

We can also apply the free cocompletion property to prove that $\text{PSh}(C)$ (and in part. sSet) is cartesian closed.

def 14 A category C with finite products is cartesian closed if for all $a \in C$, the functor

$$a \times - : C \longrightarrow C$$

has a right adjoint, which is then denoted $\underline{\text{Hom}}(a, -)$ and called the internal Hom

on exponentiation. □

Prop 15 $\text{PSh}(\mathcal{C})$ is cartesian

closed and

$$\underline{\text{Hom}}(F, G)(x) = \text{PSh}(\mathcal{C})(F \times_{\mathcal{C}} G(x), G)$$

Proof: Can be done by hand, but

also as applicat° of Thm 9.b) :

- $F \times -$ is colimit-preserving
(because {products
colimits are computed
objectwise}) \rightsquigarrow it has a right
adjoint, given by this formula. □

3) Structure of Δ and applications

- We now go into the structure of Δ and what it means for $sSet$.

Notation:

Following [Rezk], we write

$$f = \langle f_0 \dots f_n \rangle : [n] \rightarrow [m]$$

$$k \mapsto f_k$$

def 16 There are distinguished morphisms
for every $0 \leq i \leq n$:

$$\delta^i := \langle 0 \dots \hat{i} \dots n \rangle : [n-1] \hookrightarrow [n]$$

(face morphisms)

$$\sigma^i := \langle 0 \dots i \dots n \rangle : [n+1] \rightarrow [n]$$

(degeneracy morphisms)

If $X \in \text{sSet}$, we write

$$\begin{cases} d_i = (\delta^i)^*: X_n \longrightarrow X_{n-1} & \text{(face maps)} \\ \delta_i = (\epsilon^i)^*: X_n \longrightarrow X_{n+1} & \text{(degeneracy maps)} \end{cases}$$

Lemma 17:

a) We have the simplicial identities:

$$d_i d_j = d_{j-1} d_i , \quad i < j$$

$$\delta_i \delta_j = \delta_{j+1} \delta_i , \quad i \leq j$$

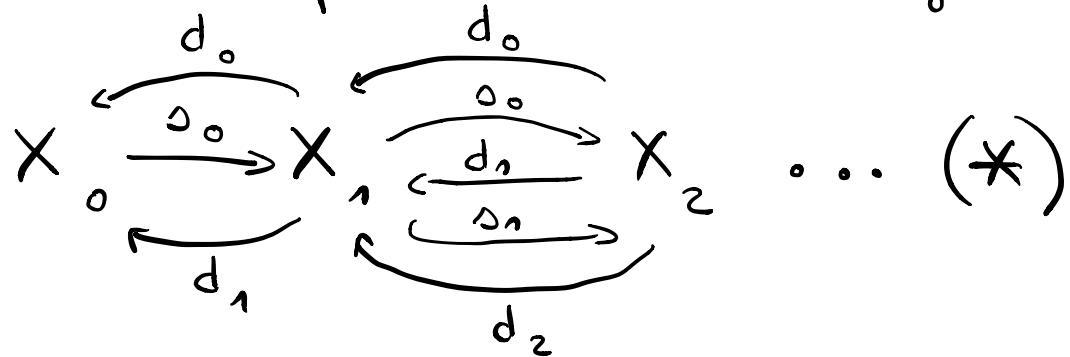
$$d_i \delta_j = \begin{cases} 1 & , i = j, j+1 \\ \delta_{j-1} d_i , & i < j \\ \delta_j d_{i-1} , & i > j+1 \end{cases}$$

b) Every morphism $g: [n] \rightarrow [m]$ can be written as

$$[n] \xrightarrow{S} [r] \xleftarrow{D} [m]$$

with $\begin{cases} S \text{ composite of deg. morphisms} \\ D \text{ face morphisms} \end{cases}$

c) The datum of a simplicial object is equivalent to a diagram



Satisfying the simplicial identities.



Proof: a) is an exercise.

b) f factors uniquely into a surjective map followed by an injective

$$\begin{array}{ccc} \text{map: } [n] & \xrightarrow{\quad} & \text{Im}(f) \hookrightarrow [m] \\ & \equiv \exists ! S \equiv & \\ S & \xrightarrow{\quad} & [n] \hookleftarrow D \end{array}$$

so it is enough to show that

- S is composition of deg. morphisms

- D ————— faces —————

Let's do the case of S (D is similar)

This is an induction on $m-n \geq 0$.

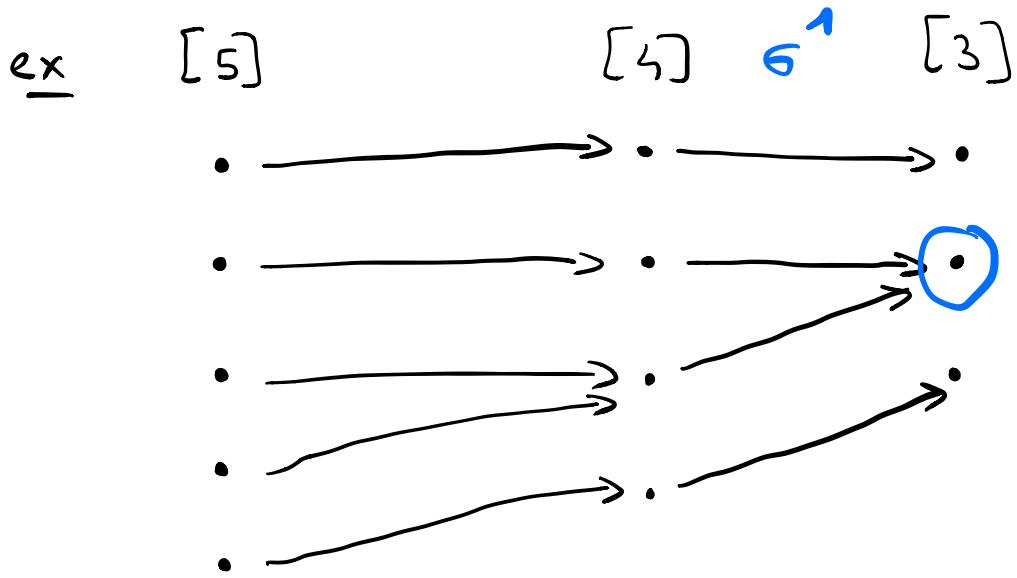
Assume $m > n$;

then $\exists i \in [n]$ such that $|D^{-1}(i)| > 1$.

Split $D^{-1}(i)$ into two non-empty sets

to get a factorisation

$$D: [n] \xrightarrow{D'} [n+1] \xrightarrow{\sigma^i} [n].$$



c) Almost follows from a) + b);

it remains to check that the

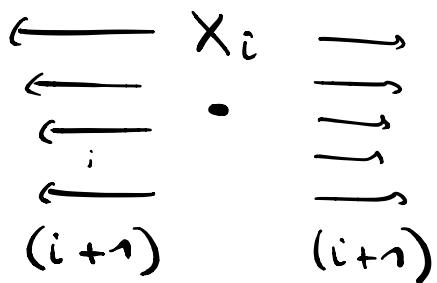
simplicial identities imply that

different choices of factorisations of

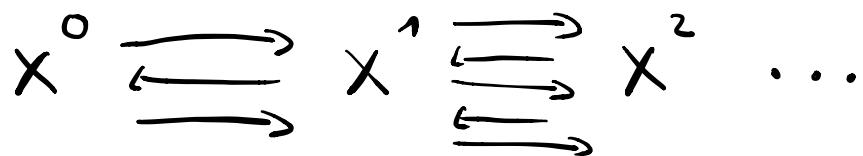
D, S yield the same map in a diagram (*). □

Rmk

a) Mnemotechnics



b) A cosimplicial object is similar:



but imbalanced # of arrows ...

def 18

- For $X_+ \in s\text{Set}$, a simplicial subset $Y_+ \subset X_+$ is the datum of $\{Y_n \subset X_n\}_n$, stable under f^* for all $f: [m] \rightarrow [n]$ in Δ .

□

Examples "Shape repertoire"!

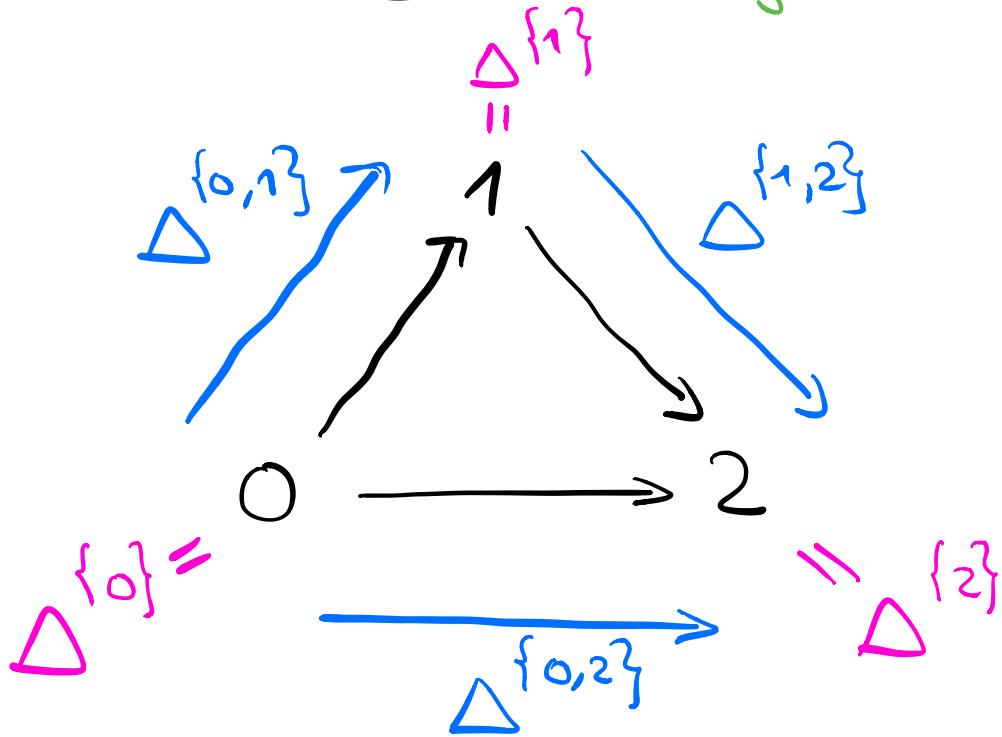
* Let $S \subseteq [n]$. We write

$\Delta^S \subseteq \Delta^n$ for the
 S -face of Δ^n :

$$(\Delta^S)_R = \left\{ f \in (\Delta^n)_R \mid \text{Im}(f) \subseteq S \right\}.$$

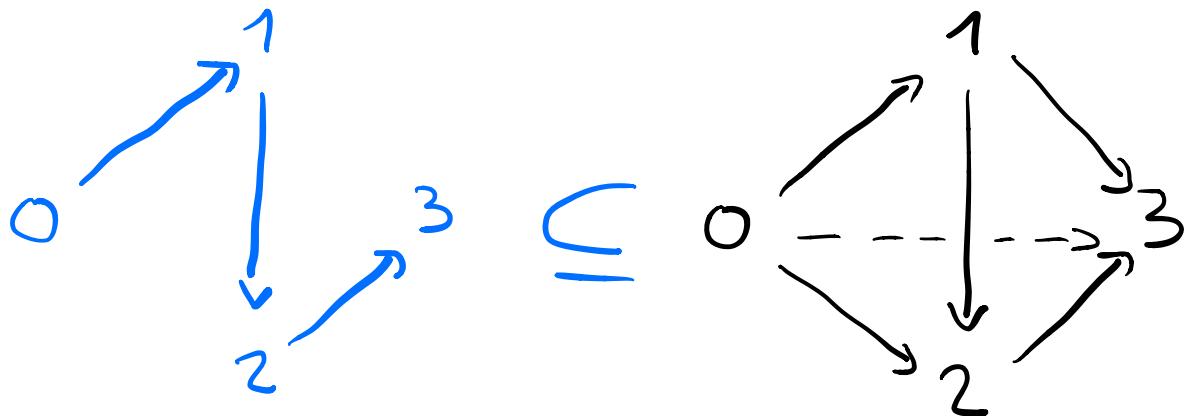
$\exists !$ isomorphism $\Delta^S \simeq \Delta^{|S|}$.

$\left\{ \begin{array}{l} \Delta^{\{i\}} \subseteq \Delta^n \text{ are vertices} \\ \Delta^{\{i < j\}} \subseteq \Delta^n \text{ are edges} \end{array} \right.$



* The spine $I^n \subseteq \Delta^n$ is

$$(I^n)_i = \left\{ \langle a_0 \dots a_i \rangle \in (\Delta^n)_i \mid a_i \leq a_0 + 1 \right\}$$



$$I^3 \subseteq \Delta^3$$

* The boundary, or

Simplicial n -sphere $\partial \Delta^n \subseteq \Delta^n$

is

$$(\partial \Delta^n)_i = \left\{ f \in (\Delta^n)_i \mid \text{Im}(f) \neq [n] \right\}$$

We have:

$$\partial \Delta^n = \bigcup_{j=0}^n \Delta^{[n] \setminus j}$$

The name comes from:

$$|\partial \Delta^n| \simeq S^n$$

* For $0 \leq k \leq n$, the k -th

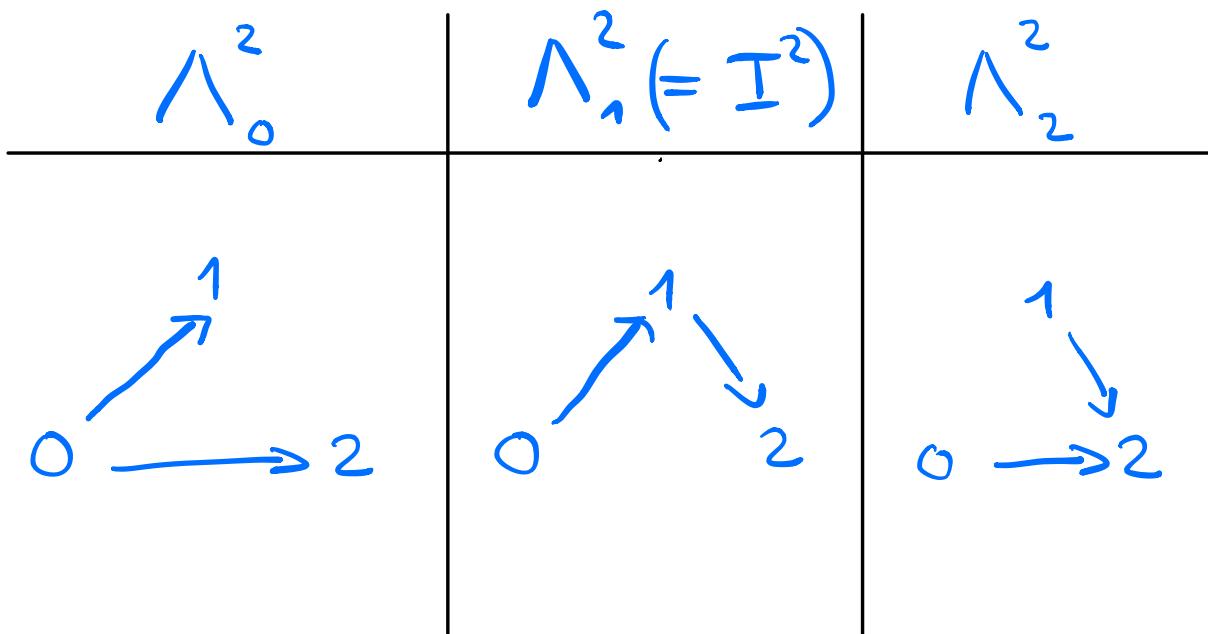
horn $\Lambda_k^n \subseteq \Delta^n$ is

$$(\Delta_k^n)_i = \left\{ g \in (\Delta_k^n) \mid i \notin \text{Im}(g) \right\}$$

So:

$$\Delta_k^n = \bigcup_{j \neq k} \Delta^{[n]} - k$$

$$|\Delta_k^n| \simeq D^n \text{ n-disk}$$



* Horns Λ_k^n with

- $0 < k < n$ are inner horns.
- $k = 0, n$ are outer horns.

Lemma 20

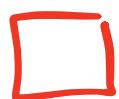
a) Inner horns contain spines:

$$0 < k < n \Rightarrow I^n \subseteq \Lambda_k^n$$

b) $\forall n \geq 3, \left\{ \begin{array}{l} I^n \subseteq \Lambda_0^n \\ I^n \subseteq \Lambda_n^n \end{array} \right.$

c) . $I^1 \notin \wedge_0^1, \wedge_1^1$

$I^2 \notin \wedge_0^2, \wedge_2^2$



4) Skeletal filtration

Presheaves are colimits of representables, but for simplicial sets we can be more precise.

Idea: Add first all the 0-simplices

then all the 1-simplices

...

def 21 Let $X_{\cdot} \in sSet$ and

$x \in X_n$. We say that x

is **degenerate** if $n > 0$ and

the following equivalent conditions

hold:

- * $x \in \text{Im}(s_i : X_{n-1} \rightarrow X_n)$ for some i .
- * x factors through Δ^m for some $m < n$:

$$x : \Delta^n \rightarrow \Delta^m \rightarrow X.$$

Otherwise we say that x is
non-degenerate. □

Notation

• $X_n = X_n^{\text{nd}} \amalg X_n^{\text{deg}}$



$X_{\cdot}^{\text{nd}}, X_{\cdot}^{\text{deg}}$ are not
simplicial subsets of X .

Example

$$(\Delta^n)_k^{\text{nd}} = \left\{ [k] \hookrightarrow [n] \right\}$$

correspond to the "true"

k -dimensional faces of $|\Delta^n|$.

Prop 22: (Eilenberg-Zilber lemma)

Let $X_* \in sSet$, $n \geq 0$ and

$x \in X_n$. Then $x: \Delta^n \rightarrow X$

can be factored uniquely as

$$\boxed{x: \Delta^n \xrightarrow{P^*} \Delta^m \xrightarrow{\tau} X_*}$$

with :

- $p: [m] \rightarrow [n]$ surjective
- τ non-degenerate m -simplex.

Proof:

Existence

Let $m \geq 0$ be minimal for
the existence of a factorisation

$$x: \Delta^n \xrightarrow{g^*} \Delta^m \xrightarrow{\tau} X.$$

Then :

- g^* is surjective (otherwise
we would have a factorisation

through $\text{Im}(g^*) \Rightarrow g$ is surjective.

- $\bar{\tau}$ is non-degenerate (otherwise we would have a factorisation

$$x: \Delta^n \rightarrow \Delta^m \rightarrow \Delta^{m'} \rightarrow X.$$

with $m' < m$).

Uniqueness:

$$\text{Let } x: \Delta^n \xrightarrow{(g')^*} \Delta^{m'} \xrightarrow{\bar{\tau}'} X.$$

be another such factorisation.

$$\text{Write } \alpha = g^*, \alpha' = (g')^*.$$

* α, α' surjective



g, g' surjective



g, g' admit sections



α, α' admit sections β, β' .

* We get

$$\tau = \tau \circ \alpha \circ \beta$$

$$= \alpha \circ \beta$$

$$= \tau' \circ \alpha' \circ \beta.$$

* Because τ is non-deg,

$$\alpha' \circ \beta \text{ inj} \Rightarrow m \leq m'.$$

By symmetry, $m = m'$.

$$\Rightarrow \alpha' \circ \beta = \text{id}_{[m]}$$

$$\Rightarrow \begin{cases} \tau = \tau' \circ \alpha' \circ \beta = \tau' \\ \alpha = \alpha' \circ \beta \circ \alpha = \alpha' \end{cases}$$



Def 23 (Skleta)

$X_+ \in sSet$, $k \geq -1$.

$$Sk_k(X_n) := \left\{ \begin{array}{l} x \in X_n, \exists \text{ fact } \circ \\ \Delta^n \rightarrow \Delta^m \rightarrow X_+ \\ \text{with } m \leq k \end{array} \right\}$$

$\text{Sk}_k(X_\cdot) \subseteq X_\cdot$ is a simplicial subset of X_\cdot ,

the k -th skeleton of X_\cdot .



By construction:

$$\left\{ \begin{array}{l} \emptyset = \text{Sk}_{-1}(X_\cdot) \subseteq \text{Sk}_0(X_\cdot) \subseteq \dots \\ \bigcup_{k \geq -1} \text{Sk}_k(X_\cdot) = X_\cdot \end{array} \right.$$

and $X_n^{\text{nd}} \cap \text{Sk}_k(X_\cdot) = \begin{cases} \emptyset, & k < n \\ X_n^{\text{nd}}, & k \geq n \end{cases}$

Rmk Sh_R induces a functor

$$\text{Sh}_R(-) : \text{sSet} \longrightarrow \text{sSet}$$

with interesting properties

(see Exercise Sheet 2).

Prop 24: Let $X \in \text{sSet}$, $k \geq 0$.

There is a pushout square

$$\begin{array}{ccc} \coprod_{X_R^{\text{nd}}} \partial \Delta^k & \longrightarrow & \coprod_{X_R^{\text{nd}}} \Delta^k \\ \downarrow & & \downarrow \\ \text{Sh}_{R-1}(X.) & \longrightarrow & \text{Sh}_R(X.) \end{array}$$

More generally, for any $A_+ \subseteq X_+$ subcomplex, there is a pushout square

$$\begin{array}{ccc} \coprod_{X_R^{nd}} & \xrightarrow{\partial \Delta^k} & \coprod_{X_R^{nd}} \\ X_R^{nd} - A_R^{nd} & & X_R^{nd} - A_R^{nd} \\ \downarrow & & \downarrow \\ & & \end{array}$$

$$A_+ \cup Sh_{h \rightarrow}^-(X_+) \longrightarrow A_+ \cup Sh_h^-(X_+)$$

Proof: Let's do the particular

case $A_+ = \emptyset$.

We have $x \in X_R^{nd} \Rightarrow x \in Sh_h^-(X_+) R$.

and the faces of $x \in X_R^{\text{nd}}$ are
 in $\text{Sk}_{R-1}(X) \Rightarrow$ we have the
 commutative square of the statement.

We observe that we have

$$\left(\coprod_{X_R^{\text{nd}}} \Delta^R \right)_n \rightarrow \left(\coprod_{X_R^{\text{nd}}} \partial \Delta^R \right)_n$$

IS

$$\left\{ g^* x \mid x \in X_R^{\text{nd}}, g: [n] \rightarrow [R] \right\}$$

IS (Eilenberg-Zilber)

$$\text{Sk}_R(X_*)_n \rightarrow \text{Sk}_{R-1}(X_*)_n .$$

- Moreover, the square is clearly a pullback.
- It remains to show

Lemma If $\begin{array}{ccc} A & \hookrightarrow & B \\ \downarrow & & \downarrow \\ C & \hookrightarrow & D \end{array}$ in \mathbf{Set}

Satisfies: $\forall n \geq 0, B_n \setminus A_n \xrightarrow{\sim} D_n \setminus C_n$

then it is a pushout.

which reduces to the same statement in \mathbf{Set} (since (ω)limits are computed objectwise) and is then an exercise. □

Cor 25 Let $X \in \text{SSet}$. Then

the geometric realisation is

a CW-complex, whose k -cells
are in bijection with X_k^{nd} .

If $A \subseteq X$ is a simplicial
subset, then $|A| \subseteq |X|$ is
a CW-subcomplex.

Proof: $|-|$ is a left adjoint

\Rightarrow commutes with colimits.

$$\text{Hence: } |X| = \bigcup_{k \geq -1} |\text{Sk}_k(X)|$$

and we have a pushout

$$\begin{array}{ccc}
 \coprod_{X_h^{\text{nd}}} S^R & \longrightarrow & \coprod_{X_h^{\text{nd}}} D^R \\
 \downarrow & & \downarrow \Gamma \\
 |\mathrm{Sh}_{h\rightarrow}^k(X.)| & \longrightarrow & |\mathrm{Sh}_h^k(X.)|
 \end{array}$$



Rmk This shows that, to model homotopy types, one could forget about degeneracies and work with $\mathrm{PSh}(\Delta^{\text{inj}})$.
 Not so for our purpose !

Examples

$$\begin{cases} \text{Sk}_0(I^n) = \coprod \Delta^{\{i\}} \\ \text{Sk}_i(I^n) = I^n \text{ for all } i \geq 1 \end{cases}$$

and $(I^n)_1^{\text{nd}} = \left\{ \Delta^{\{i, i+1\}} \mid 0 \leq i \leq n-1 \right\}$

We deduce :

$$I^n = \Delta^{\{0,1\}} \coprod_{\Delta^{\{1\}}} \Delta^{\{1,2\}} \coprod \dots \Delta^{\{n-1, n\}}$$

- Similar arguments show:

$$\partial \Delta^n = \underline{\coprod}_{\Delta^{[n] \setminus \{i,j\}}} \Delta^{[n] \setminus i}$$

$$\textstyle \bigwedge_R^n = \underline{\coprod}_{\substack{\Delta^{[n] \setminus \{i,j\}} \\ i \neq R}} \Delta^{[n] \setminus i}$$

5) Kan complexes

Simplicial sets model homotopy

types via geometric realisation.

But if one tries to develop

Homotopy theory directly in

SSet with $[0,1] \rightsquigarrow \Delta^1$,

things do not work very well:

Rmk: In general, the

relation on X_0 defined by

$$x \sim y \Leftrightarrow \exists \Delta^1 \xrightarrow{h} X,$$

$$d^0(h) = x, d^1(h) = y$$

is neither symmetric ($x = \Delta^1$)
nor transitive ($x = I^2$)

def 26 Let $X_* \in sSet$.

$$\pi_0(X_*) := X_*/\sim$$

with \sim the equivalence

relation generated by \sim . □

Rmk $\pi_0(X_*) \cong \text{Colim}(X_*; \Delta^{\text{op}} \rightarrow \text{Set})$

This is already unsatisfactory
and things get worse for
 π_n for $n \geq 1$.

This is the topic of
simplicial Homotopy theory
and reads to:

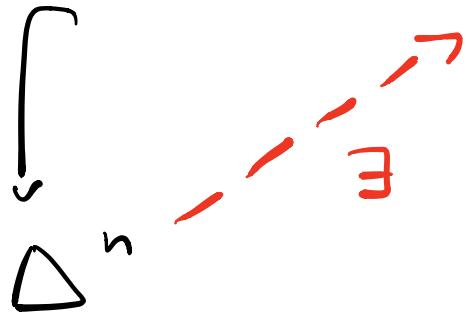
def 27 $X_{\cdot} \in sSet$ is a
Kan complex or Kan
simplicial set if it has
the Kan lifting property:

$\forall n \geq 1, \forall 0 \leq k \leq n,$

$$\begin{array}{ccc} \Delta_k^n & \longrightarrow & X_{\cdot} \\ \downarrow & \nearrow \exists & \nearrow \\ \Delta^n & & \end{array}$$

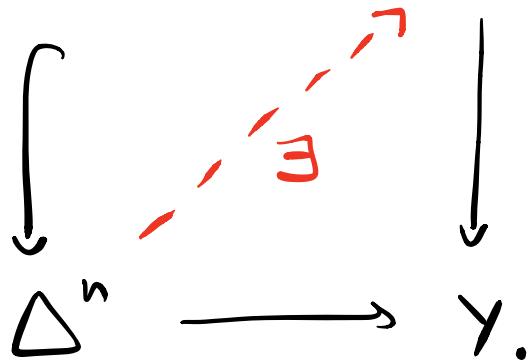
X_\cdot is a contractible Kan complex if $\forall n \geq 1$,

$$\partial \Delta^n \longrightarrow X_\cdot$$



More generally, $X_\cdot \rightarrow Y_\cdot$ is a Kan fibration if

$$\Lambda^n_R \longrightarrow X_\cdot$$



and a trivial Kan fibration

if $\partial\Delta^n \rightarrow X.$

$$\begin{array}{ccc} & \nearrow \exists & \\ \downarrow & & \downarrow \\ \Delta^n & \longrightarrow & Y. \end{array}$$

Rmk It is not easy to really motivate these defs without going into simplicial homotopy theory.

The basic idea is :

“For a Kan complex $X.$, the

Homotopical properties of $|X|$

can be expressed purely

simplicially. \Rightarrow

Prop 28 Let $A \in \text{Top}$. Then

the singular simplicial set

$S(A)$ is a Kan complex.

Proof: By adjunction, we have

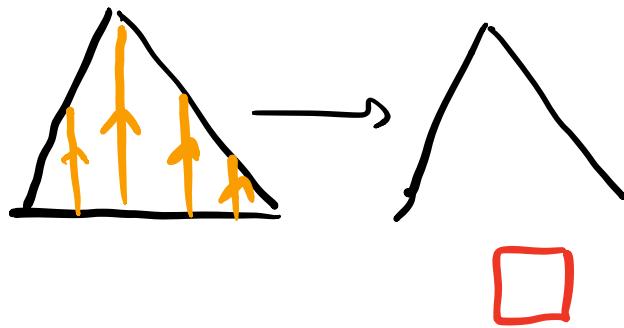
to show: $|\Delta_k^n| \longrightarrow A$

$$\begin{matrix} & & \nearrow \\ \downarrow & & \swarrow \\ |\Delta^n| & \xrightarrow{\quad} & \exists \end{matrix}$$

But $|\Delta_k^n| \hookrightarrow |\Delta^n|$ admits

a continuous retraction.

In pictures:



def 29 The Homotopy category

of Kan complexes $R\text{Kan}$

Has : - objects = Kan complexes
- morphisms = Δ^1 -Homotopy classes
 → of morphisms

(need to show composition is well-defined)

- A morphism $f: X_+ \rightarrow Y_+$ in $s\text{Set}$ is a weak homotopy eq.

if $|g|: |X.| \rightarrow |Y.|$ is an
homotopy equivalence of CW-
complexes. □

The main result of s. homotopy theory is:

Thm We have a diagram

$$\begin{array}{ccc} h\text{-Kan} & \xrightarrow{| \cdot |} & h\text{-CW} \\ \downarrow s & & \downarrow s \\ s\text{-Set}[\text{w.h.eq}] & \xrightarrow{| \cdot |} & \text{Top}[\text{w.h.eq}] \end{array}$$

Remark This still does not
explain why horns appear!

Lemma 30: Monomorphisms of

simplicial sets are

“generated” by the inclusions

$$\left\{ \partial \Delta^n \hookrightarrow \Delta^n \mid n \in \mathbb{N} \right\}$$

under • pushouts

• transfinite composition

Proof Follows from the

existence of the skeletal

filtration and

Proposition 24.



The fundamental reason that
Rrons appear in the definition

of Kan complexes is the analogous result:

Prop 31: Monomorphisms of simplicial sets which are also weak homotopy equivalences are “generated by” the horn inclusions

$$\left\{ \Lambda_k^n \hookrightarrow \Delta^n \mid \begin{array}{l} n \in \mathbb{N} \\ 0 \leq k \leq n \end{array} \right\}$$

under

- pushouts

- retracts

- transfinite composition.



4) Nerves of categories

Our model of ∞ -categories

will be a particular type of

simplicial sets. We also

want to be able to consider

ordinary categories as

∞ -categories.

→ We need a fully

faithful functor

$$N: \text{Cat} \longleftrightarrow s\text{Set} .$$

- Posets give rise to categories:

$$\text{Poset} \xrightarrow{\text{Fully faithful}} \text{Cat}$$

$$P \longrightarrow \text{Ob} : x \in P$$

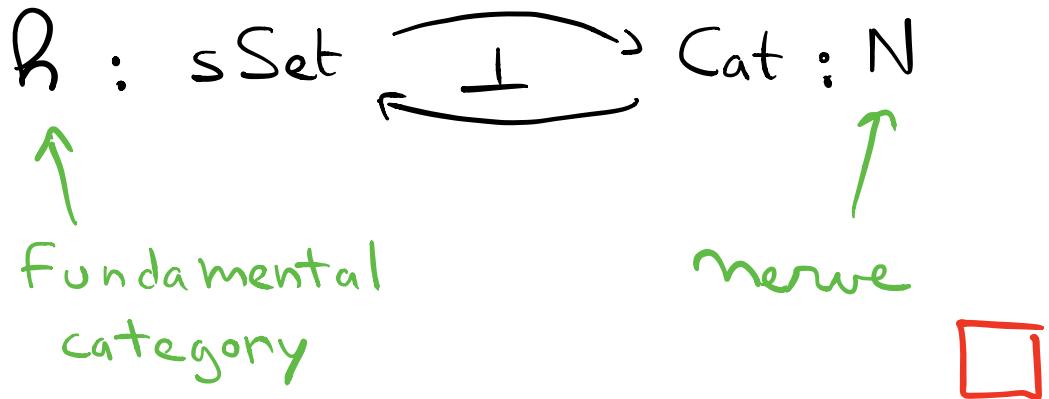
$$P(x, y) = \begin{cases} *, & x \leq y \\ \emptyset, & \text{otherwise} \end{cases}$$

In particular:

$$\Delta \xrightarrow{L} \text{Cat}$$

def 32 By the free cocompletion

property, L induces an adjunction:



By construct°:

$$\begin{cases} R(\Delta^n) = [n] \\ R \text{ commutes with colimits} \end{cases}$$

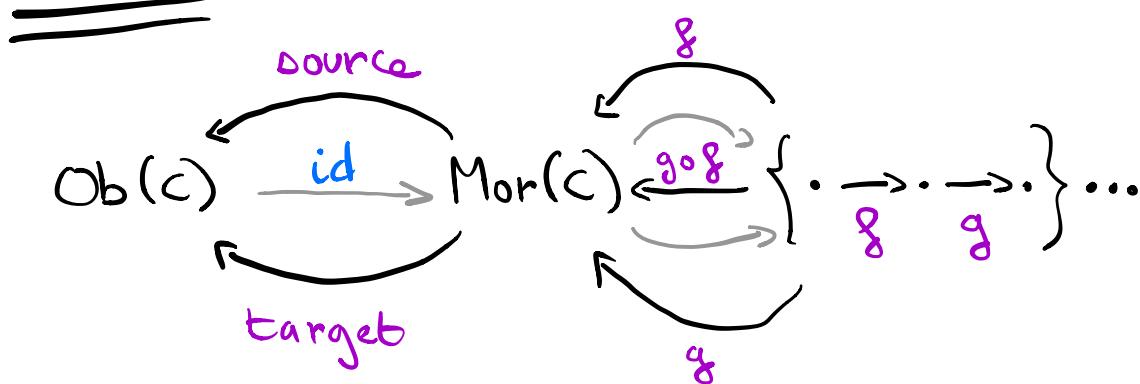
and

$$N(C)_n = \text{Cat}([n], C)$$

$$= \left\{ \begin{array}{l} n\text{-tuples of} \\ \text{composable morphisms} \end{array} \right\}$$

Rmk We have $N([n]) \simeq \Delta^n$
 $(\Leftarrow \text{ fully faithful})$.

$N(C)$:



This makes it clear that

We can reconstruct C from $N(C)$.

Prop 33 N is fully faithful

proof: For C, D in Cat , we must show

$$\text{Cat}(C, D) \xrightarrow{\sim} \text{sSet}(N_C, N_D).$$

Injectivity: A functor is determined by its effect on objects and morphisms. Since $(Nc)_0 = \text{Ob}(c)$ and $(Nc)_1 = \text{Mor}(c)$, we are done.

Surjectivity Let $\alpha: Nc \rightarrow Nd$.

We define a candidate F for the preimage functor using again

$$\begin{cases} (Nc)_0 = \text{Ob}(c) \\ (Nc)_1 = \text{Mor}(c) \end{cases}$$

To check that F is a functor, we need to see:

• F compatible with source/target:

Let $g : c \rightarrow c'$ in C

$$\left\{ \begin{array}{l} s F(g) \stackrel{\text{def}}{=} d^\circ \alpha(g) = \alpha d^\circ(g) \stackrel{\text{def}}{=} F(c) \\ t F(g) \stackrel{\text{def}}{=} d'^\circ \alpha(g) = \alpha d'^\circ(g) \stackrel{\text{def}}{=} F(c') \end{array} \right.$$

• F compatible with identities

$$F(id_c) \stackrel{\text{def}}{=} g s^\circ(c) = s^\circ g(c) \stackrel{\text{def}}{=} id_{F(c)}$$

• F compatible with compositions

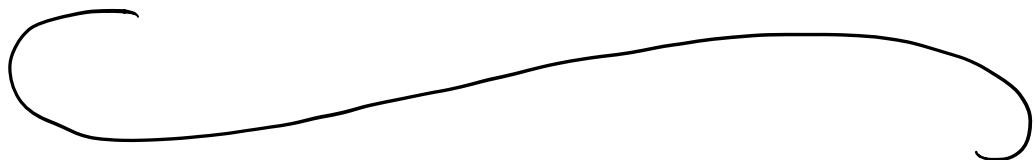
$$F(g \circ f) = \alpha(g \circ f) = \alpha(d'(g, f))$$

$$= d'^\circ \alpha(g, f) = F(g) \circ F(f)$$

$\rightsquigarrow F$ is a functor.

By construction, $N(F)$ and α coincide on 0 and 1-simplices, and it is easy to see that it forces them to be equal

(because simplices in $(NC)_i$ for $i > 2$ are uniquely determined by their 1-simplices).



Thm 35 : (characterisation
of nerves)

Let $X_\cdot \in \text{sSet}$. TFAE:

(1) $X_\cdot \simeq NC$ for some $C \in \text{Cat}$.

(2) X_\cdot satisfies the unique

Spine extension property :

For all $n \geq 2$,

$$\begin{array}{ccc} I^n & \longrightarrow & X_\cdot \\ \downarrow & \nearrow \exists! & \\ \Delta^n & & \end{array}$$

(3) X_+ satisfies the inner
horn unique extension property:

For all $n \geq 2$ and $0 < k < n$:

$$\begin{array}{ccc} \Delta^n_k & \longrightarrow & X_+ \\ \downarrow & \nearrow \text{---}^{\exists!} & \nearrow \text{---}^{\exists!} \\ \Delta^n & & \end{array}$$

Proof: