



Performance Modeling of Computer Systems and Networks

Prof. Vittoria de Nitto Personè

Generating Discrete Random Variates

Università degli studi di Roma Tor Vergata
Department of Civil Engineering and Computer Science Engineering

Copyright © Vittoria de Nitto Personè, 2021
<https://creativecommons.org/licenses/by-nc-nd/4.0/>



1

Prerequisite

We assume the knowledge of discrete random variables (sect.6.1).

In particular:

- $Equilakely(a,b)$
- $Geometric(p)$
- $Bernoulli(p)$
- $Binomial(n,p)$
- $Pascal(n,p)$
- $Poisson(\mu)$

Prof. Vittoria de Nitto Personè

2

2

sis2 è inventory system, usa equilikely.

sis2.c

```
#include <stdio.h>
#include "rng.h"

#define MINIMUM 20
#define MAXIMUM 80
#define STOP 100 /* 100 weeks = about 2 years*/
#define sqr(x) ((x) * (x))

long Equilikely(long a, long b)
{ return (a + (long) ((b - a + 1) * Random())); }

long GetDemand(void)
{
    return (Equilikely(10, 50)); }

```

3

Noi generiamo un valore tra 0 e 1, e lo trasformiamo a seconda della variabile.

ssq2.c distribution-driven simulation

```
#include <stdio.h>
#include <math.h>
#include "rng.h"
#define LAST 10000L /* number of jobs processed */
#define START 0.0

double Exponential(double m) /* -----*
                             m > 0.0
                             ----- */
{ return (-m * log(1.0 - Random())); }

double Uniform(double a, double b) /* -----*
                                   a < b
                                   * -----*/
{ return (a + (b - a) * Random()); }

double GetArrival(void)
{ static double arrival = START;
  arrival += Exponential(2.0);
  return (arrival); }

double GetService(void)
{ return (Uniform(1.0, 2.0)); }

```

Prof. Vittoria de Nitto Personè

4

4

X è v.a., F è la cumulativa (funzione di distribuzione).
 Esiste F^* che sarebbe "funzione inversa", ma formalmente non lo è. Perché?
 Perché se voglio passare dal continuo al discreto, molti punti continui potrebbero convergere nello stesso valore discreto.

Discrete Simulation
Generating Discrete Random Variates

Preliminary Definitions

X random variable, $F(\cdot)$ is the cdf of X

The inverse distribution function (idf) of X is the function

$$F^* : (0, 1) \rightarrow \mathcal{X}, \forall u \in (0, 1)$$

$$F^*(u) = \min_x \{x : u < F(x)\}$$

that is, if $F^*(u) = x$, x is the smallest possible value of X for which $F(x)$ is greater than u

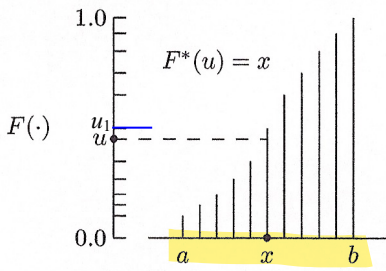
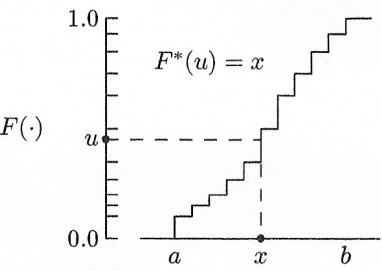
Prof. Vittoria de Nitto Personè

5

Sull'asse 'x' ho insieme discreto. Prendo un x , faccio $F(x) = u_1 > u$. Ora faccio l'inverso, $F^*(u) = "x \text{ più piccola tale che } F(x) > u"$. Ciò è vero se prendo sulle ascisse proprio 'x', perchè $F(x) = u_1 > u$.
 In pratica, quando faccio l'inverso di un certo "u", devo prendere la x sulle ascisse più piccola tale che, $F(x) < u$ minore stretto.

Discrete Simulation
Generating Discrete Random Variates

- $\mathcal{X} = \{a, a+1, \dots, b\}$, where b may be ∞ , $F(\cdot)$ is the cdf of X ,
- $F(x) = \text{Prob}\{X \leq x\} = u_1 > u \quad F^*(u) = \min_x \{x : u < F(x)\}$

Theorem

- if $u < F(a)$, $F^*(u) = a$
- else $F^*(u) = x$ where $x \in \mathcal{X}$ is the unique possible value of X for which $F(x-1) \leq u < F(x)$

Prof. Vittoria de Nitto Personè

6

Il teorema ci dice che, presa un'ascissa "a" (che non è per forza il primo punto sull'asse x).
 Se $F(a) > u$, allora $F^*(u) = a$.
 Altrimenti prendo un'ascissa tale che $F(x-1) \leq u < F(x)$

Algorithm 1

```
x = a;
while (F(x) <= u)
    x++;
return x;          /* x is F*(u) */
```

Average case analysis:

- let Y be the number of while loop passes
- $Y = X - a$
- $E[Y] = E[X - a] = E[X] - a = \mu - a$

Linear search algorithm!

Prof. Vittoria de Nitto Personè

7

7

qui partiamo dalla "moda", muovendoci
con ricerca binaria.

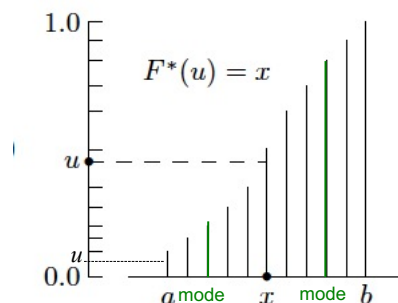
Idea: start at a more likely point

For $\mathcal{X} = \{a, a+1, \dots, b\}$, a more efficient linear search
algorithm defines $F^*(u)$

Algorithm 2

```
x = mode;          /* initialize with the mode of X */
if (F(x) <= u)
    while (F(x) <= u)
        x++;
else if (F(a) <= u)
    while (F(x-1) > u)
        x--;
else
    x = a;
return x;          /* x is F*(u) */
```

For large \mathcal{X} , consider binary search



Prof. Vittoria de Nitto Personè

8

8

alcuni esempi di inversi calcolate.

Idf Examples

- In some cases $F^*(u)$ can be determined explicitly
- If X is *Bernoulli*(p) and $F(x) = u$,
then $x=0$ iff $0 < u < 1-p$

$$F^*(u) = \begin{cases} 0 & 0 < u < 1-p \\ 1 & 1-p \leq u < 1 \end{cases}$$

Prof. Vittoria de Nitto Personè

9

9

5/03/2023

Random Variate Generation By Inversion

- X is a discrete random variable with idf $F^*(\cdot)$
- continuous random variable U is *Uniform*(0,1)
- Z is the discrete random variable defined by $Z = F^*(U)$

Theorem

Z and X are identically distributed

this Theorem allows any discrete random variable (with known idf) to be generated with one call to `Random()`

Algorithm 3

```
u = Random();  
return F*(u);
```

Prof. Vittoria de Nitto Personè

10

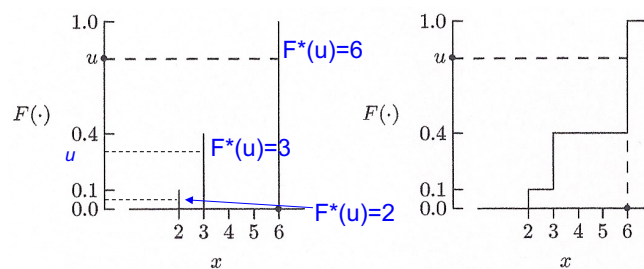
10

Inversion Examples

- Consider X with pdf

$$f(x) = \begin{cases} 0.1 & x = 2 \\ 0.3 & x = 3 \\ 0.6 & x = 6 \end{cases}$$

- The cdf for X is plotted using two formats



Prof. Vittoria de Nitto Personè

11

11

```

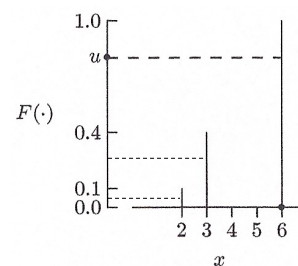
if (u < 0.1)
  return 2;
else if (u < 0.4)
  return 3;
else
  return 6;

```

This algorithm returns

2 with probability 0.1,
3 with probability 0.3 and
6 with probability 0.6.

This corresponds to the pdf of X .



more efficiency: check the ranges for u associated with $x = 6$
first (the mode), then $x = 3$, then $x = 2$

- problems may arise when $|X|$ is large or infinite

Prof. Vittoria de Nitto Personè

12

12

More inversion examples

Generating a *Bernoulli*(p) random variate

```
u = Random();
if (u < 1-p)
    return 0;
else
    return 1;
```

Generating an *Equilikely*(a,b) random variate

```
u = Random();
return a + (long) (u * (b - a + 1));
```

Library rvgs

- Includes 6 discrete random variate generators (as below) and 7 continuous random variate generators
 - `long Bernoulli(double p)`
 - `long Binomial(long n, double p)`
 - `long Equilikely(long a, long b)`
 - `long Geometric(double p)`
 - `long Pascal(long n, double p)`
 - `long Poisson(double μ)`
- Functions Bernoulli, Equilikely, Geometric use inversion; essentially ideal
- Functions Binomial, Pascal, Poisson do not use inversion

Library `rvms`

- Provides accurate pdf, cdf, idf functions for many random variates
- Idfs can be used to generate random variates by inversion
- Functions `idfBinomial`, `idfPascal`, `idfPoisson` may have high marginal execution times
- Not recommended when many observations are needed due to time inefficiency
- Array of cdf values with inversion may be preferred

Prof. Vittoria de Nitto Personè

15

15

Truncation

Sometimes, the realistic values of a variable are restricted to a subset

X random variable with possible values $\mathcal{X}=\{0, 1, 2, \dots\}$ and cdf $F(x)=\Pr(X \leq x)$

- want to restrict X to the finite range $0 \leq a \leq x \leq b < \infty$

- if $a > 0$, $\alpha = \Pr(X < a)$, $\beta = \Pr(X > b)$

$$\alpha = \Pr(X < a) = \Pr(X \leq a-1) = F(a-1)$$

$$\beta = \Pr(X > b) = 1 - \Pr(X \leq b) = 1 - F(b)$$

$$\Pr(a \leq X \leq b) = \Pr(X \leq b) - \Pr(X < a) = F(b) - F(a-1)$$

essentially, always true iff $F(b) \cong 1.0$ and $F(a-1) \cong 0.0$

Prof. Vittoria de Nitto Personè

16

16

Specifying truncation points

- if a and b are specified

Left-tail, right-tail probabilities α and β obtained using cdf

$$\alpha = \Pr(X < a) = F(a-1) \quad \text{and} \quad \beta = \Pr(X > b) = 1 - F(b)$$

transformation is exact

- if α and β are specified

idf can be used to obtain a and b

$$a = F^*(\alpha) \quad \text{and} \quad b = F^*(1 - \beta)$$

transformation is not exact because X is discrete

$$\Pr(X < a) \leq \alpha \quad \text{and} \quad \Pr(X > b) < \beta$$

Prof. Vittoria de Nitto Personè

17

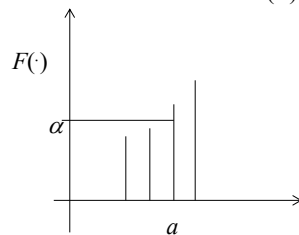
17

$$F(x-1) \leq u < F(x)$$

Specifying truncation points

- if α and β are specified

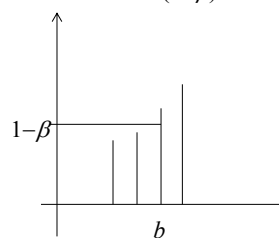
$$a = F^*(\alpha)$$



$$F(a) > \alpha$$

$$\Pr(X < a) \leq \alpha$$

$$b = F^*(1 - \beta)$$



$$F(b) > 1 - \beta$$

$$\Pr(X \leq b) > 1 - \beta$$

$$- \Pr(X \leq b) < \beta - 1$$

$$1 - \Pr(X \leq b) < \beta$$

$$\Pr(X > b) < \beta$$

Prof. Vittoria de Nitto Personè

18

18

Effects of truncation

sometimes truncation is insignificant:
truncated and un-truncated random variables have (essentially)
the same distribution

Truncation is useful for efficiency:

- When idf is complex, inversion requires cdf search
- cdf values are typically stored in an array
- Small range gives improved space/time efficiency

Truncation is useful for realism:

- Prevents arbitrarily large values possible from some variates

In some applications, truncation is significant

- Produces a new random variable
- Must be done correctly !