, y) \hookleftarrow \text{map}_C^R(x, y)$$

which are homotopy equivalences.



See [Kerodon, Prop 4.6.6.9].

- Because of this, it is also customary to write:

$$\begin{cases} h_{\alpha} = \text{map}_C(-, \alpha) : C^{\text{op}} \rightarrow \text{Spc} \\ h^{\alpha} = \text{map}_C(\alpha, -) : C \rightarrow \text{Spc} \end{cases}$$

- These new mapping spaces lead to reformulations of initial / terminal objects.

Cor : Let C be an ∞ -category and

$\alpha \in C_0$. TFAE:

(i) α is terminal.

(ii) $\forall y \in C_0$, $\text{map}_C^L(\alpha, y)$ is contractible.

(iii) map_C^R

(iv) map_C

proof: By Prop., (ii) - (iv) are equivalent. We check (i) \Leftrightarrow (iii).

- (i) is equivalent to: $C_{/x} \rightarrow C$ is a trivial fibration.
- (iii) is equivalent to: $\forall y \in C_0$,
 $(C_{/x})_y = \text{map}_C^R(y, x)$ is contractible.

Since $C_{/x} \rightarrow C$ is a right fibration,
these two conditions are eq. by Prop. \square .

- There is a recognition principle for
(c)o representable functors as in 1-cat. theory:

Lemma 20: Let $F: C^{\text{op}} \rightarrow \text{Spc}$ be a

functor. TFAE:

(i) F is isomorphic to a representable functor R_C in $\text{Fun}(C^{\text{op}}, \text{Spc})$.

(ii) $\mathcal{H}F: \mathcal{H}C^{\text{op}} \rightarrow \text{Rkan}$ is a representable

functor of R -Kan-enriched categories.

(ii) The ∞ -category $\int F$ has a terminal object \tilde{c} .

In that situation, $P_F(\tilde{c}) = c$.

Proof: [HTT, 4.4.4.1 - 4.4.4.5]. \square

Mapping Spaces

We have treated the functoriality of mapping spaces separately in each variable.

We also expect a functor

$$\text{Map}_C(-, -) : C^{\text{op}} \times C \longrightarrow \text{Spc}$$

For this, we should construct either a left or right fibration to $C^{\text{op}} \times C$ (since $(C^{\text{op}} \times C)^{\text{op}} \simeq C^{\text{op}} \times C$ canonically).

Def 21: Let C be an ∞ -category.

The twisted arrow ∞ -category $\tilde{\mathcal{O}}(C)$

is the simplicial set

$$\tilde{\mathcal{G}}(C)_n := \text{sSet}(\Delta^{n,\text{op}} * \Delta^n, C)$$

There are morphisms

$$\tilde{\mathcal{G}}(C) \xrightarrow{s} C^{\text{op}} \text{ and } \tilde{\mathcal{G}}(C) \xrightarrow{t} C$$

induced by the inclusions

$$\Delta^{n,\text{op}} \hookrightarrow \Delta^{n,\text{op}} * \Delta^n \hookrightarrow \Delta^n.$$

Rmk 22: This is a simplicial variant

of the twisted arrow category of a

1-category C , which is the category of

elements of $\text{Hom}_C: C^{\text{op}} \times C \rightarrow \text{Set}$.

$\text{Ob } \tilde{\mathcal{G}}(C)$: morphisms in C

$\text{Mor}_{\tilde{\mathcal{G}}(C)}(g, g)$: diagrams

$$\begin{array}{ccc} x & \xleftarrow{p} & y \\ \downarrow g & & \downarrow g \\ y & \xrightarrow{q} & t \end{array}$$

Exercise: Let \mathcal{C} be a 1-category.

Prove that $\tilde{\mathcal{O}}(\mathcal{N}(\mathcal{C})) \simeq \mathcal{N}(\tilde{\mathcal{O}}(\mathcal{C}))$.

Prop 3: The morphism

$$(s, t) : \tilde{\mathcal{O}}(\mathcal{C}) \longrightarrow \mathcal{C}^{\text{op}} \times \mathcal{C}$$

is a left fibration; in particular,

$\tilde{\mathcal{O}}(\mathcal{C})$ is an ∞ -category.

proof: [Higher algebra, 5.2.1.3].

- There is another nice proof of this in Barwick-Glasman, on the fibrewise effective Burnside ∞ -category, § 1. □
- We must check this has the correct fibers up to equivalence:

Lemma : Let C be an ∞ -category,
and $x \in \text{Ob}(C)$. There is a canonical diagram

$$\begin{array}{ccc} C_{/x} & \longrightarrow & \tilde{G}(C) \\ \downarrow & & \downarrow \\ \{x\} \times C & \longrightarrow & C^{\text{op}} \times C \end{array}$$

in Cat_{∞} which is Homotopy-Cartesian, i.e. the
induced map

$$C_{/x} \longrightarrow \tilde{G}(C) \times (\{x\} \times C) \\ (C^{\text{op}} \times C)$$

is an equivalence of ∞ -categories.

Proof: This is contained in the proof of
[Higher Algebra, Proposition 5.2.1.10].

—

5.2.1.11



Def : By the Grothendieck construction,

the left fibration $\tilde{\mathcal{G}}(\mathcal{C}) \rightarrow \mathcal{C}^{\text{op}} \times \mathcal{C}$

classifies a functor:

$$\mathcal{C}^{\text{op}} \times \mathcal{C} \longrightarrow \text{Spc}$$

which by adjunction gives the

Yoneda embedding:

$$y_{\mathcal{C}} : \mathcal{C} \longrightarrow \text{PSh}(\mathcal{C}) := \text{Fun}(\mathcal{C}^{\text{op}}, \text{Spc})$$

Cor : Let $x \in \mathcal{C}$. Then $y_{\mathcal{C}}(x)$

is isomorphic to h_x in $\text{PSh}(\mathcal{C})$.

proof: This follows from Lemma 

and the Grothendieck construction.

Thm : (Yoneda lemma)

Let C be an ∞ -category and $F \in \text{PSh}(C)$.

Let $x \in C_0$. There is a canonical iso.

$$\underset{\text{PSh}(C)}{\text{map}} \quad (R_x, F) \xrightarrow{\sim} F(x)$$

in Sp^C given by "evaluation at id_x ".

proof: See eg [Cisinski, Theorem 5.8.13.(ii)]

Cor : (Yoneda embedding)

Let C be an ∞ -category. The functor

y_C is fully faithful; its essential

image is spanned by representable presheaves.

proof: [Cisinski, Theorem 5.8.13.(i)] or

[Higher Algebra, Proposition 5.2.1.17]

combined with [HTT, Proposition 5.1.3.1].

Prop: The functor y_C preserves all (small) limits which exist in C .

Proof: [HTT, Proposition 5.1.3.2].

Kan fibrations

Let $p: X \rightarrow C$ be a Kan fibration to an ∞ -category. It is both a left and a right fibration, so one expects that there are two different, associated functors.

This is the case, and :

Prop: Let C be an ∞ -category and

$p: X \rightarrow C$ a functor. TFAE:

(i) p is a Kan fibration.

(ii) p is a left fibration, and for any $g \in C_n$, $g_!$ is an iso in $R\text{Kan}$.

Equivalently, if F is the functor
classified by p , then F factors through \mathbf{Spc}^{\leq} .

(ccii) dual condit° for right fibrations.

5) The $(\infty, 1)$ -Grothendieck construction

We need to identify the appropriate class of fibrations. Unfortunately this is more complicated than the definition of left/right fibrations.

Def : Let $p: X \rightarrow S$ be a morphism of simplicial set and $g: x \rightarrow y$ in X_1 . We say that g is p -Cartesian if one of the equivalent statements holds:

(i) For $n \geq 2$, any lifting problem

$$\begin{array}{ccccc} & & g & & \\ & \swarrow & & \searrow & \\ \Delta^{\{n-1,n\}} & \longrightarrow & \Lambda^n & \longrightarrow & X \\ \uparrow & & \nearrow & & \downarrow p \\ \Delta^n & \longrightarrow & S & & \end{array}$$

admits a solution.

(ii) The pullback-slice map

$$X_{/\delta} \longrightarrow X_{/y} \xrightarrow{\quad} S_{/p(y)} \quad Y_{/\rho(\delta)}$$

is a trivial fibration.

There is a dual notion of

p -cocartesian morphism.

This should remind you of the Joyal lifting theorem. Indeed we can reformulate that theorem as

Prop : With the notations above, assume moreover that p is an inner fibration between ∞ -categories. TFAE :

- δ is an isomorphism.

(ii) f is p -cartesian and $p(f)$ is an iso.

(iii) f is p -cocartesian and $p(f)$ is an iso.



Def : Let $p: X \rightarrow S$ be an inner fibration in sSet. Then p is a (co)Cartesian fibration if every morphism in S admits a p -(co)cartesian lift: for every $g \in S$, there exists $f \in X$, and p -(co)cartesian with $p(f)=g$

Lemma: TFAE:

(i) p is a left fibration

(ii) p is a cocartesian fibration

and every morphism in X is p -cocartesian.

Exercise: • Prove that (co)cartesian

fibrations between ∞ -categories are isofibrat $^\circ$.

• Prove that : right fibrat $^\circ \Leftrightarrow$ conservative cartesian
fibrat $^\circ$.

Rmk: Let $p: C \rightarrow D$ be a functor

between 1-categories. The definition

of p -cartesian morphisms in C

is the same as Def .(ii) above,

with "trivial fibration" replaced by

"equivalence of category". One can then

define the condition " p is a cartesian
fibration". This is the notion

relevant to the $(1,1)$ -Grothendieck

construction.

We then have :

p cartesian fibration
 \Leftrightarrow

$N(p)$ cartesian fibration.

One can define ∞ -categories $\text{Cart}(C)$ and $\text{CoCart}(C)$ of cartesian / cocartesian fibrations over C , and prove:

Thm (Lurie)

There are equivalences of ∞ -categories

$$\begin{cases} \mathcal{S}: \text{Fun}(C, \text{Cat}_{\infty}) \xrightarrow{\sim} \text{CoCart}(C) \\ \mathcal{S}: \text{Fun}(C^{\text{op}}, \text{Cat}_{\infty}) \xrightarrow{\sim} \text{Cart}(C) \end{cases}$$
□

Rmk:

- It is possible to write down the universal (co)cartesian fibration and thus at least describe what \mathcal{S} does on objects (pullback the universal fibrat°)
- Once again, the proof involves

deriving a Quillen equivalence of simplicial model categories. For this version of the theorem, as far as I know, there is only one reference:

[HTT, § 3.2].

Let's finish with two applications:

Ex : Let C be an ∞ -category.

Consider $\mathcal{G}(C) := \text{Fun}(\Delta^*, C) \xrightarrow{s} \text{Fun}(\Delta^{\{0\}}, C) \simeq C$

$$\xrightarrow{t} \text{Fun}(\Delta^{[n]}, C) \simeq C$$

Then $\begin{cases} s: \mathcal{G}(C) \rightarrow C \text{ is a cartesian fibration} \\ t: \mathcal{G}(C) \rightarrow C \text{ is a cocartesian fibration.} \end{cases}$

The functors classified by s, t can be written as

$$C_{\cdot /}: C^{\text{op}} \longrightarrow \text{Cat}_{\infty}$$

$$C_{/ \cdot} : C \longrightarrow \text{Cat}_{\infty}$$

and on objects and morphisms of C they recover (∞)slices together with their functoriality..

Def : An adjunction between

∞ -categories is a bicartesian

(= cartesian + cocartesian) fibration.

$$a : X \longrightarrow \Delta^1.$$



By straightening in both directions, we

get : $\begin{cases} \Delta^1 \longrightarrow \text{Cat}_{\infty} \\ (\Delta^1)^{\text{op}} \longrightarrow \text{Cat}_{\infty} \end{cases}$, that is

Functors $\begin{cases} f : X_0 \longrightarrow X_1 \\ g : X_1 \longrightarrow X_0 \end{cases}$

g is the left adjoint and g the right adjoint.

This definition is pretty strong and implies relatively easily other reasonable definitions of adjoints:

- natural equivalence

$$\mathrm{map}_{X_0}(-, g(-)) \simeq \mathrm{map}_{X_1}(g(-), -)$$

of functors to Spc .

- unit and counit + triangle identities

It is quite remarkable that all these definitions are actually equivalent!

"Adjunct $^\circ$ of ∞ -categories are automatically homotopy coherent".

[Riehl-Verity, Homotopy coherent adjunkt $^\circ$ (...)]

THE

END

?