OS-ASSIGNMENT 2 BIE7371 Radha Kerisma Garg four Conditionals of deadlock: a) mutual exclusion: All resources are non-shareable 6) No-premption: process should not be prempted. c) Hold and mait: All puocess are holding a sussemble and maiting for others. d) eircular mait Dlagram dreadfock mil not hoppen because: It has market at wither 1) The hold and weit conditions is not there because maliting for any susource 2) Circular mait is not happening There are 4 mays to deal with deadlocks 1) Deadlock Avoidance 2) peadlock dynorance 8) Scadlock prevention 4) seadleck detection and recovery () (1)

1 LA OUR STEPPE WHOND MUSICHIA	
Ans I Cultical Section: This refer to the separate on	
process or threads are accessing a shared resource on process or threads are accessing a shared resource on particul section of a code.	
no mitical section of a code.	
The state of the s	
laculy see D	
Regularments.	
and to recommend the first recommendation of the second se	
is lared the bulles seemen.	
2. Europeus: - If no purposes is currently executing in the contribut section of the	
and there are here another material to ever the	
west process to enter the best to	
11. A pt to pt 1.	
3. Bounded wait: A bound must exist on the no. of times other process on made	i.
while called the current face become soften	
a language to surel the number secretar and payor the	
guanted.	-
Solution to 2: process cuitical section is Peterron's Algorithm	E
· int fum =0;	
loool flag [2] = offalse, false };	
1/process of	
Hay [0] = time;	
fun = 1;	1
while (frag DO at and turn ===1);	
Mag. Talz Lalse;	
flog [o]= false;	7
8 11 Process I	E A
lay [v] = true;	
fun 20;	
1011 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
while Ulay CoI at fam == 0);	
Hay W= false;	7

Each purcess gives puriously to its own fleg and set it to true when it is nearly to enter for witheas section. And then give priority to others process it and maits until the other process.

Jimshus or given pulority in cuitical section. Then the fleg is set to false which indicates completion of work by process.

This Algorithm satisfies are the negativenest mentioned. Each perocess gives priority when it is meady to enter to
priority to others process id one
finishes or given priority in
Set to false which indicate
This Algorithm satisfies are the Best tit Algo 1. Process 215 kb in 215 kb partition paulition 2. Process 423 Kb in 520 Kb 3. Process 112 Kb in 25 Kb pautition in 600 kb paulition Process 428 Kb Worst Lit Algorithm 1. Process 215 Kb 600 Kb partition m 520 kb partition 2. Process 423 Kb in 360 les partition 3. Process 112 Kb in 4. Process 426 Kb 600 Kb partition in first fift Algo 1 Process 215 Kb in 215 kb partition in 600 kb partition 2 Process 423 Kb 3 Progess 112 kb In RIS Kb partition 4 Process 426kb in 600 Rb partition Best fit! - Algorithm militaires the amount of member space in

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