

CS1571
Fall 2019
10/7 In-Class Worksheet

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Where were you sitting in class today: Back left

A. Pre-Reflection

On a scale of 1-5, with 5 being most confident, how well do you think you could execute these learning objectives:

- 10.1 Compare model checking, theorem proving, and resolution as ways of inferring new knowledge from a knowledge base. 2
- 10.2 Prove that a KB entails a sentence *alpha* through logical inference rules. 2
- 10.3 Prove that a KB entails a sentence *alpha* through resolution. 2

B. Model Checking

1. Is Model Checking sound? T
2. Is Model Checking complete? T
3. What is the time complexity of Model Checking for determining entailment.
 - a. **Exponential - $O(c^n)$**
 - b. Quadratic - $O(n^2)$
 - c. Linear - $O(n)$

C. Theorem Proving

The following are a list of some logical equivalencies (from the text) and inference rules.

$$\begin{aligned}
 (\alpha \wedge \beta) &\equiv (\beta \wedge \alpha) && \text{commutativity of } \wedge \\
 (\alpha \vee \beta) &\equiv (\beta \vee \alpha) && \text{commutativity of } \vee \\
 ((\alpha \wedge \beta) \wedge \gamma) &\equiv (\alpha \wedge (\beta \wedge \gamma)) && \text{associativity of } \wedge \\
 ((\alpha \vee \beta) \vee \gamma) &\equiv (\alpha \vee (\beta \vee \gamma)) && \text{associativity of } \vee \\
 \neg(\neg\alpha) &\equiv \alpha && \text{double-negation elimination} \\
 (\alpha \Rightarrow \beta) &\equiv (\neg\beta \Rightarrow \neg\alpha) && \text{contraposition} \\
 (\alpha \Rightarrow \beta) &\equiv (\neg\alpha \vee \beta) && \text{implication elimination} \\
 (\alpha \Leftrightarrow \beta) &\equiv ((\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)) && \text{biconditional elimination} \\
 \neg(\alpha \wedge \beta) &\equiv (\neg\alpha \vee \neg\beta) && \text{de Morgan} \\
 \neg(\alpha \vee \beta) &\equiv (\neg\alpha \wedge \neg\beta) && \text{de Morgan} \\
 (\alpha \wedge (\beta \vee \gamma)) &\equiv ((\alpha \wedge \beta) \vee (\alpha \wedge \gamma)) && \text{distributivity of } \wedge \text{ over } \vee \\
 (\alpha \vee (\beta \wedge \gamma)) &\equiv ((\alpha \vee \beta) \wedge (\alpha \vee \gamma)) && \text{distributivity of } \vee \text{ over } \wedge
 \end{aligned}$$

Modus Ponens

(if $(P \Rightarrow Q)$ is true, and P is true, then we can infer that Q is true)

And Elimination

(if $(P \wedge Q)$ is true, then we can infer that P is true)

And Introduction

(if P is true, and Q is true, then we can infer that $P \wedge Q$ is true)

Or Introduction

(if P is true, then we can infer that $P \vee Q$ is true)

KB is: $P \wedge Q$, $P \Rightarrow R$, $Q \wedge R \Rightarrow S$

Does $KB \models S$? Write your proof in the table below. For each line of the proof, indicate what rule you're using to introduce a new sentence into the KB.

Sentence #	Sentence	Rule & Sentences Used
1	$P \wedge Q$ (if $P \wedge Q$ is true, P & Q is true)	N/A
2	$P \Rightarrow R$ (P is true can infer R is true)	N/A
3	$Q \wedge R \Rightarrow S$ (Q is true and R is true so S is true)	N/A
4	P is true	1, And Elimination
5	Q is true	1, And Elimination
6	R is true	4, 2 Knowledge base
7	$Q \wedge R$ is true	5, 6 And introduction
8	S is true	7, Modus Ponens

4. Is Theorem Proving sound?

Yes if the logical inference rules will give you correct truth values

5. Is Theorem Proving complete?

N but only if your set of inference rules are complete

6. What is the time complexity of Theorem Proving for determining entailment.

- a. **Exponential**
- b. Quadratic
- c. Linear
- d. Other

D. Theorem Proving with Resolution

Let's go back to the Wumpus world. Suppose our KB is currently:

$$(B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})) \wedge \neg B_{1,1}$$

We want to prove that $KB \models \alpha$, where α says "there's no pit in Room 1,2" ($\neg P_{1,2}$)

First, figure out the CNF for:

$$(B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})) \wedge \neg B_{1,1}$$

Step 1: Remove the biconditional using biconditional elimination.

Step 2: Remove the implication using implication elimination.

Step 3: Move \neg inwards using de Morgan's rules and double-negation ($\neg(\neg P) \Leftrightarrow P$):

Step 4: Apply distributivity law to get conjunctions of disjunctions.

Step 5: To prove by contradiction with resolution, we need to negate the sentence we want to prove and add it to the KB. So $KB \wedge \neg\alpha$ is:

Step 6: Apply the resolution rule to create new clauses until you get a contradiction. You can use the table to track your work.

Clause #	Clause	Clauses Used
1	$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1})$	
2	$(\neg P_{1,2} \vee B_{1,1})$	
3	$(\neg P_{2,1} \vee B_{1,1})$	
4	$\neg B_{1,1}$	
5	$P_{1,2}$	
6		
7		
8		
9		
10		
11		

E. Post-Reflection

On a scale of 1-5, with 5 being most confident, how well do you think you could execute these learning objectives:

10.1 Compare model checking, theorem proving, and resolution as ways of inferring new knowledge from a knowledge base. _____

10.2 Prove that a KB entails a sentence alpha through logical inference rules. _____

10.3 Prove that a KB entails a sentence alpha through resolution. _____