

# CS 70 Discussion 4A

September 25, 2024

## Extended Euclid's Algorithm

Finds values  $x, y \in \mathbb{Z}$  for some given  $a, b \in \mathbb{N}$  such that:

$$ax + by = \gcd(a, b)$$

The efficient algorithm is as follows (return value is of form  $(x, y, \gcd(a, b))$ ):

```
function E_GCD( $a, b$ )  
  if  $b = 0$  then return  $(1, 0, a)$   
  else  
     $(x, y, z) \leftarrow \text{E\_GCD}(b, a \bmod b)$   
    return  $(y, x - \lfloor \frac{a}{b} \rfloor y, z)$   
  end if  
end function
```

You can use EGCD to get inverses iff (if and only if)  $\gcd(a, b) = 1$ :

$$a^{-1} \equiv x \pmod{b}$$

$$b^{-1} \equiv y \pmod{a}$$

# Fermat's Little Theorem

Following is true for any prime  $p$  and  $a \not\equiv 0 \pmod{p}$ :

$$a^{p-1} \equiv 1 \pmod{p}$$

The following general variant is true for any prime  $p$  and *any*  $a \in \mathbb{Z}$ :

$$a^p \equiv a \pmod{p}$$

# Chinese Remainder Theorem

**Problem Setup:** You're given variables  $m_1, m_2, \dots, m_n \in \mathbb{N}^+$  and  $a_1, a_2, \dots, a_n \in \mathbb{Z}$  where  $(\forall i \neq j)(\gcd(m_i, m_j) = 1)$  (the set of  $m_i$ 's is pairwise *coprime*). We want to find a  $x \in \mathbb{Z}$  where:

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

$$\vdots$$

$$x \equiv a_n \pmod{m_n}$$

**Conclusion:** There always exists a *unique* solution  $x \pmod{M}$  where  $M = \prod_{i=1}^n m_i$ .

# Chinese Remainder Theorem (Cont.)

**Solution:** For the system  $(\forall i \in \{1, 2, \dots, n\})(x \equiv a_i \pmod{m_i})$   
where  $(\forall i \neq j)(\gcd(m_i, m_j) = 1)$ :

$$x \equiv \sum_{i=1}^n a_i b_i \pmod{M}$$

$$b_i = \frac{M}{m_i} \left[ \left( \frac{M}{m_i} \right)^{-1} \pmod{m_i} \right]$$

$$M = \prod_{i=1}^n m_i$$