ECE-C353: Systems Programming Final Project

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December 3, 2011

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1 Design

The design for this multi-threaded recursive text search was based mostly on provided code, but included heavy modification to prevent deadlocks.

1.1 Header File

All global instance, function, and type definitions were placed in the header file work_crew.h, to improve the readability of the main source file work_crew.c.

```
int MAX_THREADS;
int number_sleeping = 0;
int all_done = 0;
int spawned = 0;
pthread_t *threads;

pthread_cond_t wake_up;
pthread_mutex_t num_sleeping;
pthread_mutex_t done;
pthread_mutex_t done;
```

These objects were declared globally, with the following intended meanings:

MAX_THREADS Maximum number of threads to possibly create.

num_threads Number of threads that have been spawned at any time.

number_sleeping Number of spawned threads that are currently sleeping.

all_done Boolean-like variable indicating if the entire origin directory has been searched.

spawned Boolean-like variable indicating if other threads have been spawned yet.

threads Pointer to array of type pthread, containing one element for each thread spawned.

wake_up Signal to wake threads up from sleeping.

num_sleeping Mutex protecting number_sleeping.

sleepers Pointer to array of mutexes protecting number_sleeping.

done Mutex protecting all_done.

In order to provide each thread with the appropriate information, a custom datatype was created as shown below.

```
/* structure for thread arguments */
12
13
  typedef struct thread_args
14
15
       int threadID;
16
17
       queue_t *queue;
18
       pthread_mutex_t *mutex_queue;
19
20
       int *num_occurrences;
       pthread_mutex_t *mutex_count;
21
22
23
       char **argv;
24
  } THREAD_ARGS;
```

Similar to previous assignments, the following pieces of information were passed to each thread:

threadID Integer identifier for the thread. The first thread spawned will receive ID 0, and subsequent threads will receive 1, 2, ..., NUM_THREADS - 1.

queue Pointer to the queue of files and directories to search.

mutex_queue Pointer to the mutex protecting the work queue.

num_occurrences Pointer to the integer representing the number of occurrences of the target string.

mutex_count Pointer to the mutex protecting num_occurrences.

Additionally, a custom datatype was created to ease the process of generating scriptable output.

```
27  /* structure for return values
28  *
29  * these make scripting the output easier */
30  typedef struct return_type
31  {
32    int count;
33   float time;
34  } RET_TYPE;
```

While these instance variables are fairly self-explanatory, they were created with the following intended meanings:

count The number of occurrences of the target string found by the returning function.

time Amount of time required to execute the returning function.

Finally, the non-main functions were declared:

```
36 /* Function prototypes */
37 RET_TYPE* search_for_string_serial(char **);
38 RET_TYPE* search_for_string_mt(char **);
39 void *worker_thread( void * );
```

in which search_for_string_serial is the single-threaded implementation, search_for_string_mt is the multi-threaded implementation, and worker_thread is the function executed by each thread.

1.2 Main File

1.2.1 Main

The main entry point for the search test is, as is expected, the main function, shown below.

```
int main(int argc, char** argv)
24
25
       if(argc < 3)
26
27
            printf("Usage: \_\%s_{\square} < search_\_string >_{\square} < path >_{\square} [num_{\square}
28
                threads]\n", argv[0]);
            exit(0);
30
       if( argc == 4 )
31
            MAX_THREADS = atoi(argv[3]);
32
33
            MAX_THREADS = 8;
34
35
       // Perform a serial search of the file system
36
       RET_TYPE *a = search_for_string_serial(argv);
37
38
       // Perform a multi-threaded search of the file
39
            system
       RET_TYPE *b = search_for_string_mt(argv);
40
41
       printf("\n\n\n");
42
       printf("_{UUUUUUUUUUUUUUUU
43
            Threads\tSingle\t\tMulti\t\tMatch\n");
       printf("Scriptable uoutput: \t%d\t%f\t%f\t%s\n",
44
            MAX_THREADS, a->time, b->time, (a->count ==
            b->count) ? "true" : "false");
       exit(0);
45
46
```

The function first checks the command line arguments to ensure that it has sufficient information to run — a string and a target directory. If this information is not provided, the usage syntax is printed to the screen and the application exists. Additionally, if the number of threads is not provided, it defaults to eight. Once this initial validation is complete, the single- and multi-threaded implementations are executed, and a script-safe output in the format of the following example:

```
Threads Single Multi Match
Scriptable output: 8 3.292550 3.163284 true
```

This provides the number of threads, execution time for both the single- and multi-threaded implementations, and whether or not the number of occurrences for the two implementations matches in an easily-parsable format. While this was not a requirement for the assignment, it made the process of acquiring and analyzing the resulting data much easier.

1.2.2 Single-Threaded

There were no appreciable changes made to the single-threaded implementation. While the return type was changed to RET_TYPE and there were slight modifications to the post-search code, the searching portions of the function remained unmodified from the provided version.

1.2.3 Multi-Threaded

The multi-threaded approach to string search begins by allocating space for the threads array and declaring and initializing the relevant mutexes.

```
194 RET_TYPE* search_for_string_mt(char **argv)
195
        /* allocate space for threads */
196
        threads = (pthread_t
197
            *)malloc(sizeof(pthread_t)*MAX_THREADS);
198
        /* initialize mutexes */
199
       pthread_mutex_t queue_mutex, count_mutex;
200
       pthread_mutex_init( &queue_mutex, NULL );
201
       pthread_mutex_init( &count_mutex, NULL );
202
       pthread_mutex_init( &num_sleeping, NULL );
203
204
       pthread_mutex_init( &done, NULL );
205
       pthread_cond_init( &wake_up, NULL );
207
208
        /* create mutex array */
       sleepers = (pthread_mutex_t *)malloc(
209
            sizeof(pthread_mutex_t)*MAX_THREADS );
       for( int i = 0; i < MAX_THREADS; i ++ )</pre>
210
            pthread_mutex_init( &sleepers[i], NULL );
211
212
        /* initialize count */
213
       int count = 0;
214
       int *p_count = &count;
215
216
        /* create and fill thread arguments */
217
       THREAD_ARGS *t_args;
218
219
       t_{args} = (THREAD_ARGS)
220
            *) malloc(sizeof(THREAD_ARGS));
221
       t_args->threadID = 0;
222
223
       t_args->mutex_queue = &queue_mutex;
       t_args->mutex_count = &count_mutex;
224
       t_args->num_occurrences = p_count;
225
       t_args->argv = argv;
```

Next, an array of mutexes for sleeping threads is created, as well as a series of counters. Finally, an instance of THREAD_ARGS is created and filled with the appropriate values.

At this point in the execution, a timer is created and started, and an initial thread is spawned.

```
/* start timer */
228
        struct timeval start;
229
        gettimeofday(&start, NULL);
230
231
        /* spawn thread */
232
        if( pthread_create( &threads[0], NULL,
            worker_thread, (void *)t_args ) != 0 )
234
            printf("E: \_could\_not\_create\_thread!_
235
                Exiting\n");
            exit(-1);
236
        }
237
```

After spawning the first thread, the multi-threaded control function simply waits for all threads to return, at which point it stops the timer, calculates the execution time, and prints its final search results and execution time.

```
/* wait for all threads to return */
239
240
        for( int i = 0; i < MAX_THREADS; i++ )</pre>
             pthread_join( threads[i], NULL );
241
242
        /* stop timer */
243
        struct timeval stop;
244
245
        gettimeofday(&stop, NULL);
246
         /* create return value */
247
        RET_TYPE *out = (RET_TYPE
248
             *)malloc(sizeof(RET_TYPE));
249
        out->time = (float)(stop.tv_sec - start.tv_sec +
             (stop.tv_usec - start.tv_usec)/(float)1000000);
        out->count = count;
250
251
252
        /* print results */
        printf("\n_{\sqcup}\n_{\sqcup}\n");
253
        printf("Overall_execution_time_=_%fs._\n",
254
             out->time);
        printf("The_{\sqcup}string_{\sqcup}\%s_{\sqcup}was_{\sqcup}found_{\sqcup}\%d_{\sqcup}times_{\sqcup}within_{\sqcup}
255
             the | file | system. | \n", argv[1], count);
        printf("========\n\n");
256
257
258
        return out;
259 }
```

1.2.4 Worker Thread

1.2.4.1 Setup Each thread begins simply by casting the arguments of worker_thread from void (as required by the pthread libraries) back to type THREAD_ARGS.

```
THREAD_ARGS *1_args = (THREAD_ARGS *) args;
263
264
265
       queue_element_t *element, *new_element;
       struct stat file_stats;
266
267
       int status;
       DIR *directory = NULL;
268
       struct dirent *result = NULL;
269
       /* allocate memory for the directory structure */
270
       struct dirent *entry = (struct dirent
271
            *)malloc(sizeof(struct dirent) + MAX_LENGTH);
```

As in the single-threaded case, two queue elements are first declared as well as the relevant objects for any directory information that will be encountered.

If the calling thread is the first one spawned (i.e. with a threadID of zero), the shared work queue is allocated and initialized.

```
if( l_args->threadID == 0 )
274
275
            /* Create and initialize the queue data
276
                 structure. */
            1_args->queue = create_queue();
277
            element = (queue_element_t
278
                 *) malloc(sizeof(queue_element_t));
            if(element == NULL)
279
280
                 \verb|printf("E: LError| allocating| memory.L
281
                     Exiting. ...\n");
282
                 exit(-1);
283
            strcpy(element->path_name, l_args->argv[2]);
284
            element -> next = NULL:
285
286
            /* Insert the initial path name into the queue
            insert_in_queue(l_args->queue, element);
287
```

While this could have also been performed before any threads were spawned, it was included here in the interest of clarity. Testing for the thread's ID happens once per thread and requires only O(1) time, so this causes no significant slowdown.

1.2.4.2 The Loop While the implementation of the searching algorithm required modifications in order to prevent deadlocking and data corruption, its overall flow nd operation remained largely unchanged. For this reason, only the changes will be discussed here.

Each thread begins by removing the first item of the queue. Because all threads access and modify the same queue, the mutex referenced by mutex_queue must be locked prior to access and unlocked after in order to prevent conflicts.

```
while(1)
291
292
           pthread_mutex_lock( l_args->mutex_queue );
293
            /* While there is work in the queue, process
294
                it. */
           if(l_args->queue->head != NULL)
295
296
           {
297
                queue_element_t *element =
                   remove_from_queue(1_args->queue);
                pthread_mutex_unlock( l_args->mutex_queue
                    );
```

While the single-threaded implementation tested the state of the queue in the while loop's condition, the test was moved to the first operation inside the loop so that the queue could be properly protected by its mutexes.

The next change occurs when the end of a directory is reached, in which the initial thread spawns its siblings.

```
if(result == NULL)
329
330
                                 if( l_args->threadID == 0 &&
331
                                      !spawned )
332
                                     spawned = 1;
333
                                     THREAD_ARGS *t_args;
334
                                     for( int i = 1; i <
335
                                          MAX_THREADS; i++ )
336
                                          t_args = (THREAD_ARGS
337
                                               *)malloc(sizeof(THREAD_ARGS));
                                          t_args->threadID = i;
338
                                          t_args->queue =
                                               l_args->queue;
340
                                          t_args->mutex_queue =
                                               1_args->mutex_queue;
                                          t_args->mutex_count =
341
                                               1_args->mutex_count;
                                          t_args->num_occurrences
342
                                               1_args->num_occurrences;
                                          /* NULL argv implies
343
                                               it's a worker
                                               thread */
                                          t_args->argv =
344
                                               1_args->argv;
345
346
                                          if( pthread_create(
                                               &threads[i], NULL,
                                               worker_thread,
                                               (void *)t_args )
                                               != 0 )
                                          Ł
347
                                               printf("E: \_could_{\bot}
348
                                                    \mathtt{not}_{\sqcup}\mathtt{create}_{\sqcup}
                                                    thread_{\sqcup}%d!_{\sqcup}
                                                    Exiting\n");
                                               exit(-1);
349
                                          }
350
                                     }
351
                                 }
352
353
                                 break; // End of directory
354
                           }
```

Each thread tests its ID and whether or not the extra threads have been spawned. If it is thread zero and the sibling threads haven't been created (as indicated by the value of the spawned variable), then an instance of THREAD_ARGS is created and filled prior to each thread's creation.

For each new element reached by a thread, the queue must be locked and subsequently unlocked in order to contribute to the work queue.

```
/* Construct the full path name
370
                             for the directory item stored
                             in entry. */
                         strcpy(new_element->path_name,
37
                             element->path_name);
                         strcat(new_element ->path_name,
372
                             "/");
                         strcat(new_element ->path_name,
373
                             entry->d_name);
374
                         /* safely add to queue */
375
376
                         pthread_mutex_lock(
                             l_args->mutex_queue );
                         insert_in_queue(l_args->queue,
                             new_element);
                         pthread_mutex_unlock(
                             1_args->mutex_queue );
379
                         /* wake up others
                           man, pthread_broadcast never
381
                              works for me */
                         for( int i = 0; i < MAX_THREADS-1;</pre>
382
                             pthread_cond_signal( &wake_up
```

Once the newly constructed queue element has been added and the queue unlocked, all remaining threads are woken up with an iterated call to pthread_cond_signal. While pthread_cond_broadcast should be used to wake multiple threads that are waiting on one signal, it does not work in this case — this is most likely caused by the fact that each thread locks a different mutex prior to calling pthread_cond_wait.

In the event that the queue is empty, a thread simply unlocks the queue's mutex and locks the mutexes protecting the number of sleeping threads and the completion status.

```
else
434
            {
435
                 pthread_mutex_unlock( l_args->mutex_queue
436
                     );
                 pthread_mutex_lock( &num_sleeping );
437
                 pthread_mutex_lock( &done );
438
439
                 if( all_done == 1 )
440
441
                     pthread_mutex_unlock( &done );
442
                     pthread_mutex_unlock( &num_sleeping );
443
444
                     pthread_exit(0);
                 }
445
```

If the process is complete, the two mutexes are unlocked (in the reverse order of locking to prevent deadlocks) and the thread exits. If the process is not complete but the currently-executing thread is the only remaining thread awake, it sets the completion status to true before unlocking the two previously locked mutexes.

```
else if ( number_sleeping == (MAX_THREADS -
446
                     1) )
447
                     all_done = 1;
448
                     pthread_mutex_unlock( &done );
449
                     pthread_mutex_unlock( &num_sleeping );
450
                     /* wake everyone up
452
                        man, pthread_broadcast never works
453
                          for me */
                     for( int i = 0; i < MAX_THREADS-1; i++</pre>
454
                         )
                         pthread_cond_signal( &wake_up );
455
                     pthread_exit( 0 );
456
                }
457
```

As it is the only thread not sleeping, the executing thread serially wakes the remaining threads before exiting. For the case that a thread is not the last one awake and the search process is not complete, it (safely) increments the number of sleeping threads. It next locks the mutex protecting its sleeping state and safely checks the completion state to ensure that the search hasn't completed since its last check. This prevents a thread from entering an eternal sleep.

```
{
459
                     pthread_mutex_unlock( &done );
460
461
                     /* add to number of sleeping */
                     number_sleeping ++;
462
463
                     pthread_mutex_unlock( &num_sleeping );
464
465
                     /* sleep */
                     pthread_mutex_lock(
466
                         &sleepers[l_args->threadID] );
                     pthread_mutex_lock( &done ); // check
467
                         done status just before sleeping
                     if( all_done == 1 )
469
                         pthread_mutex_unlock( &done );
                         pthread_mutex_unlock(
471
                             &sleepers[l_args->threadID] );
                         pthread_exit(0);
472
473
                     pthread_mutex_unlock( &done );
474
                     pthread_cond_wait( &wake_up,
475
                         &sleepers[l_args->threadID] );
476
                     pthread_mutex_unlock(
                         &sleepers[l_args->threadID] );
                     if( all_done == 1
                         pthread_exit(0);
478
479
                     {
480
                         pthread_mutex_lock( &num_sleeping);
481
                         number_sleeping --;
                         pthread_mutex_unlock(
483
                             &num_sleeping );
484
                         continue:
                    }
485
                }
```

If the search still has not completed then the thread waits for a wake_up signal, essentially putting it to sleep until another thread wakes it. Upon waking, a thread will either exit if the search is complete or decrement the number of sleeping threads and continue from the beginning of the loop once again.

2 Results

Initial testing of the single- and multi-threaded implementations was tested on my local development machine, but the final data was acquired through the College of Engineering's Linux cluster. A recursive search beginning at /home/DREXEL/nk78 for the string "Kandasamy" was performed one hundred times for maximum thread counts of one, two, four, and eight. By acquiring the data at roughly 2:00 PM on Thanksgiving Day, the number of processes contending for the cluster's CPU's was minimized. The resulting average, minimum, and maximum execution times are plotted in Figure 1 and tabulated in Table 1.

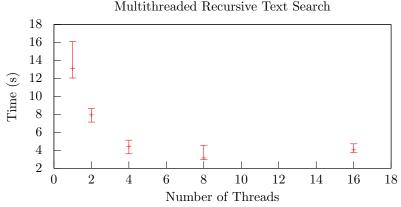


Figure 1: Mean, minimum, and maximum execution times from a set of 100 tests executed on xunil-o1.

As is shown, there is a nearly exponential decay in mean execution time over the range from one to eight threads, causing the average value to decrease from $13.09\,\mathrm{s}$ to $3.15\,\mathrm{s}$. This falls very much in line with one's expectations about parallelizing an algorithm such as the one used here. Single threaded, it takes roughly O(n) time; by adding i additional threads, this drops to $O(\frac{n}{i+1})$, as each thread can contribute $\frac{1}{i+1}$ of the work in a well-balanced program.

hline Number of Threads	Average (s)	Minimum (s)	Maximum (s)
1	13.09095947	12.047517	16.104431
2	7.94123965	7.145596	8.650321
4	4.43541181	3.615055	5.133428
8	3.1560362	2.960557	4.570848
16	4.05127182	3.758859	4.725392

Table 1: Data used to generate the plot in Figure 1.

There is however a discontinuity to this nearly exponential decay, located where the search was performed with 16 threads. Intuition states that this test case would perform better than any of the cases with fewer threads, but this is clearly not the case. Such a slow down is most likely caused by the added contention for mutexes introduced by multi-threading the process. All of the threads share a mutex protecting the completion status, one protecting the number of sleeping threads, and most importantly, a single work queue. As a result, only one thread can read from or post to the work queue at a time (similarly with checking or setting the completion status or the number of sleeping threads). As the number of threads increases, the contention for these limited resources becomes a bottleneck on the overall performance, eventually resulting in threads that spend more time waiting for shared resources than actually doing work.

A Scripts

A.1 Data Acquisition

```
1 #!/bin/bash
2 # Gets data for the ECE-C353 final project
3 SIG_TEXT="Scriptable output:"
4 STRING="kandasamy"
5 DIR="/home/DREXEL/nk78"
6 NUM_RUNS=10
  OUT = "data.txt"
9 > $OUT
  for loop in {1..100}
10
11
      echo "Running...u($100p/$NUM_RUNS)"
12
13
      for threads in 2 4 8 16
14
          out/work_crew $STRING $DIR $threads | grep "$SIG_TEXT" | sed
               "s/$SIG_TEXT\t//" >> $OUT
      done
16
17 done
18 echo "All done!"
```

A.2 Data Processing

```
1 #!/bin/bash
2 #
3
  # Re-orders data so that it can be easily plotted by gnuplot
4 # Author: Sean Barag <sjb89@drexel.edu>
5 IN="data_100.txt"
6 OUT="data_processed.txt"
7 OUT2="data_calculated.txt"
9 # split data into separate files
10 for i in 2 4 8 16
      cat $IN | grep "^$i" | awk '{print $3}' > /tmp/pData$i.txt
12
  done
14
15 # paste files together
16 paste /tmp/pData2.txt /tmp/pData4.txt /tmp/pData8.txt /tmp/pData16.txt > $OUT
18 # delete temporary files
19 rm /tmp/pData{2,4,8,16}.txt
21 # calculate min, max, and avg
python calcData.py > $OUT2
```

A.3 Data Calculation

```
1 #!/usr/bin/python
f = open('data_processed.txt', 'r')
3 lines = []
4 for line in f:
       lines.append(line.strip().split('\t'))
6 f.close()
8 cols = []
9 for i in range(4):
10
       cols.append([row[i] for row in lines])
12 f = open('data_100.txt', 'r')
13 slines = []
_{14}| for line in f:
       if line != '\n':
           slines.append(line.strip().split('\t'))
17
  cols.insert(0, [row[1] for row in slines])
19
22 print "#numuthreads\taverage\t\tmin\t\tmax"
23 for i in range(len(cols)):
      print 2**(i), "\t\t", sum(float(v) for v in cols[i])/len(cols[i]), "\t",
    min(cols[i]), "\t", max(cols[i])
24
```