# Experiment 2- Use Divide and Conquer approach to implement Develop a code for Strassen's Matrix Multiplication and analyze it.

**<u>Learning Objective:</u>** Students should be able to perform Strassen's matrix multiplication on a given data set using divide and conquer strategy.

**Tools:** C/C++/Java/Python under Windows or Linux environment.

**Theory:** Develop a code for Strassen's matrix multiplication using divide and conquer strategy and analyze it. (Menu driven program)

Following is a simple Divide and Conquer method to multiply two square matrices.

- 1. Divide matrices A and B in 4 sub-matrices of size N/2 x N/2 as shown in the below diagram.
- 2. Calculate following values recursively. ae + bg, af + bh, ce + dg and cf + dh.

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix} x \begin{bmatrix} e & f \\ g & h \end{bmatrix} = \begin{bmatrix} ae + bg & af + bh \\ ce + dg & cf + dh \end{bmatrix}$$
A B C

A, B and C are square metrices of size N x N

a, b, c and d are submatrices of A, of size N/2 x N/2

e, f, g and h are submatrices of B, of size N/2 x N/2

Strassen suggested a divide and conquer strategy-based matrix multiplication technique that requires fewer multiplications than the traditional method. The multiplication operation is defined as follows using Strassen's method:

$$\begin{bmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \times \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

C11 = P1 + P4 - P5 + P7

C12 = P3 + P5

C21 = P2 + P4

C22 = P1 + P3 - P2 + P6

```
Where,
```

This is the same as C12 derived using the conventional approach. Similarly we can derive all Cij for strassen's matrix multiplication.

#### **ALGORITHM:**

= A11 \* B12 + A12 \* B22

```
Algorithm STRESSEN MAT MUL (int *A, int *B, int *C, int n)
// A and B are input matrices
// C is the output matrix
// All matrices are of size n x n
if n == 1 then
  *C = *C + (*A) * (*B)
else
  STRESSEN MAT MUL (A, B, C, n/4)
  STRESSEN MAT MUL (A, B + (n/4), C + (n/4), n/4)
  STRESSEN MAT MUL (A + 2 * (n/4), B, C + 2 * (n/4), n/4)
  STRESSEN MAT MUL (A + 2 * (n/4), B + (n/4), C + 3 * (n/4), n/4)
  STRESSEN MAT MUL (A + (n/4), B + 2 * (n/4), C, n/4)
  STRESSEN MAT MUL (A + (n/4), B + 3 * (n/4), C + (n/4), n/4)
  STRESSEN MAT MUL (A + 3 * (n/4), B + 2 * (n/4), C + 2 * (n/4), n/4)
  STRESSEN MAT MUL (A + 3 * (n/4), B + 3 * (n/4), C + 3 * (n/4), n/4)
end
```

### **ADVANTAGE of Divide & Conquer Algorithm:**

1. The difficult problem can be solved easily.

- 2. It divides the entire problem into subproblems thus it can be solved parallelly ensuring multiprocessing
- 3. Efficiently uses cache memory without occupying much space
- 4. Reduces time complexity of the problem

#### **DISADVANTAGES of Divide & Conquer Algorithm:**

- 1. It involves recursion which is sometimes slow
- 2. Efficiency depends on the implementation of logic
- 3. It may crash the system if the recursion is performed rigorously

#### **ANALYSIS:**

#### **Time Complexity**:

Addition and Subtraction of two matrices takes O(N2) time. So time complexity can be written as

T(N) = 7T(N/2) + O(N2)

From Master's Theorem, time complexity of above method is

O(NLog7) which is approximately O(N2.8074)

## **Space Complexity:**

A new matrix is used to store the result of the multiplication. So, the space complexity is  $O(N^2)$ .

## **APPLICATION:**

Generally, Strassen's Method is not preferred for practical applications for the following reasons.

- The constants used in Strassen's method are high and for a typical application Naive method works better.
- For Sparse matrices, there are better methods especially designed for them.
- The submatrices in recursion take extra space.
- Because of the limited precision of computer arithmetic on noninteger values, larger errors accumulate in Strassen's algorithm than in Naive Method

### **<u>Learning Outcome:</u>** The student should have the ability to

LO1: understand and implement divide and conquer strategy.

LO2: perform matrix multiplication of Strassen's matrix using divide and conquer.

<u>Course Outcomes:</u> Upon completion of the course students will be able to apply divide and conquer strategy to given data.

#### **CODE:**

```
def strassen 2x2 recursive(A, B):
  n = len(A)
  if n == 1:
     return [[A[0][0] * B[0][0]]]
  else:
     a11, a12, a21, a22 = split matrix(A)
     b11, b12, b21, b22 = split matrix(B)
     # Calculate the following values recursively
     p1 = strassen 2x2 recursive(add matrices(a11, a22), add matrices(b11, b22))
     p2 = strassen 2x2 recursive(add matrices(a21, a22), b11)
     p3 = strassen 2x2 recursive(a11, subtract matrices(b12, b22))
     p4 = strassen 2x2 recursive(a22, subtract matrices(b21, b11))
     p5 = strassen 2x2 recursive(add matrices(a11, a12), b22)
     p6 = strassen 2x2 recursive(subtract matrices(a21, a11), add matrices(b11, b12))
     p7 = strassen 2x2 recursive(subtract matrices(a12, a22), add matrices(b21, b22))
     # Combine the results into a 2x2 matrix
     c11 = subtract matrices(add matrices(add matrices(p1, p4), p7), p5)
     c12 = add matrices(p3, p5)
     c21 = add matrices(p2, p4)
     c22 = subtract matrices(add matrices(add matrices(p1, p3), p6), p2)
     C = [[0, 0], [0, 0]]
     for i in range(n // 2):
       for j in range(n // 2):
         C[i][j] = c11[i][j]
         C[i][j + n // 2] = c12[i][j]
         C[i + n // 2][j] = c21[i][j]
         C[i + n // 2][j + n // 2] = c22[i][j]
     return C
def split matrix(A):
  n = len(A)
  m = n // 2
```

```
a11 = [[0] * m \text{ for i in range}(m)]
  a12 = [[0] * m \text{ for i in range}(m)]
  a21 = [[0] * m \text{ for idef strassen } 2x2 \text{ recursive(A, B):}
  n = len(A)
  if n == 1:
     return [[A[0][0] * B[0][0]]]
  else:
     a11, a12, a21, a22 = split matrix(A)
     b11, b12, b21, b22 = split matrix(B)
     # Calculate the following values recursively
     p1 = strassen 2x2 recursive(add matrices(a11, a22), add matrices(b11, b22))
     p2 = strassen 2x2 recursive(add matrices(a21, a22), b11)
     p3 = strassen 2x2 recursive(a11, subtract matrices(b12, b22))
     p4 = strassen 2x2 recursive(a22, subtract matrices(b21, b11))
     p5 = strassen 2x2 recursive(add matrices(a11, a12), b22)
     p6 = strassen 2x2 recursive(subtract matrices(a21, a11), add matrices(b11, b12))
     p7 = strassen 2x2 recursive(subtract matrices(a12, a22), add matrices(b21, b22))
     # Combine the results into a 2x2 matrix
     c11 = subtract matrices(add matrices(add matrices(p1, p4), p7), p5)
     c12 = add matrices(p3, p5)
     c21 = add matrices(p2, p4)
     c22 = subtract matrices(add matrices(add matrices(p1, p3), p6), p2)
     C = [[0, 0], [0, 0]]
     for i in range(n // 2):
       for j in range(n // 2):
          C[i][j] = c11[i][j]
          C[i][j + n // 2] = c12[i][j]
          C[i + n // 2][j] = c21[i][j]
          C[i + n // 2][j + n // 2] = c22[i][j]
     return C
def split matrix(A):
  n = len(A)
  m = n // 2
  a11 = [[0] * m \text{ for i in range}(m)]
```

```
a12 = [[0] * m for i in range(m)]
  a21 = [[0] * m \text{ for } i \text{ in range}(m)]
  a22 = [[0] * m for i in range(m)]
  for i in range(m):
     for j in range(m):
        a11[i][j] = A[i][j]
        a12[i][j] = A[i][j + m]
        a21[i][j] = A[i + m][j]
        a22[i][j] = A[i + m][j + m]
  return a11, a12, a21, a22
def add matrices(A, B):
  n = len(A)
  C = [[0] * n \text{ for } i \text{ in } range(n)]
  for i in range(n):
     for j in range(n):
        C[i][j] = A[i][j] + B[i][j]
  return C
def subtract_matrices(A, B):
  n = len(A)
  C = [[0] * n \text{ for i in range}(n)]
  for i in range(n):
     for j in range(n):
        C[i][j] = A[i][j] - B[i][j]
  return C
a = [[12,1],[23,80]]
b = [[56,14],[90,18]]
print(strassen 2x2 recursive(a,b))
```

#### **OUTPUT:**

[[762, 186], [8488, 1762]]

# **Conclusion:**

# For Faculty Use

Correction Parameters	Formative Assessment [40%]	Timely completion of Practical [ 40%]	Attendance / Learning Attitude [20%]
Marks Obtained			