* A Tree is a ‘undirected graph with no cycles’ aka a ‘connected graph with n vertices and n-1 edges’
* A rooted tree is a tree with designated root node where every node either points away (out-tree) or towards this node (in-tree).
* DAG – ‘directed graphs with no cycles’. All out trees are DAGs but not all DAGs are out trees
* A Bipartite graph – whose vertices can be split in to 2 groups U and V so that **every edge** connects between U and V. These graphs are two colourable. Graphs with ‘**no odd length cycle**’
* Complete graph – a graph where there is a ‘unique edge’ between every pair of vertices. i.e. we have edges form every vertex to another vertex. If we need to test code for worst case, test it on complete graph as it has lots of edges.
* Strongly Connected Components (SCC) = Self Contained Cycles: in a ‘directed graph’ where *every* vertex in a given cycle can reach *other every other vertex in same cycle*.
* Bridge (cut) – is any edge with its removal increases the number of connected components.
* Articulation Point (Cut vertex) – any node in the graph whose removal increases number of connected components.
* Minimum Spanning Tree – is a subset of edges of connected & edge weighted graph that ‘still’ connects all vertices together without cycles and with minimum possible edge weight.

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| **Adjacency Matrix:**   1. Edge weight lookup is O (1) 2. Requires O(V2) space though very space efficient for **dense** graphs (lots of edges) 3. Iteration over all edges takes O(V2) time | **Adjacency List:**   1. Edge weight lookup is O(E) 2. Less space efficient for dense graphs, i.e. space efficient for sparse (less edges) graphs. 3. Iteration is efficient. |

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| **Shortest Path:**   * BFS (unweighted) * Dijkstra’s * Bellman Ford * Floyd Warshall | **Connectivity:**   * DFS * Union and Find | **Negative Cycles:**   * Bellman Ford * Floyd Warshall | **Strongly Conn Comp**   * Trajan’s * Kosaraju | **Travelling Salesman**   * Held-Karp * Branch & Bound * Ant Colony Optimisa. |
| **Min Spanning Tree:**   * Kruskal * Prims * Boruvka’s | **Network Flow:**   * Ford Fulkerson * Edmonds-Karp * Dinic’s |  |  |  |

**DFS:**

Time complexity: O(V+E)

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| **private** **void** dfs(**int** at, **boolean**[] visited) {  visited[at] = **true**;  **for** (**int** adjVert : adjList.get(at)) {  **if** (visited[adjVert] == **false**) {  dfs(adjVert, visited);  }  }  } |  |
| **Connected Components by DFS:**  **void** connectedComponents() {  **boolean**[] visited = **new** **boolean**[n];  **int** color = 1;  Map<Integer, List<Integer>> map = **new** HashMap<>();  **for** (**int** i = 0; i < n; i++) {  **if** (visited[i] == **false**) {  dfs(i, visited, color++, map);  }  }  }  **private** **void** dfs(**int** at, **boolean**[] visited, **int** color, Map<Integer, List<Integer>> map) {  visited[at] = **true**;  **if** (map.get(color) == **null**) {  List<Integer> l = **new** ArrayList<>();  l.add(at);  map.put(color, l);  } **else** {  List<Integer> l = map.get(color);  l.add(at);  map.put(color, l);  }  **for** (**int** adjVert : adjList.get(at)) {  **if** (visited[adjVert] == **false**) {  dfs(adjVert, visited, color, map);  }  }  } | |

**BFS:**

Time complexity: O(V+E)

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| **void** bfs(**int** u, **int** v) {  **boolean**[] visited = **new** **boolean**[n];  **int**[] prev = **new** **int**[n];  Arrays.*fill*(prev, -1);  Queue<Integer> q = **new** ArrayDeque<>();  q.add(u);  visited[u] = **true**  **while** (!q.isEmpty()) {  **int** currVert = q.poll();  **for** (**int** adjVert : adjList.get(currVert)) {  **if** (visited[adjVert] == **false**) {  q.add(adjVert);  visited[adjVert] = **true**;  prev[adjVert] = currVert;  }  }  }  } |