

Graph Reductions with Applications to Robotics

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The Problem

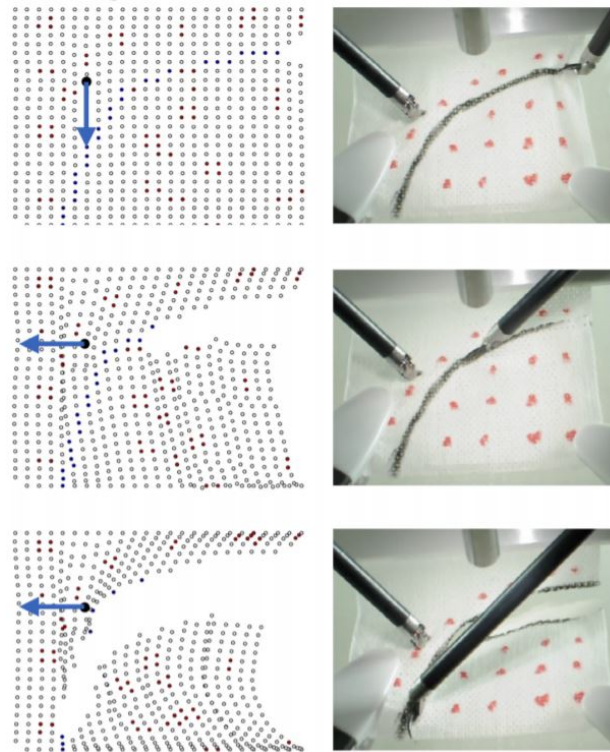


Figure 1: Demonstration of DVRK robot cutting through gauze

The Problem

- Given a piece of gauze, we'd like to cut it in the most stable way possible
- We model the gauze as a grid of points connected by springs
- We measure a cut's stability by the difference in the potential energy in the system before and after the cut

Current Methods

- A deep reinforcement learning implementation for this problem exists, but takes far too long. (5 minutes)
- Goal: come up with an approximation that is quick but does not need to be optimal.

Finding the Potential

“Consider the vector $P = X^T \mathcal{L}X$. $P \in \mathbb{R}^{1 \times 3}$, and has the following form:

$$P = [\sum_{(i,j) \in E} (x(i) - x(j))^2 \quad \sum_{(i,j) \in E} (y(i) - y(j))^2 \quad \sum_{(i,j) \in E} (z(i) - z(j))^2]$$

By Hooke’s Law, the vector P is the energy in the spring system.”

Finding the Potential

- Not too far off: actually, $P = \text{Tr}(X^T \mathcal{L} X)$
- Consider $(X^T D X)_{(0,0)} = X_0^T D X_0 = \sum_{i \in V} d(i) x(i)^2$
- And $(X^T A X)_{(0,0)} = X_0^T A X_0 = \sum_{(i,j) \in E} 2x(i)x(j)$
- $\sum_{i \in V} d(i) x(i)^2 = \sum_{(i,j) \in E} x(i)^2 + x(j)^2$

Finding the Potential

- From physics, steady state minimizes potential energy
- So need X that minimizes $\text{Tr}(X^T \mathcal{L} X)$
- Trivial solution when $X = 0$, so must enforce fixed points
- “The solution to the above optimization is a linear system of equations of the form $B(\mathcal{L} X) = Y$.”
- “Can be seen from the KKT conditions”

KKT conditions

- Have objective function $f(X) = \text{Tr}(X^T \mathcal{L}X)$
- $\nabla f(X) = 2\mathcal{L}X$
- Constraints $g_k(X)$ of the form $X_{(i)} = C_i$
- So $\nabla g_k(X)$ has ones in the i th row, zeros elsewhere
- Solution has $\nabla f(X) - \sum_k \lambda_k \nabla g_k(X) = 0$ and all constraints satisfied
- Can treat each coordinate separately

KKT conditions

- Define B' as B with all zero rows removed
- Define Y' to only contain fixed points
- $B'X = Y'$ enforces all constraints
- $\begin{bmatrix} \mathcal{L} & (-B')^T \end{bmatrix} \begin{bmatrix} X_i \\ \lambda \end{bmatrix} = 0$ enforces
 $\nabla f(X) - \sum_k \lambda_k \nabla g_k(X) = 0$
- Solution: $\begin{bmatrix} \mathcal{L} & (-B')^T \\ B' & 0 \end{bmatrix} \begin{bmatrix} X_i \\ \lambda \end{bmatrix} = \begin{bmatrix} 0 \\ Y'_i \end{bmatrix}$

Finding The Optimal Tension

- Want to maximize $\text{Tr}(X^T \mathcal{L} X) - \text{Tr}(X'^T \mathcal{L}' X')$ over choices of tensioning location
- Use gradient ascent to maximize it
- Gradient is complicated, so use finite difference approximations
- No global maximum—constrain to not move too far from resting position

Finding The Optimal Tension

- This is faster than the RL method, but still slow
- Can speed up if we make the graph smaller, but need to maintain structure

Edge Contraction

- Uses a greedy heuristic to contract edges and reduce the size of the graph.
- When 2 nodes are merged, they become connected to all the neighbors of the original 2 nodes.
- So in order to find a clustering, we continuously contract edges.
- Used the QuickUnion data structure (from 61B!) to keep track of clusters.

Heuristics

- Choose the 2 nodes in the graph with largest (or smallest) degree, find the shortest path between them, and merge them. Continue recursively
 - Works well on barbell graphs and trees.
 - Not so good on a grid graph.
- Choose the node with smallest degree in the graph, merge it with its neighbors
 - Works reasonably well!

Clusters

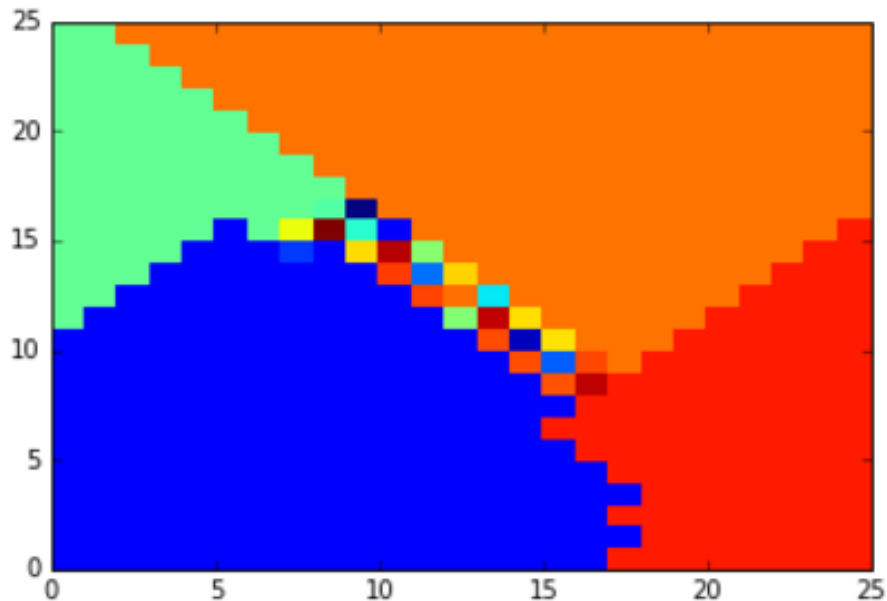


Figure 2: Heuristic: Greedily add smallest degree vertex. $k = 30$

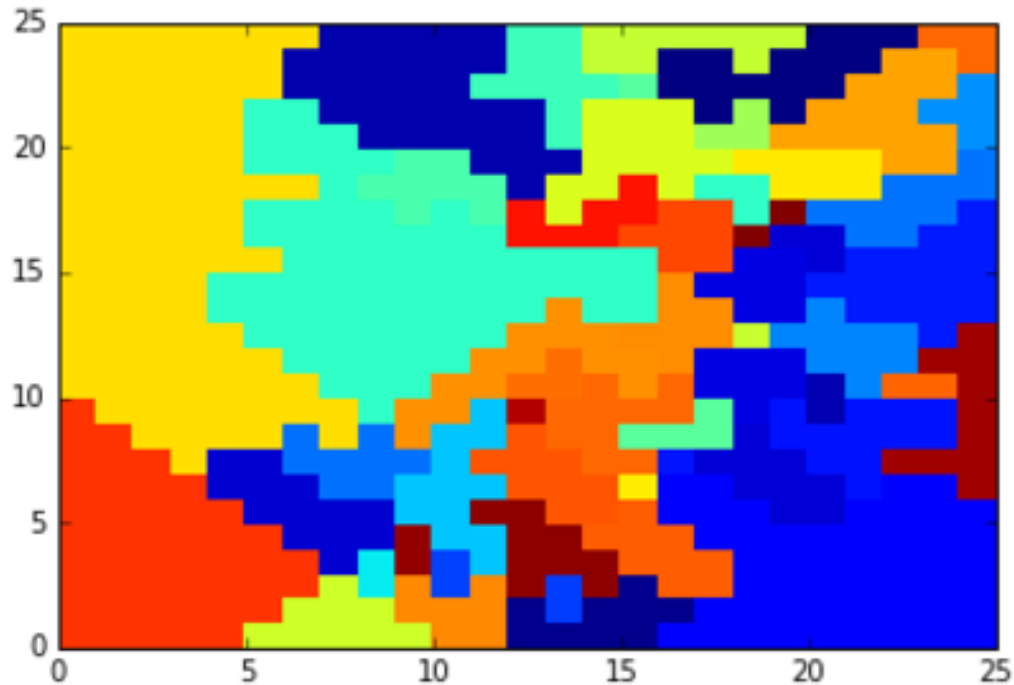


Figure 3: Heuristic: Greedily add smallest degree vertex, prefer to merge with smaller clusters. $k = 20$

Spectral Clustering

- Approximation to sparsest cut problem.
- Given \mathcal{L} of the graph, compute the unit eigenvectors corresponding to the k smallest eigenvalues to minimize Rayleigh quotient.
- Cluster these vectors (with k-means)

Spectral Clustering Visualization

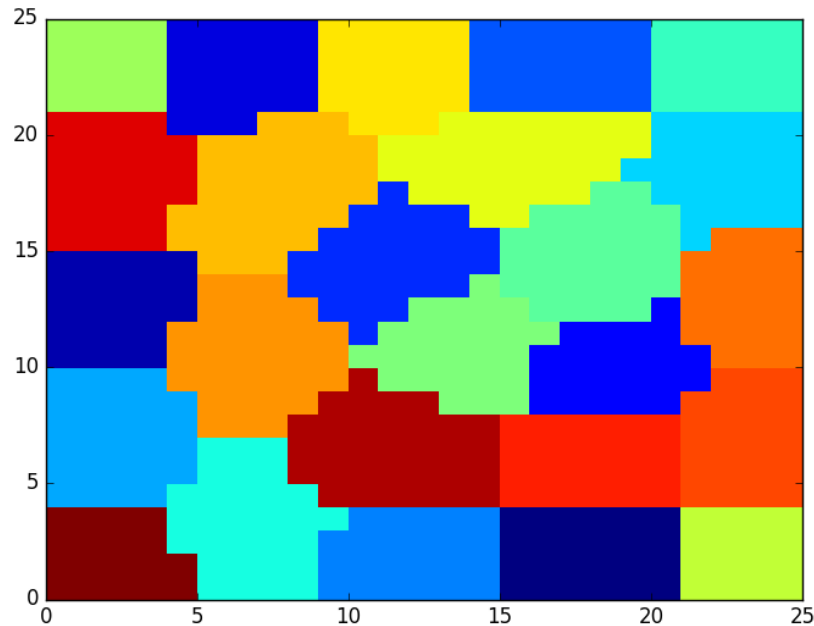


Figure 4: Clustering the grid graph

Spectral Clustering Speed

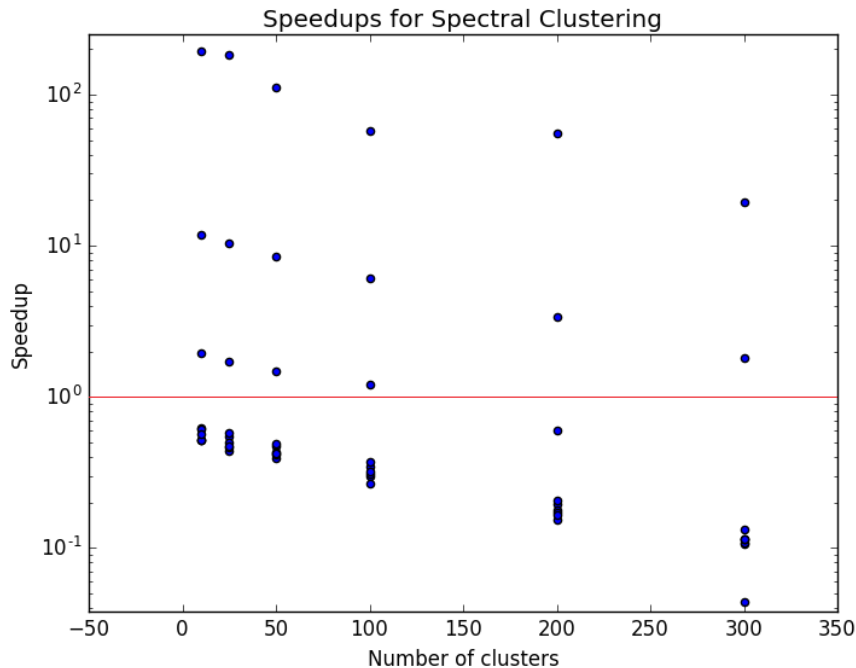


Figure 5: $Speedup = \frac{\text{Optimal tension time}}{\text{Approximate tension time}}$
Bigger is better

Spectral Clustering Performance

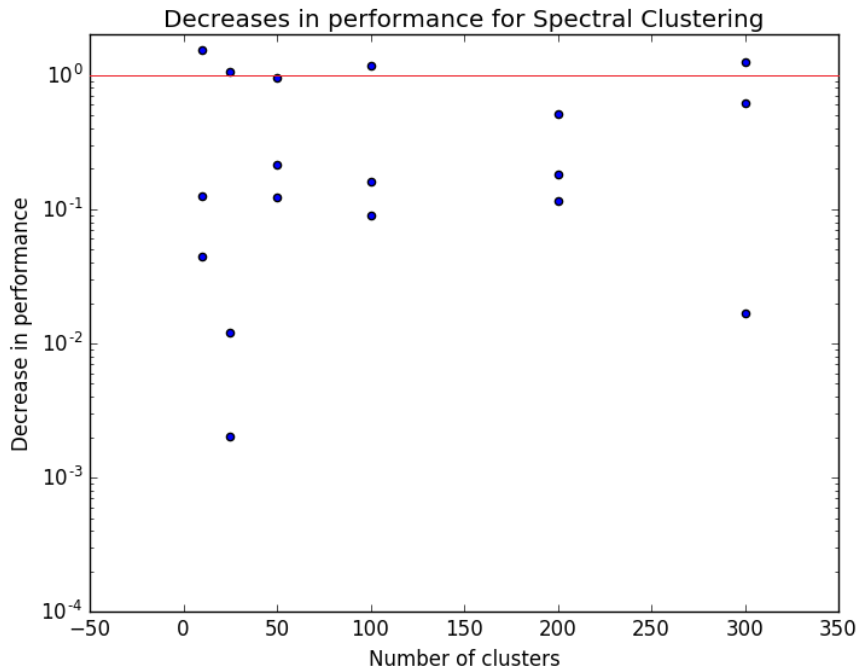


Figure 6: $\text{Performance} = \frac{\text{Optimal tension} - \text{Approximate tension}}{\text{Optimal tension} - \text{No tension}}$
Smaller is better

Our Implementation

`https://github.com/j-m-h/270-project`