

1) Prove  $f_1(n) = \log(n^3)$   $f_2(n) = (\log n)^3$   $f_1(n) = o(f_2(n))$

$$\lim_{n \rightarrow \infty} \frac{\log(n^3)}{(\log n)^3} = \lim_{n \rightarrow \infty} \frac{3 \log n}{(\log n)^3} = \lim_{n \rightarrow \infty} \frac{3 \ln n / \ln 2}{(\ln n)^3 / (\ln 2)^3} \xrightarrow{\text{L'Hôpital}} \lim_{n \rightarrow \infty} \frac{3/n}{3(\ln n)^2 / (\ln 2)^2 \cdot \frac{1}{n}} = 0$$

$$\therefore \log(n^3) = o((\log n)^3)$$

2)  $f_2(n) = (\log n)^3$ ,  $f_3(n) = 2^{\sqrt{\log n}}$

Prove  $f_2(n) = o(f_3(n))$  L'Hôpital

let  $x = \log n$

$$\lim_{n \rightarrow \infty} \frac{(\log n)^3}{2^{\sqrt{\log n}}} = \lim_{x \rightarrow \infty} \frac{x^3}{2^{\sqrt{x}}} \xrightarrow{\text{L'Hôpital}} \lim_{x \rightarrow \infty} \frac{3x^2}{2^{\sqrt{x}} \cdot \ln 2 \cdot x^{-\frac{1}{2}} \cdot \frac{1}{2}} = \lim_{x \rightarrow \infty} \frac{3 \cdot \frac{5}{2} x^{\frac{3}{2}}}{2^{\sqrt{x}} (\ln 2 / 2)^2 \cdot x^{-\frac{1}{2}}}$$

After 4 more times of apply L'Hôpital's rule,

$$= \lim_{x \rightarrow \infty} \frac{3 \cdot \frac{5}{2} \cdot 2 \cdot \frac{3}{2} \cdots \frac{1}{2} x}{2^{\sqrt{x}} (\ln 2 / 2)^6} = 0$$

$$\therefore (\log n)^3 = o(2^{\sqrt{\log n}})$$

$$3) f_3(n) = 2^{\sqrt{\log n}}, f_4(n) = 2^{\ln n}$$

Prove  $f_3(n) = o(f_4(n))$

$$\log n = \ln n / \ln 2$$

$$\text{Let } x = \ln n$$

$$\lim_{n \rightarrow \infty} \frac{2^{\sqrt{\log n}}}{2^{\ln n}} = \lim_{x \rightarrow \infty} \frac{2^{\sqrt{\ln 2} \cdot \sqrt{x}}}{2^x} = \lim_{x \rightarrow \infty} 2^{\sqrt{\ln 2} \cdot \sqrt{x} - x} = 0$$

$$\lim_{n \rightarrow \infty} \text{Thus, } f_3(n) = o(f_4(n))$$

$$4) f_4(n) = 2^{\ln n}, f_5(n) = \binom{n}{3} = \frac{n(n-1)(n-2)}{6}$$

Prove  $f_4(n) = o(f_5(n))$

$$2^{\ln n} = 2^{\log_2 n / \log_2 e} = n^{\frac{1}{\log_2 e}}$$

$$\therefore \frac{1}{\log_2 e} < 1 \quad \therefore \forall n > 1, f_4(n) < n \Rightarrow f_4(n) = o(n)$$

$$\lim_{n \rightarrow \infty} \frac{n}{f_5(n)} = \lim_{n \rightarrow \infty} \frac{n}{\frac{n(n-1)(n-2)}{6}} \xrightarrow{\text{L'Hôpital}} \lim_{n \rightarrow \infty} \frac{1}{\frac{3n^2 - 6n - 2}{6}}$$

$$\therefore n = o(f_5(n)) \quad = 0.$$

$$\text{Thus, } f_4(n) = o(f_5(n))$$

$$5). f_5(n) = \binom{n}{3} \quad f_6(n) = 5n^3$$

Prove  $f_5(n) = \theta(f_6(n))$

$$\lim_{n \rightarrow \infty} \frac{n(n-1)(n-2)/3}{5n^3}$$

$$= \lim_{n \rightarrow \infty} \frac{1}{5 \cdot 3 \cdot 2} = \frac{1}{30}$$

Apply  
L'Hôpital's rule three times

Thus,  $f_5(n) = \theta(f_6(n))$

$$6) f_6(n) = 5n^3 \quad f_7(n) = (\log n)^n$$

$$\forall n > 3, \therefore \log n > \log 3 > 1$$

$$\therefore (\log n)^n > (\log 3)^n$$

$$\text{Thus } f_7(n) = \omega((\log 3)^n)$$

$$\lim_{n \rightarrow \infty} \frac{(\log 3)^n}{5n^3} = \frac{(\log 3)^n \cdot [\ln(\log 3)]^3}{5 \cdot 3 \cdot 2} = 0$$

L'Hôpital's Rule  
for three times

$$\therefore 5n^3 = o((\log 3)^n) \text{ while } (\log 3)^n = o((\log n)^n)$$

Using the transitivity, we can know

$$f_6(n) = o(f_7(n))$$

$$7) f_7(n) = (\log n)^n \quad f_8(n) = n!$$

$$\text{Prove } f_7(n) = o(f_8(n))$$

$$n! = n(n-1)(n-2) \dots (n/2) \dots 1$$

$$\geq n(n-1)(n-2) \dots (n/2) \geq (n/2)^{n/2} \therefore (n/2)^{n/2} = O(n!)$$

$$\lim_{n \rightarrow \infty} \frac{(\log n)^n}{(n/2)^{n/2}} = \lim_{n \rightarrow \infty} \frac{[(\log n)^2]^{n/2}}{(n/2)^{n/2}} = \lim_{n \rightarrow \infty} \left( \frac{\log^2 n}{n/2} \right)^{n/2}$$

$$\text{We can find } \lim_{n \rightarrow \infty} \frac{\log^2 n}{n/2} = \lim_{n \rightarrow \infty} \frac{2 \log n}{n/2} = \lim_{n \rightarrow \infty} \frac{2}{n/2} = 0$$

$$\text{Thus, } \lim_{n \rightarrow \infty} \left( \frac{\log^2 n}{n/2} \right)^{n/2} = 0.$$

$$\therefore (\log n)^n = o((n/2)^{n/2}) \quad (n/2)^{n/2} = O(n!)$$

$$\therefore (\log n)^n = o(n!)$$

In Conclusion:

$$f_1(n) = \log(n^3)$$

$$f_2(n) = (\log n)^3$$

$$f_3(n) = 2^{\sqrt{\log n}}$$

$$f_4(n) = 2^{\ln n}$$

$$f_5(n) = \binom{n}{3}$$

$$f_6(n) = 5n^3$$

$$f_7(n) = (\log n)^n$$

$$f_8(n) = n!$$

Problem 2.

Part 1 show  $T(n) = \Omega(n)$   $T(n) = T(n/2) + T(n/4) + cn$

Base case: ~~Assume~~  $T(n) \geq \frac{1}{2}n$

when  $n=1$   $T(1) = 1 \geq \frac{1}{2}$

when  $n=2$   $T(2) = 2 > 1$

when  $n=4$   $T(4) = T(2) + T(1) + c = 3 + c > 2$  ✓

Inductive case: Assume  $\forall k \in [1, n]$   $k$  is power of 2  $T(k) \geq \frac{1}{2}k$

$$T(2n) = T(n) + T(n/2) + 2cn$$

$$\geq \frac{n}{2} + \frac{n}{4} + 2cn \geq n \quad (c \geq 1)$$

Thus,  $\exists c_0 = \frac{1}{2}$ , st.  $\forall n > 0$ ,  $n$  a power of 2,  $T(n) \geq c_0 n \Rightarrow T(n) = \Omega(n)$

Part 2 show  $T(n) = O(n)$

$k$  is power of 2.

~~Assume~~ Show  $T(k) \leq 10ck$ , where  $c$  is the given constant  $c \geq 1$

Base case. when  $n=1$   $T(1) = 1 \leq 10c$

$n=2$   $T(2) = 2 \leq 20c$

$n=4$   $T(4) = T(2) + T(1) + cn = 3 + c \leq 40c$

Inductive case: Assume  $\forall k \in [1, n]$   $T(k) \leq 10ck$  where  $k \in [1, n]$  power of 2.

$$T(2n) = T(n) + T(n/2) + 2cn$$

$$= 10cn + 5cn + 2cn \leq 20cn$$

Thus,  $\exists c_0 = 8c$ , st.  $\forall n > 0$ ,  $n$  a power of 2,  $T(n) \leq c_0 n \Rightarrow T(n) = O(n)$

Based on Part 1 and Part 2,  $T(n) = \Theta(n)$ .

### Problem 3

a.  $4^n - 3^n$

If we can choose four letters per spot, there are  $4^n$  choices. If there's no 'C' as option, there are only three options per spot,  $3^n$ .

b.  $3^{n-30} \binom{n}{30}$

Imagine there's  $n$  spots, we have  $\binom{n}{30}$  ways to select 30 spots. For the remaining  $n-30$  spots, each have 3 choices. Thus,  $3^{n-30} \binom{n}{30}$  ways in total.

c)  $H(1) = 4, H(2) = 16.$

Consider a sequence of  $n$  length.

$H(n) = 3H(n-1) + 3H(n-2)$

If the last term is not C, there are  $T(n-1)$  options for the first  $n-1$  terms, and 3 options for the last one.

If the last term is C, the second last term can not be C, there are three options.

### Problem 4

a)  $\frac{(2n)!}{2n} = (n-1)!$

For the first  $n-2$  terms,  $T(n-2)$  options are there. Two cases are mutually exclusive, we can just add the product of them altogether.

If it's to arrange people on bench, there are  $(2n)!$  ways to arrange.

Yet, there's actually  $2n$  times fewer the arranging it on bench because the table can be rotated.

b)  $2^n \frac{n!}{n} = 2^n (n-1)!$

If we  $n$  couples as  $n$  units, following the discussion of a), there are  $\frac{n!}{n}$  ways to do that. Yet, for each couple, they can switch with themselves which give rise to  $2^n$  more times arrangement. Thus, in total  $2^n (n-1)!$

### Problem 5. Consider $n$ side polygon, there will be $\alpha$ intersection,

for each distinct group of 4 points. Because we don't consider the case in which more than 2 chords intersect at one point, there are  $\binom{n}{4}$  ways to choose groups of 4 points, thus  $\binom{n}{4}$  intersections.