

## Nagarro class 1

Monday, 3 February 2020 10:48 AM

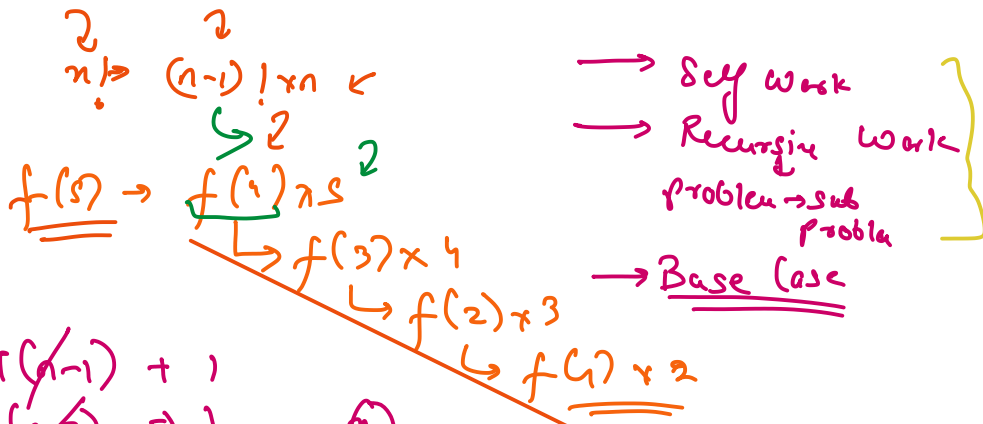
Recursion

→ When a function calls itself, to do some, repeated tasks, such that it is dividing a problem into sub-problems, along with a memory buffer.

$$f_i \rightarrow \textcircled{f_2}$$

factorial

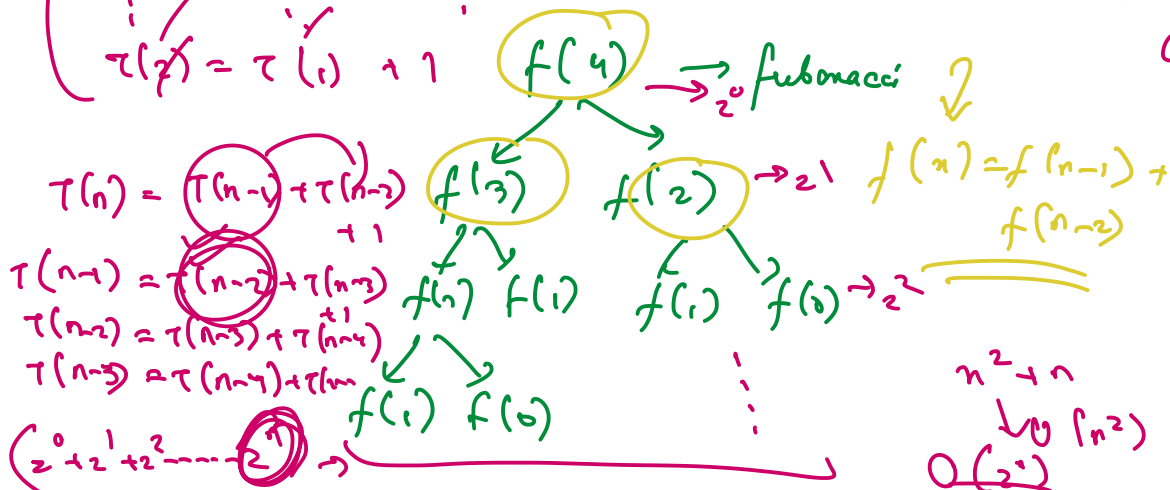
$$f(5)$$



$$\begin{cases} T(n) = T(n-1) + 1 \\ T(n-1) = T(n-2) + 1 \\ \vdots \\ T(2) = T(1) + 1 \end{cases}$$

(n)

$$T(n) = n \\ \underline{\underline{O(n)}}$$



You will be given 3 numbers, "a", "b" & "c"

return  $\rightarrow (a^b) \% c$

$a^b$

Time & Space complexity

$2^3 \Rightarrow 2 \times 2 \times 2$

$\rightarrow \textcircled{2} \times 2 \times 2$

$\rightarrow 2 \times 2$

$$a^b = a \times a^{b-1}$$

$$f(a, b) = a \times f(a, b-1)$$

2 →  $a \times (a \times f(a, b-2))$

←

←

result = 1

for (i = 1 → b) {

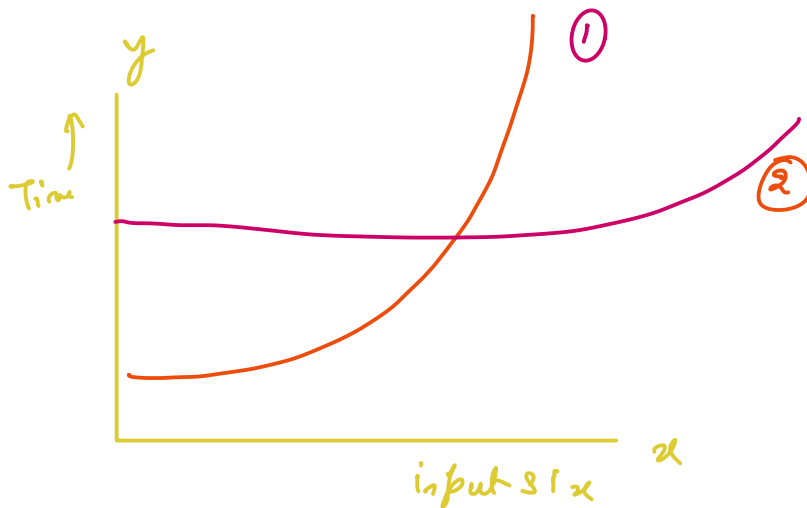
    result \*= a;

}

3

$a^b$

$O(b)$

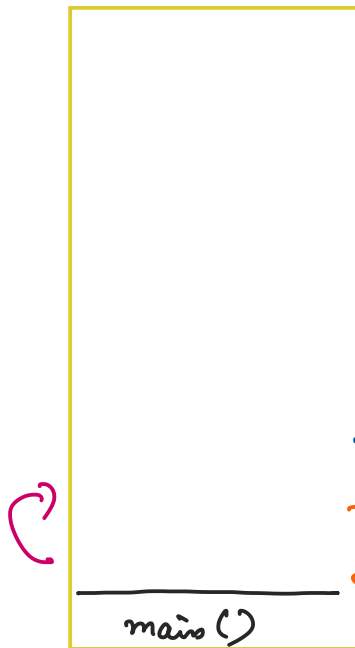


## # Modulo Arithmetic

$$\left\{ \begin{array}{l} (a+b) \% c = ((a \% c) + (b \% c)) \% c \\ (a-b) \% c = (a \% c - b \% c + c) \% c \\ (a \times b) \% c = (a \% c \times b \% c) \% c \end{array} \right\}$$

→ Inverse Modulo Arithmetic

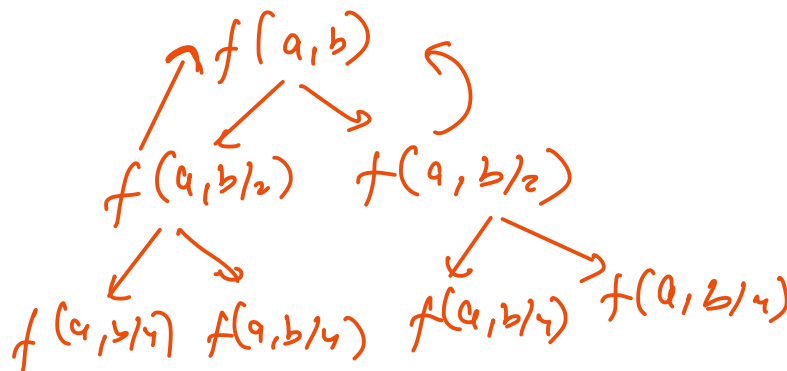
$$\underline{a^b \% c}$$



1)  $\rightarrow$  int powm(int a, int b) {  
 2) — if (b == 1) return a;  
 3)  $\rightarrow$  int result = powm(a, b-1);  
 4)  $\rightarrow$  return result \* a;  
 5) }

$2^3$

$$a^b = a^{b/2} \times a^{b/2}$$



$$f(a, b) = f(a, b/2) + f(a, b/2) + 1$$

$$= 2f(a, b/2) + 1$$

$$\hookrightarrow \underline{f(a, b) = 2f(a, b/2)}$$

$$= 2(2f(a, b/4))$$

$$= 2(2(2f(a, b/8)))$$

$$\begin{array}{c}
 a \\
 \downarrow \\
 a/2 \\
 \downarrow \\
 a/4 \\
 \downarrow \\
 a/8
 \end{array}$$

$$= \underbrace{2 \times 2 \times 2 \dots 2}_{k \text{ terms}} \times f(a, b/2^k)$$

$$\Rightarrow 2^k \times f(a, 1) = 2^{\log_2 b} \times 1$$

$$\rightarrow \underline{b}$$

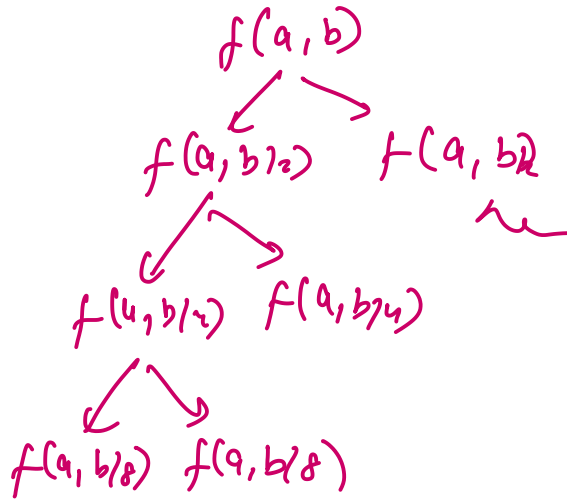
iterations

$$\hookrightarrow \frac{b}{2^k} = 1$$

$$\hookrightarrow b = 2^k \rightarrow \text{taking log both sides}$$

$$\log_2 b = (\log_2 2) \times k$$

$$\rightarrow k = \log_2 b$$



$$b \rightarrow b/2 \rightarrow b/4 \rightarrow \dots$$

$$\left. \begin{aligned} f(a, b) &= f(a, b/2) + 1 \\ f(a, b/2) &= f(a, b/4) + 1 \\ f(a, b/4) &= f(a, b/8) + 1 \\ f(a, b/8) &= f(a, 1) + 1 \end{aligned} \right\} k \rightarrow \text{no. of calls}$$

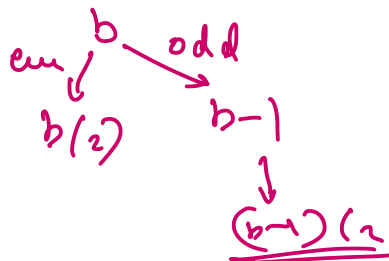
$$f(a, b) = f(a, 1) + k + 1$$

$$T(f(a, b)) = \underbrace{O(1) + \log_2 b}_{O(\log_2 b)}$$

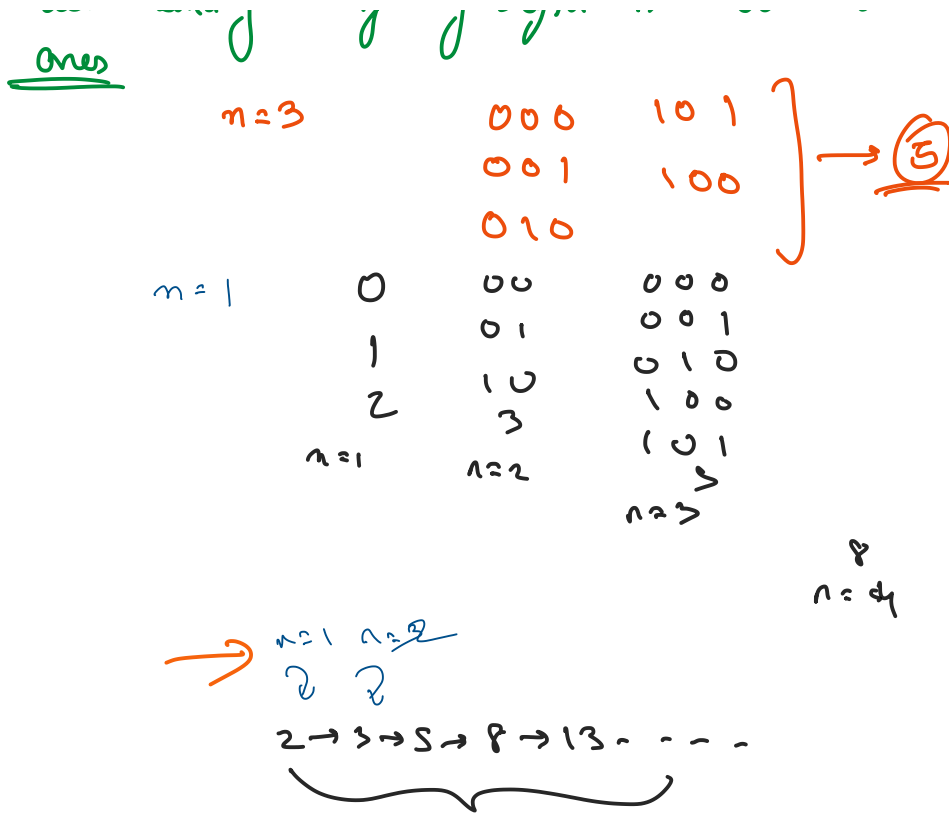
$$2^{10} = 1024 \approx 10^3$$

$$2^{20} \approx 10^6$$

$$2^{30} \approx 10^9$$

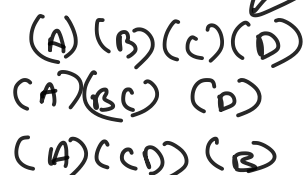
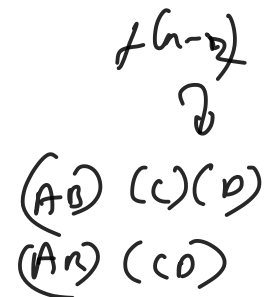
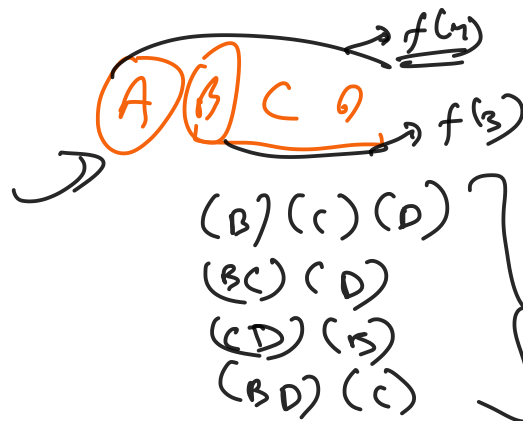
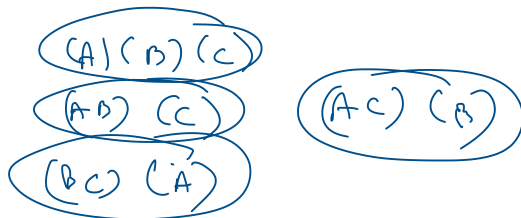


Q You have been given a number "n". Count all binary strings of length "n" with no consecutive



Q friends pairing → 0

$n=3$  A B C



(A) (B) (C)

$$f(n) = f(n-1) + f(n-2) + f(n-3)$$

f(3)

BCD

(C) (D) → f(2)

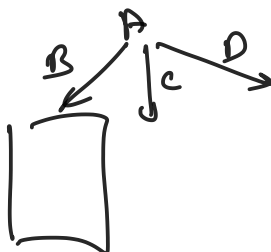
(B) (C) (D)

(B) (C) (D)

$$f(3) = f(2) + \underline{\hspace{2cm}}$$

f(4)

(n-1)



(A) (B) (C) (D)

(A) (B) (C) (D)

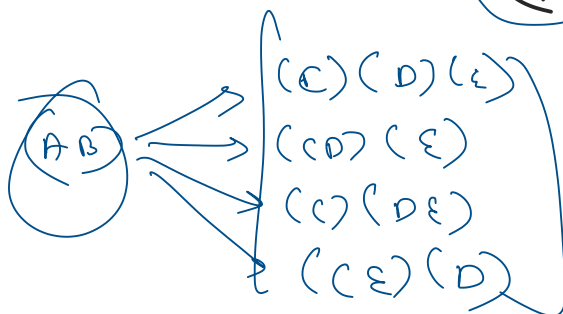
(A) (B) (C) (D)

(A) (C) (B) (D)

(A) (C) (B) (D)

(A) (D) (B) (C)

(A) (D) (B) (C)



abc

abc

→

" "

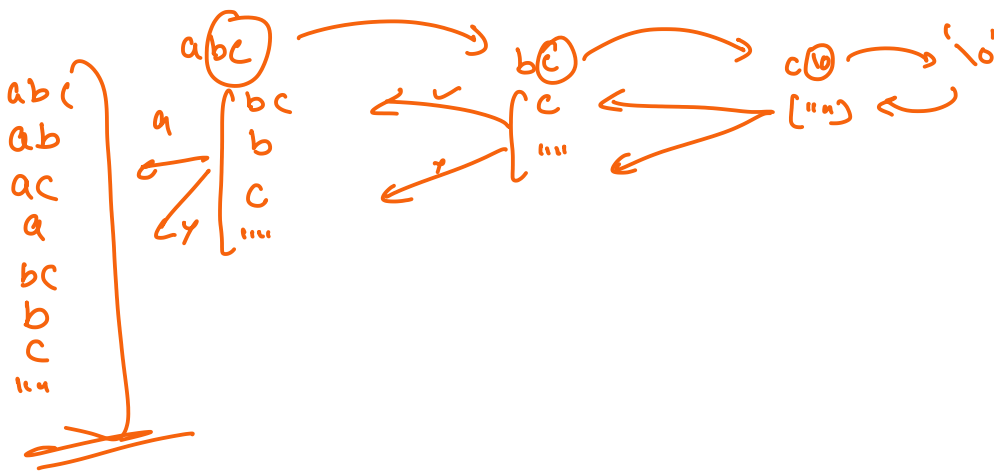
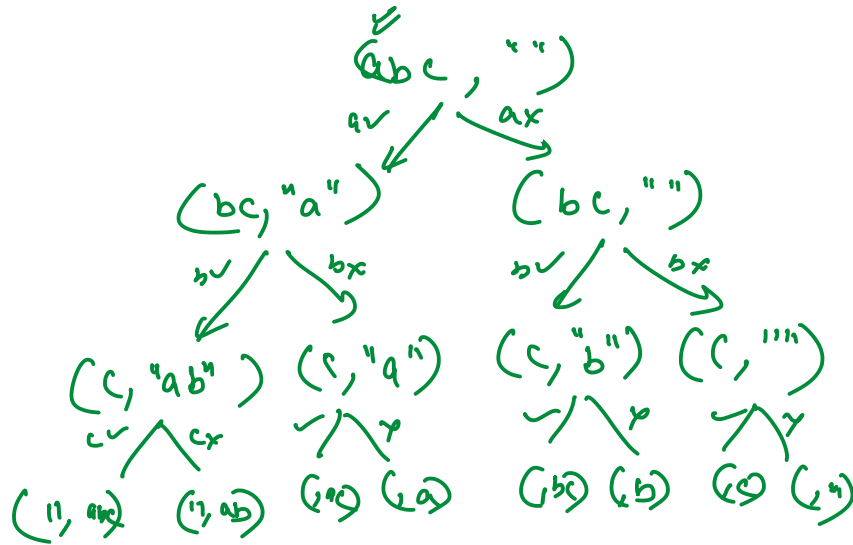
a

✓

✗

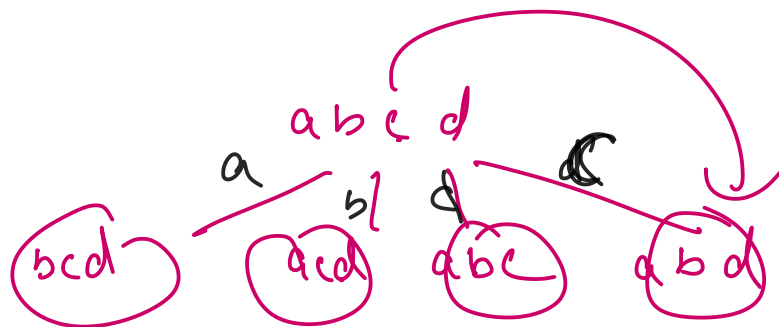
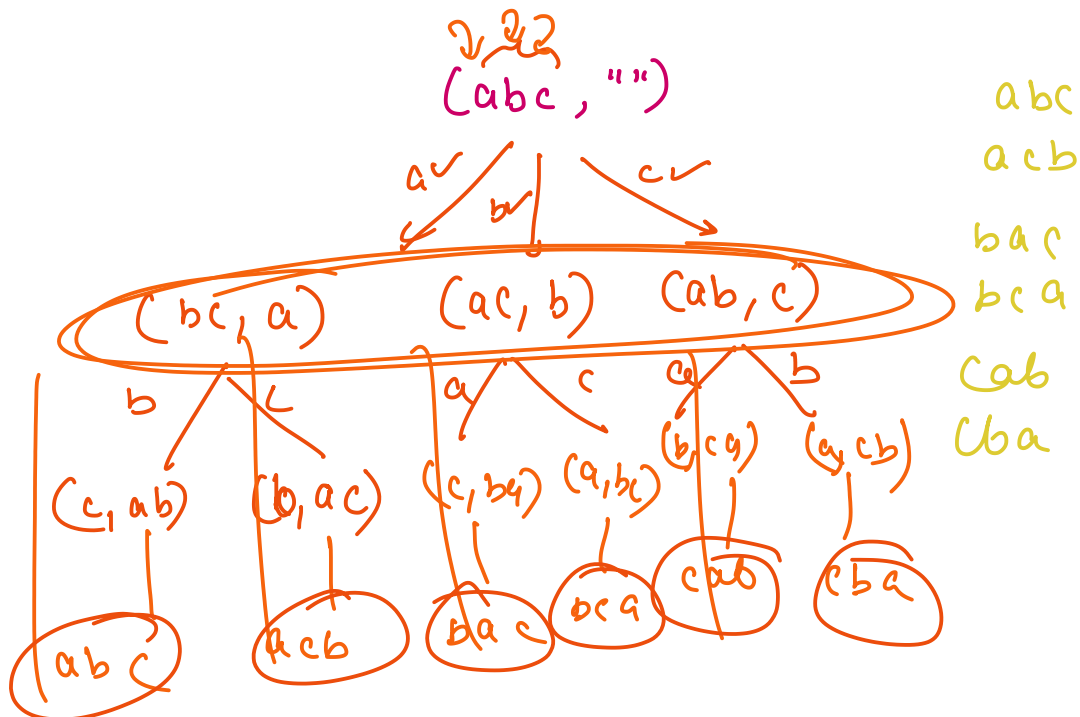
b	↓	↓	↓
c	a	b	c
ab	↓	↓	↓
bc	a	a	a
abc	→ 2		

$\{a\}$        $\{abc\}$        $\{a, b, c\}$



abc  
ab  
ac

<u>ss(' ', ac)</u>	
<u>ss(c, a)</u>	ch = c (s1)
<u>ss(bc, a)</u>	ch = b (s1)
<u>ss(abc, '')</u>	ch = a (s1)
main()	



str.substrig(0, 2) + str.substrig(2, 10)

↓ ↓

ab ←→ d

abd



$$T(n) = T(n-1) * n \quad \text{for loop}$$

↓  
occurs:

$$T(n) = n \times [2n + \cancel{n! \times n} + T(n-1)]$$

for loop

$$T(n) = n \times (T(n-1) + O(1))$$

