Assemblers, Linkers, and Loaders

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See: P&H Appendix B.3-4

C

compiler

MIPS assembly

assembler

machine code





Gates



```
calc.c
vector v = malloc(8);
v->x = prompt("enter x");
v->y = prompt("enter y");
int c = pi + tnorm(v);
print("result", c);
math.c
int tnorm(vector v) {
 return abs(v->x)+abs(v->y);
lib3410.o
  global variable: pi
  entry point: prompt
  entry point: print
```

entry point: malloc

```
int n = 100;
int main (int argc, char* argv[]) {
      int i;
      int m = n;
      int count = 0;
      for (i = 1; i <= m; i++)
             count += i;
      printf ("Sum 1 to %d is %d\n", n, count);
[csug01] mipsel-linux-gcc -S add1To100.c
```

n:	<pre>.data .globl .align .word .rdata .align</pre>	n 2 100 2	\$L2:	lw lw slt bne lw	\$2,24(\$fp) \$3,28(\$fp) \$2,\$3,\$2 \$2,\$0,\$L3 \$3,32(\$fp)
\$str0:	.asciiz "Sum .text .align .globl	<pre>1 to %d is %d\n" 2 main \$sp,\$sp,-48 \$31,44(\$sp) \$fp.40(\$sp)</pre>		lw addu sw lw addiu	\$2,24(\$fp) \$2,\$3,\$2 \$2,32(\$fp) \$2,24(\$fp) \$2,\$2,1
	addiu sw sw move sw sw la lw sw sw li sw		\$L3:	sw b la lw jal move lw lw addiu j	\$2,\$2,1 \$2,24(\$fp) \$L2 \$4,\$str0 \$5,28(\$fp) \$6,32(\$fp) printf \$sp,\$fp \$31,44(\$sp) \$fp,40(\$sp) \$sp,\$sp,48 \$31

Variables

Visibility Lifetime Location

Function-Local

Global

Dynamic

```
int n = 100;
int main (int argc, char* argv[]) {
    int i, m = n, count = 0, *A = malloc(4 * m);
    for (i = 1; i <= m; i++) { count += i; A[i] = count; }
    printf ("Sum 1 to %d is %d\n", n, count);</pre>
```

Variables

Visibility Lifetime Location

Function-Local

Global

Dynamic

C Pointers can be trouble

```
int *trouble()
{ int a; ...; return &a; }
char *evil()
{ char s[20]; gets(s); return s; }
int *bad()
{ s = malloc(20); ... free(s); ... return s; }
(Can't do this in Java, C#, ...)
```

Variables

Visibility Lifetime Location

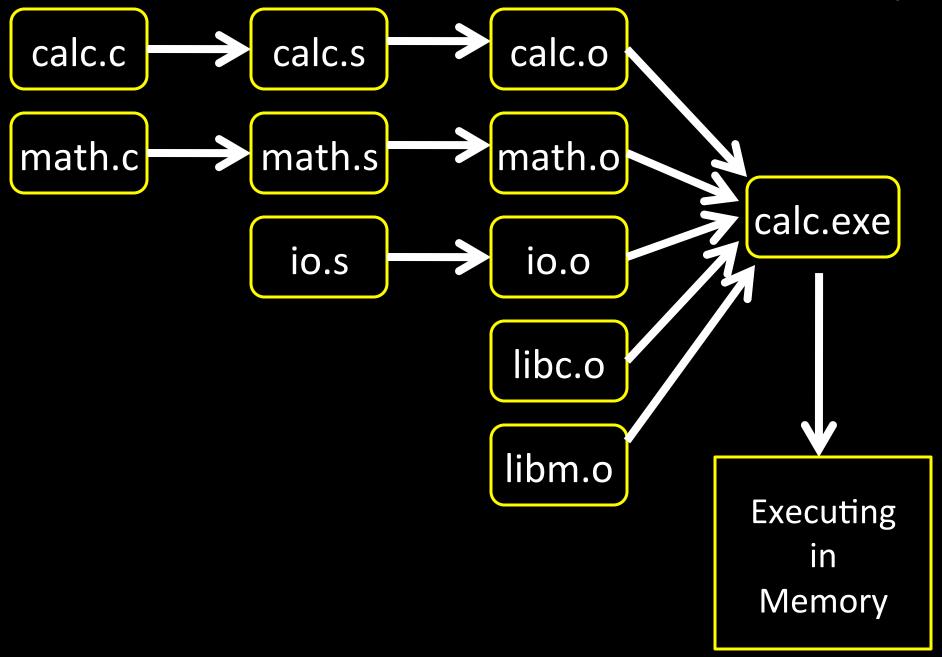
Function-Local

Global

Dynamic

C Pointers can be trouble

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int *trouble()
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{ char s[20]; gets(s); return s; }
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(Can't do this in Java, C#, ...)
```



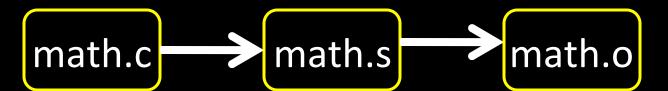
Compiler output is assembly files

Assembler output is obj files

Linker joins object files into one executable

Loader brings it into memory and starts execution

Compilers and Assemblers



Output is obj files

- Binary machine code, but not executable
- May refer to external symbols
- Each object file has illusion of its own address space
 - Addresses will need to be fixed later

Global labels: Externally visible "exported" symbols

- Can be referenced from other object files
- Exported functions, global variables

Local labels: Internal visible only symbols

- Only used within this object file
- static functions, static variables, loop labels, ...

Header

Size and position of pieces of file

Text Segment

instructions

Data Segment

static data (local/global vars, strings, constants)

Debugging Information

line number → code address map, etc.

Symbol Table

- External (exported) references
- Unresolved (imported) references

math.c

```
int pi = 3;
int e = 2;
static int randomval = 7;
extern char *username;
extern int printf(char *str, ...);
int square(int x) { ... }
static int is_prime(int x) { ... }
int pick prime() { ... }
int pick_random() {
   return randomval;
```

gcc -S ... math.c gcc -c ... math.s objdump --disassemble math.o objdump --syms math.o

csug01 ~\$ mipsel-linux-objdump --disassemble math.o

math.o: file format elf32-tradlittlemips

Disassembly of section .text:

00000000 <pick_random>:

|--|

4: afbe0000 sw s8,0(sp)

8: 03a0f021 move s8,sp

c: 3c020000 lui v0,0x0

10: 8c420008 lw v0,8(v0)

14: 03c0e821 move sp,s8

18: 8fbe0000 lw s8,0(sp)

1c: 27bd0008 addiu sp,sp,8

20: 03e00008 jr ra

24: 00000000 nop

00000028 <square>:

28: 27bdfff8 addiu sp,sp,-8

2c: afbe0000 sw s8,0(sp)

30: 03a0f021 move s8,sp

34: afc40008 sw a0,8(s8)

•••

csug01 ~\$ mipsel-linux-objdump --syms math.o

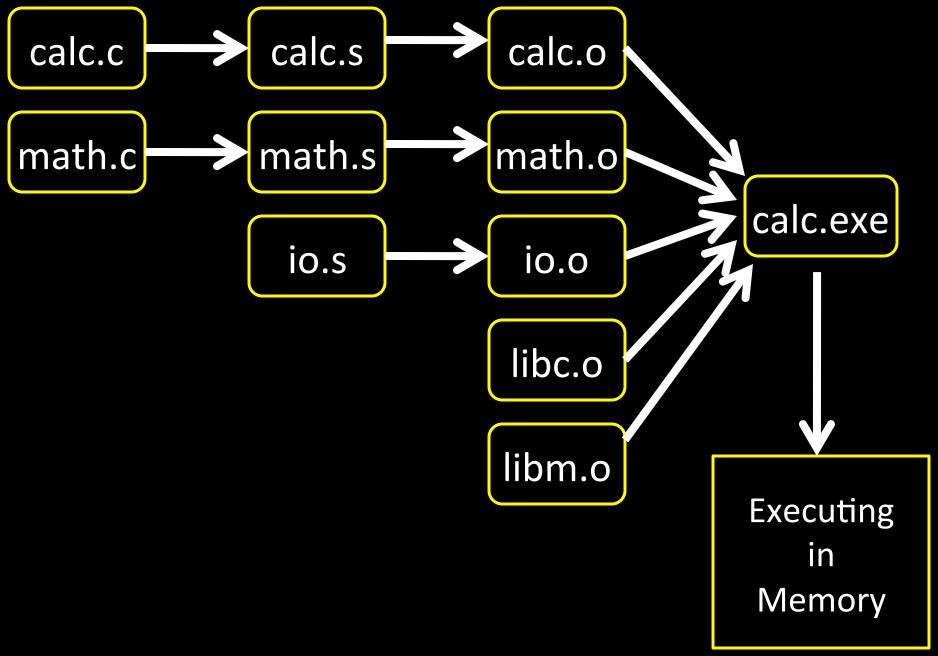
math.o: file format elf32-tradlittlemips

```
SYMBOL TABLE:
00000000 1
              df *ABS*
                                 00000000 math.c
00000000 1
                                 00000000 .text
              d
                 .text
                                 00000000 .data
00000000 1
          d
                 .data
00000000 1
              d
                 .bss
                                 00000000 .bss
00000000 1
              d
                 .mdebug.abi32
                                 00000000 .mdebug.abi32
                 .data
8000000
               0
                                 00000004 randomval
00000060 1
                .text
                                 00000028 is prime
00000000 1
              d
                 .rodata
                                 00000000 .rodata
00000000
                                 00000000 .comment
                 .comment
00000000 g
                .data
                                 00000004 pi
00000004 g
                .data
                                 00000004 e
00000000
                                 00000028 pick random
               F .text
                                 00000038 square
00000028
                .text
00000088 g
                .text
                                 0000004c pick prime
               F
0000000
                 *UND*
                                 00000000 username
                                 00000000 printf
00000000
                 *UND*
```

Q: Why separate compile/assemble and linking steps?

A: Can recompile one object, then just relink.

Linkers



Linker combines object files into an executable file

- Relocate each object's text and data segments
- Resolve as-yet-unresolved symbols
- Record top-level entry point in executable file

End result: a program on disk, ready to execute

main.o

→ 0C000000 21035000 1b80050C → 4C040000 21047002 → 0C000000

00 T main
00 D uname
UND printf
UND pi

40, JL, printf 4C, LW/gp, pi 54, JL, square

math.o

21032040 **→** 0C000000 1b301402 → 3C040000 → 34040000 20 T square 00 D рi printf *UND* *UND* uname 28, JL, printf 30, LUI, uname 34, LA, uname

main.o

→ 0C000000 21035000 1b80050C → 4C040000 21047002 → 0C000000

00 T main
00 D uname
UND printf
UND pi

40, JL, printf 4C, LW/gp, pi 54, JL, square

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21032040 **→** 0C000000 1b301402 → 3C040000 → 34040000 20 T square 00 D рi printf *UND* *UND* uname 28, JL, printf 30, LUI, uname 34, LA, uname

printf.o

3C T printf

main.o

→ 0C000000 21035000 1b80050C → 4C040000 21047002 → 0C000000

00 T main 00 D uname *UND* printf *UND* pi

40, JL, printf 4C, LW/gp, pi 54, JL, square

math.o

21032040 → 0C000000 1b301402 → 3C040000 → 34040000

20 T square 00 D pi *UND* printf *UND* uname

28, JL, printf 30, LUI, uname 34, LA, uname

printf.o -

3C T printf

calc.exe -

21032040 0C40023C 1b301402 3C041000 34040004

0C40023C 21035000 1b80050c 4C048004 21047002 0C400020

10201000 21040330 22500102

00000003 0077616B

entry:400100 text: 400000

data:1000000

Header

location of main entry point (if any)

Text Segment

instructions

Data Segment

static data (local/global vars, strings, constants)

Relocation Information

- Instructions and data that depend on actual addresses
- Linker patches these bits after relocating segments

Symbol Table

Exported and imported references

Debugging Information

Unix

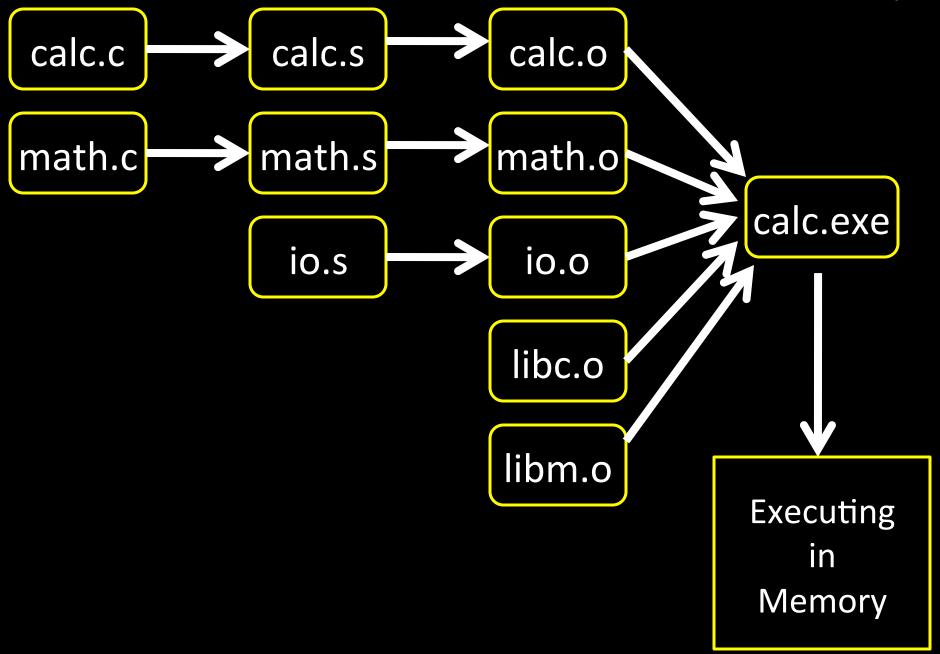
- a.out
- COFF: Common Object File Format
- ELF: Executable and Linking Format
- •

Windows

PE: Portable Executable

All support both executable and object files

Loaders and Libraries



Loader reads executable from disk into memory

- Initializes registers, stack, arguments to first function
- Jumps to entry-point

Part of the Operating System (OS)

Static Library: Collection of object files (think: like a zip archive)

Q: But every program contains entire library!

A: Linker picks only object files needed to resolve undefined references at link time

e.g. libc.a contains many objects:

- printf.o, fprintf.o, vprintf.o, sprintf.o, snprintf.o, ...
- read.o, write.o, open.o, close.o, mkdir.o, readdir.o, ...
- rand.o, exit.o, sleep.o, time.o,

Q: But every program still contains part of library!

A: shared libraries

- executable files all point to single shared library on disk
- final linking (and relocations) done by the loader

Optimizations:

- Library compiled at fixed non-zero address
- Jump table in each program instead of relocations
- Can even patch jumps on-the-fly

Direct call:

```
00400010 <main>:
     0x00400330
     0x00400620
     0x00400330
00400330 <printf>:
00400620 <gets>:
```

Drawbacks:

Linker or loader must edit every use of a symbol (call site, global var use, ...)

Idea:

Put all symbols in a single "global offset table"

Code does lookup as needed

```
00400010 <main>:
 jal 0x00400330
 jal 0x00400620
 jal 0x00400330
00400330 <printf>:
00400620 <gets>:
```

GOT	global offset table

Indirect call:

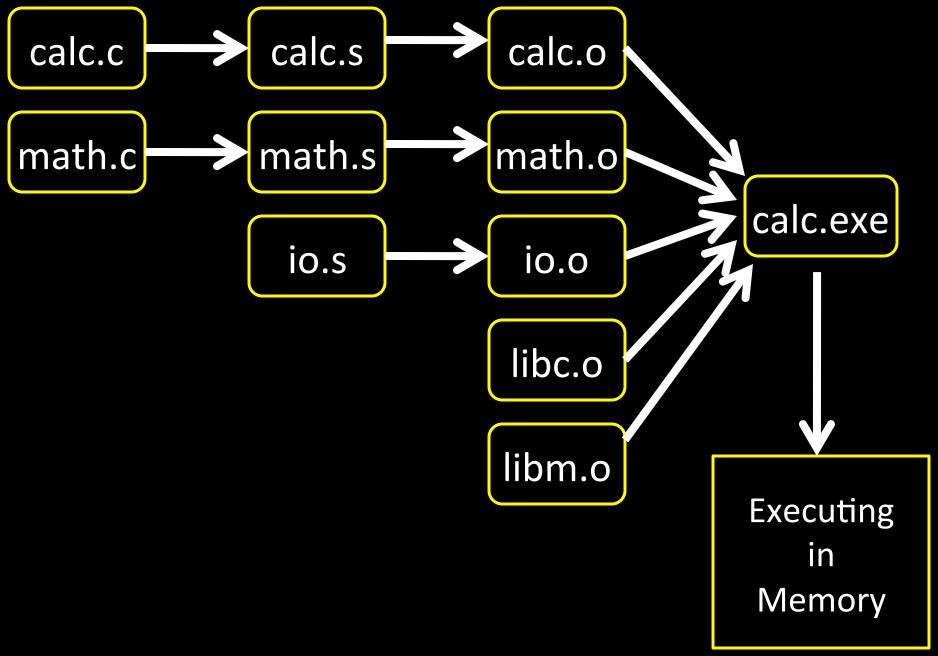
```
00400010 <main>:
 lw t9, ? # printf
 jalr t9
 lw t9, ? # gets
 jalr t9
00400330 <printf>:
00400620 <gets>:
```

```
# data segment
# global offset table
# to be loaded
\# at -32712(gp)
.got
.word 00400010 # main
.word 00400330 # printf
.word 00400620 # gets
```

Indirect call with on-demand dynamic linking:

```
00400010 <main>:
  # load address of prints
  # from .got[1]
  lw t9, -32708(gp)
  # also load the index 1
  li t8, 1
  # now call it
  jalr t9
.got
  .word 00400888 # open
  .word 00400888 # prints
  .word 00400888 # gets
  .word 00400888 # foo
```

```
00400888 <dlresolve>:
  # t9 = 0x400888
  # t8 = index of func that
       needs to be loaded
  # load that func
  \dots # t7 = loadfromdisk(t8)
  # save func's address so
  # so next call goes direct
  ... # got[t8] = t7
  # also jump to func
  jr t7
  # it will return directly
  # to main, not here
```



Windows: dynamically loaded library (DLL)

PE format

Unix: dynamic shared object (DSO)

ELF format

Unix also supports Position Independent Code (PIC)

- Program determines its current address whenever needed (no absolute jumps!)
- Local data: access via offset from current PC, etc.
- External data: indirection through Global Offset Table (GOT)
- ... which in turn is accessed via offset from current PC

Static linking

Dynamic linking