The Art of Presenting Science

Sivaram Ambikasaran

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Greatest communicators

Who?

Bullshit rule: 1 Work will speak for itself

On August 30, 2012 Mochizuki released four preprints, whose total size was about 500 pages, that develop inter-universal Teichmüller theory and apply it to attempt to prove several very famous problems in Diophantine geometry.^[8] These include the strong Szpiro conjecture, the hyperbolic Vojta conjecture and the abc conjecture over every number field. The preprints have not been published. In September 2018, Mochizuki posted a report on his work by Peter Scholze and Jakob Stix asserting that the third preprint contains an irreparable flaw; he also posted several documents containing his rebuttal of their criticism.^[9] The maiority of number theorists have found Mochizuki's preprints very difficult to follow and have not accepted the conjectures as settled, although there are a few prominent exceptions, including Go Yamashita, Ivan Fesenko, and Yuichiro Hoshi, who vouch for the work and have written expositions of the theory.[10][11]

Reality rule: 1

Your work: your product

Talk: advertisement of your work

Good product with bad adverstisement ©

Bad product with good adverstisement ©

Bullshit rule: 2
Every slide must be packed from top left to bottom right with text

More in paper

Specific problem To determine the safety of infining the lambda term law at the call site I(f ...)], we need to know that for every environment a in which this call is evaluated, that $\rho[f] = (\text{faw}, \rho')$ and $\rho(v) = \rho'(v)$ for each free variable vin the term fees.2

$$\eta(b) = \hat{b} \text{ iff } \eta(g(b)) = \hat{g}(\hat{b}).$$

$$\hat{\beta}(e_i) \in \widehat{Bind}_1$$
 $\hat{b}_i \in \widehat{Bind}_1$
 $\hat{\beta}(e_i) \equiv' \hat{b}_i$,

Theorem 4. It is safe to remeterialize the expression e' in place of the expression e in the call site call iff for every reachable compound abstract state of the form $((cull, \hat{\beta}^a, \hat{vc}, \hat{l}), \equiv)$, it is the case that $\hat{E}(e', \hat{\beta}^a, \hat{vc}) = (lam', \hat{\beta}')$ and $\hat{\mathcal{E}}(e, \hat{\beta}^o, \hat{s}e) = (lsm, \hat{\beta})$ and the relation $\sigma \subseteq Var \times Var$ is a substitution that unifirst the free variables of law with lam and for each $(v', v) \in \sigma$, $\hat{\beta}'(v') \equiv \hat{\beta}(v)$.

$$\alpha^{*}(\operatorname{call}, \beta, w, t) = (\alpha^{*}(V), \alpha^{*}(\beta), \alpha^{*}(\operatorname{cw}), \psi(t))$$
 $\zeta \in \widehat{\Sigma} = \mathbb{C} \operatorname{sill} \times \widehat{BEav} \times \widehat{VEnv} \times \widehat{Time}$
 $\alpha^{*}_{BEnv}(\mathcal{G}) = \lambda \operatorname{tr} \operatorname{Gl}(\sigma)$ $\beta \in \widehat{BEav} = Var \to \widehat{Bind}$
 $\alpha^{*}_{VEnv}(\operatorname{cv}) = \lambda \operatorname{tr} \operatorname{Gl}(\sigma)$ $\psi_{v, t+1}$ $\psi_{v} \in \widehat{VEnv} = \widehat{Bind} \to \widehat{D}$
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 $q_{\mu}^{-1}(b) = b$ $(([(f e_1 ... e_n)^\ell], \hat{\beta}, \widehat{ve}, \hat{t}), \equiv) \rightarrow ((call, \hat{\beta}^{\prime\prime}, \widehat{ve}^\prime, \hat{t}^\prime), \equiv'), \text{ where:}$ $g_B^{-1}(g(b)) = \begin{cases} b & \eta(b) = \eta(b') \text{ for some } g(b') \in B \\ g(b) & \text{otherwise} \end{cases}$ $q_n^{-1}(lam, \beta) = (lam, q_n^{-1}(\beta))$ Theorem 3. Given a compound abstract state $((call, \hat{\beta}, \hat{ve}, \hat{t}), \equiv)$ and two abstruct bindings, \hat{b} and \hat{b}' , if $\alpha^{iq}(call, \beta, ve, t) \subseteq ((call, \hat{\beta}, \hat{ve}, \hat{t}), \equiv)$ and $\eta(b) = \hat{b}$ $g_D^{-1}(\beta) = \lambda v. g_D^{-1}(\beta(v))$ $g_B^{-1}(ve) = \lambda b. g_B^{-1}(ve(b)).$

and $\eta(b') = \hat{b}'$ and $\hat{b} \equiv \hat{b}'$, then ve(b) = ve(b').

Theorem 2. If $\alpha^{\eta}(\beta_1) = \hat{\beta}_1$ and $\alpha^{\eta}(\beta_2) = \hat{\beta}_2$, and $\hat{\beta}_1(v) = \hat{\beta}_2(v)$ and $\hat{\beta}_1(v) \in$ \widehat{Bind}_1 , then $\beta_1(v) = \beta_2(v)$.

Theorem 1. If
$$\alpha^{o}(\zeta) \subseteq \zeta$$
 and $\zeta \mapsto \zeta'$, then there exists a state ζ' such that
$$\begin{cases} -\zeta = \zeta' \text{ and } \alpha^{o}(\zeta) \subseteq \zeta \text{ and } \zeta \mapsto \zeta', \text{ then there exists a state } \zeta' \text{ such that} \\ -\zeta = \zeta' \text{ and } \alpha^{o}(\zeta') \subseteq \zeta'. \end{cases}$$

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 $([(f e_1 ... e_n)^\ell], \beta, ve, t) \Rightarrow (call, \beta'', ve', t'), where:$ $d_i = \mathcal{E}(e_i, \beta, \pi e)$ $d_0 = (\lceil (\lambda^{\ell'} \mid (v_1 \dots v_n) \mid cull) \rceil, \beta')$ t' = tick(call, t) $b_r = alloc(v_r, t')$ $B = \{b_i : b_i \in Bind_1\}$ $\beta'' = (q_D^{-1}\beta')|_{V_i} \mapsto f_i|_{i=1}^{n}$ $ve' = (g_R^{-1}ve)[b_i \mapsto (g_R^{-1}d_i)],$

 $\hat{b} \in \widehat{Bind}$ is a finite set of bindings

 $\hat{t} \in \widehat{T_{imc}}$ is a finite set of times

 $\hat{d}_i = \hat{\mathcal{E}}(e_i, \hat{\beta}, \hat{\psi}_e)$ $\hat{d}_0 \ni ([((\lambda^{\ell'} (v_1 ... v_n) coll)], \hat{\beta}'))$

 $\hat{t}' = \widehat{tick}(call, \hat{t})$

 $\hat{b}_i = \widehat{alloc}(v_i, \hat{t}')$ $\hat{B} = \{\hat{b}_i : \hat{b}_i \in \widehat{Bind}_1\}$ $\hat{\beta}^{\prime\prime} = (\hat{g}_{\hat{\alpha}}^{-1}\hat{\beta}^{\prime})[v_i \mapsto \hat{b}_i]$

 $\widehat{ve'} = (\widehat{g}_{\widehat{a}}^{-1}\widehat{ve}) \sqcup [\widehat{b}_i \mapsto (\widehat{g}_{\widehat{a}}^{-1}\widehat{d}_i)],$

 $\zeta \sim \zeta'$ and $\alpha^{q}(\zeta') \sqsubseteq \zeta'$

 $\tilde{m}' = (\tilde{g}_{\tilde{m}}^{-1}\tilde{m}_{\tilde{e}}) \sqcup [\tilde{b}_{i} \mapsto (\tilde{g}_{\tilde{m}}^{-1}\tilde{d}_{i})],$

Reality rule: 2 Minimal text; Each slide should showcase only one thing **YOU** need to talk about the content

Bullshit rule: 3
Will convince my audience,
I have done LOTS of work

Reality rule: 3
Quantity of work doesn't matter
Quality does

Bullshit rule: 4 Pictures and Figures are pointless

Reality rule: 4 Picture is worth a thousand words

Even in mathematical communication

Bullshit rule: 5
Present every single step in proof so that people will know I have done it correctly

Reality rule: 5
Writing an article is different from presenting your work

Reality rule: 5 Writing an article is different from presenting your work Encapsulate the ideas of the theorems and proofs

Reality rule: 5 Writing an article is different from presenting your work Encapsulate the ideas of the theorems and proofs True understanding

Bullshit rule: 6
Don't talk about failures or incorrect attempts

Reality rule: 6
The path you took is more important than result

Bullshit rule: 7
Be serious;
No jokes, no deviations;
Only content

Reality rule: 7
Your audience are not ROBOTS

Bullshit rule: 8 Same content works for all audience

Reality rule: 8 Audience are our customers

Bullshit rule: 9 No need to motivate the talk

Reality rule: 9
Spend almost a quarter of the talk
on motivation

Bullshit rule: 10
More slides \Longrightarrow More work

Reality rule: 10

If I Had More Time, I Would Have
Written a Shorter Letter.

Bullshit rule: 11

Talk fast \Longrightarrow More words \Longrightarrow More work

Reality rule: 11
Talk short sentences at the right pace.

Bullshit rule: 12
Presentations should be monologue.

Reality rule: 12 Engage your audience.

Have a catchy title

Do not have an outline slide It is boring; Suspense is lost

Start with a question or an interesting fact

Start with a question or an interesting fact
Use that as a motivation

Tell a story
Build up the suspense

Do not memorize your talk!

Do not memorize your talk! Try to adapt on the fly!

Do not memorize your talk!

Try to adapt on the fly!

Extempore!

You are the hero(ine)!

You are the hero(ine)!

Don't become a villain!

Show current slide/total slides

Show current slide/total slides Have a good climax

Show current slide/total slides

Have a good climax

Finish on time with a bang!