



## *Practical No. 9*

**Aim:** To implement Chinese Remainder Theorem to perform primality testing.

**Theory:**

The Chinese Remainder Theorem (which will be referred to as CRT in the rest of this article) was discovered by Chinese mathematician Sun Zi.

### Formulation

Let  $p = p_1 \cdot p_2 \cdots p_k$ , where  $p_i$  are pairwise relatively prime. In addition to  $p_i$ , we are also given a set of congruence equations

$$a \equiv a_1 \pmod{p_1}$$

$$a \equiv a_2 \pmod{p_2} \dots$$

$$a \equiv a_k \pmod{p_k}$$

where  $a_i$  are some given constants. The original form of CRT then states that the given set of congruence equations always has one and exactly one solution modulo  $p$ .

### Corollary

A consequence of the CRT is that the equation

$$x \equiv a \pmod{p}$$

is equivalent to the system of equations

$$x \equiv a_1 \pmod{p_1}$$

...

$$x \equiv a_k \pmod{p_k}$$

(As above, assume that  $p = p_1 p_2 \cdots p_k$  and  $p_i$  are pairwise relatively prime).

### Garner's Algorithm

Another consequence of the CRT is that we can represent big numbers using an array of small integers. For example, let  $p$  be the product of the first 1000 primes. From calculations we can see that  $p$  has around 3000 digits.



**PRIYADARSHINI COLLEGE OF ENGINEERING**  
(Recognised by A.I.C.T.E., New Delhi & Govt. of Maharashtra, Affiliated to  
R.T.M.Nagpur University) Near CRPF Campus, Hingna Road,  
Nagpur-440 019, Maharashtra (India)  
Phone : 07104 - 236381, 237307, Fax : 07104 - 237681,



Any number  $a$  less than  $p$  can be represented as an array  $a_1, \dots, a_k$ , where  $a_i \equiv a \pmod{p_i}$ . But to do this we obviously need to know how to get back the number  $a$  from its representation. In this section, we discuss Garner's Algorithm, which can be used for this purpose. We seek a representation on the form

$$a = x_1 + x_2 \cdot p_1 + x_3 \cdot p_1 \cdot p_2 + \dots + x_k \cdot p_1 \cdot p_2 \cdot \dots \cdot p_{k-1}$$

which is called the mixed radix representation of  $a$ . Garner's algorithm computes the coefficients  $x_1, \dots, x_k$ .

Let  $r_{ij}$  denote the inverse of  $p_i$  modulo  $p_j$

$$r_{ij} = (p_i)^{-1} \pmod{p_j}$$

which can be found using the algorithm described in Modular Inverse. Substituting  $a$  from the mixed radix representation into the first congruence equation we obtain

$$a_1 \equiv x_1 \pmod{p_1}.$$

Substituting into the second equation yields

$$a_2 \equiv x_1 + x_2 p_1 \pmod{p_2}.$$

which can be rewritten by subtracting  $x_1$  and dividing by  $p_1$  to get

$$a_2 - x_1 \equiv x_2 p_1 \pmod{p_2}$$

$$(a_2 - x_1) r_{12} \equiv x_2 \pmod{p_2}$$

$$x_2 \equiv (a_2 - x_1) r_{12} \pmod{p_2}$$

Similarly we get that

$$x_3 \equiv ((a_3 - x_1) r_{13} - x_2) r_{23} \pmod{p_3}.$$

**Conclusion:** Chinese Remainder Theorem to perform primality testing is implemented successfully.

**Viva Questions:**

- Q. 1 What is Concept of co-prime?
- Q. 2 What is use of Chinese Remainder Theorem in Cryptography?
- Q. 3 What is concept Chinese Remainder Theorem?