1 Concrete Semantics of Scheme CESKt*

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Semantic Domains:
              Syntax Domains:
                                                                                      \varsigma \in \Sigma \triangleq E\langle Eval \rangle + A\langle Apply \rangle
      e \in \mathsf{Exp} ::= x
                    | (if e e e)
                                                                                        Eval \triangleq \mathsf{Exp} \times Env \times Store
                    | (let ([x \ e] ...) \ e) |
                                                                                                  \times KStore \times Addr \times Time
                    |(\operatorname{call/cc} e)|
                                                                                     Apply \triangleq Val \times Store
                    |(\mathtt{set!}\ x\ e)|
                                                                                                  \times KStore \times Addr \times Time
                    | (prim op e ...) |
                                                                                 \rho \in Env \triangleq \mathsf{Var} \rightharpoonup Addr
                    | (apply-prim \ op \ e) |
                                                                               \sigma \in Store \triangleq Addr \rightarrow Val
                    | (apply e e) |
                                                                               \sigma \in Store \triangleq Addr \rightarrow Kont
                    |(e e ...)|
                                                                                   v \in Val \triangleq Clo + Kont + \mathbb{Z}
  \mathbf{x} \in \mathsf{AExp} ::= x \mid lam \mid \mathbb{Z} \mid \mathsf{\#t} \mid \mathsf{\#f}
                                                                                                 + {#t, #f, Null, Void}
                               | (quote e)
                                                                                                 + \{ \mathbf{quote}(e), \mathbf{cons}(v, v) \}
lam \in Lam ::= (\lambda (x...) e) | (\lambda x e)
                                                                               clo \in Clo \triangleq \mathsf{Lam} \times Env
       x \in \mathsf{Var}
                           A set of identifiers
                                                                         a, a_{\kappa} \in Addr \triangleq \mathbb{N}
  op \in \mathsf{Prim}
                           A set of primitives
                                                                               t \in Time \triangleq \mathbb{N}
          Atomic Evaluation:
                                                                               \kappa \in \mathit{Kont} ::= \mathbf{mt}
                     \mathcal{A}: \Sigma_E \times \sigma \rightharpoonup Val
                                                                                                | ifk(e, e, \rho, a)
          \mathcal{A}(\langle lam, \rho, \_, \_ \rangle, \_) \triangleq (lam, \rho)
                                                                                                | \mathbf{callcck}(a) |
\mathcal{A}(\langle (\mathtt{quote}\ e), \_, \_, \_ \rangle, \_) \triangleq \mathtt{quote}(e)
                                                                                                 | setk(x, \rho, a)
             \mathcal{A}(\langle x, \rho, \neg, \neg \rangle, \sigma) \triangleq \sigma(\rho(x))
                                                                                                 | appappk(val?, e, \rho, a)
              \mathcal{A}(\langle x_{2}, x_{2}, x_{2}, x_{2} \rangle, x_{2}) \triangleq x_{2}
                                                                                                 | \mathbf{appk}(done, todo, \rho, a) |
                   Tick/Alloc:
                                                                                                 | appprimk(op, \rho, a)
             tick : Eval \times \mathbb{N} \to Time
                                                                                                | \mathbf{primk}(op, done, todo,
       tick(E\langle -, -, -, t \rangle, n) \triangleq (t + n)
                                                                                                                    \rho, a
                 alloc: \Sigma \times \mathbb{N} \triangleq Addr
                                                                                                | letk(vars, done, todo,
     alloc(E\langle -, -, -, t \rangle, n) \triangleq (t + n)
                                                                                                                    e, \rho, a
                                                                                       done \triangleq Val^*
                      Injection:
                     \mathcal{I}: \mathsf{Exp} \to \Sigma
                                                                                        todo \triangleq \mathsf{Exp}^*
    \mathcal{I}(e) \triangleq (e, \varnothing, \varnothing, \{0 : \mathbf{mt}\}, 0, 1)
                                                                                        vars \triangleq \mathsf{Var}^*
                    Transition:
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Collecting Semantics:

Eval Rules

Rules for when the control is an expression

$$E\langle \mathfrak{X}, \rho, \sigma, \sigma_{\kappa}, a_{\kappa}, t \rangle \rightsquigarrow A\langle v, \sigma, \sigma_{\kappa}, a_{\kappa}, t \rangle$$

where $v \triangleq \mathcal{A}(\varsigma)$

$$E\langle(\text{if }e_c\ e_t\ e_f),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e_c,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ \text{where }a'_\kappa \triangleq ailoc(\varsigma,0)\\ \text{of } \alpha'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{let }()\ e,\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma,a',a'_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma,\alpha_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma,a',a'_\kappa,t'\rangle\\ \text{where }bs = ([x_0\ e_0]\ [x_s\ e_s]\ ...)\\ a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{letk}(x_0::x_s,[\],e_s,e_b,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{call}/\text{cc }e),\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{callcck}(a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{cset}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{callcck}(a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{cappppimk}(\emptyset,e_r,\rho,a_\kappa)\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(\text{set}!\ x\ e),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(e_f\ e_s\ ...),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]\\ E\langle(e_f\ e_s\ ...),\rho,\sigma,\sigma_\kappa,a_\kappa,t\rangle\\ \sim E\langle e,\rho,\sigma,\sigma'_\kappa,a'_\kappa,t'\rangle\\ \text{where }a'_\kappa \triangleq alloc(\varsigma,0)\\ t' \triangleq tick(\varsigma,1)\\ \kappa \triangleq \text{appappk}(\emptyset,e_x,\rho,a_\kappa)\\ \sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$$

Apply Rules

Rules for when the control is a value $\varsigma = A \langle v, \sigma, \sigma_{\kappa}, a_{\kappa}, t \rangle$ $\kappa \triangleq \sigma_{\kappa}(a_{\kappa})$ Proceed by matching on κ

$\mathbf{mt} \leadsto \varsigma$

More Apply Rules

Rules for when the control is a value

$$\begin{aligned} & \mathbf{appapk}(v_f, ., ., a_\kappa') \\ & \sim E\langle e_b, \rho_\lambda', \sigma', \sigma_\kappa, a_\kappa', t' \rangle \\ & \text{where } v_f = ((\lambda \langle x_s...) e_b), \rho_\lambda) \\ & a_i \triangleq alloc(\varsigma, i) \\ & t' \triangleq tick(\varsigma, n) \\ & \rho_\lambda' \triangleq \rho_\lambda [x_0 \mapsto a_0 \dots \\ & a_{n-1} \mapsto a_{n-1}] \\ & \sigma' \triangleq \sigma [a_0 \mapsto v_0 \dots \\ & a_{n-1} \mapsto v_{n-1}] \\ & \Rightarrow E\langle e_b, \rho_\lambda', \sigma', \sigma_\kappa, a_\kappa', t' \rangle \\ & \text{where } v_f = ((\lambda \langle x_s...) e_b), \rho_\lambda) \\ & \Rightarrow \sigma [a_0 \mapsto v_0 \dots \\ & a_{n-1} \mapsto v_{n-1}] \\ & \Rightarrow E\langle e_b, \rho_\lambda', \sigma', \sigma_\kappa, a_\kappa', t' \rangle \\ & \text{where } v_f = ((\lambda \langle x_se_b), \rho_\lambda) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \rho_\lambda' \triangleq \rho_\lambda [x \mapsto a] \\ & \sigma' \triangleq \sigma [a \mapsto v] \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma, 1) \\ & \Rightarrow e^* = alloc(\varsigma, 0) \\ & t' \triangleq tick(\varsigma,$$

2 Abstract Semantics of Scheme CESKt*

Abstract Semantic Domains:

$$\begin{array}{c} \hat{\varsigma} \in \hat{\Sigma} \triangleq E \langle Eval \rangle + E \langle Apply \rangle \\ Eval \triangleq \operatorname{Exp} \times \widehat{Env} \times \widehat{Store} \\ \times \widehat{KStore} \times \widehat{Addr} \times \widehat{Time} \\ Apply \triangleq \widehat{Val} \times \widehat{Store} \\ \times \widehat{KStore} \times \widehat{Addr} \times \widehat{Time} \\ \hat{\rho} \in \widehat{Env} \triangleq \operatorname{Var} \longrightarrow \widehat{Addr} \\ \hat{\sigma} \in \widehat{Store} \triangleq \widehat{Addr} \longrightarrow \mathcal{P}(Val) \\ \hat{\sigma}_{\kappa} \in \widehat{KStore} \triangleq \widehat{Addr} \longrightarrow \mathcal{P}(Kont) \\ \hat{v} \in \widehat{Val} \triangleq \widehat{Clo} + \widehat{Kont} + \mathbb{Z} \\ + \{ \#t, \#f, Null, Void \} \\ + \{ \mathbf{quote}(e), \mathbf{cons}(v, v) \} \\ \widehat{clo} \in \widehat{Clo} \triangleq \operatorname{Lam} \times \widehat{Env} \\ \hat{alloc} : \hat{\Sigma} \times \widehat{Kont} \rightarrow \widehat{Addr} \\ \hat{alloc} : \hat{\Sigma} \times \widehat{Kont} \rightarrow \widehat{Addr} \\ \widehat{alloc} : \hat{\Sigma} \times \widehat{Kont} \rightarrow \widehat{Addr}$$

$$\mathcal{A}: Eval \to \widehat{Val}$$

$$\mathcal{A}(E\langle lam, \hat{\rho}, ... \rangle) \triangleq \{(lam, \hat{\rho})\}$$

$$\mathcal{A}(\mathfrak{X}) \triangleq \{\mathfrak{X}\}$$

Abstract Eval Rules

Rules for when the control is an expression

$$E\langle \mathbf{x}, \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle \leadsto A\langle \hat{v}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$$
where $\hat{v} \triangleq \hat{\mathcal{A}}(\hat{\varsigma}, \hat{\sigma})$

$$E\langle(\mathbf{if}\ e_{c}\ e_{t}\ e_{f}),\hat{\rho},\hat{\sigma},\hat{\sigma}_{\hat{\kappa}},\hat{a}_{\hat{\kappa}},\hat{t}\rangle \\ \sim E\langle e_{c},\hat{\rho},\hat{\sigma},\hat{\sigma}'_{k},\hat{a}'_{k},\hat{t}'\rangle \\ \text{where } \hat{a}'_{k} \triangleq \widehat{alloc}(\hat{\varsigma},0,\hat{\kappa}) \\ \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\kappa} \triangleq \widehat{ifk}(e_{t},e_{f},\hat{\rho},\hat{a}_{\hat{\kappa}}) \\ \hat{\sigma}'_{k} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{k} \mapsto \hat{\kappa}] \\ E\langle(\mathbf{let}\ ()\ e),\hat{\rho},\hat{\sigma},\hat{\sigma}_{k},\hat{a}_{k},\hat{t}\rangle \\ \sim E\langle e,\hat{\rho},\hat{\sigma},\hat{\sigma}'_{k},\hat{a}'_{k},\hat{t}'\rangle \\ \text{where } bids = ([x_{0}\ e_{0})[x_{s}\ e_{s}] \ldots) \\ \hat{a}'_{k} \triangleq \widehat{alloc}(\hat{\varsigma},0,\hat{\kappa}) \\ \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\kappa} \triangleq \mathbf{letk}(x_{0}\ ::x_{s},[],e_{s},e_{b},\hat{\rho},\hat{a}_{\hat{\kappa}}) \\ \hat{\sigma}'_{k} \triangleq \widehat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{k} \mapsto \hat{\kappa}] \\ E\langle(\mathbf{call}/\mathbf{cc}\ e),\hat{\rho},\hat{\sigma},\hat{\sigma}_{k},\hat{a}_{k},\hat{t}\rangle \\ \sim E\langle e,\hat{\rho},\hat{\sigma},\hat{\sigma}'_{k},\hat{a}'_{k},\hat{t}'\rangle \\ \text{where } \hat{a}'_{k} \triangleq \widehat{alloc}(\hat{\varsigma},0,\hat{\kappa}) \\ \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\kappa} \triangleq \mathbf{callcck}(\hat{\varsigma},\hat{\alpha},\hat{\kappa}) \\ \hat{\sigma}'_{k} \triangleq \widehat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{k} \mapsto \hat{\kappa}] \\ E\langle(\mathbf{call}/\mathbf{cc}\ e),\hat{\rho},\hat{\sigma},\hat{\sigma}_{k},\hat{a}_{k},\hat{t}\rangle \\ \sim E\langle e,\hat{\rho},\hat{\sigma},\hat{\sigma}'_{k},\hat{a}'_{k},\hat{t}'\rangle \\ \text{where } \hat{a}'_{k} \triangleq \widehat{alloc}(\hat{\varsigma},0,\hat{\kappa}) \\ \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k} \mapsto \hat{\tau}] \\ E\langle(\mathbf{call}/\mathbf{cc}\ e),\hat{\rho},\hat{\sigma},\hat{\sigma}_{k},\hat{a}_{k},\hat{t}\rangle \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{alloc}(\hat{\varsigma},0,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma},1,\hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k} \mapsto \hat{\tau}] \\ E\langle(\mathbf{call}/\mathbf{cc}\ e,\hat{\tau},\hat{\tau},\hat{\tau},\hat{\tau},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau},\hat{\tau}') \\ \hat{\tau}' \triangleq \widehat{\tau}_{k} \sqcup [\widehat{a}'_{k},\hat{\tau},\hat{\tau},\hat{\tau},$$

Apply Rules

Rules for when the control is a value $\hat{\varsigma} = A \langle \hat{v}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\hat{\kappa} \in \hat{\sigma}_{\hat{\kappa}}(\hat{a}_{\hat{\kappa}})$ Proceed by matching on $\hat{\kappa}$

$\mathbf{mt} \leadsto \hat{\varsigma}$

$$\text{ifk}(e_t,e_f,\hat{\rho}_{\hat{\kappa}},\hat{a}'_{\kappa}) \rightsquigarrow E\langle e_f,\hat{\rho}_{\hat{\kappa}},\hat{\sigma},\hat{\sigma}_{\hat{\kappa}},\hat{a}'_{\kappa},\hat{t} \rangle \\ \text{where } \hat{v} = \#\mathbf{f} \\ \text{where } \hat{v} = ((\lambda(x) e), \hat{\rho}_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}) \\ \text{where } \hat{v} \neq \#\mathbf{f} \\ \text{letk}(vars, \widehat{done}, [], e_b, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}'_{\hat{\kappa}}) \\ \sim E\langle e_b, \hat{\rho}'_{\hat{\kappa}}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ \text{where } \hat{a}_i \triangleq \widehat{alloc}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\rho}'_k \triangleq \widehat{alloc}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\rho}'_k \triangleq \widehat{alloc}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\sigma}' \triangleq \widehat{\sigma} \sqcup [\hat{a}_0 \mapsto \widehat{done}'_0 \dots \\ \widehat{a}_{n-1} \mapsto \widehat{done}'_{n-1}] \\ \text{letk}(vars, \widehat{done}, e_h :: e_t, e_b, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \\ \hat{\sigma}' \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\sigma}' \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\sigma}'_k \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{a}'_{\hat{\kappa}} \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{a}'_{\hat{\kappa}} \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{\tau}' \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{\tau}' \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{\tau}' \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{\tau}' \sqcup [\widehat{\tau}' \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{\tau} \sqcup [\widehat{\tau}' \mapsto \hat{\tau}'] \\ \hat{\tau}' \triangleq \widehat{\tau}' \sqcup [\widehat{\tau}' \mapsto \hat{\tau}$$

More Apply Rules

Rules for when the control is a value

$$\begin{aligned} & \mathbf{appappk}(\hat{v}_{f_1},...,a'_{k}) \\ & \sim E\langle e_b, \rho'_{\lambda},\hat{\sigma}',\hat{\sigma}_{\hat{\kappa}},\hat{a}'_{k};\hat{t}' \rangle \\ & \text{where } \hat{v}_{f} = ((\lambda \left(x_{s...}\right) e_b), \hat{\rho}_{\lambda}) \\ & \hat{a}_{i} \triangleq \widehat{alloc}(\hat{c},i,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},n,\hat{\kappa}) \\ & \hat{b}'_{\lambda} \triangleq \hat{\rho}_{\lambda}[x_{0} \mapsto \hat{a}_{0} \dots \\ & x_{n-1} \mapsto \hat{a}_{n-1}] \\ & \hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{a}_{0} \mapsto \hat{v}_{0} \dots \\ & \hat{a}_{n-1} \mapsto \hat{v}_{n-1}] \\ & \Rightarrow E\langle e_b, \rho'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{k}; \hat{t}' \rangle \\ & \text{where } \hat{v}_{f} = ((\lambda \left(x_{s...}\right) e_b), \hat{\rho}_{\lambda}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},n,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},n,\hat{\kappa}) \\ & \hat{\rho}'_{\lambda} \triangleq \hat{\rho}_{\lambda}[x_{0} \mapsto \hat{a}_{0} \dots \\ & x_{n-1} \mapsto \hat{a}_{n-1}] \\ & \hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{a}_{0} \mapsto \hat{v}_{0} \dots \\ & \hat{a}_{n-1} \mapsto \hat{v}_{n-1}] \\ & \Rightarrow E\langle e_b, \rho'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}; \hat{t}' \rangle \\ & \text{where } \hat{v}_{f} = ((\lambda \left(x_{s...}\right) e_b), \hat{\rho}_{\lambda}) \\ & \hat{a} \triangleq \widehat{alloc}(\hat{c},n,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},n,\hat{\kappa}) \\ & \hat{\rho}'_{\lambda} \triangleq \hat{\rho}_{\lambda}[x \mapsto \hat{a}] \\ & \hat{\tau}' \triangleq \hat{\sigma} \sqcup [\hat{a}_{k} \mapsto \hat{v}] \\ & \hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{a}'_{k} \mapsto \hat{v}] \\ & \Rightarrow E\langle e_b, \rho'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{k}, \hat{t}' \rangle \\ & \text{where } \hat{a}''_{\kappa} \triangleq \widehat{alloc}(\hat{c},0,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},1,\hat{\kappa}) \\ & \hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{a}'_{\kappa} \mapsto \hat{\kappa}'] \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \text{where } \hat{a}''_{\kappa} \triangleq \widehat{alloc}(\hat{c},0,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},1,\hat{\kappa}) \\ & \hat{\sigma}'_{\kappa} \triangleq \hat{\sigma}_{\kappa} \sqcup [\hat{a}'_{\kappa} \mapsto \hat{v}'] \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \text{where } \hat{a}''_{\kappa} \triangleq \widehat{alloc}(\hat{c},0,\hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{c},1,\hat{\kappa}) \\ & \hat{\sigma}'_{\kappa} \triangleq \hat{\sigma}_{\kappa} \sqcup [\hat{a}'_{\kappa} \mapsto \hat{\kappa}'] \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho}_{\lambda}, \hat{\sigma}, \hat{a}'_{\kappa}, \hat{t}' \rangle \\ & \Rightarrow E\langle e_b, \hat{\rho$$