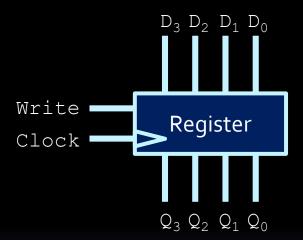
Week 6 Lectorial

Question #1

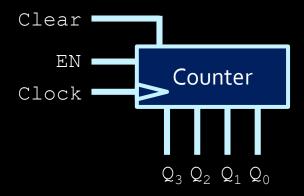
Imagine you have access to a 4-bit register.



What does the Write signal do?

Question #2

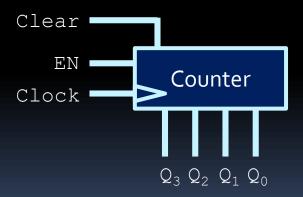
Assume that you have access to a counter circuit:



- How do you make a signal that goes high after 10 clock cycles?
- How do you make a signal that goes high every 10 clock cycles?

Question #2 (cont'd)

• How do you make a signal that goes high every 100 clock cycles, only using 4-bit counters like the one below (and a few additional gates)?



Question #3

• How many flipflops would you need to implement the following finite state machine (FSM)?

- 11 states
- # flip-flops = 4

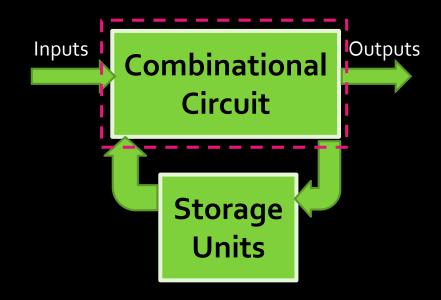


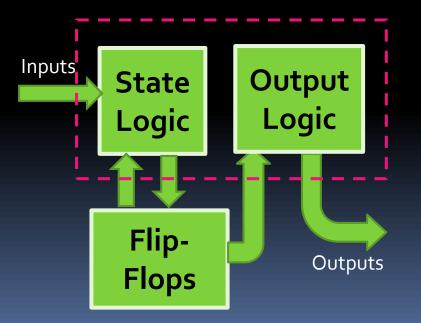
Reminder: How to Design FSM

- As a brief reminder:
 - Draw state diagram
 - Derive state table from state diagram
 - 3. Assign flip-flop configuration to each state
 - Number of flip-flops needed is: \[\log(# of states) \]
 - 4. Redraw state table with flip-flop values
 - Derive combinational circuit for output and for each flip-flop input.

Review of FSMs

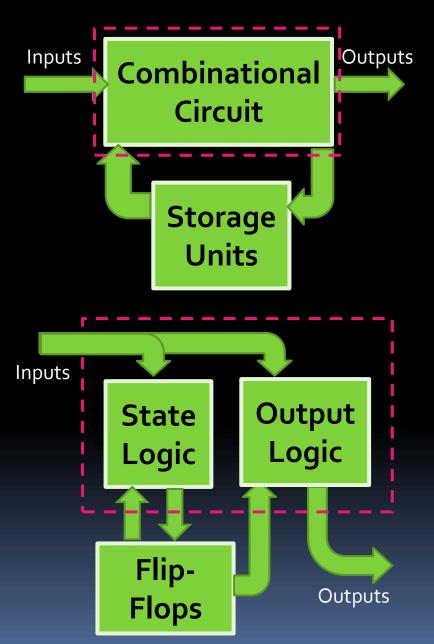
- Step 5 requires two combinational circuit design tasks.
 - For Moore machines
 (pictured bottom right),
 output is determined
 solely based on current
 state (i.e. flip-flop
 values).





Review of FSMs

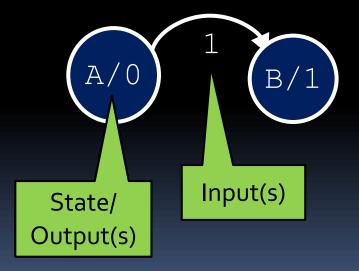
- For Mealy machines, output is determined by both the current state and the current input values.
 - For simplicity, most of our examples will focus on Moore machines.



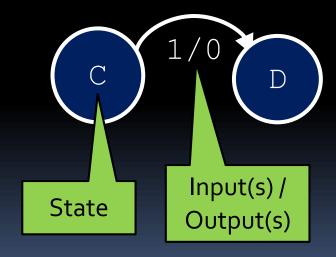
State diagrams with output

 Output values are incorporated into the state diagram, depending on the machine used.

Moore Machine



Mealy Machine



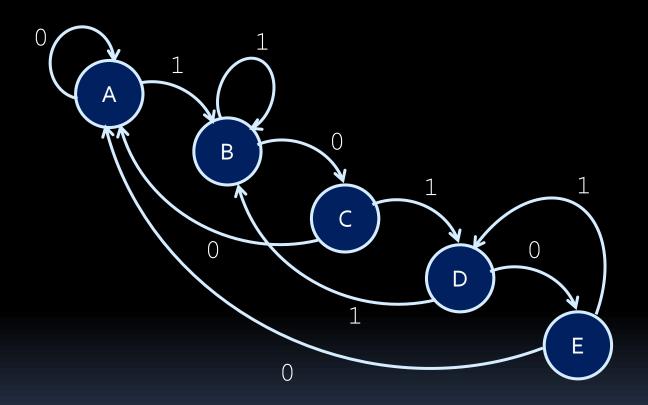
FSM Example: Barcode Reader

When scanning UPC
 barcodes, the laser
 scanner looks for black
 and white bars that
 indicate the start of the



- indicate the start of the code.
- If black is read as a 1 and white is read as a 0, the start of the code (from either direction) has a 1010 pattern.
 - Can you create a state machine that detects this pattern?

Step #1: Draw state diagram



Step #2: State Table

- Output Z is determined by the current state.
 - Denotes Moore machine.
- Next step: allocate flipflops values to each state.
 - How many flip-flops will we need for 5 states?
 - Recall:
 - # flip-flops = \[\log(\psi \text{ of states}) \]

Present State	Z	x	Next State
А	0	0	A
А	0	1	В
В	0	0	С
В	0	1	В
С	0	0	A
С	0	1	D
D	0	0	E
D	0	1	В
E	1	0	A
E	1	1	D

Step #3: Flip-Flop Assignment

 3 flip-flops needed here.

Assign states carefully though!

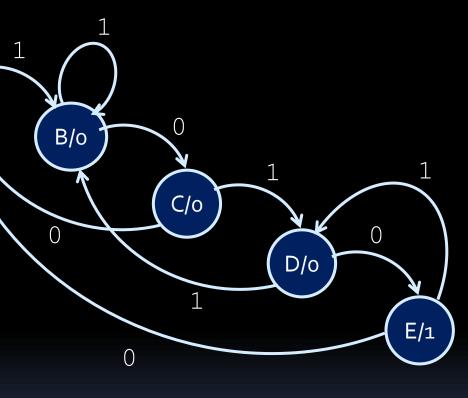
Can't simply do this:

$$> A = 100 > B = 011$$

$$\triangleright$$
 C = 010 \triangleright D = 001

Why not?

A/o



Step #3: Flip-Flop Assignment

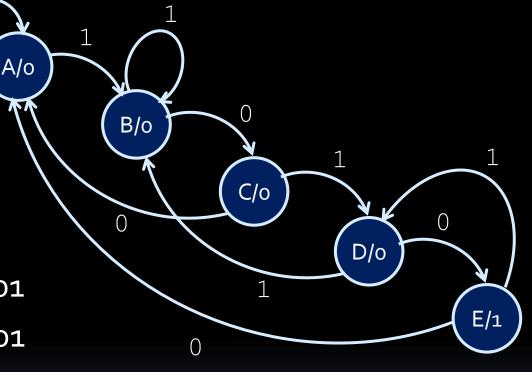
Be careful of race conditions.

Better solution:

$$A = 000 > B = 001$$



Sometimes, extra flip-flops may be necessary.



Step #4: Redraw State Table

- From here, we can construct the K-maps for the state logic combinational circuit.
 - Derive equations for each flip-flop value, given the previous values and the input X.
 - Three equations total, plus one more for $\mathbb Z$ (trivial for Moore machines).

Present State		Z	x	Next State			
0	0	0	0	0	0	0	0
0	0	0	0	1	0	0	1
0	0	1	0	0	0	1	1
0	0	1	0	1	0	0	1
0	1	1	0	0	0	0	0
0	1	1	0	1	1	0	1
1	0	1	0	0	1	0	0
1	0	1	0	1	0	0	1
1	0	0	1	0	0	0	0
1	0	0	1	1	1	0	1

• Karnaugh map for F₂:

	$\overline{\mathbf{F}}_0 \cdot \overline{\mathbf{X}}$	F ₀ ⋅ x	F ₀ · X	$\mathbf{F}_0 \cdot \overline{\mathbf{X}}$	
$\overline{\mathbf{F}}_2 \cdot \overline{\mathbf{F}}_1$	0	0	0	0	
$\overline{\mathbf{F}}_2 \cdot \mathbf{F}_1$	X	X	1	0	
$\mathbf{F}_2 \cdot \mathbf{F}_1$	X	X	X	X	
$\mathbf{F}_2 \cdot \overline{\mathbf{F}}_1$	0	1	0	1	

$$F_2 = F_1X + F_2\overline{F}_0X + F_2F_0\overline{X}$$

• Karnaugh map for F₁:

	$\overline{\mathbf{F}}_0 \cdot \overline{\mathbf{X}}$	F ₀ ⋅ x	F ₀ · X	$\mathbf{F}_0 \cdot \overline{\mathbf{X}}$
$\overline{\mathbf{F}}_2 \cdot \overline{\mathbf{F}}_1$	0	0	0	1
$\overline{\mathbf{F}}_2 \cdot \mathbf{F}_1$	X	X	0	0
$\mathbf{F}_2 \cdot \mathbf{F}_1$	X	X	X	X
$\mathbf{F}_2 \cdot \overline{\mathbf{F}}_1$	0	0	0	0

$$F_1 = F_2 F_1 F_0 \overline{X}$$

• Karnaugh map for F_o:

	$\mathbf{\overline{F}}_0 \cdot \mathbf{\overline{X}}$	F ₀ · x	F ₀ ·X	$\mathbf{F}_0 \cdot \overline{\mathbf{X}}$
$\overline{\mathbf{F}}_2 \cdot \overline{\mathbf{F}}_1$	0	1	1	1
$\overline{\mathbf{F}}_2 \cdot \mathbf{F}_1$	X	X	1	0
$\mathbf{F}_2 \cdot \mathbf{F}_1$	X	X	X	X
$\mathbf{F}_2 \cdot \overline{\mathbf{F}}_1$	0	1	1	0

$$F_0 = X + \overline{F}_2 \overline{F}_1 F_0$$

Output value Z goes high based on the following output equation:

$$Z = F_2 \overline{F}_1 \overline{F}_0$$

- Note: All of these equations would be different, given different flip-flop assignments!
 - lacktriangle Practice alternate assignment for the midterm lacktriangle