

# Hopfield networks

# Associative memories

What were you doing on 22 November 1963? (Or 20 January 2009)?

Human memories are **content addressable**. Given some partial clue, retrieve the whole item. (Does this describe Google?)

Compare this with **direct access** computer memories where you have to know the location to find the contents.

Which memory techniques are (a) distributed? (b) fault tolerant?

Types of associative memory:

**Autoassociative** Given part of item X (partial, noisy), recall whole item X.

**Heteroassociative** Given part of item X (partial, noisy), recall whole item Y.

## Basins of attraction



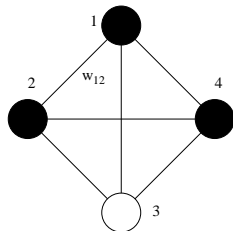
# Hopfield (1982): Equations 1

Term	meaning
$N$	number of units in network
$w_{ij}$	strength of connection from unit $i$ to unit $j$
$\mu$	pattern number, $\mu = 1 \dots p$
$x_i^\mu$	$i$ th component of pattern $\mu \{ +1, -1 \}$
$v_i$	state of unit $i \{ +1, -1 \}$
$\theta_i$	threshold value of unit $i$ (often zero).

- Fully-connected network: Each neuron connects to all other neurons; typically  $w_{ii} = 0$ . Symmetric weights:  $w_{ij} = w_{ji}$ .

e.g. Store three patterns:

pattern 1	-1	-1	+1	-1
pattern 2	-1	+1	+1	+1
pattern 3	+1	+1	-1	-1



## Equations 2

- State of unit  $i$ :

$$v_i = \text{sgn}\left(\sum_{j=1}^N w_{ij}v_j - \theta_i\right)$$

$$\text{where } \text{sgn}(x) = \begin{cases} +1 & \text{if } x \geq 0 \\ -1 & \text{if } x < 0 \end{cases}$$

- (Assume weights have been chosen.)
- Start network by setting state of nodes to values of input pattern.
- Asynchronous update: Select one unit at random, and update its state. (Gibbs sampling.)
- Repeat until network is stable (the state of all units in the network is unchanging).
- Synchronous update is unstable ... (ok for continuous time Hopfield model).

## Stability of one pattern

Can the network remember just one pattern?

Assume we are just storing one pattern ( $p = 1$ )  $x_i$ .

For a stable network,  $v_i = x_i \quad \forall i$ . When the state of a unit is updated, it should be the same as the input pattern:

$$v_i = x_i$$

$$\operatorname{sgn}\left(\sum_{j=1}^N w_{ij}x_j\right) = x_i$$

What should  $w_{ij}$  be to satisfy this constraint? If we “guess”  $w_{ij} = x_i x_j$  then:

$$\begin{aligned}\sum_{j=1}^N w_{ij}x_j &= (x_i x_1)x_1 + (x_i x_2)x_2 + \dots + (x_i x_N)x_N \\ &= x_i(x_1^2 + x_2^2 + \dots + x_N^2) \\ &= Nx_i \quad \text{since } x_i = \pm 1, x_i^2 = 1\end{aligned}$$

And since  $N$  is positive, we find that  $\operatorname{sgn}(\sum_{j=1}^N w_{ij}x_j) = x_i$ .

## Hopfield network: learning rule

$$w_{ij} = \frac{1}{N} \sum_{\mu=1}^p x_i^{\mu} x_j^{\mu}$$

Constant  $1/N$  is not needed for discrete updates.

Here we are measuring pairwise correlations.

After setting weights as above, enforce  $w_{ii} = 0$ . Useful to show convergence.  
Symmetric weights needed to ensure convergence.

This could be regarded as a **Hebbian** learning rule. More of them later.

# Example (Mackay)



. 0 0 0 0 -2 2 -2 2 2 -2 0 0 0 2 0 0 -2 0 2 2 0 0 -2 -2  
0 . 4 4 0 -2 -2 -2 -2 -2 0 -4 0 -2 0 0 -2 0 -2 -2 4 4 2 -2  
0 4 . 4 0 -2 -2 -2 -2 -2 0 -4 0 -2 0 0 -2 0 -2 -2 4 4 2 -2  
0 4 4 . 0 -2 -2 -2 -2 -2 0 -4 0 -2 0 0 -2 0 -2 -2 4 4 2 -2  
0 0 0 0 . 2 -2 -2 2 -2 2 -4 0 0 -2 4 -4 -2 0 -2 2 0 0 -2 2  
-2 -2 -2 -2 2 . 0 0 0 0 4 -2 2 -2 0 2 -2 0 -2 0 0 -2 -2 0 4  
2 -2 -2 -2 -2 0 . 0 0 4 0 2 2 -2 4 -2 2 0 -2 4 0 -2 -2 0 0  
-2 -2 -2 -2 -2 0 0 . 0 0 0 2 2 2 0 -2 2 4 2 0 0 -2 -2 0 0  
2 -2 -2 -2 2 0 0 0 . 0 0 -2 2 2 0 2 -2 0 2 0 4 -2 -2 -4 0  
2 -2 -2 -2 -2 0 4 0 0 . 0 2 2 -2 4 -2 2 0 -2 4 0 -2 -2 0 0  
-2 -2 -2 -2 2 4 0 0 0 0 . -2 2 -2 0 2 -2 0 -2 0 0 -2 -2 0 4  
0 0 0 0 0 -4 -2 2 2 -2 2 -2 . 0 0 2 -4 4 2 0 2 -2 0 0 2 -2  
0 -4 -4 -4 0 2 2 2 2 2 2 0 . 0 2 0 0 2 0 2 2 -4 -4 -2 2  
0 0 0 0 0 -2 -2 2 2 -2 -2 0 0 . -2 0 0 2 4 -2 2 0 0 -2 -2  
2 -2 -2 -2 -2 0 4 0 0 0 4 0 2 2 -2 . -2 2 0 -2 4 0 -2 -2 0 0  
0 0 0 0 0 4 2 -2 -2 2 -2 2 -4 0 0 -2 . -4 -2 0 -2 2 0 0 -2 2  
0 0 0 0 -4 -2 2 2 -2 2 -2 4 0 0 2 -4 . 2 0 2 -2 0 0 2 -2  
-2 -2 -2 -2 -2 0 0 4 0 0 0 2 2 2 0 -2 2 . 2 0 0 -2 -2 0 0  
0 0 0 0 0 -2 -2 2 2 -2 -2 0 0 4 -2 0 0 2 . -2 2 0 0 -2 -2  
2 -2 -2 -2 -2 0 4 0 0 4 0 2 2 -2 4 -2 2 0 -2 . 0 -2 -2 0 0  
2 -2 -2 -2 2 0 0 0 4 0 0 -2 2 2 0 2 -2 0 2 0 . -2 -2 -4 0  
0 4 4 4 0 -2 -2 -2 -2 -2 0 -4 0 -2 0 0 -2 0 -2 -2 . 4 2 -2  
0 4 4 4 0 -2 -2 -2 -2 -2 0 -4 0 -2 0 0 -2 0 -2 -2 4 . 2 -2  
-2 2 2 2 -2 0 0 0 -4 0 0 2 -2 -2 0 -2 2 -2 0 -4 2 2 . 0  
-2 -2 -2 -2 2 4 0 0 0 0 4 -2 2 -2 0 2 -2 0 -2 0 0 -2 -2 0 .





# Fault tolerance and capacity (Mackay)

Left: 20/300 connections deleted and extra pattern added.

Right: One further pattern causes more breakdown.



```

..1 1-1 1 x x-3 3 x x-1 1-1 x-1 1-3 x 1 3-1 1 x-1
-1 1 3 5-1-1-3-1-3-1-3 1 x 1-3 1-1-1-1-1-3 5 3 3-3
1 3 3 1-3-1 x-1-3-1-3-1-3 x-1-1-1 1-3 1-3-1 3 5 1-1
-1 5 3 -1-1-3-1-3-1-3 1-5 1-3 1-1-1-1-1-3 5 x 3-3
1-1 1-1 1-1-3 x x 3-5 1-1-1 3 x-3 1-3 3-1 1-3 3
x-1-3-1 1 -1 1-1 1 3-1 1-1-1 3-3 1 x 1 x-1-3 1 3
x-3-1-3-1-1 -1 1 3 1 1 3-3 5-3 3-1-1 x 1-3-1-1 1
-3-1 x-1-1-1-1 -1 1-1 1 3 1 x-1-1 1 5 1-1-1 x-3 1-1
3-3-1-3 x-1 1-1 -1-1-3 3 1 1 1-1-1-1 3-1-3-3-1 x 1
x-1-3-1 x 1 3 1-1 -1 3 1-1 3-1 x 1-3 5-1-1-3 1-1
x-3-1-3 3 3 1-1 1-1 -3 3-3 1 1-1-1-1-1-1-3-1 1 5
1 1-1 1-5-1 1 3-3 3-3 -1 1 1-3 3 x-1 3-3 1-1 3-3
1 x x-5 1 3 1 3 1 3-1 -1 3-1 1 1 1 3-5-3-3 3
-1 1-1 1-1-1-3 x 1-1-3 1-1 x 1-1 3 3-1 1 1-1-1-3
x-3-1-3-1-3 5-1 1 3 1 1 3 x x 3-1-1 3 1-3-1-1 1
-1 1 1 1 3 3-3 1 1-1 1-3-1 1 x -5-1-1-1 1 1-1 1 1
1-1 1-1 x-3 1-1 x-1 3 1-1 3-5 1 1 1-1-1 1 1-1
-3-1-3-3 1-1 5-1 1-1 x 3-1-1 1 1 1-1-1-3 1-1
x-1 1-1 1 x-1 1 3-3-1-1 1 3-1-1 1 1 -3 3-1 1-3-1
1-1-3-1-3 1 x 1-1 5-1 3 1-1 3-1 1 1-3 x-1-3 1-1
3-3-1-3 3 x 1-1 5-1 1-3 3 1 1 1-1 3 x -3-1-5 1
-1 5 3 5-1-1-3 x-3-1-3 1-5 1-3 1-1-1-1-1-3 3 x-3
1 3 5 1-3-1-3-1-3-1-3-1-3-1-1-1-3 1-3-1 3 1-1
x 3 1 3-3 1-1 1 x 1-1 3-3-1-1-1 1-3 1-5 x 1 -1
1-3-1-3 3 3 1-1 1-1 5-3 3-3 1 1-1-1-1-1 1-3-1-1
    
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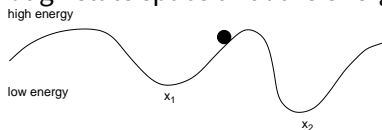


## Energy function

- Two units are in low energy state (stable) if states agree with sign of connection weight.

weight	$v_i$	$v_j$	energy
+	+1	+1	low
+	+1	-1	high
-	+1	+1	high
-	+1	-1	low

- Captured by  $e_{ij} = -w_{ij}v_i v_j$ .
- Energy of whole system:  $E = -1/2 \sum_{ij} w_{ij}v_i v_j$
- System travels through state space until the energy reaches a minimum:



- Each minima (“attractor”) corresponds to a stable state (stored pattern).
- Despite recurrency, system always settles in a stable state since energy monotonically decreases.

## Energy minimised (1)

$E$  is an energy function that always decreases to a minimum.

Proof: When updating the state of a given unit  $k$ , separate energy function into terms that depend on  $k$  and those that do not:

$$\begin{aligned} E &= -\frac{1}{2} \sum_{i \neq k, j \neq k} w_{ij} v_i v_j - \frac{1}{2} \sum_{j \neq k} w_{kj} v_k v_j - \frac{1}{2} \sum_{i \neq k} w_{ik} v_i v_k \\ &= -\frac{1}{2} \sum_{i \neq k, j \neq k} w_{ij} v_i v_j - \sum_{i \neq k} w_{ik} v_i v_k \end{aligned}$$

Let first term be constant ( $C$ ) since it does not involve  $k$ :

$$E = C - v_k \sum_{i \neq k} w_{ik} v_i$$

$$E = C - v_k a_k \quad (a_k \text{ is the total input to unit } k)$$

## Energy minimised (2)

After unit  $k$  is updated, new state =  $v'_k$ , new energy =  $E'$

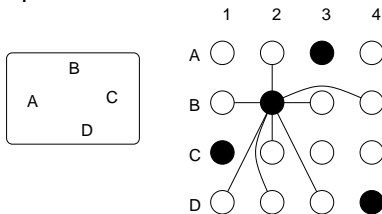
$$E' = C - v'_k a_k$$
$$\Delta E = E' - E = a_k(v_k - v'_k)$$

$a_k$	$v_k$	$v'_k$	$\Delta E$
+	-1		
+	+1		
-	-1		
-	+1		

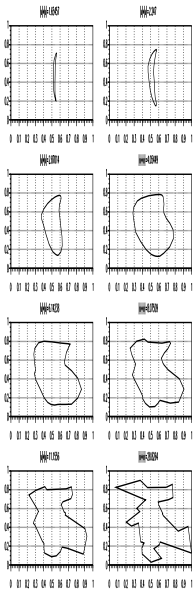
$\Delta E \leq 0$ . Therefore  $E$  always stays the same or decreases.

## Example: Travelling salesman problem (TSP)

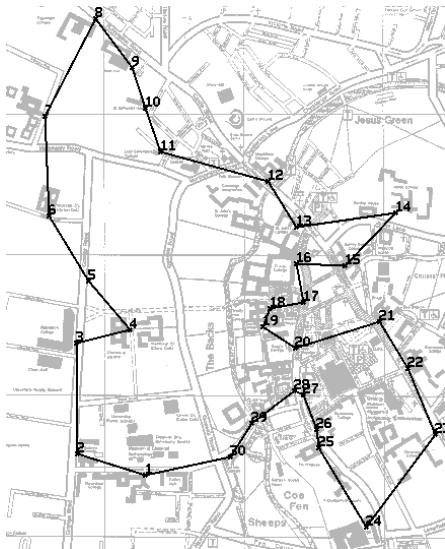
- Related to “N-Rooks” and “N-Queens” chess problems.
- Hopfield and Tank (1985) used network to solve TSP. Example use of “relaxation network”.
- Analogue rather than binary states  $v_i \in [-1, 1] : v_j = \tanh(\sum_i w_{ij}v_i)$ .
- Weights encode knowledge of problem:
  1. Cannot visit two cities at once (inhibition within column)
  2. Cannot visit a city more than once (inhibition within row)
  3. Inhibition between neighbouring columns proportional to distance between cities.
- Finds valid tours. Optimal tour hard to find!



# TSP, Cambridge style (Mackay, orig. Aiyer)



(a)



(b)

## Limitations: capacity and alternative states

- Empirical observations from Hopfield (1982) suggested about  $0.15N$  patterns could be recalled before recall was severely impaired.
- Theoretical limits (assuming random inputs) suggest around  $0.138N$  (Hertz, Krogh & Palmer, 1991).
- **Palimpsest** memories: erase old memories to store new memories. Clip weights or exponential decay of old weights.
- Can use supervised learning algorithms to improve capacity of Hopfield network.
- “spurious-states” as well as input patterns:
  1. Inverse of pattern also stable.  
e.g.  $[+1 + 1 - 1] \Leftrightarrow [-1 - 1 + 1]$   
Detect inverse by adding sign bit.
  2. Mixture states: combination of an odd number of input patterns.
  3. “Spin glass states” not correlated with any original pattern (when  $p$  is large).

# Summary

1. Associative memories.
2. Hopfield as a **generative** network.
3. Energy minimisation.
4. Travelling Salesman Problem (TSP).
5. Limitations. Future: (Restricted) Boltzmann machines, Generative Adversarial Networks.
6. See Stone Chapters 5–7.