

粉笔网架构

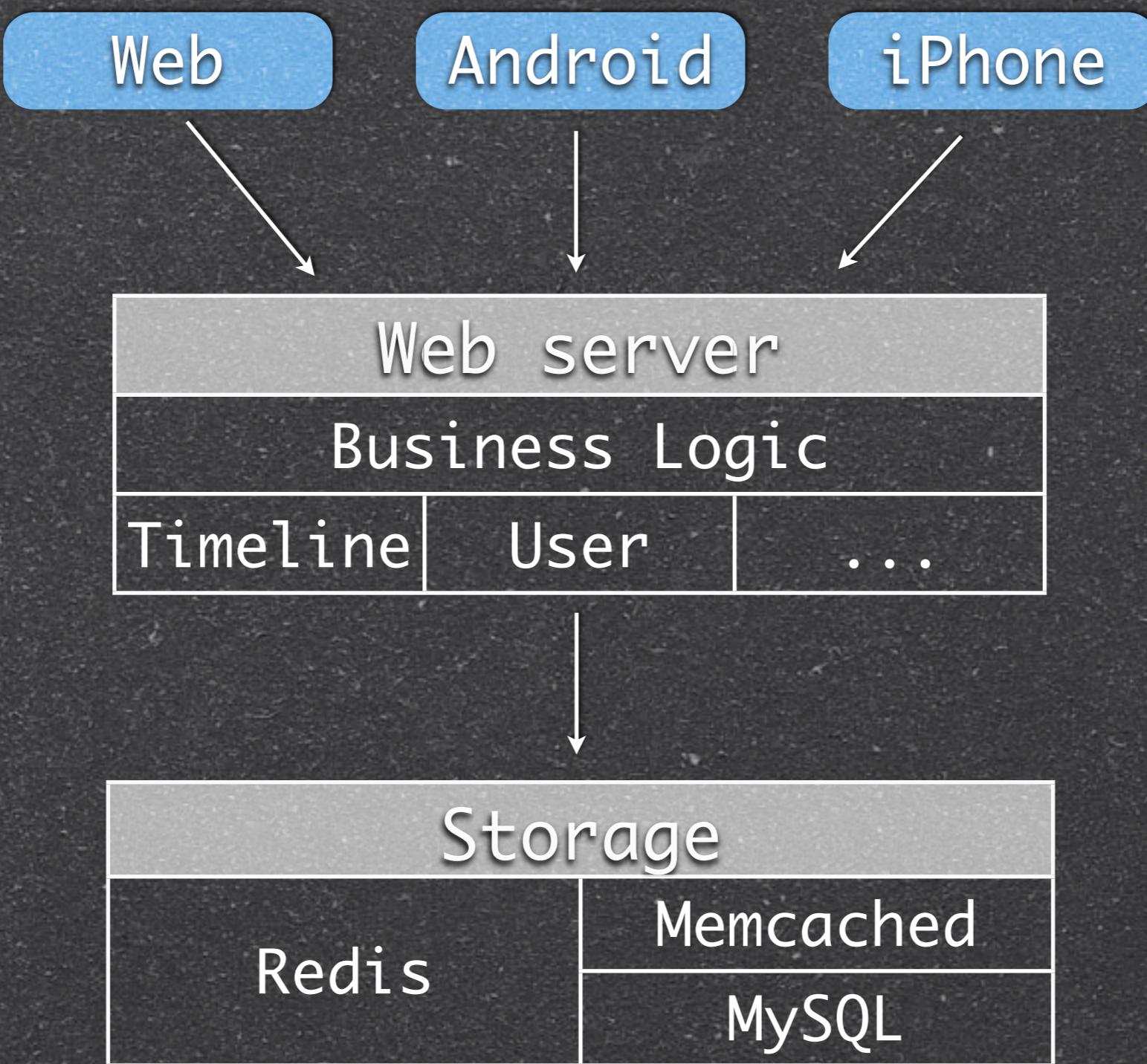
支持多终端的服务器实现

杨元祖@粉笔网

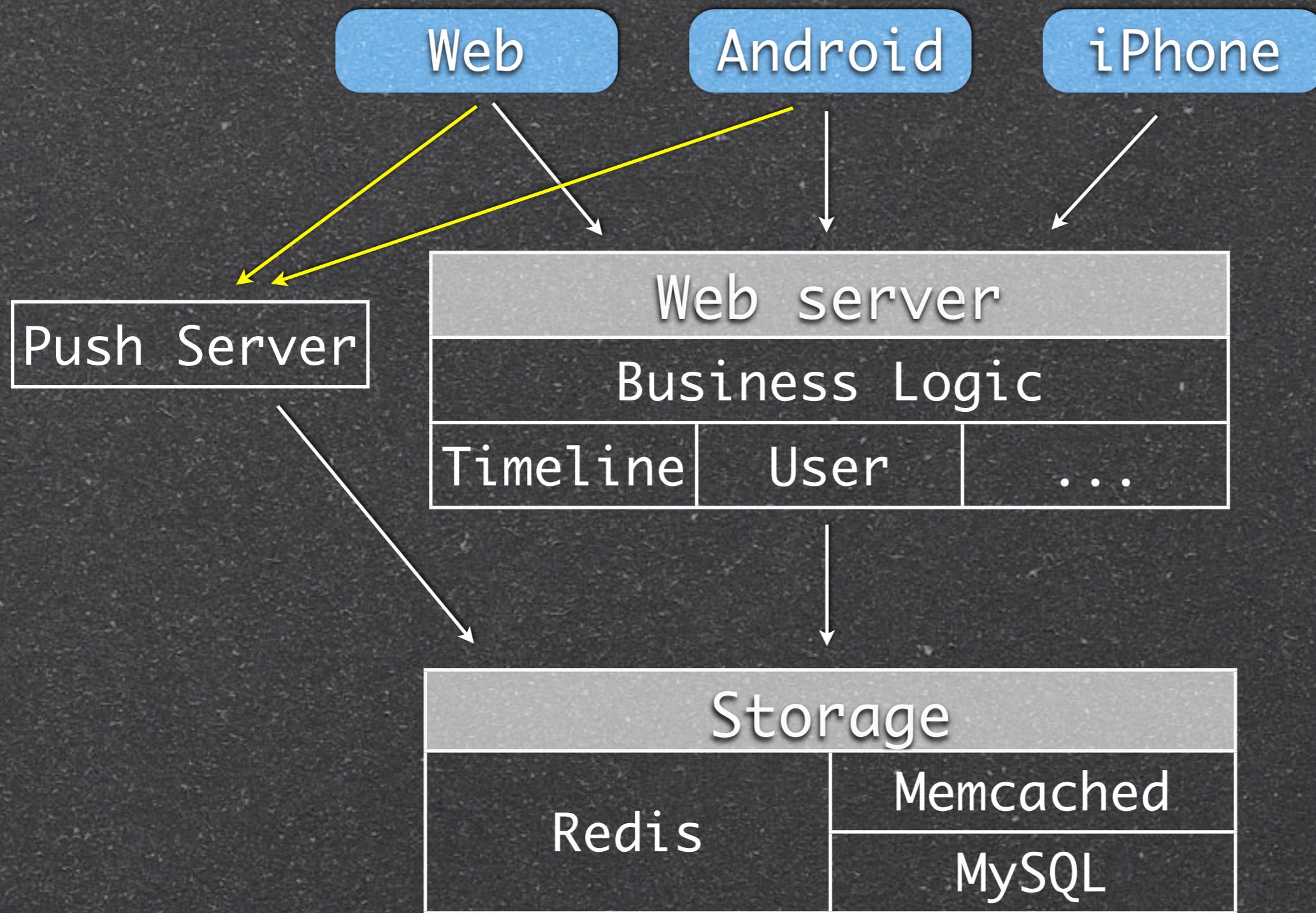
主要技术

- 语言：Java、Spring
 - 开发速度、熟练程度、性能
- 存储：MySQL、Redis
- 消息推送：Jetty Continuation
- RESTful

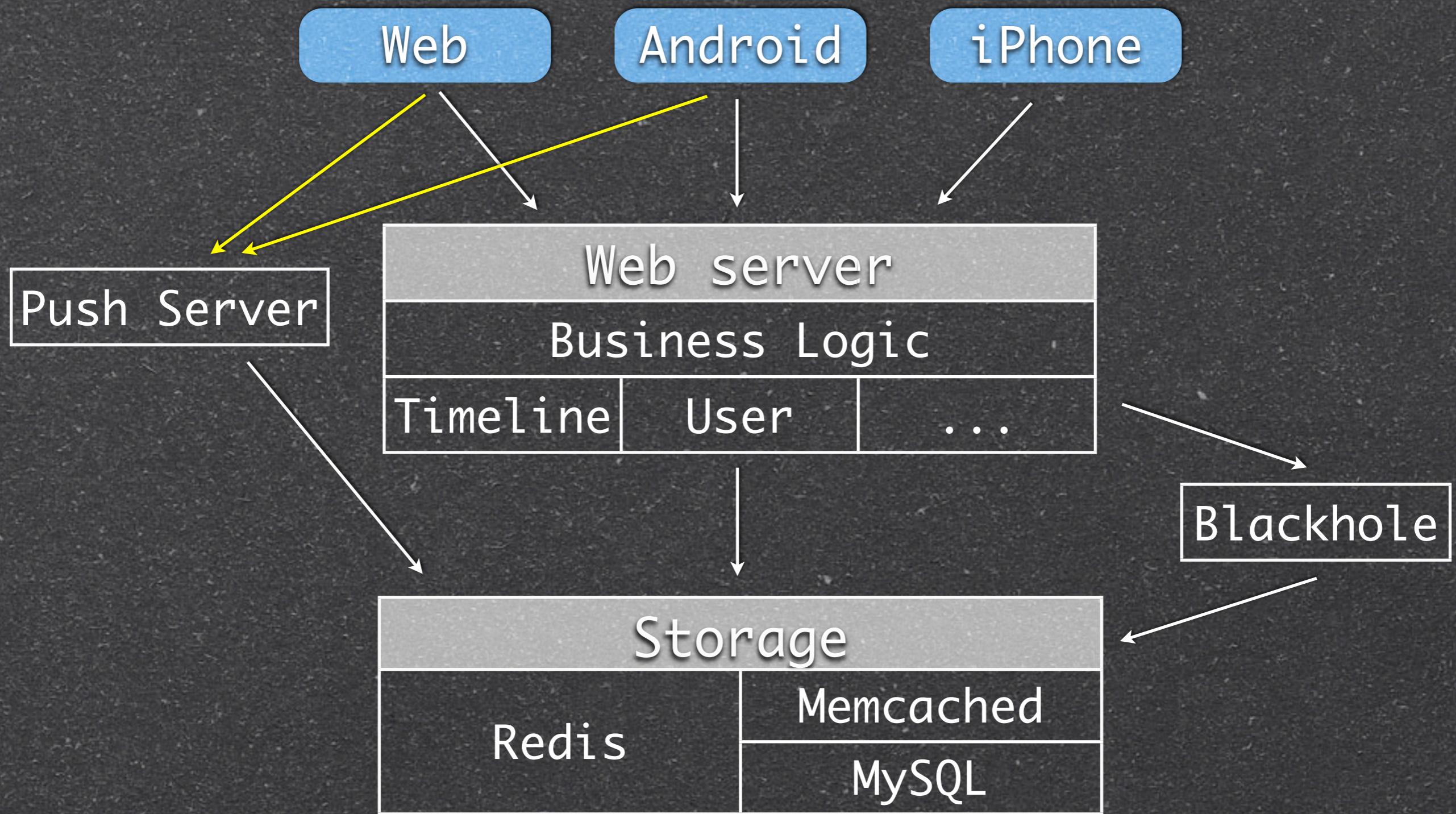
粉笔网体系结构



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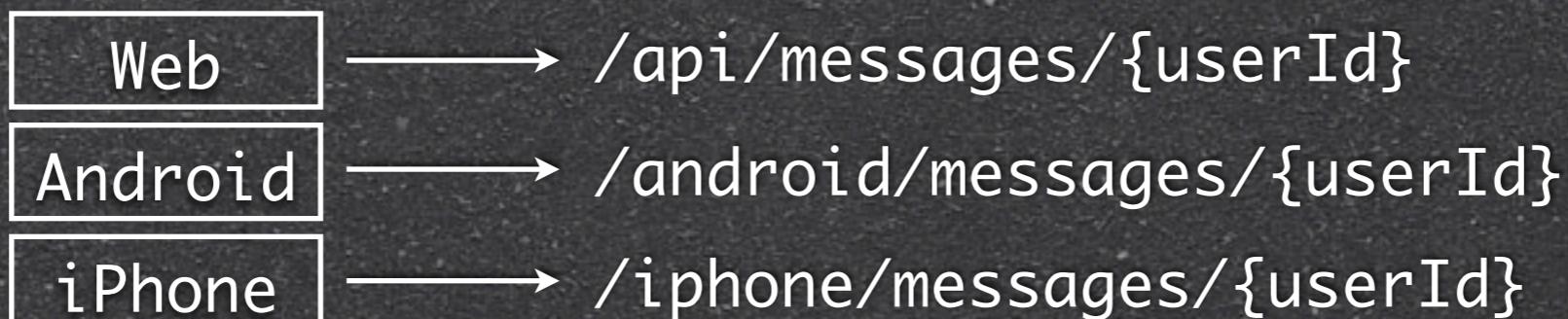


粉笔网体系结构



Web Server

- 接口定义--支持多个终端
 - 方案一：各个终端共用一个接口

```
graph LR; Web[Web] --> API["/api/messages/{userId}"]; Android[Android] --> API; iPhone[iPhone] --> API;
```
 - 方案二：每个终端有各自的接口

```
graph LR; Web[Web] --> API1["/api/messages/{userId}"]; Android[Android] --> API2["/android/messages/{userId}"]; iPhone[iPhone] --> API3["/iphone/messages/{userId}"];
```

Client URL Controller

Android → /android/messages/{userId} → getMessages4Android

Web → /api/messages/{userId} → getMessages

iPhone → /iphone/messages/{userId} → getMessages4iPhone

Client URL Controller

Android → /android/messages/{userId} →

Web → /api/messages/{userId} → getMessages

iPhone → /iphone/messages/{userId} →

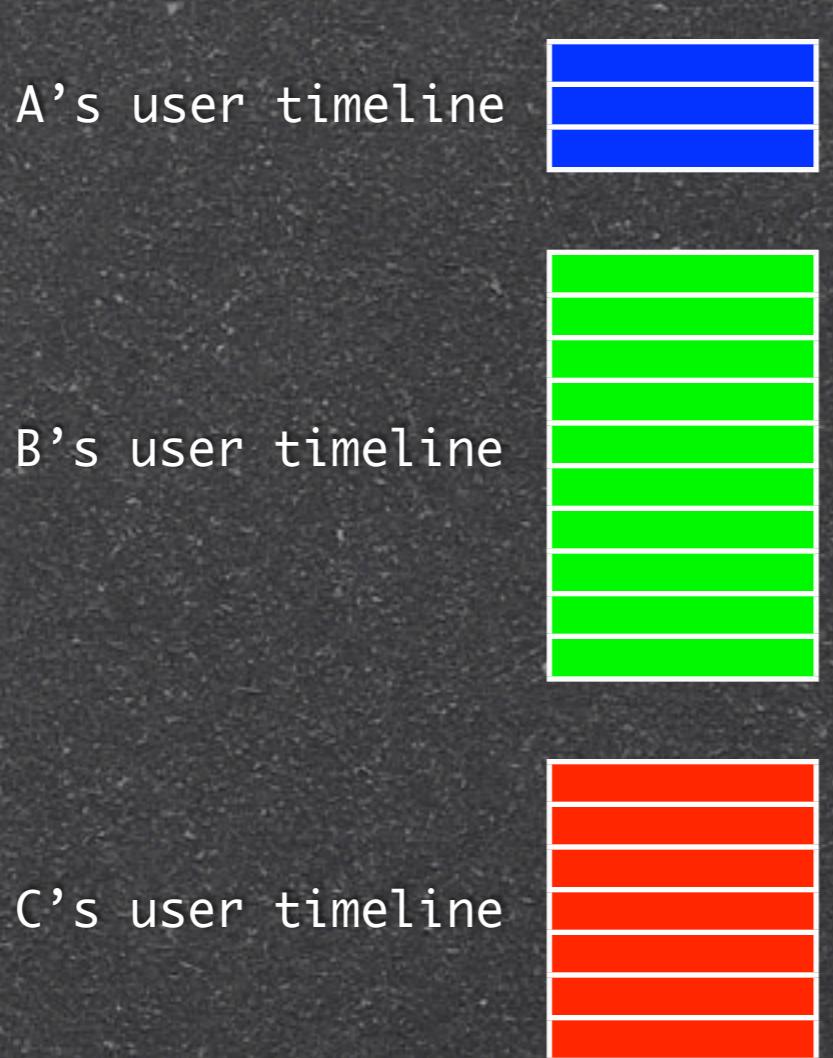
Client URL Controller

Android → /android/messages/{userId} →

Web → /api/messages/{userId} → getMessages

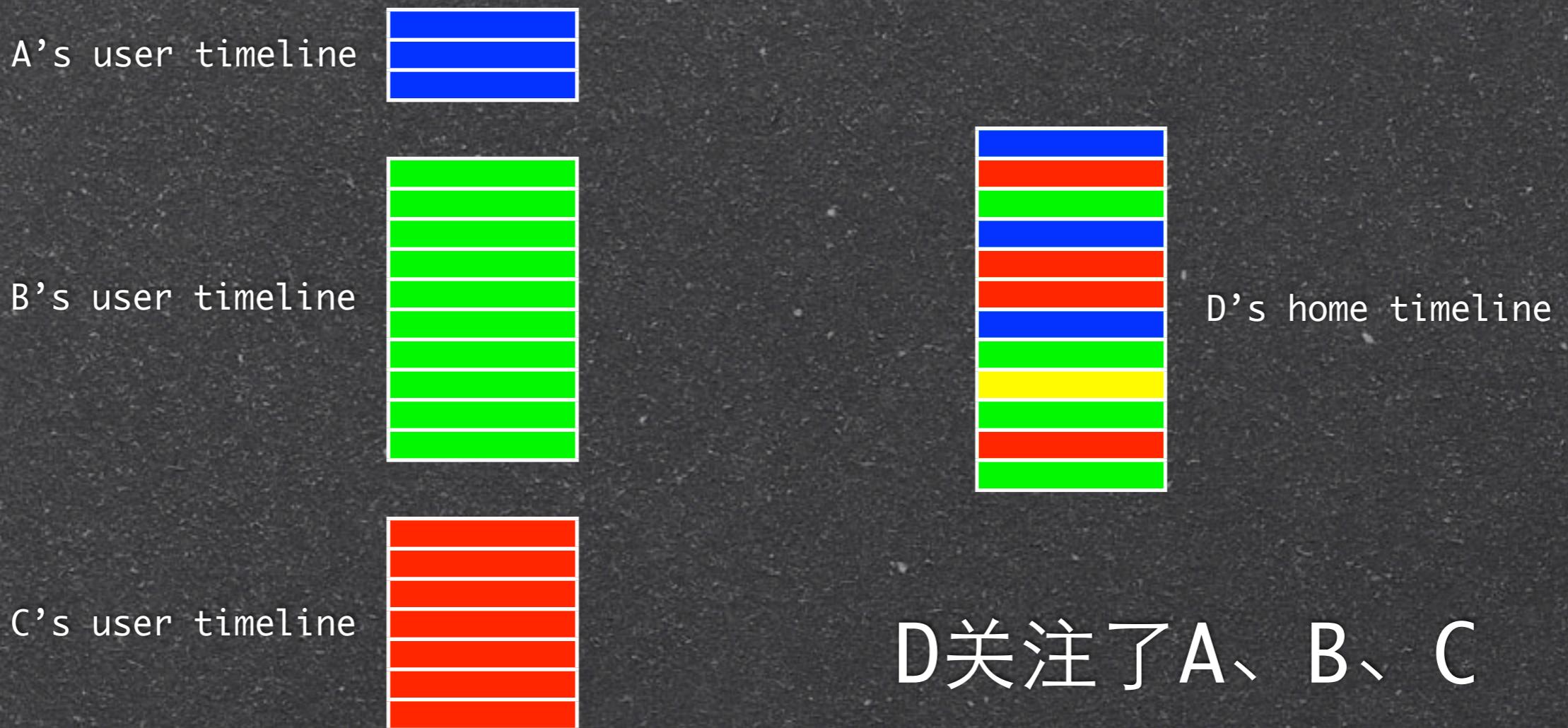
iPhone → /iphone/messages/{userId} → getMessages4iPhone

Timeline



D关注了A、B、C

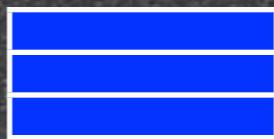
Timeline



一个SQL语句就能搞定？

```
SELECT id, time FROM timeline_t  
WHERE userId in (?,?,?,?,?)  
AND time>? AND time<=?  
ORDER BY time DESC limit ?
```

Timeline 前传一

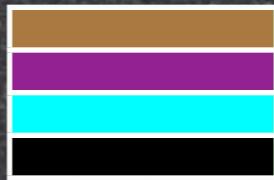


A's user timeline

D、E、F关注A



D's home timeline



E's home timeline



F's home timeline

Timeline 前传一

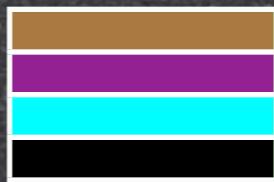


A's user timeline

D、E、F关注A



D's home timeline



E's home timeline



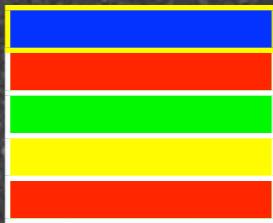
F's home timeline

Timeline 前传一



A's user timeline

D、E、F关注A



D's home timeline



E's home timeline



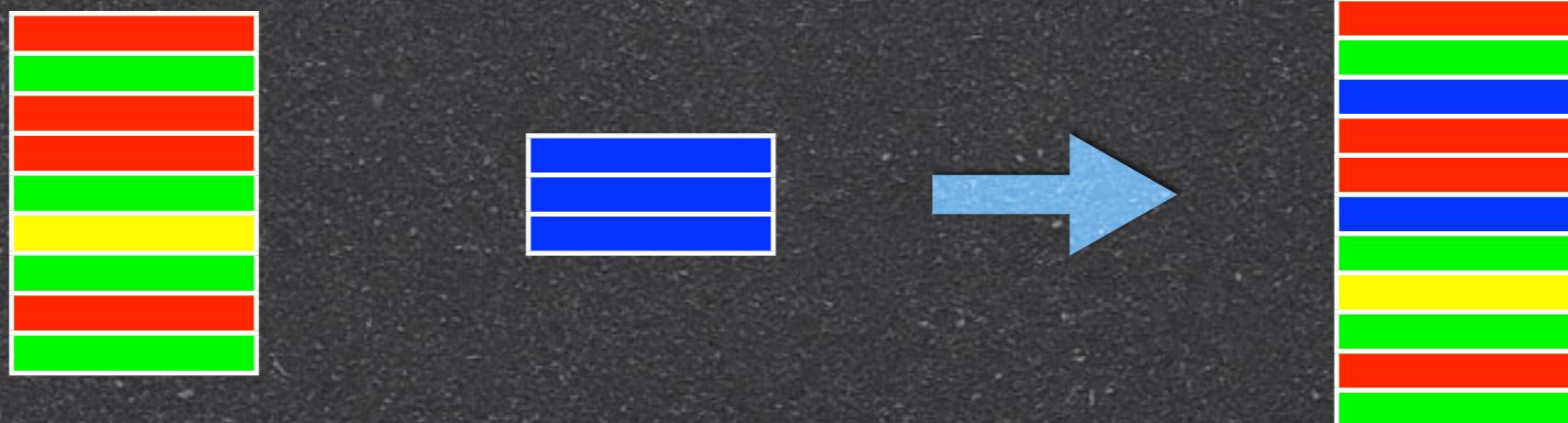
F's home timeline

- 读性能优秀——空间换时间
- 写时异步分发
- 数据冗余、操作时延
- 关注、取消关注

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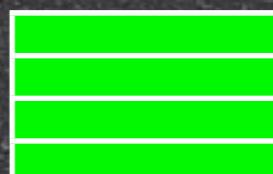


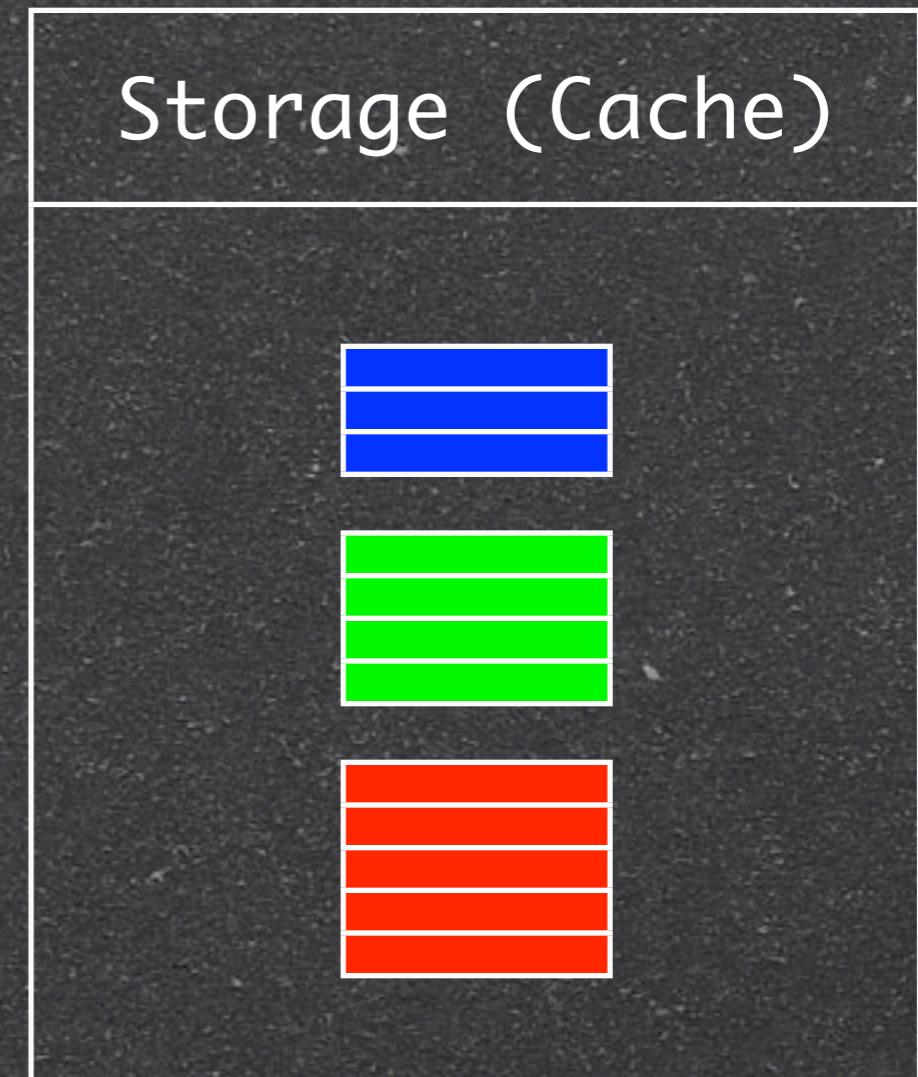
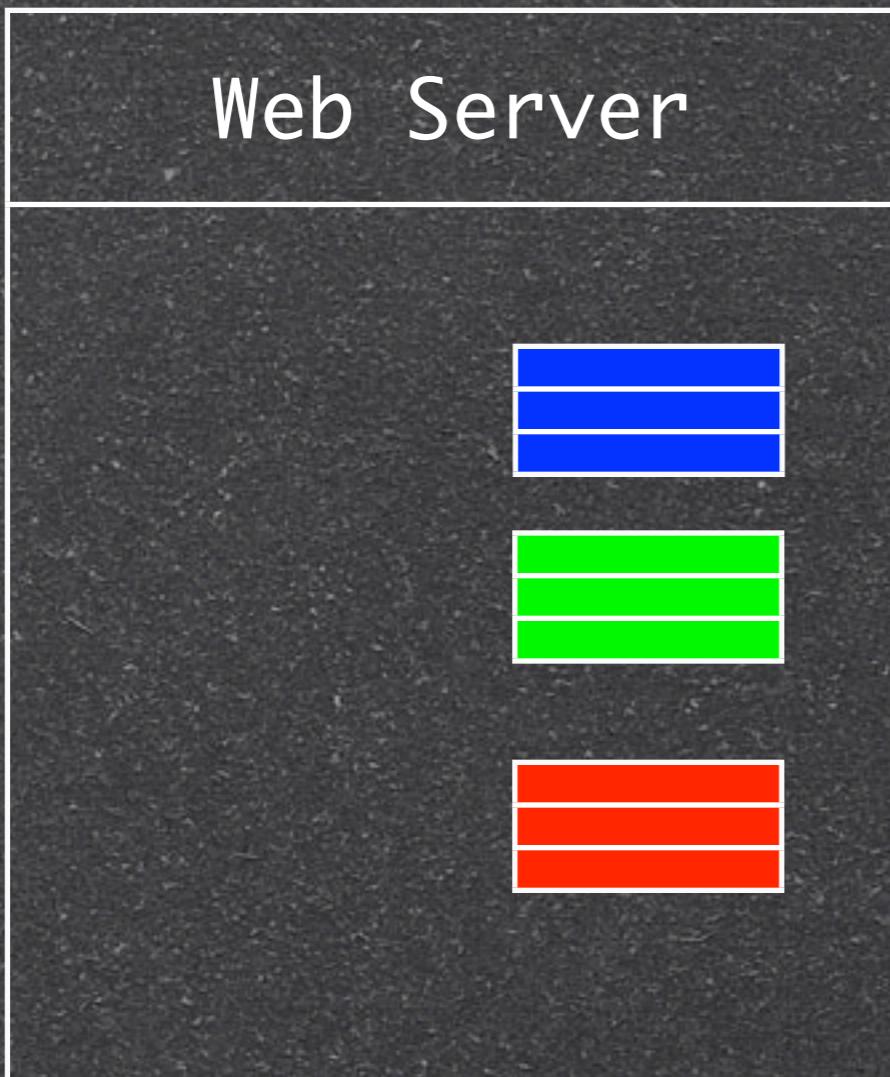
Timeline前传二

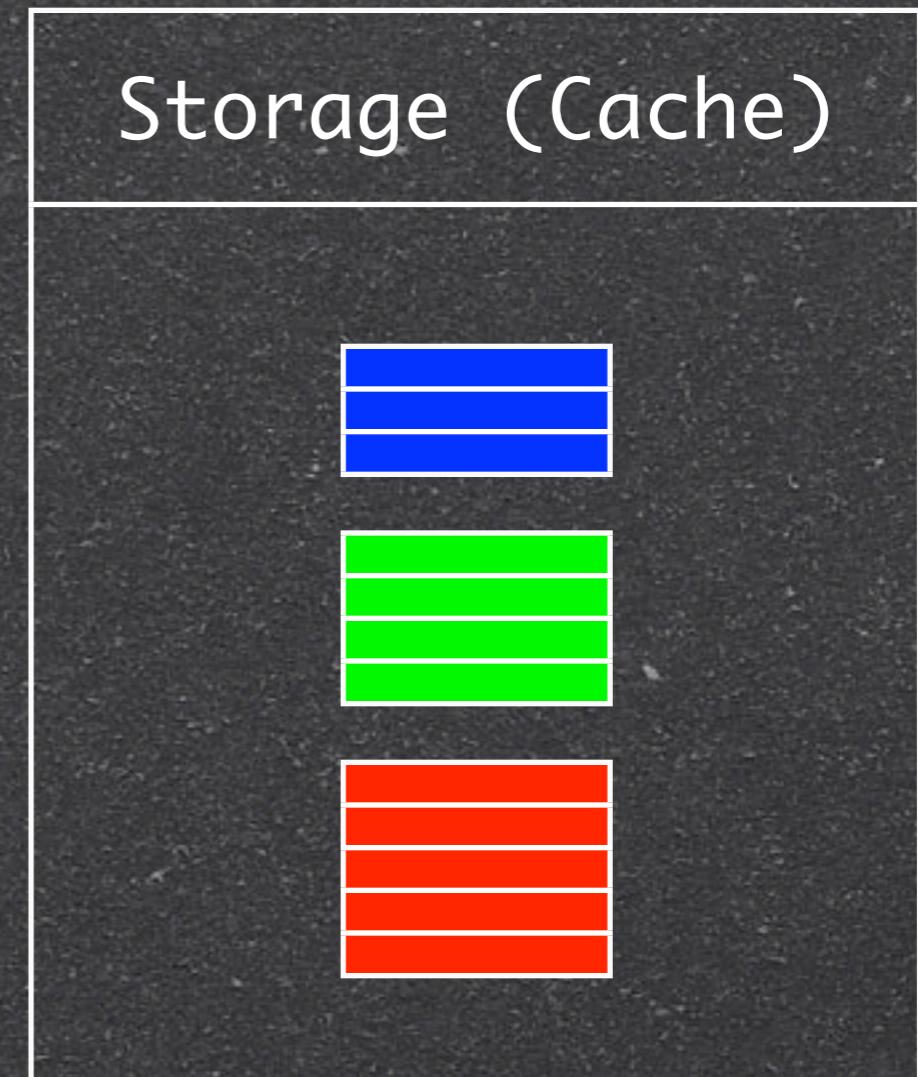
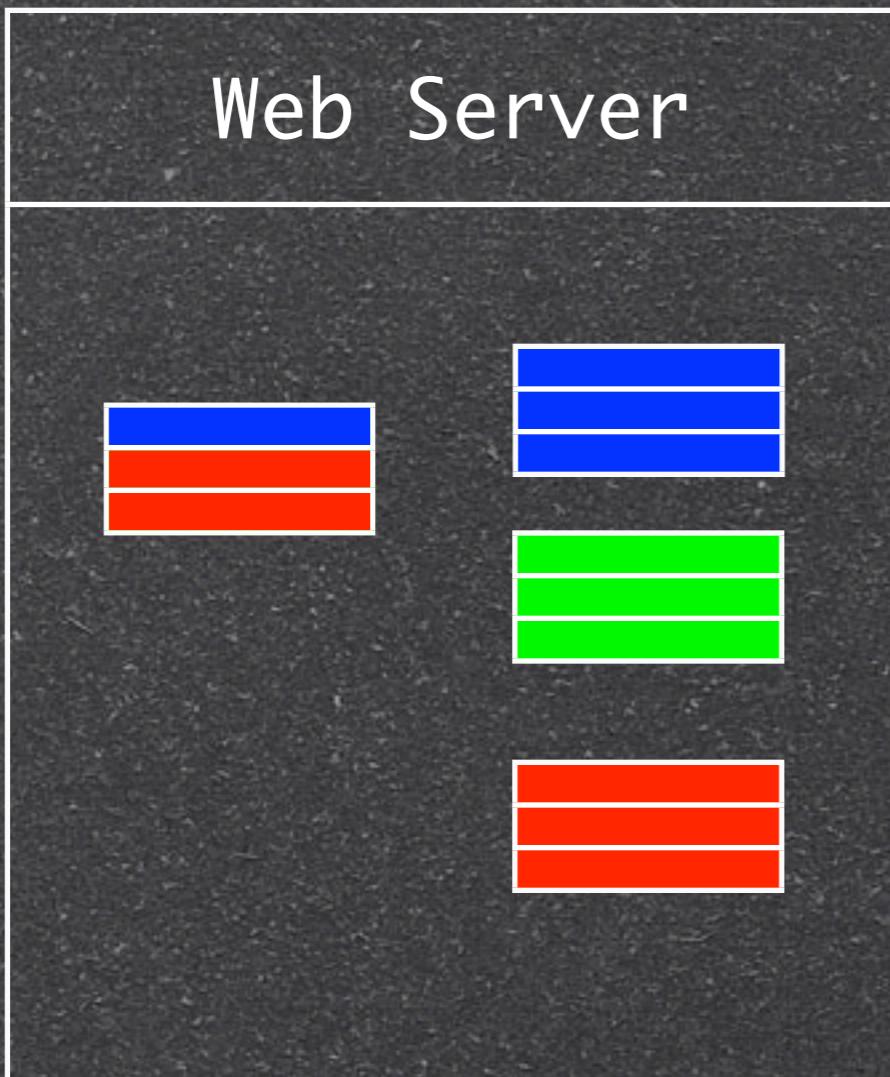
- 前提——关注上限
- 缓存所有用户的user timeline
- 获取所有关注的人的user timeline，在应用层归并，实时生成home timeline

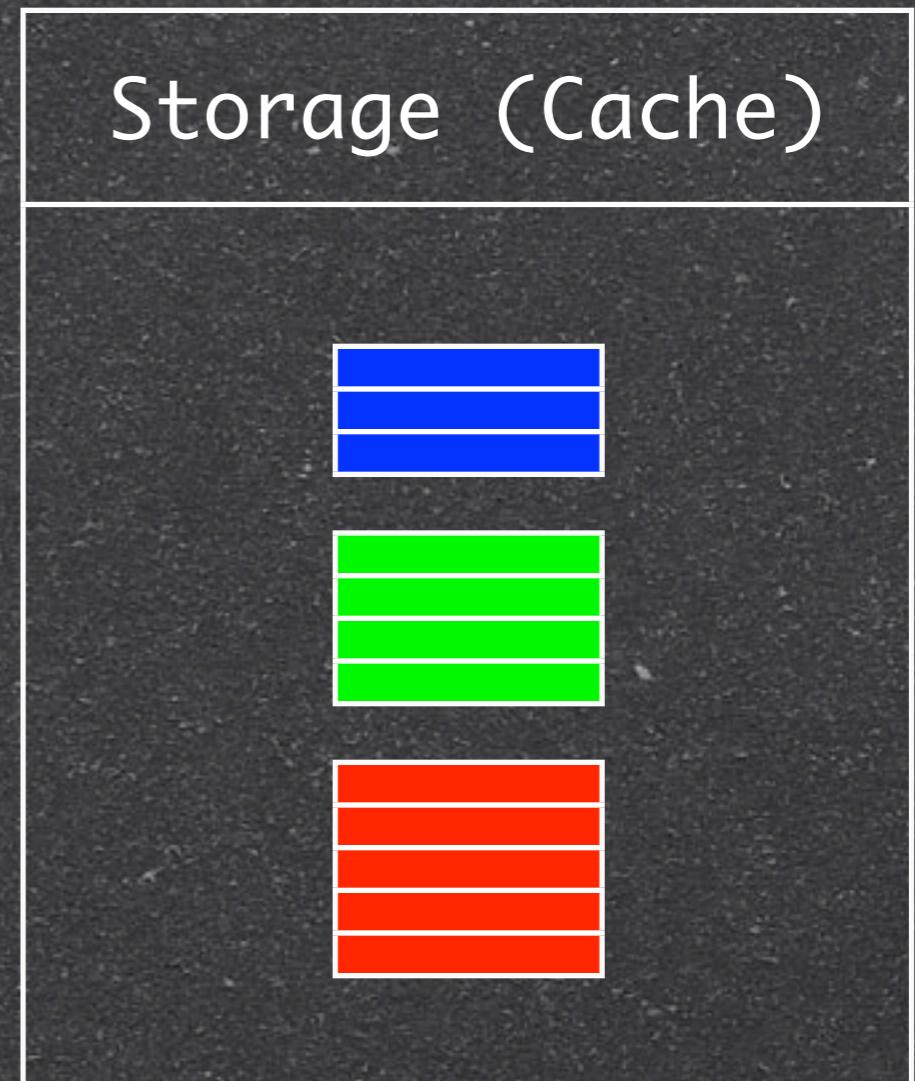
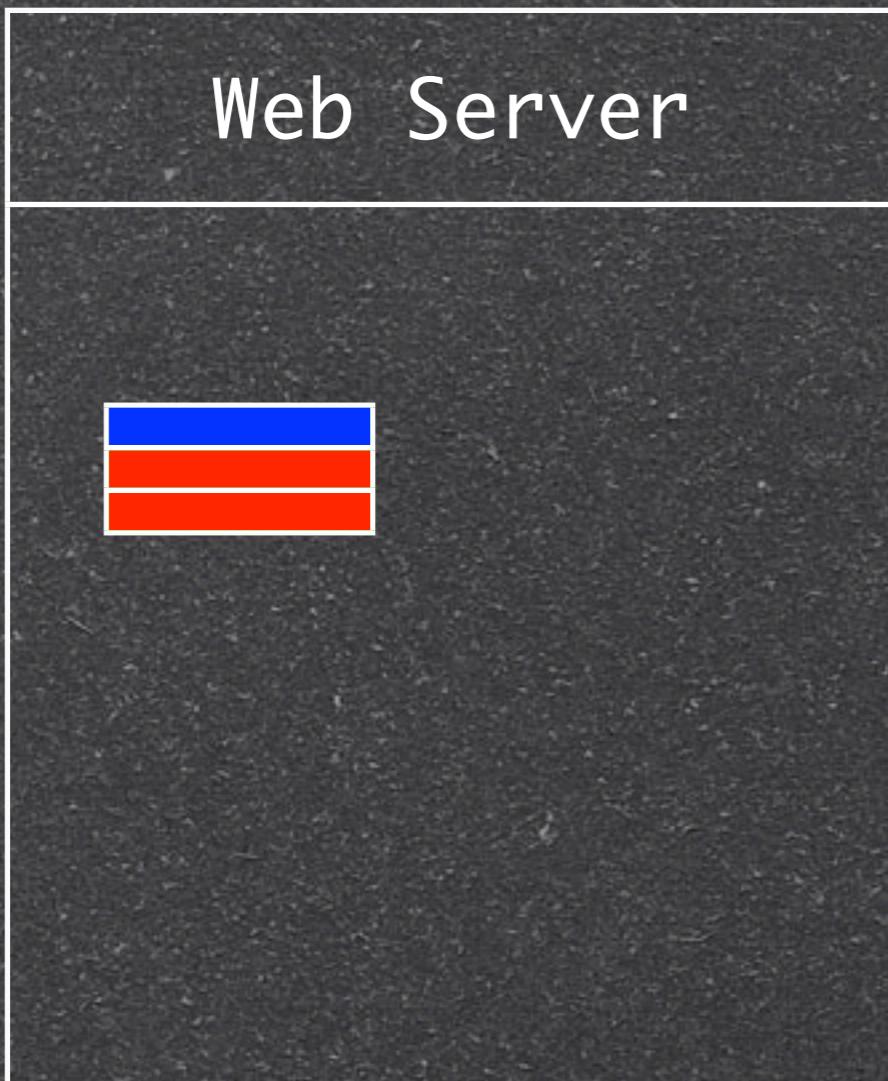
Web Server

Storage (Cache)









逻辑简单、灵活--关注、取消关注

写性能突出--没有分发、没有时延

数据没有冗余

网络I/O--获取所有关注用户的user timeline

逻辑简单、灵活--关注、取消关注

写性能突出--没有分发、没有时延

数据没有冗余

网络I/O--获取所有关注用户的user timeline



爱转角
izhuangjiao.com

如何减小网络I/O?

Web Server

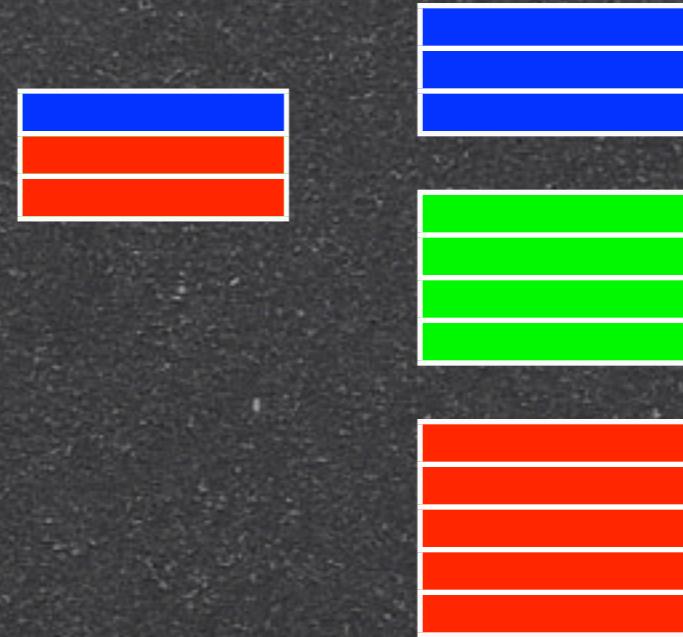
Storage



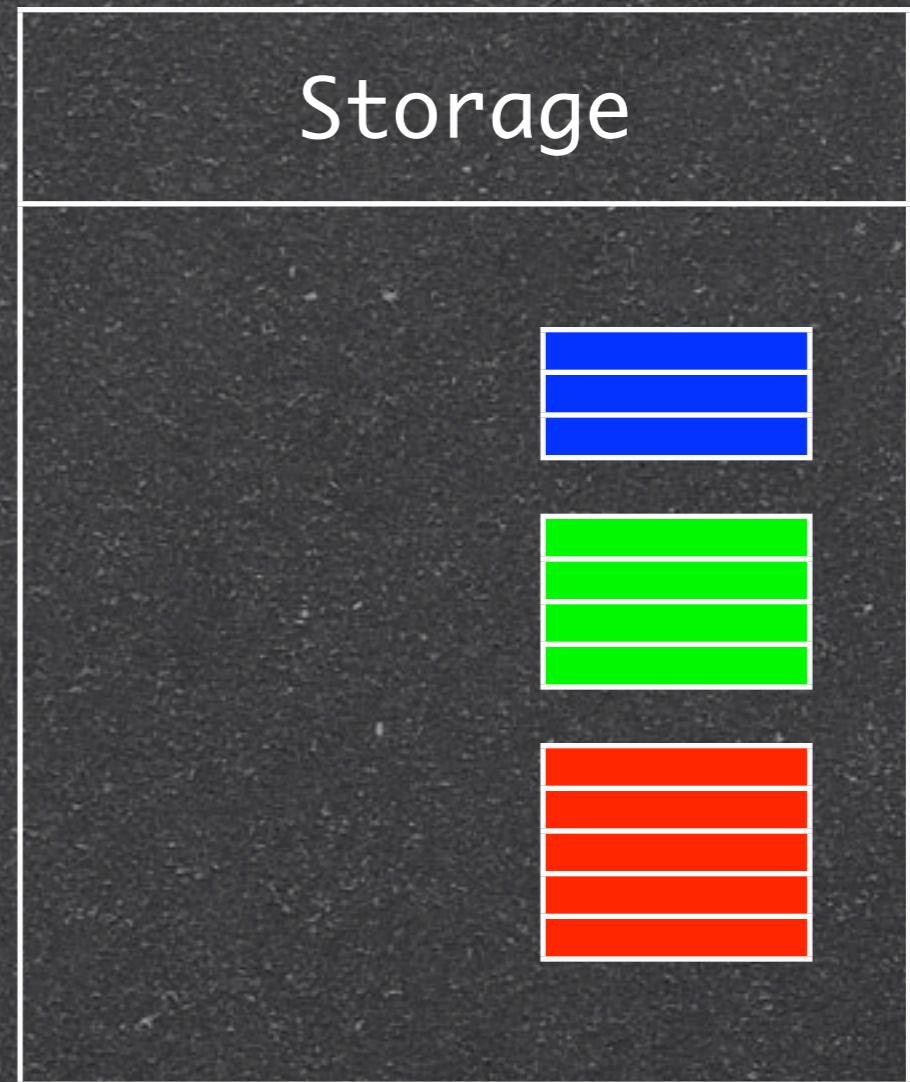
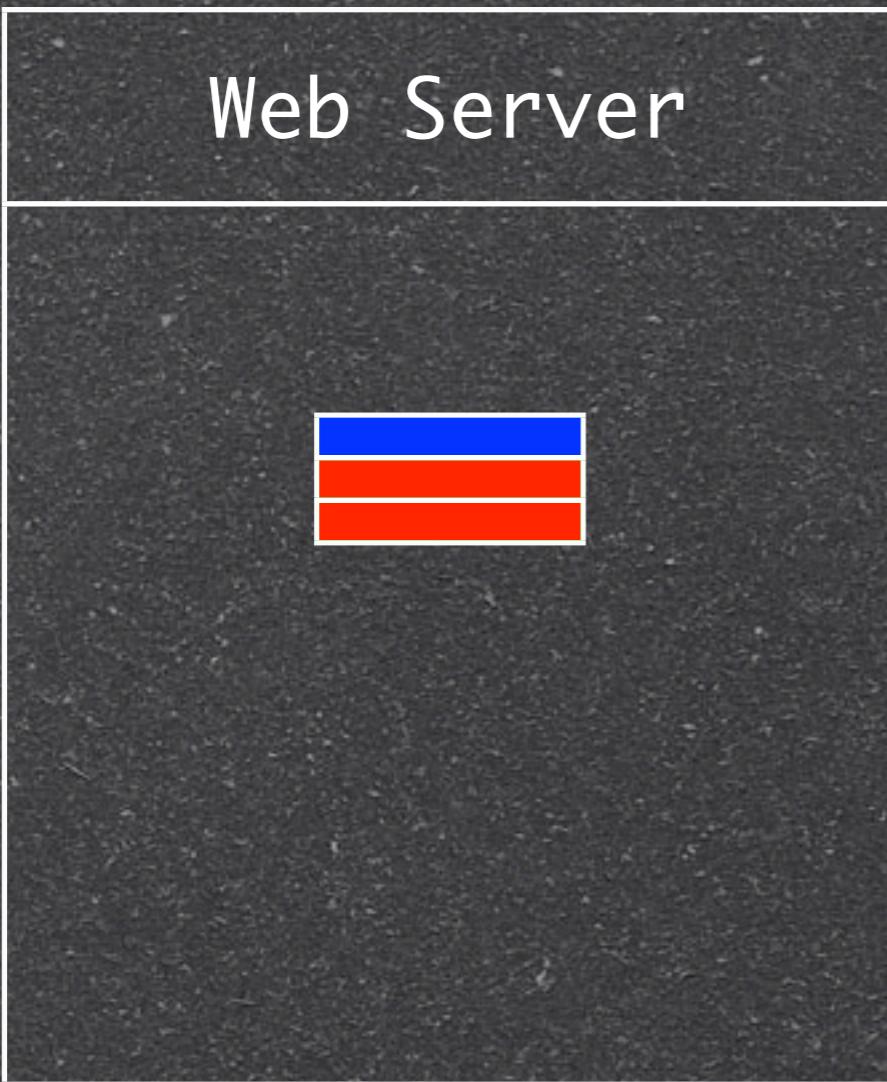
在存储层完成归并操作

Web Server

Storage



在存储层完成归并操作



在存储层完成归并操作

在Redis中实现多路归并

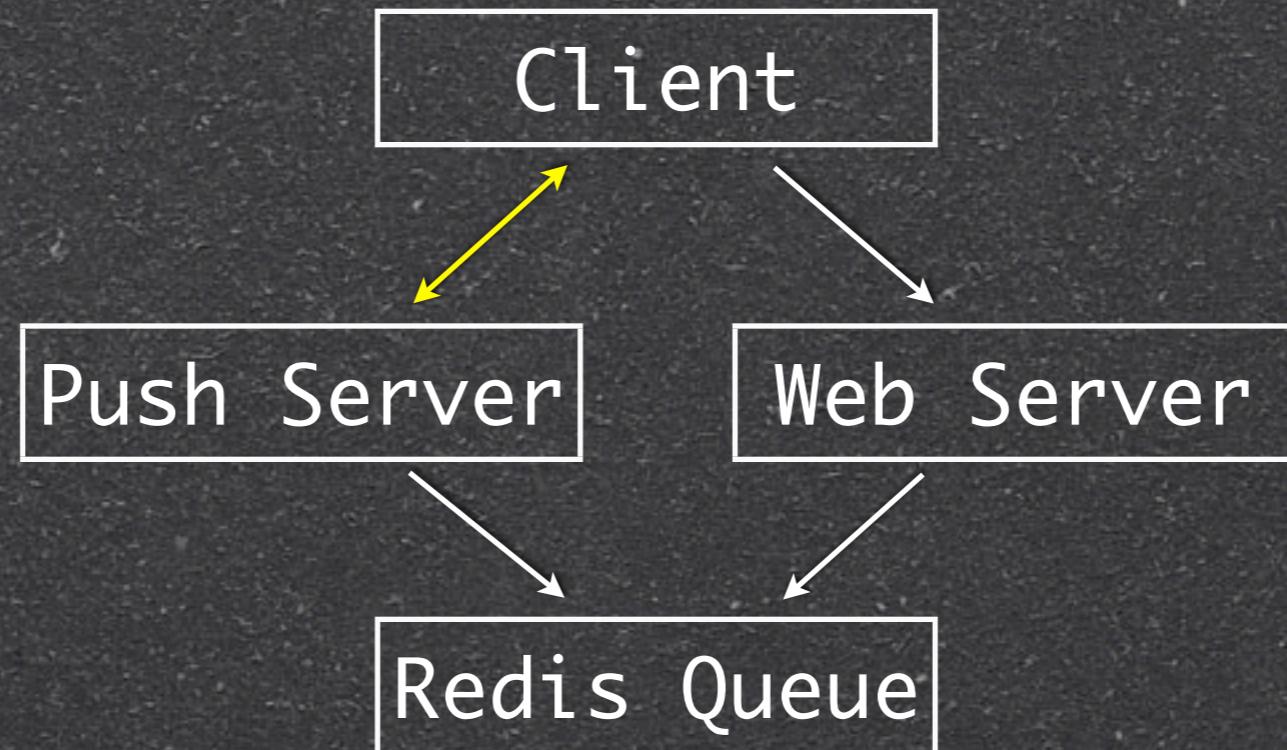
- 全内存
- ZSet存储user timeline, 按时间排序
- 增加zmerge接口

持久化队列

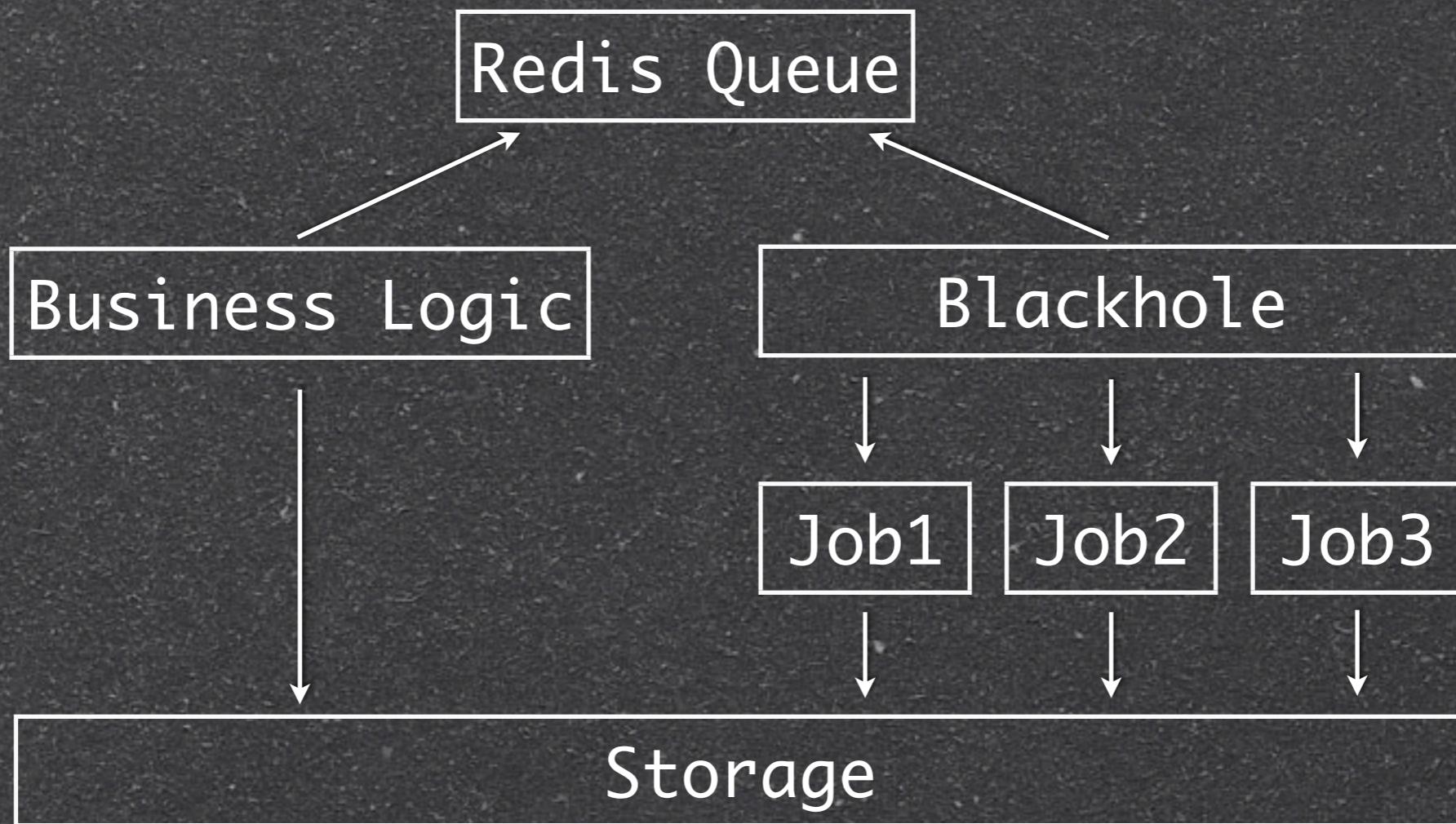
- Redis中List的几个基本操作
 - LPUSH、RPUSH、LPOP、RPOP
- 实现FIFO持久化队列
 - offer --- RPUSH
 - take --- LPOP

消息推送

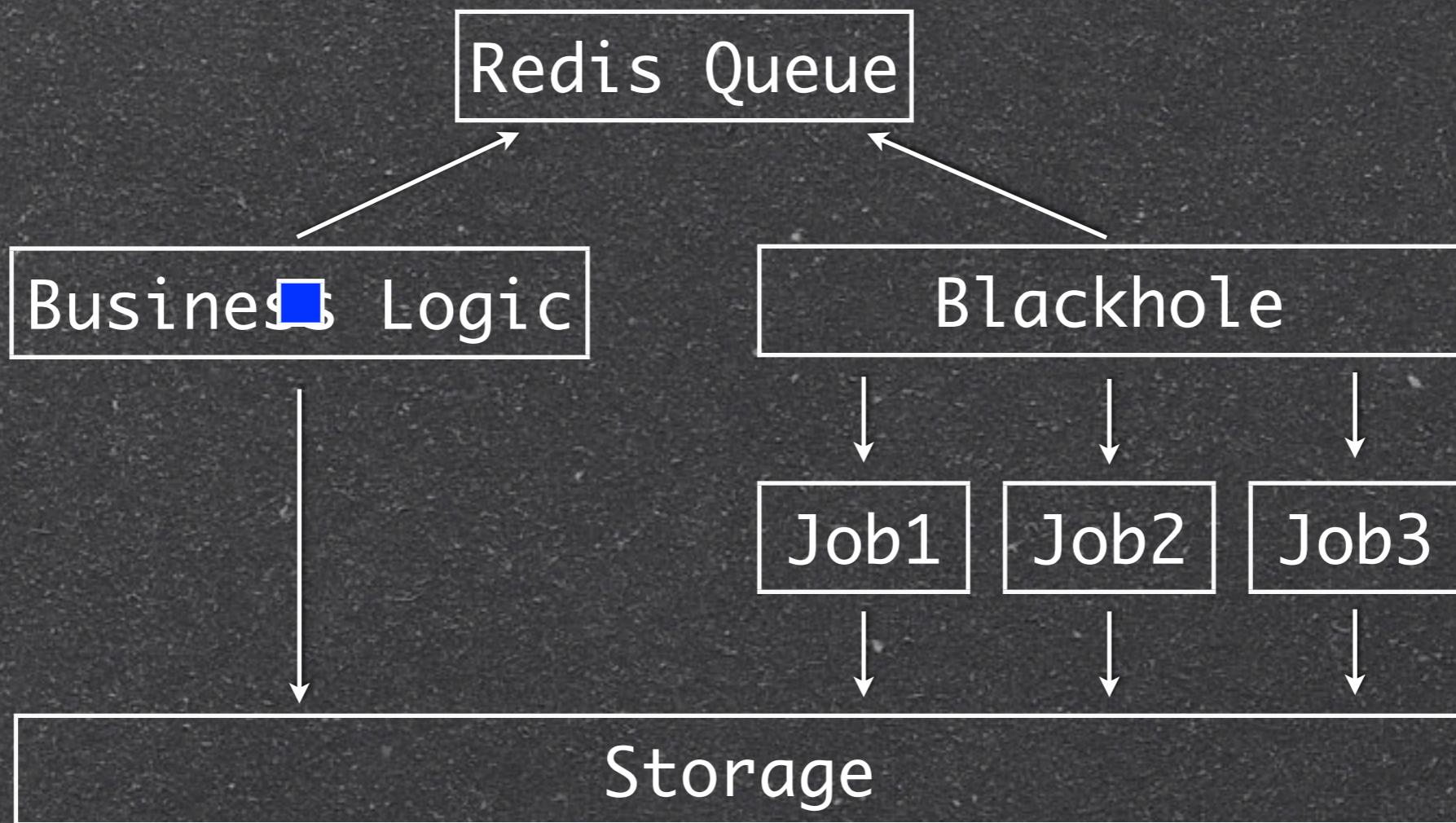
- 通用的PushServer，与业务逻辑无关
- 长连接--Jetty Continuation



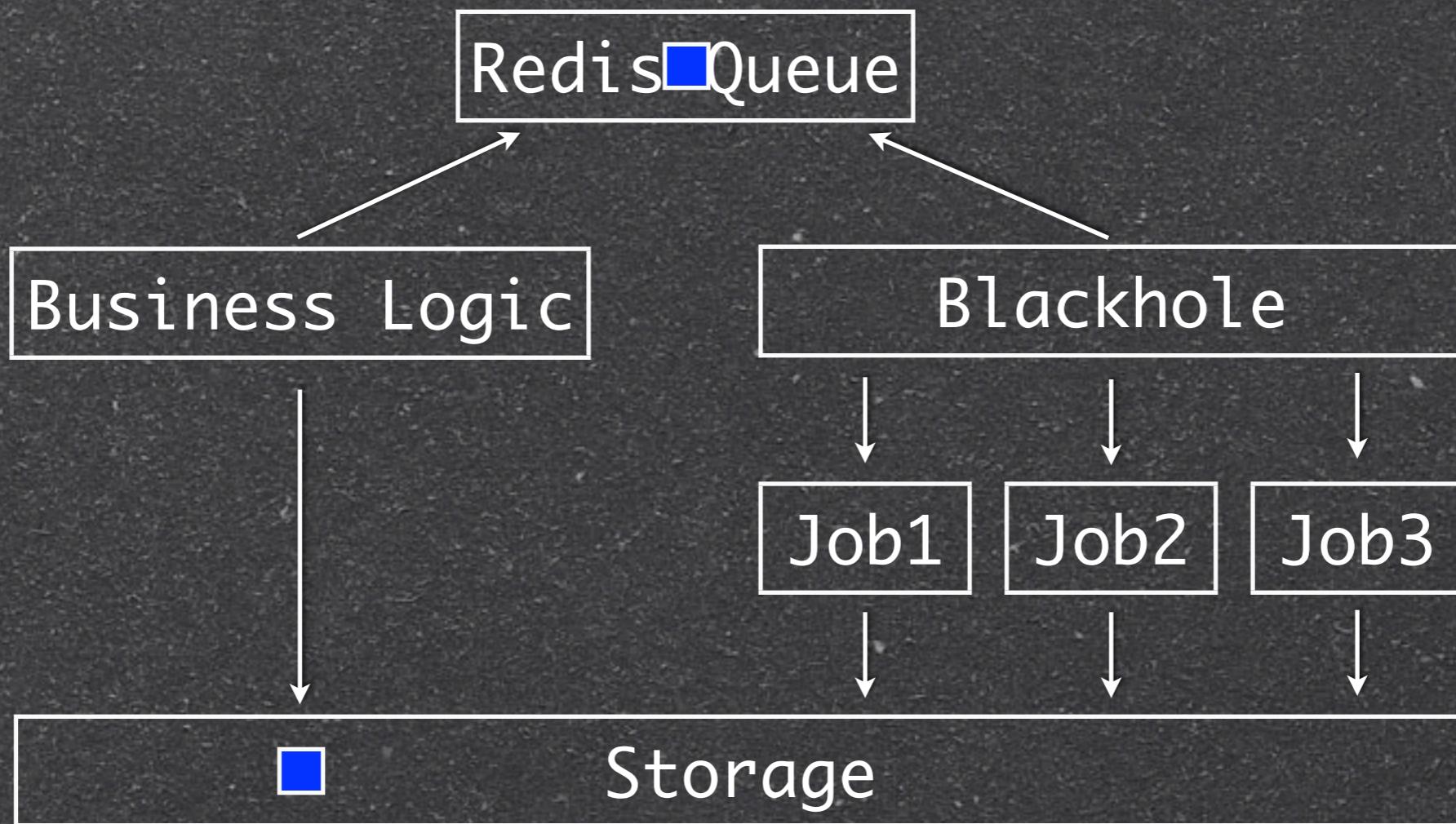
Blackhole



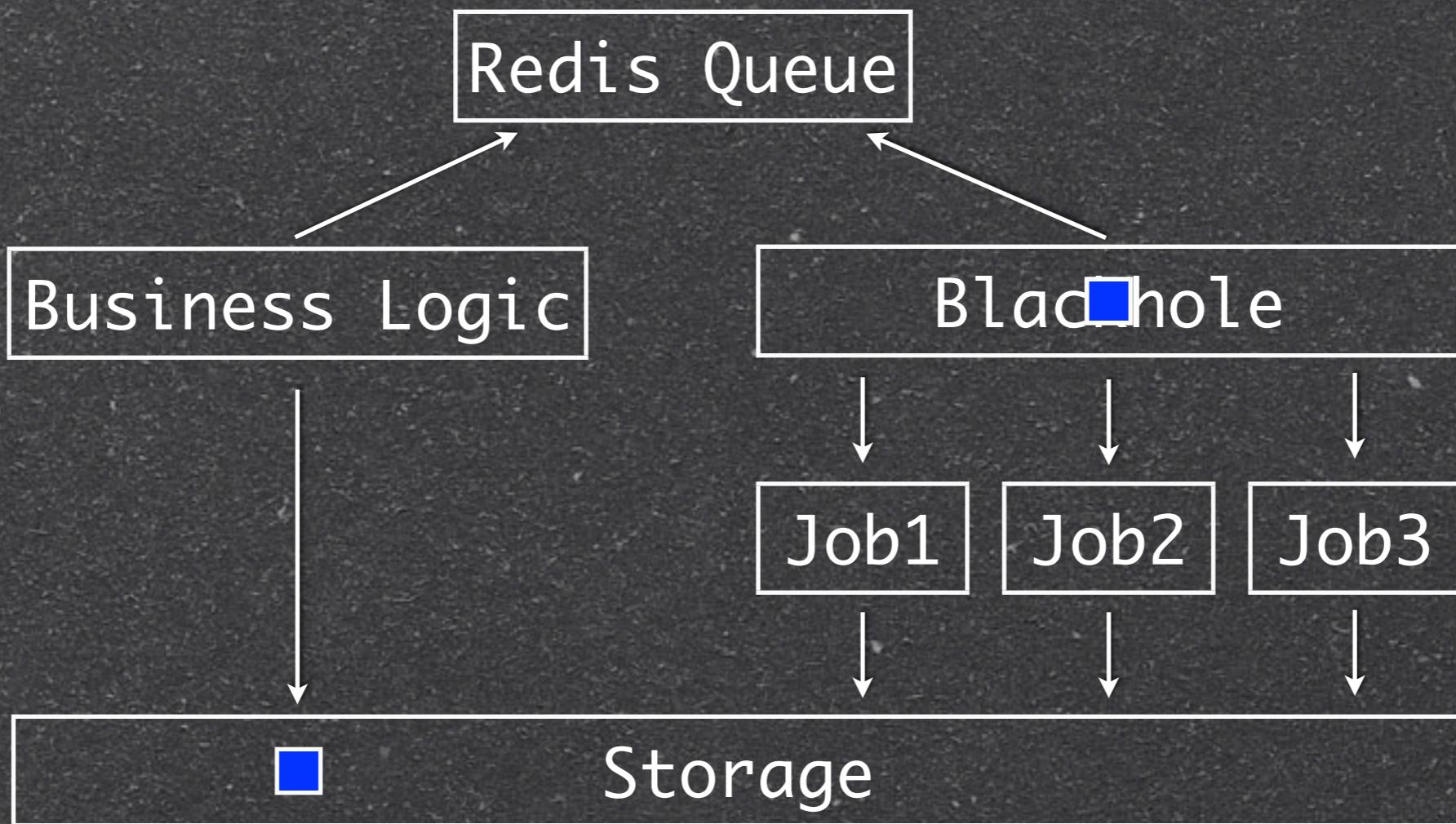
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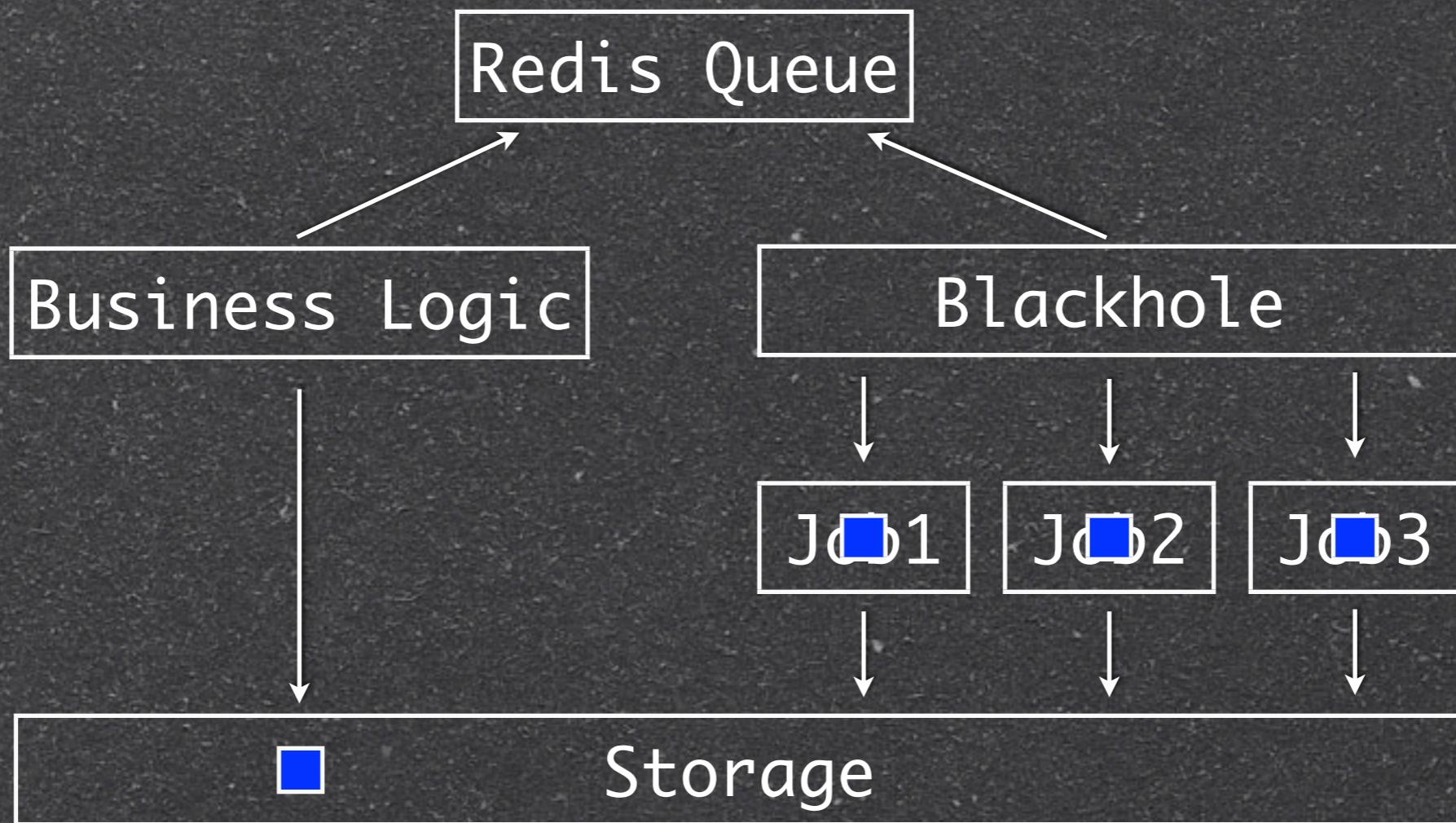
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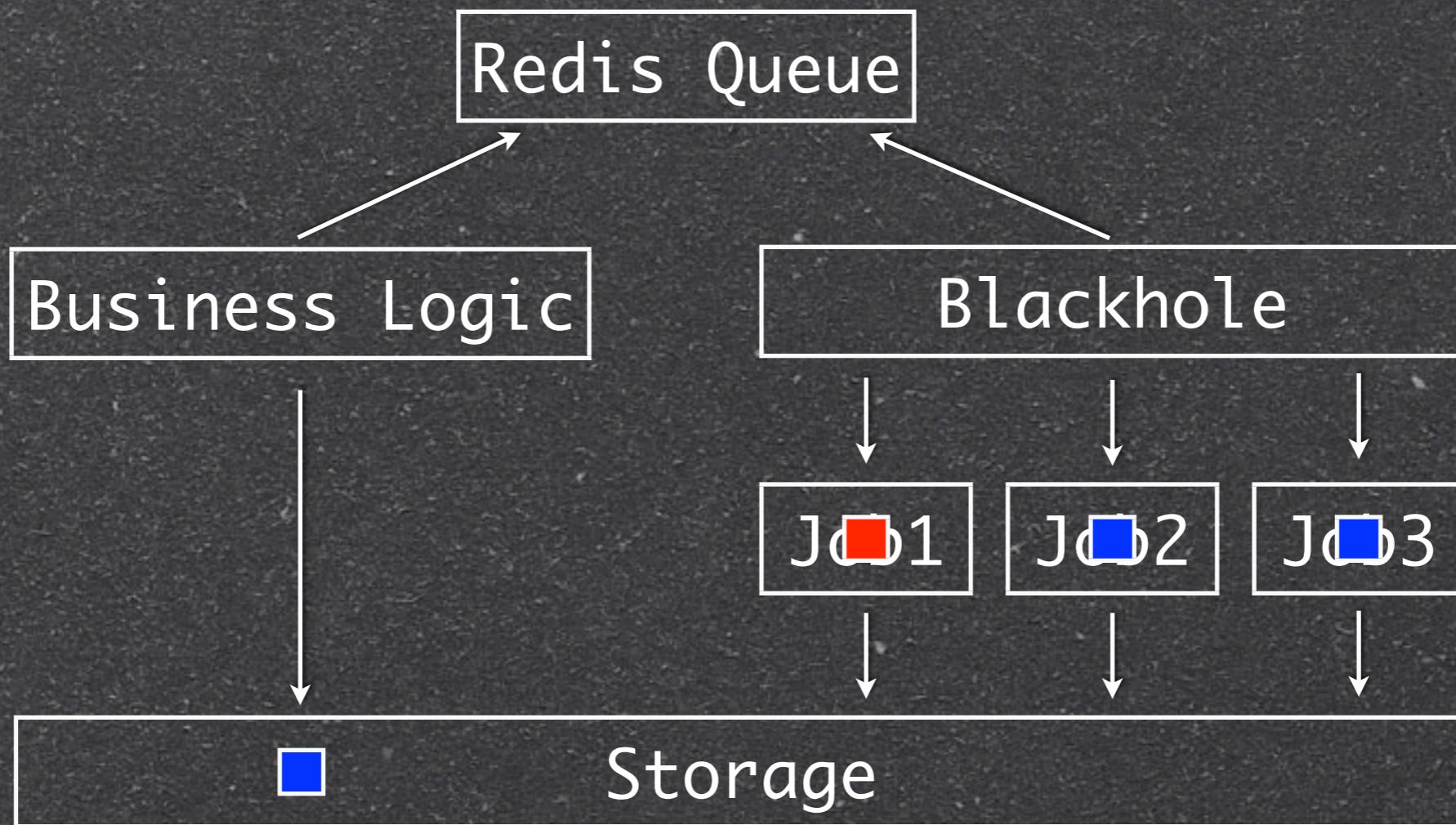
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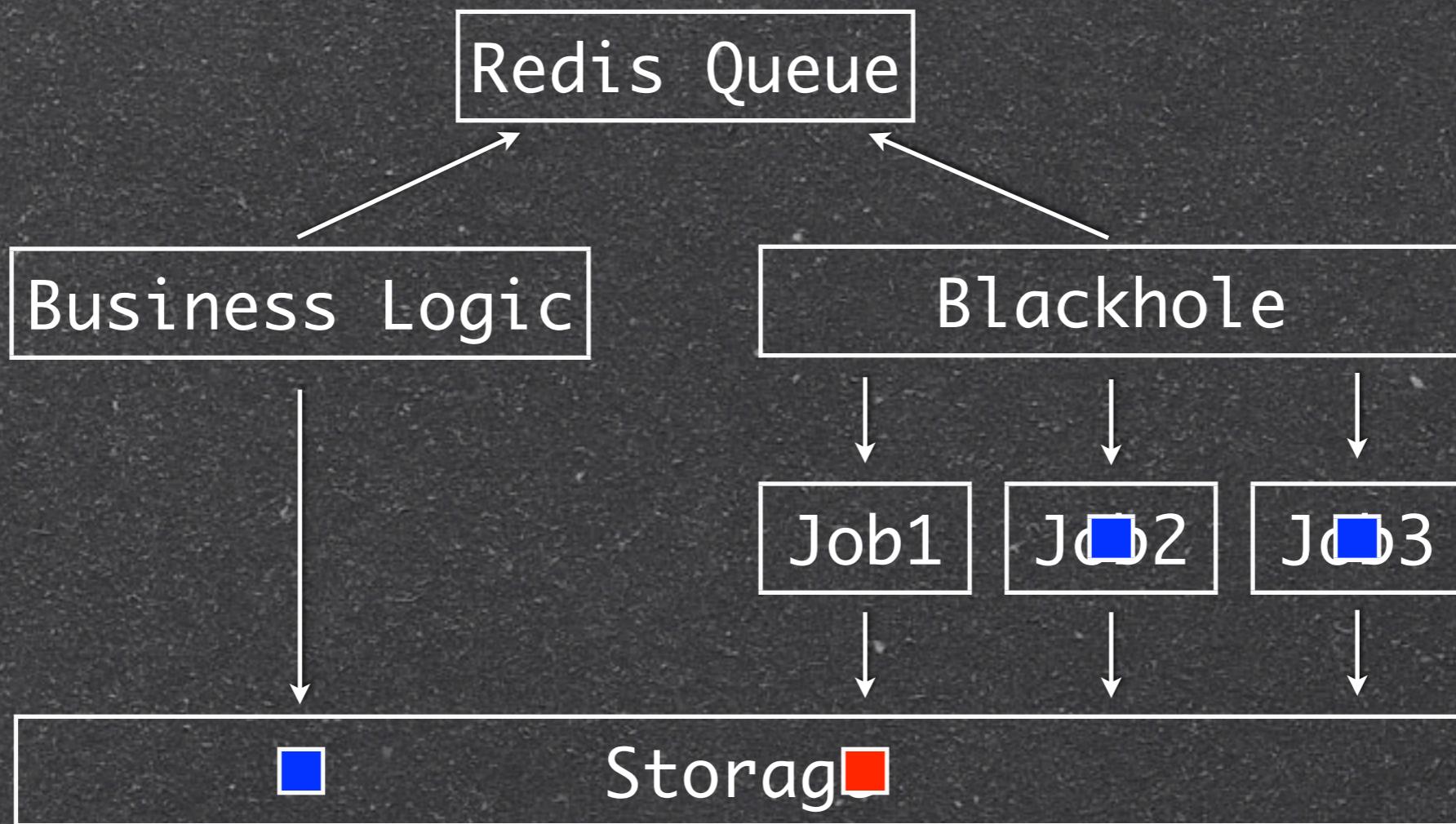
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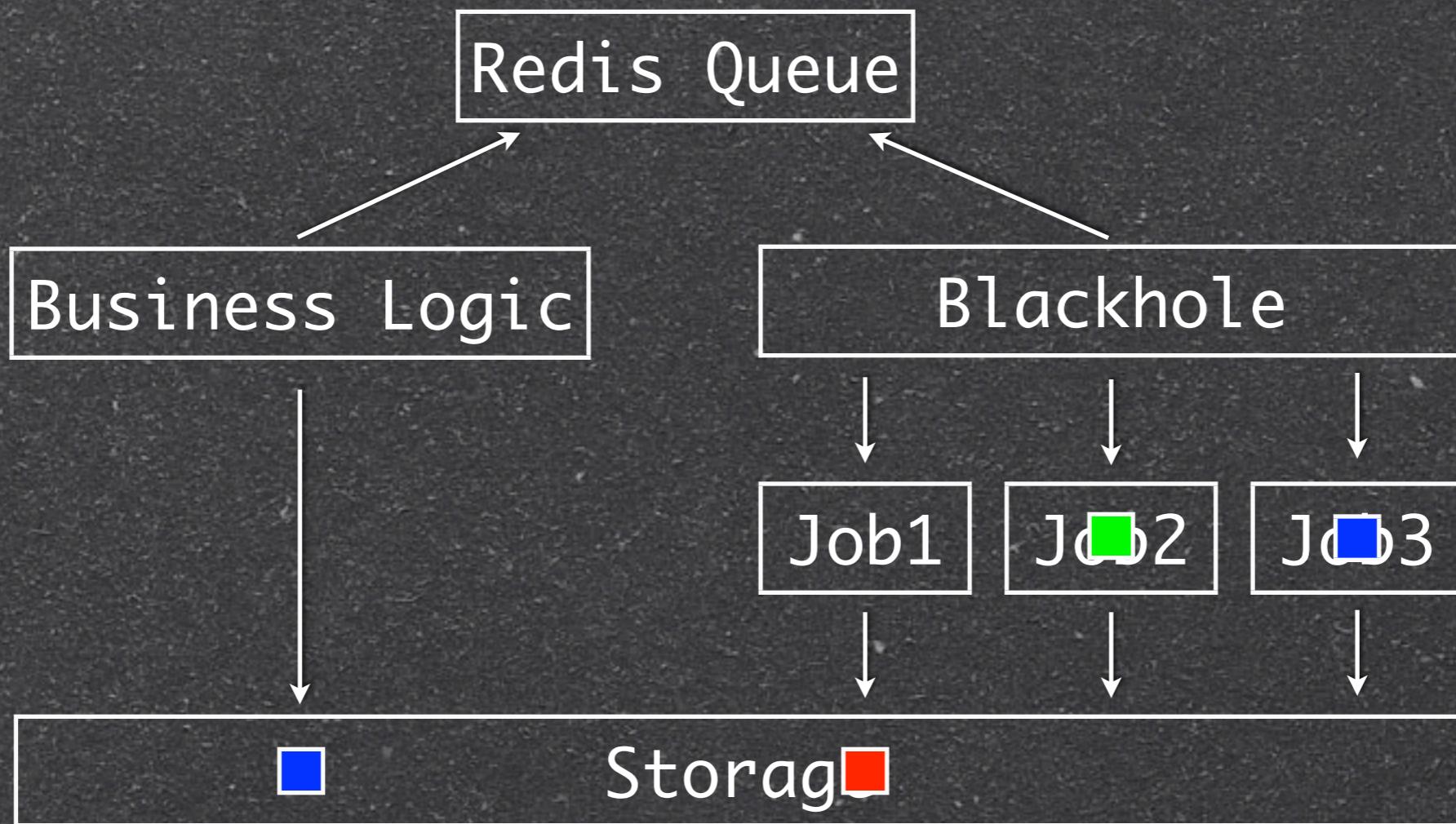
Blackhole



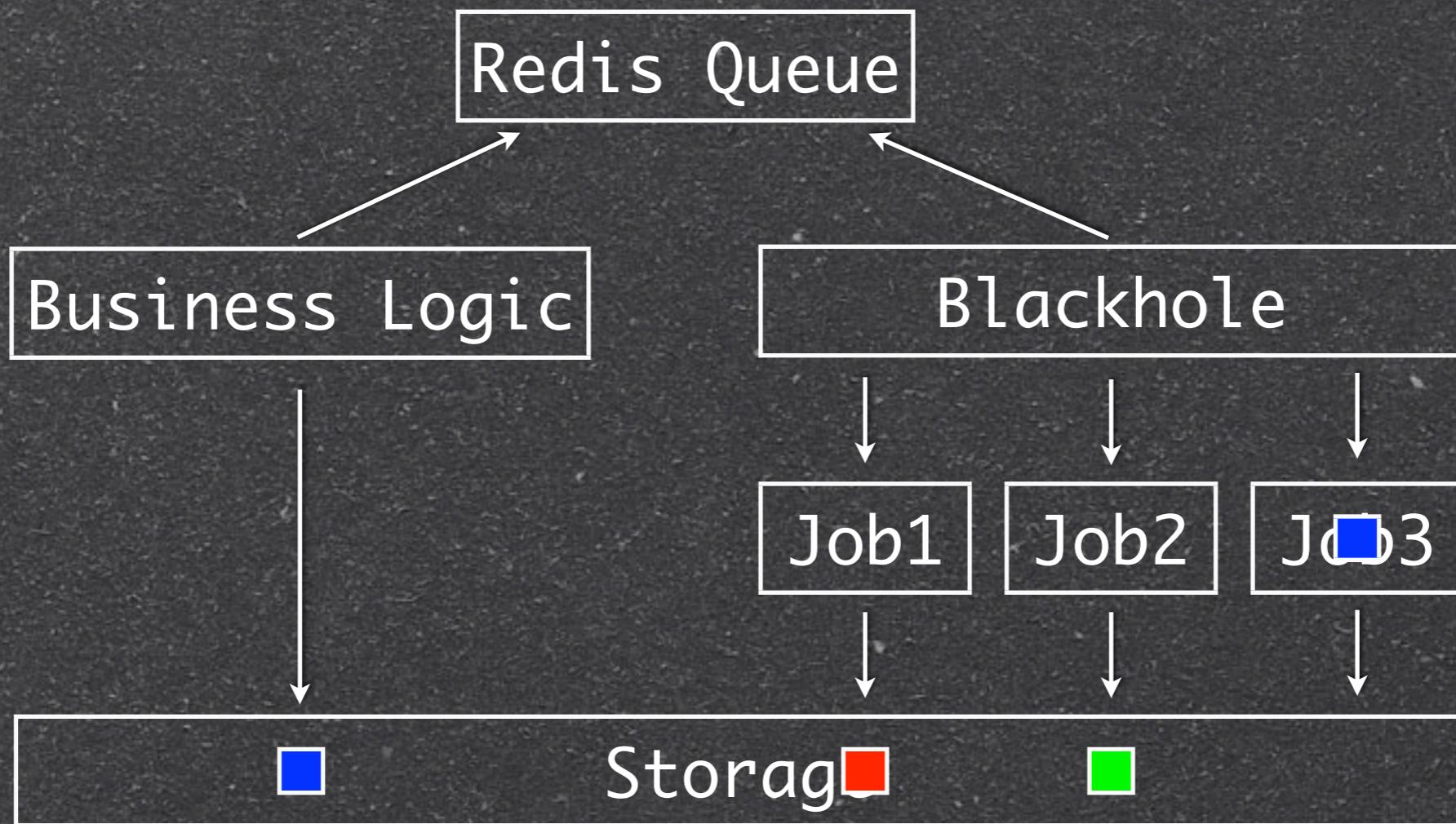
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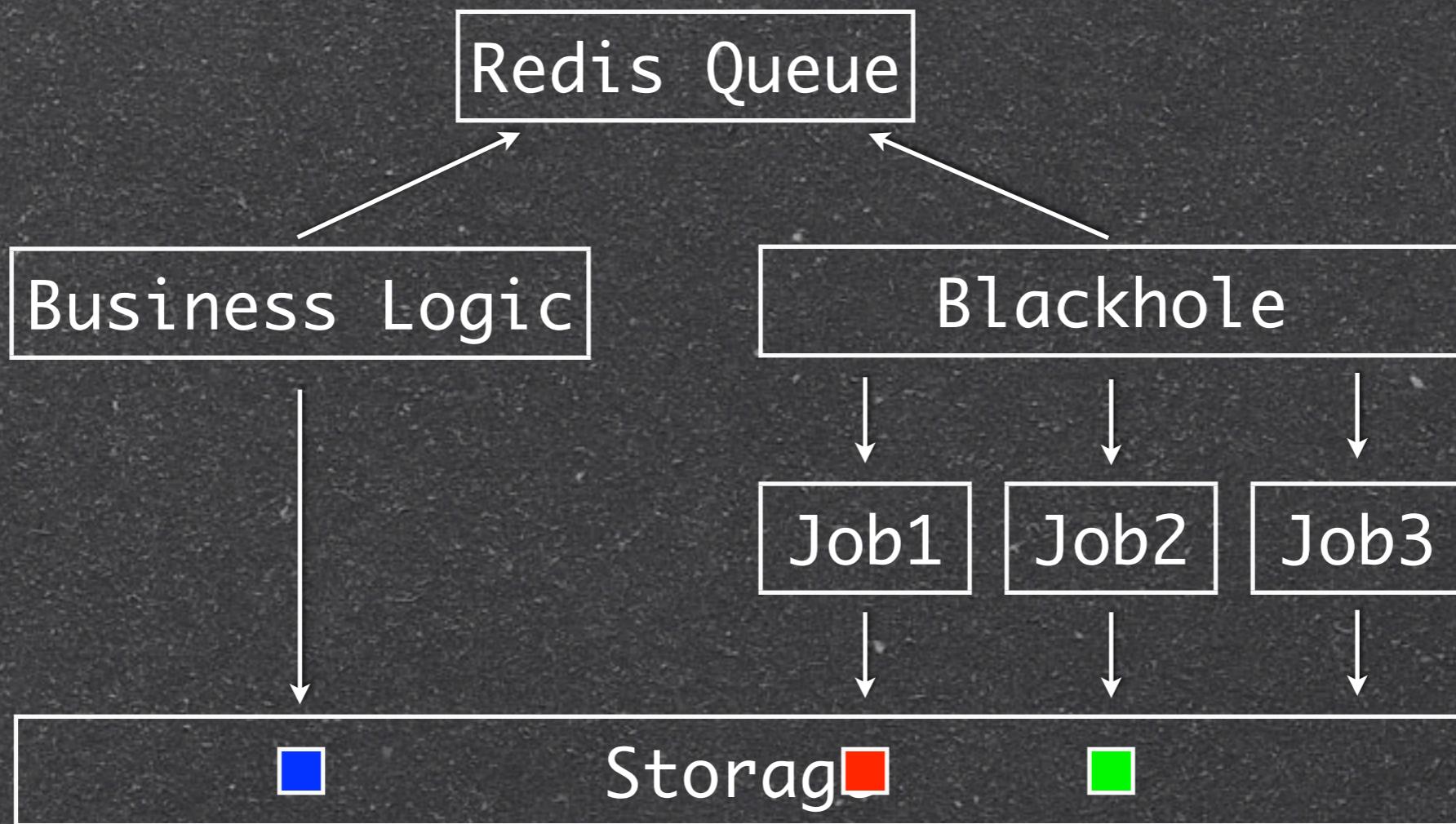
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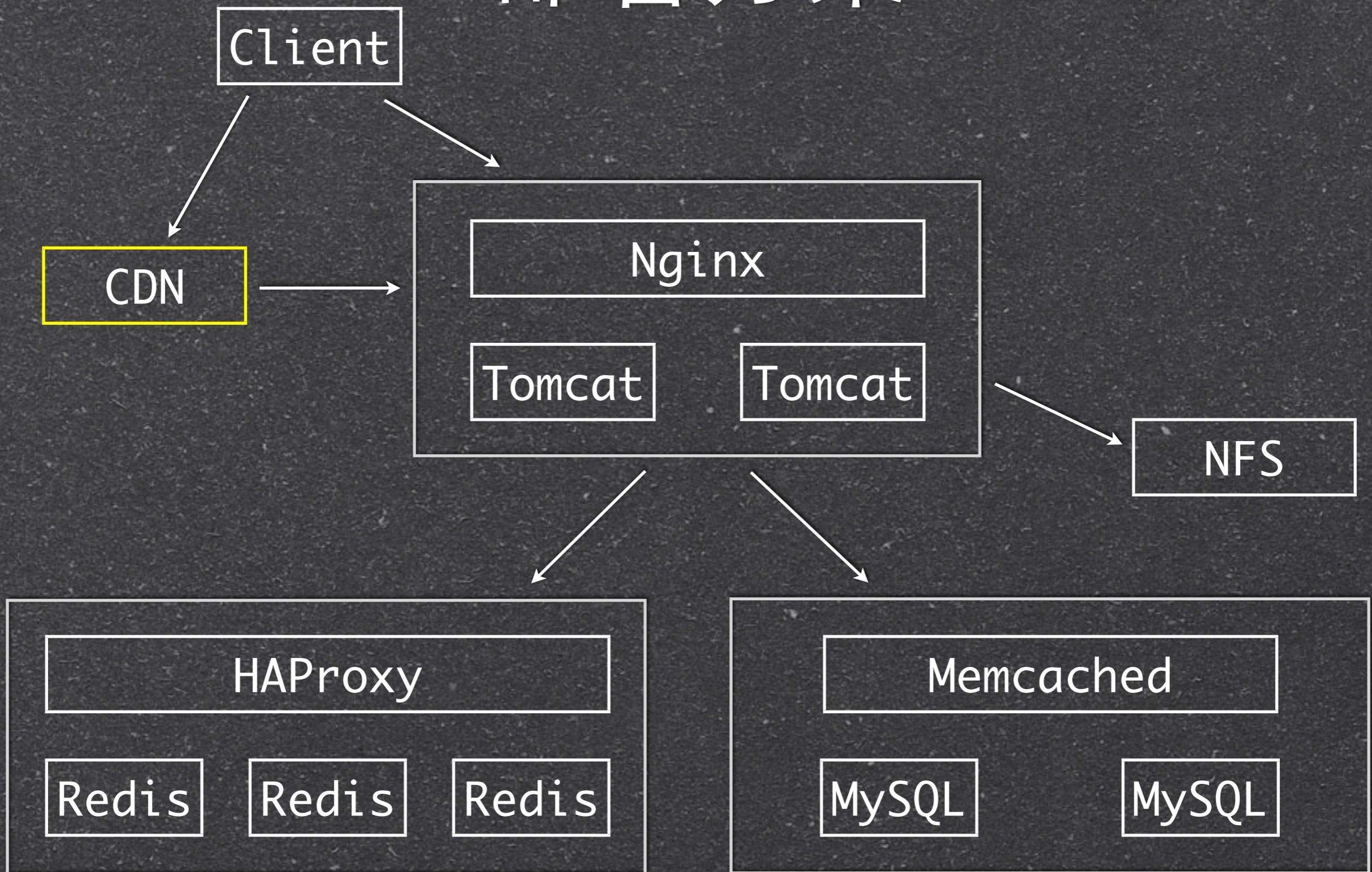
Blackhole



Blackhole



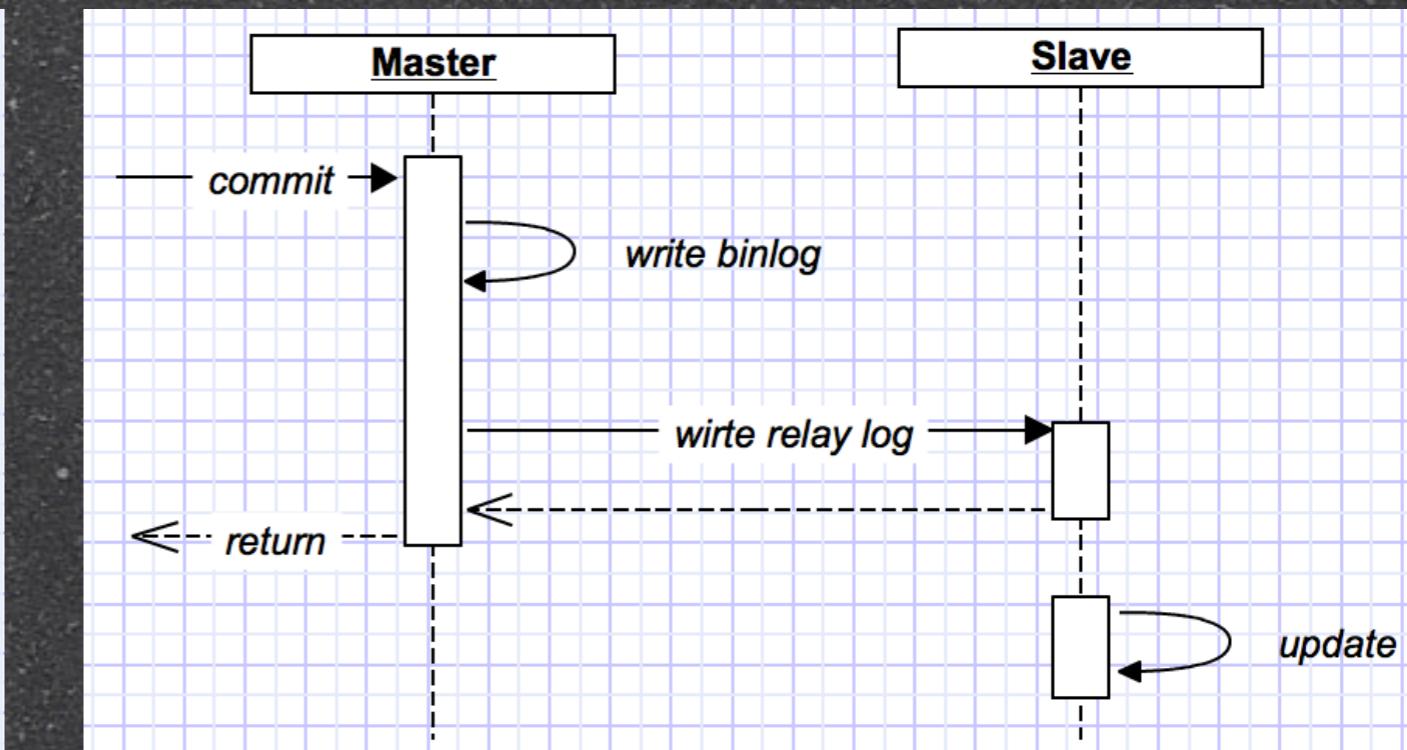
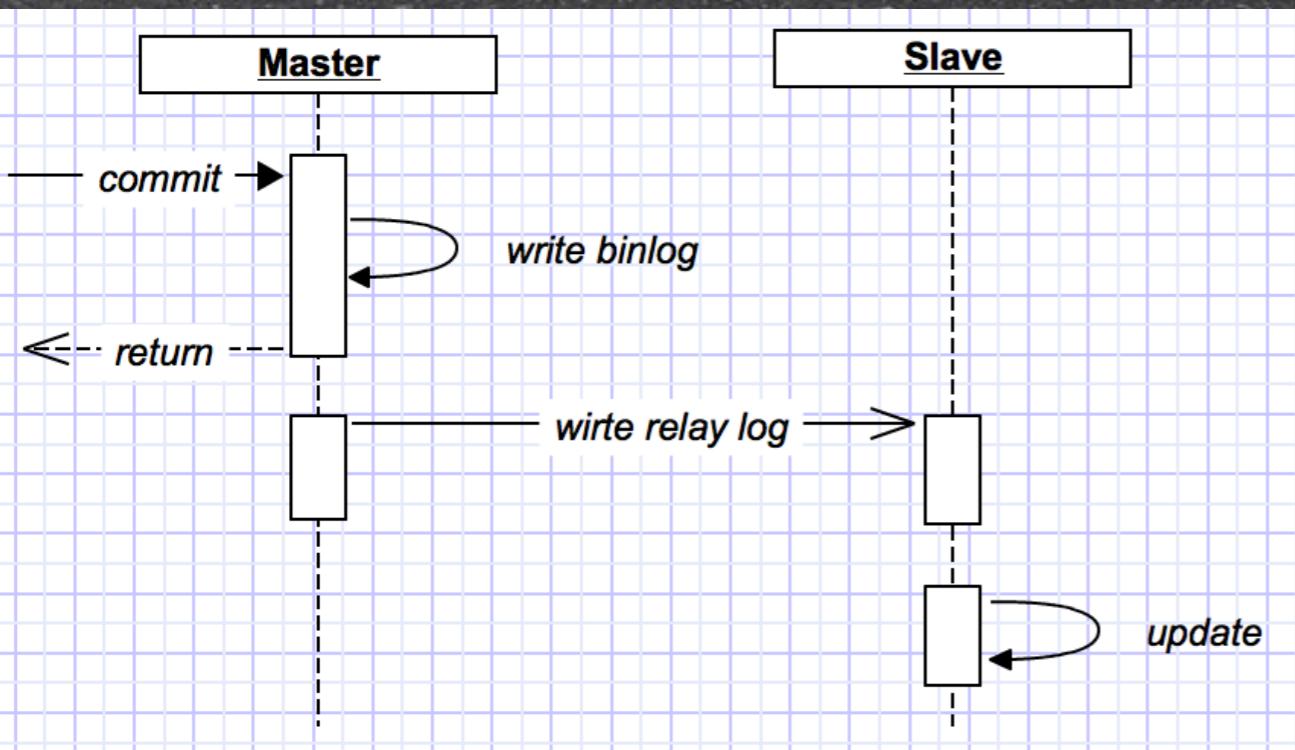
部署方案



NFS

- 图片和文件服务器
- 分布式文件系统 vs NFS
- 多个NFS Server, 分散磁盘读写请求
- RAID + 定时备份

MySQL

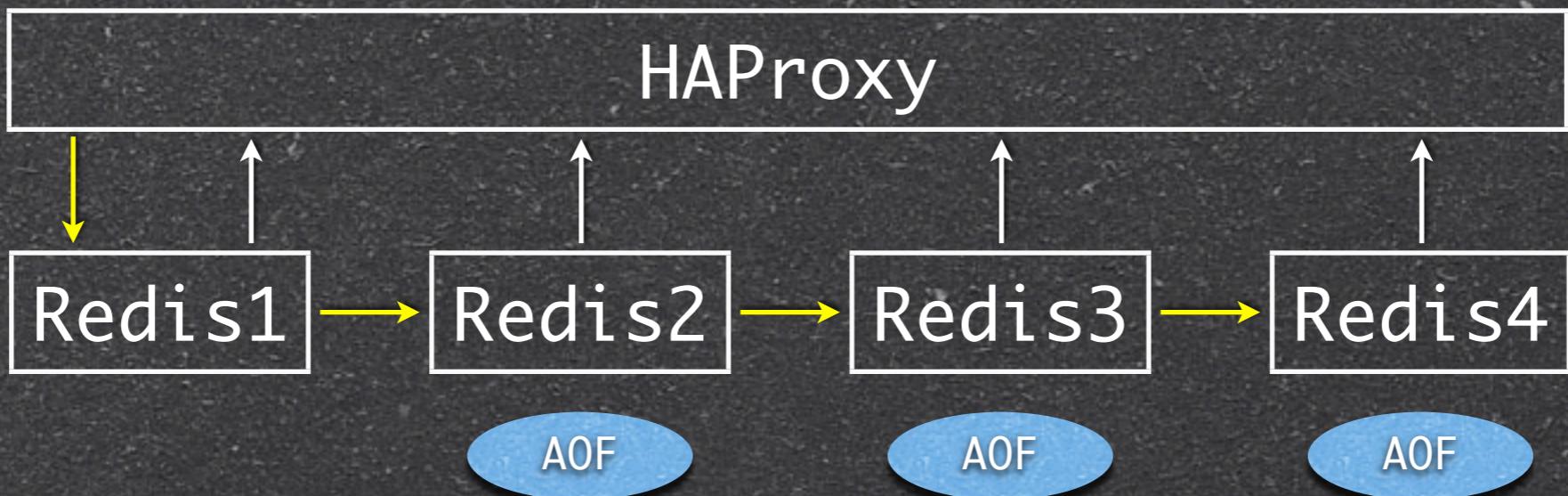


MySQL默认为异步
主从复制

semi-sync半同步
主从复制

Redis

- HAProxy
- 链式Master Slave方案



改进方向

- 自动化、高可用
- 可扩展性

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