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Tugas 4

# Studi Kasus 5: Mencari Pasangan Tititk Terdekat (Closest Pair of Points) Identifikasi Problem:

Diberikan array n poin dalam bidang kartesius, dan problemnya adalah mencari tahu pasangan poin terdekat dalam bidang tersebut dengan merepresentasikannya ke dalam array. Masalah ini muncul di sejumlah aplikasi. Misalnya, dalam kontrol lalu lintas udara, kita mungkin ingin memantau pesawat yang terlalu berdekatan karena ini mungkin menunjukkan kemungkinan tabrakan. Ingat rumus berikut untuk jarak antara dua titik p dan q.

$$||pq|| = \sqrt{(p_x - q_x)^2 + (p_y - q_y)^2}$$

#### Solusi

Solusi umum dari permasalahan tersebut adalah menggunakan algoritma **Brute force** dengan O (n ^ 2), hitung jarak antara setiap pasangan dan kembalikan yang terkecil. Namun, kita dapat menghitung jarak terkecil dalam waktu **O** (nLogn) menggunakan strategi Divide and Conquer.

## Tugas:

- 1) Buatlah program untuk menyelesaikan problem closest pair of points menggunakan algoritma divide & conquer yang diberikan. Gunakan bahasa C++
- 2) Tentukan rekurensi dari algoritma tersebut, dan selesaikan rekurensinya menggunakan metode recursion tree untuk membuktikan bahwa algoritma tersebut memiliki Big-O (n lg n)

## Jawaban Studi Kasus 5

## 1. Program C++: cpp.sh/5eb52

```
: Sitti Ufairoh Azzahra
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    Kelas
    Deskripsi : Program Closest Pair of Points
#include <bits/stdc++.h>
#include <chrono>
using namespace std;
using namespace std::chrono;
class Point {
    public:
    int x, y;
};
int compareX(const void* a, const void* b){
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->x - p2->x);
int compareY(const void* a, const void* b){
    Point *p1 = (Point *)a, *p2 = (Point *)b;
    return (p1->y - p2->y);
float dist(Point p1, Point p2){
    return sqrt( (p1.x - p2.x)*(p1.x - p2.x) + (p1.y - p2.y)*(p1.y - p2.y));
```

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```
float bruteForce(Point P[], int n){
    float min = FLT MAX;
    for (int i = 0; i < n; ++i)
        for (int j = i+1; j < n; ++j)
            if (dist(P[i], P[j]) < min)</pre>
                min = dist(P[i], P[j]);
    return min;
float min(float x, float y){
    return (x < y)? x : y;
float stripClosest(Point strip[], int size, float d){
    float min = d; //Inisiasi jarak minimum = d
    qsort(strip, size, sizeof(Point), compareY);
    for (int i = 0; i < size; ++i)
        for (int j = i+1; j < size && (strip[j].y - strip[i].y) < min; ++j)
            if (dist(strip[i],strip[j]) < min)</pre>
                min = dist(strip[i], strip[j]);
    return min;
float closestUtil(Point P[], int n){
    //Jika ada 2 atau 3 points, gunakan brute force
    if (n <= 3)
        return bruteForce(P, n);
    int mid = n/2;
    Point midPoint = P[mid];
    float dl = closestUtil(P, mid);
    float dr = closestUtil(P + mid, n - mid);
    float d = min(dl, dr);
    Point strip[n];
    int j = 0;
for (int i = 0; i < n; i++)</pre>
        if (abs(P[i].x - midPoint.x) < d)</pre>
            strip[j] = P[i], j++;
    return min(d, stripClosest(strip, j, d) );
float closest(Point P[], int n){
    qsort(P, n, sizeof(Point), compareX);
    return closestUtil(P, n);
```

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```
int main(){
    high_resolution_clock::time_point t1 = high_resolution_clock::now();
    Point P[] = {{2, 3}, {12, 30}, {40, 50}, {5, 1}, {12, 10}, {3, 4}};
    int n = sizeof(P) / sizeof(P[0]);

    cout<<"P[] = {{2, 3}, {12, 30}, {40, 50}, {5, 1}, {12, 10}, {3, 4}};"<<end
l<<endl;
    cout<<"Jarak terkecil = "<<closest(P, n);

    high_resolution_clock::time_point t2 = high_resolution_clock::now();
    auto duration = duration_cast<microseconds>( t2 - t1 ).count();
    cout<<endl<<endl<<<duration</pre>
```

### **Hasil Program:**

```
P[] = {{2, 3}, {12, 30}, {40, 50}, {5, 1}, {12, 10}, {3, 4}};

Jarak terkecil = 1.41421

78 microseconds
```

#### 2. Rekurensi

```
T(n) = 2T(n/2) + O(n) + O(n * log n) + O(n)
T(n) = 2T(n/2) + O(n * log n)
T(n) = T(n * log n * log n)
```

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Tugas 4

# Studi Kasus 6: Algoritma Karatsuba untuk Perkalian Cepat Identifikasi Problem:

Diberikan dua string biner yang mewakili nilai dua bilangan bulat, cari produk (hasil kali) dari dua string. Misalnya, jika string bit pertama adalah "1100" dan string bit kedua adalah "1010", output harus 120. Supaya lebih sederhana, panjang dua string sama dan menjadi n. Tugas:

- 1) Buatlah program untuk menyelesaikan problem fast multiplication menggunakan algoritma divide & conquer yang diberikan (Algoritma Karatsuba). Gunakan bahasa C++
- 2) Rekurensi dari algoritma tersebut adalah T(n) = 3T(n/2) + O(n), dan selesaikan rekurensinya menggunakan metode substitusi untuk membuktikan bahwa algoritma tersebut memiliki Big-O (n lg n)

## Jawaban Studi Kasus 6

# 1. Program C++: cpp.sh/7ivm5y

```
: Sitti Ufairoh Azzahra
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    NPM
    Kelas
    Deskripsi
               : Program Karatsuba
#include<iostream>
#include<chrono>
#include<stdio.h>
using namespace std;
using namespace std::chrono;
int makeEqualLength(string &str1, string &str2){
    int len1 = str1.size();
    int len2 = str2.size();
    if (len1 < len2){
        for (int i = 0; i < len2 - len1; i++)
            str1 = '0' + str1;
        return len2;
    else if (len1 > len2){
        for (int i = 0; i < len1 - len2; i++)
            str2 = '0' + str2;
    return len1; // If len1 >= len2
string addBitStrings( string first, string second ){
    string result;
    int length = makeEqualLength(first, second);
    int carry = 0;
    for (int i = length-1; i >= 0; i--){
        int firstBit = first.at(i) - '0';
        int secondBit = second.at(i) - '0';
```

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```
int sum = (firstBit ^ secondBit ^ carry)+'0';
        result = (char)sum + result;
        carry = (firstBit&secondBit) | (secondBit&carry) | (firstBit&carry);
    }
    if (carry) result = '1' + result;
    return result;
int multiplyiSingleBit(string a, string b){
    return (a[0] - '0')*(b[0] - '0');
long int multiply(string X, string Y){
    int n = makeEqualLength(X, Y);
    if (n == 0) return 0;
    if (n == 1) return multiplyiSingleBit(X, Y);
    int fh = n/2;
    int sh = (n-fh);
    string Xl = X.substr(0, fh);
    string Xr = X.substr(fh, sh);
    string Yl = Y.substr(0, fh);
    string Yr = Y.substr(fh, sh);
    long int P1 = multiply(X1, Y1);
    long int P2 = multiply(Xr, Yr);
    long int P3 = multiply(addBitStrings(X1, Xr), addBitStrings(Y1, Yr));
    return P1*(1<<(2*sh)) + (P3 - P1 - P2)*(1<<sh) + P2;
int main(){
    high_resolution_clock::time_point t1 = high_resolution_clock::now();
    cout<<"String 1: 1100, String 2: 1010"<<endl;</pre>
    cout<<"String 1: 11, String 2: 11"<<endl;</pre>
    cout<<"\nHasil kali: "<<multiply("1100", "1010");</pre>
    cout<<"\nHasil kali: "<<multiply("11", "11");</pre>
    high_resolution_clock::time_point t2 = high_resolution_clock::now();
    auto duration = duration_cast<microseconds>( t2 - t1 ).count();
    cout<<endl<<endl<<duration<<" microseconds" <<endl;</pre>
```

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## Hasil program:

String 1: 1100, String 2: 1010 String 1: 11, String 2: 11

Hasil kali: 120 Hasil kali: 9

105 microseconds

### 2. Rekurensi

- Let's try divide and conquer.
  - Divide each number into two halves.

• 
$$x = x_H r^{n/2} + x_L$$

• 
$$y = y_H r^{n/2} + y_L$$

- Then:

$$\begin{aligned} xy &= \left( x_{H} \, r^{n/2} + x_{L} \right) \, y_{H} \, r^{n/2 \, +} \, y_{L} \\ &= x_{H} y_{H} r^{n} + \left( x_{H} y_{L} + x_{L} y_{H} \right) r^{n/2} + x_{L} y_{L} \end{aligned}$$

- Runtime?
  - T(n) = 4 T(n/2) + O(n)
  - T(n) = O(n^2)
- Instead of 4 subproblems, we only need 3 (with the help of clever insight).
- Three subproblems:

$$-a = x_H y_H$$

$$-d = x_L y_L$$

$$- e = (x_H + x_L) (y_H + y_L) - a - d$$

- Then xy =  $a r^n + e r^{n/2} + d$
- T(n) = 3 T(n/2) + O(n)
- $T(n) = O(n^{\log 3}) = O(n^{1.584...})$

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Tugas 4

## Studi Kasus 7: Permasalahan Tata Letak Keramik Lantai (Tilling Problem) Identifikasi Problem:

Diberikan papan berukutan n x n dimana n adalah dari bentuk 2k dimana k  $\ge$  1 (Pada dasarnya n adalah pangkat dari 2 dengan nilai minimumnya 2). Papan memiliki satu sel yang hilang (ukuran 1 x 1). Isi papan menggunakan ubin berbentuk L. Ubin berbentuk L berukuran 2 x 2 persegi dengan satu sel berukuran 1 x 1 hilang.

#### Tugas:

- 1) Buatlah program untuk menyelesaikan problem tilling menggunakan algoritma divide & conquer yang diberikan. Gunakan bahasa C++
- 2) Relasi rekurensi untuk algoritma rekursif di atas dapat ditulis seperti di bawah ini. C adalah konstanta. T(n) = 4T(n/2) + C. Selesaikan rekurensi tersebut dengan Metode Master

# <u>Jawaban Studi Kasus 7</u>

## 1. Program c++

```
: Sitti Ufairoh Azzahra
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    Kelas
    Deskripsi
               : Program Tilling
#include <iostream>
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
using namespace std;
int board[1000][1000];
int no = 0;
int quadrant = 0;
void trominoTile(int xBoard,
                 int yBoard,
                 int x_hole,
                 int y_hole,
                 int boardSize);
void trominoTile(int xBoard, int yBoard, int x_hole, int y_hole, int boardSize
){
    int halfSize = boardSize / 2,
        xCenter = 0,
        yCenter = 0;
    xCenter = xBoard + halfSize - 1;
    yCenter = yBoard + halfSize - 1;
    if (boardSize == 2){
        if (board[xBoard][yBoard + 1] == 0 && board[xBoard + 1][yBoard] == 0 &
& board[xBoard + 1][yBoard + 1] == 0){
            no++;
```

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```
board[xBoard][yBoard + 1] = no;
            board[xBoard + 1][yBoard] = no;
            board[xBoard + 1][yBoard + 1] = no;
        }
        if (board[xBoard][yBoard] == 0 && board[xBoard + 1][yBoard] == 0 && bo
ard[xBoard + 1][yBoard + 1] == 0){
            no++;
            board[xBoard][yBoard] = no;
            board[xBoard + 1][yBoard] = no;
            board[xBoard + 1][yBoard + 1] = no;
        }
        if (board[xBoard][yBoard + 1] == 0 && board[xBoard][yBoard] == 0 && bo
ard[xBoard + 1][yBoard + 1] == 0){
            no++;
            board[xBoard + 1][yBoard + 1] = no;
            board[xBoard][yBoard + 1] = no;
            board[xBoard][yBoard] = no;
        }
        if (board[xBoard][yBoard + 1] == 0 && board[xBoard + 1][yBoard] == 0 &
& board[xBoard][yBoard] == 0){
            no++;
            board[xBoard][yBoard] = no;
            board[xBoard][yBoard + 1] = no;
            board[xBoard + 1][yBoard] = no;
        }
        return;
    }
    if (x_hole <= xCenter){</pre>
        if (y_hole <= yCenter){</pre>
            if (board[xCenter][yCenter + 1] == 0 && board[xCenter + 1][yCenter
] == 0 && board[xCenter + 1][yCenter + 1] == 0){
                board[xCenter][yCenter + 1] = no;
                board[xCenter + 1][yCenter] = no;
                board[xCenter + 1][yCenter + 1] = no;
                quadrant = 1;
        else{
            if (board[xCenter][yCenter] == 0 && board[xCenter + 1][yCenter] ==
0 && board[xCenter + 1][yCenter + 1] == 0){
                no++;
                board[xCenter][yCenter] = no;
                board[xCenter + 1][yCenter + 1] = no;
                board[xCenter + 1][yCenter] = no;
                quadrant = 2;
            }
        }
```

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```
else{
       if (y_hole <= yCenter){</pre>
           if (board[xCenter][yCenter + 1] == 0 && board[xCenter][yCenter] ==
0 && board[xCenter + 1][yCenter + 1] == 0){
               no++;
               board[xCenter][yCenter] = no;
               board[xCenter][yCenter + 1] = no;
               board[xCenter + 1][yCenter + 1] = no;
               quadrant = 3;
           }
       }
       else{
           if (board[xCenter + 1][yCenter] == 0 && board[xCenter][yCenter] ==
0 && board[xCenter][yCenter + 1] == 0){
               no++;
               board[xCenter][yCenter] = no;
               board[xCenter][yCenter + 1] = no;
               board[xCenter + 1][yCenter] = no;
               quadrant = 4;
       }
   }
   if (quadrant == 1){
       trominoTile(xBoard, yBoard, x_hole, y_hole, halfSize);
       trominoTile(xBoard, yCenter + 1, xCenter, yCenter + 1, halfSize);
       trominoTile(xCenter + 1, yBoard, xCenter + 1, yCenter, halfSize);
       trominoTile(xCenter + 1, yCenter + 1, xCenter + 1, yCenter + 1,
                   halfSize);
   }
   if (quadrant == 2){
       trominoTile(xBoard, yBoard, xCenter, yCenter, halfSize);
       trominoTile(xBoard, yCenter + 1, x_hole, y_hole, halfSize);
       trominoTile(xCenter + 1, yBoard, xCenter + 1, yCenter, halfSize);
       trominoTile(xCenter + 1, yCenter + 1, xCenter + 1, yCenter + 1,
                   halfSize);
   }
   if (quadrant == 3){
       trominoTile(xBoard, yBoard, xCenter, yCenter, halfSize);
       trominoTile(xBoard, yCenter + 1, xCenter, yCenter + 1, halfSize);
       trominoTile(xCenter + 1, yBoard, x_hole, y_hole, halfSize);
       trominoTile(xCenter + 1, yCenter + 1, xCenter + 1, yCenter + 1,
                   halfSize);
   }
   if (quadrant == 4){
       trominoTile(xBoard, yBoard, xCenter, yCenter, halfSize);
       trominoTile(xBoard, yCenter + 1, xCenter, yCenter + 1, halfSize);
       trominoTile(xCenter + 1, yBoard, xCenter + 1, yCenter, halfSize);
       trominoTile(xCenter + 1, yCenter + 1, x hole, y hole, halfSize);
```

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```
int main(){
    int boardSize, x_hole, y_hole;
        printf("\n------
        printf("\nEnter size of board (0 to quit): ");
        scanf("%d", &boardSize);
        if (boardSize){
             printf("\nEnter coordinates of missing hole: ");
             scanf("%d%d", &x_hole, &y_hole);
            for (int i = 1; i \leftarrow pow(2, boardSize); i++){
                 for (int j = 1; j \leftarrow pow(2, boardSize); j++){
                     board[i][j] = 0;
             board[x_hole][y_hole] = -1;
            trominoTile(1, 1, x_hole, y_hole, pow(2, boardSize));
            for (int i = 1; i \le pow(2, boardSize); i++){
                 for (int j = 1; j <= pow(2, boardSize); j++){
   if (i == x_hole && j == y_hole){</pre>
                         board[i][j] == -1;
                          printf("%4s", "X");
                     else{
                          printf("%4d", board[i][j]);
                 }cout << endl;</pre>
        no = 0;
    while (boardSize);
    return EXIT_SUCCESS;
```

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Tugas 4

## 2. Rekurensi

Relasi perulangan untuk algoritma rekursif di atas dapat ditulis seperti di bawah ini. C adalah konstanta.

$$T(n) = 4T(n/2) + C$$

Rekursi di atas dapat diselesaikan dengan menggunakan Metode Master dan kompleksitas waktu adalah O (n2)

$$T(n) = 4T (n 2)$$
  
 $) + C$   
 $a = 4, b = 2, f(n) = 1$   
 $n \log b \ a = n \log 2 \ 4 = n 2$   
 $f(n) = 1 = O(n \log 2 \ 4 - \epsilon)$  untuk  $\epsilon = 1$   
Case 1 applies

Maka, solusinya adalah

$$T(n) = O(n \log b \ a) = O(n \log 2 \ 4) = O(n2)$$