

# **RECURSION**

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# ASSUMPTIONS

## ❖ familiarity with

- simple Python programs
  - including function calls, and
  - list comprehensions
- tracing programs
- activation record (a.k.a. stack frame)
  - and its application in tracing function calls
- PyCharm IDE

# objectives

- ❖ introduce recursive functions/programs
- ❖ describe how recursive programs work
- ❖ compare recursion vs iteration

# before we dive into it

- ❖ **recursion** occurs when a thing is defined in terms of itself (wikipedia)
- ❖ **recur** means to come up again (merriam-webster)

# recursive examples

## ❖ factorial function

$$\begin{cases} \text{factorial}(n) = n * \text{factorial}(n-1) \\ \text{factorial}(0) = 1 \end{cases}$$

**recursive case**

**base case**

## ❖ fibonacci function

$$\begin{cases} \text{fibonacci}(n) = \text{fibonacci}(n-1) + \text{fibonacci}(n-2) \\ \text{fibonacci}(1) = 1 \\ \text{fibonacci}(0) = 1 \end{cases}$$

**recursive cases**

**base cases**

A recursive function has  
at least one **base case** and at least one **recursive case**

# another example

a recursive definition: **balanced\_string**

❖ **recursive cases:**

- (x) is balanced if x is a **balanced\_string**
- xy is balanced if x and y are **balanced\_string**

❖ **base case:**

- a string containing no parentheses is balanced

# how about these functions?

- ❖  $f(n) = n^2 + n - 1$
- ❖  $f(n) = g(n-1) + 1, \quad g(n) = n/2$
- ❖  $f(n) = 5, \quad f(n-1) = 4$
- ❖  $f(n) = n * (n-1) * (n-2) * \dots * 2 * 1$
- ❖  $f(n) = f(\lfloor n/2 \rfloor) + 1, \quad f(1) = 1$
- ❖  $f(n) = f(n-1)$

# recursive programs

- ❖ solution defined in terms of solutions for smaller problems

```
def solve (n):
```

```
    ...
```

```
    value = solve(n-1) + solve(n/2)
```

```
    ...
```

- ❖ one or more base cases

```
    if n < 10:
```

```
        value = 1
```

- ❖ some base case is always reached eventually; otherwise, it's an infinite recursion



# general form of recursion

**if** (condition to detect a base case):

{do something without recursion}

**else:** (general case)

{do something that involves recursive call(s)}



devise it such that some base case is always reached eventually

# recursive programs example 1

$0! = 1$  ,  $n! = n * (n-1)!$

```
def factorial(n):
```

```
    # pre: n ≥ 0
```

```
    # post: returns n!
```

```
    if n==0:  
        return 1
```

```
    else:  
        return n * factorial(n-1)
```

- ❖ structure of code typically parallels structure of definition

# recursive programs example 2

$\text{fib}(0) = 1, \text{fib}(1) = 1,$

$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$

```
def fib(n):
```

```
    # pre: n ≥ 0
```

```
    # post: returns the nth Fibonacci number
```

```
    if n < 2:
```

```
        return 1
```

```
    else:
```

```
        return fib(n-1) + fib(n-2)
```

- ❖ structure of code typically parallels structure of definition

# max\_list() example 3

# tracing factorial: intuitively

**trace** fact (3)

# tracing max\_list()

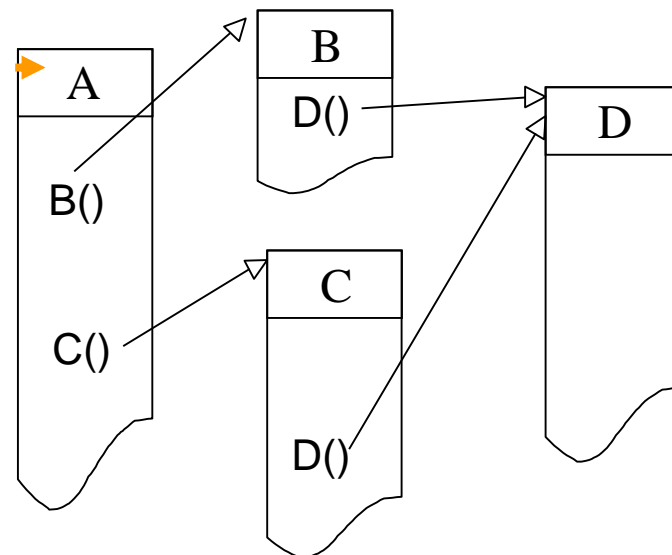
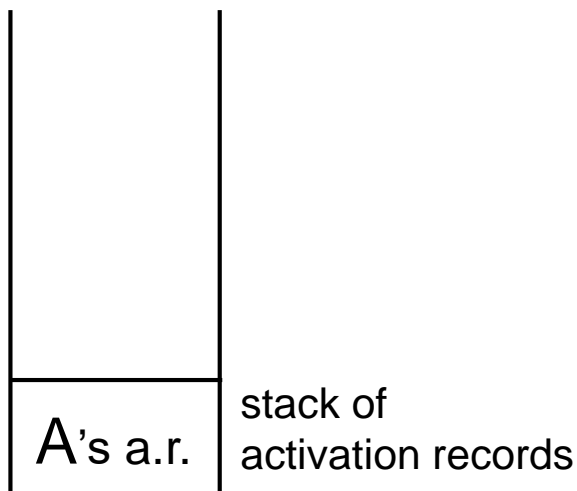
```
trace max_list([4, 2, [[4, 7], 5], 8])
```

# stacks and tracing calls review

- ❖ recall:
  - stack applications in tracing function calls
- ❖ **activation record** (a.k.a stack frame)
  - all information necessary for tracing a function call
  - such as *parameters, local variables, return address*, etc.
- ❖ **when function called:**
  - activation record is created, initialized, and **pushed** onto the stack
- ❖ **when function finishes:**
  - its activation record (that is on top of the stack) is **popped** from the stack

# tracing calls review

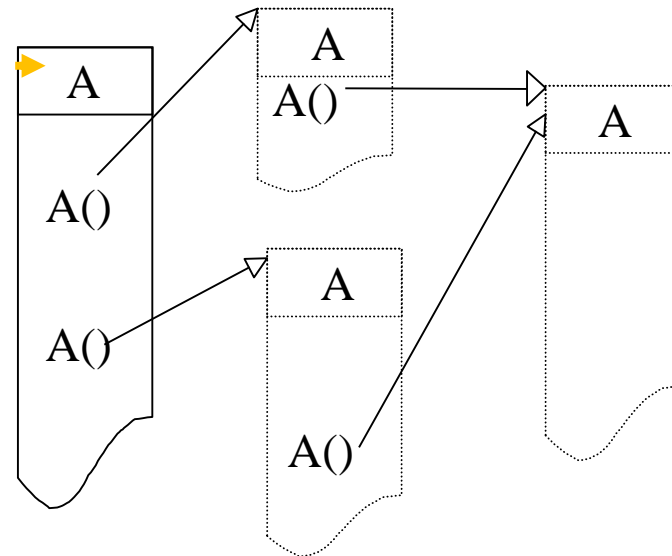
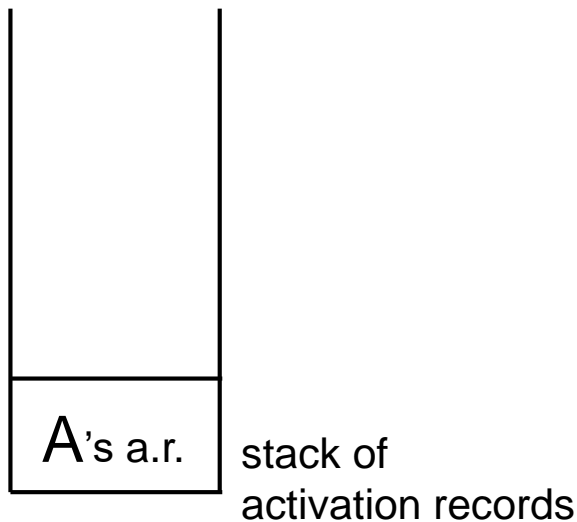
- ❖ recall: stack of activation records
  - **when function called:**
    - activation record created, initialized, and **pushed** onto the stack
  - **when function finishes,**
    - its activation record is **popped**





# tracing recursive calls

- ❖ same mechanism for recursive programs



# tracing factorial

```
1.  def f(n) :  
2.      # pre: n ≥ 0  
3.      # post: returns n!  
4.      if (n==0) : return 1  
5.      else: return n * f(n-1)
```

```
1. → def main() :
```

```
    . . .  
8.    print(f(6))  
    . . .
```

5,f,0	Return 1
5,f,1	Return 1
5,f,2	Return 2
8,m,3	Return 6
-----	

line#	func.	n
-------	-------	---

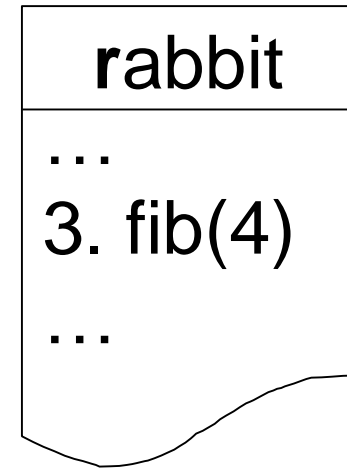
activation record

# tracing fibonacci

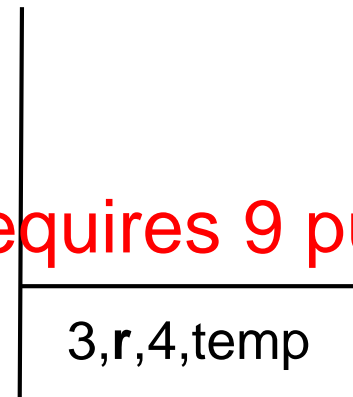
```
1. def fib(n):  
2.     # pre: n ≥ 0  
3.     # post: returns the  
4.     # nth Fibonacci number  
4.     if (n < 2): return 1  
5.     else: return fib(n-1) +  
6.             fib(n-2)
```

line#	func.	n	temp
-------	-------	---	------

activation record



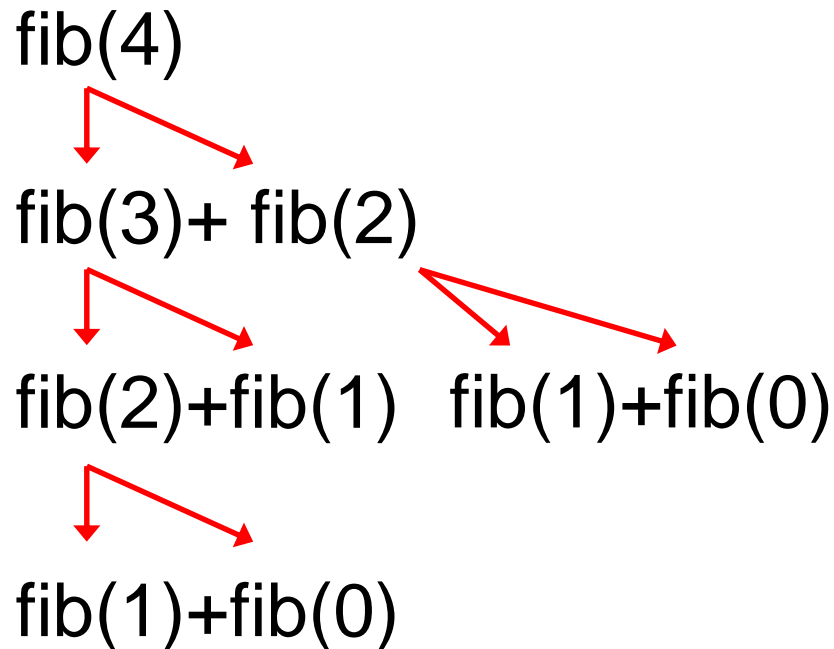
hint: requires 9 pushes



stack of  
activation records

# why 9?

- ❖ using rewriting (aka call tree, intuitively)



# recursive vs iterative

- ❖ recursive functions impose a repeated task (i.e. loop)
  - ❖ the loop is implicit and Python takes care of it
  - ❖ this comes at a price: time & memory
  - ❖ the price may be negligible in many cases
- 
- ❖ after all, no recursive function is more efficient than its iterative equivalent

# recursive vs iterative cont'ed

- ❖ every recursive function can be written iteratively (by explicit loops)
  - may require stacks too
- ❖ yet, when the nature of a problem is recursive, writing it iteratively can be
  - time consuming, and
  - less readable
- ❖ so, recursion is a very powerful technique for problems that are naturally recursive

# more examples

- ❖ balanced strings
- ❖ more functions on nested lists
- ❖ merge sort
- ❖ quick sort
- ❖ traversing trees
  
- ❖ in general,
  - properties of recursive definitions/structures

# summary

❖ **recursion** is

`fun(fun(fun(...fun(this is the base)...)))`

❖ **recursion** is **powerful**

❖ questions on today's lecture?

❖ don't forget to

- `trace fib(4)`
- `trace max_list([4, [[4, 7], 9], 8])`
- implement this lecture's examples

❖ prior to start working on this week's lab