

## WEEKLY TEST – 03

## Subject : Database Management System

## Topic : Normal Form and Serializability



Maximum Marks 12

## Q.1 to 4 Carry ONE Mark Each

## [MCQ]

1. Consider the relation  $R(ABCDE)$  with the following set of functional dependencies (FD's)  
 $F: \{AB \rightarrow C, C \rightarrow D, D \rightarrow E, E \rightarrow A, D \rightarrow B\}$   
 The which of the following Highest Normal form satisfied by the above relation  $R$ ?  
 (a) BCNF (b) 3NF  
 (c) 2NF (d) 1NF

## [MSQ]

2. Consider the following relation with given functional dependencies.  
 $A(R, S, T, U, V)$   
 $FD's: \{RS \rightarrow T, RS \rightarrow U, U \rightarrow R, ST \rightarrow U, ST \rightarrow V\}$   
 Then Relation is in which Normal Form?  
 (a) 1NF (b) 2NF  
 (c) 3NF (d) BCNF

## 3. [NAT]

Consider 3 transactions  $T_1, T_2$  and  $T_3$  having 2, 3 and 4 operations respectively.

The total number of non-serial schedule is \_\_\_\_\_.

## [MSQ]

4. Consider the following schedule:  
 $S_1: r_1(x) \ w_2(x) \ w_1(x) \ w_3(x) \ r_3(x)$   
 $S_2: r_1(x) \ w_2(x) \ w_3(x) \ w_1(x) \ r_3(x)$   
 Which of the following is correct about the Schedule  $S$ ?  
 (a)  $S_1$  is conflict serializable.  
 (b)  $S_1$  is not conflict serializable.  
 (c)  $S_2$  is conflict serializable.  
 (d)  $S_2$  is not conflict serializable.

## Q.5 to 8 Carry TWO Mark Each

## 5. [NAT]

Consider the following Relation:

$R(ABCDEFGH)$  with FD set of Relation  $R \ \{A \rightarrow B, C \rightarrow D, E \rightarrow FGH\}$  what is the minimum number of relations required to decompose into BCNF which satisfy lossless join and Dependency preserving decomposition \_\_\_\_\_.

## 6. [MSQ]

Consider the following two schedules  $S_1$  and  $S_2$  involving three transactions  $T_1, T_2$  and  $T_3$  as:

$S_1: R_2(x)W_1(y)W_1(z) \ R_1(w) \ R_3(y) \ R_2(y) \ R_1(x) \ R_3(w) \ R_2(z) \ W_3(x) \ R_2(w) \ W_1(x)$

$S_2: R_1(x) \ R_3(w) \ W_1(y) \ W_1(z) \ R_1(w) \ W_3(y) \ R_3(y) \ R_2(y) \ R_2(x) \ R_2(z) \ R_2(w)$

Which of the following options is/are correct regarding the above two schedules  $S_1$  and  $S_2$ ?

- (a)  $S_1$  is conflict serializable.  
 (b)  $S_1$  is not conflict serializable.  
 (c)  $S_2$  is conflict serializable.  
 (d)  $S_2$  is not conflict serializable.

## 7. [NAT]

Consider the following relational schema with given FD's

**Schema I:**  $R(ABCD)$  and FD's are  $[AB \rightarrow C, C \rightarrow D, D \rightarrow A]$

**Schema II:** R(ABCDE) and FD's are [ $AB \rightarrow C$ ,  
 $C \rightarrow D$ ,  $D \rightarrow B$ ,  $D \rightarrow E$ ]

**Schema III:** R(ABCDE) and FD's are [ $A \rightarrow B$ ,  
 $B \rightarrow C$ ,  $C \rightarrow D$ ,  $D \rightarrow E$ ,  $D \rightarrow A$ ]

**Schema IV:** R(ABCDE) and FD's are [ $B \rightarrow C$ ,  
 $B \rightarrow D$ ,  $D \rightarrow E$ ]

How many above schema satisfy the third normal form (3NF) is\_\_\_\_\_?

**8. [MSQ]**

Consider the relation R(ABCDE) with the following set of functional dependencies

**F:** [ $A \rightarrow D$ ,  $AB \rightarrow C$ ,  $AD \rightarrow CE$ ,  $B \rightarrow C$ ,  $D \rightarrow A$ ,  
 $D \rightarrow B$ ]

Then which of the following is/are True for R?

- (a) R is in 1NF
- (b) R is in 2NF
- (c) R is in 3NF
- (d) R is Not in 3NF & Not in BCNF

## Answer Key

- |    |           |    |           |
|----|-----------|----|-----------|
| 1. | (b)       | 6. | (b, c)    |
| 2. | (a, b, c) | 7. | (2)       |
| 3. | (1254)    | 8. | (a, b, d) |
| 4. | (b, d)    |    |           |
| 5. | (4)       |    |           |

## Hints and Solutions

1. (b)

F: { $AB \rightarrow C$ ,  $C \rightarrow D$ ,  $D \rightarrow E$ ,  $E \rightarrow A$ ,  $D \rightarrow B$ }

Candidate keys = [AB, EB, C, D]

Prime / key Attribute = [A, B, C, D, E]

Here all attribute of relation R is key/prime attribute.

So R is in 3NF

Here R is not in BCNF because in FD

$E \rightarrow A$ : E is not super key.

So R is in 1NF, 2NF and 3NF.

So highest normal form satisfy is 3NF.

2. (a, b, c)

FD's: { $RS \rightarrow T$ ,  $RS \rightarrow U$ ,  $U \rightarrow R$ ,  $ST \rightarrow U$ ,  $ST \rightarrow V$ }

In this question candidate keys = [RS, SU, ST]

In the FD  $U \rightarrow R$ , U is not a super key.

So Relation is Not in BCNF but in 3NF.

So, this relation is in 1NF, 2NF & 3NF also.

3. (1254)

3 Transaction  $T_1$ : 2 operation ( $n_1$ )

$T_2$ : 3 operation ( $n_2$ )

$T_3$ : 4 operation ( $n_3$ )

Total number of Concurrent Schedule

$$= \frac{(n_1 + n_2 + n_3)!}{(n_1!)(n_2!)(n_3!)} \\ = \frac{(2 + 3 + 4)!}{(2!)(3!)(4!)} = \frac{9!}{2 \times 6 \times 4!} \\ = \frac{9 \times 8 \times 7 \times 6 \times 5 \times 4!}{2 \times 6 \times 4!}$$

$$\boxed{\text{Total No. of Concurrent Schedule} = 1260}$$

Total No. of serial schedule =  $3! = 6$  Serial Schedule.

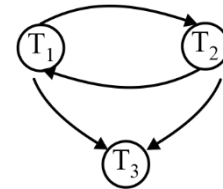
Total Number of =  $1260 - 6$

Non-Serial Schedule = 1254

4. (b, d)

Schedule: S1

$T_1$	$T_2$	$T_3$
r(x)		
	w(x)	
w(x)		
		w(x)
		r(x)



Cycle exists, therefore S1 not conflict serializable.

Schedule S2: S2 is also not conflict serializable as precedence graph contains cycle.

5. (4)

R (ABCDEFGH) { $A \rightarrow B$ ,  $C \rightarrow D$ ,  $E \rightarrow FGH$ }

Candidate key = [ACE]

Check BCNF = ?

$A \rightarrow B$ ,  $C \rightarrow D$ ,  $E \rightarrow FGH$  violates BCNF

Because  $\underline{x} \rightarrow y$  : x is Not super key

So, R Not in BCNF.

BCNF Deconposition

$R_1(AB)$   $R_2(CD)$   $R_3(EFGH)$   $R_4(ACE)$

Now its in BCNF + losseless join + Dependency Preserving.

6. (b, c)

Sol.

S1:  $R_2(x)W_1(y)W_1(z)$   $R_1(w)$   $R_3(y)$   $R_2(y)$   $R_1(x)$   $R_3(w)$   
 $R_2(z)$   $W_3(x)$   $R_2(w)$   $W_1(x)$

S2:  $R_1(x)$   $R_3(w)$   $W_1(y)$   $W_1(z)$   $R_1(w)$   $W_3(y)$   $R_3(y)$   $R_2(y)$   
 $R_2(x)$   $R_2(z)$   $R_2(w)$

2)

**Schema I:** R(ABCD) and FD's are  $[AB \rightarrow C, C \rightarrow D, D \rightarrow A]$

Candidate keys =  $[AB, DB, CB]$

Relation R is in 1NF, 2NF & 3NF But Not in BCNF because  $C \rightarrow D, D \rightarrow A$  violate BCNF definition.

**Schema II:** R(ABCDE) and FD's are  $[AB \rightarrow C, C \rightarrow D, D \rightarrow B, D \rightarrow E]$

Candidate keys =  $[AB, AD, AC]$

Here  $D \rightarrow E$  is a Partial Dependency so R Not in 2NF.

**Schema III:** R(ABCDE) and FD's are  $[A \rightarrow B, B \rightarrow C, C \rightarrow D, D \rightarrow E, D \rightarrow A]$

Candidate key =  $[A, B, C, D]$

So, R is in 2NF But Not in 3NF & Not in BCNF

