

Q.1 Explain additive cipher? Use additive cipher with key 2 to encrypt message "DEPARTMENT".

$\rightarrow A \sim 0, B \sim 1, C \sim 2, D \sim 3, E \sim 4, F \sim 5, G \sim 6, H \sim 7, I \sim 8, J \sim 9$

$K \sim 10, L \sim 11, M \sim 12, N \sim 13, O \sim 14, P \sim 15, Q \sim 16, R \sim 17, S \sim 18, T \sim 19$

$U \sim 20, V \sim 21, W \sim 22, X \sim 23, Y \sim 24, Z \sim 25$

plaintxt = 0 3 4 15 0 17 19 12 4 13 19
 key = 2 2 2 2 2 2 2 2 2 2 2

key = 2

Encryption Using additive cipher : →

Ciphertext = (plaintxt + key) mod 26

$$\textcircled{1} \quad C = (P+K) \mod 26 \quad \text{as } P = 0, K = 2 \Rightarrow C = (0+2) \mod 26$$

$$C = (0+2) \mod 26 \quad \text{as } 0+2 = 2 \mod 26$$

$$= 2 \mod 26$$

$$= \textcircled{2} \text{ (Ans)}$$

$$\text{plaintxt} = 5 \sim F \text{ from } \text{table} \rightarrow \text{Ans} \text{ is } 14 \sim O$$

$$\text{as } 5+9 = 14 \mod 26$$

$$\textcircled{2} \quad C = (E+K) \mod 26$$

$$= (4+2) \mod 26$$

$$= 6 \mod 26$$

$$= 6 \sim G$$

$$\textcircled{3} \quad C = (A+K) \mod 26$$

$$as \ 1+2 = 3 \mod 26$$

$$as 1+2 = 3 \mod 26$$

$$= 3 \sim C$$

$$\textcircled{5} \quad C = (R+K) \bmod 26$$

$$= (19+2) \bmod 26$$

$$= 14 \bmod 26$$

$$\textcircled{7} \quad C = (m+k)c \bmod 26$$

$$= (12+2) \bmod 26$$

$$= 14 \bmod 26$$

$$= 14$$

$$\textcircled{6} \quad C = (T+K) \bmod 26$$

$$= (19+2) \bmod 26$$

$$\textcircled{8} \quad C = (N+K) \bmod 26$$

$$= (13+2) \bmod 26$$

$$P = 215 \bmod 167$$

$$= 21 \sim V$$

$$= 15 \bmod 26$$

$$= 15 \sim P$$

S = 193

 \therefore Cipher Text = FLRLTVODHPV.

* Additive cipher \rightarrow (just + extanting) = fixed length

① An additive cipher is a type of mono-alphabetic

cipher that shifts the plaintext alphabets by a fixed amount to obtain the cipher alphabet.

② It is also referred as 'shift cipher' or 'Caesar cipher'.

③ As name suggests, addition modulus 2 operation is performed on plain text to obtain a cipher

as text \rightarrow ④

as form (x_1, x_2, \dots) ⑤

⑥ $C = (P+K) \bmod 26$

⑦ $C = (x_1 + k_1, x_2 + k_2, \dots)$

$P = (C-K) \bmod 26$

⑧ $P = (x_1 - k_1, x_2 - k_2, \dots)$

where,

P = plaintext / message in the natural language or plain text

C = ciphertext / ciphered text formed "111111111"

K = key

③ It is not very secure. It can be broken by brute force attack.

④ It means the message encrypted by this method can easily be decrypted.

⑤ It is the weak method of Cryptography.

⑥ Advantages:

1) Very easy to implement

2) Simplest method, and only one key involved in its entire process.

3) It requires only few computing resources.

⑦ Disadvantages:

1) It can be easily hacked.

2) It provides very little security.

3) By looking at the pattern of letters in it, the entire message can be decrypted easily.

A B C D E F G H I J K L M N O P Q R S T
F1 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

A B C D E F G H I J K L M N O P Q R S T
F1 11 12 13 14 15 16 17 18 19 10 11 12 13 14 15

Question 2. Explain Substitution techniques? Using the key "COMPUTER" encrypt the message "TRANSFER" by using Vigenere cipher.

→ * Substitution technique → ~~using four methods~~

- (1) Substitution technique is a classical encryption technique where the characters present in the original message are replaced by other characters or numbers or symbols.
- (2) If the plain text is considered as the string of bits, then the substitution technique would replace bit patterns of plain text with the bit pattern of cipher text.
- (3) Some substitution techniques are:
 - a) Caesar cipher.
 - b) Monoalphabetic cipher.
 - c) Polyalphabetic cipher.
 - d) Playfair cipher.
 - e) Hill cipher.
 - f) One-Time pad.

*

Key = C O M P U T E R
2 14 12 15 20 19 4 17

plaintext = T R A N S F E R
14 17 0 13 18 5 4 17

Encryption : $(N-3) \cdot 9 \oplus$

$$\begin{aligned} \textcircled{1} \quad c &= (p+k) \bmod 26 \\ &= (2+19) \bmod 26 \\ &= 21 \bmod 26 \quad \text{Eq. } \textcircled{2} \\ &= 21 \sim V \quad p = 2 \end{aligned}$$

$\rightarrow \text{J. I. K. M. L. N. T. P.}$

$$\begin{aligned} \textcircled{6} \quad c &= (p+k) \bmod 26 \\ &= (19+5) \bmod 26 \\ &= 24 \bmod 26 \\ &= 24 \sim Y \end{aligned}$$

$\rightarrow \text{J. I. K. M. L. N. T. P.} \sim Y \text{ Eq. } \textcircled{3}$

$$\begin{aligned} \textcircled{2} \quad c &= (p+k) \bmod 26 \quad \text{Eq. } \textcircled{4} \\ &= (14+17) \bmod 26 \\ &= 31 \bmod 26 \\ &= 5 \sim F \end{aligned}$$

$$\begin{aligned} \textcircled{7} \quad c &= (p+k) \bmod 26 \\ &= 0 \sim \text{mod 26} \\ &= 8 \sim I \end{aligned}$$

$$\begin{aligned} \textcircled{3} \quad c &= (p+k) \bmod 26 \\ &= 12 \bmod 26 \\ &= 12 \sim M \end{aligned}$$

$$\begin{aligned} \textcircled{8} \quad c &= (p+k) \bmod 26 \\ &= (17+17) \bmod 26 \\ &= 34 \bmod 26 \\ &= 8 \sim I \end{aligned}$$

$$\begin{aligned} \textcircled{4} \quad c &= (p+k) \bmod 26 \\ &= (5+13) \bmod 26 \\ &= 28 \bmod 26 \\ &= 2 \sim C \end{aligned}$$

ciphertext = VFNMCLM4II

$$\begin{aligned} \textcircled{5} \quad c &= (p+k) \bmod 26 \\ &= (20+18) \bmod 26 \\ &= 38 \bmod 26 \\ &= 12 \sim M \end{aligned}$$

$\rightarrow \text{J. I. K. M. L. N. T. P.} \sim Y \text{ Eq. } \textcircled{1}$

$\rightarrow \text{J. I. K. M. L. N. T. P.} \sim Y \text{ Eq. } \textcircled{2}$

$\rightarrow \text{J. I. K. M. L. N. T. P.} \sim Y \text{ Eq. } \textcircled{3}$

Description :

$$\begin{aligned} \textcircled{1} \quad p &= (c - k) \bmod 26 \\ &= (21 - 2) \bmod 26 \\ &= 19 \bmod 26 \\ &= 19 \sim T \end{aligned}$$

$$\begin{aligned} \textcircled{8} \quad p &= (8 - 17) \bmod 26 \\ &= 17 \sim R \end{aligned}$$

\therefore plaintext = TRANSFER.

$$\begin{aligned} \textcircled{2} \quad p &= (5 - 14) \bmod 26 \\ &= 17 \sim R \end{aligned}$$

$$\begin{aligned} \textcircled{3} \quad p &= (12 - 6) \bmod 26 \\ &= 0 \bmod 26 \\ &= 0 \sim A \end{aligned}$$

$$\begin{aligned} \textcircled{4} \quad p &= (2 - 15) \bmod 26 \\ &= 13 \sim N \end{aligned}$$

$$\begin{aligned} \textcircled{5} \quad p &= (12 - 20) \bmod 26 \\ &= 18 \sim S \end{aligned}$$

$$\begin{aligned} \textcircled{6} \quad p &= (24 - 19) \bmod 26 \\ &= 5 \bmod 26 \\ &= 5 \sim F \end{aligned}$$

$$\begin{aligned} \textcircled{7} \quad p &= (8 - 4) \bmod 26 \\ &= 4 \bmod 26 \\ &= 4 \sim E \end{aligned}$$

Question 3. Explain the substitution Techniques? if the cipher is keyed by a word "COMPUTER" and plaintext is "please transfer one million dollars to my account six six two obtain the cipher text by playfair substitution methods?"



* Substitution Techniques : →

- (1) A Substitution Cipher simply means that each letters in the plaintext is substituted with another letter to form the ciphertext.
- (2) Substitution technique is a classical encryption approach where the characters present in the initial message are restored by the other characters or numbers or by symbol.
- (3) There are various types of substitution ciphers :
 - a) monalphabetic cipher.
 - b) polyalphabetic cipher.

Plaintext : PLEASE TRANSFER ONE MILLION DOLLARS
TO MY ACCOUNT SIX SIX TWO .

Key = COMPUTER

Q	O	P	N	V	W	H	I	J	K	L	M	F	G	B	R	A	T	U	S	Z
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

← : машина Г сеит виноград

PL	EA	SE	TR	AN	SF	ER	OM	EM	JL	LI
OF	RB	LB	EA	RQ	LI	RA	ML	RO	FS	SF

ON	DO	LEX	EA	RS	HTB	МУ	AL	LO	UN	TS
ML	FC	MW	QE	BN	EC	PX	TR	Om	MS	BK

IX	SI	XT	W	W	W	W	W	W	W	W
GZ	ZJ	VR	OE	W	W	W	W	W	W	W

Слідчий текст: ОФРBLB EAQLERA MIR DFSSFML

Інформація: FCNWFEBN EC PX TROMMSB

Код: GZ ZJ VR OE

Q.4 One time pad :

plaintext = SUBJECT \oplus key = \oplus 230 \oplus 11000 (0. 1000)

key = abcdefg \oplus 230 \oplus 11000 (0. 1000)

\rightarrow plaintext = SUBJECT \oplus key = \oplus 230 \oplus 11000 (0. 1000)

plaintext = SUBJECT \oplus key = \oplus 230 \oplus 11000 (0. 1000)

key = abcdefg \oplus adding g to key, because length

is 6 digits in 230 for plaintext so key must be
6 digits in 230 for ciphertext

A B C D E F G H I J K L M N O P
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

Q R S T U V W X Y Z
16 17 18 19 20 21 22 23 24 25

plain text = S U B T E C T
18 20 1 9 4 2 19

key = a b c d e f g
0 1 2 3 4 5 6

first 8 pt

Add: 18 18 21 25 8 12 8 78 25

(plaintext + key)
(addition)

first 8 pt

all 6 numbers

Subtract: 18 21 3 12 8 7 25

(subtract if no. is
greater than 26)

all 6 numbers

Cipher text = S V D M T H Z
18 21 3 12 8 7 25

270 27

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Ques. 5) Explain DES in detail.

→ ① DES stands for Data Encryption Standard.

② The DES algorithm uses a key of 56 bit size.

③ Using the key, the DES takes a block of 64 bit plain text as input and generates a block of 64 bit cipher text.

④ General structure of DES is depicted in the following illustration:

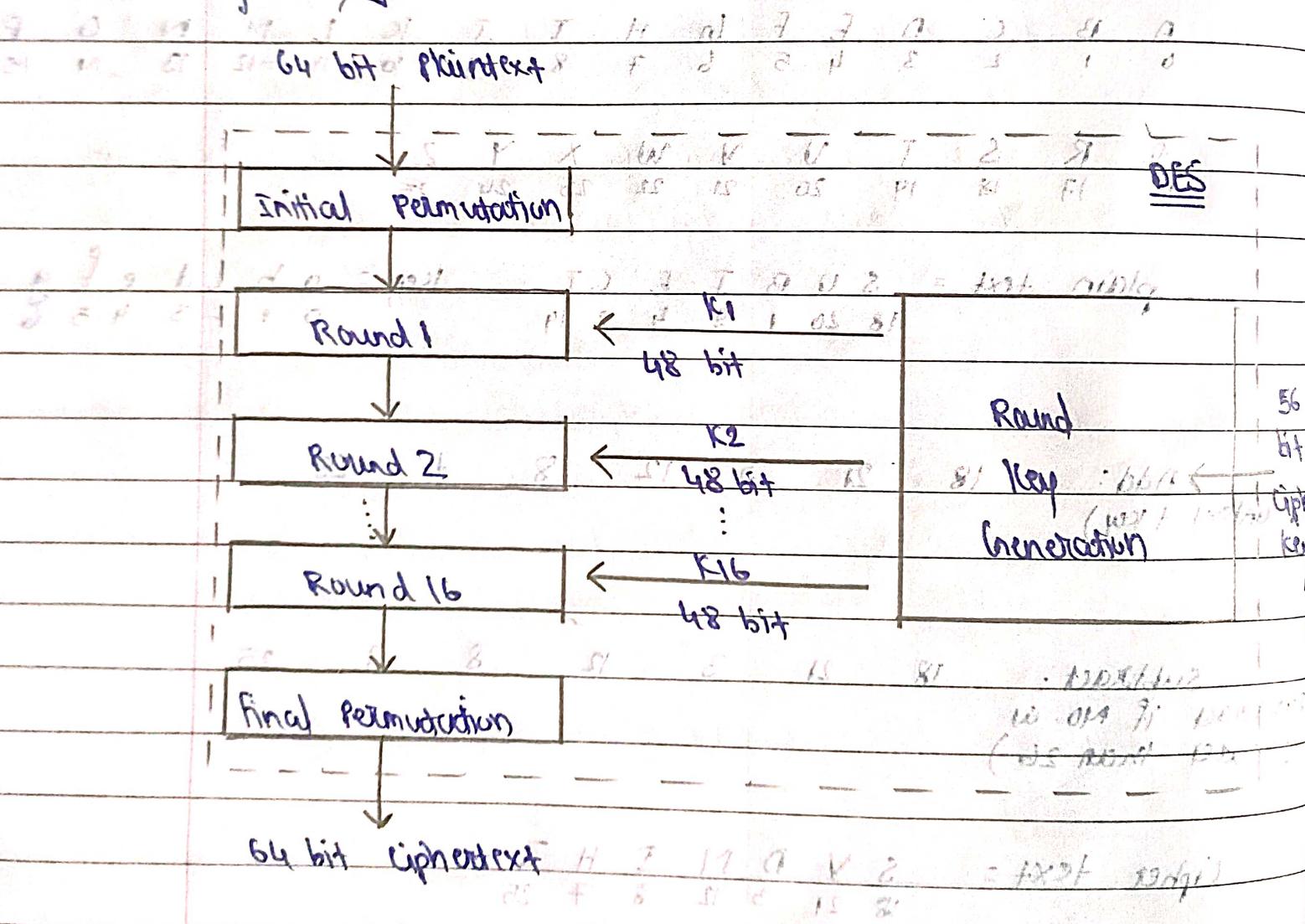


Fig: DES

- (5) DES is an implementation of Feistel cipher. So it uses 16 rounds Feistel structure. It consists of -
- (6) Since DES is based on the Feistel cipher, all that is required to specify DES is -
- Round functions
 - Key Schedule
 - Any additional processing (initial and final) permutations.
- (7) The encryptor and the decryptor need to use the same key. Otherwise they will work not be able to communicate together.
- (8) The decryption process is the logical opposite of the encryption. It takes in a 64 bit block of ciphertext and produces the 64-bit block of plaintext using the same 48 bit key during encryption.

Q.6 Difference between →

- Private & Public key in cryptography.
- Conventional vs public key encryption.
- Symmetric vs asymmetric cryptography.

→ In 230 pages of this book

Private key cryptography

① Private key is faster than the public key.

② In this, the same key & algo. are used to encrypt and decrypt in the message.

③ private keys in symmetrical because there is only one key that is called secret key.

④ It is efficient technology.

⑤ It is used for large amount of text.

Public key cryptography

① It is slower than private key.

② In this, two keys are used.

③ public keys in asymmetric because there are two keys.

④ It is an inefficient technology.

⑤ It is used for only short message.

Q.7 Explain a AES algorithm ? How is different from DES algorithm ? in details.

→ AES algorithm : →

① AES : Advance Encryption standard.

② AES is block cipher, encrypts data in blocks of 128 bit each.

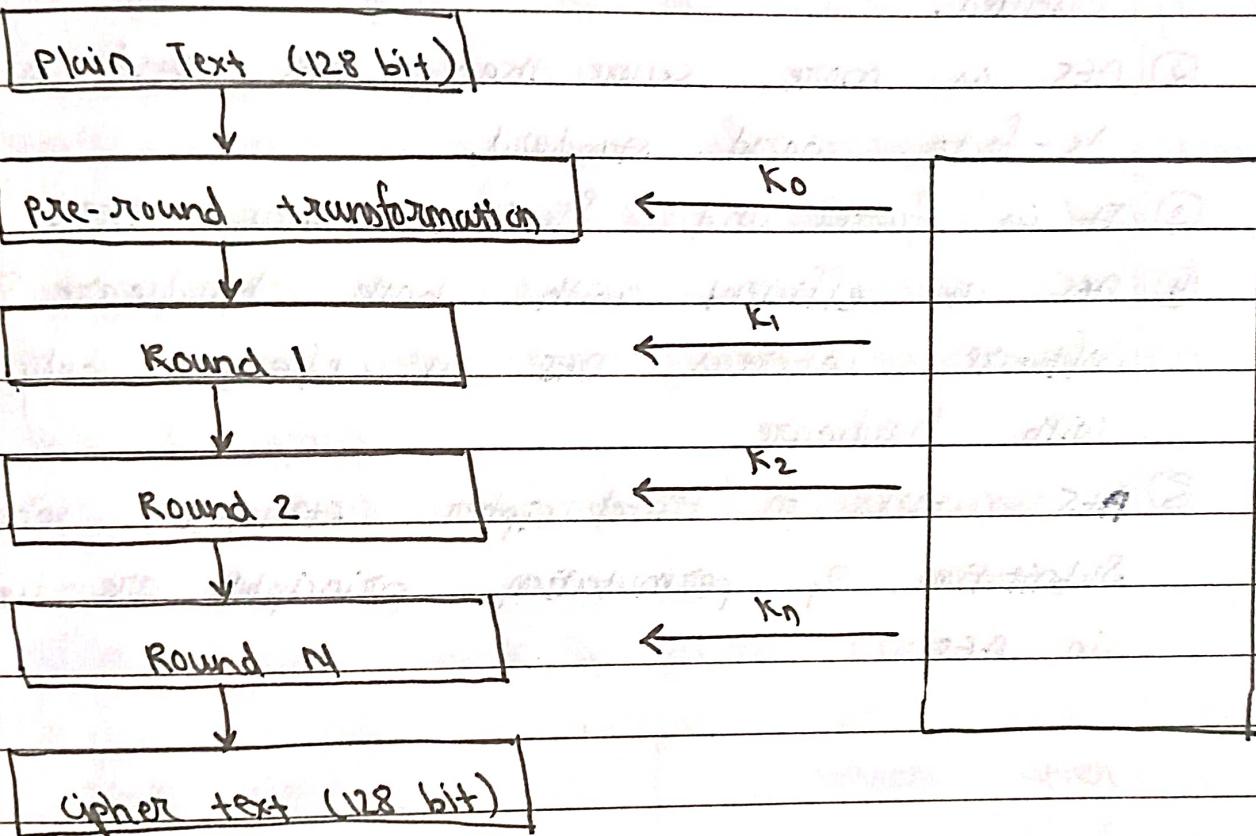


Fig: Creation of Round keys.

- ③ It converts individual block using key of 128, 192 & 256 bits.
- ④ Once it encrypts these blocks, it joins them together to form the ciphertext.

⑤ AES implemented in software throughout the world to encrypt sensitive data.

* AES and DES comparison:

① AES is byte oriented whereas DES is bit oriented.

② AES is more secure than DES and is the de-facto world standard.

③ It is faster and flexible than DES.

④ AES is efficient with both hardware & Software whereas DES is efficient only with hardware.

⑤ DES works on Feistel cipher structure whereas Substitution & permutation principles are used in AES.

Q.8 Explain a block cipher principle? what is difference between stream cipher and block cipher?

→ * Block cipher : →

- ① Block ciphers are built in the Feistel Cipher structure.
- ② It has a specific no. of rounds and keys for generating ciphertext.
- ③ Block cipher is a type of encryption algorithm that processes fixed-size blocks of data, usually 64 or 128 bits, to produce ciphertext.
- ④ Block cipher uses same key for encryption and decryption.
- ⑤ DES is best example of block cipher.

* Difference betⁿ block & stream cipher.

	Block cipher	stream cipher
①	It is kind of encryption that converts plaintext by taking each block individual by.	It is the kind of encryption that converts plaintext by taking one bytes of the plaintext at a time.
②	feistel cipher is the most popular block cipher implementation.	vernam cipher is the main implementation of stream cipher.

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③ It is difficult to reverse encrypted text

It uses XOR encryption, which is easily reversed to plaintext.

④ uses both confusion & diffusion

Relies on confusion only

⑤ ECB and CBC modes are used

CFB and OFB modes are used.

⑥ Complexity is simple

Complexity is complex

⑦ 64 bits or more used

Used 8 bits

Q.9) write RSA algorithm to perform Encryption using RSA algo., for the following
 $p=7$
 $q=11$ $e=3$ $M=9$.

→ (1) RSA algorithm is an asymmetric cryptography algorithm.

(2) Asymmetric actually means that it works on two different keys i.e. public and private key.

(3) public key is given to everyone and the private key kept private.

- ④ The public key is used for encryption and the private key is used for decryption.
- ⑤ RSA is most common public-key algorithm, named after its inventors Rivest, Shamir and Adleman (RSA).

(6)

$$p = 7 \quad q = 11 \quad r = 3 \quad M = 9$$

$$n = p * q$$

$$\therefore n = 7 * 11 = 77$$

$$\phi(n) = (p-1) * (q-1)$$

$$\therefore \phi(n) = (7-1) * (11-1) = 6 * 10 = 60$$

$$d = e^{-1} \pmod{\phi(n)}$$

$$= 3^{-1} \pmod{60}$$

$$\therefore d = 27 \pmod{60}$$

$$\therefore d = 27$$

To encrypt a message,

$$c = m^n \pmod{n}$$

$$= 9^n 3 \pmod{77}$$

$$= 729 \pmod{77}$$

$$\therefore c = 44$$

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Q.10) write note on : ECC

- (1) ECC means Elliptic curve cryptography.
- (2) It is a public-key encryption technique that uses pairs of public and private keys to encrypt and decrypt data.
- (3) ECC is based on the mathematical theory of elliptic curves.
- (4) It uses smaller keys than other cryptography methods to provide the same level of security.
- (5) The encryption process takes less time in ECC and decryption process takes more time.

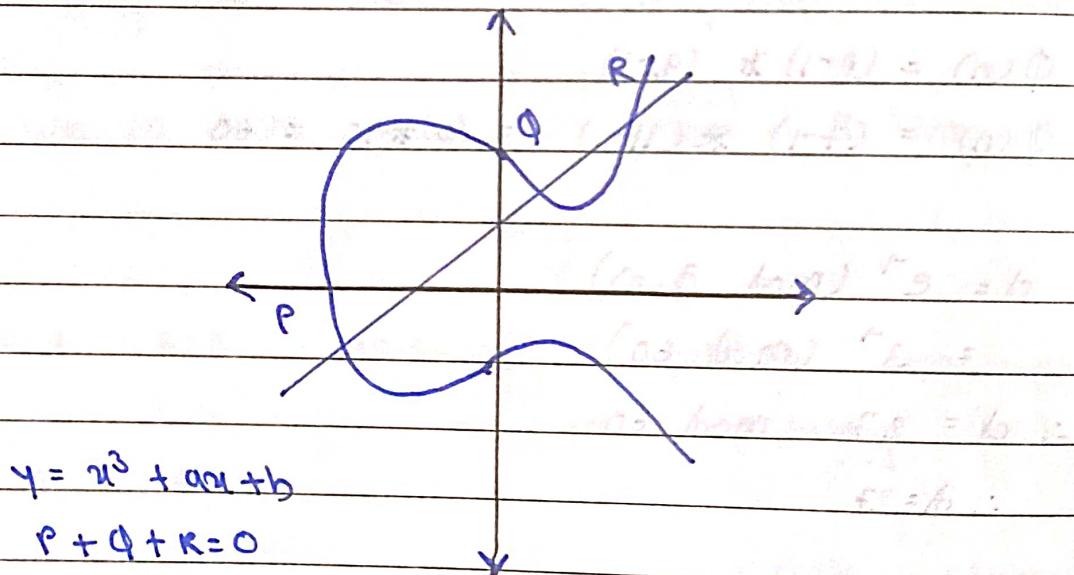


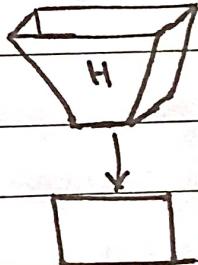
Fig : Elliptic curve cryptography

- (6) websites make extensive use of ECC to secure consumers hypertext transfer protocol connection.
- (7) ECC encryption is smaller, safer and faster.

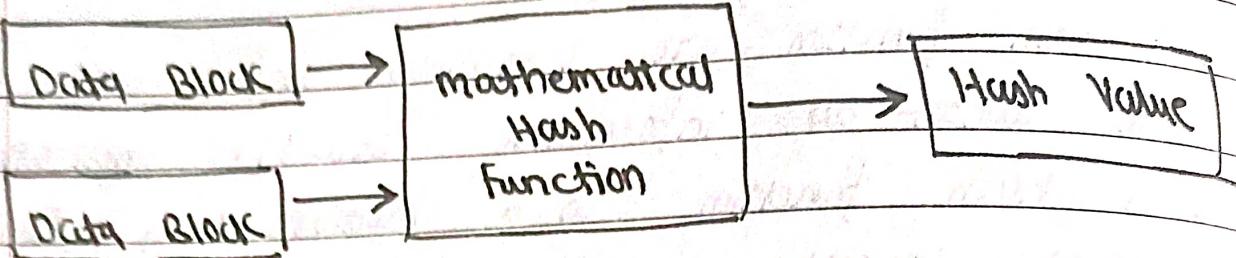
Q.11) Explain Hash function:

- ① Hash function are extremely useful and appear in almost all information security application.
- ② A hash function is a mathematical function that converts a numerical input value into another compressed numerical value.
- ③ The input to the hash function is of arbitrary length but output has fixed length always.
- ④ Values returned by a hash function are called message digest or simply hash values.

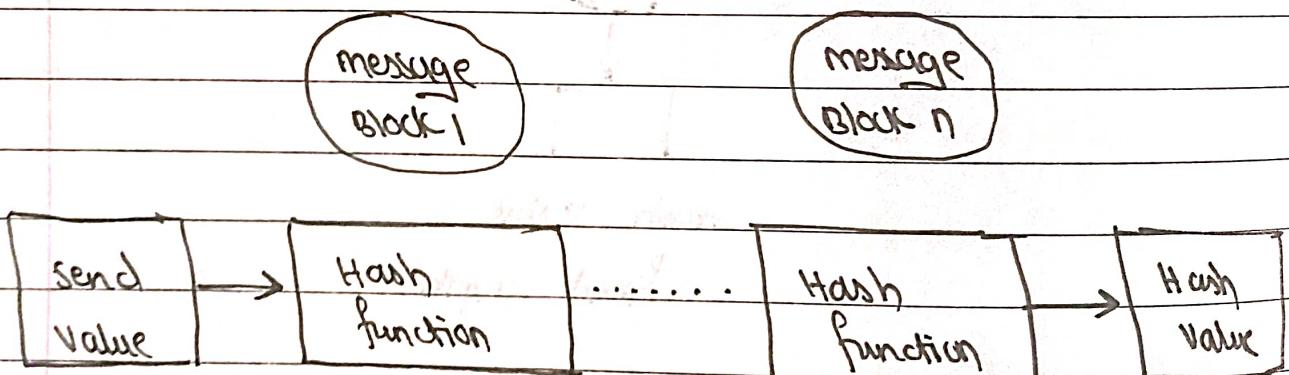
Message M (arbitrary length)



- ⑤ Converting hash codes to their original value is an impossible task to perform. This is the main difference b/w encryption & hash function.



- ⑥ The size of each block varies depending on the algo.
- ⑦ Typically the block sizes are from 128 bits to 512 bits.
- ⑧ Hashing algo, involves rounds of above hash function like a block cipher.
- ⑨ The process is repeated for as many round as we required to hash the entire message..



Q.12) Describe Diffie - Hellman key exchange algorithm.

- ① The diffie - hellman key exchange algorithm is also known as exponential key exchange.
- ② It is method for securely exchanging cryptographic keys over an insecure channel.
- ③ It is a fundamental building block of many secure communication protocols, including SSL/TLS and SSH.
- ④ It is method for securely exchanging and used to establish a shared communications while exchanging data over a public Network using the elliptic curve to generate points and get the secret key using the parameter.
- ⑤ The key exchange involves the following steps:-
 - i) Alice and Bob agree on two large prime no. p and g , and a public key exchange alg.
 - ii) Alice chooses a secret integer a and computes $A = g^a \text{ mod } p$. She sends A to Bob.
 - iii) Bob chooses a secret integer b and computes $B = g^b \text{ mod } p$. He sends B to Alice.
 - iv) Alice computes $S = B^a \text{ mod } p$. Bob computes $S = A^b \text{ mod } p$.
 - v) Alice and Bob now both have a shared secret keys, which they can use to established a secured communication channel.

Q. 13) Solve modular exponentiation algorithm:

(1) $23^{\text{power 3}} \text{ mod } 30$.

$$\rightarrow 23^3 \text{ mod } 30$$

$$= -7^3 \text{ mod } 30 \quad || \quad 23 \text{ mod } 30 \text{ can be } 23 \text{ or } -7$$

$$= -7^3 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= -7^2 \times -7 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= 49 \times -7 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= -133 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= -133 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= 17 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$= 17 \text{ mod } 30 \quad || \quad \text{not possible}$$

$$(2) 11^{\text{power 23}} \text{ mod } 187$$

$$\rightarrow 11^{23} \text{ mod } 187$$

$$= 11^{23} = 11 * 11 * 11 * 11^{12}$$

$11 * 11 * 11 = 1331$ need to find this modulo 187 we

look at 1331 / 187 which is > 7 . so we can take

$$1331 - 7 * 187$$

which is 22. so $11^3 \text{ mod } 187$ is 22. we can

simplify $11^{23} = (11^3)^7 * 11^2$ because we know

$$11^3 = 22$$

we can write this as $22 * 7 * 11^2$ and 11^2 mod

187 is 121. so now have a no. to modulo.

$$11 * 7 * 121 \text{ mod } 187 = 937 \text{ mod } 187$$

$$= \underline{\underline{88}}$$

Q.14) Explain the following terms: →

i) Trap door and cross site scripting:

→ A back door or trap door, is a hidden entry to a computer device or site that bypasses security measures such as logins and password protections. It works in background and hides from the user. It is very similar to a virus and therefore is quite difficult to detect and completely disable.

Cross-site Scripting (XSS) is a type of computer security vulnerability typically found in web application. XSS enables attackers to inject client-side script into web pages viewed by other users. Attackers inject malicious script into a website. Attackers often initiate an XSS attack by sending a malicious link to user.

② Host based v/s Network based IDS.



HIDS

NIDS

- | ① Host Intrusion Detection system | ① Network Intrusion Detection system. |
|--|--|
| ② It doesn't work in real-time. | ② It operates in real time. |
| ③ Response time is slow. | ③ Response time is fast. |
| ④ As it needs to be installed on every host, the installation process can be tiresome. | ④ Few installation points made it easier to install NIDS. |
| ⑤ HIDS can be installed on each and every computer server. | ⑤ NIDS being concerned with network is installed at places like router or servers as there are the main intersection points in the network system. |

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Q.15 Short Notes on :

① PGP : →

- ① PGP stands for pretty good privacy, which is invented by phil zimmermann.
- ② PGP was design to provide all four aspect of security ie. privacy, integrity, authentication & non-repudiation in sending of emails.
- ③ PGP uses digital signatures to provide integrity, non-repudiation and authentication.
- ④ PGP uses a combination of secret key encryption & public key encryption to provide privacy.
- ⑤ Therefore, we can say digital signature uses one hash function, one secret key and two private public key pairs.
- ⑥ PGP is an open source and freely available software package for email security.
- ⑦ Services offered by PGP are:
 - a) Authentication
 - b) Confidentiality
 - c) compression
 - d) Email compatibility
 - e) Segmentation.

② SQL injection: →

① SQL injection is a type of security vulnerability that occurs when an attacker is able to manipulate an SQL query in a way that allows them to gain unauthorized access to a database.

② This can lead to unauthorised access, data-manipulation and potentially even data deletion.

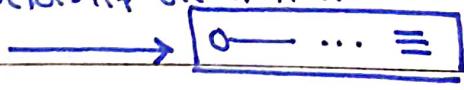
③ SQL injection is a common and dangerous attack vector, which web developers and database administrators need to be aware of it and take steps to prevent it.

④ A successful SQL injection attack can result in unauthorised access to sensitive data, such as:

password, credit score details, personal user info.

`http://teacher.com?`
`teacherid=117 OR 1=1...`

`Select * from teachers
where teacherid = 117 OR 1=1;`



Data for any
teacher is returned
to the attacker

Return data for
all teachers

③ Chip Card Transaction :→

- ① chip cards use cryptography to generate unique transaction codes that allow the terminal to authenticate the card.
- ② The card uses its private key to generate a digital signature of the transaction details and send this back to the reader.
- ③ The card's chip contains a digital certificate that is used to verify the authenticity of card.
- ④ During a transaction, the card generates a cryptogram that is based on the transaction data and a secret key stored on the card.
- ⑤ A chip card is a debit or credit card that contains an embedded microchip along with the traditional magnetic stripe.

Q.16) Chinese problem:

Apply the Chinese Remainder Theorem to solve
following congruent equation.

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

Up

$$m_1 = 3, m_2 = 5, m_3 = 7$$

$$a_1 = 2, a_2 = 3, a_3 = 2$$

i) Common modulus M:

$$M = m_1 * m_2 * m_3$$

$$= 3 * 5 * 7$$

$$\Delta = 105$$

$$ii) M_1 = \frac{m}{m_1} = \frac{105}{3} = 35$$

$$M_2 = \frac{m}{m_2} = \frac{105}{5} = 21$$

$$M_3 = \frac{m}{m_3} = \frac{105}{7} = 15$$

3) $m_1^{-1} = 35^{-1} \pmod{3} = 2$

$m_2^{-1} = 35^{-1} \pmod{5} = 1$

$m_3^{-1} = 35^{-1} \pmod{7} = 1$

4)

$$x = (a_1 * m_1 * m_1^{-1}) + (a_2 * m_2 * m_2^{-1}) + (a_3 * m_3 * m_3^{-1}) \pmod{105}$$

$\pmod{105}$

$$= ((2 * 35 * 2) + (0 * 3 * 21 * 1) + (2 * 15 * 1)) \pmod{105}$$

$$= (140 + 63 + 30) \pmod{105}$$

$$= 233 \pmod{105}$$

$$= 23$$

$$\therefore 23 \pmod{3} = 2$$

$$\therefore 23 \pmod{5} = 3$$

$$\therefore 23 \pmod{7} = 2$$

} cross checking ...