

# Spanner: Google's Globally-Distributed Database

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representing a host of authors

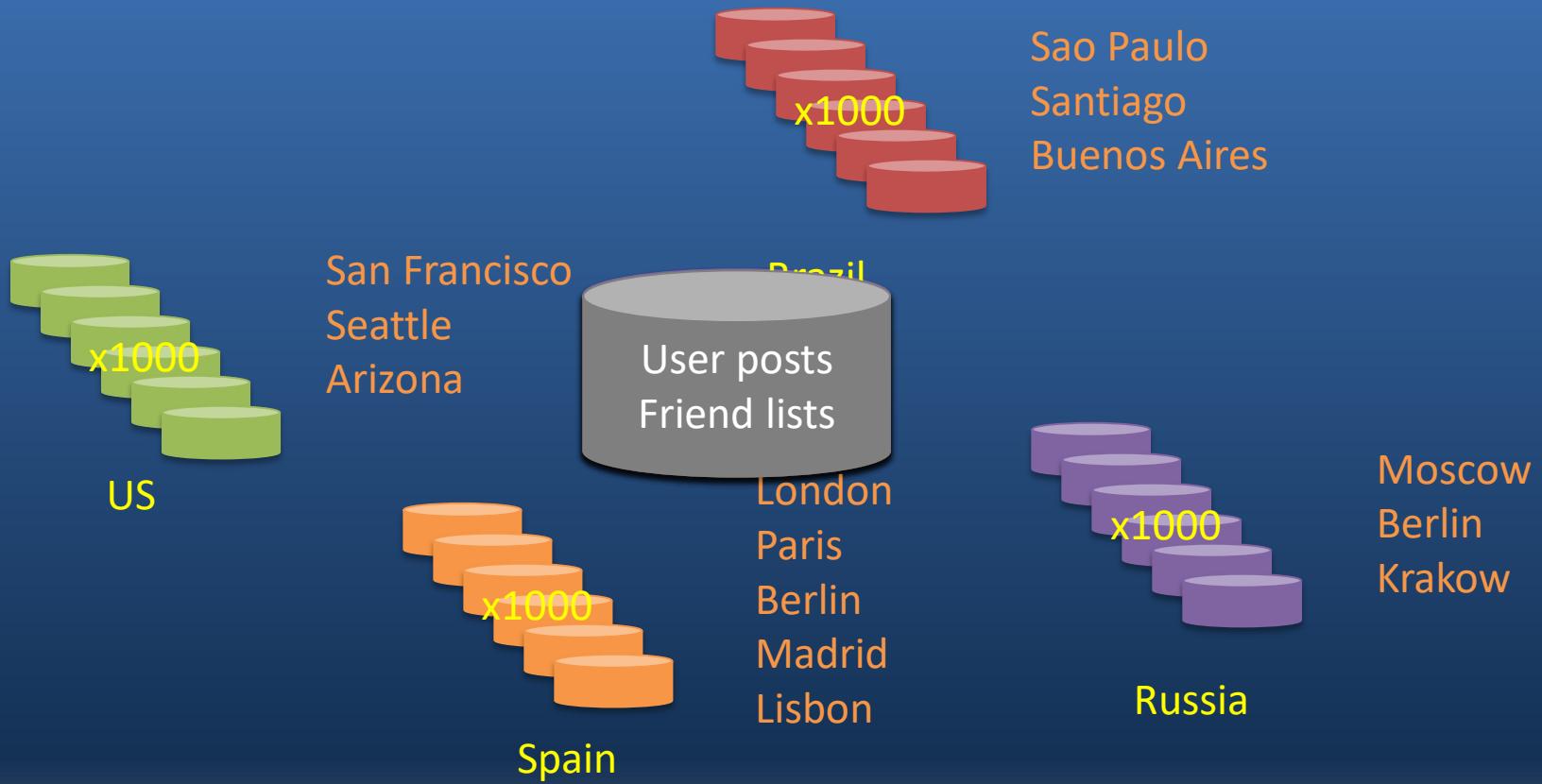
OSDI 2012



# What is Spanner?

- Distributed multiversion database
  - General-purpose transactions (ACID)
  - SQL query language
  - Schematized tables
  - Semi-relational data model
- Running in production
  - Storage for Google's ad data
  - Replaced a sharded MySQL database

# Example: Social Network



# Overview

- Feature: Lock-free distributed read transactions
- Property: External consistency of distributed transactions
  - First system at global scale
- Implementation: Integration of concurrency control, replication, and 2PC
  - Correctness and performance
- Enabling technology: TrueTime
  - Interval-based global time

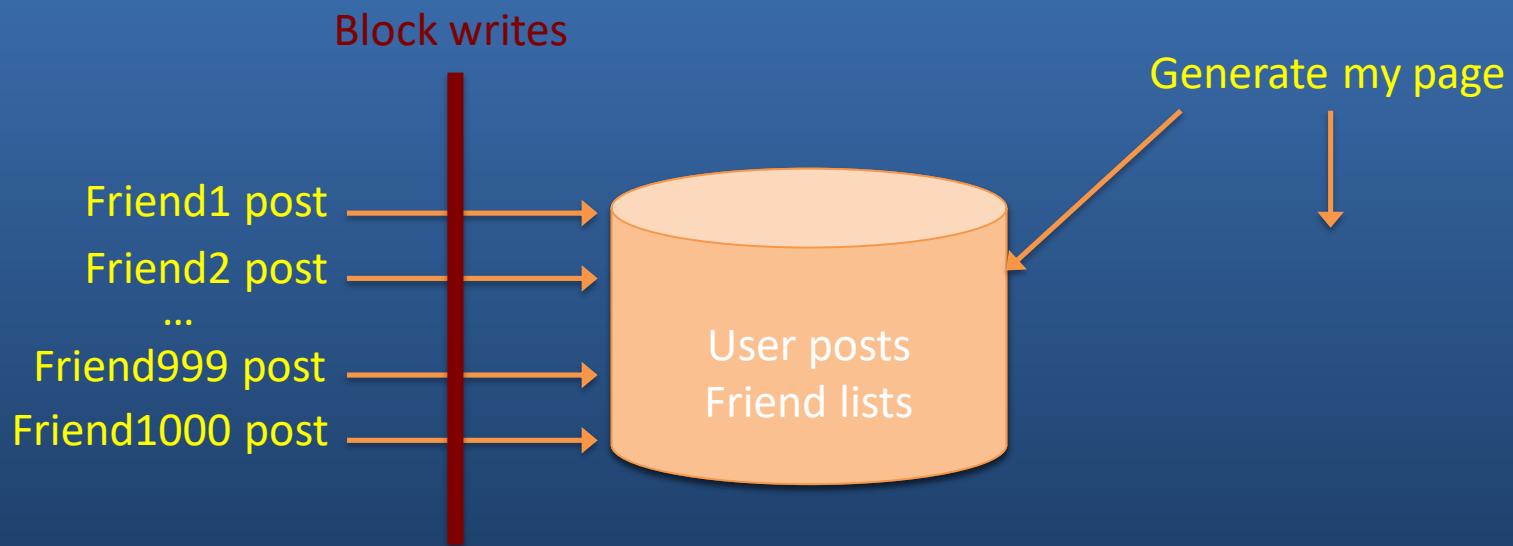
# Read Transactions

- Generate a page of friends' recent posts
  - Consistent view of friend list and their posts

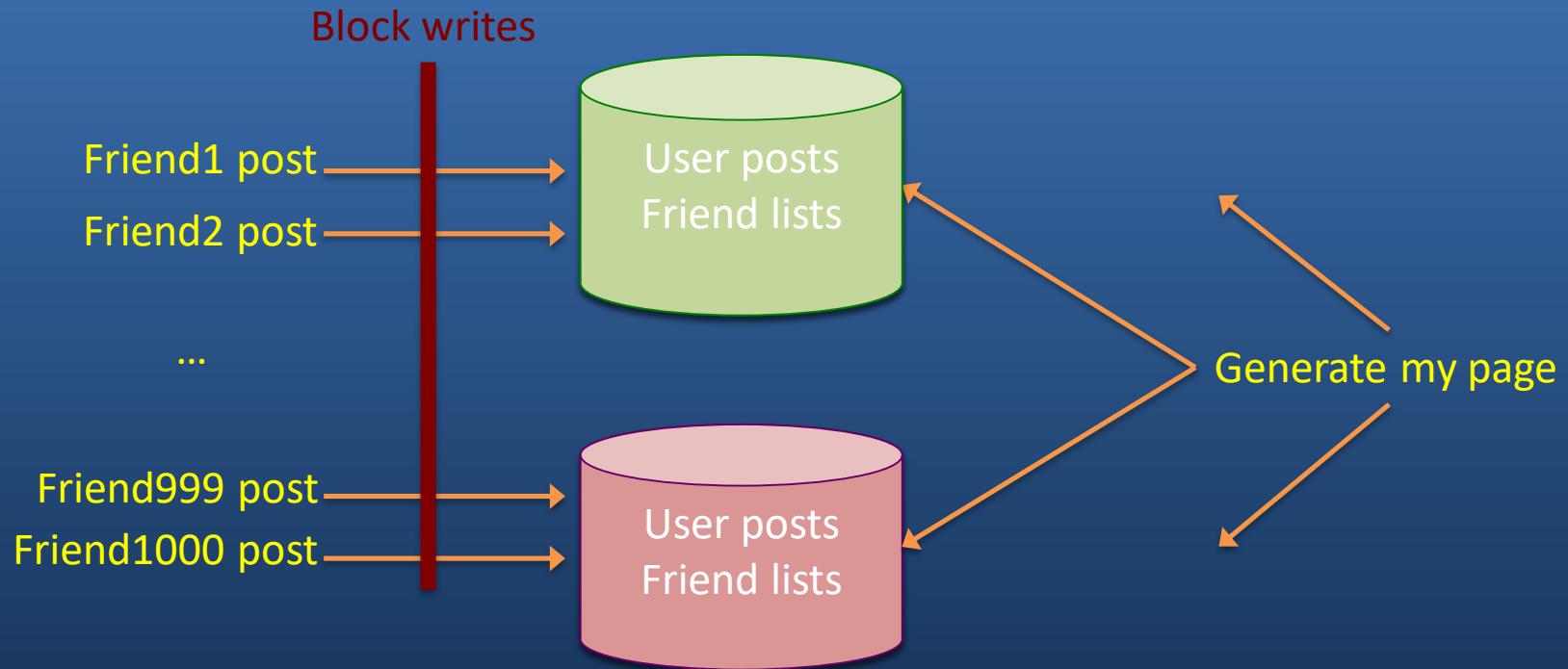
Why consistency matters

1. Remove untrustworthy person X as friend
2. Post P: “My government is repressive...”

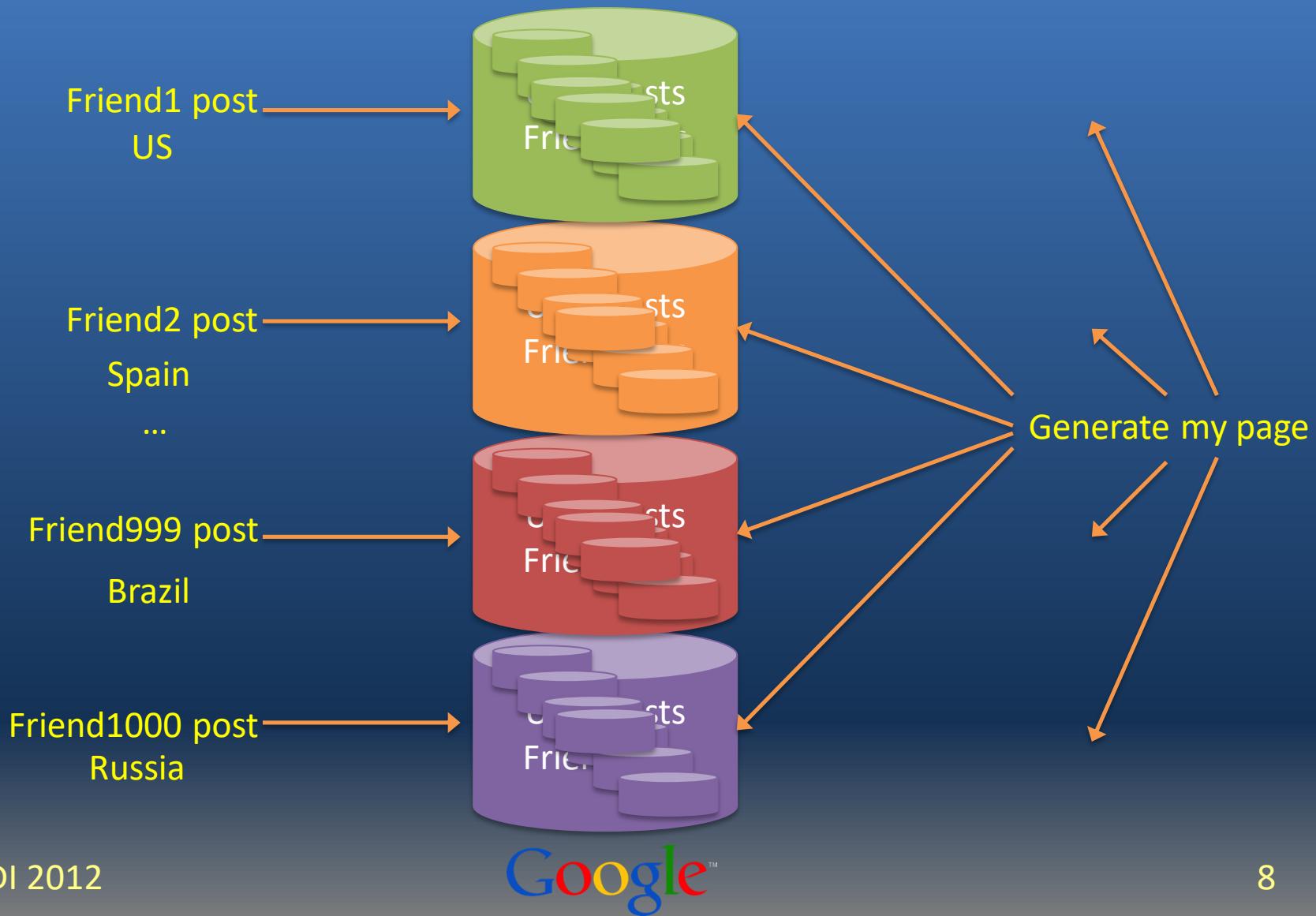
# Single Machine



# Multiple Machines



# Multiple Datacenters



# Version Management

- Transactions that write use strict 2PL
  - Each transaction  $T$  is assigned a timestamp  $s$
  - Data written by  $T$  is timestamped with  $s$

Time	<8	8	15
My friends	[X]	[]	
My posts			[P]
X's friends	[me]	[]	

# Synchronizing Snapshots

Global wall-clock time

==

External Consistency:

Commit order respects global wall-time order

==

Timestamp order respects global wall-time order

given

timestamp order == commit order

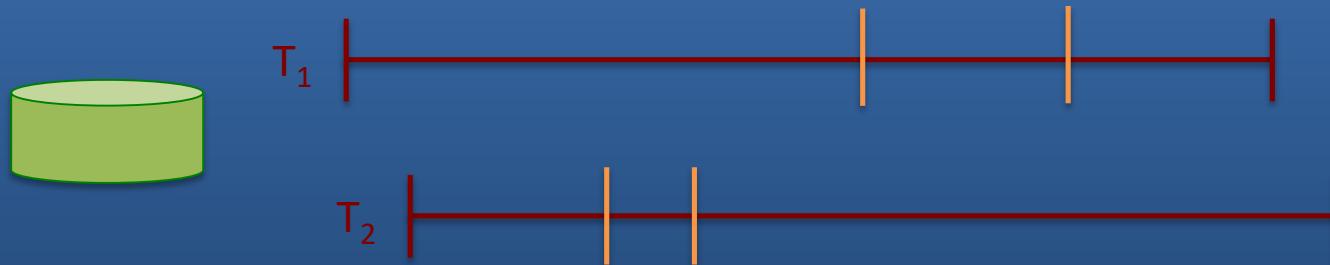
# Timestamps, Global Clock

- Strict two-phase locking for write transactions
- Assign timestamp while locks are held



# Timestamp Invariants

- Timestamp order == commit order

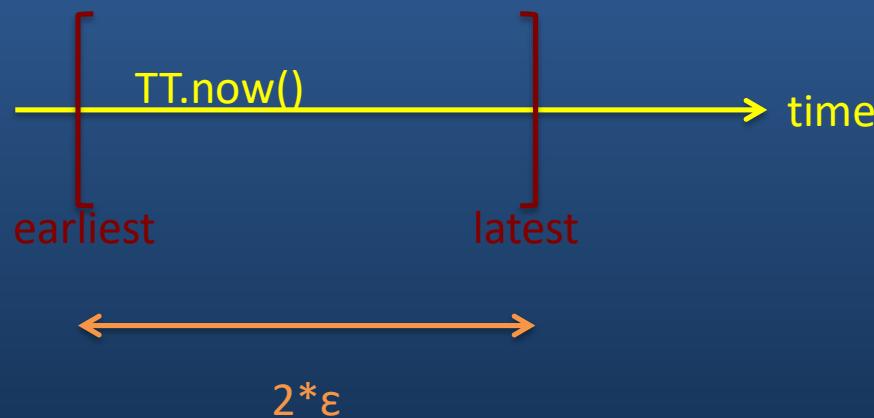


- Timestamp order respects global wall-time order

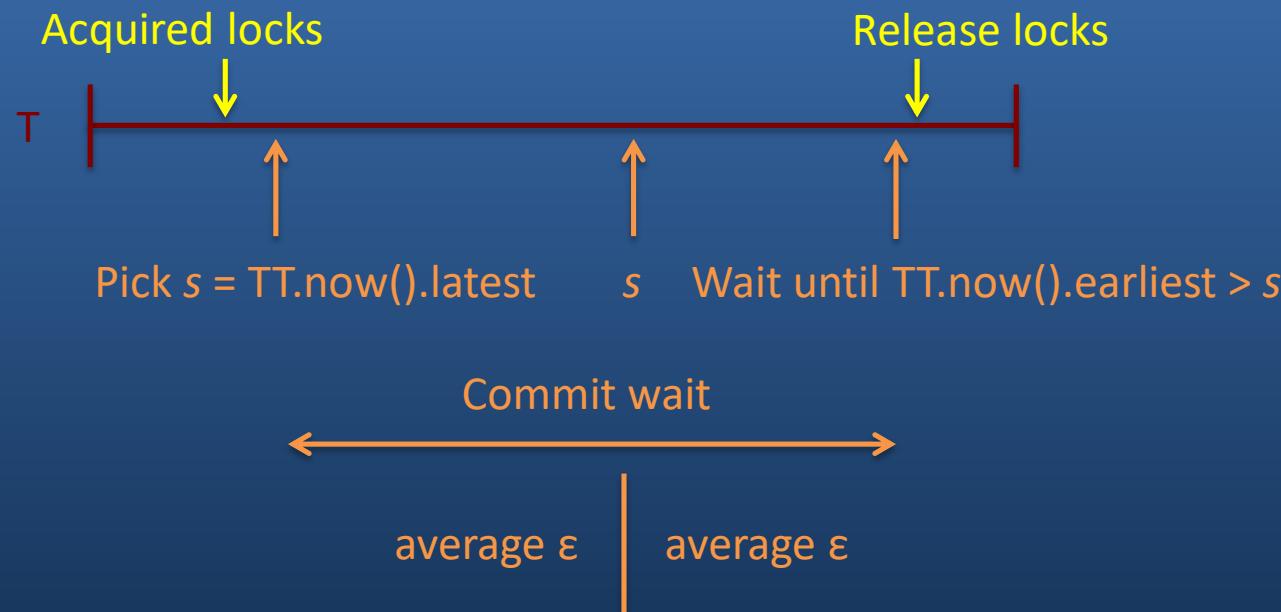


# TrueTime

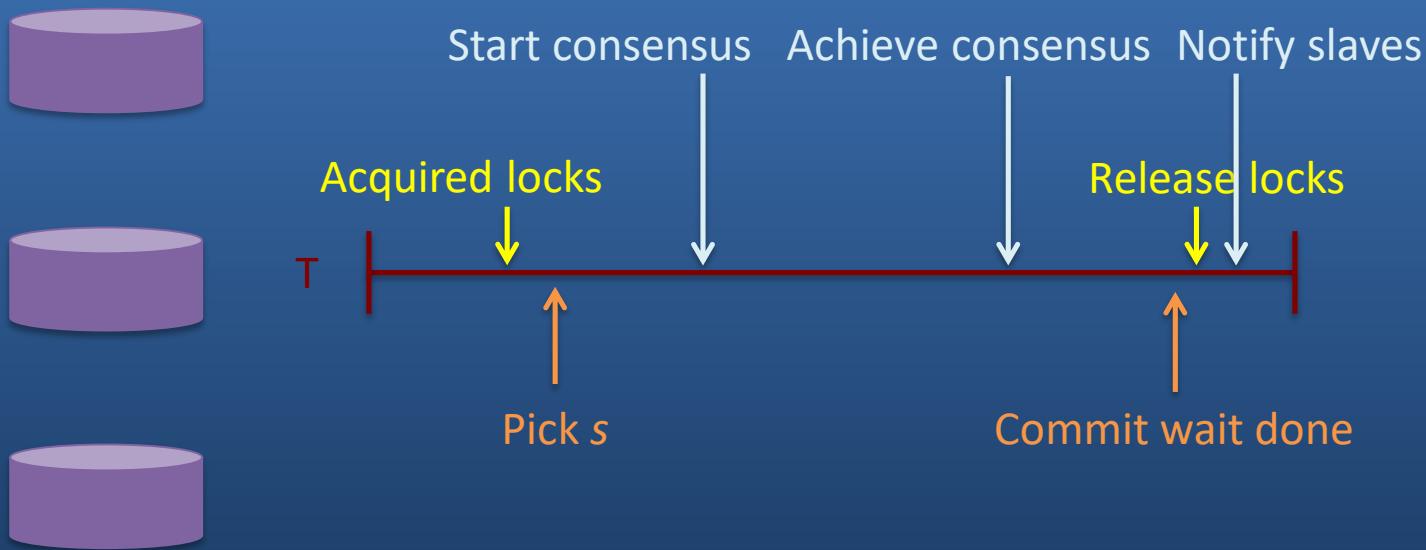
- “Global wall-clock time” with bounded uncertainty



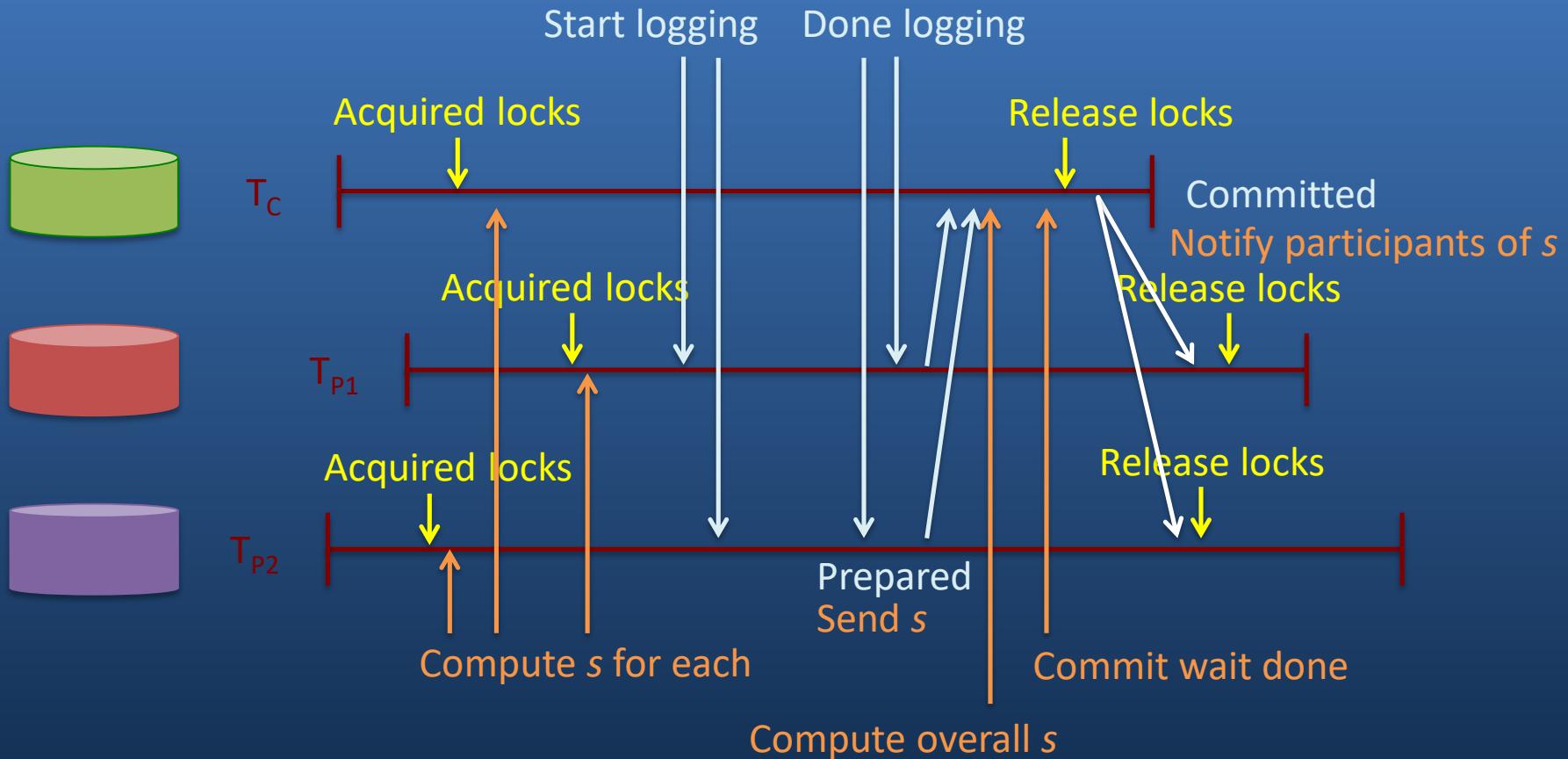
# Timestamps and TrueTime



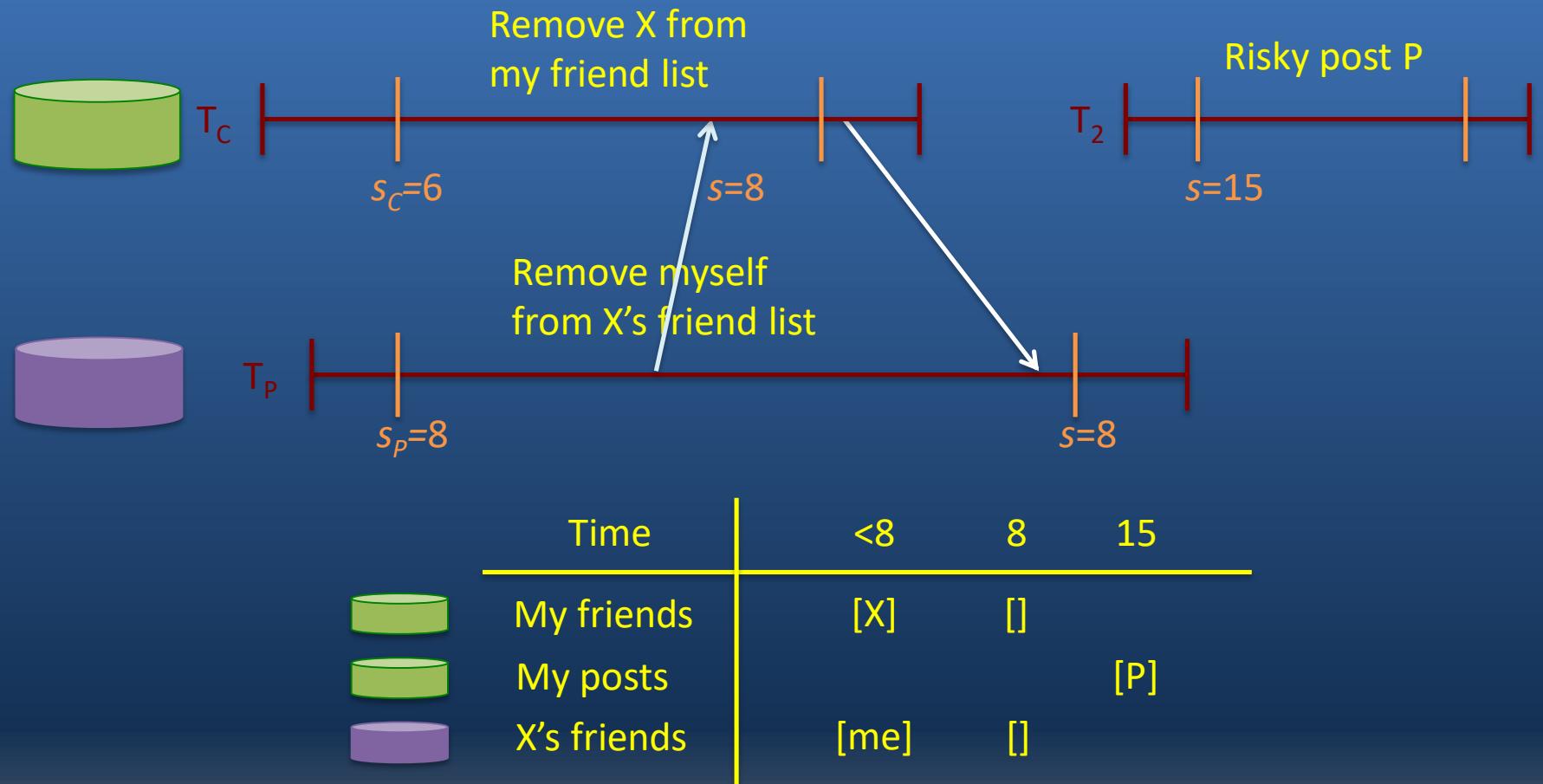
# Commit Wait and Replication



# Commit Wait and 2-Phase Commit



# Example



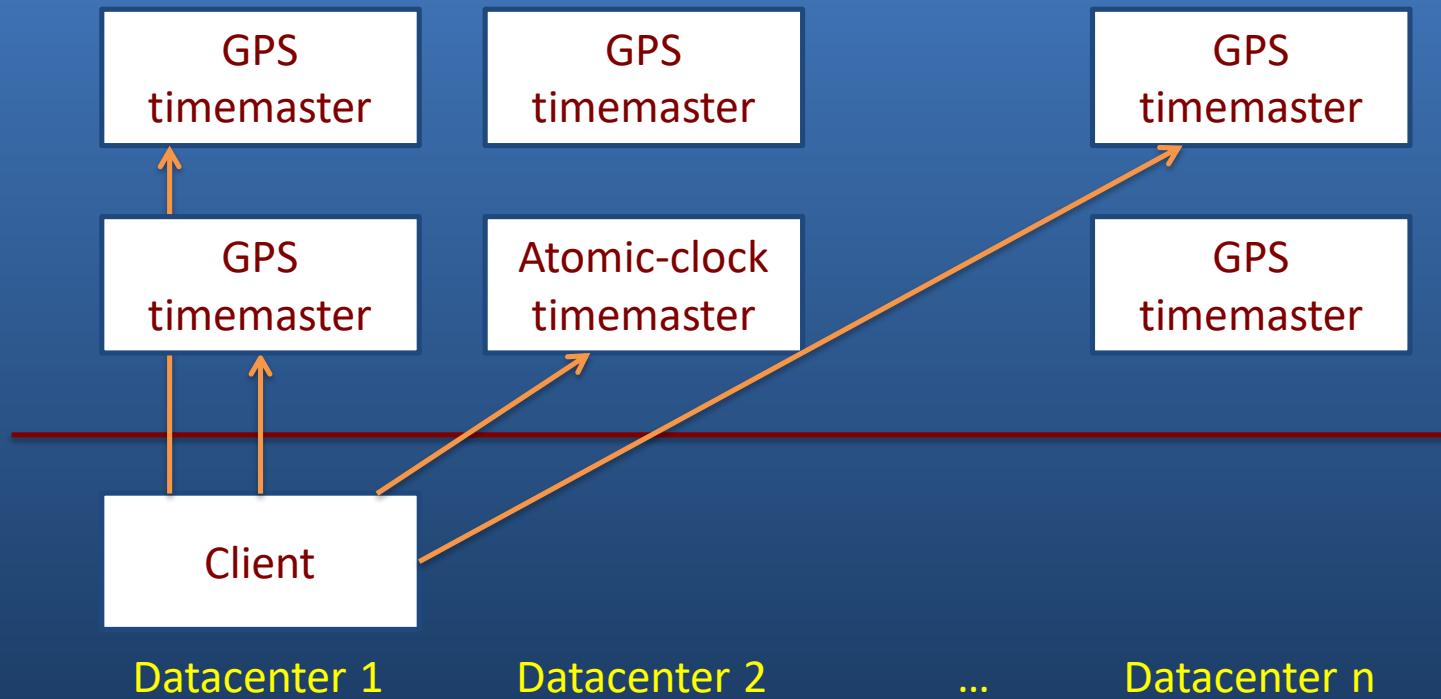
# What Have We Covered?

- Lock-free read transactions across datacenters
- External consistency
- Timestamp assignment
- TrueTime
  - Uncertainty in time can be waited out

# What Haven't We Covered?

- How to read at the present time
- Atomic schema changes
  - Mostly non-blocking
  - Commit in the future
- Non-blocking reads in the past
  - At any sufficiently up-to-date replica

# TrueTime Architecture

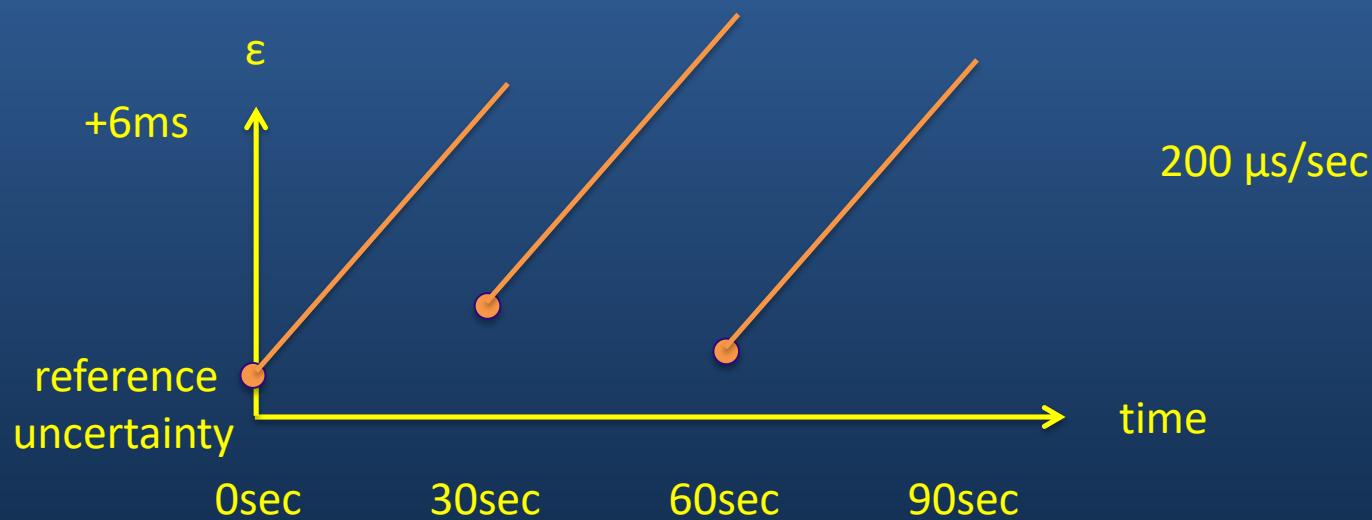


Compute reference [earliest, latest] = now  $\pm \varepsilon$

# TrueTime implementation

$\text{now} = \text{reference now} + \text{local-clock offset}$

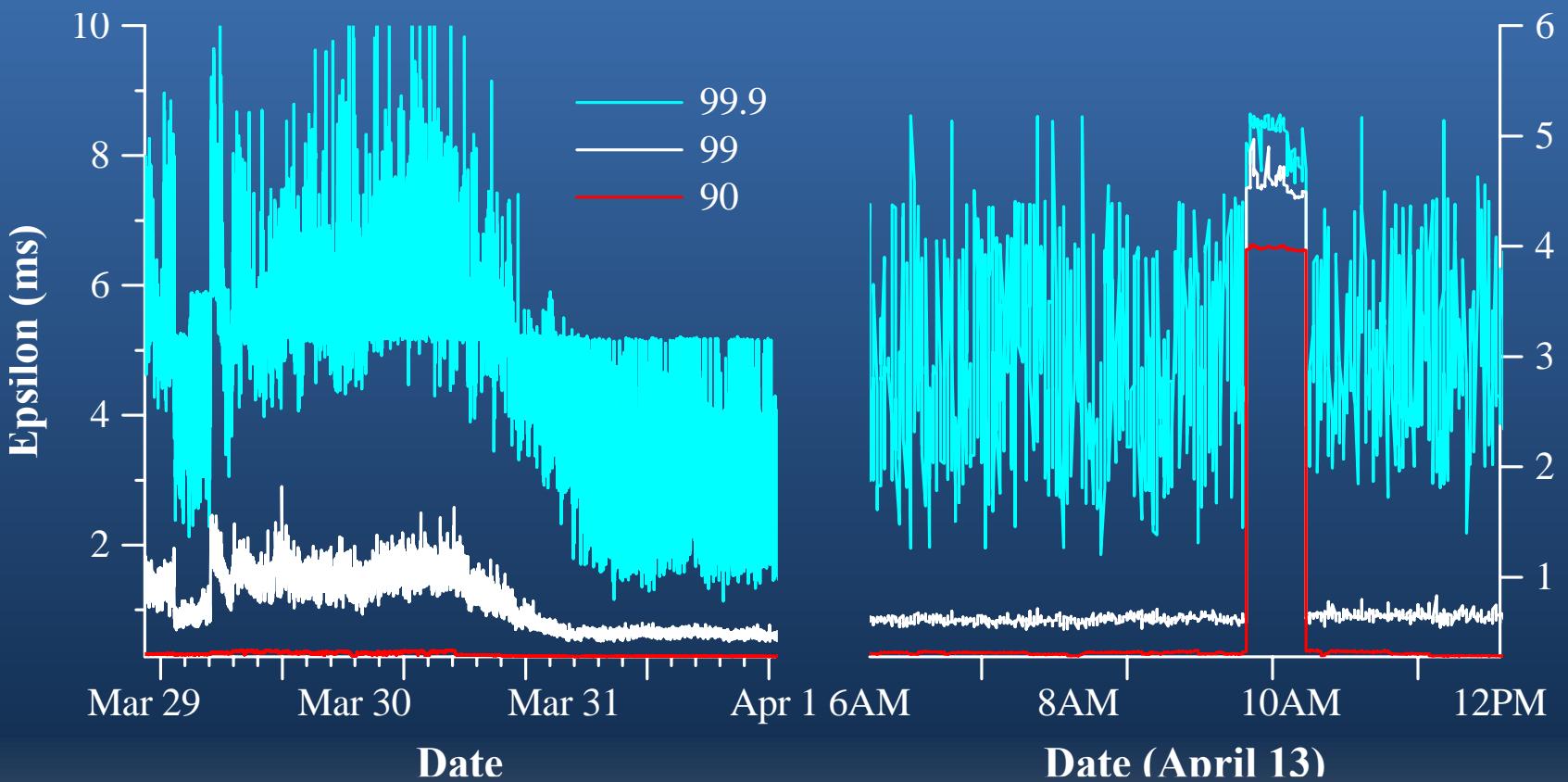
$\varepsilon = \text{reference } \varepsilon + \text{worst-case local-clock drift}$



# What If a Clock Goes Rogue?

- Timestamp assignment would violate external consistency
- Empirically unlikely based on 1 year of data
  - Bad CPUs 6 times more likely than bad clocks

# Network-Induced Uncertainty



# What's in the Literature

- External consistency/linearizability
- Distributed databases
- Concurrency control
- Replication
- Time (NTP, Marzullo)

# Future Work

- Improving TrueTime
  - Lower  $\varepsilon < 1$  ms
- Building out database features
  - Finish implementing basic features
  - Efficiently support rich query patterns

# Conclusions

- Reify clock uncertainty in time APIs
  - Known unknowns are better than unknown unknowns
  - Rethink algorithms to make use of uncertainty
- Stronger semantics are achievable
  - Greater scale  $\neq$  weaker semantics

# Thanks

- To the Spanner team and customers
- To our shepherd and reviewers
- To lots of Googlers for feedback
- To you for listening!
- Questions?