

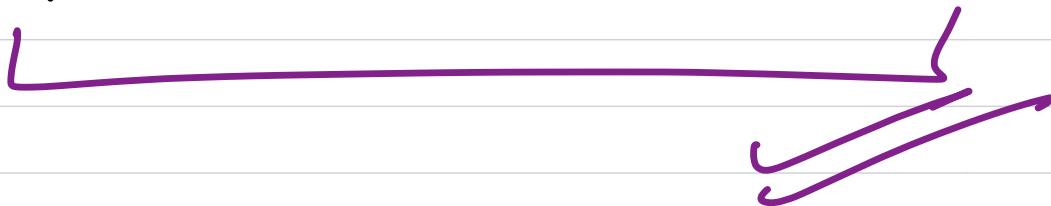
# Recursion

→ Recursion is not just a computer science topic  
It correlates with maths as well:

Mathematically, recursion means a function

calling itself.

$$f(x) = f(f(x'))$$



$$f(a, b) = a \times f(a, b-1)$$

this function

calculates

$$\underline{\underline{a^b}}$$

func<sup>i</sup> definition

$$\boxed{a^b = a \times a^{b-1}}$$

Ex  $f(2, 3) \rightarrow 2^3 \rightarrow 8$

$$2^4 = 2 \times 2^3$$

$f(2, 4) \rightarrow 2^4 \rightarrow 16$

In computer science, recursion means function

calling itself.

function fun(x) {  
...  
...  
...  
fun(x-1);  
...  
}

```
function gem(x) {  
    console.log(x);  
    y  
}
```

```
function fum(x) {  
    ...  
    ...  
    gem(x-1)  
    fum(x-1);  
    ...  
    ...
```



This is not recursion



This is recursion

}

Example-1

Factorial

factorial of any value  $n \Rightarrow n \times (n-1) \times (n-2) \dots \times 2 \times 1$

$\overbrace{\hspace{10em}}$   $\rightarrow n!$   
 $\overbrace{\hspace{1em}}$

$$5! = 5 \times 4 \times 3 \times 2 \times 1 \rightarrow 120$$

$$4! = 4 \times 3 \times 2 \times 1 \rightarrow 24$$

$$3! = 3 \times 2 \times 1 \rightarrow 6$$

$$f(n) = n \times f(n-1)$$

this function

returns  $n!$

The mathematical representation of a recursive function is called Recurrence Relation

$$\begin{aligned}
 5! &\xrightarrow{120} \\
 \downarrow & \quad \downarrow \\
 s \times 4! &\approx 2^4 \\
 \downarrow & \quad \downarrow \\
 4 \times 3! &\approx 6 \\
 \downarrow & \quad \downarrow \\
 3 \times 2! &\approx 2 \\
 \downarrow & \quad \downarrow \\
 2 \times 1! &\approx 1
 \end{aligned}$$

we already  
know them

What is  $11! \Rightarrow 1$

Recursion → In recursion , a function calls itself , where it tries to solve a smaller problem and then calc value of a larger problem

# P M I (principle of Mathematical Induction)

Q. What is the sum of first  $N$  natural numbers ??

$$\text{Ans} \rightarrow \frac{n \times (n+1)}{2}$$

## 3 step process

- 1) Base Case → the smallest value for which we can directly verify the ans-
- 2) Assumption
- 3) Verification / self task

for  $n=1$   $\rightarrow \frac{1 \times (n+1)}{2} \rightarrow$  is correct  
Base Case

for some  $n=k \rightarrow$  the formula is correct

i.e.  $\frac{k \times (k+1)}{2}$

Let's prove ourselves that formula is correct

for  $n = k+1$

$$1 + 2 + 3 + \dots + k + k + 1$$


$$\rightarrow \frac{k \times (k+1)}{2} + (k+1)$$

$$\frac{k \times (k+1) + 2(k+1)}{2} \rightarrow \frac{(k+1)(k+2)}{2}$$

By formula

$$\rightarrow n = (k+1)$$

$$\frac{n \times (n+1)}{2} \rightarrow \frac{(k+1)(k+1+1)}{2}$$

$$= \frac{(k+1)(k+2)}{2}$$

# # Components Of Recursion

- 1) Base Case
- 2) Assumption
- 3) Self work

Write a recursive sol' for Factorial of  $n$

We will write a func<sup>i</sup>

$f(n)$



returns  $n!$

Base Case  $\rightarrow (n == 1)$   $f(1) \Rightarrow 1$   
if  $(n == 1)$  return 1

Assumption  $\rightarrow$  assume the function works  
fine for  $n-1$

$f(n-1)$  correctly calls  $(n-1)!$

~~Self work~~  $\rightarrow f(n) = n \times f(n-1)$

Power function  $\underline{\underline{a^b}}$

$f(a, b)$



returns  $\underline{\underline{a^b}}$

Base Case → if ( $b == 0$ ) then ans is 1

Assumption → assume func<sup>n</sup> works fine for  $b-1$   
i.e.  $f(a, b-1)$  is correct.

Self work  $\rightarrow$   $a^b = \underbrace{a \times a^{b-1}}$

$$f(a, b) \leftarrow a \times f(a, b-1)$$

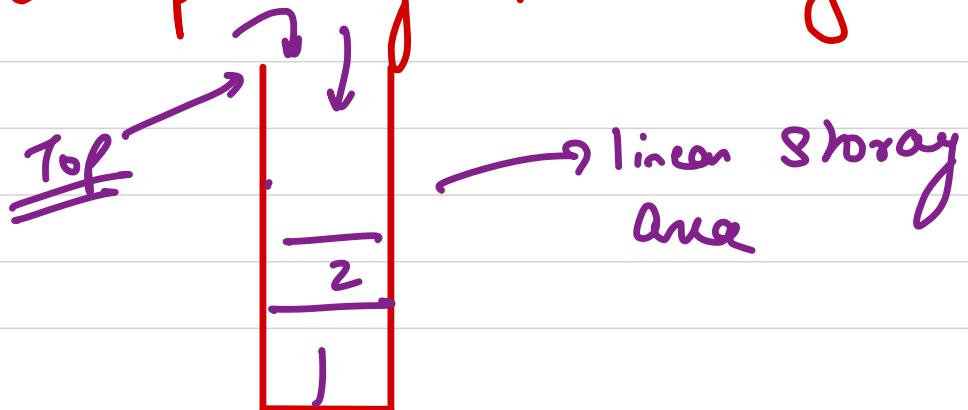
↓  
Conut

$$\underline{\underline{27}}^{51} \quad \underline{\underline{27}} \times \underline{\underline{27}}^{50} \Rightarrow y$$

Q What happens when we call a function ??

In our memory, we have a lot of things going on.

One part of the memory is call stack



1 function fun(x)

2 let a = x + 0;

3 return a;

4 }

5 function gen(y)

6 let b = y + 20

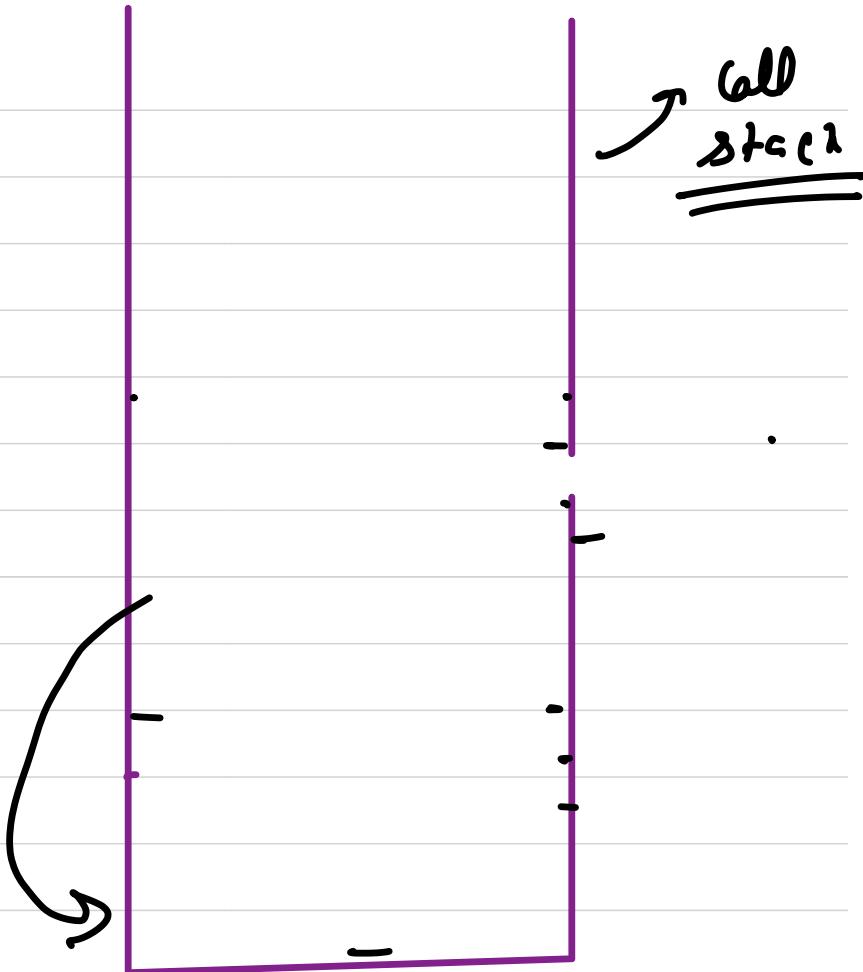
7 let z = fun(b);

8 return z;

9 }

10 console.log(gen(10))

→ call stack



Whenever we call a new function it adds a new entry in the stack.

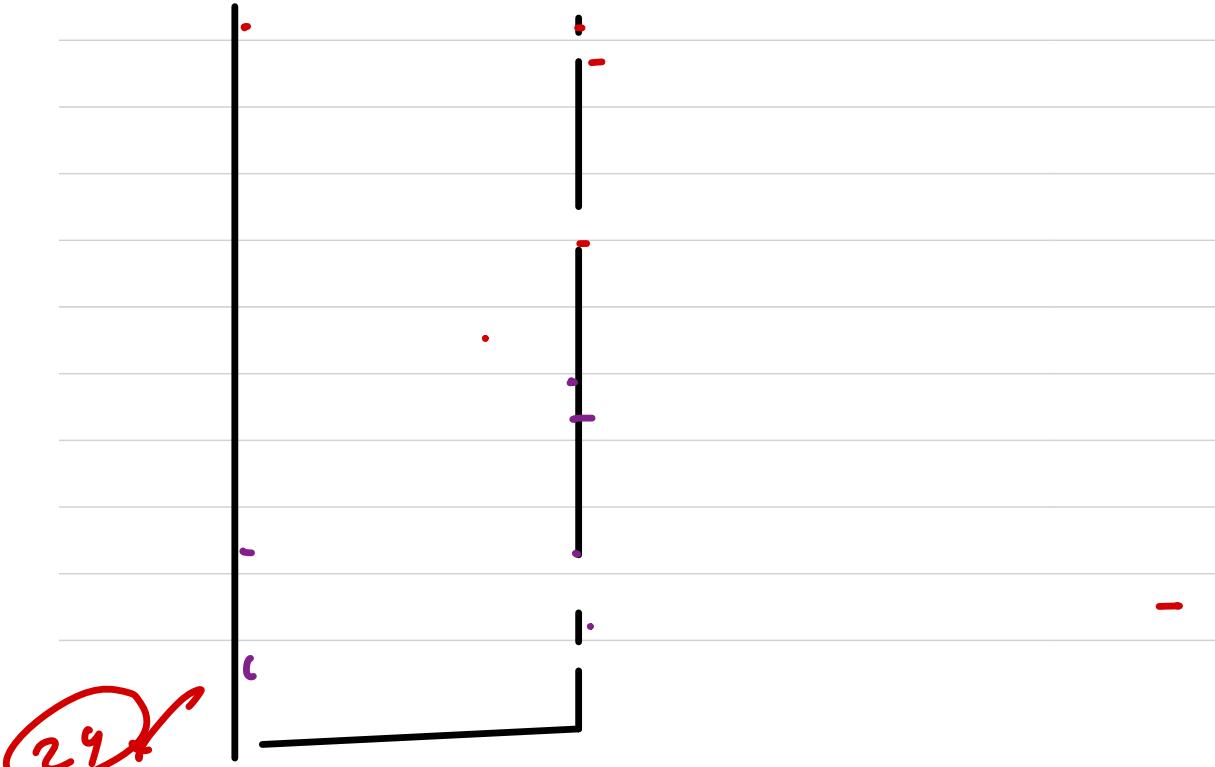
This new entry is called stack trace.

Inside a trace we store all the variables and more data like current line of execution of function.

When we hit a return , the stack trace is removed & value is given back to the caller.

```
1 function fact(n) { // this function should calculate n!
2     // Base case
3     if(n == 1) return 1;
4     let assume = fact(n-1); // that fact of n-1 is correct
5     return n*assume;; // self work
6 }
```

fact(4)



~~$\phi^n$~~  Given a value  $n$ , calc  $n^{\text{th}}$  fibonacci.

starting value

→ 0, 1, 1, 2, 3, 5, 8, 13, 21 ...  
0<sup>th</sup> fib 1<sup>st</sup> fib 2<sup>nd</sup> fib 3<sup>rd</sup> fib 4<sup>th</sup> fib 5<sup>th</sup> fib 6<sup>th</sup> fib ... - - - - -

$n = ?$

ans = 13

Solve it recursively

$n^{\text{th}}$  fib

$f(n)$



return  $n^{\text{th}}$

fib

Base Case

$$\begin{aligned}f(0) &\rightarrow 0 \\f(1) &\rightarrow 1\end{aligned}\quad \left.\right\} \xrightarrow{\text{start}} \underline{\text{start}}$$

Assumption → we assume  $f^{(n-1)}$  works  
fine &  $f^{(n-2)}$  works fine

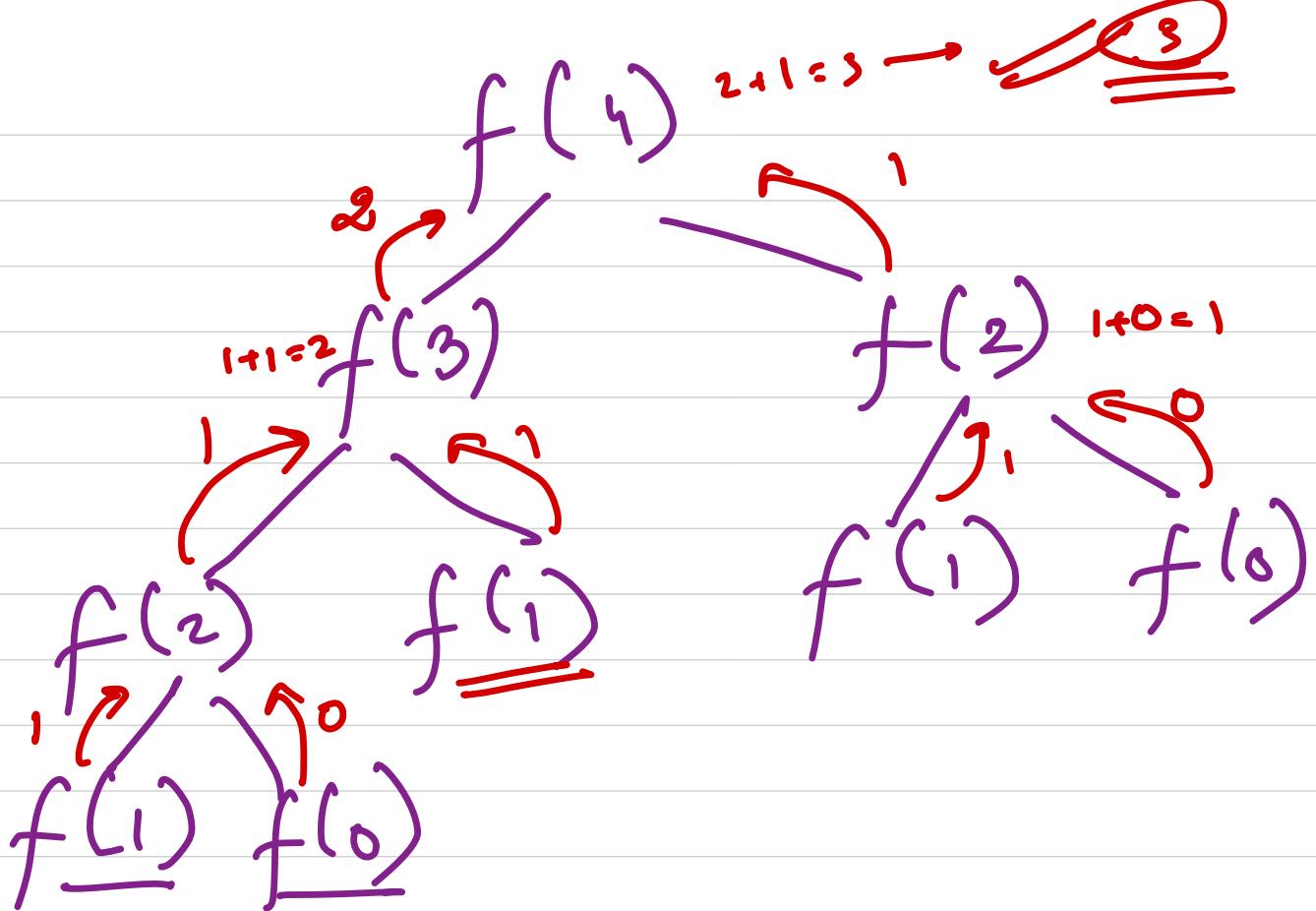
Self work →  $f^{(n-1)} + f^{(n-2)}$

$$\boxed{f(n) = f(n-1) + f(n-2)}$$

```
1 function fib(n) {  
2     // base case  
3     if(n == 0) return 0;  
4     if(n == 1) return 1;  
5     // assumption  
6     let a = fib(n-1);  
7     → let b = fib(n-2);  
8     // self work  
9     return a+b;  
10 }  
11  
12 console.log(fib(3))
```

Recursion takes  
Extra Space  
in memory

~~tree~~



$$\underline{f(n)} + f(n+1) = f(n+2)$$

$$f(n-1) + \underline{f(n)} = f(n+1)$$

$$f(n-2) + f(n-1) = f(n)$$

Say  $n+2 = x$

$$\underline{f(x-2) + f(x-1)} = f(x)$$

~~Q.~~ Given a number  $n$ , print the first  $N$  natural no. in increasing order, recursively

$$n = 5$$

Base  $\rightarrow n = 1 \rightarrow \text{console.} \log(1)$

Ans  $\rightarrow$  1  
2  
3  
4  
5

assume  $\rightarrow$  if someone can give me the first  $4$  natural no.'s then we can do something

Self  $\rightarrow$  point 5

$f(n)$  $f(n-1)$ 

console.log(n)

which print

first  $n$  natural

no.s

Assume  $f(n-1)$  works correctly

s  
↓  
→

s  
4  
3  
2  
1  
~~1~~  
~~1~~

$f(n)$

= `console.log(n)`

$f(n-1)$

which prints  
first  $n$  natural no  
in dec order

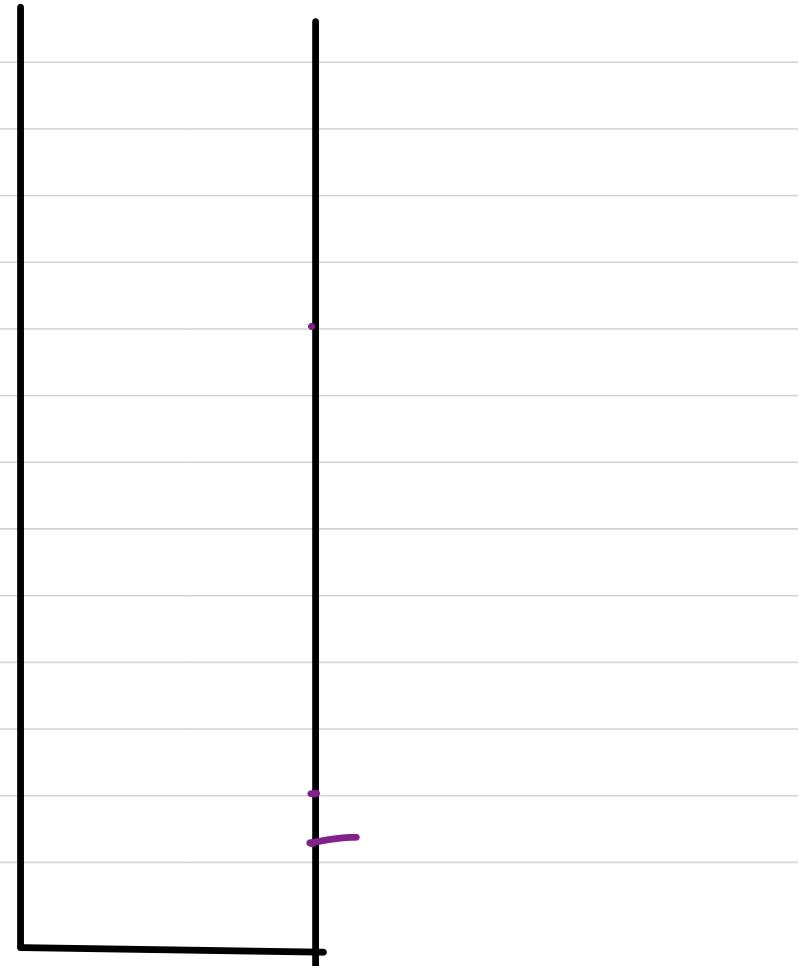
$f(n)$

5  
4  
3  
2  
1 }  
  

`console.log( $f_n$ )`  
 $f(n-1) \rightarrow$  assume it works few

```
1  function print_inc(n) {  
2      // base case  
3      if(n == 1) {  
4          console.log(1);  
5          return;  
6      }  
7      print_inc(n-1);  
8      console.log(n);  
9  }  
10  
11 → print_inc[3]||
```

1  
2  
3



```
1 function print_dec(n) {  
2     // base case  
3     if(n == 1) {  
4         console.log(1);  
5         return;  
6     }  
7     console.log(n); // Self task  
8     print_dec(n-1); // recursive  
9 }  
10  
11 print_dec(3)
```

3  
2  
1

Q. There are  $n$  friends, who want to go to a party. They can either go alone or go in a pair. Calc the no. of ways in which  $n$  friends can go to party. (Try Recursivity)

Ex  $n=3$

ans  $\rightarrow 4$

- |             |     |
|-------------|-----|
| (A) (B) (C) | → ✓ |
| (A B) (C)   | → ✓ |
| (A C) (B)   | → ✓ |
| (A) (B C)   | → ✓ |

$n=9$

A

B

C

D

Base Case ] →  $n \geq 1$  → Ans → 1  
                        ] →  $n = \leq 2$  → Ans → 2 ←  
A B → (A) (B) ]  
                        (A B) ]

$n=9$

A      B      C      D

$$f(n) = f(n-1) + (n-1) \times f(n-2)$$

return the  
no. of ways

in which  
 $n$  friends can go

$n$  friend  
said to go  
alone

$n$  friend  
said to  
make a pair

$$\overline{\overline{n=9}}$$

A

B

C

D

n<sup>th</sup> foiled

alone fair

what if D goes alone ??

- |      |     |     |     |
|------|-----|-----|-----|
| (A)  | (B) | (c) | (D) |
| (AB) | (c) | (D) |     |
| (AC) | (B) | (D) |     |
| (BC) | (A) | (D) |     |

What if  $n^{\text{th}}$  friend asked to make a pair ??

$$n=4$$

A

B

C

D

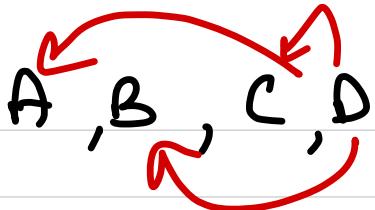
$\rightarrow \underline{n^{\text{th}} \text{ friend}}$

In how many ways  $n^{\text{th}}$  friend  
can make a pair ??  $\rightarrow (n-1)$

$$\underline{(n-1) \times f(n-2)}$$

$$3 \times 2 \\ = 6$$

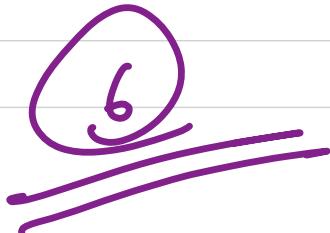
After making a pair how many elements left ??  $\rightarrow (n-2)$



$$\begin{bmatrix} (D)(C) & (A)(B) \\ (D)(C) & (A)(B) \end{bmatrix}$$

$$\begin{bmatrix} (D)(B) & (A)(C) \\ (D)(B) & (A)(C) \end{bmatrix}$$

$$\begin{bmatrix} (A)(D) & (B)(C) \\ (A)(D) & (B)(C) \end{bmatrix}$$



$f^{(n-2)} + f^{(n-2)}$   
 $+ f^{(n-2)}$   
 $(n-1) \times \underline{\underline{f^{(n-2)}}}$

$\Omega^n$  Given a number  $n$ , calculate the no. of  
binary strings of length  $n$  with no  
Consecutive One.

$$\underline{\underline{n=3}}$$

$$\text{Ans} \rightarrow \underline{\underline{5}}$$

000  
001  
010  
100  
101



$n$

1

2

3

4

5

ans

2

3

5

8

13

fibonacci

strings  
0, 1

00, 01, 10

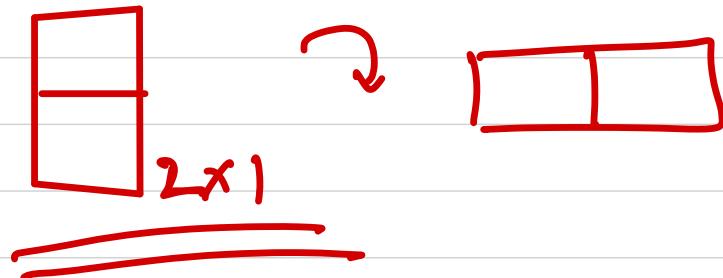
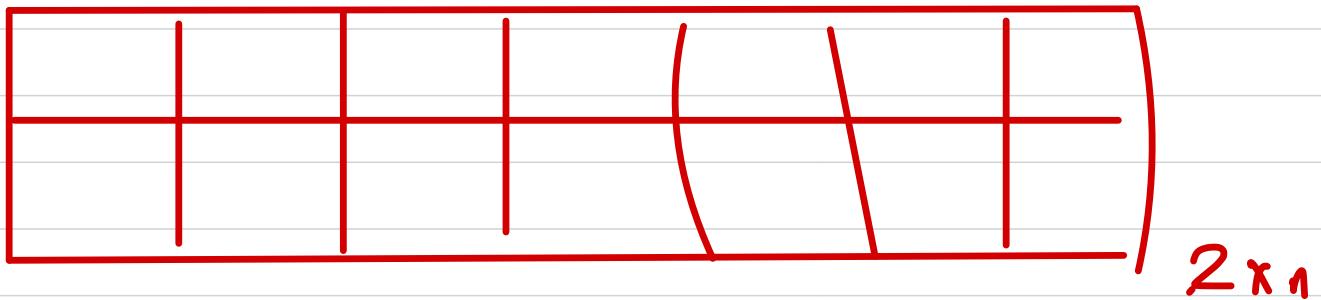
000, 001, 010,  
100, 101

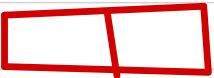
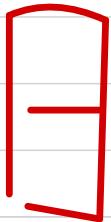
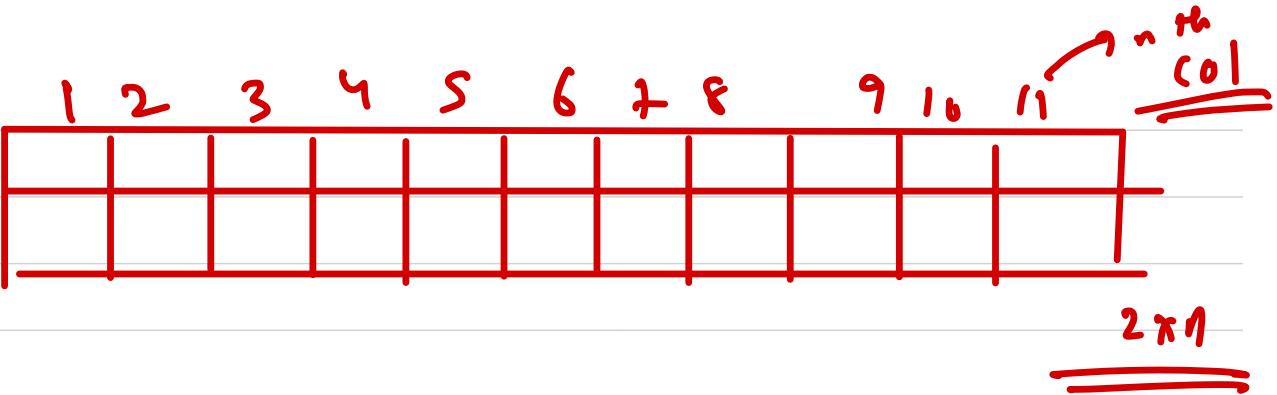
0000, 0001, 0010  
0100, 1000, 1001  
1010, 0101

Base Case  $\rightarrow$

if ( $n \geq 1$ ) return 2  
if ( $n \leq 2$ ) return 3

$$f(n) = \underline{\underline{f(n-1) + f(n-2)}}$$





On the  $n^{th}$  column you can put a tile either vertically or horizontally

$$f(n) = f(n-1) + f(n-2)$$

value of  
no. of ways to  
put tiles on

$2 \times n$  board

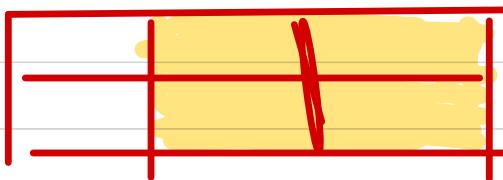
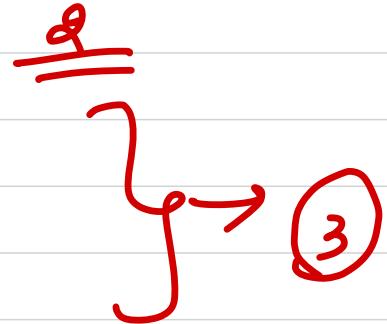
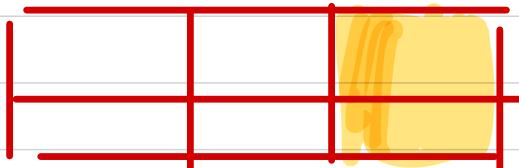
$n^{\text{th}}$  board  
via putting  
vertical  
tile

+

$n^{\text{th}}$  board with  
 $(n-1)^{\text{th}}$  board filled  
with horizontal  
tile

Base Case

$$\begin{aligned} n &= 1 \rightarrow 1 \\ n &= 2 \rightarrow 2 \end{aligned}$$



~~Q~~

Given an array , print all subsets of  
the array .

→ [1, 2, 3]

1

2

3

1 2

1 3

2 3

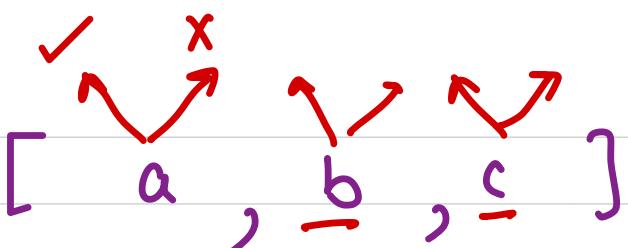
1 2 3

empty

if there is a set of length  $n$ ,  
how many subsets are possible ??

$$\underline{\underline{2^n}}$$

Permutation Combination

 [ a , b , c ]

$2 \times 2 \times 2 \dots$         $\rightarrow b c$

  $\rightarrow a c$

  $\rightarrow b$

  $\rightarrow a b c$

  $\rightarrow \text{empty}$

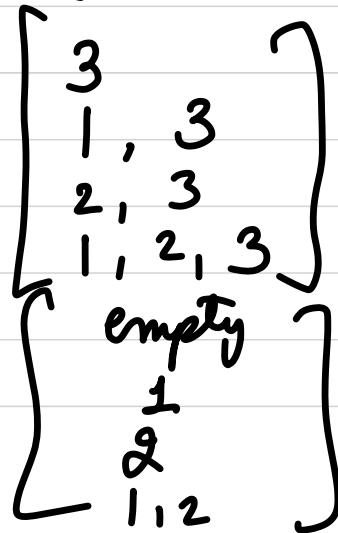
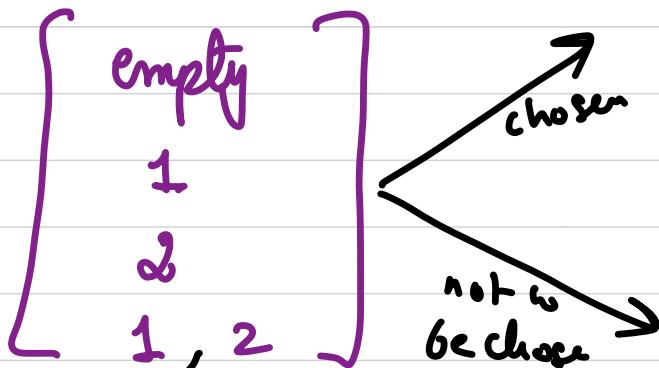
  $\rightarrow a b$

  $\rightarrow c$

We can use the same formula to build a  
recursion solution.

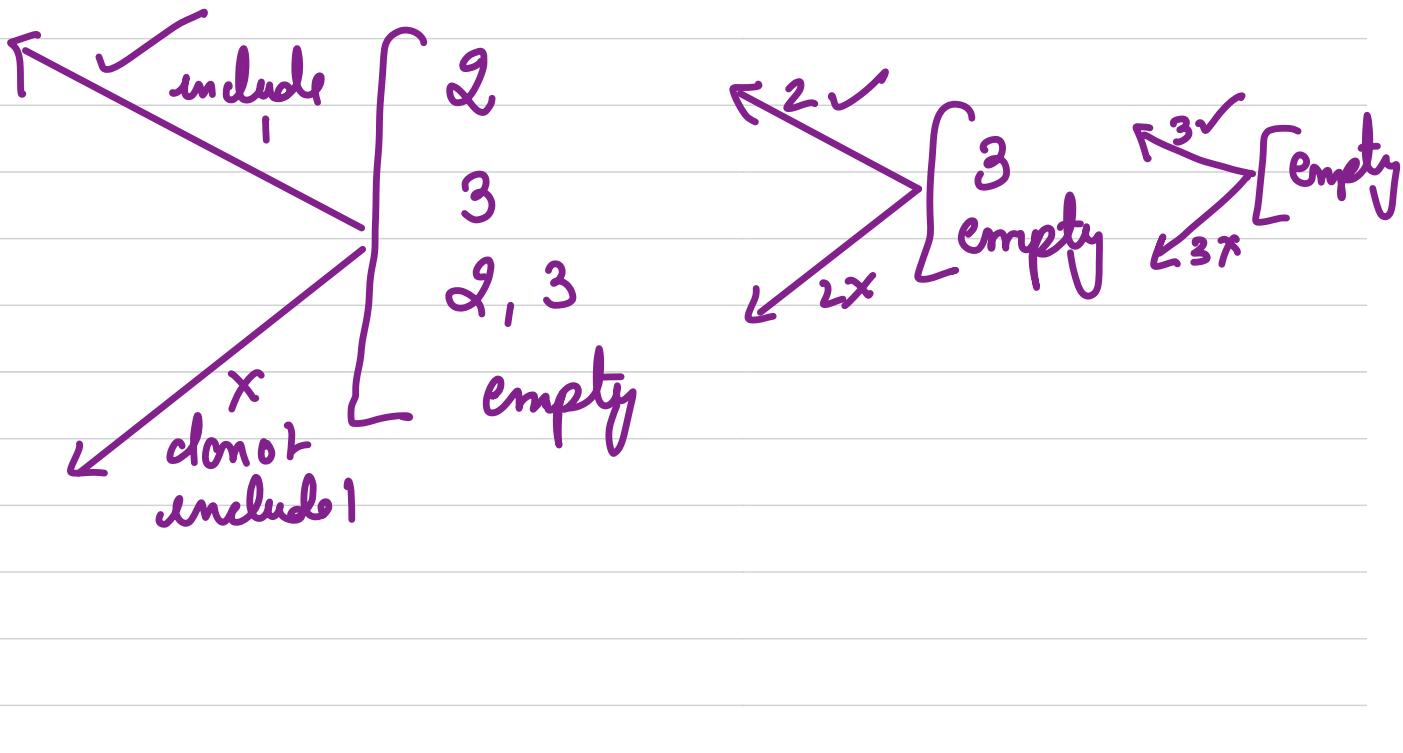
$$[1, 2, 3]_n$$

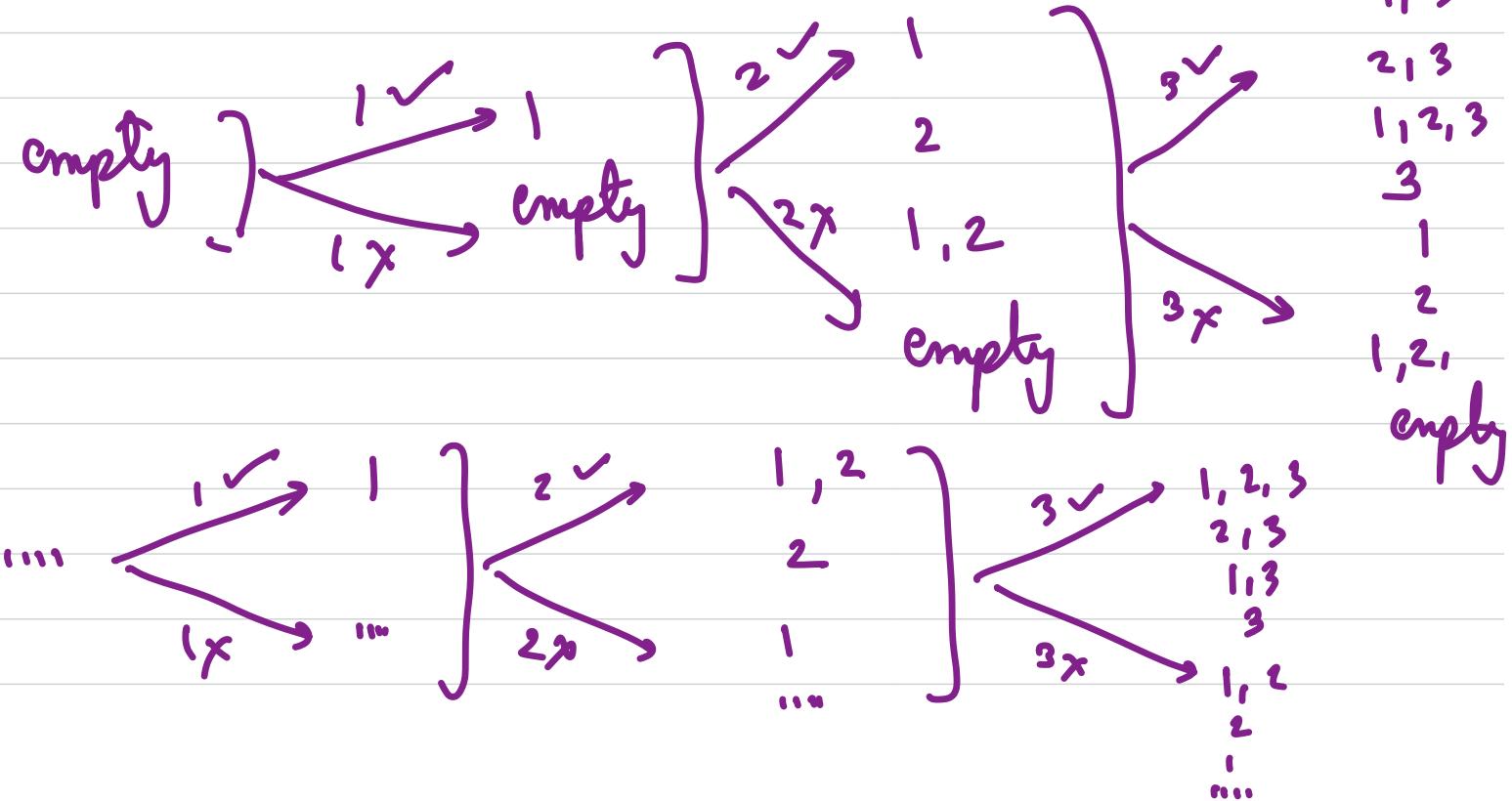
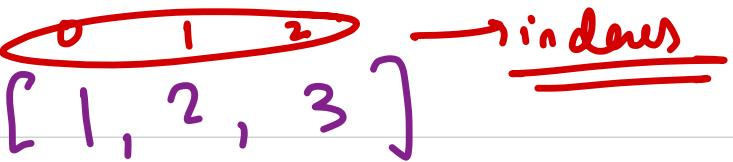
→ we can take the  
choices of  $n^{\text{th}}$   
element.



1st client

[ 1, 2, 3 ]





$$f(\text{arr}, \text{res}, i) = \begin{cases} f(\text{arr}, \text{res} + \overset{\text{append}}{\text{arr}[0]}, i+1) \\ f(\text{arr}, \text{res}, i+1) \end{cases}$$

two functions

prints all the  
Subsets of  
arr.

Console.log(f([1, 2, 3], ""))

•  
"3"  
"52"

( [S,6,7] , " " , 0 )

( [S,6,7] , "S" , 1 )

( [S,6,7] , " " , 1 )

( [S,6,7] , "S 6" , 2 )

( [S,6,7] , "S" , 2 )

( [S,6,7] , "6" , 2 )

( [S,6,7] , " " , 2 )

$\alpha\in [S, b, ?]$

$(\alpha, "", 0)$

$(\alpha, "S"., 1)$

$1v$

$1x$

$(\alpha, "S 6", 2)$

$(\alpha, "S", 2)$

$2v$

$2x$

$(\alpha, "S 6 7", 3)$

$(\alpha, "S 7", 3)$

$(\alpha, S, 3)$

$(\alpha, 6 7, 3)$

$(\alpha, 7, 3)$

$(\alpha, 6, 3)$

$(\alpha, "", 1)$

$1v$

$1x$

$(\alpha, "6", 2)$

$(\alpha, "", 2)$

$(\alpha, "7", 3)$

$(\alpha, "3", 3)$

