

1 Concrete Semantics of Scheme CESKt*

Syntax Domains:

$e \in \text{Exp} ::= \mathfrak{x}$

| (if e e e)
 | (let ($[x$ $e]$...) e)
 | (call/cc e)
 | (set! x e)
 | (prim op e ...)
 | (apply-prim op e)
 | (apply e e)
 | (e e ...)

$\mathfrak{x} \in \text{AExp} ::= x \mid lam \mid \mathbb{Z} \mid \#t \mid \#f$
 | (quote e)

$lam \in \text{Lam} ::= (\lambda (x \dots) e) \mid (\lambda x e)$

$x \in \text{Var}$ A set of identifiers

$op \in \text{Prim}$ A set of primitives

Atomic Evaluation:

$\mathcal{A} : \Sigma_E \times \sigma \rightarrow \text{Val}$

$\mathcal{A}(\langle lam, \rho, -, - \rangle, -) \triangleq (lam, \rho)$

$\mathcal{A}(\langle \langle \text{quote } e \rangle, -, -, - \rangle, -) \triangleq \text{quote}(e)$

$\mathcal{A}(\langle x, \rho, -, - \rangle, \sigma) \triangleq \sigma(\rho(x))$

$\mathcal{A}(\langle \mathfrak{x}, -, -, - \rangle, -) \triangleq \mathfrak{x}$

Tick/Alloc:

$tick : \text{Eval} \times \mathbb{N} \rightarrow \text{Time}$

$tick(E \langle -, -, -, t \rangle, n) \triangleq (t + n)$

$alloc : \Sigma \times \mathbb{N} \triangleq \text{Addr}$

$alloc(E \langle -, -, -, t \rangle, n) \triangleq (t + n)$

Injection:

$\mathcal{I} : \text{Exp} \rightarrow \Sigma$

$\mathcal{I}(e) \triangleq (e, \emptyset, \emptyset, \{0 : \text{mt}\}, 0, 1)$

Transition:

Collecting Semantics:

Semantic Domains:

$\varsigma \in \Sigma \triangleq E \langle \text{Eval} \rangle + A \langle \text{Apply} \rangle$

$\text{Eval} \triangleq \text{Exp} \times \text{Env} \times \text{Store}$
 $\times K\text{Store} \times \text{Addr} \times \text{Time}$

$\text{Apply} \triangleq \text{Val} \times \text{Store}$
 $\times K\text{Store} \times \text{Addr} \times \text{Time}$

$\rho \in \text{Env} \triangleq \text{Var} \rightarrow \text{Addr}$

$\sigma \in \text{Store} \triangleq \text{Addr} \rightarrow \text{Val}$

$\sigma \in \text{Store} \triangleq \text{Addr} \rightarrow \text{Kont}$

$v \in \text{Val} \triangleq \text{Clo} + \text{Kont} + \mathbb{Z}$
 $+ \{\#t, \#f, \text{Null}, \text{Void}\}$
 $+ \{\text{quote}(e), \text{cons}(v, v)\}$

$clo \in \text{Clo} \triangleq \text{Lam} \times \text{Env}$

$a, b, c \in \text{Addr} \triangleq \mathbb{N}$

$t, u \in \text{Time} \triangleq \mathbb{N}$

$\kappa \in \text{Kont} ::= \text{mt}$

| ifk(e, e, ρ, a)

| calck(a)

| setk(x, ρ, a)

| appappk($val?, e, \rho, a$)

| appk($done, todo, \rho, a$)

| appprimk(op, ρ, a)

| primk($op, done, todo,$
 ρ, a)

| letk($vars, done, todo,$
 e, ρ, a)

$done \triangleq \text{Val}^*$

$todo \triangleq \text{Exp}^*$

$vars \triangleq \text{Var}^*$

Eval Rules

Rules for when the control is an expression

$$E\langle \mathfrak{x}, \rho, \sigma, \sigma_\kappa, a, t \rangle \rightsquigarrow A\langle v, \sigma, \sigma_\kappa, a, t \rangle$$

where $v \triangleq \mathcal{A}(\varsigma)$

$E\langle (\text{if } e_c \ e_t \ e_f), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e_c, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{ifk}(e_t, e_f, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (\text{let } () \ e), \rho, \sigma, \sigma_\kappa, a, t \rangle$ $\rightsquigarrow E\langle e, \rho, \sigma, \sigma_\kappa, a, t \rangle$ $E\langle (\text{let } ([x_0 \ e_0] \ [x_s \ e_s] \ \dots) \ e_b),$ $\rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle \rightsquigarrow E\langle e_0, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{letk}(x_0 :: x_s, [], e_s, e_b, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (\text{call/cc } e), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{callcck}(a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (\text{set! } x \ e), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{setk}(x, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$	$E\langle (\text{prim } op), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow A\langle v, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ <p style="text-align: center;">where $v = op$ applied to 0 arguments</p> $E\langle (\text{prim } op \ e_0 \ e_s \ \dots), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e_0, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{primk}(op, [], e_s, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (\text{apply-prim } op \ e), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e, \rho, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{appprimk}(op, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (\text{apply } e_f \ e_x), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e_f, \rho, \sigma, \sigma_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{appappk}(\emptyset, e_x, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$ $E\langle (e_f \ e_s \ \dots), \rho, \sigma, \sigma_\kappa, a_\kappa, t \rangle$ $\rightsquigarrow E\langle e_f, \rho, \sigma, \sigma_\kappa, a'_\kappa, t' \rangle$ <p style="text-align: center;">where $a'_\kappa \triangleq \text{alloc}(\varsigma, 0)$</p> $t' \triangleq \text{tick}(\varsigma, 1)$ $\kappa \triangleq \mathbf{appk}([], e_s, \rho, a_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa]$
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Apply Rules

Rules for when the control is a value

$$\begin{aligned}\varsigma &= A\langle v, \sigma, \sigma_\kappa, a_\kappa, t \rangle \\ \kappa &\triangleq \sigma_\kappa(a_\kappa)\end{aligned}$$

Proceed by matching on κ

$$\mathbf{mt} \rightsquigarrow \varsigma$$

$\begin{aligned}\mathbf{ifk}(e_t, e_f, \rho_\kappa, a'_\kappa) &\rightsquigarrow E\langle e_f, \rho_\kappa, \sigma, \sigma_\kappa, a'_\kappa, t \rangle \\ &\text{where } v = \#f \\ \mathbf{ifk}(e_t, e_f, \rho_\kappa, a'_\kappa) &\rightsquigarrow E\langle e_t, \rho_\kappa, \sigma, \sigma_\kappa, a'_\kappa, t \rangle \\ &\text{where } v \neq \#f \\ \mathbf{letk}(vars, done, [], e_b, \rho_\kappa, a'_\kappa) &\rightsquigarrow E\langle e_b, \rho'_\kappa, \sigma', \sigma_\kappa, a'_\kappa, t' \rangle \\ \text{where } a_i &\triangleq alloc(\varsigma, i) \\ t' &\triangleq tick(\varsigma, n) \\ \rho'_\kappa &\triangleq \rho_\kappa[vars_0 \mapsto a_0 \dots \\ &\quad vars_{n-1} \mapsto a_{n-1}] \\ done' &\triangleq done \# [v] \\ \sigma' &\triangleq \sigma[a_0 \mapsto done'_0 \dots \\ &\quad a_{n-1} \mapsto done'_{n-1}] \\ \mathbf{letk}(vars, done, e_h :: e_t, e_b, \rho_\kappa, a'_\kappa) &\rightsquigarrow E\langle e_h, \rho_\kappa, \sigma, \sigma'_\kappa, a''_\kappa, t' \rangle \\ \text{where } a''_\kappa &\triangleq alloc(\varsigma, 0) \\ t' &\triangleq tick(\varsigma, 1) \\ \kappa' &\triangleq \mathbf{letk}(vars, done \# [v], \\ &\quad e_t, e_b, \rho_\kappa, a'_\kappa) \\ \sigma'_\kappa &\triangleq \sigma_\kappa[a''_\kappa \mapsto \kappa'] \\ \mathbf{setk}(x, \rho_\kappa, a'_\kappa) &\rightsquigarrow A\langle Void, \sigma', \sigma_\kappa, a'_\kappa, t \rangle \\ \text{where } \sigma' &\triangleq \sigma[\rho_\kappa(x) \mapsto v]\end{aligned}$	$\begin{aligned}\mathbf{callock}(a'_\kappa) &\rightsquigarrow E\langle e, \rho'_\lambda, \sigma, \sigma_\kappa, a'_\kappa, t \rangle \\ \text{where } v &= ((\lambda (x) e), \rho_\lambda) \\ a &\triangleq tick(\varsigma, 0) \\ t' &\triangleq alloc(\varsigma, 1) \\ \kappa' &\triangleq \sigma_\kappa(a'_\kappa) \\ \rho'_\lambda &\triangleq \rho_\lambda[x \mapsto a] \\ \sigma' &\triangleq \sigma[a \mapsto \kappa'] \\ \mathbf{callock}(-) &\rightsquigarrow A\langle \kappa, \sigma, \sigma'_\kappa, a'_\kappa, t' \rangle \\ \text{where } v &= \kappa' \\ a'_\kappa &\triangleq alloc(\varsigma, 0) \\ t' &\triangleq tick(\varsigma, 1) \\ \sigma'_\kappa &\triangleq \sigma_\kappa[a'_\kappa \mapsto \kappa'] \\ \mathbf{appprimk}(op, \rho_\kappa, a'_\kappa) &\rightsquigarrow A\langle v', \sigma, \sigma_\kappa, a'_\kappa, t \rangle \\ \text{where } v' &\triangleq op \text{ applied to } v \\ \mathbf{primk}(op, done, [], \rho_\kappa, a'_\kappa) &\rightsquigarrow A\langle v', \sigma, \sigma_\kappa, a'_\kappa, t \rangle \\ \text{where } v' &\triangleq op \text{ applied to } (done \# [v]) \\ \mathbf{primk}(op, done, e_h :: e_t, \rho_\kappa, a'_\kappa) &\rightsquigarrow E\langle e_h, \rho_\kappa, \sigma, \sigma'_\kappa, a''_\kappa, t' \rangle \\ \text{where } a''_\kappa &\triangleq alloc(\varsigma, 0) \\ t' &\triangleq tick(\varsigma, 1) \\ \kappa' &\triangleq \mathbf{primk}(op, done \# [v], \\ &\quad e_t, \rho_\kappa, a'_\kappa) \\ \sigma'_\kappa &\triangleq \sigma_\kappa[a''_\kappa \mapsto \kappa']\end{aligned}$
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More Apply Rules

Rules for when the control is a value

$\mathbf{appappk}(v_f, -, -, a'_\kappa)$ $\rightsquigarrow E\langle e_b, \rho'_\lambda, \sigma', \sigma_\kappa, a'_\kappa, t' \rangle$ <p>where $v_f = ((\lambda (x_s \dots) e_b), \rho_\lambda)$</p> $a_i \triangleq \mathit{alloc}(\varsigma, i)$ $t' \triangleq \mathit{tick}(\varsigma, n)$ $\rho'_\lambda \triangleq \rho_\lambda[x_0 \mapsto a_0 \dots$ $x_{n-1} \mapsto a_{n-1}]$ $\sigma' \triangleq \sigma[a_0 \mapsto v_0 \dots$ $a_{n-1} \mapsto v_{n-1}]$ $\mathbf{appappk}(v_f, -, -, a'_\kappa)$ $\rightsquigarrow E\langle e_b, \rho'_\lambda, \sigma', \sigma_\kappa, a'_\kappa, t' \rangle$ <p>where $v_f = ((\lambda x e_b), \rho_\lambda)$</p> $a \triangleq \mathit{alloc}(\varsigma, 0)$ $t' \triangleq \mathit{tick}(\varsigma, 1)$ $\rho'_\lambda \triangleq \rho_\lambda[x \mapsto a]$ $\sigma' \triangleq \sigma[a \mapsto v]$ $\mathbf{appappk}(\emptyset, e, \rho_\kappa, a'_\kappa)$ $\rightsquigarrow E\langle e, \rho_\kappa, \sigma, \sigma'_\kappa, a''_\kappa, t' \rangle$ <p>where $a''_\kappa \triangleq \mathit{alloc}(\varsigma, 0)$</p> $t' \triangleq \mathit{tick}(\varsigma, 1)$ $\kappa' \triangleq \mathbf{appappk}(v, e, \rho_\kappa, a'_\kappa)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a''_\kappa \mapsto \kappa']$	$\mathbf{appk}(\mathit{done}, [], -, a'_\kappa)$ $\rightsquigarrow E\langle e_b, \rho'_\lambda, \sigma', \sigma_\kappa, a'_\kappa, t' \rangle$ <p>where $\mathit{done} = ((\lambda (x_s \dots) e_b), \rho_\lambda) :: v_s$</p> $a_i \triangleq \mathit{alloc}(\varsigma, i)$ $t' \triangleq \mathit{tick}(\varsigma, n)$ $v'_s \triangleq v_s \# [v]$ $\rho'_\lambda \triangleq \rho_\lambda[x_0 \mapsto a_0 \dots$ $x_{n-1} \mapsto a_{n-1}]$ $\sigma' \triangleq \sigma[a_0 \mapsto v'_0 \dots$ $a_{n-1} \mapsto v'_{n-1}]$ $\mathbf{appk}(\mathit{done}, [], -, a'_\kappa)$ $\rightsquigarrow E\langle e_b, \rho'_\lambda, \sigma', \sigma_\kappa, a'_\kappa, t' \rangle$ <p>where $\mathit{done} = ((\lambda x e_b), \rho_\lambda) :: v_s$</p> $a \triangleq \mathit{alloc}(\varsigma, 0)$ $t' \triangleq \mathit{tick}(\varsigma, 1)$ $v'_s \triangleq (v_s \# [v])$ $\rho'_\lambda \triangleq \rho_\lambda[x \mapsto a]$ $\sigma' \triangleq \sigma[a \mapsto v'_s]$ $\mathbf{appk}([\kappa_\lambda], [], \rho_\kappa, a'_\kappa) \rightsquigarrow A\langle v, \sigma, \sigma'_\kappa, a''_\kappa, t' \rangle$ <p>where $a''_\kappa \triangleq \mathit{alloc}(\varsigma, 0)$</p> $t' \triangleq \mathit{tick}(\varsigma, 1)$ $\sigma'_\kappa \triangleq \sigma_\kappa[a''_\kappa \mapsto \kappa_\lambda]$ $\mathbf{appk}(\mathit{done}, e_h :: e_t, \rho_\kappa, a'_\kappa)$ $\rightsquigarrow E\langle e_h, \rho_\kappa, \sigma, \sigma'_\kappa, a''_\kappa, t' \rangle$ <p>where $a''_\kappa \triangleq \mathit{alloc}(\varsigma, 0)$</p> $t' \triangleq \mathit{tick}(\varsigma, 1)$ $\kappa' \triangleq \mathbf{appk}(\mathit{done} \# [v], e_t, \rho_\kappa, a'_\kappa)$ $\sigma'_\kappa \triangleq \sigma[a''_\kappa \mapsto \kappa']$
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2 Abstract Semantics of Scheme CESKt*

Abstract Semantic Domains:

$$\begin{aligned}
\hat{\varsigma} \in \widehat{\Sigma} &\triangleq E\langle Eval \rangle + E\langle Apply \rangle & \hat{\kappa} \in \widehat{Kont} &::= \mathbf{mt} \\
Eval &\triangleq \mathbf{Exp} \times \widehat{Env} \times \widehat{Store} & & | \mathbf{ifk}(e, e, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
&\times \widehat{KStore} \times \widehat{Addr} \times \widehat{Time} & & | \mathbf{callcck}(\hat{a}_{\hat{\kappa}}) \\
Apply &\triangleq \widehat{Val} \times \widehat{Store} & & | \mathbf{setk}(x, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
&\times \widehat{KStore} \times \widehat{Addr} \times \widehat{Time} & & | \mathbf{appappk}(\widehat{val?}, e, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
\hat{\rho} \in \widehat{Env} &\triangleq \mathbf{Var} \rightarrow \widehat{Addr} & & | \mathbf{appk}(\widehat{done}, \widehat{todo}, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
\hat{\sigma} \in \widehat{Store} &\triangleq \widehat{Addr} \rightarrow \mathcal{P}(\widehat{Val}) & & | \mathbf{appprimk}(op, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
\hat{\sigma}_{\kappa} \in \widehat{KStore} &\triangleq \widehat{Addr} \rightarrow \mathcal{P}(\widehat{Kont}) & & | \mathbf{primk}(op, \widehat{done}, \widehat{todo}, \\
&&&& \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
\hat{v} \in \widehat{Val} &\triangleq \widehat{Clo} + \widehat{Kont} + \mathbb{Z} & & | \mathbf{letk}(\widehat{vars}, \widehat{done}, \widehat{todo}, \\
&+ \{\mathbf{\#t}, \mathbf{\#f}, \mathbf{Null}, \mathbf{Void}\} & & e, \hat{\rho}, \hat{a}_{\hat{\kappa}}) \\
&+ \{\mathbf{quote}(e), \mathbf{cons}(v, v)\} \\
\widehat{clo} \in \widehat{Clo} &\triangleq \mathbf{Lam} \times \widehat{Env} \\
\hat{a}, \hat{a}_{\hat{\kappa}} \in \widehat{Addr} &\triangleq \mathbb{N} \\
\hat{t} \in \widehat{Time} &\triangleq \mathbb{N} \\
\widehat{done} &\triangleq \widehat{Val}^*
\end{aligned}$$

Abstract Atomic Evaluation:

$$\begin{aligned}
\mathcal{A} &: Eval \rightarrow \widehat{Val} \\
\mathcal{A}(E\langle lam, \hat{\rho}, \dots \rangle) &\triangleq \{(lam, \hat{\rho})\} \\
\mathcal{A}(\mathbf{\ae}) &\triangleq \{\mathbf{\ae}\}
\end{aligned}$$

Tick/Alloc:

$$\begin{aligned}
\widehat{tick} &: \widehat{\Sigma} \times \widehat{Kont} \rightarrow \widehat{Time} \\
\widehat{tick}(\hat{\varsigma}, \hat{\kappa}) &\triangleq 0 \\
\widehat{alloc} &: \widehat{\Sigma} \times \widehat{Kont} \rightarrow \widehat{Addr} \\
\widehat{alloc}(\hat{\varsigma}, \hat{\kappa}) &\triangleq 0
\end{aligned}$$

Abstract Eval Rules
Rules for when the control is an expression

$$E\langle \mathfrak{a}, \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle \rightsquigarrow A\langle \hat{v}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$$

where $\hat{v} \triangleq \hat{\mathcal{A}}(\hat{\zeta}, \hat{\sigma})$

$E\langle (\text{if } e_c \ e_t \ e_f), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e_c, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{ifk}(e_t, e_f, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (\text{let } () \ e), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e, \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $E\langle (\text{let } ([x_0 \ e_0] \ [x_s \ e_s] \ \dots) \ e_b),$ $\hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle \rightsquigarrow E\langle e_0, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{letk}(x_0 :: x_s, [], e_s, e_b, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (\text{call/cc } e), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{callcck}(\hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (\text{set! } x \ e), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{setk}(x, a)$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$	$E\langle (\text{prim op}), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow A\langle \hat{v}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ <p>where $\hat{v} = \text{op}$ applied to 0 arguments</p> $E\langle (\text{prim op } e_0 \ e_s \ \dots), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e_0, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{primk}(\text{op}, [], e_s, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (\text{apply-prim op } e), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{appprimk}(\text{op}, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (\text{apply } e_f \ e_x), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e_f, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{appappk}(\emptyset, e_x, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$ $E\langle (e_f \ e_s \ \dots), \hat{\rho}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$ $\rightsquigarrow E\langle e_f, \hat{\rho}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$ <p>where $\hat{a}'_{\hat{\kappa}} \triangleq \widehat{\text{alloc}}(\hat{\zeta}, 0, \hat{\kappa})$</p> $\hat{t}' \triangleq \widehat{\text{tick}}(\hat{\zeta}, 1, \hat{\kappa})$ $\hat{\kappa} \triangleq \mathbf{appk}([], e_s, \hat{\rho}, \hat{a}_{\hat{\kappa}})$ $\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}]$
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Apply Rules

Rules for when the control is a value

$$\hat{\varsigma} = A\langle \hat{v}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}_{\hat{\kappa}}, \hat{t} \rangle$$

$$\kappa \in \hat{\sigma}_{\hat{\kappa}}(\hat{a}_{\hat{\kappa}})$$

Proceed by matching on κ

$$\mathbf{mt} \rightsquigarrow \hat{\varsigma}$$

$$\mathbf{ifk}(e_t, e_f, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow E\langle e_f, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

$$\text{where } \hat{v} = \#f$$

$$\mathbf{ifk}(e_t, e_f, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow E\langle e_t, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

$$\text{where } \hat{v} \neq \#f$$

$$\mathbf{letk}(\text{vars}, \widehat{done}, [], e_b, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow E\langle e_b, \hat{\rho}'_{\hat{\kappa}}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$$

$$\text{where } \hat{a}_i \triangleq \widehat{alloc}(\hat{\varsigma}, i, \hat{\kappa})$$

$$\hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, n, \hat{\kappa})$$

$$\hat{\rho}'_{\hat{\kappa}} \triangleq \hat{\rho}_{\hat{\kappa}}[vars_0 \mapsto \hat{a}_0 \dots \\ vars_{n-1} \mapsto \hat{a}_{n-1}]$$

$$\widehat{done}' \triangleq \widehat{done} + [\hat{v}]$$

$$\hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{a}_0 \mapsto \widehat{done}'_0 \dots \\ \hat{a}_{n-1} \mapsto \widehat{done}'_{n-1}]$$

$$\mathbf{letk}(\text{vars}, \widehat{done}, e_h :: e_t, e_b, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow E\langle e_h, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}''_{\hat{\kappa}}, \hat{t}' \rangle$$

$$\text{where } \hat{a}''_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa})$$

$$\hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa})$$

$$\hat{\kappa}' \triangleq \mathbf{letk}(\text{vars}, \widehat{done}, +[\hat{v}], \\ e_t, e_b, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}})$$

$$\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}''_{\hat{\kappa}} \mapsto \hat{\kappa}']$$

$$\mathbf{setk}(x, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow A\langle \text{Void}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

$$\text{where } \hat{\sigma}' \triangleq \hat{\sigma} \sqcup [\hat{\rho}_{\hat{\kappa}}(x) \mapsto \hat{v}]$$

$$\mathbf{calccck}(\hat{a}'_{\hat{\kappa}}) \rightsquigarrow E\langle e, \hat{\rho}'_{\hat{\lambda}}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

THIS IS BROKEN, see concrete version

$$\text{where } \hat{v} = ((\lambda (x) e), \hat{\rho}_{\hat{\lambda}})$$

$$\hat{\rho}'_{\hat{\lambda}} \triangleq \hat{\rho}_{\hat{\lambda}}[x \mapsto \hat{a}'_{\hat{\kappa}}]$$

$$\mathbf{calccck}(-) \rightsquigarrow A\langle \hat{\kappa}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle$$

$$\text{where } v = \kappa'$$

$$\hat{a}'_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa})$$

$$\hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa})$$

$$\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \kappa']$$

$$\mathbf{appprimk}(op, -, \hat{a}'_{\hat{\kappa}}) \rightsquigarrow A\langle \hat{v}', \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

$$\text{where } \hat{v}' \triangleq op \text{ applied to } \hat{v}$$

$$\mathbf{primk}(op, \widehat{done}, [], -, \hat{a}'_{\hat{\kappa}})$$

$$\rightsquigarrow A\langle \hat{v}', \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t} \rangle$$

$$\text{where } \hat{v}' \triangleq op \text{ applied to } (\widehat{done} + [\hat{v}])$$

$$\mathbf{primk}(op, \widehat{done}, e_h :: e_t, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}})$$

$$\rightsquigarrow E\langle e_h, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}''_{\hat{\kappa}}, \hat{t}' \rangle$$

$$\text{where } \hat{a}''_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa})$$

$$\hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa})$$

$$\hat{\kappa}' \triangleq \mathbf{primk}(op, \widehat{done} + [\hat{v}], \\ e_t, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}})$$

$$\hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}''_{\hat{\kappa}} \mapsto \hat{\kappa}']$$

More Apply Rules

Rules for when the control is a value

$$\begin{aligned} & \mathbf{appappk}(\hat{v}_f, -, -, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e_b, \hat{\rho}'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \hat{v}_f = ((\lambda (x_s \dots) e_b), \hat{\rho}_{\lambda}) \end{aligned}$$

$$\begin{aligned} \hat{a}_i & \triangleq \widehat{alloc}(\hat{\varsigma}, i, \hat{\kappa}) \\ \hat{t}' & \triangleq \widehat{tick}(\hat{\varsigma}, n, \hat{\kappa}) \\ \hat{\rho}'_{\lambda} & \triangleq \hat{\rho}_{\lambda}[x_0 \mapsto \hat{a}_0 \dots \\ & \quad x_{n-1} \mapsto \hat{a}_{n-1}] \\ \hat{\sigma}' & \triangleq \hat{\sigma} \sqcup [\hat{a}_0 \mapsto \hat{v}_0 \dots \\ & \quad \hat{a}_{n-1} \mapsto \hat{v}_{n-1}] \end{aligned}$$

$$\begin{aligned} & \mathbf{appappk}(\hat{v}_f, -, -, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e_b, \hat{\rho}'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \hat{v}_f = ((\lambda x e_b), \hat{\rho}_{\lambda}) \end{aligned}$$

$$\begin{aligned} \hat{a} & \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ \hat{t}' & \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\rho}'_{\lambda} & \triangleq \hat{\rho}_{\lambda}[x \mapsto \hat{a}] \\ \hat{\sigma}' & \triangleq \hat{\sigma} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{v}] \end{aligned}$$

$$\begin{aligned} & \mathbf{appappk}(\emptyset, e, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \hat{a}''_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ & \hat{\kappa}' \triangleq \mathbf{appappk}(\hat{v}, e, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \\ & \hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}''_{\hat{\kappa}} \mapsto \hat{\kappa}'] \end{aligned}$$

$$\begin{aligned} & \mathbf{appk}(\widehat{done}, [], -, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e_b, \hat{\rho}'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \end{aligned}$$

$$\text{where } \widehat{done} = ((\lambda (x_s \dots) e_b), \hat{\rho}_{\lambda}) :: \hat{v}_s$$

$$\begin{aligned} \hat{a}_i & \triangleq \widehat{alloc}(\hat{\varsigma}, i, \hat{\kappa}) \\ \hat{t}' & \triangleq \widehat{tick}(\hat{\varsigma}, n, \hat{\kappa}) \\ \hat{\rho}'_{\lambda} & \triangleq \hat{\rho}_{\lambda}[x_0 \mapsto \hat{a}_0 \dots \\ & \quad x_{n-1} \mapsto \hat{a}_{n-1}] \\ \hat{v}'_s & \triangleq \hat{v}_s \mathbin{++} [\hat{v}] \\ \hat{\sigma}' & \triangleq \hat{\sigma} \sqcup [\hat{a}_0 \mapsto \hat{v}'_0 \dots \\ & \quad \hat{a}_{n-1} \mapsto \hat{v}'_{n-1}] \end{aligned}$$

$$\begin{aligned} & \mathbf{appk}(\widehat{done}, [], -, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e_b, \hat{\rho}'_{\lambda}, \hat{\sigma}', \hat{\sigma}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \widehat{done} = ((\lambda x e_b), \hat{\rho}_{\lambda}) :: \hat{v}_s \end{aligned}$$

$$\begin{aligned} \hat{a} & \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ \hat{t}' & \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ \hat{\rho}'_{\lambda} & \triangleq \hat{\rho}_{\lambda}[x \mapsto \hat{a}] \\ \hat{v}'_s & \triangleq \hat{v}_s \mathbin{++} [\hat{v}] \\ \hat{\sigma}' & \triangleq \hat{\sigma} \sqcup [\hat{a} \mapsto \hat{v}'_s] \end{aligned}$$

$$\begin{aligned} & \mathbf{appk}([\hat{\kappa}_{\lambda}], [], \hat{\rho}_{\hat{\kappa}}, -) \rightsquigarrow A\langle \hat{v}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \hat{a}'_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ & \hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma} \sqcup [\hat{a}'_{\hat{\kappa}} \mapsto \hat{\kappa}_{\lambda}] \end{aligned}$$

$$\begin{aligned} & \mathbf{appk}(\widehat{done}, e_h :: e_t, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \\ & \rightsquigarrow E\langle e_h, \hat{\rho}_{\hat{\kappa}}, \hat{\sigma}, \hat{\sigma}'_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}, \hat{t}' \rangle \\ & \text{where } \hat{a}''_{\hat{\kappa}} \triangleq \widehat{alloc}(\hat{\varsigma}, 0, \hat{\kappa}) \\ & \hat{t}' \triangleq \widehat{tick}(\hat{\varsigma}, 1, \hat{\kappa}) \\ & \hat{\kappa}' \triangleq \mathbf{appk}(\widehat{done} \mathbin{++} [\hat{v}], e_t, \hat{\rho}_{\hat{\kappa}}, \hat{a}'_{\hat{\kappa}}) \\ & \hat{\sigma}'_{\hat{\kappa}} \triangleq \hat{\sigma}_{\hat{\kappa}} \sqcup [\hat{a}''_{\hat{\kappa}} \mapsto \hat{\kappa}'] \end{aligned}$$