

Adaptive Control-Based Clock Synchronization in Wireless Sensor Networks

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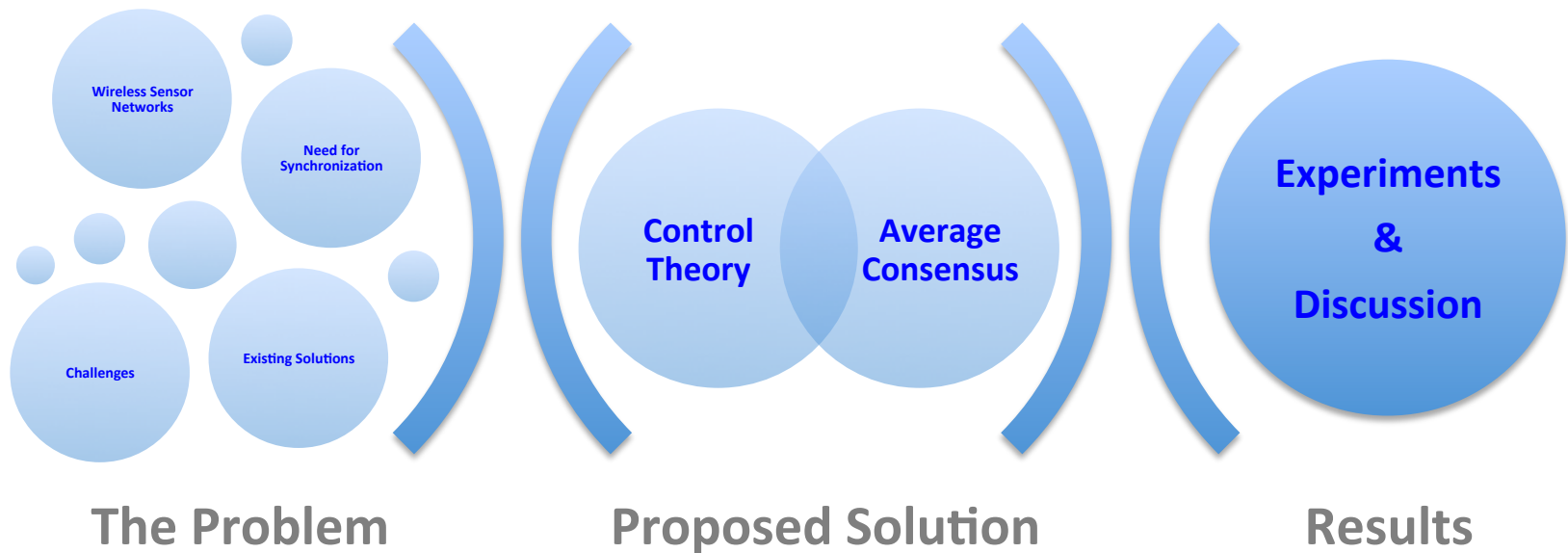


Design and Analysis of Communication Systems (DACs)

University of Twente

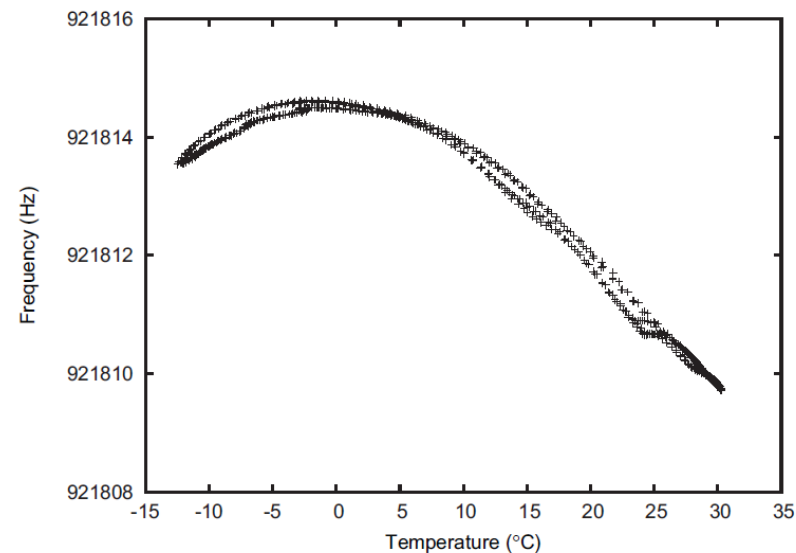
March 19, 2015

Outline



Wireless Sensor Networks

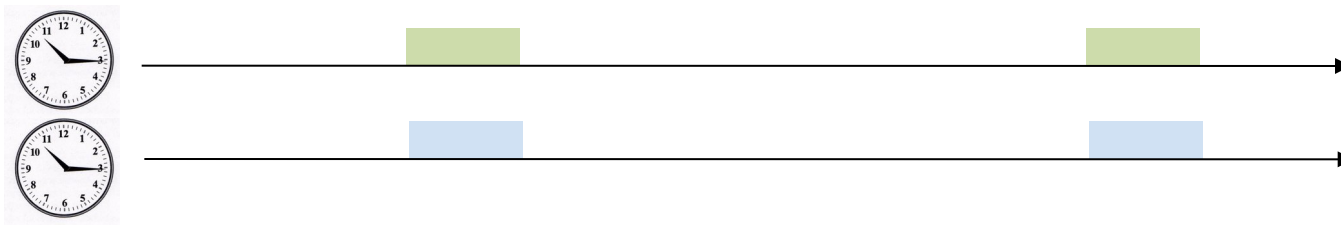
- **Hardware clock** (built-in clock)
 - Counter register
 - Clocked with external crystal
 - 32kHz, 7.37 MHz nominal rate
 - Read-only
- **Clock drift**
 - Deviation from the nominal rate
 - 30-100 parts per million (ppm)
 - dependent on environmental factors such as:
 - temperature, power supply, aging...



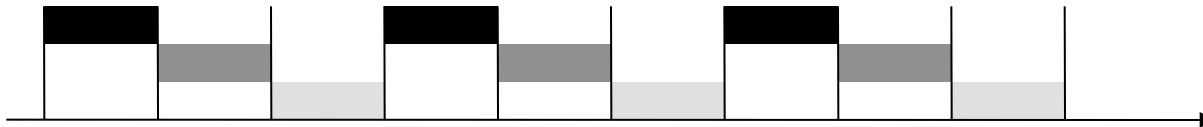
[Sommer et al. 2009]

Need for Synchronization

- Assigning **global timestamps**
 - sensed data and events
- **Coordinated actions**
 - Duty-cycling of the nodes for energy efficient operation



- Low-power, **TDMA-based** MAC layer

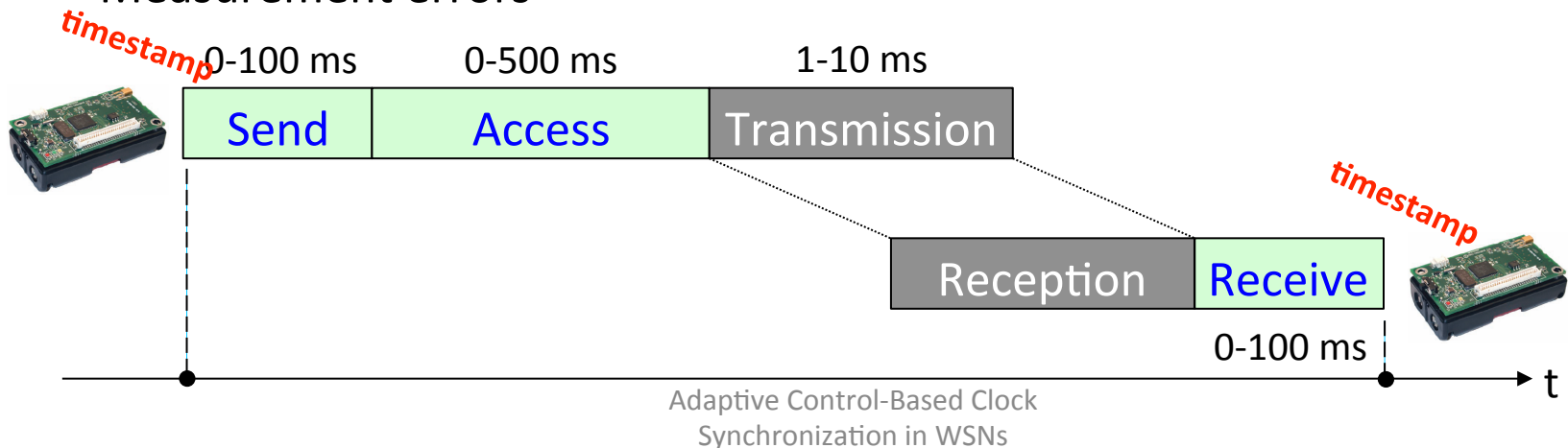
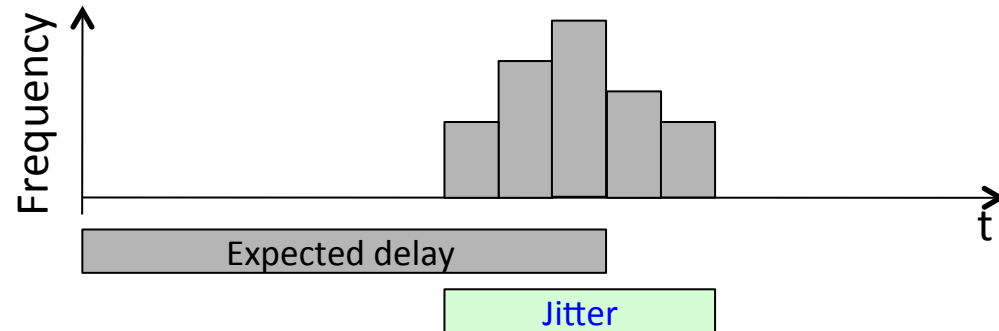


- **Applications**
 - Localization via time-of-flight measurements
 - Tracking of moving objects

Time Synchronization

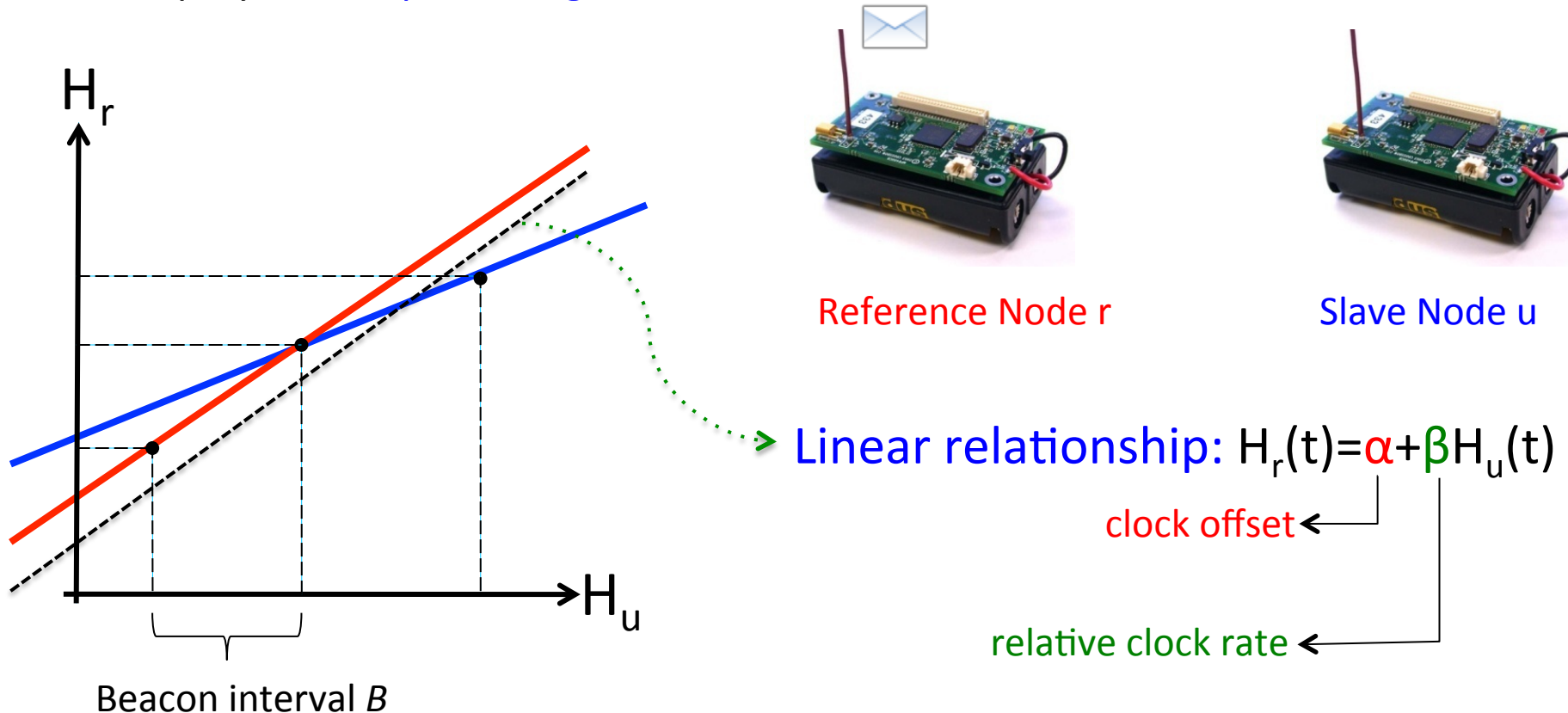
- Communicate
 - Exchange current time information periodically
 - Hardware clock value
- Compute
 - Calculate a **logical clock**
 - Software function
 - represents the global time
- Challenge
 - Message delay
 - Deterministic and **non-deterministic** components
 - Measurement errors

MAC Layer timestamping



Pairwise Synchronization in Practice

- Master – Slave Synchronization
 - Collect (H_u, H_r) timestamp pairs periodically
 - Employ **least-squares regression**



Least-Squares with Two Pairs

Measurement errors

Slope: $\hat{\beta} = \frac{H_r(t_1) - H_r(t_0)}{H_u(t_1) - H_u(t_0)}$

Poor multi-hop
performance
scalability

cont. by
sion delays

Error
contributed by
the transmission
delays

Measure

Measurement errors

Intercept:

$$\hat{\alpha} = \frac{H_r(t_1) + H_r(t_0)}{2}$$


Errors enter non-linearly to the equations!

$$\hat{H}_r(t) = \hat{\alpha} + \hat{\beta}H_u(t)$$

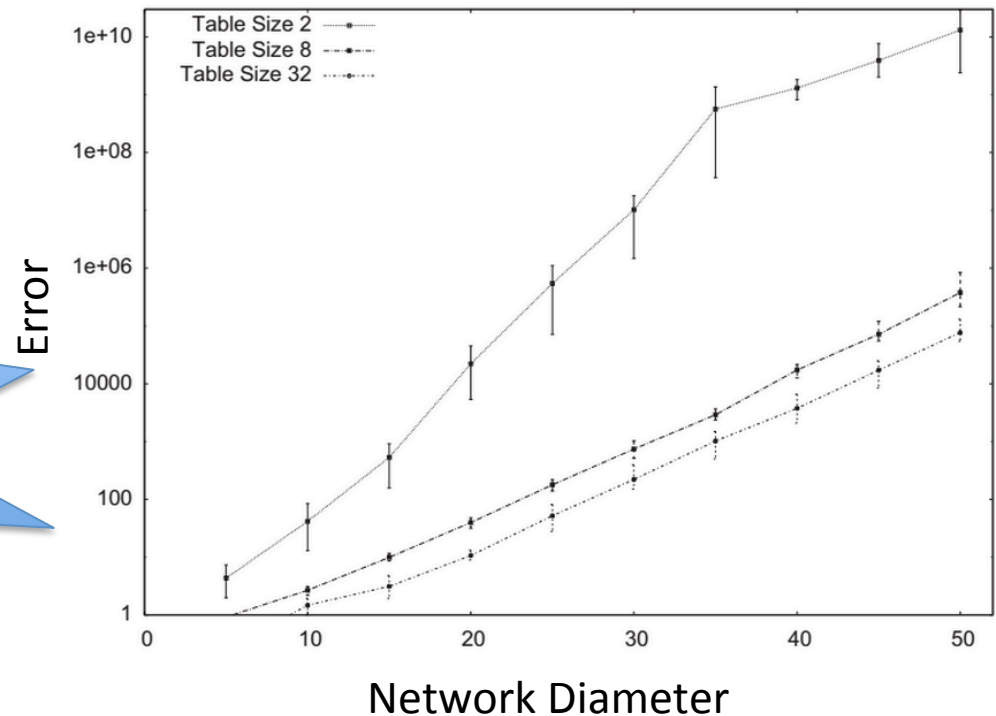
The effect of various error sources appear as **multiplicative noise!**

Multi-hop Performance of Least-Squares

- Simulation of multi-hop least-squares based time synchronization with different regression tables sizes.
 - Amplification of errors at each hop
 - **Multiplicative noise!**

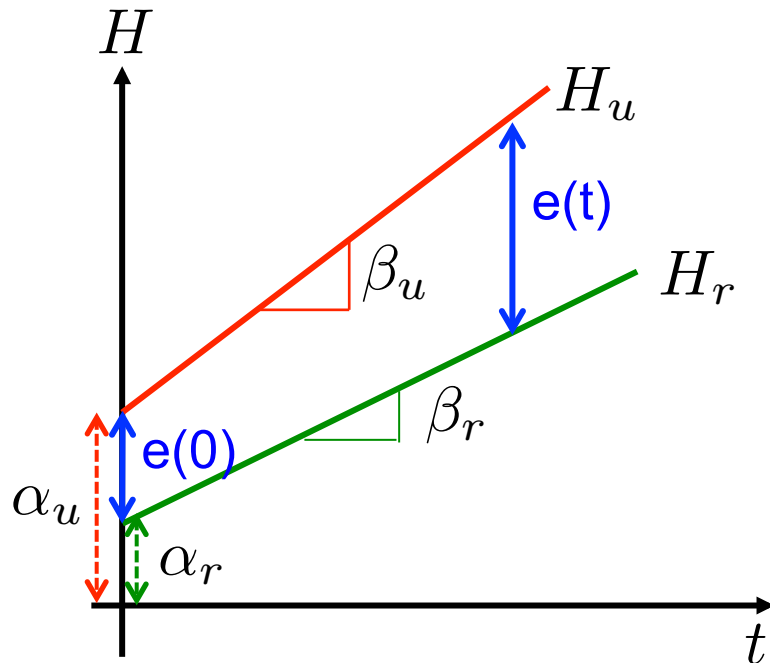


Exponentially
Increasing
Synchronization
Error!



[Lenzen et al. 2009]

Time Synchronization as a Control Problem

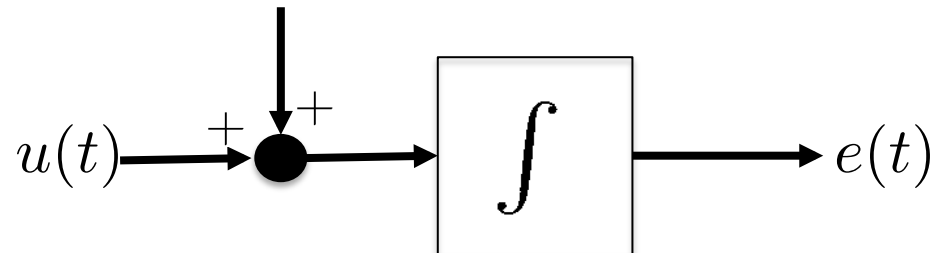


$$e(t) = H_u(t) - H_r(t)$$

$$e(t + B) = e(t) + \underbrace{(\beta_u - \beta_r)B}_{\text{Speed difference}} \underbrace{1}_{\text{input}}$$

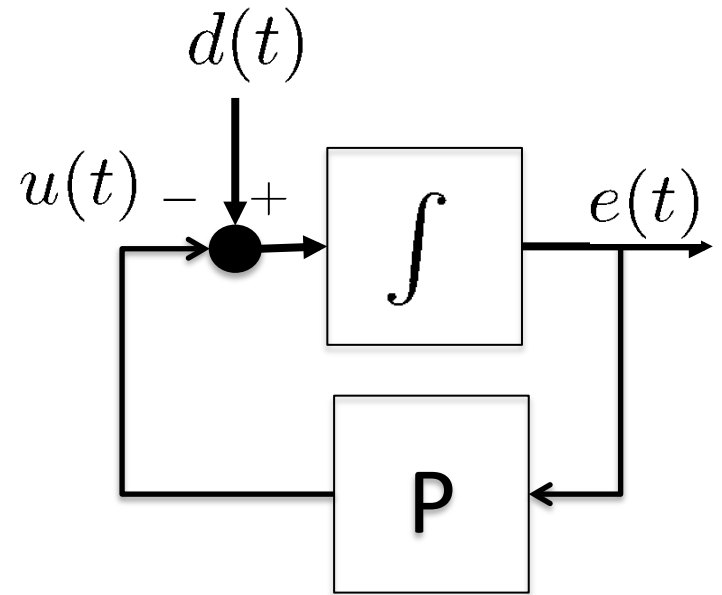
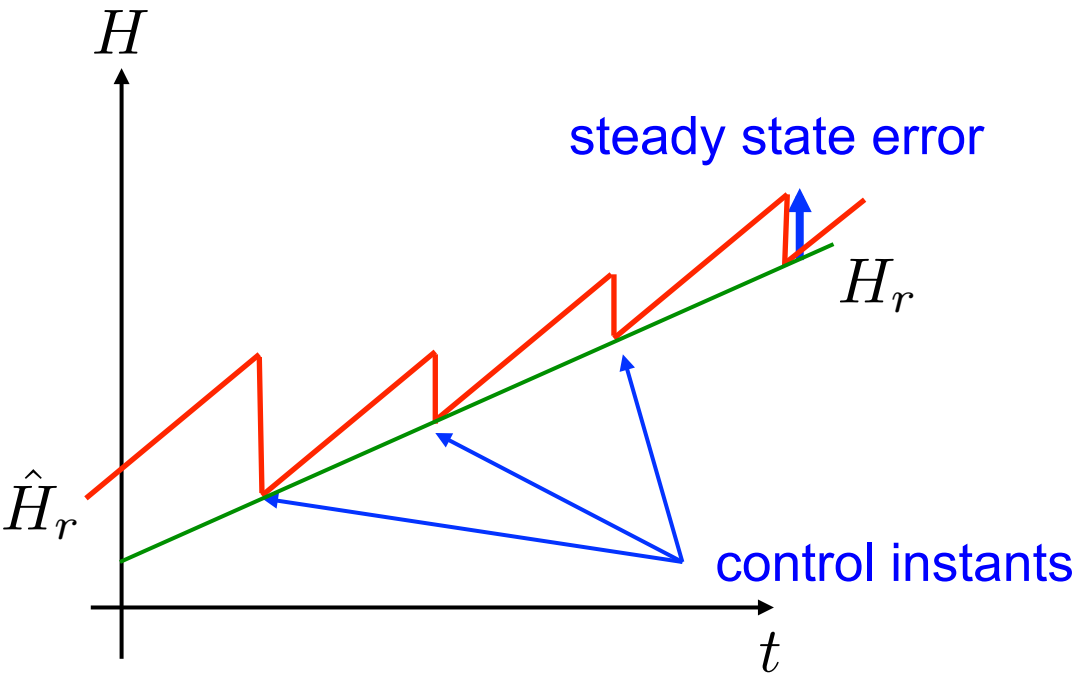
constant disturbance

$$d(t) = (\beta_u - \beta_r)$$



Proportional Control

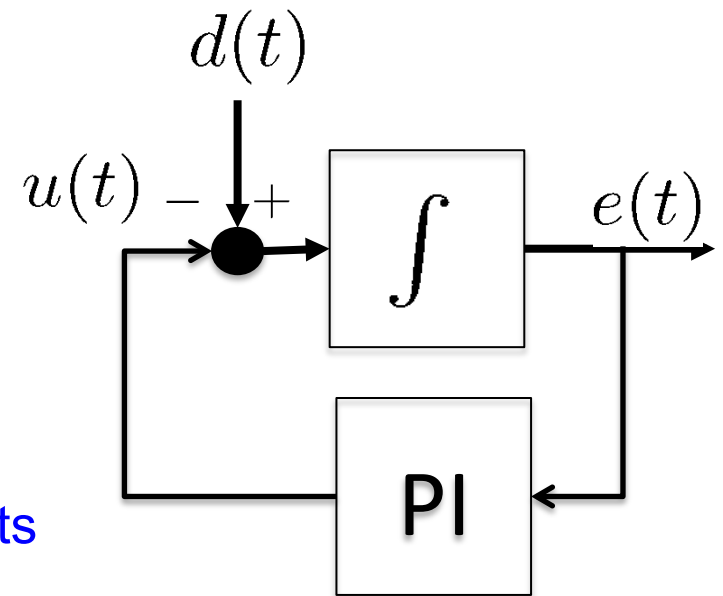
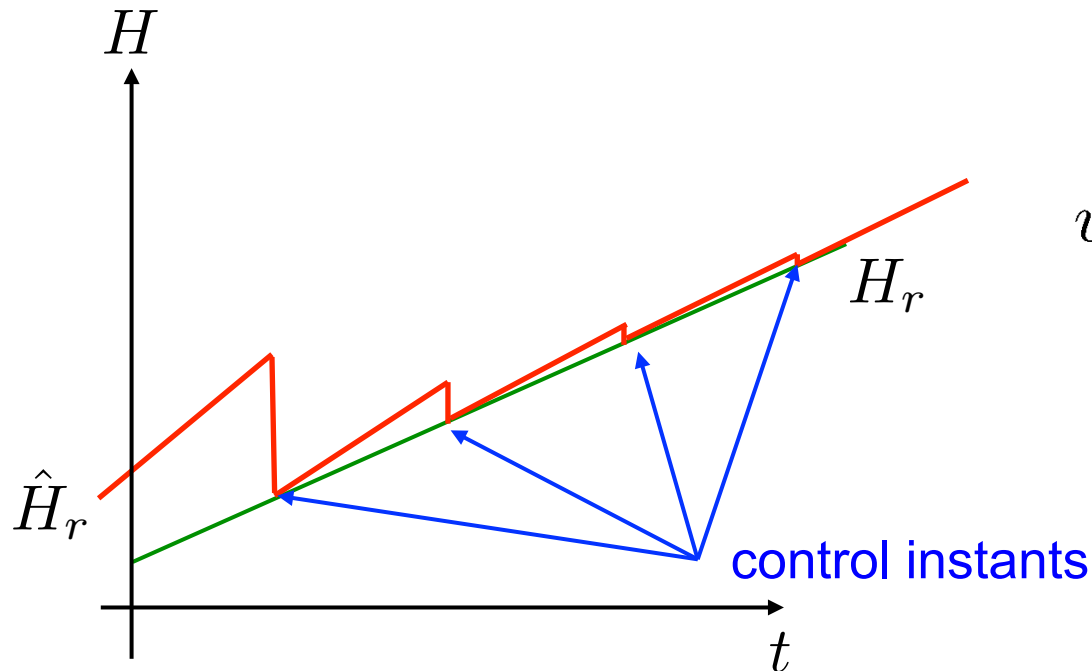
$$u(t) = -k_P e(t)$$



Compensates only initial offset differences

Proportional-Integral Control

$$u(t) = -k_P e(t) - k_I \int e(t) dt$$



Compensates both offset and clock speed differences

PI Sync Algorithm - Pairwise Synchronization

Reference Clock Estimate

$$\hat{H}_r(t) = \underbrace{\hat{H}_r(t_{up})}_{\text{Offset}} + \underbrace{\Delta(t_{up})}_{\text{Speed}} \underbrace{(H_u(t) - H_u(t_{up}))}_{\text{Local time passed since update}}$$

Update Rule

- Receive $\langle H_r(t_{up}), t_{up} \rangle$
- Calculate Error
- Update Offset
- Update Speed

No need to store received timestamps!
 No regression table!
 Very simple arithmetic operations!

$$\begin{aligned} e(t_{up}) &= H_r(t_{up}) - H_u(t_{up}) \\ \Delta(t_{up}) &= k_P e(t_{up}) \\ \hat{H}_r(t_{up}) &= \hat{H}_r(t_{up}) - k_I e(t_{up}) \end{aligned}$$

Convergence Conditions

$$\begin{aligned} 0 < k_P < 2 \\ 0 < k_I < \frac{2(2-k_P)}{f_{nom}B} \end{aligned}$$

Optimal Convergence Rate

$$k_I^* = \frac{1}{f_{nom}B} \text{ for } k_P = 1$$

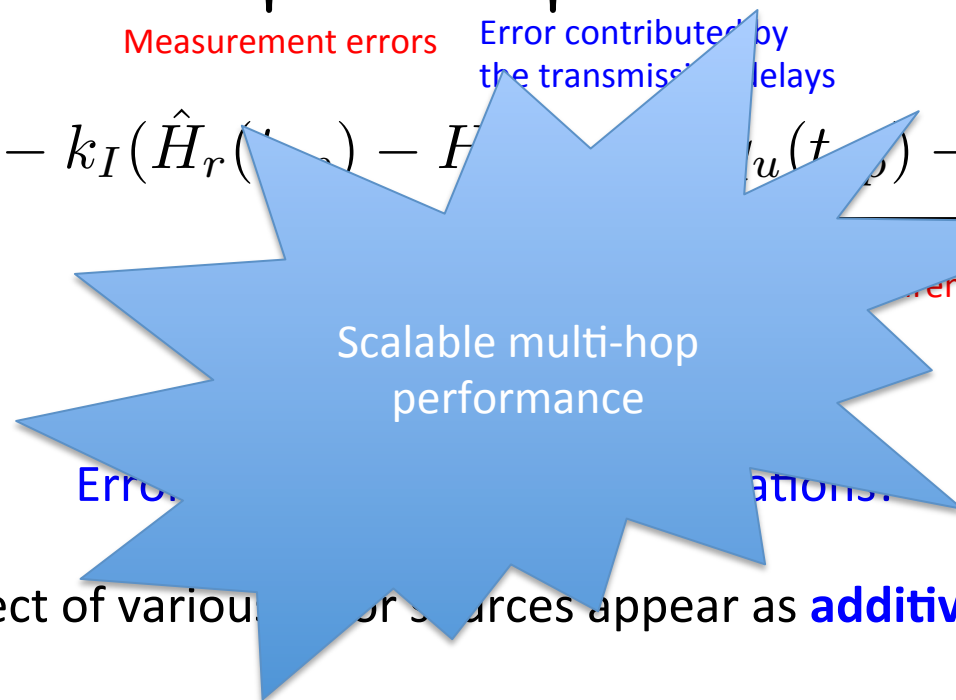
Not the smallest steady-state error!
 Integral gain should be adjusted adaptively!

PISync Algorithm – Error Dynamics

For $k_p=1$, PSync algorithm becomes

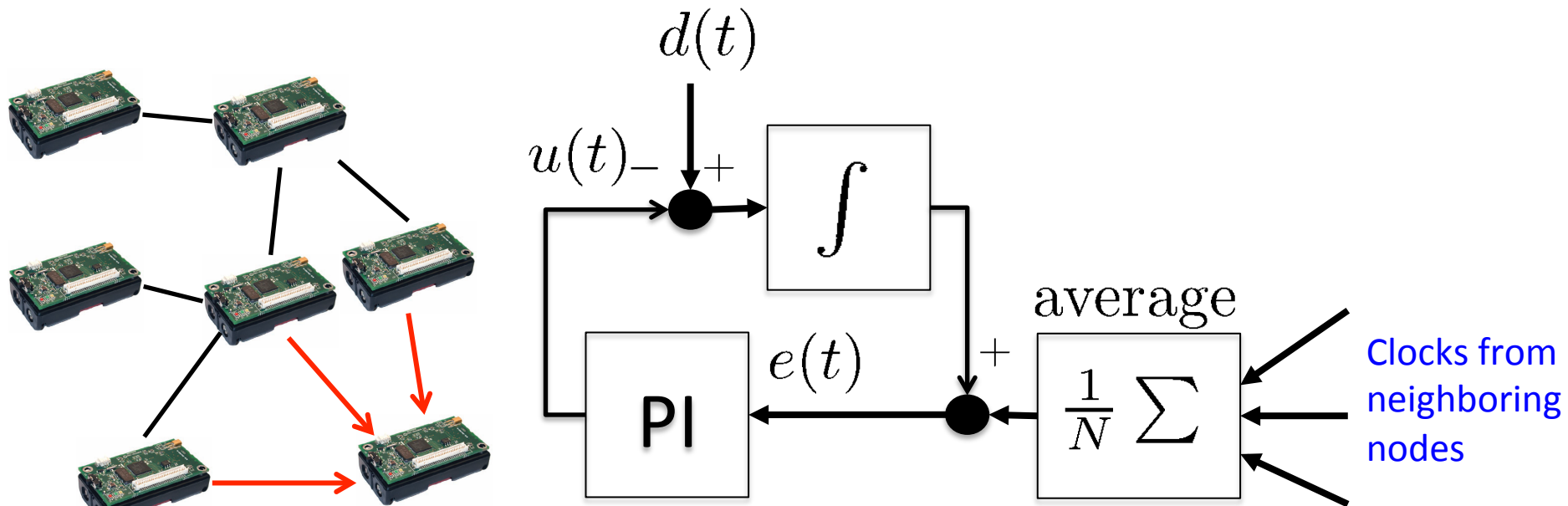
$$\hat{H}_r(t_{up}^+) = H_r(t_{up}) + \underbrace{q_r(t_{up})}_{\text{Measurement errors}} + \underbrace{\varepsilon(t_{up})}_{\text{Error contributed by the transmission delays}}$$

$$\Delta(t_{up}^+) = \Delta(t_{up}) - k_I \left(\hat{H}_r(t_{up}^+) - H_r(t_{up}) - \underbrace{q_r(t_{up}) + \varepsilon(t_{up})}_{\text{Measurement errors}} \right)$$



The effect of various error sources appear as **additive noise!**

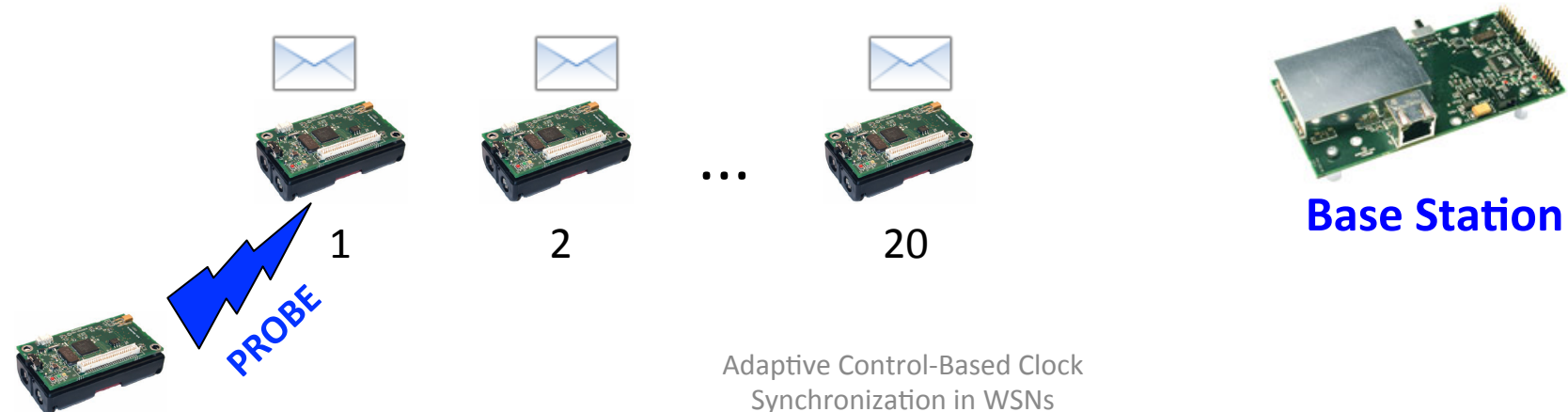
PI control + Average Consensus



No special reference node!
Nodes synchronize directly to their neighbors!

Testbed

- Testbed setup
 - 20 MICAz sensor nodes
 - FTSP (Flooding Time Synchronization Protocol [\[Marotti et al. 2004\]](#)) and PISync implemented in TinyOS 2.1
- Network topology
 - Single-hop setup
 - Nodes are communication range of each other
 - Virtual [line](#) and [grid](#) topologies (in software)



Flooding Time Synchronization Protocol

Marotti et al. 2004

- In every 30 seconds, each node
 - Receives a new reference clock estimate
 - Stores the timestamp pair on the regression table of size 8
 - Employs least-squares regression on the pairs in the regression table
 - Forwards its estimate of the reference clock
 - each estimate is affected by
 - Message delays
 - Measurement errors
 - Estimation errors

reference



1



2

...



20

PISync - Integral Gain Adaptation

$$k_I^* = \frac{1}{f_{nom} B} \text{ for } k_P = 1$$

As k_I gets smaller, we obtain smaller steady-state error but longer convergence time!

$$e_{max} = 2B \underbrace{\Delta f}_{\text{Maximum frequency error}} / f_{nom}$$

Maximum frequency
error

Adaptation Rule

If $e(h) > e_{max}$ then $\alpha = 0$

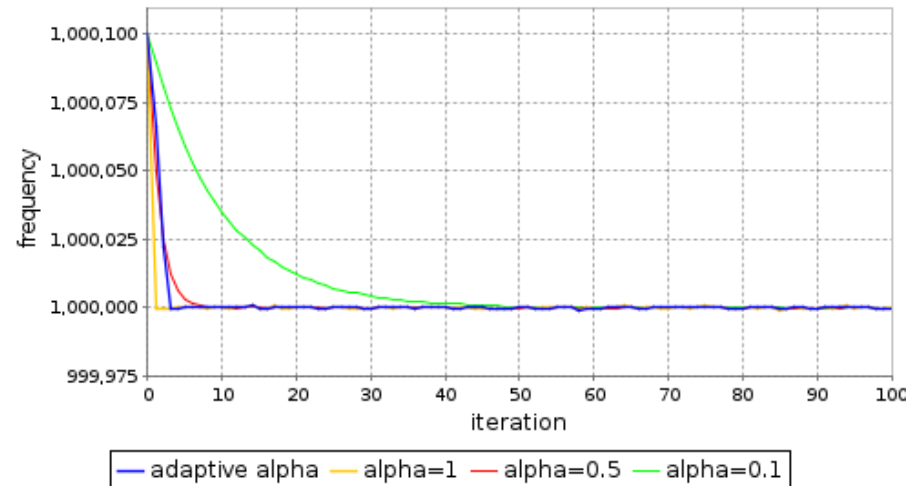
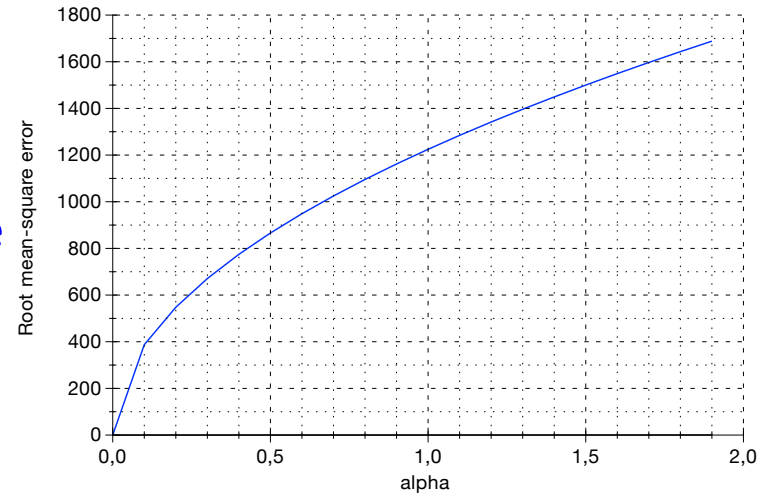
Else $k_I =$

Else $k_I = \max\{2k_I, k_I^*\}$

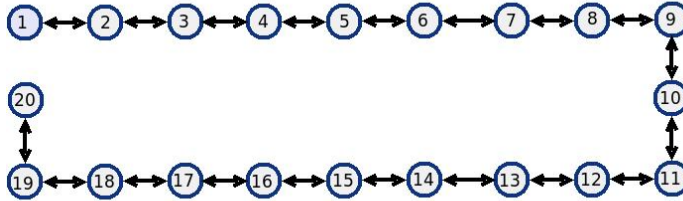
Else $k_I =$

Gradually increase
or decrease

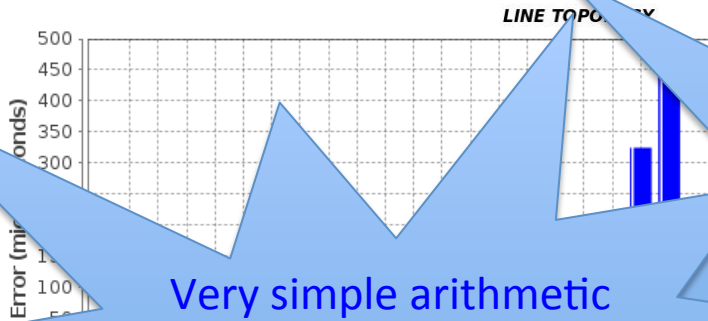
Inspired from [1] and [2] (Turcan et al. 2014)



FTSP vs PISync



FTSP

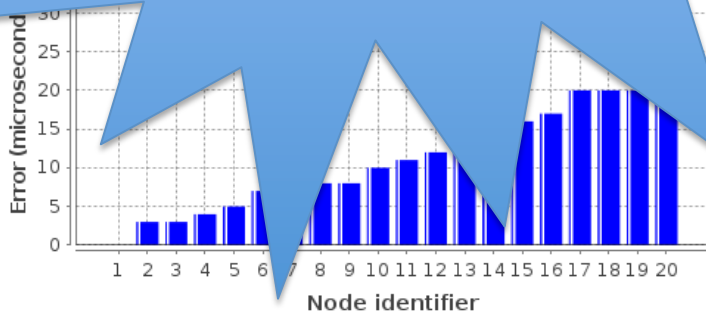


Very simple arithmetic operations

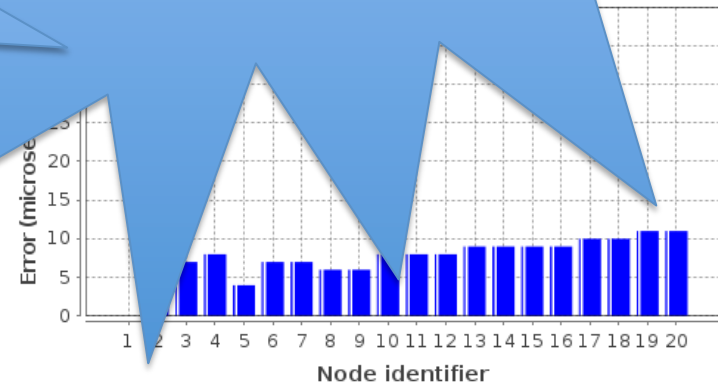
No need to store time information explicitly!

Scalable!

FloodPISync



AvgPISync



Completely blind operation!

No need to store neighbors time information explicitly!

Suitable for mobile networks!

Robust to packet losses

PISync Complexity

```
void synchronize(TimeSyncMsg *msg)
{
    int32_t timeError;
    float newSkew = skew;

    /* calculate offset difference */
    timeError = msg->localTime;
    call GlobalTime.local2Global((uint32_t *)&timeError);
    timeError = msg->globalTime - timeError;

    /* adjust the speed of the logical clock */
    if( timeError < E_MAX && timeError > -E_MAX )
    {
        /* calculate adaptive alpha */
        if(lastError != 0 && currentAlpha != 0)
            currentAlpha *= ((float)timeError / lastError);

        currentAlpha = fabs(currentAlpha);

        if(currentAlpha > ALPHA_MAX)
            currentAlpha = ALPHA_MAX;
        /* adjust rate multiplier */
        newSkew += currentAlpha * timeError;
    }

    lastError = timeError;

    /* update logical clock parameters */
    atomic{
        skew = newSkew;
        clock = msg->globalTime;
        lastUpdate = msg->localTime;
    }
}
```

Very easy to implement & code

20 times better performance!
50 times fewer operations!
4 times less RAM requirements!
Low power consumption!

	PISync
	20 microseconds
	16 bytes
5.5 ms	15432 bytes
	145 microseconds

Just 15 Lines of TinyOS code!

Conclusions

- We considered time synchronization as a control problem
- We solved this problem with a very simple and practical technique
 - Proportional-Integral Controller
- We introduced a new time synchronization protocol
 - PISync
- We observed
 - Better performance scalability
 - Less resource consumption
 - Lower CPU and main memory overhead
 - Lower power consumption

Future Research Directions

- Performance evaluation on real mobile networks
- Implementation & evaluation of time-synchronized MAC layer
 - Adaptive receiving window
 - Evaluation of power consumption
- Other algorithmic techniques?
 - Better steady-state error & convergence
 - Even lower resource requirements?