

Soft Inference and Posterior Marginals

September 19, 2013

Soft vs. Hard Inference

- Hard inference
 - “Give me a single solution”
 - Viterbi algorithm
 - Maximum spanning tree (Chu-Liu-Edmonds alg.)
- Soft inference
 - Task 1: Compute a distribution over outputs
 - Task 2: Compute functions on distribution
 - marginal probabilities, expected values, entropies, divergences

Why Soft Inference?

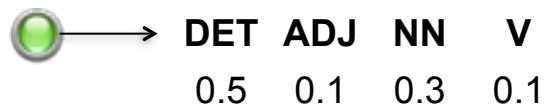
- Useful applications of posterior distributions
 - **Entropy**: how confused is the model?
 - **Entropy**: how confused is the model of its prediction at time i ?
 - **Expectations**
 - What is the expected number of words in a translation of this sentence?
 - What is the expected number of times a word ending in -ed was tagged as something other than a verb?
 - **Posterior marginals**: given some input, how likely is it that some (*latent*) event of interest happened?

String Marginals

- Inference question for HMMs
 - What is the probability of a string w ?
Answer: generate all possible tag sequences and explicitly *marginalize*

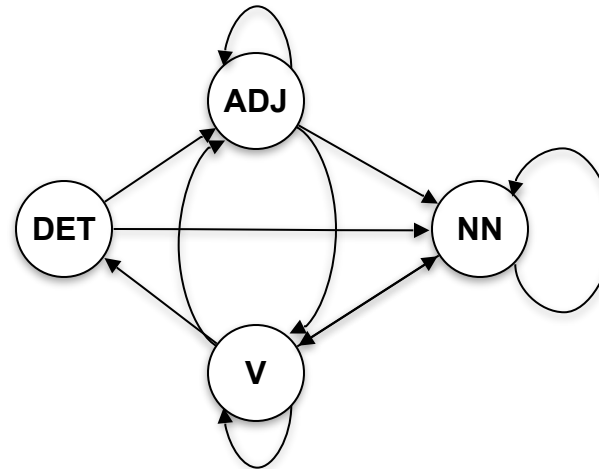
$$O(|\Omega|^{|w|}) \text{ time}$$

Initial Probabilities:



η Transition Probabilities:

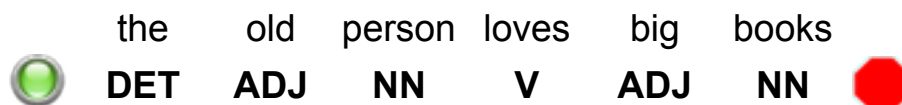
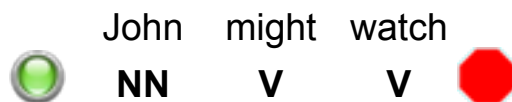
	DET	ADJ	NN	V
DET	0.0	0.0	0.0	0.5
ADJ	0.3	0.2	0.1	0.1
NN	0.7	0.7	0.3	0.2
V	0.0	0.1	0.4	0.1
Red Circle	0.0	0.0	0.2	0.1



γ Emission Probabilities:

	DET		ADJ		NN		V	
the	0.7	green	0.1	book	0.3	might	0.2	
a	0.3	big	0.4	plants	0.2	watch	0.3	
		old	0.4	people	0.2	watches	0.2	
		might	0.1	person	0.1	loves	0.1	
				John	0.1	reads	0.19	
				watch	0.1	books	0.01	

Examples:



John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$
DET	DET	DET	0.0	ADJ	DET	DET	0.0	NN	DET	DET	0.0	V	DET	DET	0.0
DET	DET	ADJ	0.0	ADJ	DET	ADJ	0.0	NN	DET	ADJ	0.0	V	DET	ADJ	0.0
DET	DET	NN	0.0	ADJ	DET	NN	0.0	NN	DET	NN	0.0	V	DET	NN	0.0
DET	DET	V	0.0	ADJ	DET	V	0.0	NN	DET	V	0.0	V	DET	V	0.0
DET	ADJ	DET	0.0	ADJ	ADJ	DET	0.0	NN	ADJ	DET	0.0	V	ADJ	DET	0.0
DET	ADJ	ADJ	0.0	ADJ	ADJ	ADJ	0.0	NN	ADJ	ADJ	0.0	V	ADJ	ADJ	0.0
DET	ADJ	NN	0.0	ADJ	ADJ	NN	0.0	NN	ADJ	NN	0.0000042	V	ADJ	NN	0.0
DET	ADJ	V	0.0	ADJ	ADJ	V	0.0	NN	ADJ	V	0.0000009	V	ADJ	V	0.0
DET	NN	DET	0.0	ADJ	NN	DET	0.0	NN	NN	DET	0.0	V	NN	DET	0.0
DET	NN	ADJ	0.0	ADJ	NN	ADJ	0.0	NN	NN	ADJ	0.0	V	NN	ADJ	0.0
DET	NN	NN	0.0	ADJ	NN	NN	0.0	NN	NN	NN	0.0	V	NN	NN	0.0
DET	NN	V	0.0	ADJ	NN	V	0.0	NN	NN	V	0.0	V	NN	V	0.0
DET	V	DET	0.0	ADJ	V	DET	0.0	NN	V	DET	0.0	V	V	DET	0.0
DET	V	ADJ	0.0	ADJ	V	ADJ	0.0	NN	V	ADJ	0.0	V	V	ADJ	0.0
DET	V	NN	0.0	ADJ	V	NN	0.0	NN	V	NN	0.0000096	V	V	NN	0.0
DET	V	V	0.0	ADJ	V	V	0.0	NN	V	V	0.0000072	V	V	V	0.0

John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$	John	migh	watc	$\Pr(x, y)$
DET	DET	DET	0.0	ADJ	DET	DET	0.0	NN	DET	DET	0.0	V	DET	DET	0.0
DET	DET	ADJ	0.0	ADJ	DET	ADJ	0.0	NN	DET	ADJ	0.0	V	DET	ADJ	0.0
DET	DET	NN	0.0	ADJ	DET	NN	0.0	NN	DET	NN	0.0	V	DET	NN	0.0
DET	DET	V	0.0	ADJ	DET	V	0.0	NN	DET	V	0.0	V	DET	V	0.0
DET	ADJ	DET	0.0	ADJ	ADJ	DET	0.0	NN	ADJ	DET	0.0	V	ADJ	DET	0.0
DET	ADJ	ADJ	0.0	ADJ	ADJ	ADJ	0.0	NN	ADJ	ADJ	0.0	V	ADJ	ADJ	0.0
DET	ADJ	NN	0.0	ADJ	ADJ	NN	0.0	NN	ADJ	NN	0.0000042	V	ADJ	NN	0.0
DET	ADJ	V	0.0	ADJ	ADJ	V	0.0	NN	ADJ	V	0.0000009	V	ADJ	V	0.0
DET	NN	DET	0.0	ADJ	NN	DET	0.0	NN	NN	DET	0.0	V	NN	DET	0.0
DET	NN	ADJ	0.0	ADJ	NN	ADJ	0.0	NN	NN	ADJ	0.0	V	NN	ADJ	0.0
DET	NN	NN	0.0	ADJ	NN	NN	0.0	NN	NN	NN	0.0	V	NN	NN	0.0
DET	NN	V	0.0	ADJ	NN	V	0.0	NN	NN	V	0.0	V	NN	V	0.0
DET	V	DET	0.0	ADJ	V	DET	0.0	NN	V	DET	0.0	V	V	DET	0.0
DET	V	ADJ	0.0	ADJ	V	ADJ	0.0	NN	V	ADJ	0.0	V	V	ADJ	0.0
DET	V	NN	0.0	ADJ	V	NN	0.0	NN	V	NN	0.0000096	V	V	NN	0.0
DET	V	V	0.0	ADJ	V	V	0.0	NN	V	V	0.0000072	V	V	V	0.0

$$p = 0.0000219$$

Weighted Logic Programming

- Slightly different notation than the textbook, but you will see it in the literature
- WLP is useful here because it lets us **build hypergraphs**

$$\frac{I_1 : w_1 \quad I_2 : w_2 \quad \cdots \quad I_k : w_k}{I : w} \quad \phi$$

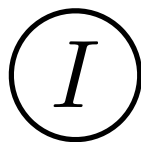
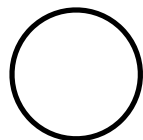
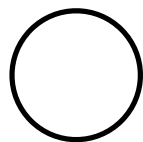
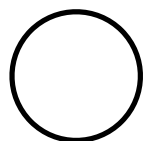
Weighted Logic Programming

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- WLP is useful here because it lets us **build hypergraphs**

$$\begin{array}{c}
 \frac{I_1 : w_1 \quad I_2 : w_2 \quad \cdots \quad I_k : w_k}{I : w} \quad \phi \\
 \bigcirc \\
 \vdots
 \end{array}$$

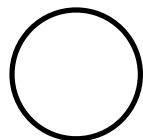
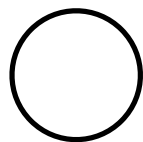
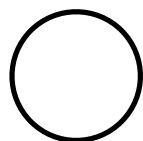
Hypergraphs

A	B	C
<hr/>		
	I	

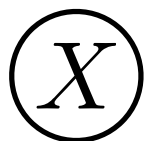
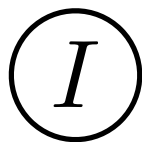


Hypergraphs

$$\frac{A \quad B \quad C}{I}$$



$$\frac{X \quad Y}{I}$$

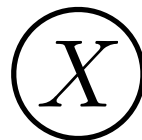
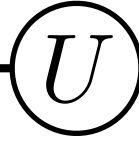
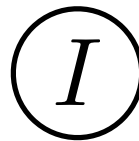
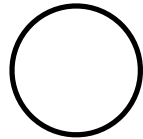
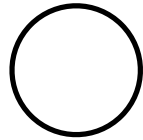
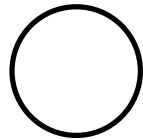


Hypergraphs

$$\frac{A \quad B \quad C}{I}$$

$$\frac{X \quad Y}{I}$$

$$\frac{U}{\bar{I}}$$



Viterbi Algorithm

Item form

$$[q, i]$$

Viterbi Algorithm

Item form

$$[q, i]$$

Axioms

$$[\text{START}, 0] : 1$$

Viterbi Algorithm

Item form

$$[q, i]$$

Axioms

$$[\text{START}, 0] : 1$$

Goals

$$[\text{STOP}, |\mathbf{x}| + 1]$$

Viterbi Algorithm

Item form

$$[q, i]$$

Axioms

$$\overline{[\text{START}, 0] : 1}$$

Goals

$$[\text{STOP}, |\mathbf{x}| + 1]$$

Inference rules

$$\frac{[q, i] : w}{[r, i + 1] : w \otimes \eta(q \rightarrow r) \otimes \gamma(r \downarrow x_{i+1})}$$

Viterbi Algorithm

Item form

$$[q, i]$$

Axioms

$$\overline{[\text{START}, 0] : 1}$$

Goals

$$[\text{STOP}, |\mathbf{x}| + 1]$$

Inference rules

$$\frac{[q, i] : w}{[r, i + 1] : w \otimes \eta(q \rightarrow r) \otimes \gamma(r \downarrow x_{i+1})}$$

$$\frac{[q, |\mathbf{x}|] : w}{[\text{STOP}, |\mathbf{x}| + 1] : w \otimes \eta(q \rightarrow \text{STOP})}$$

Viterbi Algorithm

w=(John, might, watch) Goal: [STOP, 4]

String Marginals

- Inference question for HMMs
 - What is the probability of a string \mathbf{w} ?
Answer: generate all possible tag sequences and explicitly *marginalize*

$$O(|\Omega|^{|\mathbf{w}|}) \text{ time}$$

Answer: use the **forward algorithm**

$$O(|\Omega|^2 \times |\mathbf{w}|) \text{ time}$$

$$O(|\Omega|) \text{ space}$$

Forward Algorithm

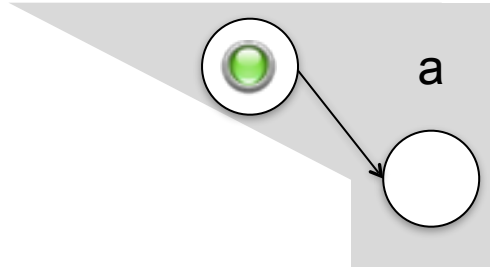
- Instead of computing a **max** of inputs at each node, use **addition**
- Same run-time, same space requirements
- Viterbi cell interpretation
 - What is the score of the best path through the lattice ending in state q at time i ?
- **What does a forward node weight correspond to?**

Forward Algorithm Recurrence

$$\alpha_0(\text{START}) = 1$$

$$\alpha_t(y) = \sum_{q \in \Omega} \eta(q \rightarrow y) \times \gamma(y \downarrow x_i) \times \alpha_{t-1}(q)$$

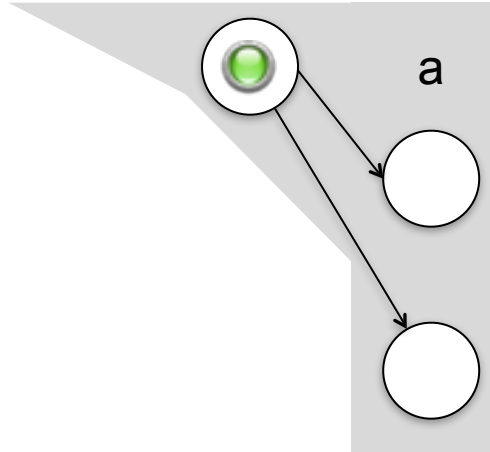
Forward Chart



i=1

$$\alpha_t(q) = p(\text{START}, x_1, \dots, x_t, y_t = q)$$

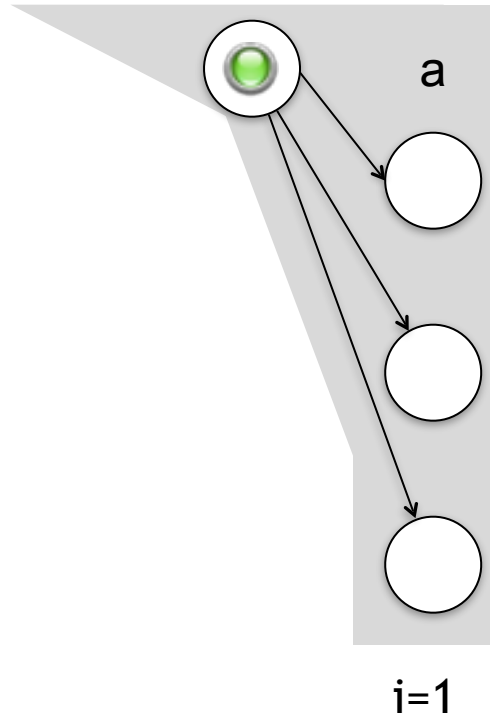
Forward Chart



i=1

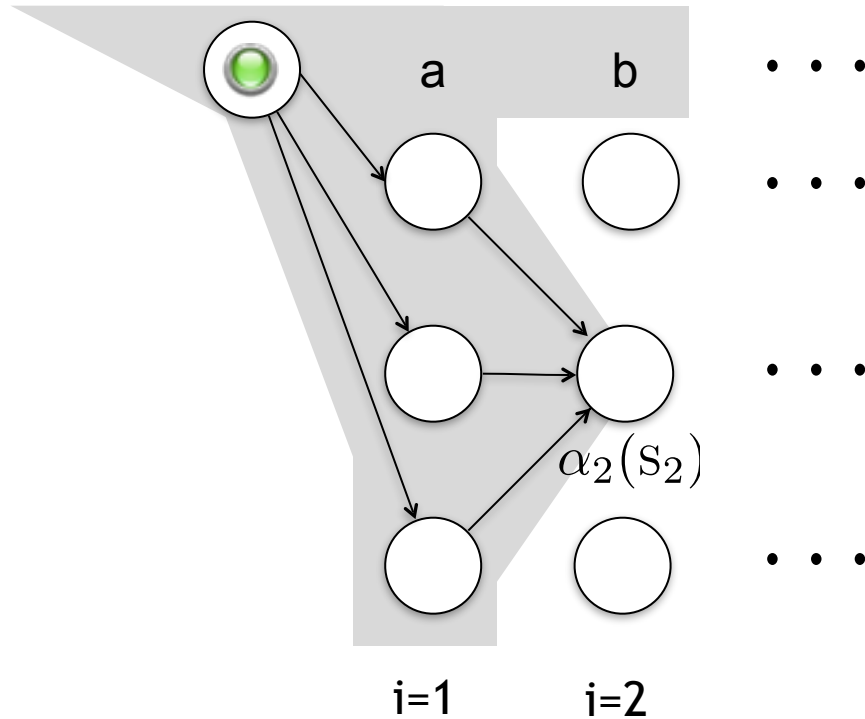
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Forward Chart

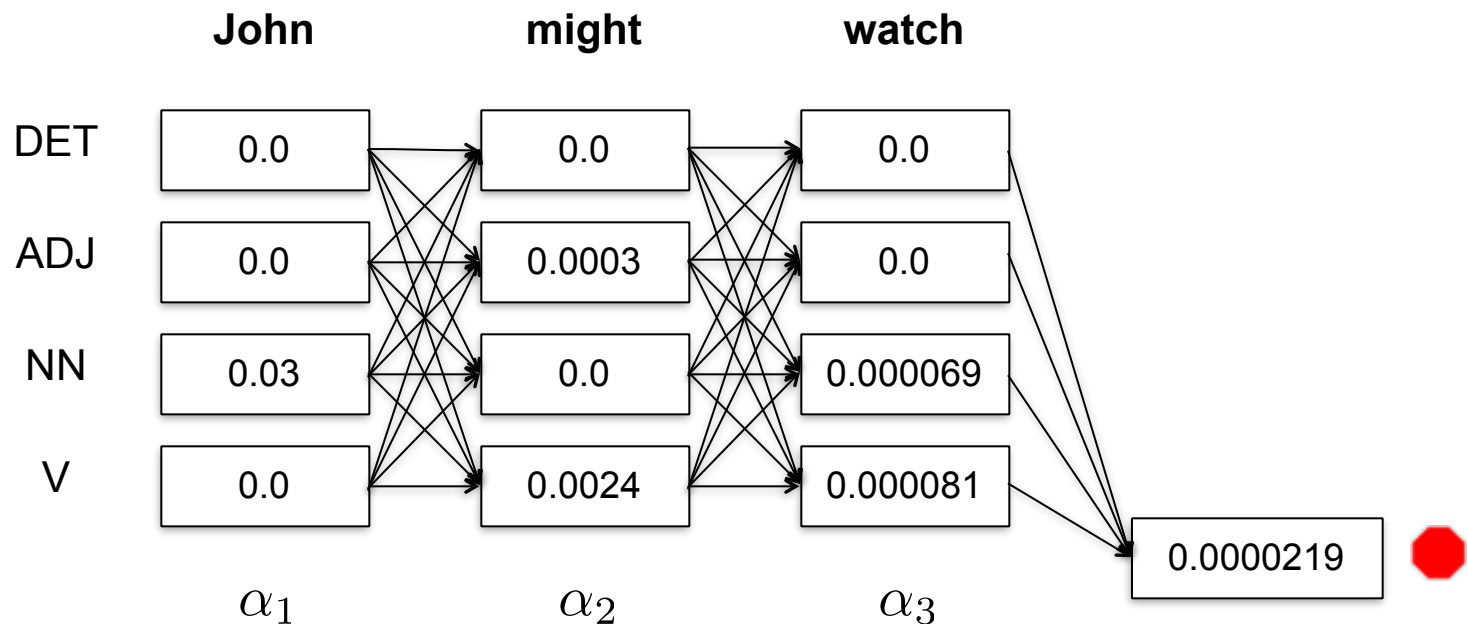


$$\alpha_t(q) = p(\text{START}, x_1, \dots, x_t, y_t = q)$$

Forward Chart



$$\alpha_t(q) = p(\text{START}, x_1, \dots, x_t, y_t = q)$$



$$p = 0.0000219$$

Posterior Marginals

- Marginal inference question for HMMs
 - Given \mathbf{x} , what is the probability of being in a state q at time i ?

$$p(x_1, \dots, x_i, y_i = q \mid y_0 = \text{START}) \times$$

$$p(x_{i+1}, \dots, x_{|\mathbf{x}|} \mid y_i = q)$$

- Given \mathbf{x} , what is the probability of transitioning from state q to r at time i ?

$$p(x_1, \dots, x_i, y_i = q \mid y_0 = \text{START}) \times$$

$$\eta(q \rightarrow r) \times \gamma(r \downarrow x_{i+1}) \times$$

$$p(x_{i+2}, \dots, x_{|\mathbf{x}|} \mid y_{i+1} = r)$$

Posterior Marginals

- Marginal inference question for HMMs
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$$\eta(q \rightarrow r) \times \gamma(r \downarrow x_{i+1}) \times$$

$$p(x_{i+2}, \dots, x_{|\mathbf{x}|} \mid y_{i+1} = r)$$

Backward Algorithm

- Start at the goal node(s) and work **backwards** through the hypergraph
- What is the probability in the goal node cell?
- What if there is more than one cell?
- What is the value of the axiom cell?

Backward Recurrence

$$\beta_{|\mathbf{x}|+1}(\text{STOP}) = 1$$

$$\beta_i(q) = \sum_{r \in \Omega} \beta_{i+1}(r) \times \gamma(r \downarrow x_{i+1}) \times \eta(q \rightarrow r)$$

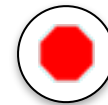
Backward Chart

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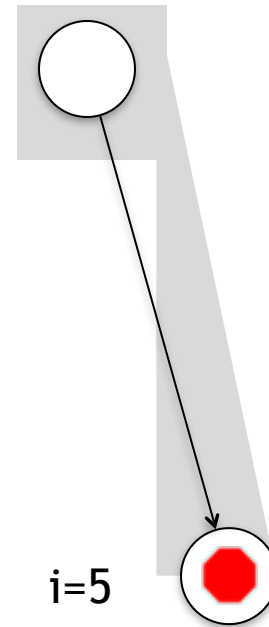
Backward Chart

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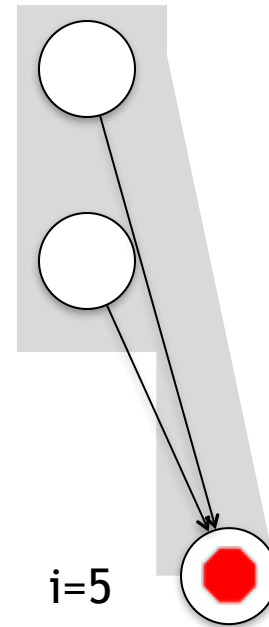
Backward Chart

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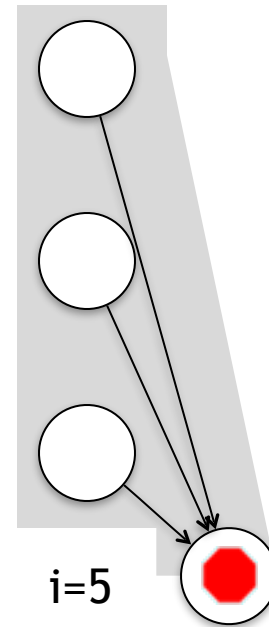
Backward Chart

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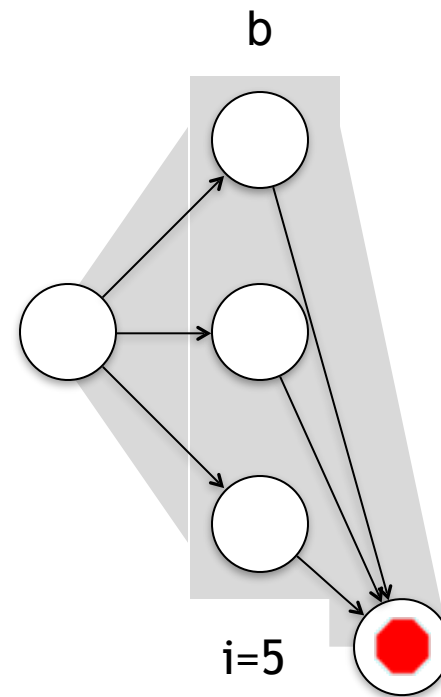
Backward Chart

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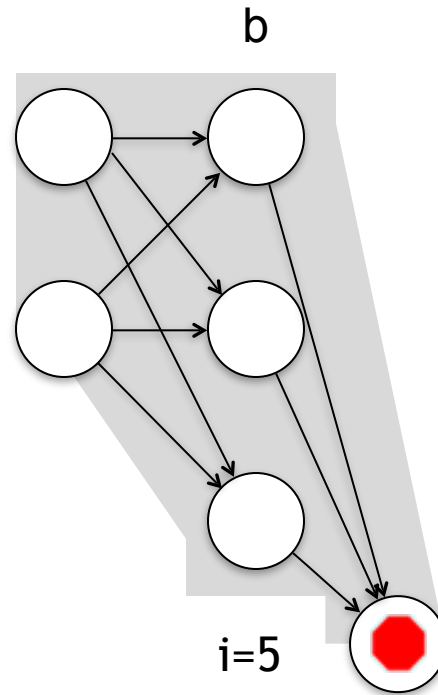
Backward Chart

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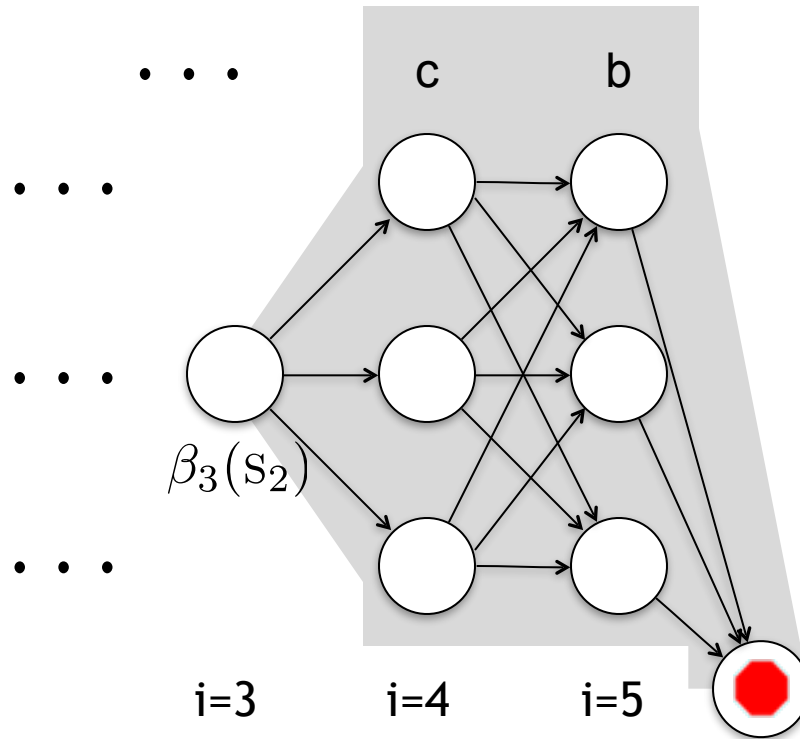
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Backward Chart



$$\beta_t(q) = p(x_{t+1}, \dots, x_{|\mathbf{x}|} \mid y_t = q)$$

Forward-Backward

- Compute forward chart

$$\alpha_t(q) = p(\text{START}, x_1, \dots, x_t, y_t = q)$$

- Compute backward chart

$$\beta_t(q) = p(x_{t+1}, \dots, x_{|\mathbf{x}|}, \text{STOP} \mid y_t = q)$$

What is $\alpha_t(q) \times \beta_t(q)$?

Forward-Backward

- Compute forward chart

$$\alpha_t(q) = p(\text{START}, x_1, \dots, x_t, y_t = q)$$

- Compute backward chart

$$\beta_t(q) = p(x_{t+1}, \dots, x_{|\mathbf{x}|}, \text{STOP} \mid y_t = q)$$

What is $\alpha_t(q) \times \beta_t(q)$ **?**

$$p(\mathbf{x}, y_t = q) = \alpha_t(q) \times \beta_t(q)$$

Edge Marginals

- What is the probability that \mathbf{x} was generated and $q \rightarrow r$ happened at time t ?

$$\begin{aligned} & p(x_1, \dots, x_i, y_i = q \mid y_0 = \text{START}) \times \\ & \quad \eta(q \rightarrow r) \times \gamma(r \downarrow x_{i+1}) \times \\ & \quad p(x_{i+2}, \dots, x_{|\mathbf{x}|} \mid y_{i+1} = r) \end{aligned}$$

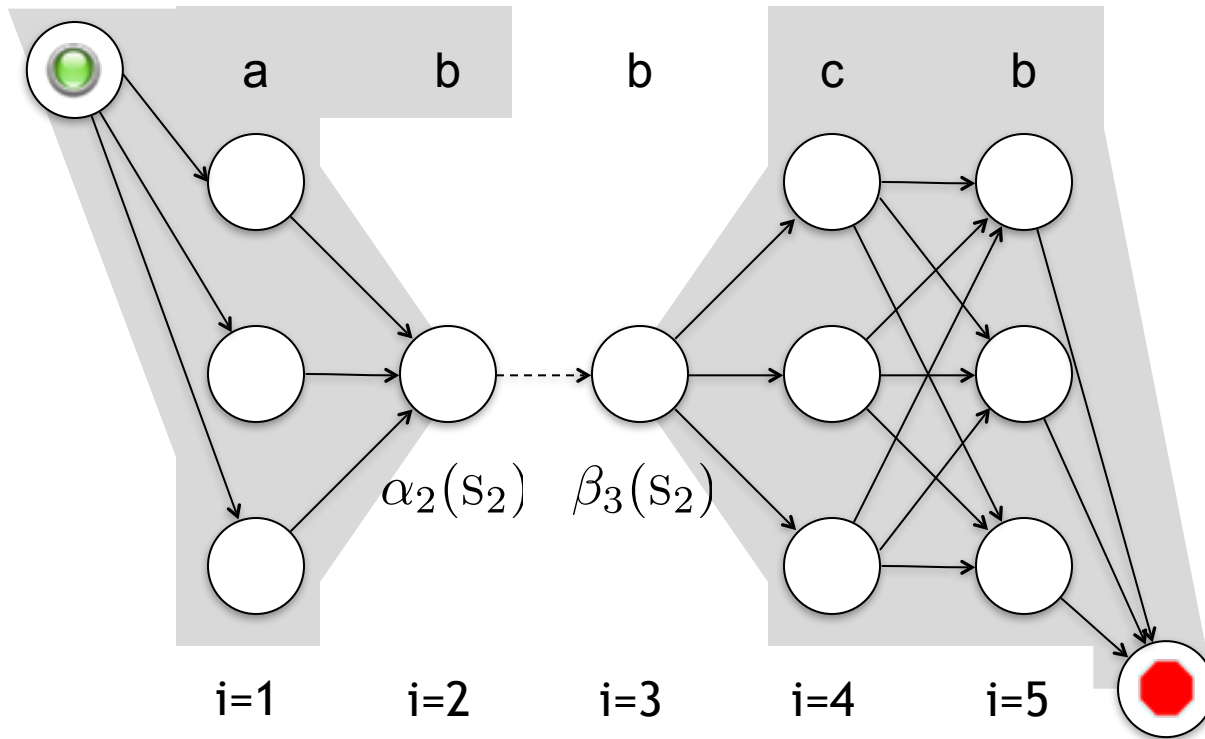
Edge Marginals

- What is the probability that \mathbf{x} was generated and $q \rightarrow r$ happened at time t ?

$$p(x_1, \dots, x_i, y_i = q \mid y_0 = \text{START}) \times \\ \eta(q \rightarrow r) \times \gamma(r \downarrow x_{i+1}) \times \\ p(x_{i+2}, \dots, x_{|\mathbf{x}|} \mid y_{i+1} = r)$$

$$\alpha_t(q) \times \\ \eta(q \rightarrow r) \times \gamma(r \downarrow x_{t+1}) \times \\ \beta_{t+1}(r)$$

Forward-Backward



Generic Inference

- **Semirings** are useful structures in abstract algebra
 - Set of values
 - *Addition*, with additive identity 0: ($a + 0 = a$)
 - *Multiplication*, with mult identity 1: ($a * 1 = a$)
 - Also: $a * 0 = 0$
 - *Distributivity*: $a * (b + c) = a * b + a * c$
 - **Not required**: commutativity, inverses

So What?

- You can unify Forward and Viterbi by changing the semiring

$$\text{FORWARD}(\mathcal{G}) = \bigoplus_{\pi \in \mathcal{G}} \bigotimes_{e \in \pi} w[e]$$

Table 2.1: Elements of common semirings.

semiring	\mathbb{K}	\oplus	\otimes	$\bar{0}$	$\bar{1}$	notes
Boolean	$\{0,1\}$	\vee	\wedge	0	1	idempotent
count	$\mathbb{N}_0 \cup \{\infty\}$	$+$	\times	0	1	
probability	$\mathbb{R}_+ \cup \{\infty\}$	$+$	\times	0	1	
tropical	$\mathbb{R} \cup \{-\infty, \infty\}$	\max	$+$	$-\infty$	0	idempotent
log	$\mathbb{R} \cup \{-\infty, \infty\}$	\oplus_{\log}	$+$	$-\infty$	0	

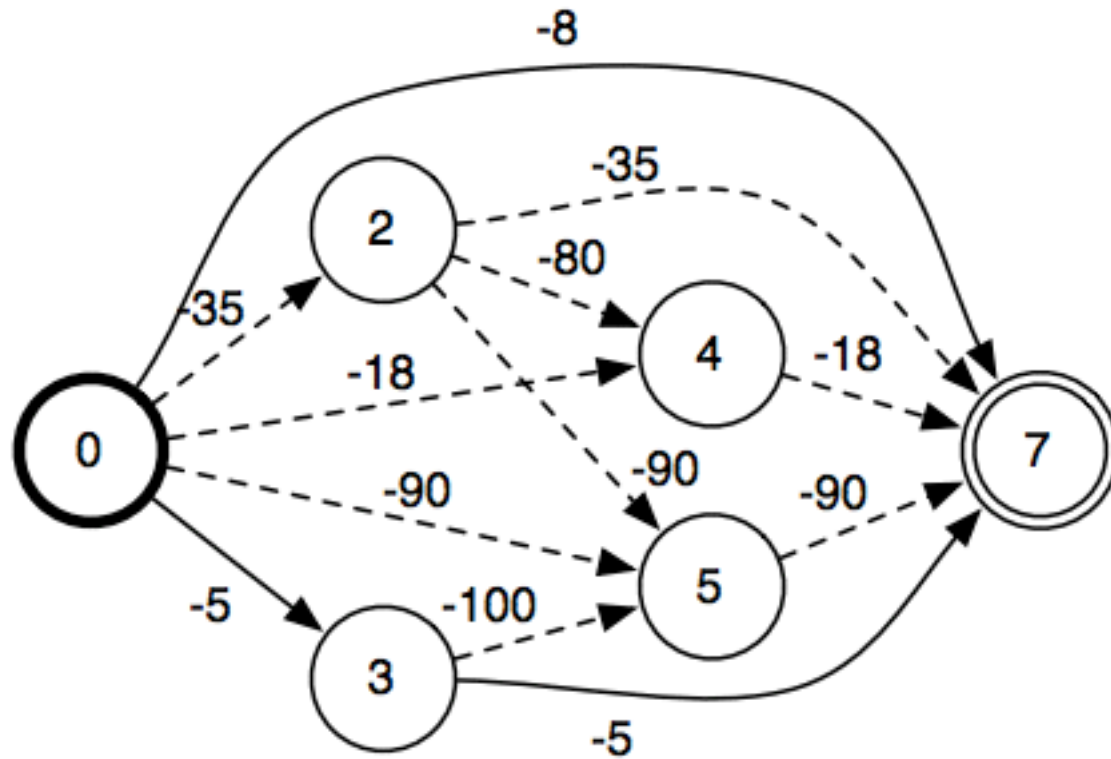
Semiring Inside

- **Probability semiring**
 - marginal probability of output
- **Counting semiring**
 - number of paths (“taggings”)
- **Viterbi semiring**
 - best scoring derivation
- **Log semiring** $w[e] = \mathbf{w}^T \mathbf{f}(e)$
 - $\log(Z) = \log$ partition function

Semiring Edge-Marginals

- **Probability semiring**
 - posterior marginal probability of each edge
- **Counting semiring**
 - number of paths going through each edge
- **Viterbi semiring**
 - score of best path going through each edge
- **Log semiring**
 - $\log(\text{sum of all exp path weights of all paths with } e)$
= $\log(\text{posterior marginal probability}) + \log(Z)$

Max-Marginal Pruning



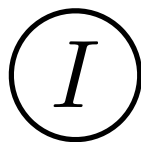
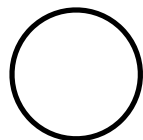
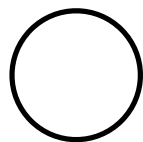
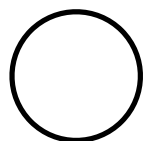
Weighted Logic Programming

- Slightly different notation than the textbook, but you will see it in the literature
- WLP is useful here because it lets us **build hypergraphs**

$$\begin{array}{c}
 \frac{I_1 : w_1 \quad I_2 : w_2 \quad \cdots \quad I_k : w_k}{I : w} \quad \phi \\
 \bigcirc \\
 \vdots
 \end{array}$$

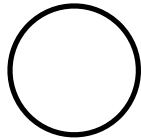
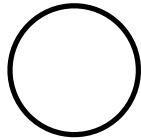
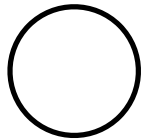
Hypergraphs

A	B	C
<hr/>		
	I	

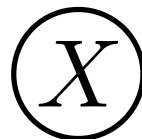
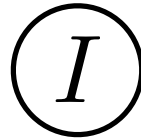


Hypergraphs

$$\frac{A \quad B \quad C}{I}$$



$$\frac{X \quad Y}{I}$$

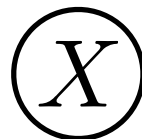
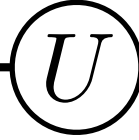
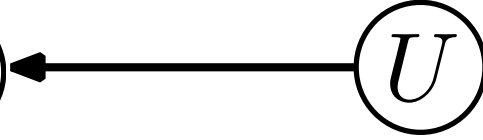
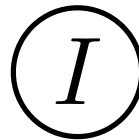
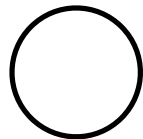
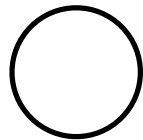
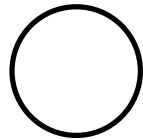


Hypergraphs

$$\frac{A \quad B \quad C}{I}$$

$$\frac{X \quad Y}{I}$$

$$\frac{U}{I}$$



Generalizing Forward-Backward

- Forward/Backward algorithms are a special case of **Inside/Outside algorithms**
- It's helpful to think of I/O as algorithms on PCFG parse forests, but it's more general
 - Recall the 5 views of decoding: decoding is parsing
 - **More specifically, decoding is a weighted proof forest**

CKY Algorithm

Item form

$[X, i, j]$

CKY Algorithm

Item form

$[X, i, j]$

Goals

$[S, 1, |\mathbf{x}| + 1]$

CKY Algorithm

Item form

$[X, i, j]$

Goals

$[S, 1, |\mathbf{x}| + 1]$

Axioms

$$\frac{}{[N, i, i + 1] : w} \quad (N \xrightarrow{w} x_i) \in G$$

CKY Algorithm

Item form

$[X, i, j]$

Goals

$[S, 1, |\mathbf{x}| + 1]$

Axioms

$$\frac{}{[N, i, i + 1] : w} \quad (N \xrightarrow{w} x_i) \in G$$

Inference rules

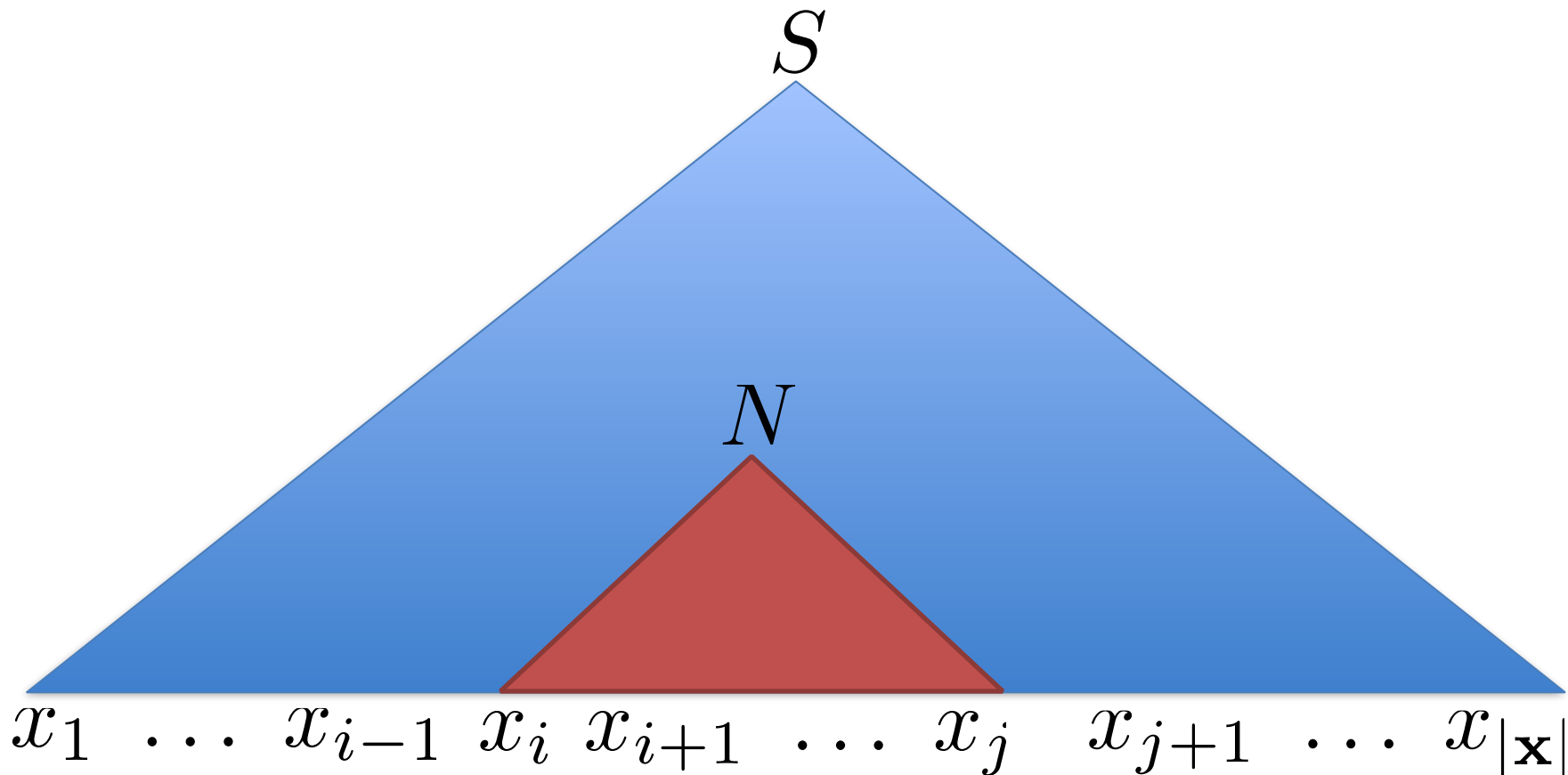
$$\frac{[X, i, k] : u \quad [Y, k, j] : v}{[Z, i, j] : u \otimes v \otimes w} \quad (Z \xrightarrow{w} X Y) \in G$$

Posterior Marginals

- Marginal inference question for PCFGs
 - Given w , what is the probability of having a constituent of type Z from i to j ?
 - Given w , what is the probability of having a constituent of *any* type from i to j ?
 - Given w , what is the probability of using rule $Z \rightarrow XY$ to derive the span from i to j ?

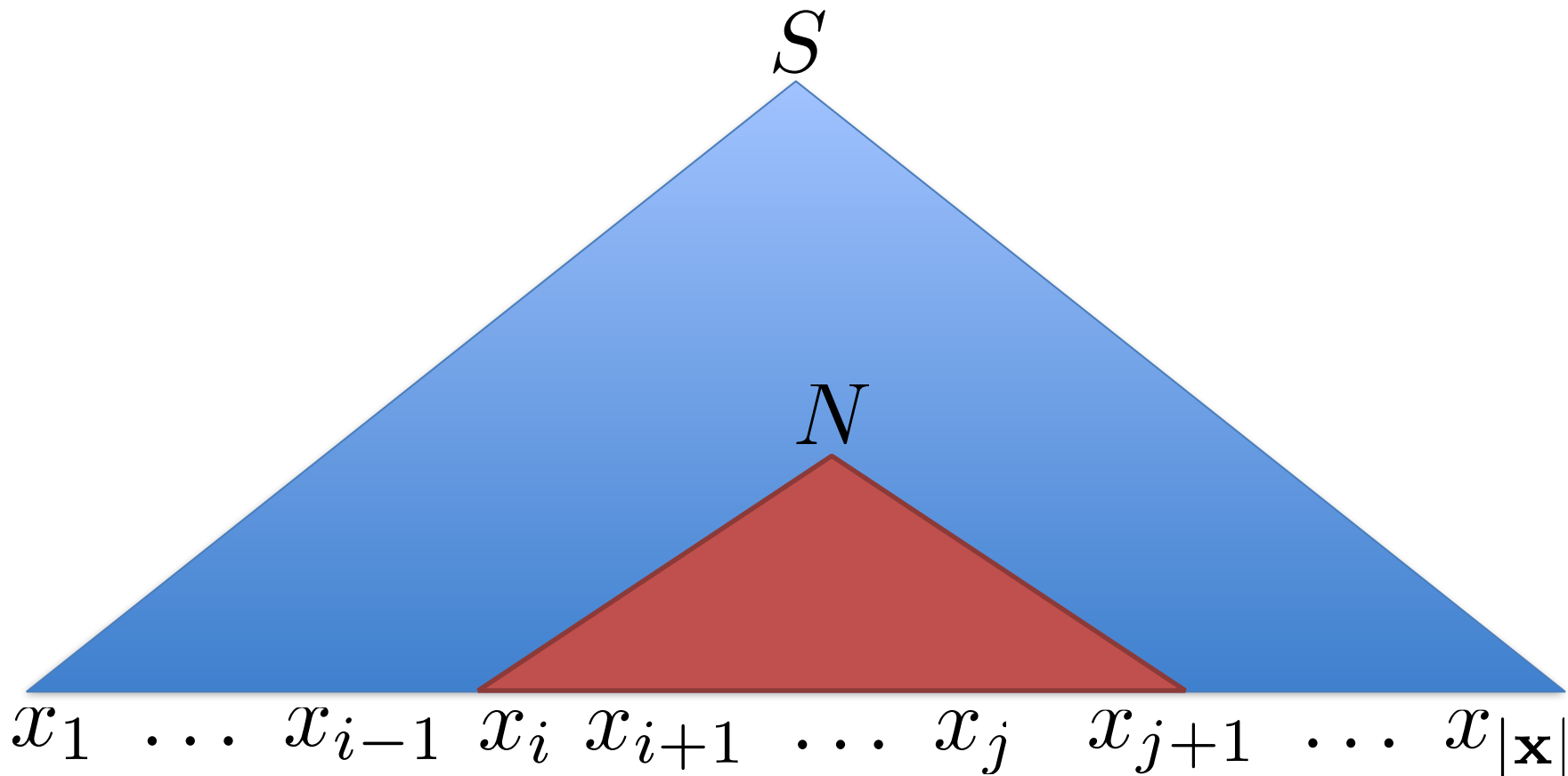
Inside Algorithm

$$\alpha_{[i,j]}(N) = p(x_i, x_{i+1}, \dots, x_j \mid N; \mathcal{G})$$



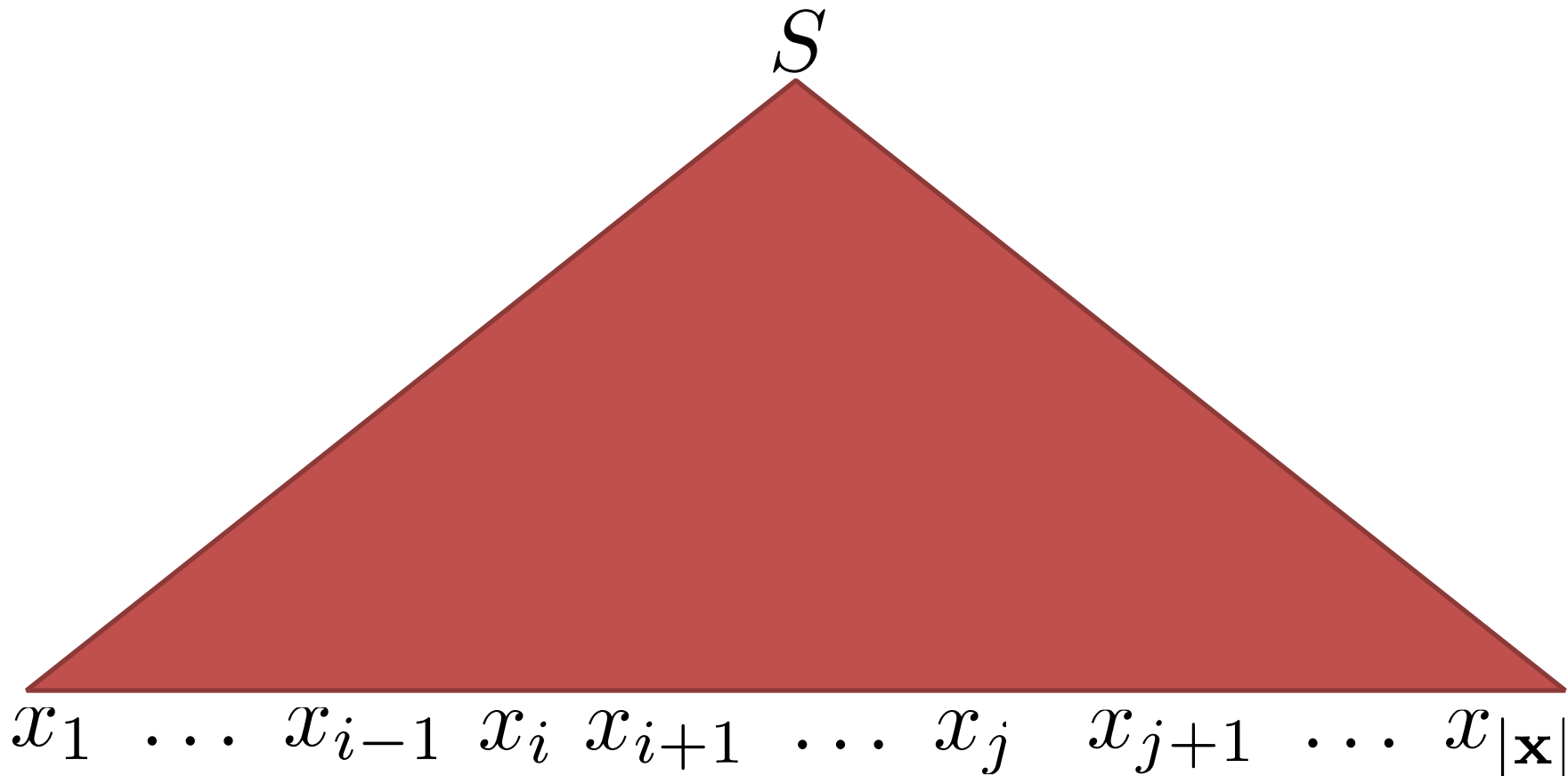
Inside Algorithm

$$\alpha_{[i,j]}(N) = p(x_i, x_{i+1}, \dots, x_j \mid N; \mathcal{G})$$



Inside Algorithm

$$\alpha_{[i,j]}(N) = p(x_i, x_{i+1}, \dots, x_j \mid N; \mathcal{G})$$



CKY Inside Algorithm

Base case(s)

$$\alpha_{[i,i+1]}(Z) = p(Z \rightarrow x_i)$$

Recurrence

$$\alpha_{[i,j]}(Z) = \sum_{k=i+1}^{j-1} \sum_{(Z \rightarrow XY) \in G} \alpha_{[i,k]}(X) \times \alpha_{[k,j]}(Y) \times p(Z \rightarrow XY)$$

Generic Inside

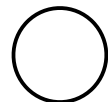
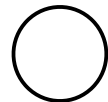
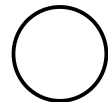
```

1: function INSIDE( $\mathcal{G}, K$ )                                 $\triangleright \mathcal{G}$  is an acyclic hypergraph and  $K$  is a semiring
2:   for  $q$  in topological order in  $\mathcal{G}$  do
3:     if  $B(q) = \emptyset$  then
4:        $\alpha(q) \leftarrow \bar{1}$                                  $\triangleright$  assume states with no in-edges are axioms
5:     else
6:        $\alpha(q) \leftarrow \bar{0}$ 
7:       for all  $e \in B(q)$  do                                 $\triangleright$  all in-coming edges to node  $q$ 
8:          $k \leftarrow w(e)$ 
9:         for all  $r \in \mathbf{t}(e)$  do                                 $\triangleright$  all tail (previous) nodes of edge  $e$ 
10:           $k \leftarrow k \otimes \alpha(r)$ 
11:           $\alpha(q) \leftarrow \alpha(q) \oplus k$ 
12:   return  $\alpha$ 

```

$$B(I) = \{e_1, e_2, e_3\}$$

$$\mathbf{t}(e_1) = \langle A, B, C \rangle$$


 e_1

 e_2

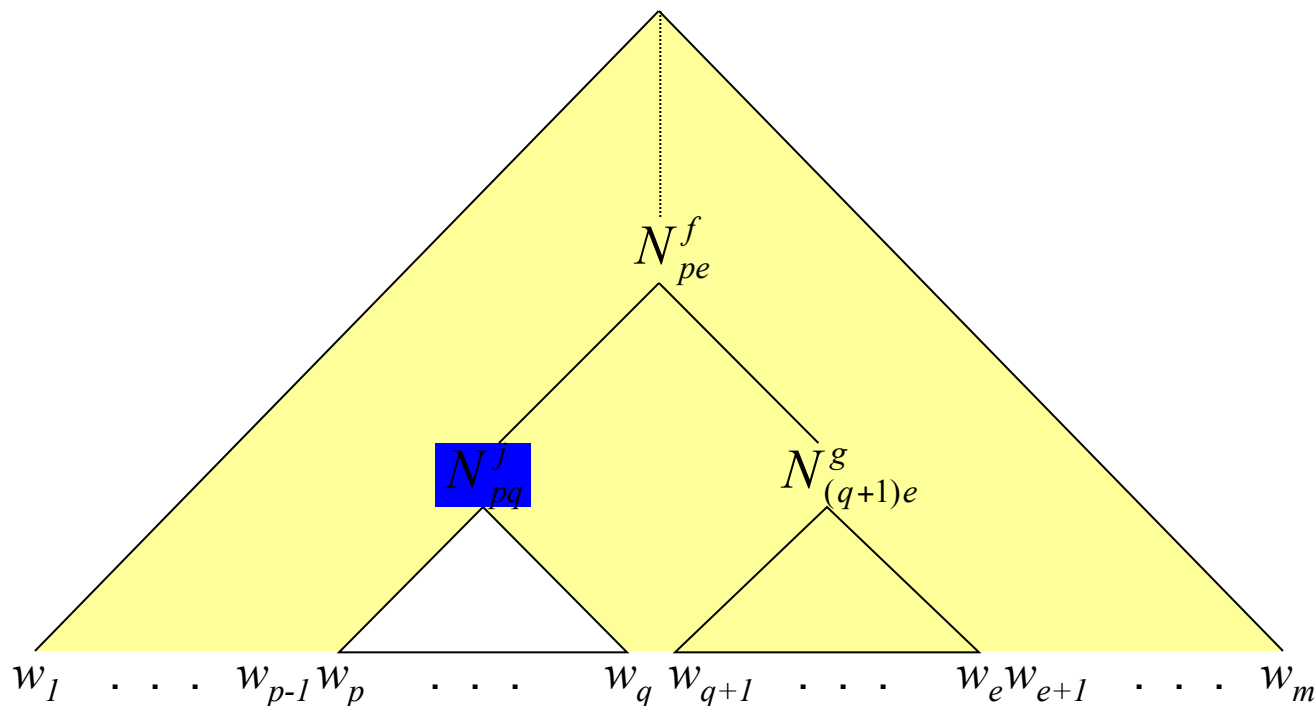
 e_3


Questions for Generic Inside

- Probability semiring
 - Marginal probability of input
- Counting semiring
 - Number of paths (pares, labels, etc)
- Viterbi semiring
 - Viterbi probability (max joint probability)
- Log semiring
 - $\log Z(\text{input})$

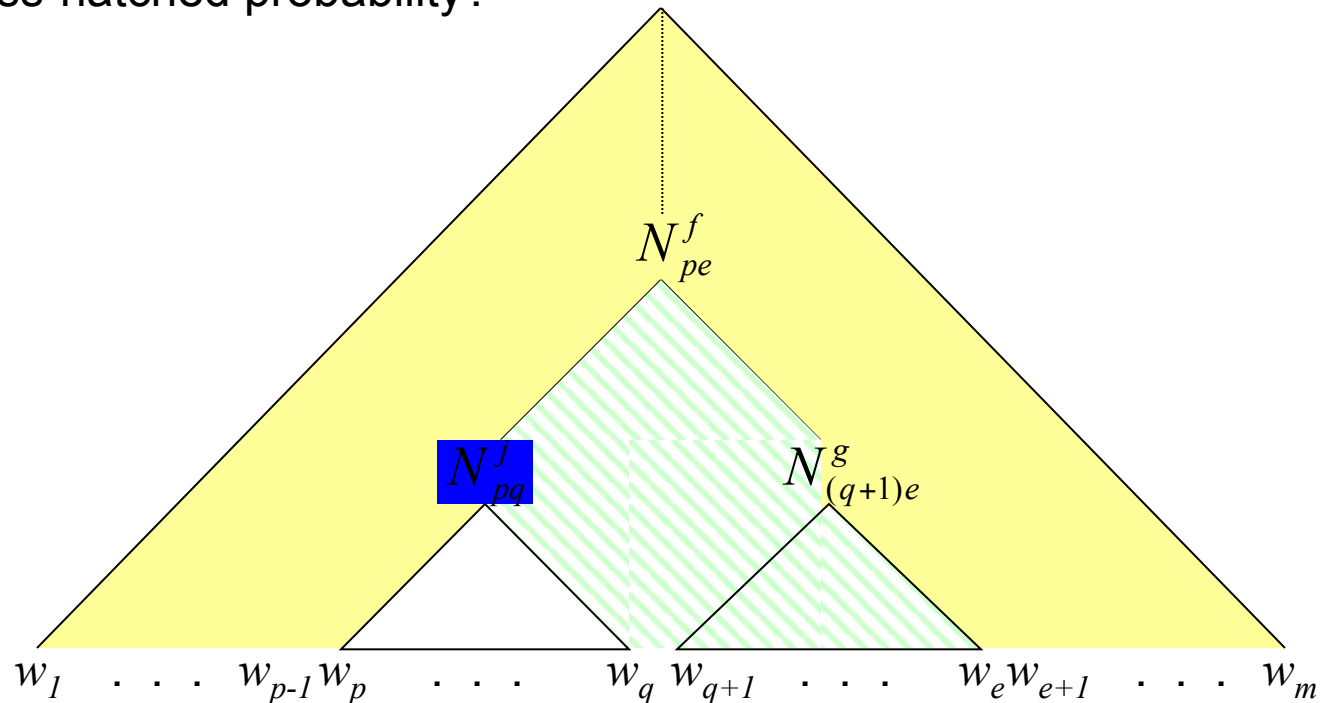
Outside probabilities: decomposing the problem

The shaded area represents the outside probability $\alpha_j(p, q)$ which we need to calculate. How can this be decomposed?



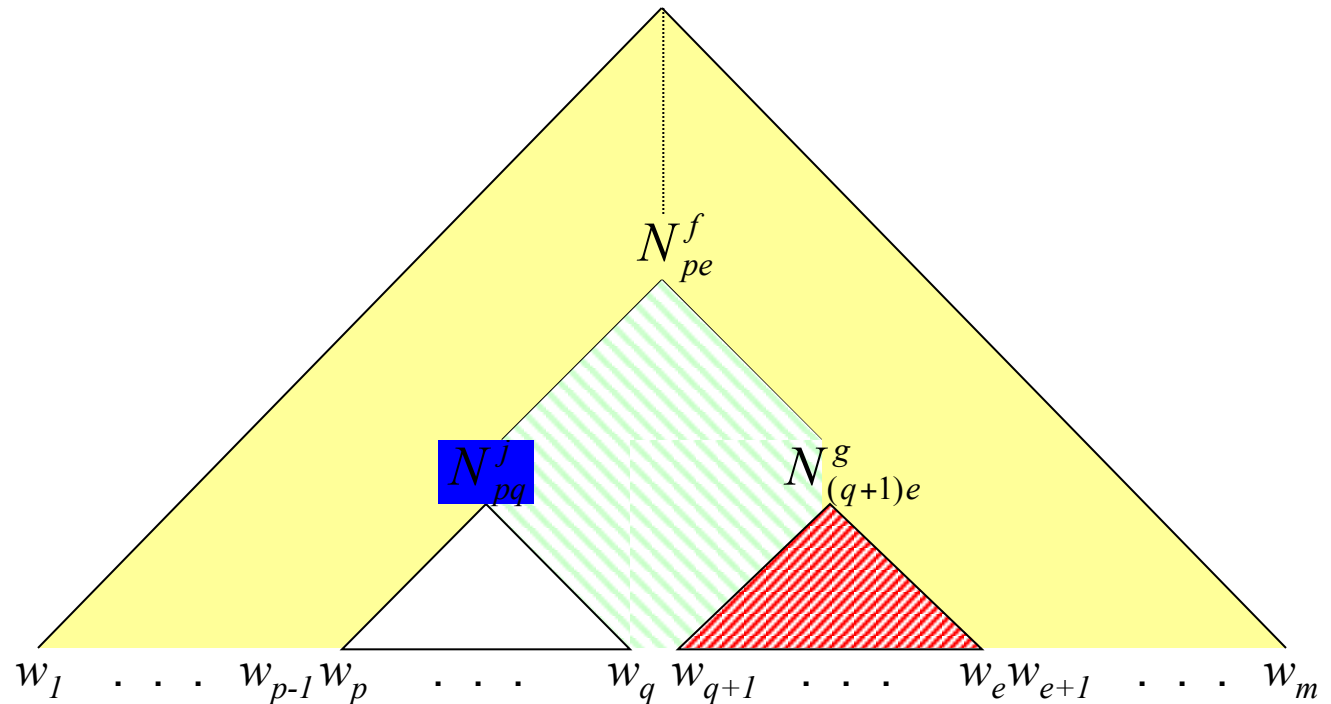
Outside probabilities: decomposing the problem

Step 1: We assume that N_{pe}^f is the parent of N_{pq}^j . Its outside probability, $\alpha_f(p, e)$, (represented by the yellow shading) is available recursively. How do we calculate the cross-hatched probability?



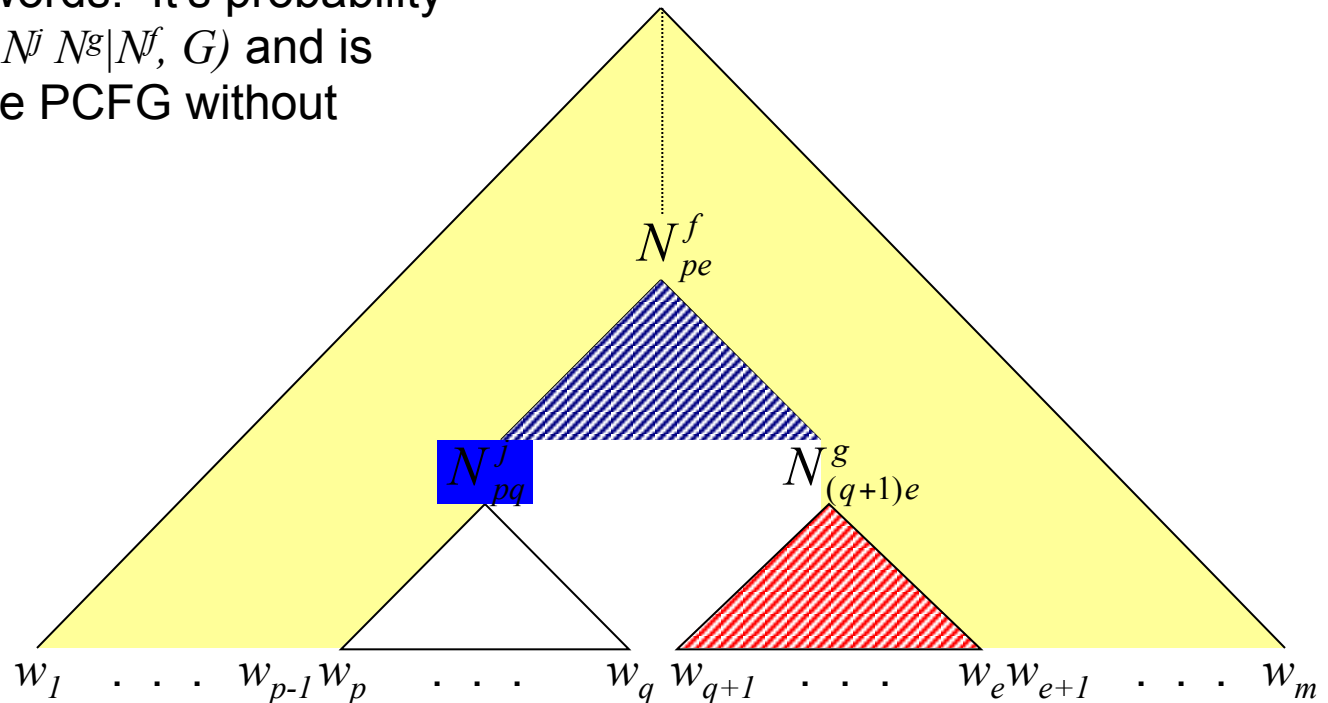
Outside probabilities: decomposing the problem

Step 2: The red shaded area is the inside probability of $N_{(q+1)e}^g$, which is available as $\beta_g(q+1, e)$.



Outside probabilities: decomposing the problem

Step 3: The blue shaded part corresponds to the production $N^f \rightarrow N^j N^g$, which because of the context-freeness of the grammar, is not dependent on the positions of the words. It's probability is simply $P(N^f \rightarrow N^j N^g | N^f, G)$ and is available from the PCFG without calculation.



Generic Outside

```
1: function OUTSIDE( $\mathcal{G}, K, \alpha$ )                                ▷  $\alpha$  is the result of INSIDE( $\mathcal{G}, K$ )
2:   for all  $q \in \mathcal{G}$  do
3:      $\beta(q) \leftarrow \bar{0}$ 
4:      $\beta(q_{goal}) = \bar{1}$ 
5:   for  $q$  in reverse topological order in  $\mathcal{G}$  do
6:     for all  $e \in B(q)$  do                                     ▷ all in-coming edges to node  $q$ 
7:       for all  $r \in t(e)$  do                                     ▷ all tail (previous) nodes of edge  $e$ 
8:          $k \leftarrow w(e) \otimes \beta(q)$ 
9:         for all  $s \in t(e)$  do                                     ▷ all tail (previous) nodes of edge  $e$ , again
10:          if  $r \neq s$  then
11:             $k \leftarrow k \otimes \alpha(s)$                        ▷ incorporate inside score
12:           $\beta(r) \leftarrow \beta(r) \oplus k$ 
13:   return  $\beta$ 
```

Generic Inside-Outside

```
1: function INSIDEOUTSIDE( $\mathcal{G}, K$ )                                ▷ compute edge marginals
2:    $\alpha \leftarrow \text{INSIDE}(\mathcal{G}, K)$ 
3:    $\beta \leftarrow \text{OUTSIDE}(\mathcal{G}, K, \alpha)$ 
4:   for edge  $e$  in  $\mathcal{G}$  do
5:      $\gamma(e) \leftarrow w(e) \otimes \beta(n(e))$     ▷ edge weight and outside score of edge's head node
6:     for all  $q \in \mathbf{t}(e)$  do
7:        $\gamma(e) \leftarrow \gamma(e) \otimes \alpha(q)$     ▷ inside score of tail nodes
8:   return  $\gamma$                                               ▷  $\gamma(e)$  is the edge marginal of  $e$ 
```

Inside-Outside

- Inside probabilities are required to compute Outside probabilities
- Inside-Outside works where Forward-Backward does, but not vice-versa
- Implementation considerations
 - Building a hypergraph explicitly simplifies code, but it can be expensive in terms of memory