Analysis Report

v1

Duration	2.65266 ms (2,652,656 ns)
Grid Size	[1024,1,1]
Block Size	[1024,1,1]
Registers/Thread	16
Shared Memory/Block	0 B
Shared Memory Requested	64 KiB
Shared Memory Executed	64 KiB
Shared Memory Bank Size	4 B

[0] Quadro K620

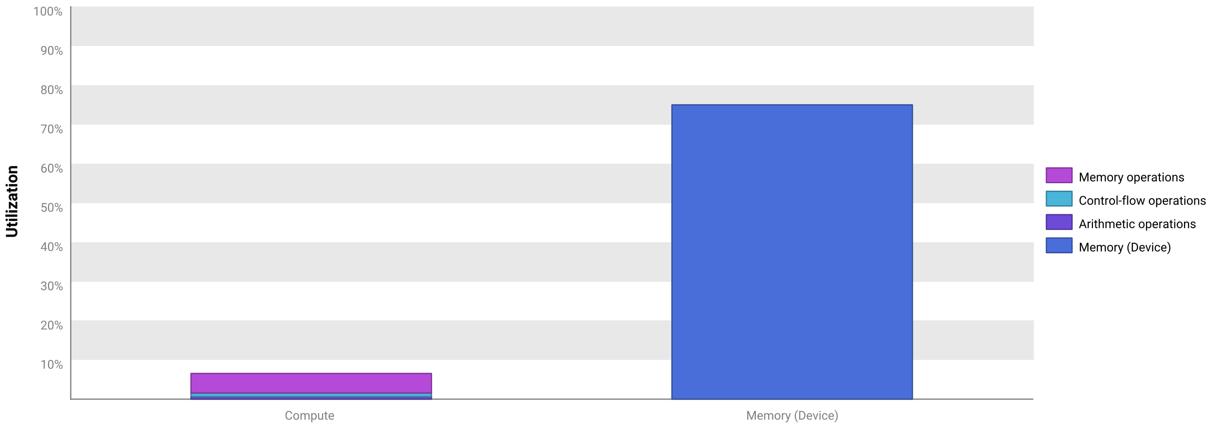
GPU UUID	GPU-1ef8bc5c-b0a8-9167-636d-16994a50ac47
Compute Capability	5.0
Max. Threads per Block	1024
Max. Threads per Multiprocessor	2048
Max. Shared Memory per Block	48 KiB
Max. Shared Memory per Multiprocessor	64 KiB
Max. Registers per Block	65536
Max. Registers per Multiprocessor	65536
Max. Grid Dimensions	[2147483647, 65535, 65535]
Max. Block Dimensions	[1024, 1024, 64]
Max. Warps per Multiprocessor	64
Max. Blocks per Multiprocessor	32
Single Precision FLOP/s	863.232 GigaFLOP/s
Double Precision FLOP/s	26.976 GigaFLOP/s
Number of Multiprocessors	3
Multiprocessor Clock Rate	1.124 GHz
Concurrent Kernel	true
Max IPC	6
Threads per Warp	32
Global Memory Bandwidth	28.8 GB/s
Global Memory Size	1.946 GiB
Constant Memory Size	64 KiB
L2 Cache Size	2 MiB
Memcpy Engines	1
PCIe Generation	2
PCIe Link Rate	5 Gbit/s
PCIe Link Width	16

1. Compute, Bandwidth, or Latency Bound

The first step in analyzing an individual kernel is to determine if the performance of the kernel is bounded by computation, memory bandwidth, or instruction/memory latency. The results below indicate that the performance of kernel "v1" is most likely limited by memory bandwidth. You should first examine the information in the "Memory Bandwidth" section to determine how it is limiting performance.

1.1. Kernel Performance Is Bound By Memory Bandwidth

For device "Quadro K620" the kernel's compute utilization is significantly lower than its memory utilization. These utilization levels indicate that the performance of the kernel is most likely being limited by the memory system. For this kernel the limiting factor in the memory system is the bandwidth of the Device memory.



2. Memory Bandwidth

Memory bandwidth limits the performance of a kernel when one or more memories in the GPU cannot provide data at the rate requested by the kernel. The results below indicate that the kernel is limited by the bandwidth available to the device memory.

2.1. Global Memory Alignment and Access Pattern

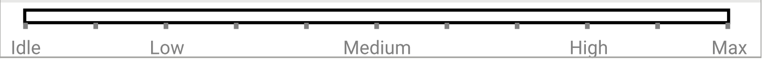
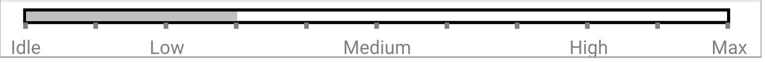
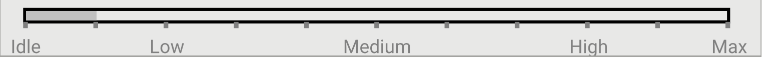
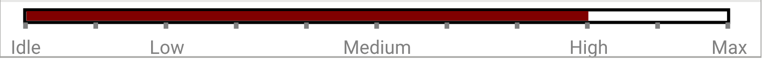
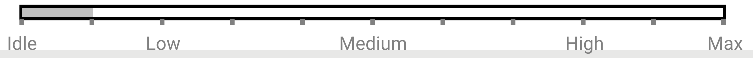
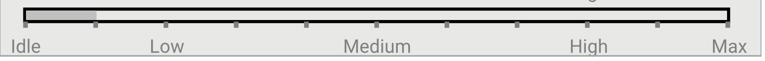
Memory bandwidth is used most efficiently when each global memory load and store has proper alignment and access pattern. The analysis is per assembly instruction.

Optimization: Each entry below points to a global load or store within the kernel with an inefficient alignment or access pattern. For each load or store improve the alignment and access pattern of the memory access.

2.2. GPU Utilization Is Limited By Memory Bandwidth

The following table shows the memory bandwidth used by this kernel for the various types of memory on the device. The table also shows the utilization of each memory type relative to the maximum throughput supported by the memory. The results show that the kernel's performance is potentially limited by the bandwidth available from one or more of the memories on the device.

Optimization: Try the following optimizations for the memory with high bandwidth utilization.
Shared Memory - If possible use 64-bit accesses to shared memory and 8-byte bank mode to achieved 2x throughput.
L2 Cache - Align and block kernel data to maximize L2 cache efficiency.
Unified Cache - Reallocate texture data to shared or global memory. Resolve alignment and access pattern issues for global loads and stores.
Device Memory - Resolve alignment and access pattern issues for global loads and stores.
System Memory (via PCIe) - Make sure performance critical data is placed in device or shared memory.

Transactions	Bandwidth	Utilization	
Shared Memory			
Shared Loads	0	0 B/s	
Shared Stores	0	0 B/s	
Shared Total	0	0 B/s	
L2 Cache			
Reads	7960240	101.77 GB/s	
Writes	3932166	50.272 GB/s	
Total	11892406	152.042 GB/s	
Unified Cache			
Local Loads	0	0 B/s	
Local Stores	0	0 B/s	
Global Loads	4063232	50.691 GB/s	
Global Stores	0	0 B/s	
Texture Reads	786432	10.054 GB/s	
Unified Total	4849664	60.745 GB/s	
Device Memory			
Reads	813358	10.399 GB/s	
Writes	786458	10.055 GB/s	
Total	1599816	20.453 GB/s	
System Memory [PCIe configuration: Gen2 x16, 5 Gbit/s]			
Reads	16	204.556 kB/s	
Writes	5	63.923 kB/s	

3. Instruction and Memory Latency

Instruction and memory latency limit the performance of a kernel when the GPU does not have enough work to keep busy. The results below indicate that the GPU does not have enough work because instruction execution is stalling excessively.

3.1. Instruction Latencies May Be Limiting Performance

Instruction stall reasons indicate the condition that prevents warps from executing on any given cycle. The following chart shows the break-down of stalls reasons averaged over the entire execution of the kernel. The kernel has good theoretical and achieved occupancy indicating that there are likely sufficient warps executing on each SM. Since occupancy is not an issue it is likely that performance is limited by the instruction stall reasons described below.

Constant - A constant load is blocked due to a miss in the constants cache.

Instruction Fetch - The next assembly instruction has not yet been fetched.

Memory Dependency - A load/store cannot be made because the required resources are not available or are fully utilized, or too many requests of a given type are outstanding. Data request stalls can potentially be reduced by optimizing memory alignment and access patterns.

Not Selected - Warp was ready to issue, but some other warp issued instead. You may be able to sacrifice occupancy without impacting latency hiding and doing so may help improve cache hit rates.

Execution Dependency - An input required by the instruction is not yet available. Execution dependency stalls can potentially be reduced by increasing instruction-level parallelism.

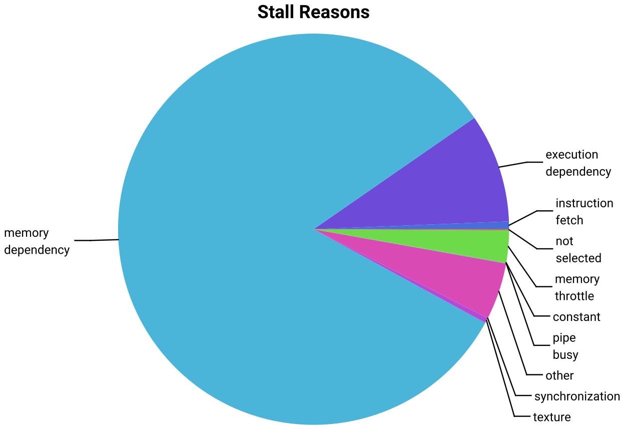
Pipeline Busy - The compute resource(s) required by the instruction is not yet available.

Synchronization - The warp is blocked at a __syncthreads() call.

Memory Throttle - Large number of pending memory operations prevent further forward progress. These can be reduced by combining several memory transactions into one.

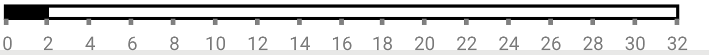
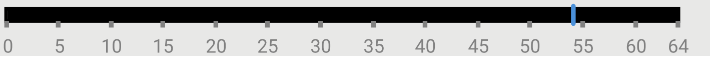
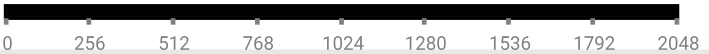
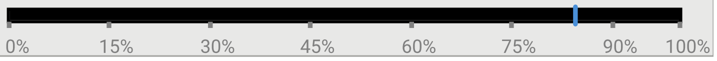
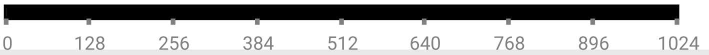
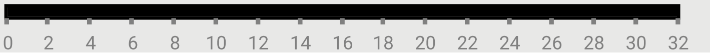
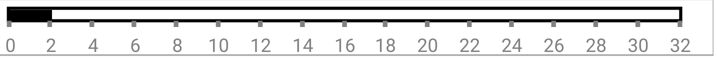
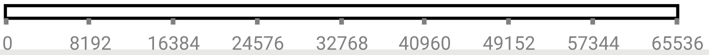
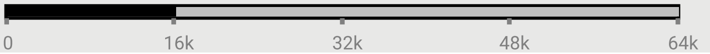
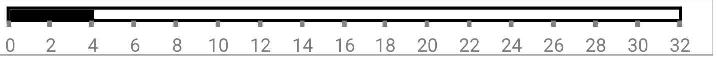
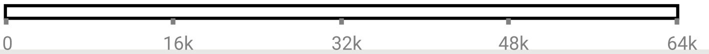
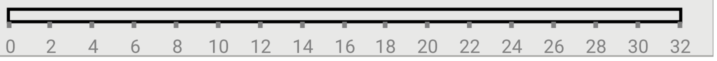
Texture - The texture sub-system is fully utilized or has too many outstanding requests.

Optimization: Resolve the primary stall issue; memory dependency.



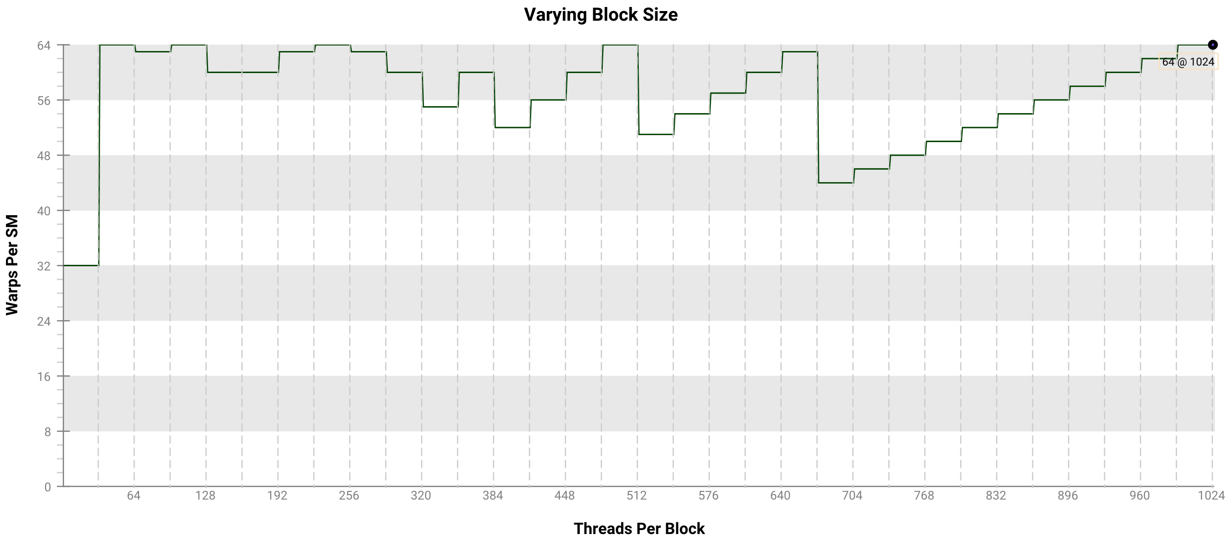
3.2. Occupancy Is Not Limiting Kernel Performance

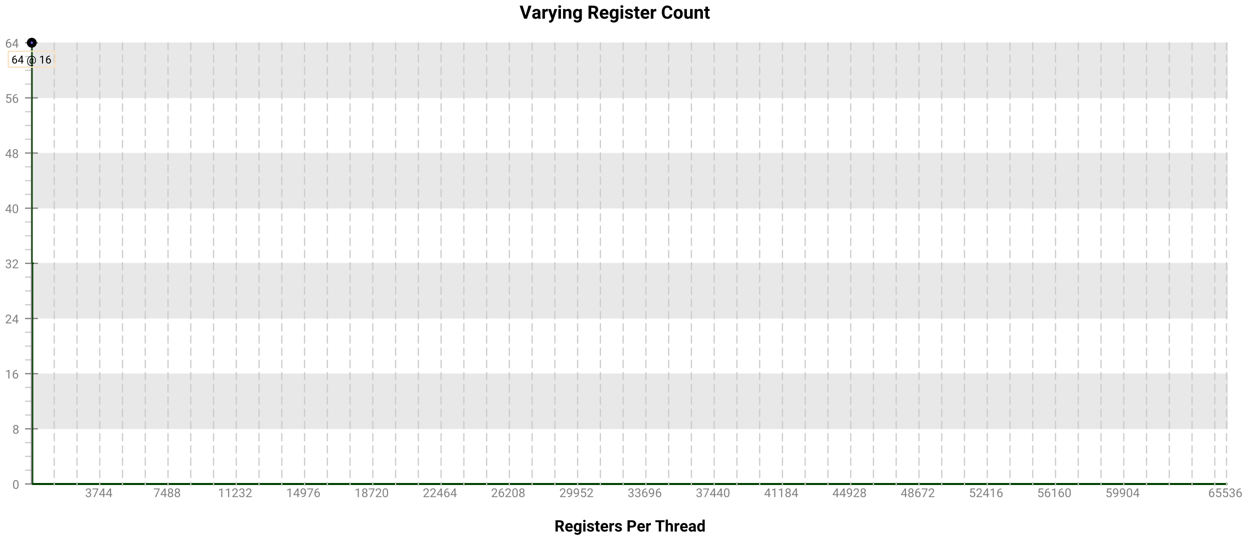
The kernel's block size, register usage, and shared memory usage allow it to fully utilize all warps on the GPU.

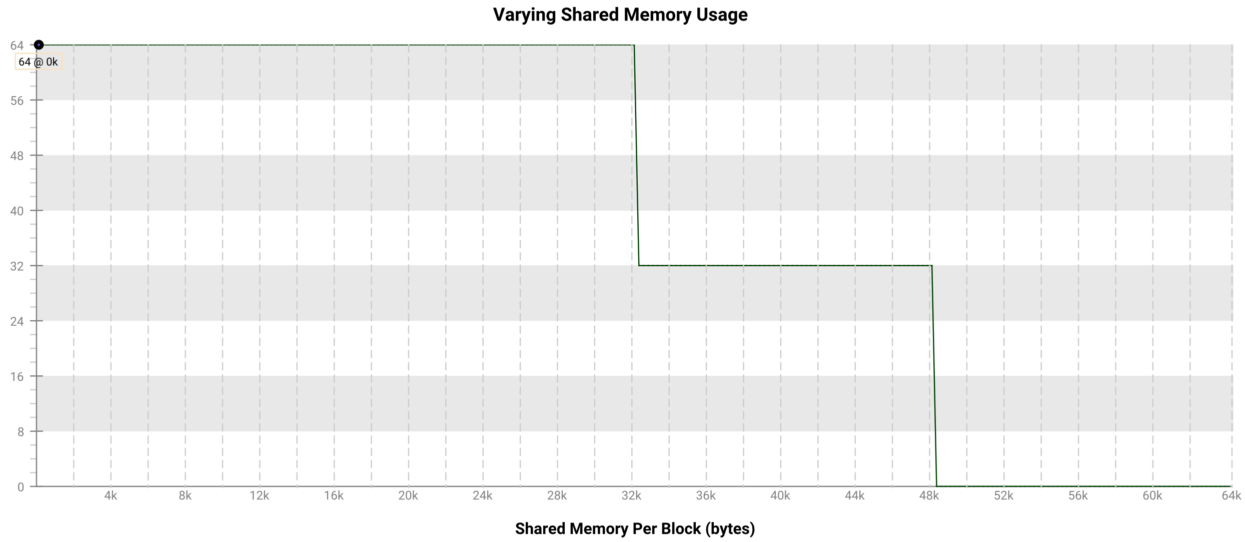
Variable	Achieved	Theoretical	Device Limit	Grid Size: [1024,1,1] (1024 blocks) Block Size: [1024,1,1] (1024 threads)
Occupancy Per SM				
Active Blocks		2	32	
Active Warps	53.79	64	64	
Active Threads		2048	2048	
Occupancy	84%	100%	100%	
Warps				
Threads/Block		1024	1024	
Warps/Block		32	32	
Block Limit		2	32	
Registers				
Registers/Thread		16	65536	
Registers/Block		16384	65536	
Block Limit		4	32	
Shared Memory				
Shared Memory/Block		0	65536	
Block Limit		0	32	

3.3. Occupancy Charts

The following charts show how varying different components of the kernel will impact theoretical occupancy.

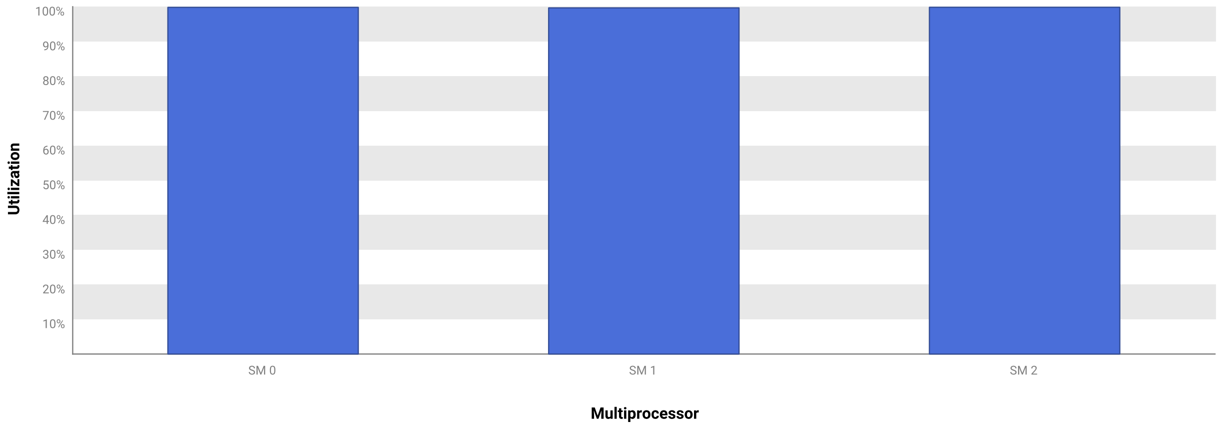

Varying Block Size


Varying Register Count



3.4. Multiprocessor Utilization

The kernel's blocks are distributed across the GPU's multiprocessors for execution. Depending on the number of blocks and the execution duration of each block some multiprocessors may be more highly utilized than others during execution of the kernel. The following chart shows the utilization of each multiprocessor during execution of the kernel.



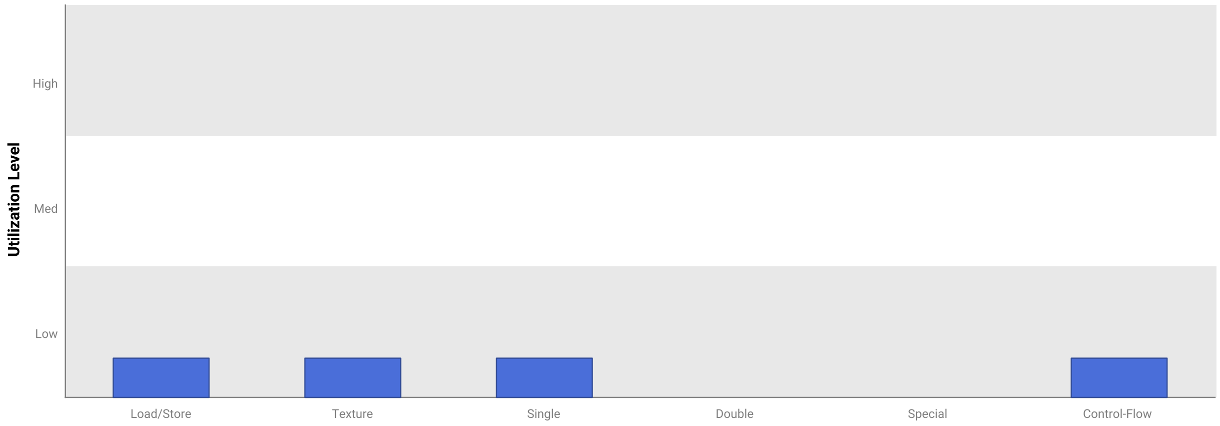
4. Compute Resources

GPU compute resources limit the performance of a kernel when those resources are insufficient or poorly utilized.

4.1. Function Unit Utilization

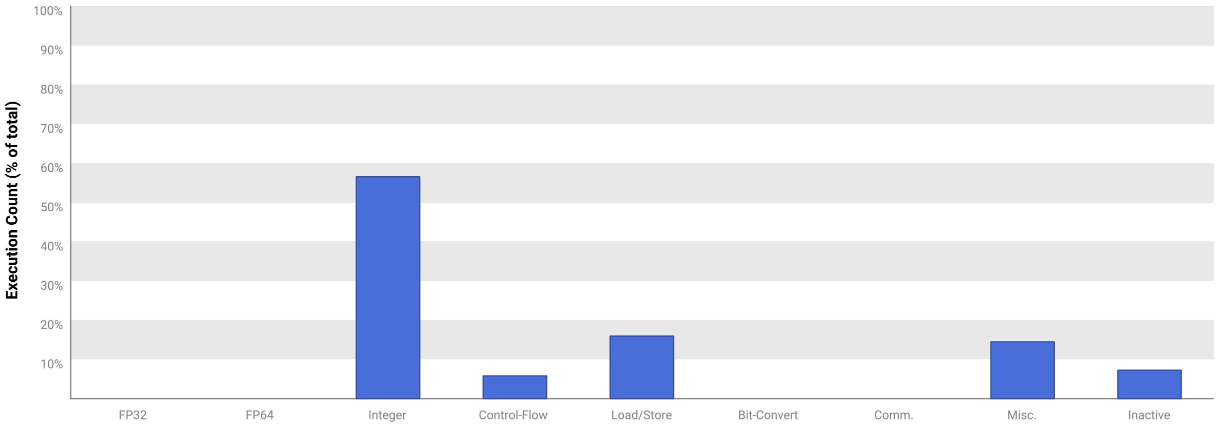
Different types of instructions are executed on different function units within each SM. Performance can be limited if a function unit is over-used by the instructions executed by the kernel. The following results show that the kernel's performance is not limited by overuse of any function unit.

Load/Store - Load and store instructions for shared and constant memory.
Texture - Load and store instructions for local, global, and texture memory.
Single - Single-precision integer and floating-point arithmetic instructions.
Double - Double-precision floating-point arithmetic instructions.
Special - Special arithmetic instructions such as sin, cos, popc, etc.
Control-Flow - Direct and indirect branches, jumps, and calls.



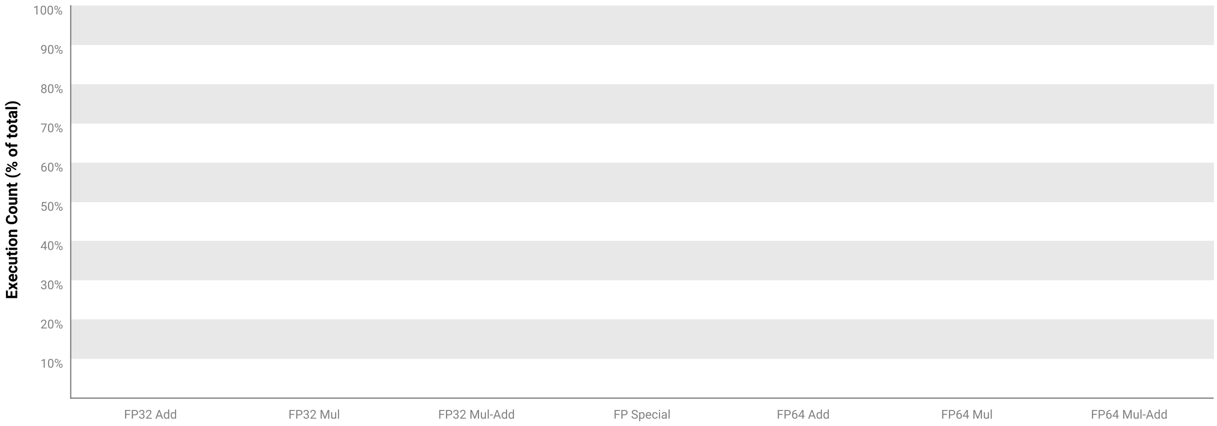
4.2. Instruction Execution Counts

The following chart shows the mix of instructions executed by the kernel. The instructions are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing instructions in that class. The "Inactive" result shows the thread executions that did not execute any instruction because the thread was predicated or inactive due to divergence.



4.3. Floating-Point Operation Counts

The following chart shows the mix of floating-point operations executed by the kernel. The operations are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing operations in that class. The results do not sum to 100% because non-floating-point operations executed by the kernel are not shown in this chart.


Execution Count (% of total)
FP32 Add
FP32 Mul
FP32 Mul-Add
FP Special
FP64 Add
FP64 Mul
FP64 Mul-Add