# संगणक विज्ञान एवं अभियांत्रि की विभाग



## Computer Science & Engineering Department

National Institute of Technology Sikkim, Ravangla, South Sikkim-737139 (AN INSTITUTE OF NATIONAL IMPORTANCE, GOVT. OF INDIA)

## **Second Online Examination (Type I)**

Marks: 10 Time: 45 minutes

Advance Algorithm (CS15102) for B.Tech (5th Semester), CSE

### **Instructions:**

- Online Google form: You have to write the answer in the online google form. Each online submitted answer is limited to maximum 20 words only. Submit the google form by 11:45AM.
- Submission of justification/rough works:
  - Take A4/blank paper and write your Name, Roll No., Subject Name with Subject Code and Date at the top of the first paper.
  - You have to write justification behind your answer or a rough work for each answer. **The submitted answer in the google form without justification will not be considered.**
  - Sign at the end of each page.
  - Submit the **scan/photocopy** of the justification/rough works through **google class**.
  - Submit the rough works by 12:00Noon.

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#### **Answer all the questions:**

- 1. Assume a *dynamic table* where size of the table is dynamically changed as per requirement.
  - TABLE-INSERT inserts into the table an item that occupies a single slot, that is, a space for one item. *Expansion required for full table*.
  - TABLE-DELETE removes an item from the table, thereby freeing a slot. *Contraction required for sufficiently empty table*.

To analyze the amortized cost of the  $i^{th}$  TABLE-INSERT and TABLE-DELETE operation, let

- $num_i$  denote the number of items stored in the table after the  $i^{th}$  operation.
- $size_i$  denote the total size of the table after the  $i^{th}$  operation, and
- $\Phi_i$  denote the potential after the  $i^{th}$  operation.
- $\alpha(T)$  be the load factor of the table T.

The potential function for such operation is designed as follow.

Potential function 
$$\Phi(T) = \begin{cases} num[T] - \frac{size[T]}{2}, & \text{If } \alpha(T) \ge \frac{1}{2} \\ \frac{size[T]}{2} - num[T], & \text{If } \alpha(T) < \frac{1}{2} \end{cases}$$

What is the amortized cost of  $i^{\text{th}}$  TABLE-INSERT operation when  $\alpha_{i-1} < \frac{1}{2}$  and  $\alpha_i \ge \frac{1}{2}$ ?

- 2. Consider the Q1. If  $\alpha_{i-1} \ge \frac{1}{2}$ , the amortized cost of the TABLE-INSERT operation is at most \_\_\_\_.
- 3. Consider the Q1. Assume,  $\alpha_{i-1} < \frac{1}{2}$  and the  $i^{th}$  TABLE-DELETE operation does trigger a contraction. What is the amortized cost of such  $i^{th}$  operation?
- 4. Use Masters theorem:  $T(n) = 6T (n/3) + n^2 \log n$
- 5. An algorithm has time complexity  $T(n) = 2n^3 + 33n^2 + 64$ . Asymptotically analyze it based on Theta  $(\Theta)$ .

- 6. An algorithm has time complexity  $T(n) = 2n^2 + 25n + 5n\log n$ . Asymptotically analyze it based on Big-omega ( $\Omega$ ).
- 7. Consider the following recursive function.

```
fun (n)

{

If (n \le 0) then
Print "NIT"

Else

{

x = fun(n-1) + fun(n-1)
Print "Sikkim"

}
```

What is the time complexity of fun(n)?

- 8. The problem of determining whether any element from a given list greater than *k* (for a given *k*) exists or not is in NP (Write *True* or *False*).
- 9. 3SAT is NP-Complete problem. If 3SAT is polynomial time reducible to the problem *B* then *B* is NP-Complete (Write *True* or *False*).
- 10. Let T(n) be a polynomial and O(T(n)) be the time complexity of non-deterministic Turing machine for the given input of length n. What is the time complexity of the equivalent deterministic Turing machine for the same input?

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