## 11. Prove and demonstrate that binary search takes O(log(n)).

Binary Search calls itself/iterates through the while loop a number of times based on what power of 2 is it less than or equal to.

If its  $2^1 = 2$  elements it's found in one loop

If its  $2^2 = 4$  elements it'll take 2 loops to find the individual elements when breaking up by 2

If its  $2^3 = 8$  elements it'll take 3 loops splitting it into two to find each individual

And so on

We're looking for what power of 2, 2<sup>x</sup> can be used to divide N to be less than or equal to 1 Aka

 $1 = n / (2^{x})$ Multiply both sides  $2^{x} = n$ The definition of logarithms would let you do  $\log(Base2)(n) = x$ 

So the amount of times the function loops is based on Log(Base2)(n)

Which falls under the O(log(n)) value