

## Module-3 :> Java + DSA

### Collection Frameworks

① ArrayList

② Stack

③ Queue

④ Priority Queue { heap }

⑤ HashMap

### Object Oriented Programming

→ Classes & Objects

→ 5 pillars

① encapsulation

② data hiding

③ inheritance

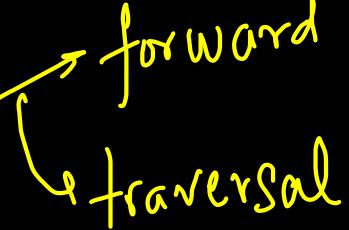
④ polymorphism

⑤ abstraction

Recursion

Elite Batch

## Arrays

- ① linear 
- ② contiguous memory allocation
  - ↳ random index access  $O(1)$
- ③ similar type of data (primitive array)
  - ↳ homogenous

## ④ Fixed Size Memory Allocation

`int[] arr = new int[5]`

`arr = new int[10]`

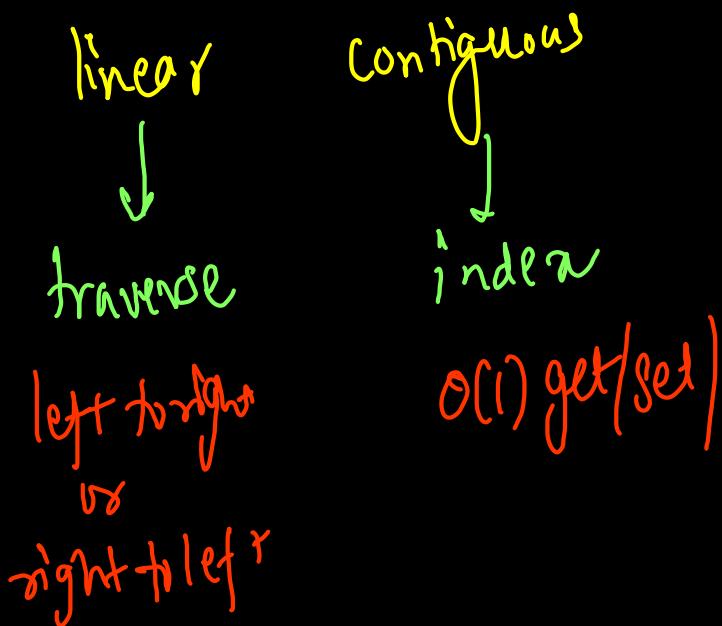
`arr = new int[n];`

↓  
it will remain

no addition  
no removal  
fixed size



internal implementation ~ array



`ArrayList<datatype> arr = new ArrayList<>();`

↳ `size` `arraylist` → `empty`

```
ArrayList<Integer> arr = new ArrayList<>();
ArrayList<String> arr2 = new ArrayList<>();
ArrayList<Character> arr3 = new ArrayList<>();
ArrayList<Boolean> arr4 = new ArrayList<>();
ArrayList<Double> arr5 = new ArrayList<>();
ArrayList<Long> arr6 = new ArrayList<>();
```

↑  
generics

Wrapper  
class

object

`int` → `Integer`

`char` → `Character`

`boolean` → `Boolean`

`long` → `Long`

`double` → `Double`

```
ArrayList<Integer> arr = new ArrayList<>();  
System.out.println(arr);  
  
arr.add(10);  
System.out.println(arr);  
arr.add(20);  
System.out.println(arr);  
arr.add(30);  
System.out.println(arr);  
arr.add(40);  
System.out.println(arr);
```

Insert in the last of  
the arraylist!  
→ O(1) constant

[]
[10]
[10, 20]
[10, 20, 30]
[10, 20, 30, 40]

array → length

String → length()

ArrayList → size()

```
ArrayList<Integer> arr = new ArrayList<>();
System.out.println(arr);
System.out.println(arr.size());

arr.add(10);
System.out.println(arr);
System.out.println(arr.size());

arr.add(20);
System.out.println(arr);
System.out.println(arr.size());

arr.add(30);
System.out.println(arr);
System.out.println(arr.size());

arr.add(40);
System.out.println(arr);
System.out.println(arr.size());
```

Finished in 74 ms

[]  
0  
[10]  
1  
[10, 20]  
2  
[10, 20, 30]  
3  
[10, 20, 30, 40]  
4

```
ArrayList<Integer> arr = new ArrayList<>();  
Scanner scn = new Scanner(System.in);  
int size = scn.nextInt();  
  
for(int count = 0; count < size; count++){  
    int var = scn.nextInt();  
    arr.add(var);  
    System.out.println(arr);  
}
```

user inputted arraylist

Finished in 147 ms

```
[10]  
[10, 20]  
[10, 20, 30]  
[10, 20, 30, 40]  
[10, 20, 30, 40, 50]  
[10, 20, 30, 40, 50, 60]  
[10, 20, 30, 40, 50, 60, 70]
```

```
ArrayList<Integer> arr = new ArrayList<>();  
arr.add(10);  
arr.add(20);  
arr.add(30);  
arr.add(40);  
arr.add(50);  
  
// Traversal left to right traversal  
for(int idx = 0; idx < arr.size(); idx++){  
    System.out.println(arr.get(idx));  
}
```

Finished in 89 ms

```
10  
20  
30  
40  
50
```

*arr.length → arr.size()*

*arr[idx] → arr.get(idx)*

```

public static void main(String[] args) {
    ArrayList<Integer> arr = new ArrayList<>();

    Scanner scn = new Scanner(System.in);
    int size = scn.nextInt();

    for(int idx = 0; idx < size; idx++){
        int val = scn.nextInt();
        arr.add(val);
    }

    // For Loop: Forward Traversal: Indices
    for(int idx = 0; idx < arr.size(); idx++){
        System.out.print(arr.get(idx) + " ");
    }

    System.out.println();

    // For Each Loop: Forward Traversal: Values
    for(int val : arr){
        System.out.print(val + " ");
    }
}

```

$10 \ 20 \ 30 \ 40 \ 50$   
 0 1 2 3 4

$\text{index} \rightarrow \text{value} = \text{arr.get}(idx)$

$\text{value} \rightarrow \cancel{\text{index}}$

// (OPTIONAL: EXTRA GYAAN) For Each Method: Lambda Expression  
arr.forEach((val) -> System.out.print(val + " "));

```
public static void main(String[] args) {  
    ArrayList<Integer> arr = new ArrayList<>();  
  
    Scanner scn = new Scanner(System.in);  
    int size = scn.nextInt();  
  
    for(int idx = 0; idx < size; idx++){  
        int val = scn.nextInt();  
        arr.add(val);  
    }  
  
    // For Loop: Backward Traversal: Indices  
    for(int idx = arr.size() - 1; idx >= 0; idx--){  
        System.out.print(arr.get(idx) + " ");  
    }  
  
    System.out.println();  
    |  
    Collections.reverse(arr);  
    // Reverse The ArrayList: Two Pointer  
  
    // For Each Loop: Forward Traversal: Values  
    for(int val : arr){  
        System.out.print(val + " ");  
    }  
}
```

ArrayList Reverse Printing

# ArrayList if - else

10  $\downarrow^L$  20  $\downarrow^R$  30 40 50

Format for next T Lines : (case, x(optional))

- case 1: Print the size of the ArrayList in a separate line.  $\text{size}()$
- case 2: Print and Remove element from the last index of the ArrayList.  $\text{remove}(\text{arr.size()}-1);$
- case 3: Print x and Add x in last index of the ArrayList.  $\text{add}(x);$
- case 4: Print and Remove an element from the starting (index = 0) of the ArrayList.  $\text{remove}(0);$
- case 5: Print x and Add x at beginning (index = 0) of the ArrayList.  $\text{add}(x, 0);$
- case 6: Print all the elements from left to right that are there inside the ArrayList.  $\rightarrow \text{for } \text{int } \text{id}x = 0 \dots ; \text{id}x < \text{arr.size();} \text{System.out.println(arr.get(id}x));}$

①  $\text{add}()$   $\rightarrow$  insert

②  $\text{size}()$   $\rightarrow$  count

③  $\text{get}()$   $\rightarrow$  read

## Sample Input 0

```

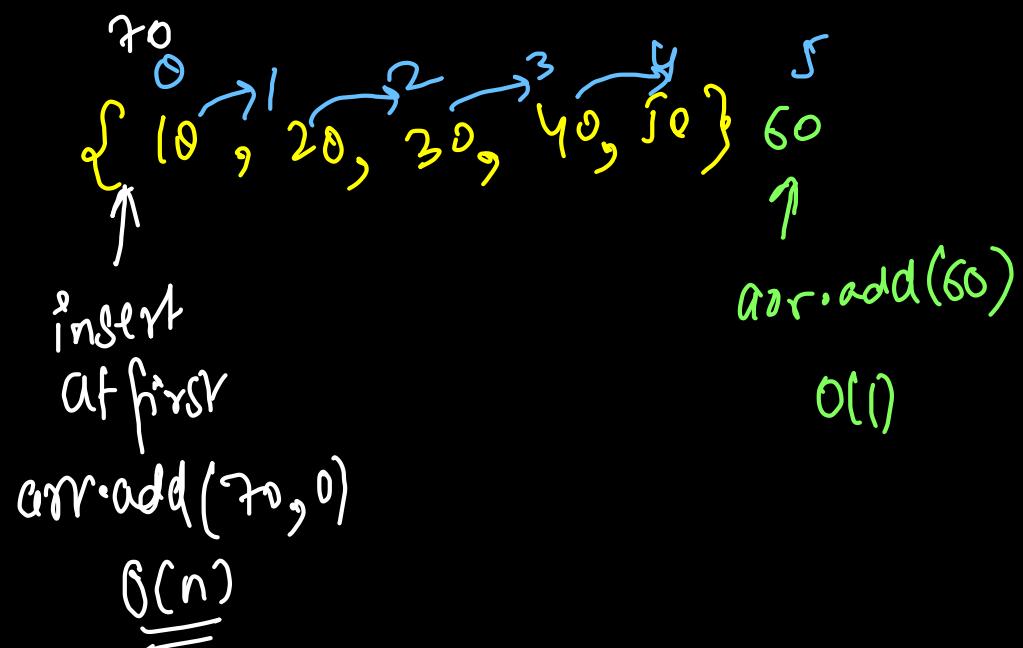
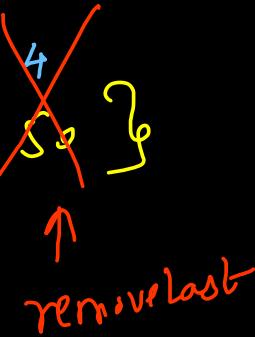
8
→ 2 : removeLast
→ 6 : traversal
→ 3 2 : addLast(2)
→ 5 1 : addFirst(1)
→ 6 : traversal
→ 1 : size
→ 3 3 addLast(3)
→ 2 : removeLast()

```

$\{ 1, 2, \cancel{3} \}$

invalid-move  
invalid-move

2  
1  
7 2  
2  
3  
3



```

public static void main(String[] args) {
    ArrayList<Integer> arr = new ArrayList<>();

    Scanner scn = new Scanner(System.in);
    int tests = scn.nextInt();

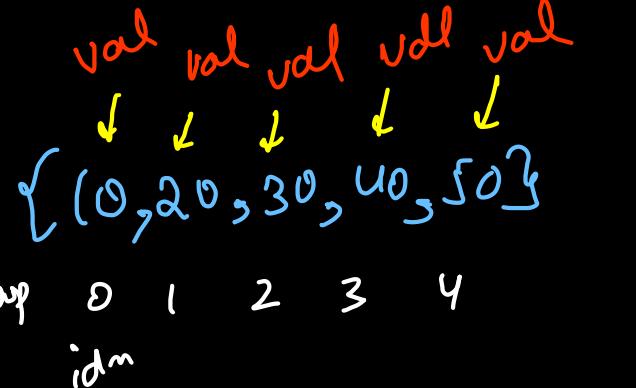
    for(int t = 0; t < tests; t++){
        int caseNo = scn.nextInt();

        if(caseNo == 1){
            System.out.println(arr.size()); // print size : O(1)
        }
        else if(caseNo == 2){

            if(arr.size() == 0) System.out.println("invalid-move"); // empty
            else System.out.println(arr.remove(arr.size() - 1)); // remove from last : O(1)
        }
        else if(caseNo == 3){

            int x = scn.nextInt();
            System.out.println(x);
            arr.add(x); // insert at last : O(1)
        }
    }
}

```


  
 val val val val val  
 ↓ ↓ ↓ ↓ ↓  
 { (0, 20, 30, 40, 50)  
 for (wp 0 1 2 3 4  
 idm

```

        else if(caseNo == 4){

            if(arr.size() == 0) System.out.println("invalid-move"); // empty
            else System.out.println(arr.remove(0)); // remove from front : O(n) shifting
        }
        else if(caseNo == 5){

            int x = scn.nextInt();
            System.out.println(x);
            arr.add(0, x); // insert at front : O(n) shifting
        }
    }
    else {

        if(arr.size() == 0) System.out.println("invalid-move"); // empty
        else {
            for(int val: arr){ // for each loop traversal
                System.out.print(val + " ");
            }
            System.out.println();
        }
    }
}

```

# Merge 2 Sorted ArrayLists

arr1:	10	30	<del>30</del>	50	60	80
	<del>p1</del>	<del>p1</del>	<del>p1</del>	<del>p1</del>	<del>p1</del>	
arr2:	20	40	<del>50</del>	60	60	90
	<del>p2</del>	<del>p2</del>	<del>p2</del>	<del>p2</del>	<del>p2</del>	

res: { 10 20 30 40 50 } }

```
public static ArrayList<Integer> merge(int[] a1, int[] a2){  
    int p1 = 0, p2 = 0;  
    ArrayList<Integer> res = new ArrayList<>();  
  
    while(p1 < a1.length && p2 < a2.length){  
        if(a1[p1] <= a2[p2]){  
            if(res.size() == 0 || res.get(res.size() - 1) != a1[p1])  
                res.add(a1[p1]);  
            p1++;  
        } else {  
            if(res.size() == 0 || res.get(res.size() - 1) != a2[p2])  
                res.add(a2[p2]);  
            p2++;  
        }  
    }  
  
    while(p1 < a1.length){  
        if(res.size() == 0 || res.get(res.size() - 1) != a1[p1])  
            res.add(a1[p1]);  
        p1++;  
    }  
  
    while(p2 < a2.length){  
        if(res.size() == 0 || res.get(res.size() - 1) != a2[p2])  
            res.add(a2[p2]);  
        p2++;  
    }  
  
    return res;  
}
```

```
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
  
    int n = scn.nextInt();  
    int[] arr1 = new int[n];  
  
    for(int idx = 0; idx < n; idx++){  
        arr1[idx] = scn.nextInt();  
    }  
  
    int m = scn.nextInt();  
    int[] arr2 = new int[m];  
  
    for(int idx = 0; idx < m; idx++){  
        arr2[idx] = scn.nextInt();  
    }  
  
    ArrayList<Integer> res = merge(arr1, arr2);  
    for(int val: res){  
        System.out.print(val + " ");  
    }  
}
```

# Stack Data Structure

1) pile of plates

2) Undo-redo

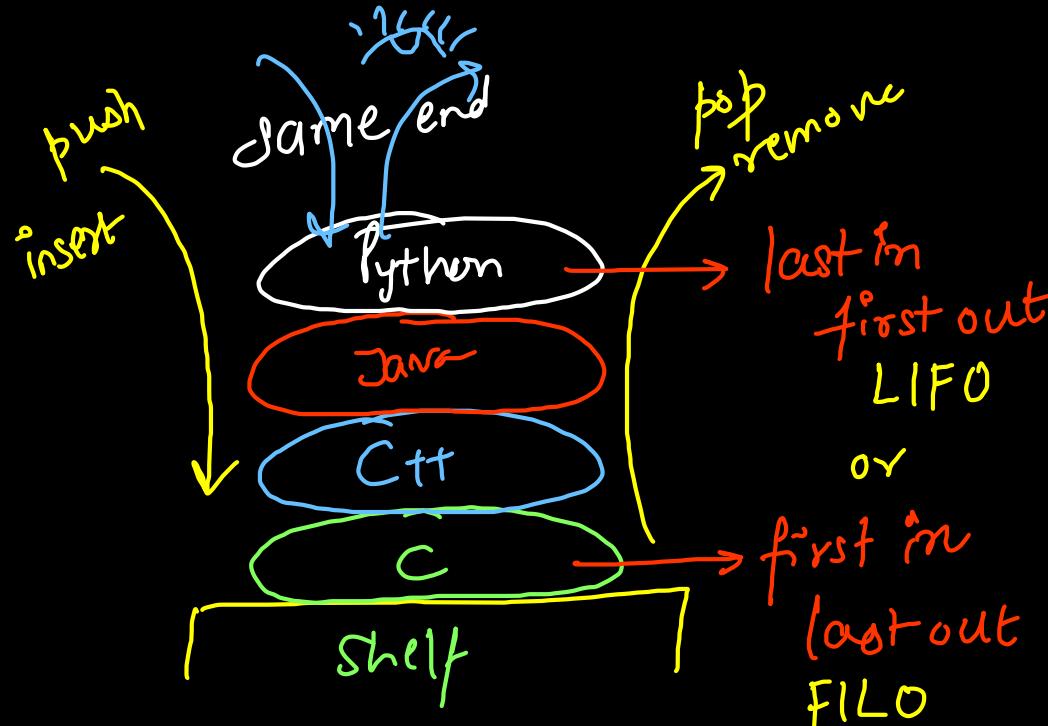
3) pile of books

4) cd holder

5) webpages

6) <sup>function/recursion</sup>  
call stack

~~ArrayList~~  
~~Stack~~



last in  
first out  
LIFO  
or  
first in  
last out  
FILO



## Stack Data Structure

LIFO or FILO policy

Stack < Integer >    stk = new Stack <>();

(1) push() : insert O(1)

(2) pop() : last O(1)

(3) peek() : get the top element O(1)

(4) empty()  $\begin{cases} \text{true: 0} \\ \text{false: } \neq 0 \end{cases}$  } O(1)

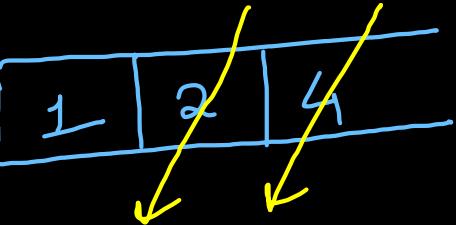
(5) size() : count of elements  $\rightarrow$  O(1)

```
public static void main(String[] args){  
    Stack<Integer> stk = new Stack<>();  
  
    System.out.println(stk.empty());  
    System.out.println(stk.size());  
  
    stk.push(10);  
    stk.push(20);  
    stk.push(30);  
    stk.push(40);  
  
    System.out.println(stk.empty());  
    System.out.println(stk.size());  
    System.out.println(stk.peek());  
    System.out.println(stk);  
  
    while(stk.size() > 0){  
        int val = stk.pop();  
        System.out.print(val + " ");  
    }  
}
```

Finished in 102 ms  
true  
0  
false  
4  
40  
[10, 20, 30, 40]  
40 30 20 10

1. Declare an Empty *stack s*.
2. Take Single Integer *T* as input.
3. For next *T* Lines format (*case, x(optional)*)
  - case 1. *Print* the *size* of the *stack* in a separate line.  $\rightarrow \text{size}()$
  - case 2. *Remove* an element from the stack. If the stack is empty then print  $-1$  in a separate line.  $\rightarrow \text{empty}() ? -1$   
 $\therefore \text{stk.pop()}$
  - case 3. *Add* Integer *x* to the *stack s*.  $\sim \text{stk.push}(x);$
  - case 4. *Print* an element at the *top* of the *stack*. If stack is empty print  $-1$  in a seperate line.  
 $\left( \begin{array}{l} \text{empty}() ? -1 \\ \therefore \text{stk.peek()} \end{array} \right)$

10
3 1 ✓
3 2 ✓
4 ✓
4 ✓
2 ✓
4 ✓
3 4 ✓
2 ✓
4 ✓
1 ✓



stk.add(1)  
stk.add(2)  
stk.peek(): ②  
stk.peek(): ②  
stk.pop()  
stk.peek(): ①  
stk.add(4)  
stk.pop()  
stk.peek(): ③  
stk.size(), 1

```
public class Solution {  
  
    public static void main(String[] args) {  
        Scanner scn = new Scanner(System.in);  
  
        Stack<Integer> stk = new Stack<>();  
  
        int tests = scn.nextInt();  
        for(int t = 0; t < tests; t++){  
            int choice = scn.nextInt();  
  
            if(choice == 1){  
                System.out.println(stk.size()); // size  
  
            } else if(choice == 2){  
                if(stk.empty()) System.out.println(-1);  
                else stk.pop(); // remove  
            } else if(choice == 3){  
                int x = scn.nextInt();  
                stk.push(x); // insert  
            } else {  
                if(stk.empty()) System.out.println(-1);  
                else System.out.println(stk.peek()); // get top  
            }  
        }  
    }  
}
```

Annotations:

- $O(T \text{ test cases})$  (red handwritten note with arrows pointing to the loop header and the `choice` variable)
- $O(1)$  (red handwritten note with arrows pointing to the `System.out.println(stk.size())` line and the `if(stk.empty())` condition)
- $O(1)$  (red handwritten note with arrows pointing to the `stk.pop()` call and the `System.out.println(stk.peek())` call)
- $O(1)$  (red handwritten note with arrows pointing to the `stk.push(x)` call)

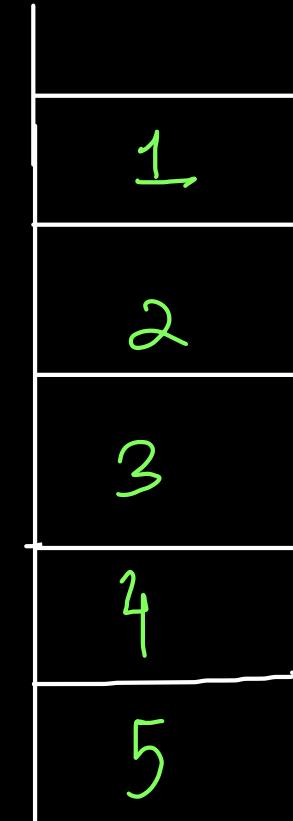
# Reverse Integers using Stack

int 12345 → 54321

int 70000 → 7

int 70707 → 70707

98765 → 56789



12345  
↓  $\%10 \rightarrow 5$

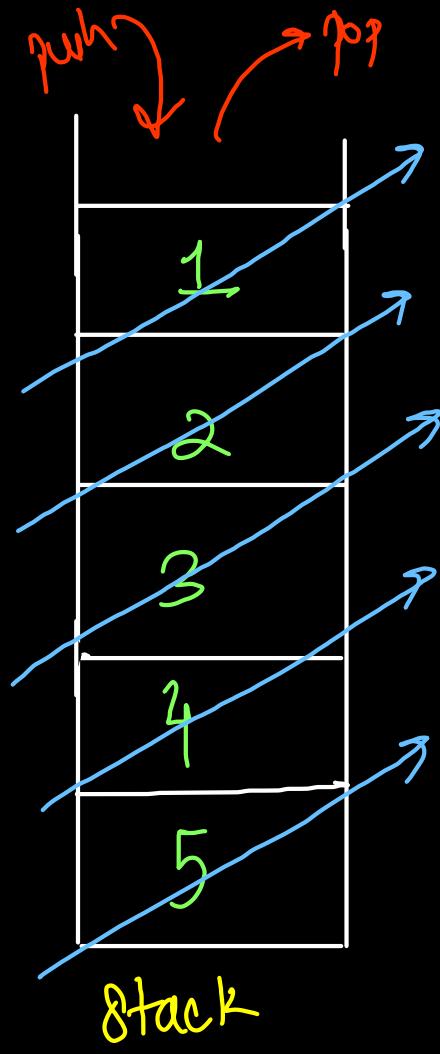
1234  
↓  $\%10 \rightarrow 4$

123  
↓  $\%10 \rightarrow 3$   
/10

12  
↓  $\%10 \rightarrow 2$

1  
↓  $\%10 \rightarrow 1$   
/10

0



54321

$$= 5 \times 10^4 + 4 \times 10^3 + 3 \times 10^2 + 2 \times 10^1 + 1 \times 10^0$$

$$708 = 0 + 1 \times 10^0 = 1 + 2 \times 10 = 21 + 3 \times 10^0$$

$$= 321 \\ + 4 \times 10^0$$

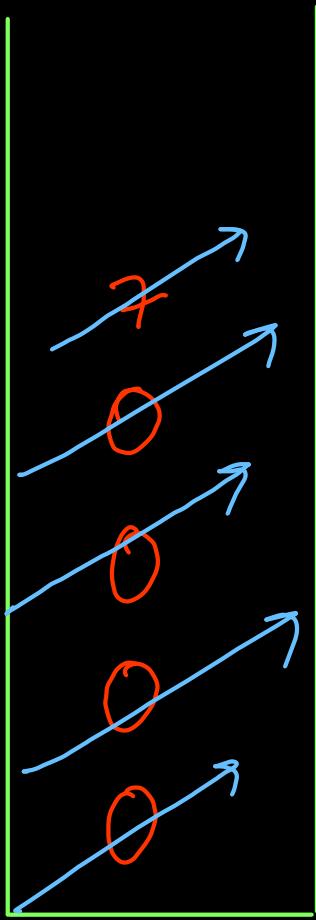
$$= 4321$$

$$+ 5 \times 10000$$

$$= 54321$$

Required  
and  
=

$$\text{power} = 1 \times 10 = 10 \times 10 = 100 \times 10 = 1000 \times 10 = 10000$$



① Fill the stack

70000

$$\downarrow \begin{matrix} \%10 = 0 \\ /10 \end{matrix}$$

7000

$$\downarrow \begin{matrix} \%10 = 0 \\ /10 \end{matrix}$$

700

$$\downarrow \begin{matrix} \%10 = 0 \\ /10 \end{matrix}$$

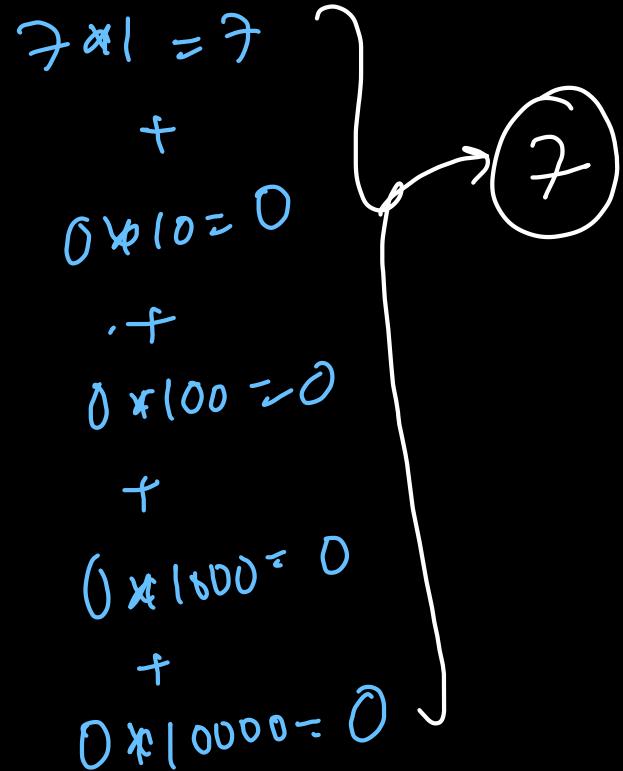
70

$$\downarrow \begin{matrix} \%10 = 0 \\ /10 \end{matrix}$$

$$\downarrow \begin{matrix} \%10 = 7 \\ /10 \end{matrix}$$

⑥

$$\text{power} = 10^0 = 1 \times 10 \times 10 \times 10 \times 10 \times 10$$



```

public static int reverse(int num){
    Stack<Integer> stk = new Stack<>();

    // 1. fill the stack
    while(num > 0){
        int digit = num % 10;
        stk.push(digit);
        num = num / 10;
    }

    // 2. clear the stack
    int res = 0, power = 1;

    while(stk.size() > 0){
        int digit = stk.pop();

        res = res + digit * power;
        power = power * 10;
    }

    return res;
}

```

digits

```

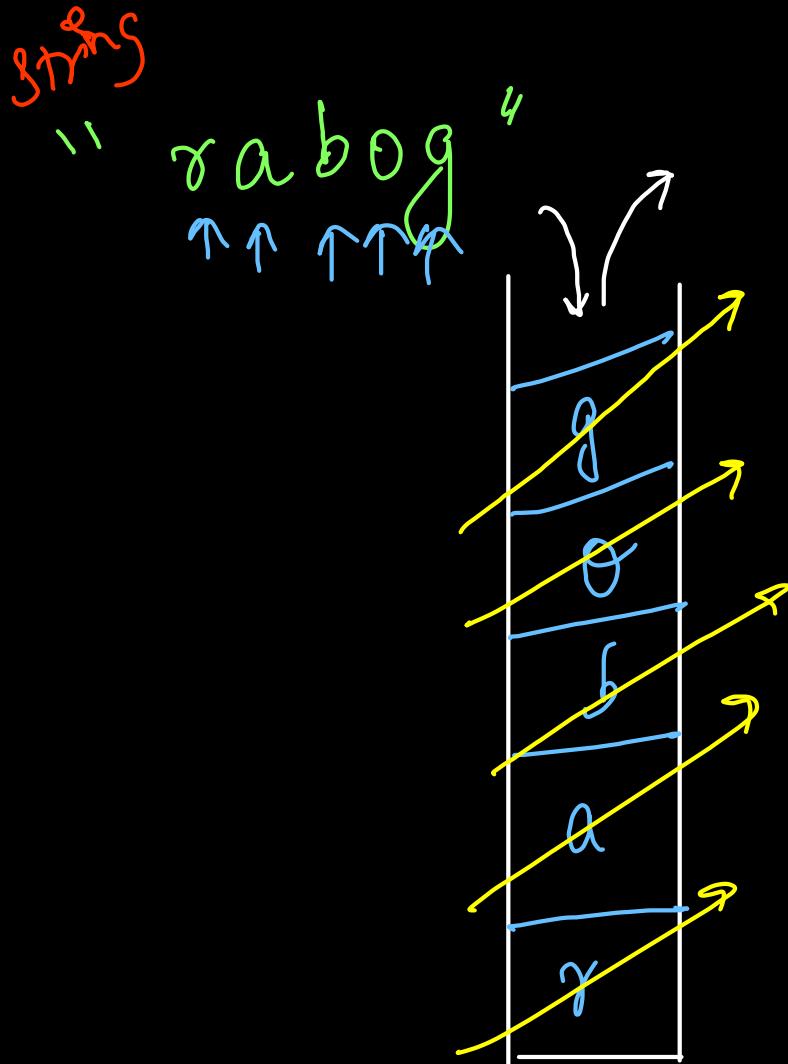
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int tests = scn.nextInt();

    for(int t = 0; t < tests; t++){
        int num = scn.nextInt();
        System.out.println(reverse(num));
    }
}

```

$$\begin{aligned}
 & O(T * (10+10)) \\
 & O(T) \quad O(T * \cancel{\frac{\log n}{10}}) \\
 & \leq 10
 \end{aligned}$$

## Reverse Stack Using Stack



res(StringBuilder)

" " + 'g' → "g"

"g" + 'o' → "go"

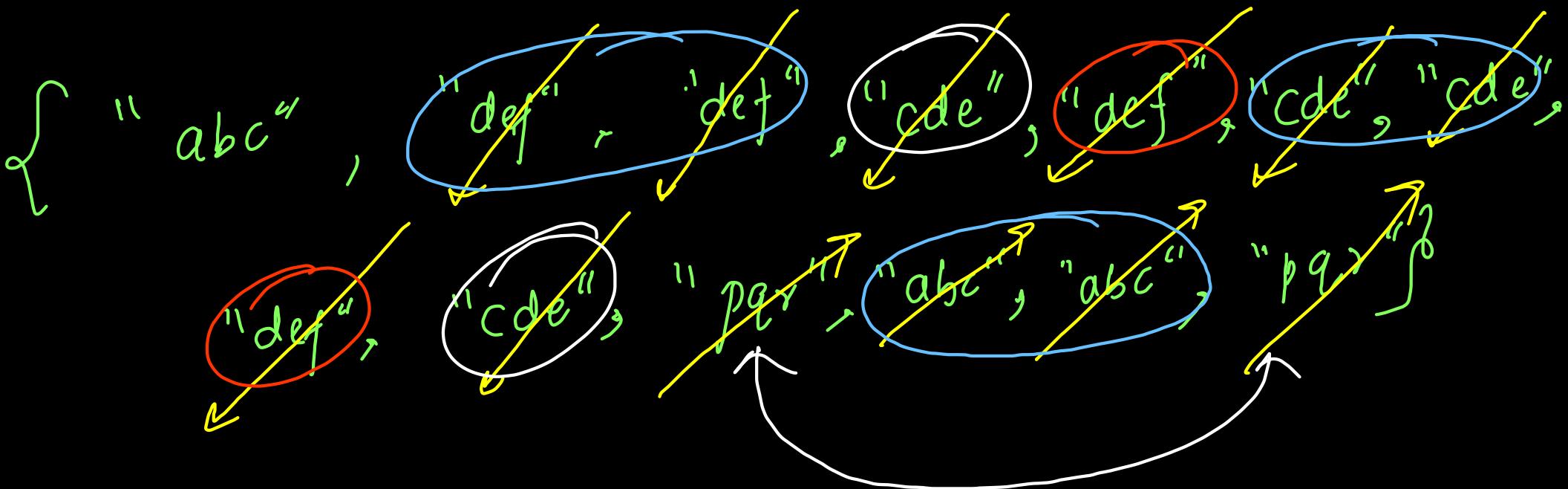
"go" + 'b' → "gob"

"gob" + 'a' → "goba"

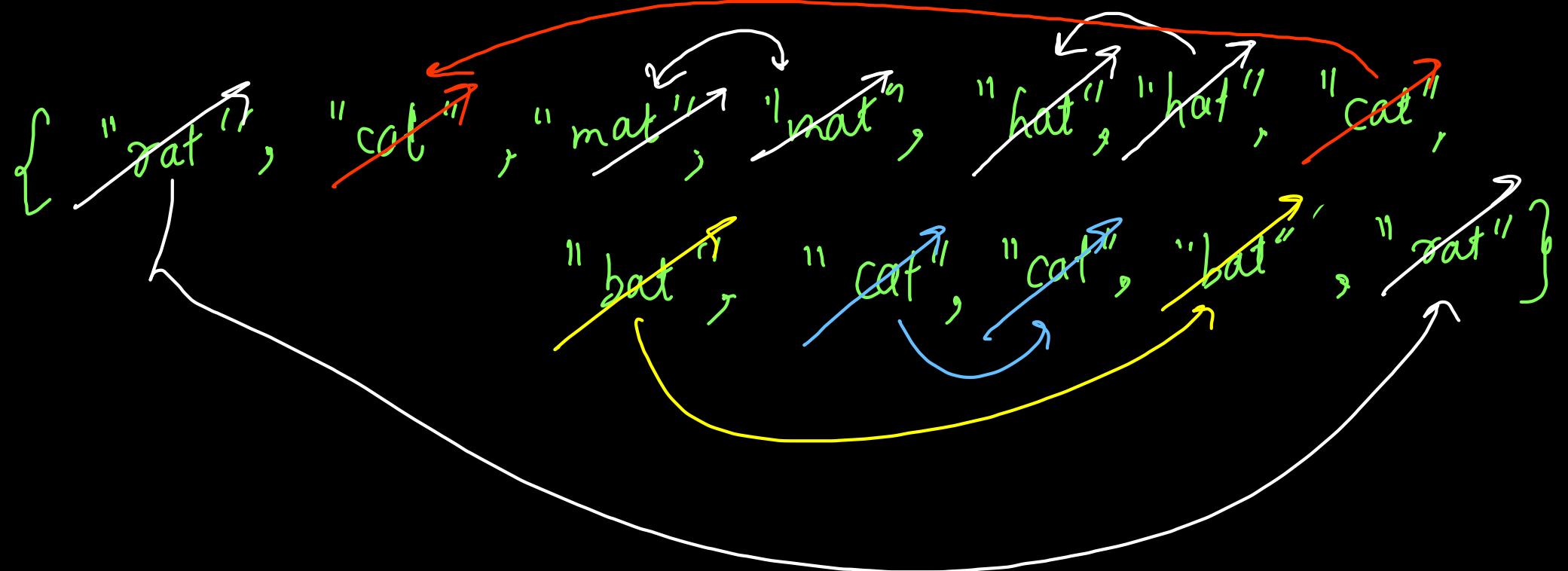
"goba" + 'g' → "gobag"

```
public static String reverse(String str){  
    Stack<Character> stk = new Stack<>();  
  
    for(int idx = 0; idx < str.length(); idx++){  
        char ch = str.charAt(idx);  
        stk.push(ch);  
    }  
  
    StringBuilder res = new StringBuilder();  
  
    while(stk.size() > 0){  
        char ch = stk.pop();  
        res.append(ch);  
    }  
  
    return res.toString();  
}
```

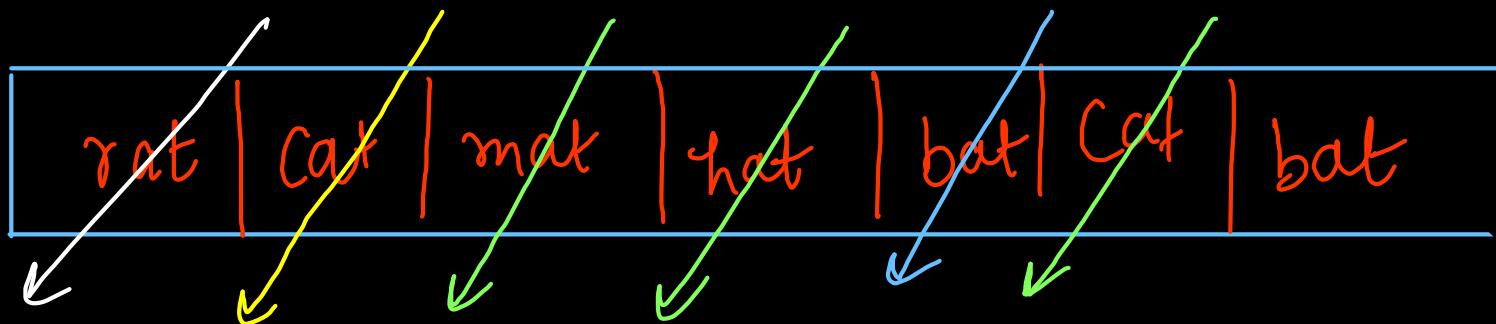
## ~~Remove All Adjacent Duplicates Recursively~~



1 remaining



∅ string remaining



stk.size()  $\Rightarrow$  ①

if ( $stk.size() == 0$  ||  $str.equals(stk.peek()) == false$ )

$stk.push(str);$

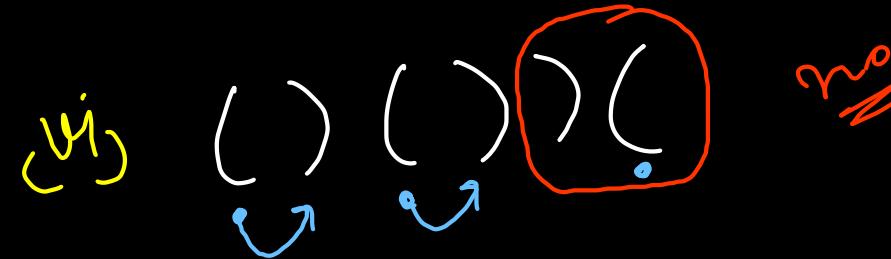
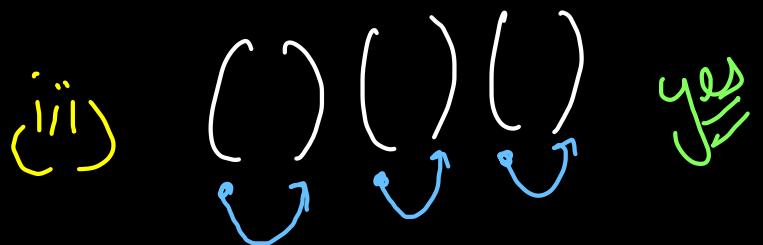
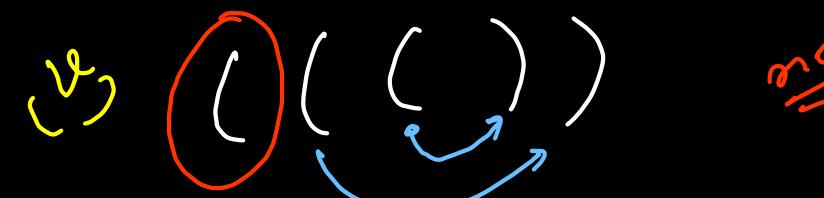
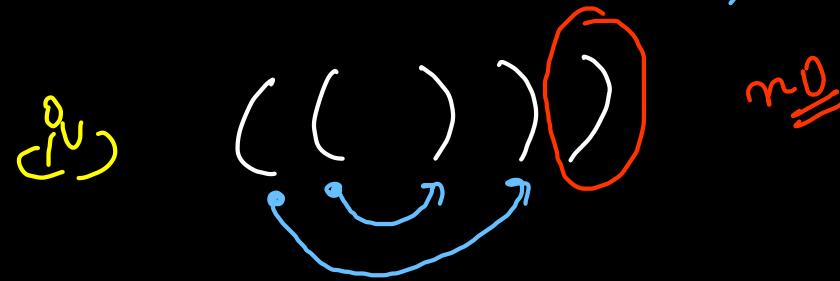
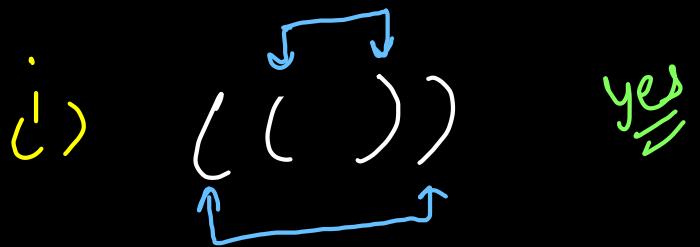
else  $stk.pop();$

```
public static void main(String[] args) {  
    Stack<String> stk = new Stack<>();  
  
    Scanner scn = new Scanner(System.in);  
    int n = scn.nextInt();  
  
    for(int idx = 0; idx < n; idx++){  
        String str = scn.next();  
  
        if(stk.size() == 0 || str.equals(stk.peek()) == false){  
            stk.push(str);  
        } else {  
            stk.pop();  
        }  
    }  
  
    System.out.println(stk.size());  
}
```

$\rightarrow O(n)$

Leetcode

## Balanced / Valid parentheses {stack}



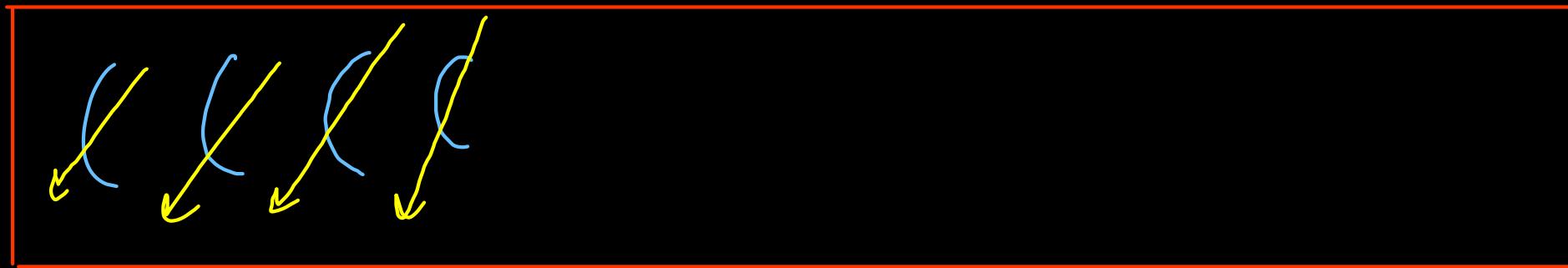
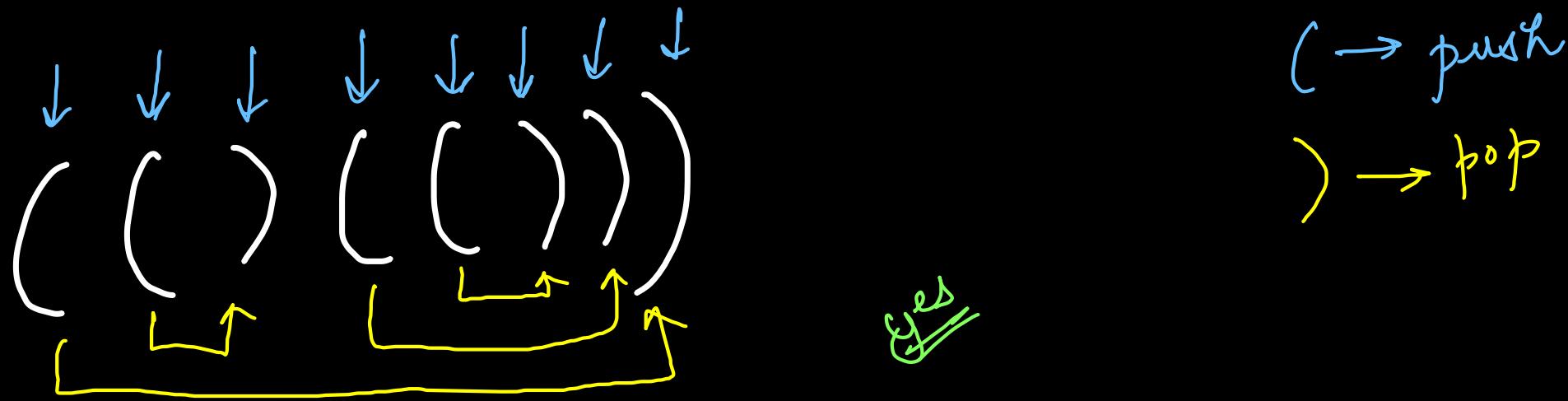
push → ?

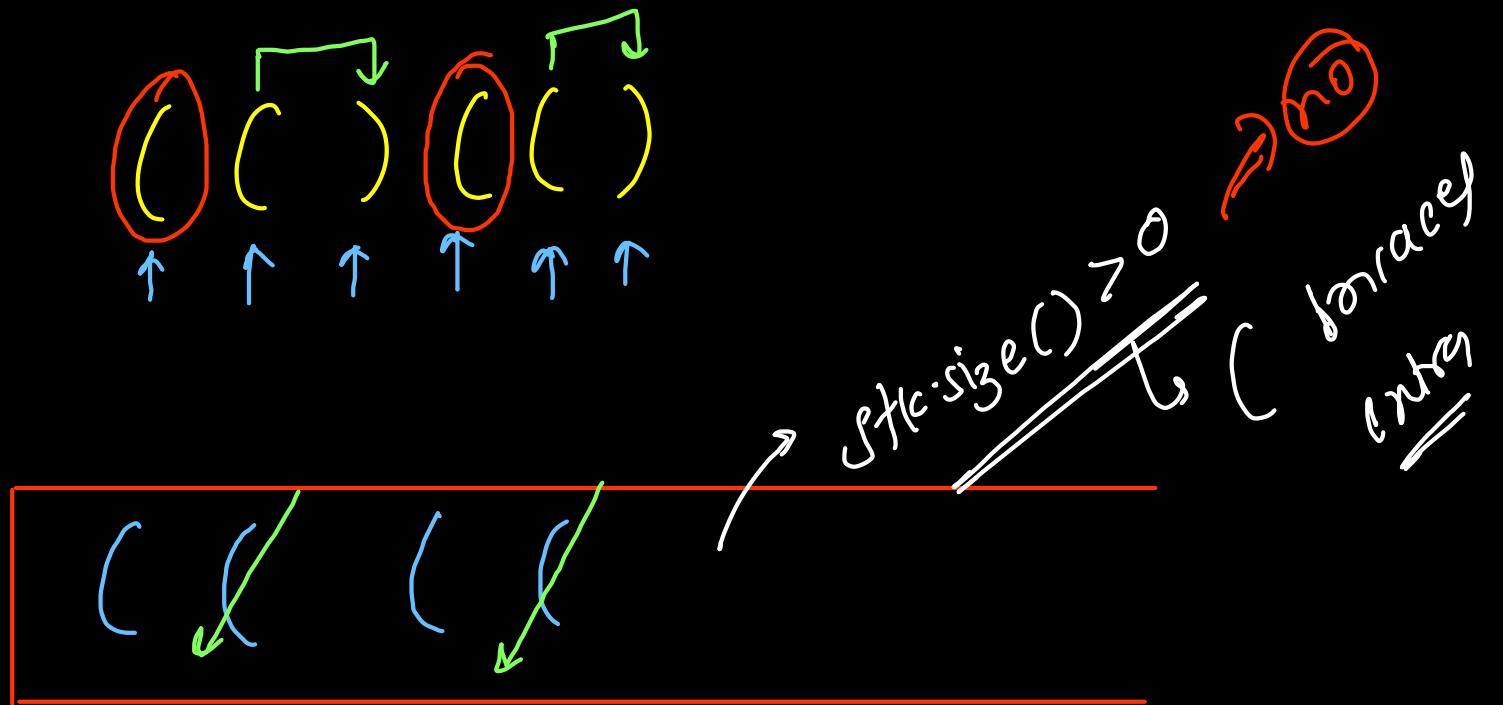
pop → ?

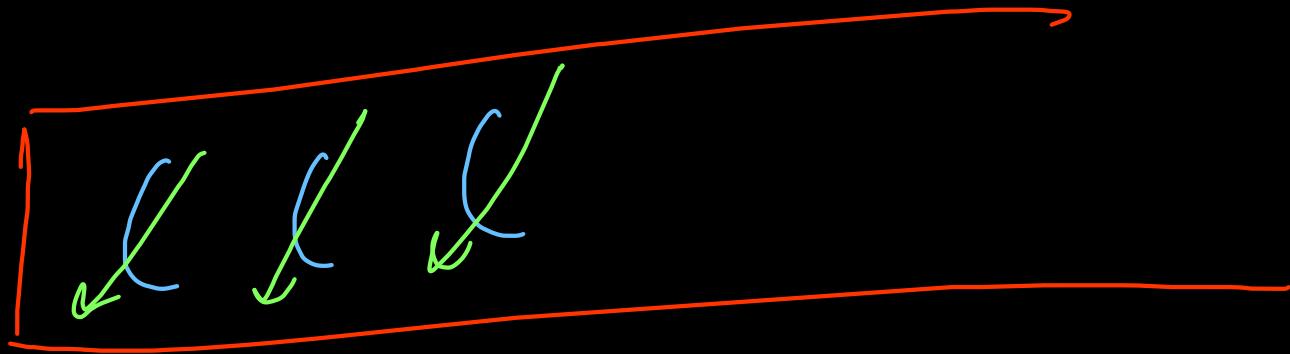
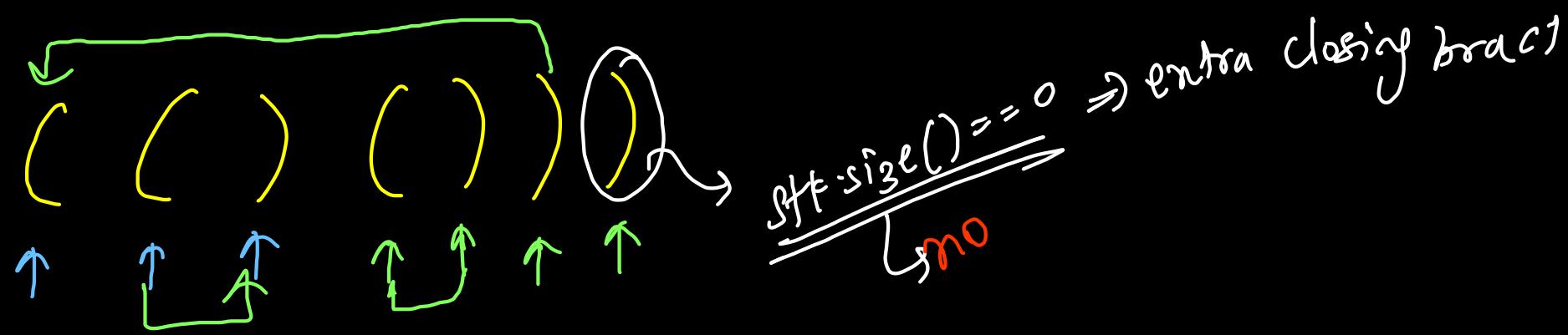
$$((a+b) * (c+d))$$

p  
q

p\*q







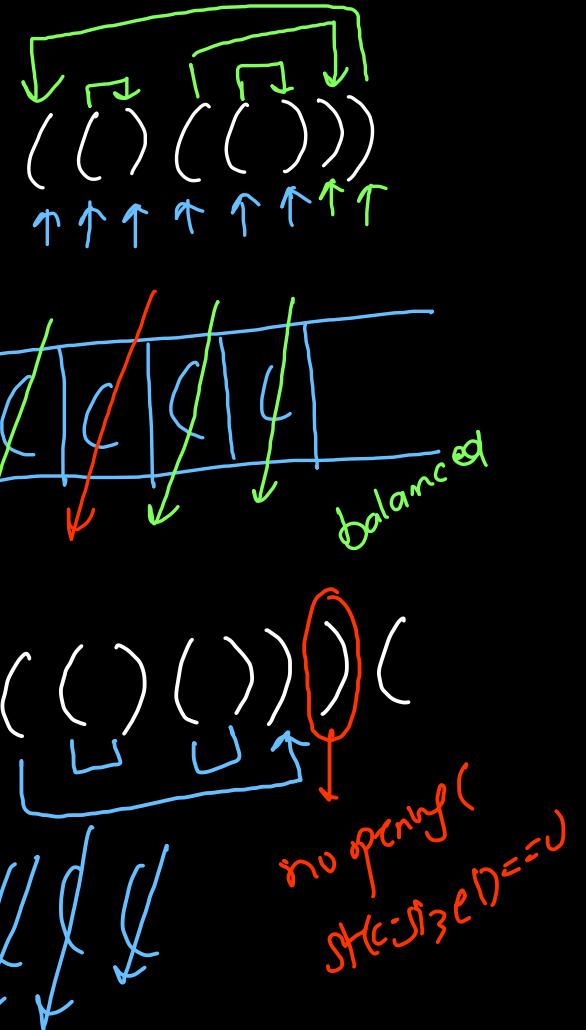
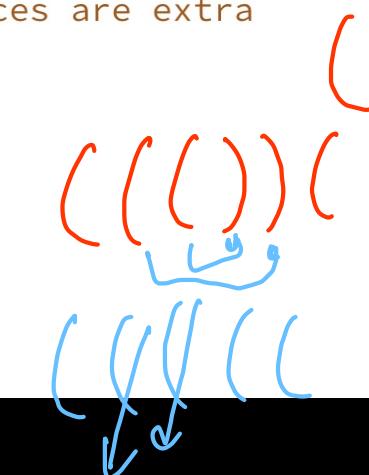
```
public static boolean validParantheses(String str){  
    Stack<Character> stk = new Stack<>();  
  
    for(int idx = 0; idx < str.length(); idx++){  
        char ch = str.charAt(idx);  
  
        if(ch == '('){  
            stk.push('(');  
        } else if(ch == ')'){  
            if(stk.size() == 0) return false; // closing braces are extra  
            stk.pop();  
        }  
    }  
  
    if(stk.size() > 0) return false; // opening braces are extra  
    return true;  
}  
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
    String str = scn.next();  
  
    System.out.println(validParantheses(str));  
}
```

) a+b(

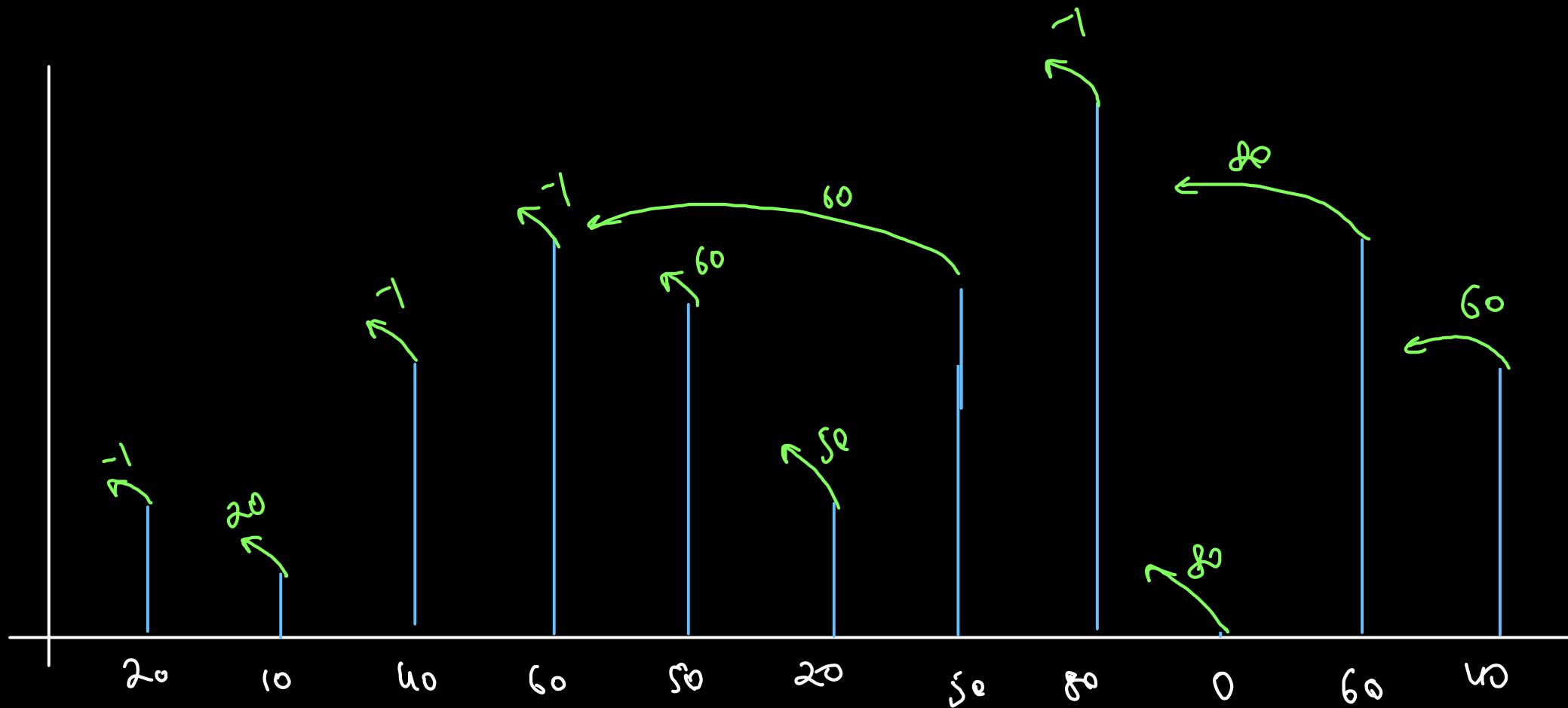
```

public static boolean validParantheses(String str){
    Stack<Character> stk = new Stack<>();
    for(int idx = 0; idx < str.length(); idx++){
        char ch = str.charAt(idx);
        if(ch == '('){
            stk.push('(');
        } else if(ch == ')'){
            if(stk.size() == 0) return false; // closing braces are extra
            stk.pop();
        }
        if(stk.size() > 0) return false; // opening braces are extra
    }
    return true;
}
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    String str = scn.next();
    System.out.println(validParantheses(str));
}

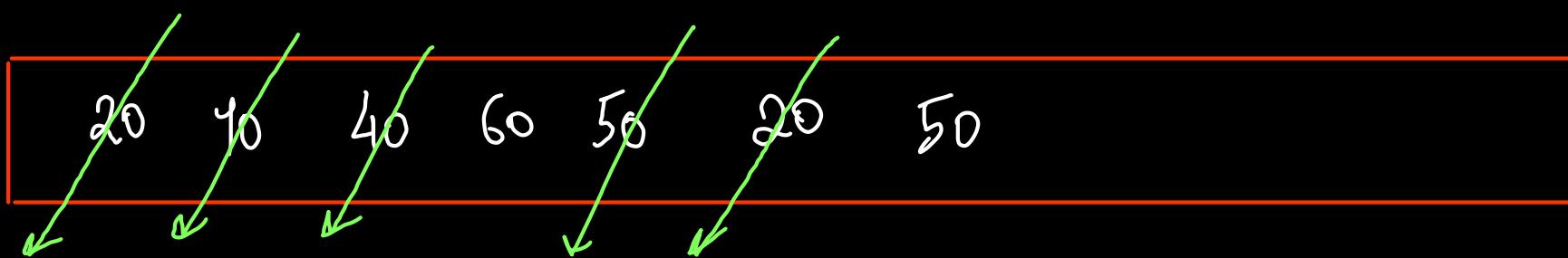
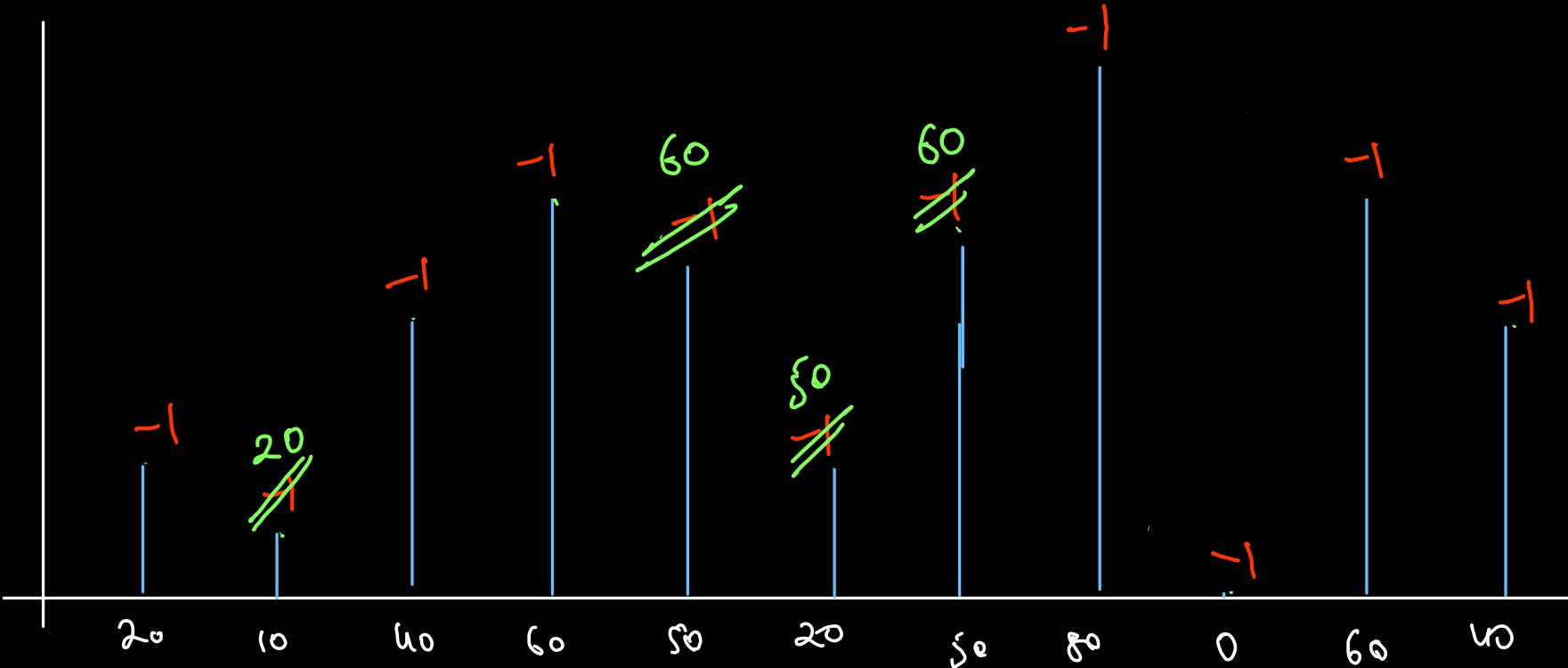
```



Merge Greater Elements to Left

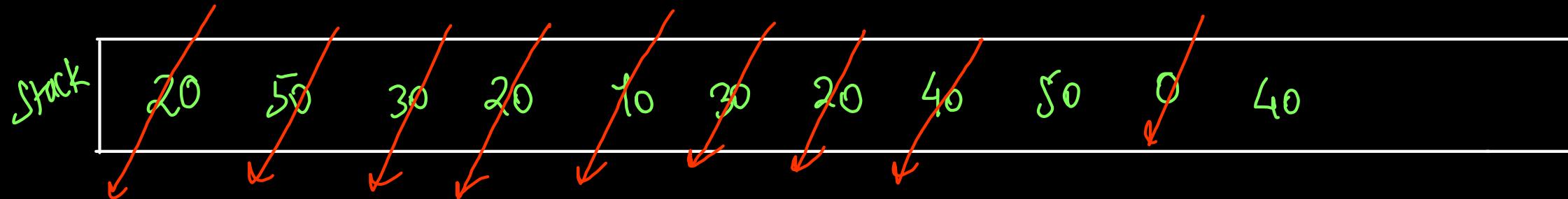
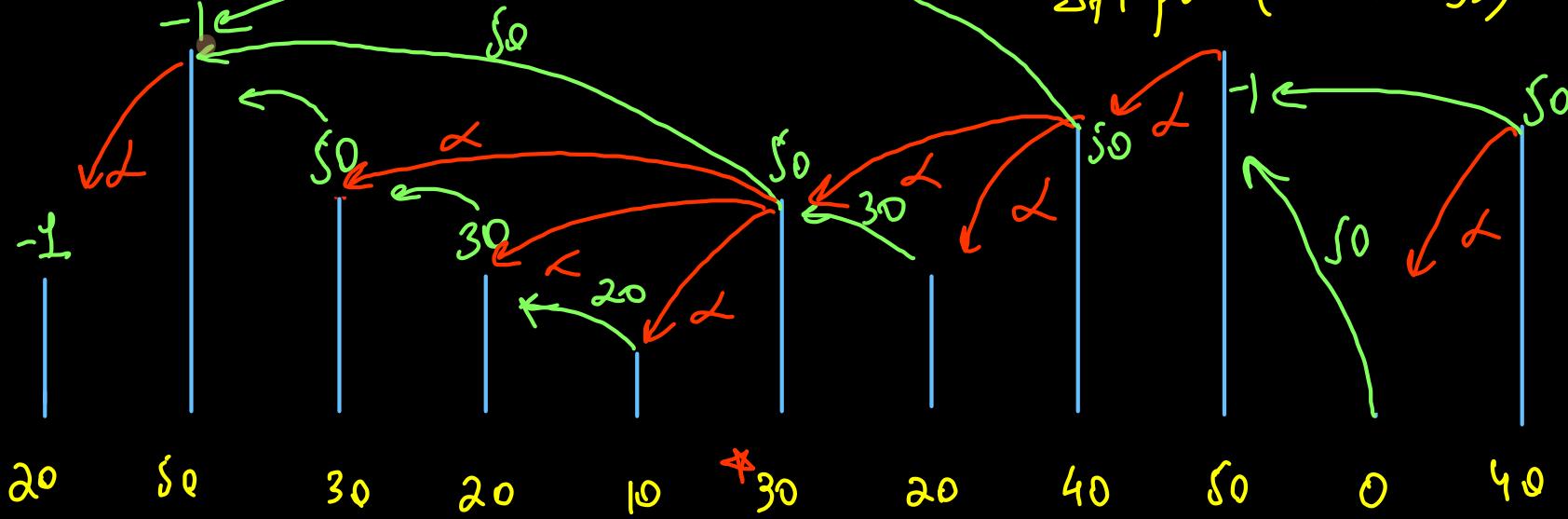


```
public static int[] nextGreaterElement(int[] arr){  
    int[] nge = new int[arr.length];  
  
    // Nested Loop: O(N ^ 2)  
    for(int i = 0; i < arr.length; i++){  
        nge[i] = -1;  
        for(int j = i - 1; j >= 0; j--){  
            if(arr[j] > arr[i]){  
                nge[i] = arr[j];  
                break;  
            }  
        }  
    }  
  
    return nge;  
}
```



Next Greater Element in Left

```
for(int idn=0; i<n ; i++) {  
    while(stk.size()>0 && arr[i]>=arr[idn])  
        stk.pop();  
    if(stk.size()==0) res[idn]=-1; else res[idn]  
        =stk.top();  
    stk.push(arr[idn]);
```



```

public static int[] nextGreaterElement(int[] arr){
    Stack<Integer> stk = new Stack<>();
    int[] res = new int[arr.length];

    for(int idx = 0; idx < arr.length; idx++){
        // Smaller or Equal Elements in left should be popped
        while(stk.size() > 0 && stk.peek() <= arr[idx]){
            stk.pop();
        }
        // Finding the Next Greater Element
        res[idx] = (stk.size() == 0) ? -1 : stk.peek();

        // Push the current element
        stk.push(arr[idx]);
    }

    return res;
}

```

$$O(n+n) = \underline{\underline{O(n)}}$$

linear

1 2 3 4

1,

4 3 2 1 5

✓ ✓ ✓ ✓

Amazon  
& Microsoft

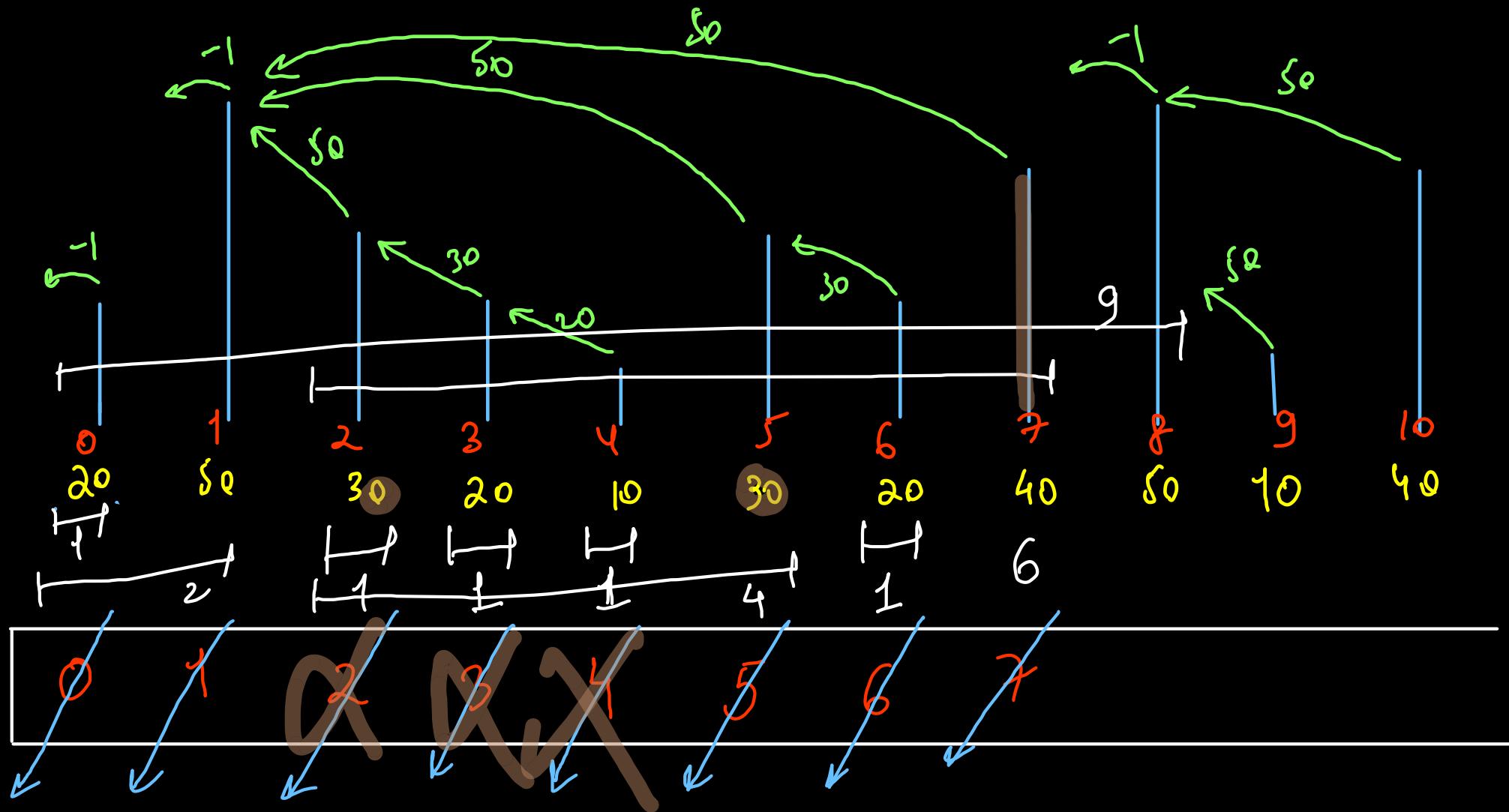
## Application of NGE Online Stock Spanner

Design an algorithm that collects daily price quotes for some stock and returns the **span** of that stock's price for the current day.

The **span** of the stock's price today is defined as the maximum number of consecutive days (starting from today and going backward) for which the stock price was less than or equal to today's price.

- For example, if the price of a stock over the next 7 days were **prices** = [100,80,60,70,60,75,85], then the stock spans would be **result** = [1,1,1,2,1,4,6].





Span = no. of consecutive days for which current height is max<sup>m</sup>  
= distance b/w current building & the NGE in left.

```
public static int[] stockSpan(int[] arr){
    Stack<Integer> stk = new Stack<>();
    int[] res = new int[arr.length];

    for(int idx = 0; idx < arr.length; idx++){
        // Pop the smaller or equal elements
        while(stk.size() > 0 && arr[stk.peek()] <= arr[idx]){
            stk.pop();
        }

        // Form the resultant (stock span)
        res[idx] = (stk.size() == 0) ? idx + 1 : idx - stk.peek();

        // Push the current index
        stk.push(idx);
    }

    return res;
}
```

# Simple Text Editor

Implement a simple text editor. The editor initially contains an empty string,  $S$ . Perform  $Q$  operations of the following 4 types:

1.  $\text{append}(W)$  - Append string  $W$  to the end of  $S$ .
2.  $\text{delete}(k)$  - Delete the last  $k$  characters of  $S$ .
3.  $\text{print}(k)$  - Print the  $k^{th}$  character of  $S$ .
4.  $\text{undo}()$  - Undo the last (not previously undone) operation of type 1 or 2, reverting  $S$  to the state it was in prior to that operation.

" "  $\stackrel{\textcircled{1}}{+}$   $\text{append}()$  "archit" = "archit"  $\stackrel{\textcircled{1}}{+}$  "aggarwal" = "Architaggarwal"  
↓  $\textcircled{2}$   $\text{delete}(5)$

" "  $\xleftarrow{\text{undo}}$  "archit"  $\xleftarrow{\text{undo}}$  "architaggarwal"  $\xleftarrow{\text{undo}}$  "architagg"  
 $\textcircled{3}$   $\text{print}(5) \Rightarrow t$

## Example

$S = \text{'abcde'}$

$ops = [1 \text{ fg}, 3 \text{ print}, 2 \text{ delete}, 4 \text{ undo}, 3 \text{ print}, 4 \text{ undo}, 3 \text{ print}]$

index	S	ops[index]	explanation
0	abcde	1 fg	append fg
1	abcdefg	3 6	print the 6th letter - f
2	abcdefg	2 5	delete the last 5 letters
3	ab	4	undo the last operation, index 2
4	abcdefg	3 7	print the 7th character - g
5	abcdefg	4	undo the last operation, index 0
6	abcde	3 4	print the 4th character - d

## Example

$S = 'abcde'$

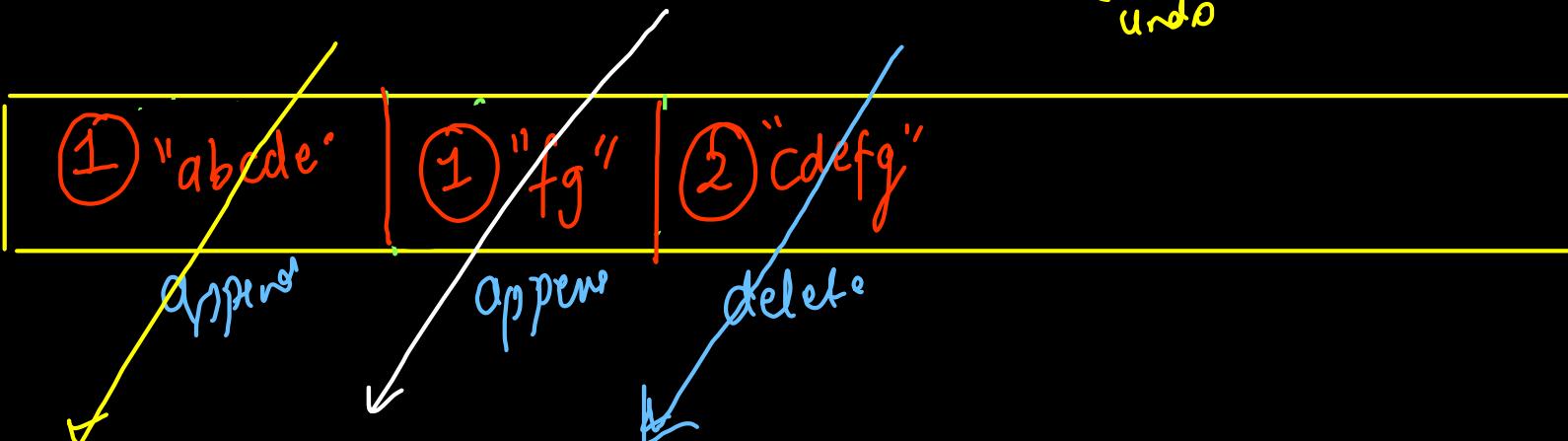
$ops = [1 \text{ fg}, 3 \text{ } \underline{6}, 2 \text{ } \underline{5}, 4, 3 \text{ } \underline{7}, 4, 3 \text{ } \underline{4}]$

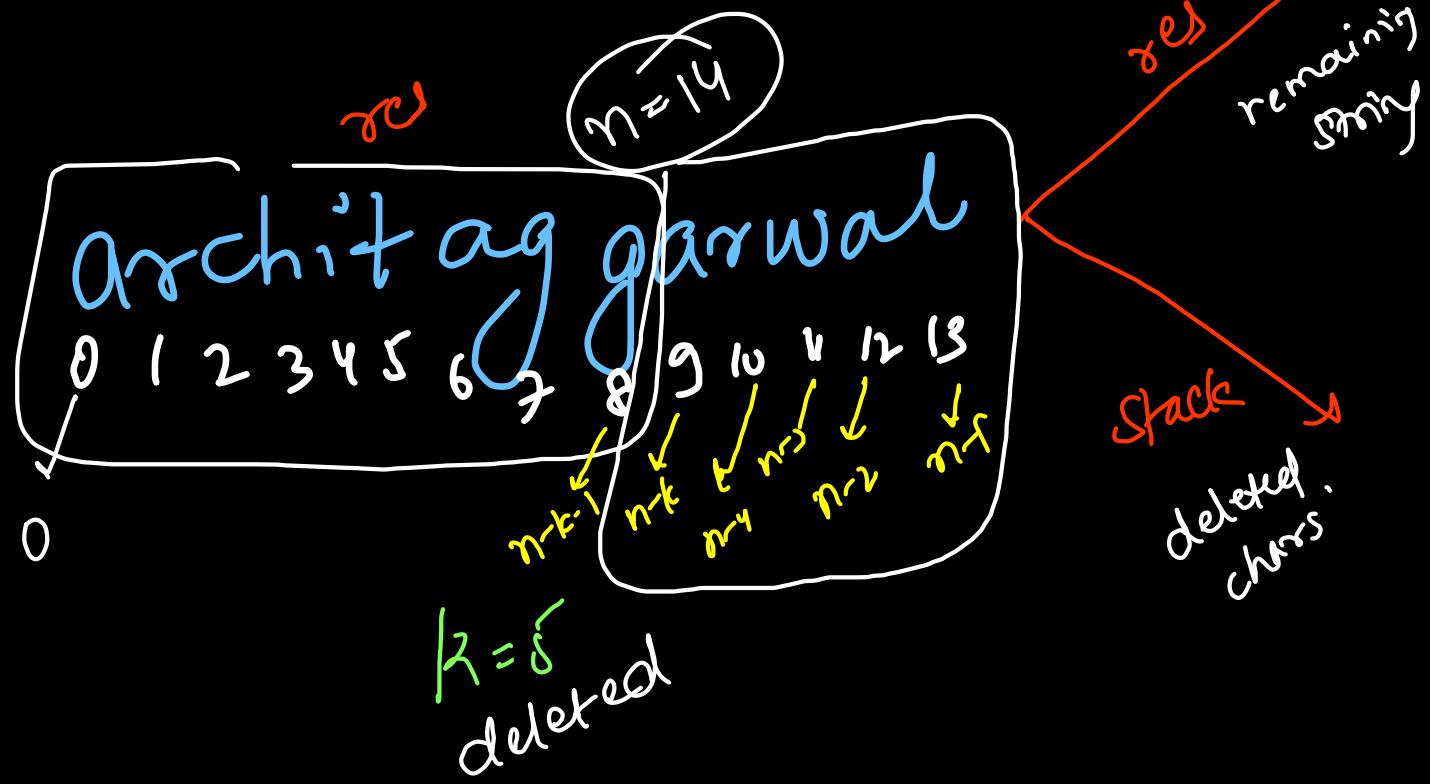
*print  
sb.charAt(i-1)  
delete*

StringBuilder {Stack → delete from last or append at last}

↳ "abcde" + "fg" = "abcdefg"  $\xrightarrow{\text{undo}}$  "ab" + "cdefg"  $\xrightarrow{\text{undo}}$  "abcde"  $\xrightarrow{\text{undo}}$  "abcde"

Stack:  
operations:





"architagg"

yes.substring( $0, n-k$ )

② "arw<sup>5</sup>al"

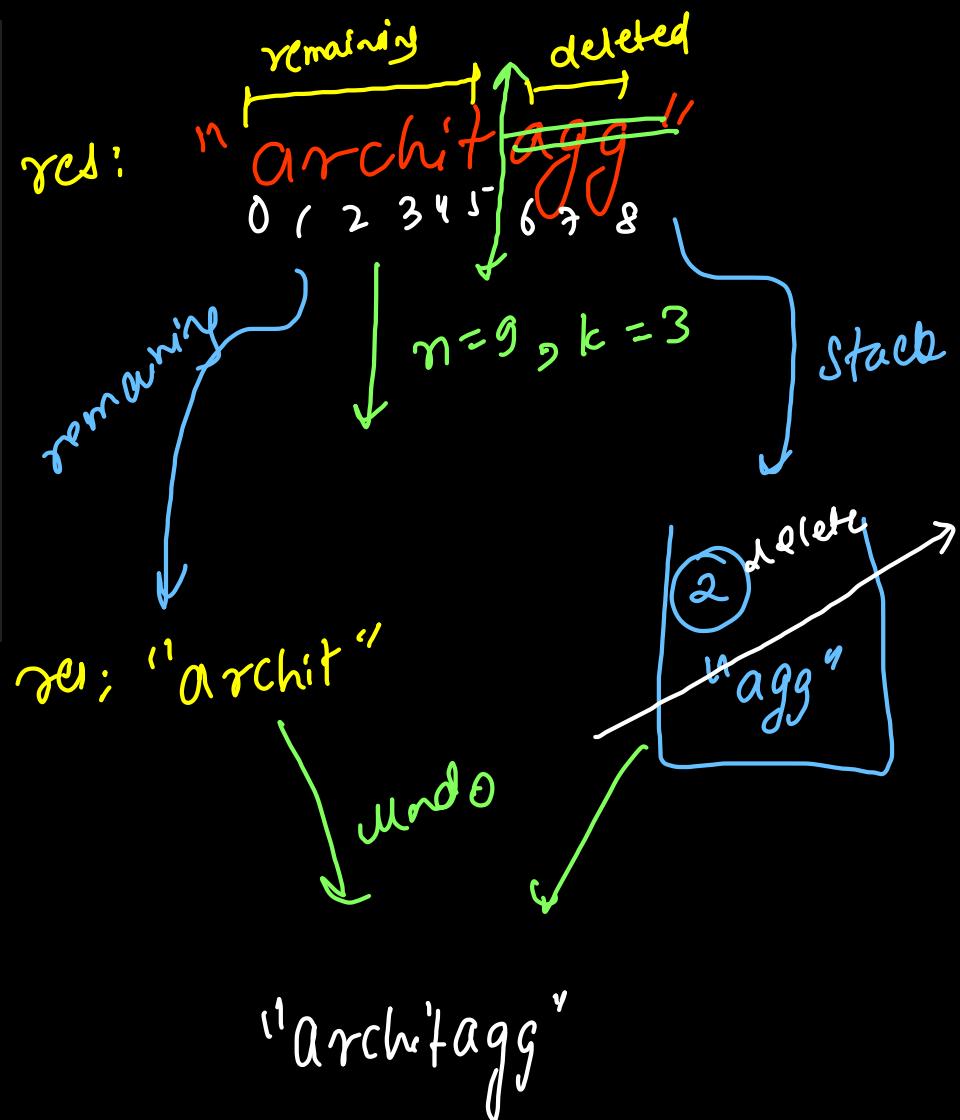
yes.substring( $n-k, n$ )

```

else if(choice == 2){
    // delete
    int k = scn.nextInt();
    int n = res.length();

    type.push(2); // delete operation
    undo.push(res.substring(n - k, n));
    stack
    res = res.substring(0, n - k);
}

```



```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int tests = scn.nextInt();

    String res = new String();
    Stack<Integer> type = new Stack<>(); // Type of Operation
    Stack<String> undo = new Stack<>(); // String which we want to do undo

    for(int t = 0; t < tests; t++){
        int choice = scn.nextInt();

        if(choice == 1){
            // append
            String str = scn.next();
            res += str;
            type.push(1); // insert operation

            undo.push(str);
        }
        else if(choice == 2){
            // delete
            int k = scn.nextInt();
            int n = res.length();

            type.push(2); // delete operation
            undo.push(res.substring(n - k, n));

            res = res.substring(0, n - k);
        }
    }
}

```

```

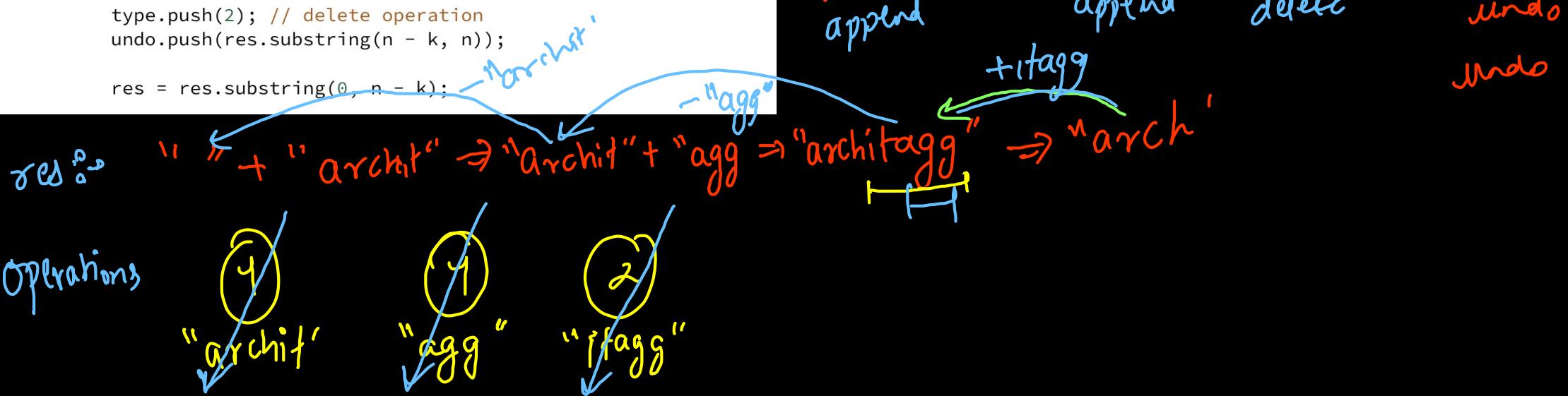
} else if(choice == 3){
    // print
    int k = scn.nextInt(); // 1-based indexing
    System.out.println(res.charAt(k - 1));

} else if(type.pop() == 1){
    // undo insert: delete
    String str = undo.pop();
    int n = res.length();
    int k = str.length();

    res = res.substring(0, n - k);
}

else {
    // undo delete: append
    String str = undo.pop();
    res += str;
}

```



`str "archit aggarwal is educator"`

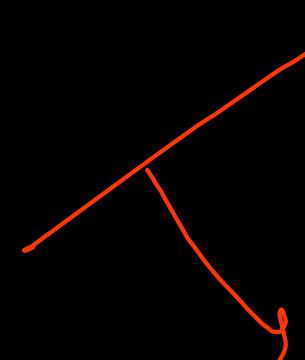
$\downarrow$  `str.split(" ")`;

`{ "archit", "aggarwal", "is", "educator" }`

`str: " / / / archit - aggarwal / / / - is / / / "`

$\downarrow$  `str.lstrip()`

`" archit - aggarwal - is "`

Prefix	Infix	Postfix	Evaluation
$((5 + 2 * 8) - (9 / 4) + 7))$ infix expression brackets Operand1    operator    Operand2 2              *              8		 prefix operator op1 op2 * 2 8	postfix op1 op2 operator 2 8 *

infix

"a+b\*(c+d-e)+f+g\*h)-i"

postfix

abcd+e-fgh\*+\*+i-



humans ↗

$6 + 5 * (7 + 8 - 2) + (1 / 2 * 3) - 5$



6, 5, 7, 8, +, 2, -, 1, 2, 3, \*, /, +, \*, +, -,

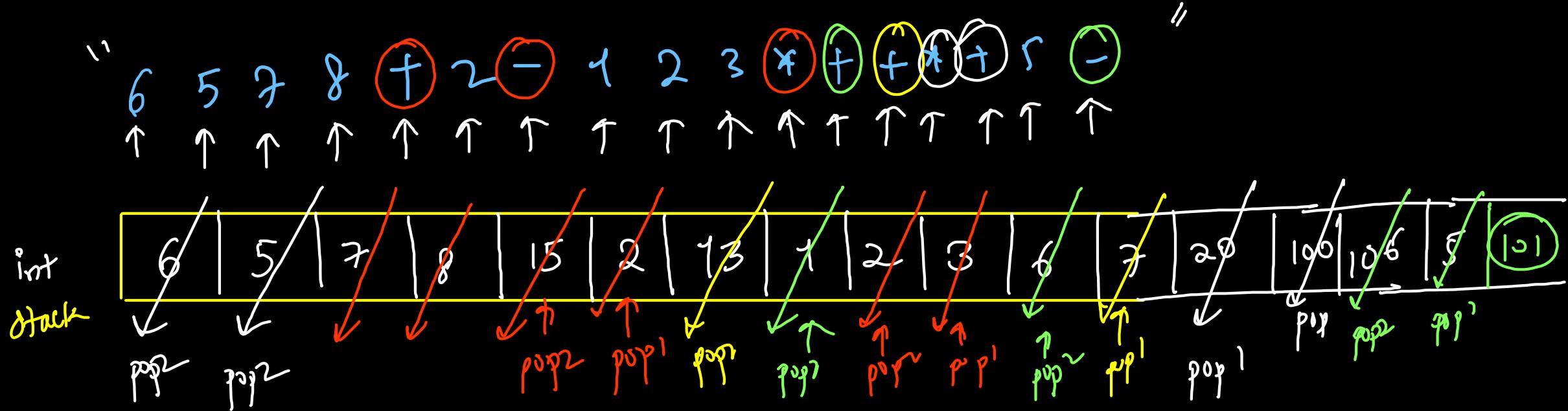


algo(stack)

Reverse Polish Notation

# Postfix Expression Evaluation

type → 0-9



$$6 + 5 * ((7 + 8 - 2) + (1 + 2 * 3)) - 5$$

$\gamma = 20$

Diagram showing the expression evaluated in two main parts, each enclosed in an oval. The first part is  $(7 + 8 - 2)$  and the second part is  $(1 + 2 * 3)$ . The result of the first part is labeled "humans" with an arrow pointing to it.

```

if(ch > '0' & ch <= '9')
    Stk.push(ch - '0');

else
    int pop1 = Stk.pop();
    int pop2 = Stk.pop();
    Stk.push(pop2 op pop1)

```

```
public static int evaluatePostfix(String str){  
    Stack<Integer> stk = new Stack<>();  
    Extra space  
    for(int idx = 0; idx < str.length(); idx++){  
        char ch = str.charAt(idx);  
  
        if(ch >= '0' && ch <= '9'){  
            stk.push(ch - '0');  
        } else {  
            int pop1 = stk.pop();  
            int pop2 = stk.pop();  
  
            if(ch == '+') stk.push(pop2 + pop1);  
            else if(ch == '-') stk.push(pop2 - pop1);  
            else if(ch == '/') stk.push(pop2 / pop1);  
            else stk.push(pop2 * pop1);  
        }  
    }  
    return stk.peek();  
}
```

Time :  $\rightarrow O(n)$   
Space :  $\rightarrow O(n)$

# HashMap Data Structure

key → no multiple values

key : value

String      Integer  
IPLTeam : IPLTrophies

CSK : 4

KKR : 2

MI : 5

RCB : 0

SRH : 2

String      String  
Country : Capital  
India : Delhi

Pakistan : karachi

UK : London

US : Washington

Russia : moscow

Japan : Tokyo

Integer      String  
CountryCode : Country

g1 : India

g2 : USA

g3 : Pakistan

g4 : UAE

g5 : Malaysia

Not possible

{ India : Delhi, Kolkata }  
US  
{ India : Delhi, Kolkatta }  
{ India : Kolkata, Kolkatta }

① 2 keys cannot be same  
All keys will be unique /  
distinct keys |

② 2 different tetx can have  
same values

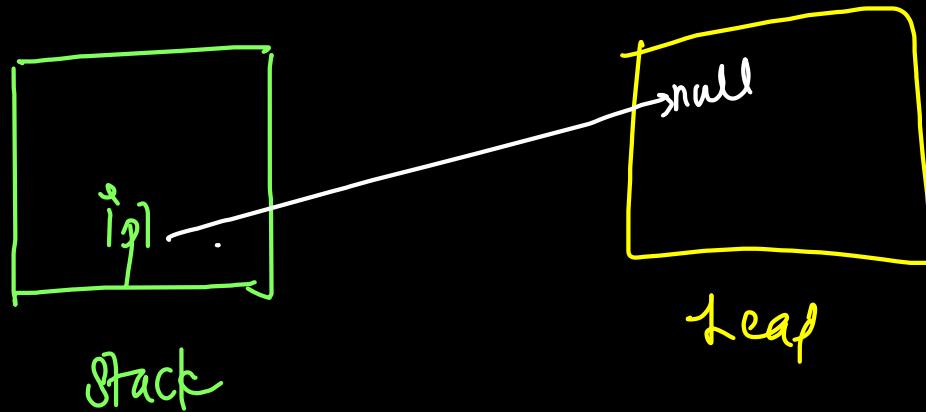
③  $\text{hashmap.get(key)}$   
 $\text{key} \rightarrow \text{value}; O(1) \text{ avg}$   $\Rightarrow$  What is capital time of India?  
" " " " " Japan?  
  
 $\text{hashmap.getkey(value)}$   
 $\text{value} \rightarrow \text{key}; O(n)$   $\Rightarrow$  what is country with capital Karadu?  
" " " " " Moscow?

```
// Read: Get the Value of a Key  
System.out.println(ipl.get("CSK"));  
System.out.println(ipl.get("MI"));  
System.out.println(ipl.get("RCB"));  
System.out.println(ipl.get("PWI")); // null
```

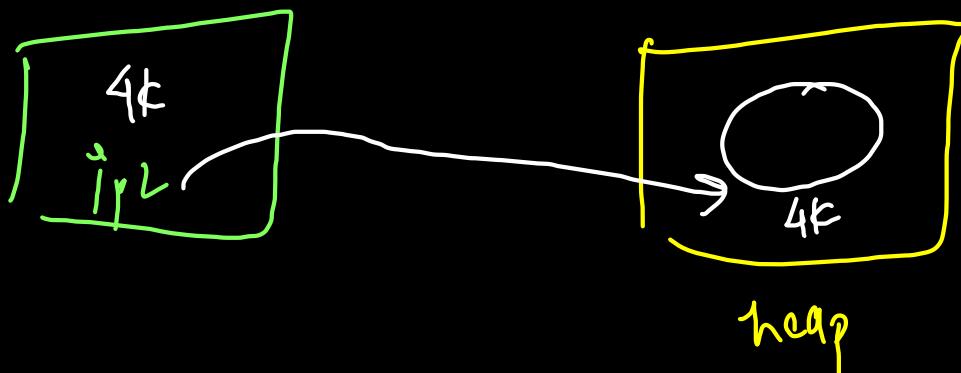
key is present  
↳ value

→ If Key is not present  
↳ null

```
HashMap<String, Integer> ipl = new HashMap<>();
```



```
HashMap<String, Integer> ipl = new HashMap<>();
```



```
// Creation
HashMap<String, Integer> ipl = new HashMap<>();

System.out.println(ipl.size());
System.out.println(ipl);

// Insertion: Key-Value Pair
ipl.put("RCB", 0);
ipl.put("LSG", 0);
ipl.put("MI", 5);
ipl.put("RR", 1);
ipl.put("CSK", 4);
ipl.put("PK", 0);
ipl.put("DC", 0);
ipl.put("GT", 1);
ipl.put("KKR", 2);
ipl.put("SRH", 2);

System.out.println(ipl);
// 1. Insertion Order: Incorrect
// 2. Sorted Order on Keys: Incorrect
// 3. Sorted Order on Values: Incorrect
// 4. I Dont Know: HashCode: Random Order: Correct

System.out.println(ipl);
System.out.println(ipl.size());

// Read: Get the Value of a Key
System.out.println(ipl.get("CSK"));
System.out.println(ipl.get("MI"));
System.out.println(ipl.get("RCB"));
System.out.println(ipl.get("PWI")); // null
System.out.println(ipl.get(null));

// ipl.put(null, 1); // Key: null Value: 1
// System.out.println(ipl.get(null));
```

`ipl.put("CSk", 4)  $\Rightarrow$  Insert`

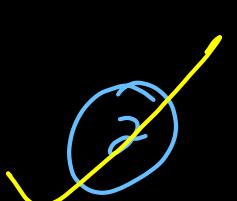
"CSk" : 4

`ipl.put("CSk", 5)  $\Rightarrow$  Update`



"CSk": 4

"CSk": 5



"CSk" : 5

String  
state

ArrayList<String>  
cities

Punjab

Ambala, Ludhiana, Jalandhar,  
Chandigarh, --

Haryana

Kohitak, Sirsa, Bhiwani, Faridabad,--  
Gurugram,

UP

Gorakhpur, Lucknow, Ayodhya

Hyderabad

Gachibowli, Raidurg, Jubilee Hills  
Kukatpally,

```
// Traversal: For Each on Key  
  
for(String team : ipl.keySet()){  
    System.out.print(team + " : ");  
    System.out.println(ipl.get(team));  
}
```

## Word meaning

You are required to create a **dictionary** consisting of word and its meaning.

Take an integer N as input and **Continue** the process until **Case 4** is not achieved.

- If **N==1**, take **word** and **meaning** as input from user and add it to the dictionary.  
*//insert : dict.put(word, meaning)*
- If **N==2**, take a **word** as input from the user and print its **meaning**, if the word is not found print -1.
- If **N==3**, take a **word** as input from the user and delete it from the dictionary.  
*dict.remove(word);*
- If **N==4**, Close the dictionary(Exit the program).

"apple" : "a fruit which is red"  
String : String

{  
dict.getOrDefault("apple", "-1"),  
↳ "a fruit which is red"  
dict.getOrDefault("orange", "-1");  
↳ "-1",

```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);

    // EMPTY HASHMAP CREATION
    HashMap<String, String> dict = new HashMap<>();

    while(true){
        int choice = scn.nextInt();

        if(choice == 1){
            // insert
            String word = scn.next();
            String meaning = scn.next();
            dict.put(word, meaning); →  $\Theta(1)$  avg  $O(n)$  worst
        } else if(choice == 2){
            // get or default
            String word = scn.next();
            System.out.println(dict.getOrDefault(word, "-1"));
        } else if(choice == 3){
            // delete
            String word = scn.next();
            dict.remove(word);
        } else break; // exit the while loop or program
    }
}

```

menu driven

1 ✓
Geekster ✓
Coding ✓
1 ✓
Geek ✓
Coder ✓
2 ✓
Geek ✓
3 ✓
Geekster ✓
2 ✓
Geekster ✓
4 ✓

String : String  
~~Geekster is Coding~~  
 Geek & Coder

Output:  
 Coder  
 ==  
 -1

key present  
↳ returns value

get(key)

ipl.get("CSk") : 4

ipl.get("M1") : 5

key not present  
↳ null

ipl.get("Pw1") : null

ipl.get("garbage") : null

key present  
↳ returns value

getOrDefault(key, def)

ipl.getOrDefault("CSk", def) : 4

ipl.getOrDefault("M1", def) : 5

key not present  
↳ default value

ipl.getOrDefault("Pw1", 0) : 0

TreeMap : Hashmap with sorted keys  
 ↪ lexicographical order

A r c h i t   A g g a r w a l @ 2 0 2 2  
 ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑ ↑↑

if not present

freq.put(ch, 1)

else

freq.put(ch, freq.get(ch) + 1);

Character, Integer

key                  value

A : 1+1      g : 1+1

r : 1+1      a : 1+1

c : 1      w : 1

h : 1      l : 1 r

i : 1      @ : 1

t : 1      '2' : 1+1+1  
 '0' : 1

```

public static void printFrequency(char[] arr){
    TreeMap<Character, Integer> freq = new TreeMap<>();
    for(char ch: arr){ O(n)
        if(freq.containsKey(ch) == false){ O(1)
            // if(freq.get(ch) == null){
                freq.put(ch, 1); //insert
            } else { O(1)
                int oldFreq = freq.get(ch);
                freq.put(ch, oldFreq + 1); //update
            }
        }
        O(unique keys) = O(128) → ASCII code
    for(char ch: freq.keySet()){
        System.out.println(ch + " " + freq.get(ch));
    }
}

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);

    int n = scn.nextInt();
    |
    char[] arr = new char[n];
    for(int idx = 0; idx < n; idx++){
        arr[idx] = scn.next().charAt(0);
    }
    printFrequency(arr);
}

```

ignore TreeMap (assume hashmap)

Total Time

Time g:07 - g:08

$O(n)$

containsKey(ch)

↳ true: key is present in HM  
 $\text{impl.ck("Cjk") = true}$

↳ false: key is not present in HM  
 $\text{impl.ck("pw1") = false}$

```

public static void printFrequency(char[] arr){
    TreeMap<Character, Integer> freq = new TreeMap<>();

    for(char ch: arr){
        if(freq.containsKey(ch) == false){
            // if(freq.get(ch) == null){
                freq.put(ch, 1);
            } else {
                int oldFreq = freq.get(ch);
                freq.put(ch, oldFreq + 1);
            }
        }

        for(char ch: freq.keySet()){
            System.out.println(ch + " " + freq.get(ch));
        }
    }

    public static void main(String[] args) {
        Scanner scn = new Scanner(System.in);

        int n = scn.nextInt();
        |

        char[] arr = new char[n];
        for(int idx = 0; idx < n; idx++){
            arr[idx] = scn.next().charAt(0);
        }

        printFrequency(arr);
    }
}

```

$\text{freq.put}(ch)$   
 $\text{freq.getOrDefault}(ch, 0) + 1$

$A : \text{not present} :$   
 $1 \Rightarrow 0 + 1 = 1$

$A : \text{inserted already}$   
 $1 \Rightarrow 1 + 1 = 2$

```
public static void printFrequency(char[] arr){  
    TreeMap<Character, Integer> freq = new TreeMap<>();  
  
    for(char ch: arr)          ↗ O(1) avg  
        freq.put(ch, freq.getOrDefault(ch, 0) + 1);  
    ↘ O(n) avg  
    for(char ch: freq.keySet())  
        System.out.println(ch + " " + freq.get(ch));  
}  
  
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
  
    int n = scn.nextInt();  
  
    char[] arr = new char[n];  
    for(int idx = 0; idx < n; idx++){  
        arr[idx] = scn.next().charAt(0);  
    }  
  
    printFrequency(arr);  
}
```

# Same No Same Frequency

4	-3	2	0	-3	4	2	2	4	-3	4	3
↑	↑	↑	↑	↑	↑	↑	↑	↑	↑	↑	↑

4 : 4 ✓

-3 : 3 ✓

0 : 1 ✗

3 : 1 ✗

2 : 3 ✗

hashmap → random keys

treenmap → sorted keys

↳ Ascending  
order

int : int

```
public static void printFrequency(int[] arr){  
    TreeMap<Integer, Integer> freq = new TreeMap<>();  
  
    for(int val: arr){  
        freq.put(val, freq.getOrDefault(val, 0) + 1);  
    }  
    for(int key: freq.keySet()){  
        if(freq.get(key) == Math.abs(key)){  
            System.out.println(key);  
        }  
    }  
}
```

// matching  
the  
expected  
output ]  
(sorted order)

\* Longest Substring w/o Repeating  
Characters

Salesforce

18 base

25k USD → 4 years

Spinifex

35 LPA base  
      

Stock X

Largest Substring w/o Repeating Characters

"architsharmadev" 

Character → Integer  
 $a \rightarrow 1$

- ① contiguous
  - ② max length
  - ③ no duplicate
- length

to exclude characters  $\leftarrow$  left  $\leq$  right  $\rightarrow$  to include characters

Sliding window  
↳ Two Pointer

+

hashmap

Character → Integer

a → 4012121

g → 4010

c → 401

h → 4210

$j \rightarrow 1/0$   
 $t \rightarrow 1/0$   
 $S \rightarrow 1/0$   
 $m \rightarrow 1/0$

$$t \rightarrow 1$$

d → 1

从

$$\begin{matrix} \mu \rightarrow 1 \\ \theta \rightarrow 1 \\ \gamma \rightarrow 1 \end{matrix}$$

```
public int lengthOfLongestSubstring(String s) {  
    int left = 0, maxlen = 0;  
    HashMap<Character, Integer> freq = new HashMap<>();  
  
    for(int right = 0; right < s.length(); right++){  
        char ch = s.charAt(right);  
        freq.put(ch, freq.getOrDefault(ch, 0) + 1); // right -> include  
  
        while(freq.get(ch) > 1){  $\rightarrow O(n)$   
            char chl = s.charAt(left);  
            freq.put(chl, freq.getOrDefault(chl, 0) - 1); // left -> delete  
            left++;  
        }  
  
        maxlen = Math.max(maxlen, right - left + 1);  
    }  
  
    return maxlen;  
}
```

including

discarding

$O(n)$

Time :  $O(n)$

$n+n$   
↑  
independent loops  
discard

Geekster  $\Rightarrow$  Jobs  $\rightarrow$  Indian startups  
low headcount ( $\leq 100$  employees)

Recession

$\downarrow$

US based companies

$\nearrow$  hiring  
 $\nearrow$  freeze  
 $\nearrow$  fire

$\nearrow$

Service based companies

$\nearrow$  Wipro  
 $\neararrow$  Cognizant

First Unique Character

leftmost index      non-repeating

"geek s far geeks" : 5<sup>th</sup> index

0 1 2 3 4 5 6 7 8 9 10 11 12

↑↑↑↑↑↑↑

"geeks geeks" (-1)

↑↑↑↑↑↑↑↑↑↑↑↑

Two Traversal Approach

① Fill frequency hashmap

g → 1/2      f → 1

e → 1/2/3/4      a → 1

k → 1/2      r → 1

s → 1/2

② Leftmost char: freq(ch) = 1

```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    String str = scn.next();
    extra space
    HashMap<Character, Integer> freq = new HashMap<>();

    for(int idx = 0; idx < str.length(); idx++){
        char ch = str.charAt(idx);
        freq.put(ch, freq.getOrDefault(ch, 0) + 1);
    }

    for(int idx = 0; idx < str.length(); idx++){ // leftmost
        char ch = str.charAt(idx);
        if(freq.get(ch) == 1){ // unique
            System.out.println(idx);
            return;
        }
    }

    System.out.println(-1);
}

```

Two Traversal  
Approach

Time  $\Rightarrow O(n+n) = O(n)$

Space  $\Rightarrow O(128)$

hashmap

# First Repeating Character

```
char firstRep(String str)
{
    HashMap<Character, Integer> freq = new HashMap<>();
    char ans = '#';

    for(int idx = str.length() - 1; idx >= 0; idx--){
        char ch = str.charAt(idx);
        freq.put(ch, freq.getOrDefault(ch, 0) + 1);
        if(freq.get(ch) > 1) ans = ch;
    }

    return ans;
}
```

"abccaa"  
↑↑↖↖↑  
b:1  
a:2  
c:2

"abcddccef"  
↑↑↑↑↑↑↑↑↑  
Ans = '#', 'd', 'c'

f: 1	b: 1
c: 1	a: 1
d: 2	
e: 1	

## Highest Palindrome Length

- pick any character ✓ (need not be contiguous)
- rearrange characters ✓ (order may not be maintained)

a b c d a c e f

a b c d a c e f b a c

"a c (B) c a"

a b c (d) c b a

"aaa bbb ccc ddd"

$$a: k_2 \not\sim k_4 \quad b: k_2 \sim k_3 \quad c: k_2 \not\sim k_4 \quad d: k_2 \not\sim k_3$$

"↓↓↓↓↓ ( ↓↓↓↓↓"

a: 1 ✓  
b: 1 ✓  
c: 1 ✓

a<sup>~</sup> a<sup>-</sup> b<sup>\backslash</sup> c<sup>\backslash</sup> d<sup>b</sup> d<sup>\backslash\backslash</sup> c<sup>\backslash</sup> b<sup>\backslash\backslash</sup> a<sup>a</sup>  
odd length

$$\text{len} = 0 + 4 + 2 + 4 + 2 + 1$$

a b c d d d c b a  
✓ ✓ ✓ ✓ ✓ ✓ ✓ ✓ ✓ ✓  
even length

$$\text{len} = 0 + 2 + 2 + 2 + 4 \\ = \textcircled{10} = n$$

geeks rfy skeeg  
geeks for geeks rr  
↑↑↑↑↑↑↑↑↑↑↑↑↑↑

- ✓  $g : y_2$
- ✓  $c : y_2 \neq 4$
- ✓  $k : x_2$
- ✓  $s : x_2$
- ✓  $f : 1$
- 0 : 1
- v :  $\beta_2$

$$\{e_n = 0 + 2 + 4 + 2 + 2 + 2 + 1\}$$

# AVRL Company

g CTC { go + + → Base }

Front End Engineer (React)

1st Interviews

Web Dev → 2 hours

↳ API → CRUD

Javascript

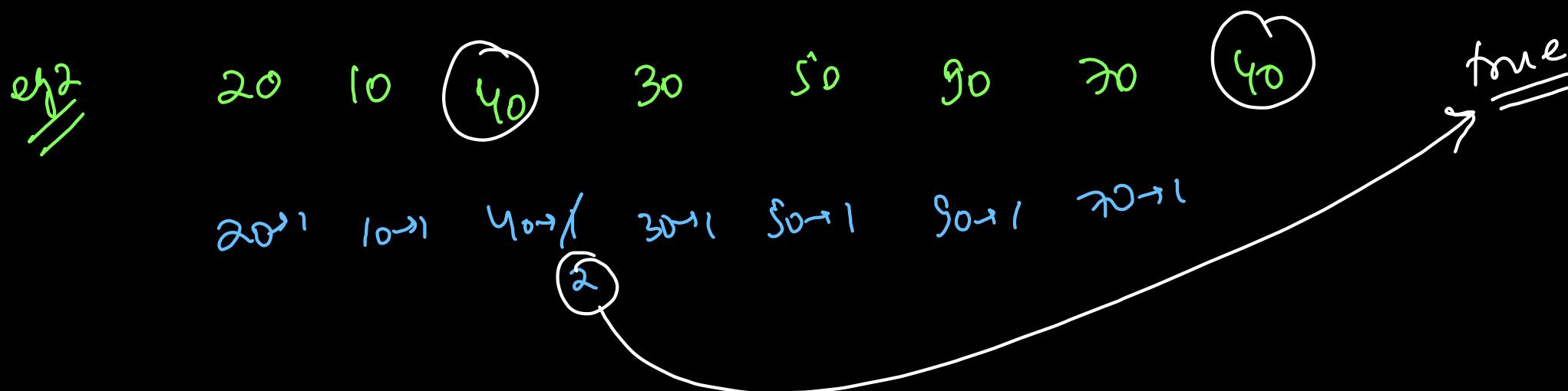
2nd Interviews

→ Reverse words leetcode

→ Merge 2 sorted Array

HR → Interviewbit HR Questions

Contains Duplicate ?



```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();

    int[] arr = new int[n];
    for(int idx = 0; idx < n; idx++){
        arr[idx] = scn.nextInt();
    }

    HashMap<Integer, Integer> freq = new HashMap<>();

    for(int key: arr){
        freq.put(key, freq.getOrDefault(key, 0) + 1);
        if(freq.get(key) >= 2){
            System.out.println(true);
            return;
        }
    }
    System.out.println(false);
}
```

Time  $\Rightarrow O(n)$   
Linear

Space  $\Rightarrow O(n)$   
~~hashmap~~

## Unique Number of Occurrences

ex:

10 20 30 10 10 10 20 30 40 20

10: 5  
20: 3  
30: 2  
40: 1

Value is unique

hashmap : freq

there are no 2 keys

having same frequency

⇒ true

10

20

30

40

map[10]

20

30

40

Integers | Integer  
in hm

5 → 10  
3 → 20  
2 → 30  
1 → 40

eg 2

10 10 10 10 20 30 30 40 50 40 20 20

10: 4

20: 3

30: 2

40: 2

50: 1

same frequency  
or false

4 → 10

3 → 20

2 → 30, 40

false

```

public static boolean isDuplicate(int[] arr){
    HashMap<Integer, Integer> freq = new HashMap<>();
    for(int key: arr){ → O(n)
        freq.put(key, freq.getOrDefault(key, 0) + 1);
    }

    for(int key1: freq.keySet()){
        int val1 = freq.get(key1);

        for(int key2: freq.keySet()){
            int val2 = freq.get(key2);
            if(key1 != key2 && val1 == val2)
                return false;
        }
    }

    return true;
}

```

$O(n^2)$

26 10 20 10 20 30 40



key1 = 20  
val1 = 3  
key2 = 10  
val2 = 2

key1 = 10  
val = 2

key2 = 20, 30

key1 = 30  
val1 = 1  
key2 = 40  
val2 = 1

```

public static boolean isDuplicate(int[] arr){
    HashMap<Integer, Integer> freq = new HashMap<>();
    for(int key: arr){
        freq.put(key, freq.getOrDefault(key, 0) + 1);
    }

    HashMap<Integer, Integer> inv = new HashMap<>();

    for(int key: freq.keySet()){
        int val = freq.get(key);
        if(inv.get(val) != null) return false;
        inv.put(val, key);
    }

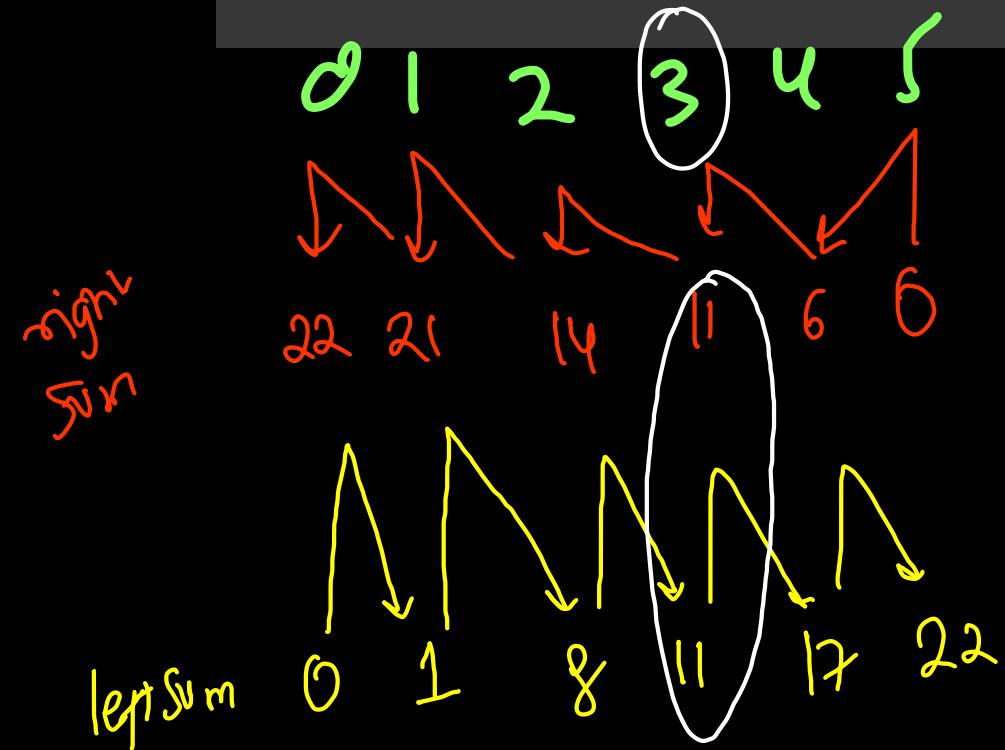
    return true;
}

```



Find Pivot Index

[1, 7, 3, 6, 5, 6]



## Queue Data Structure

- Array (1D, 2D) ⇒ similar data in contiguous ⇒ random access (index)  
⇒ fixed size
- Strings & String Builder ⇒ array of characters
- ArrayList ⇒ array of dynamic size ⇒ random access
- Stack ⇒ FILO (First-In Last-Out) or LIFO (Last In - First Out)  
⇒ Both insertion & deletion from same end (top) 
- HashMap & TreeMap ⇒ store key-value pair  
⇒ unique keys ⇒  $\text{hashmap}(\text{key} \& \text{value})$  ⇒ random  
⇒ search/insert/remove ((key)) ⇒ constant avg

# Queue Data Structure



Sharma



First In



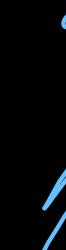
Verna



Aggarwal



Arman



Rishal



Last In

Last Out

(LIFO)

Removal

First Out

(FIFO)

#

First Come

First Serve Basis

→ Circular Queue  
(Using Array)

→ Linked List

insertion

# Insertion & Removal  
from the different  
ends

```
Queue<String> q = new ArrayDeque<>();  
  
// Insert  $\rightarrow O(1)$   
q.add("Sharma");  
System.out.println(q); [Sharma]  
  
q.add("Verma");  
System.out.println(q); [S, V]  
  
q.add("Arman");  
System.out.println(q); [S, V, A]  
  
q.add("Vishal");  
System.out.println(q); [S, V, A, Vi]  
  
// Get (Read): Front Element  $\rightarrow O(1)$   
System.out.println(q.peek());  $\rightarrow$  Sharma  
System.out.println(q);  $\rightarrow$  [S, V, A, Vi]
```

$O(1)$

Sharma

[Ve, A, Vi]

Verma

System.out.println(q.remove());

System.out.println(q); [A, Vishal]

Arman

System.out.println(q.remove());

System.out.println(q); (Vishal)

Vishal

System.out.println(q.remove());

System.out.println(q); []

System.out.println(q.size())

$O(1)$

queue<integer> q = new ArrayQueue<X>;

1. Declare an Empty queue  $s$ .
2. Take Single Integer  $T$  as input.
3. For next  $T$  Lines format (*case, x(optional)*)

- case 1. *Print* the size of the queue in a separate line.  $\rightarrow q.size()$
- case 2. *Remove* an element from the queue. If the queue is empty then print  $-1$  in a separate line.  $\rightarrow q.remove();$   $\rightarrow$  queue underflow
- case 3. *Add* Integer  $x$  to the queue  $s$ .  $\rightarrow q.add(x);$
- case 4. *Print* an element at the *front* of the queue. If queue is empty print  $-1$  in a seperate line.  $\rightarrow q.peek();$   $\rightarrow$  queue underflow

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    Queue<Integer> q = new ArrayDeque<>();

    int test = scn.nextInt();
    while(test-- > 0){
        int choice = scn.nextInt();

        if(choice == 1){
            // size
            System.out.println(q.size());

        } else if(choice == 2){
            // remove
            if(q.size() == 0) System.out.println(-1);
            else q.remove();

        } else if(choice == 3){
            // insert
            int val = scn.nextInt();
            q.add(val);

        } else {
            // get/peek
            if(q.size() == 0) System.out.println(-1);
            else System.out.println(q.peek());
        }
    }
}
```

test=5 for (int idn=0; idn < r; idn++)

0 → 4

while (test -- > 0)

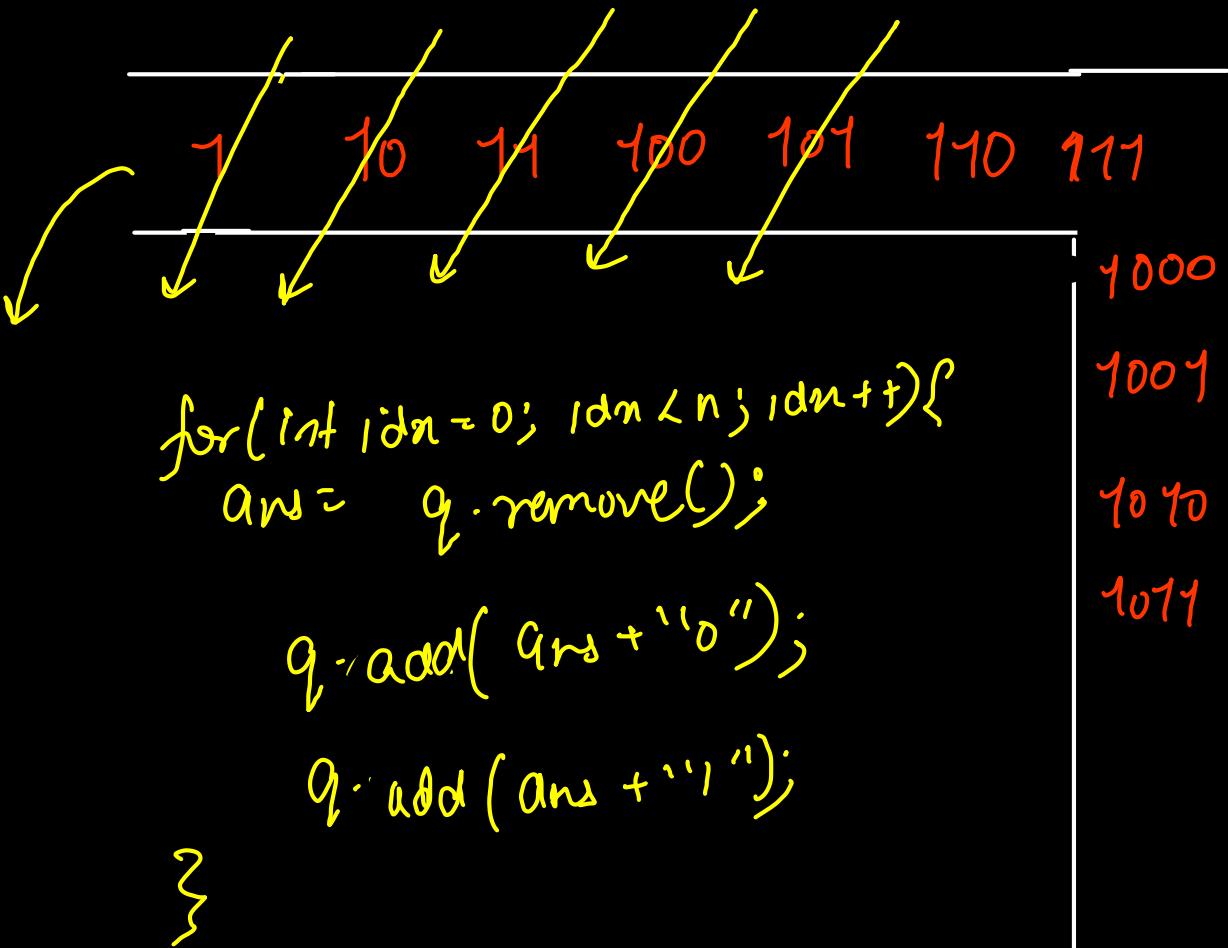
5 → 1

int []  
char []  
String []

Queue Print Binary { BFS }

All from 1 to N in increasing order

1	1 ✓	9	1001
2	10 ✓	10	1010
3	11 ✓	11	1011
4	100 ✓	12	1100
5	101 ✓	13	1101
6	110 -	14	1110
7	111 -	15	1111
8	1000 -	16	10000
		=n	



```

static ArrayList<String> generate(int N)
{
    ArrayList<String> res = new ArrayList<>();
    Queue<String> q = new ArrayDeque<>();
    q.add("1");

    while(N-- > 0){
        String ans = q.remove(); → O(1)
        res.add(ans); → O(1)

        q.add(ans + "0"); → O(1)
        q.add(ans + "1"); → O(1)
    }

    return res;
}

```

Time  
 $O(n)$   
 ignoring  
concatenation

Queue Problem

Generate Binary  
 Nos

from 1 to N.

{ GFG }

[1]
[10, 11]
[11, 100, 101]
[100, 101, 110, 111]
[101, 110, 111, 1000, 1001]
[110, 111, 1000, 1001, 1010, 1011]
[111, 1000, 1001, 1010, 1011, 1100, 1101]

## Implement a Queue { Array Implementation }



q.add(10);

q.add(20);

q.add(30);

q.add(40);

q.add(50)

q.add(60)

q.add(70);

q.add(80)

q.add(90);

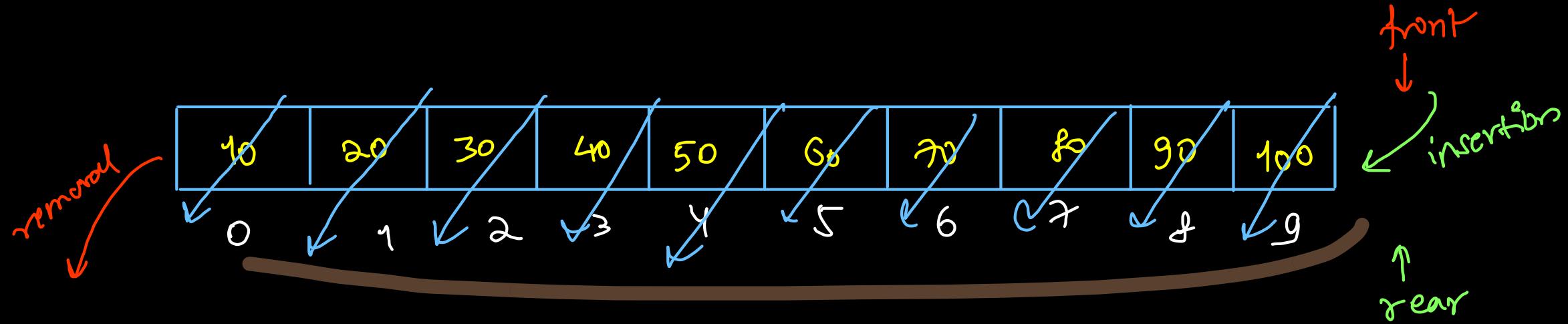
q.add(100)

q.add(110) {  
    **rear == n** }  
    ↳ queue overflow }

void add(int x){

    arr[rear] = x;

**rear++**;



$q.\text{remove}() \Rightarrow 10$

$q.\text{remove}() \Rightarrow 20$

$q.\text{remove}() \Rightarrow 30$

.

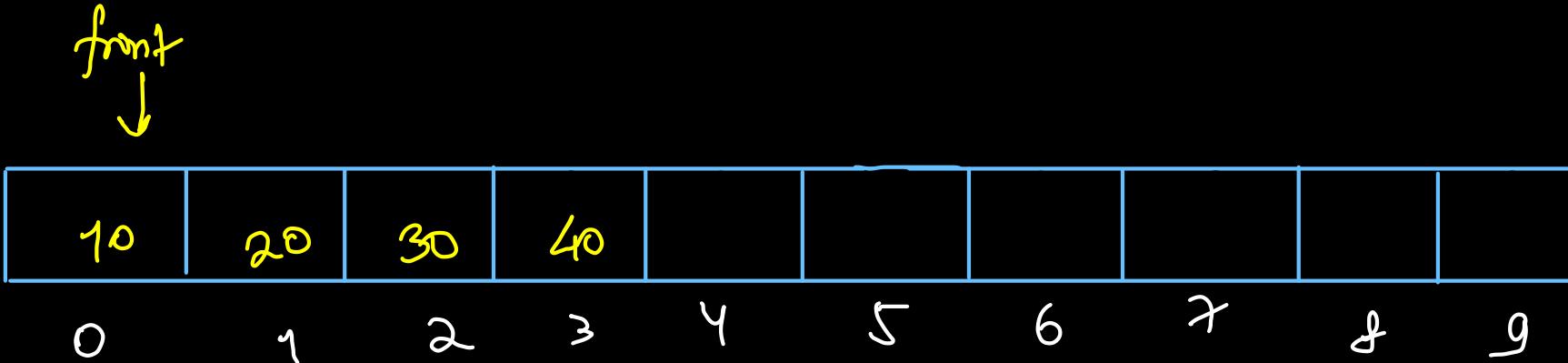
;

$q.\text{remove}() \Rightarrow 90$

$q.\text{remove}() \Rightarrow 100$

$q.\text{remove}()$   
↳ Queue Underflow

```
int remove() {
    int res = arr[front];
    front++;
    return res;
}
```



↑  
rear

`q.remove()` :  $\Rightarrow f = r$   
queue underflow

`q.add(10)`

`q.remove()`  $\Rightarrow 10$

`q.add(20)`

`q.remove()`  $\Rightarrow 20$

`q.add(30)`

`q.remove()`  $\Rightarrow 30$

`q.add(40)`

`q.remove()`  $\Rightarrow 40$

`q.remove()`  
 $\Rightarrow f = r$

queue underflow

```

int[] q = new int[100000];
int front = 0, rear = 0;

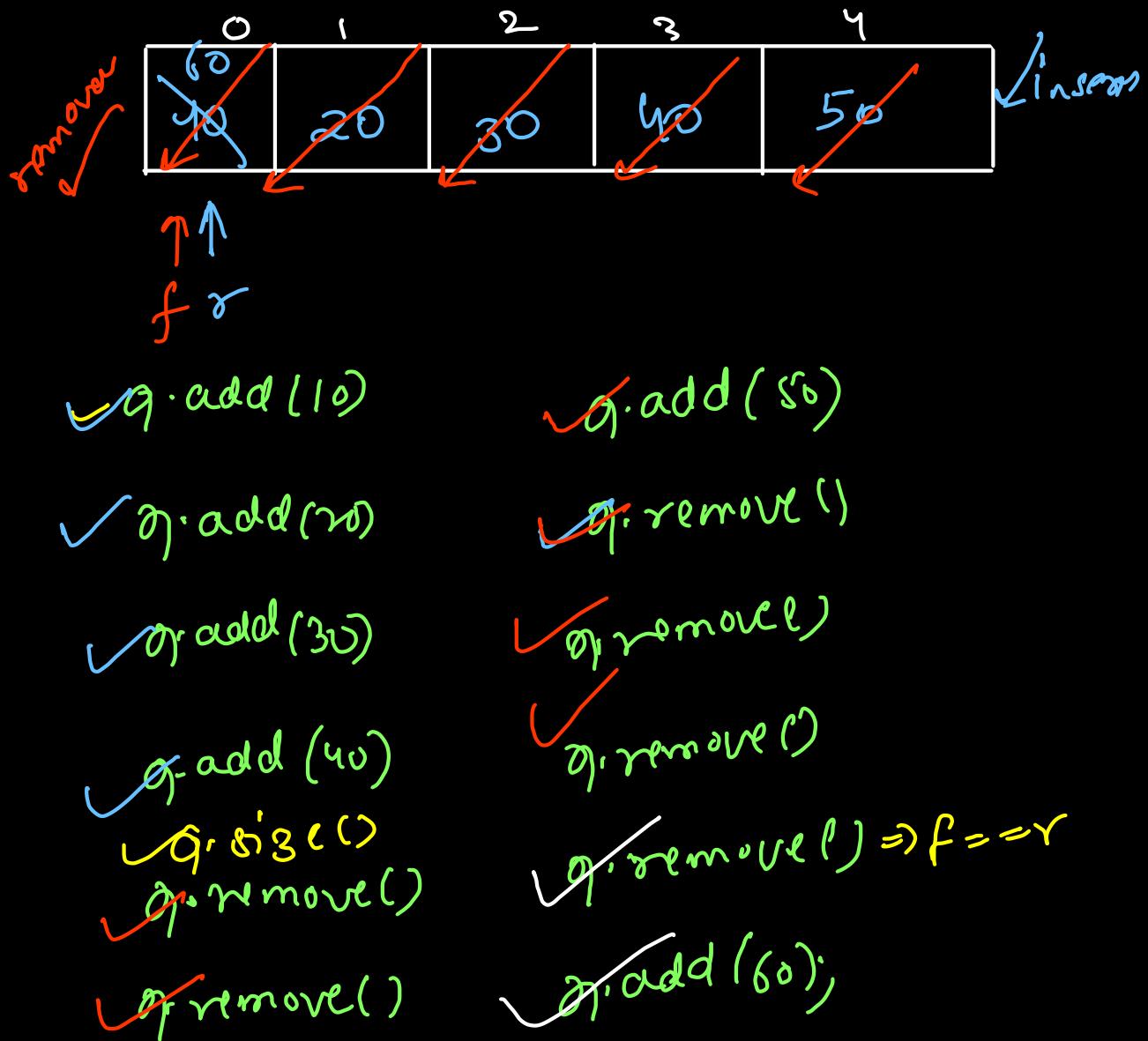
Scanner scn = new Scanner(System.in);
int tests = scn.nextInt();

while(tests-- > 0){
    String str = scn.next();

    if(str.equals("enqueue")){
        int x = scn.nextInt();
        if(rear < q.length){ //to ensure queue overflow
            q[rear] = x;
            rear++;
        }
    }

    } else if(str.equals("dequeue")){
        if(front == rear){
            // Queue Underflow: Empty Queue
            front = rear = 0;
        } else {
            front++;
        }
    } else if(str.equals("display")){
        for(int idx = front; idx < rear; idx++){
            System.out.print(q[idx] + " ");
        }
        System.out.println();
    } else if(str.equals("size")){
        System.out.println(rear - front);
    }
}

```



Implement queue using stack

FIFO

→ FIFO/LIFO

q.remove()

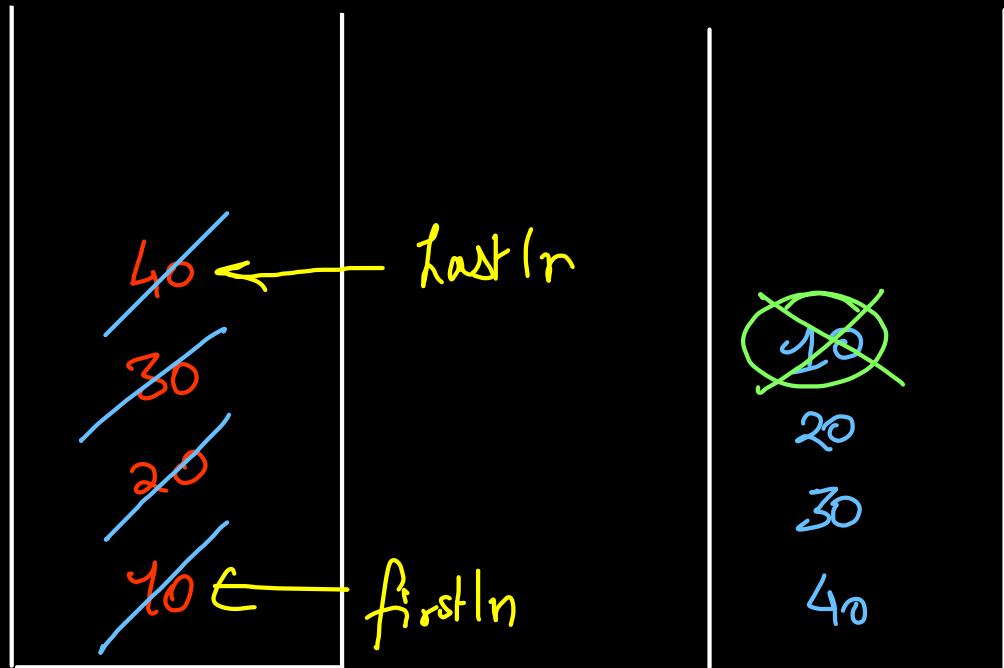
q.add(10)  $\Rightarrow$  stack.push(10)

q.add(20)  $\Rightarrow$  stack.push(20)

q.add(30)  $\Rightarrow$  stack.push(30)

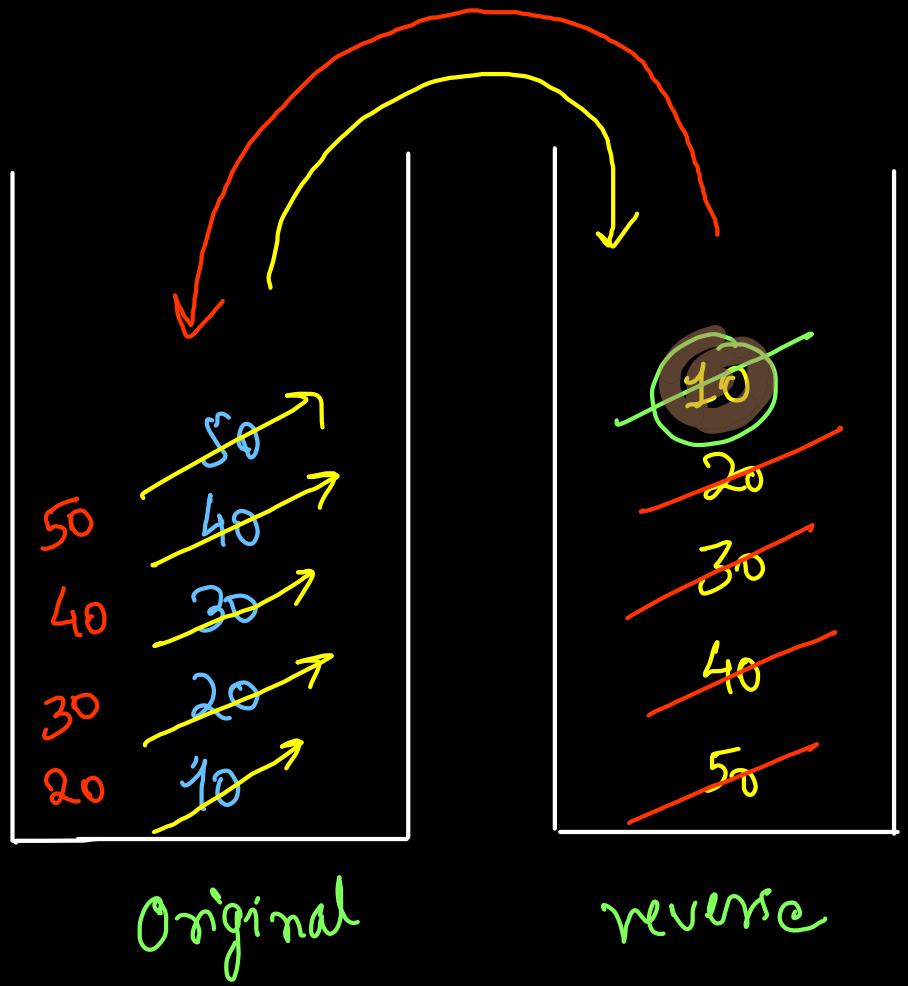
q.add(40)  $\Rightarrow$  stack.push(40)

} each  
O(1)  
operation



Stack

revStack



```

public int pop() {
    Stack<Integer> rev = new Stack<>();
    while(org.size() > 0){
        rev.push(org.pop());
    }

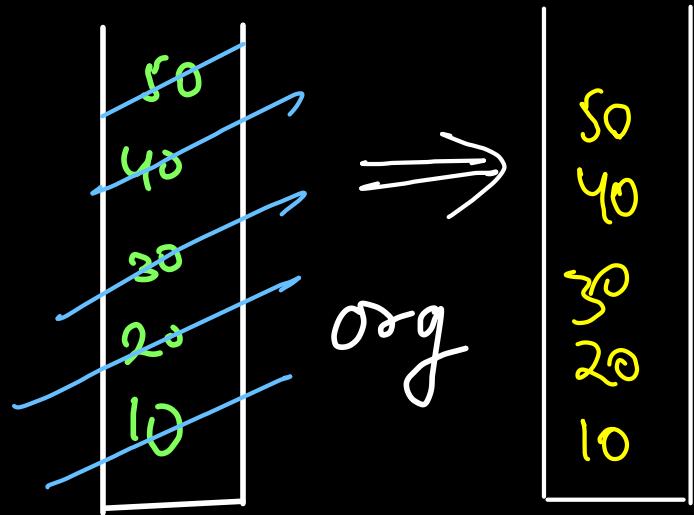
    int res = rev.pop();

    while(rev.size() > 0){
        org.push(rev.pop());
    }

    return res;
}

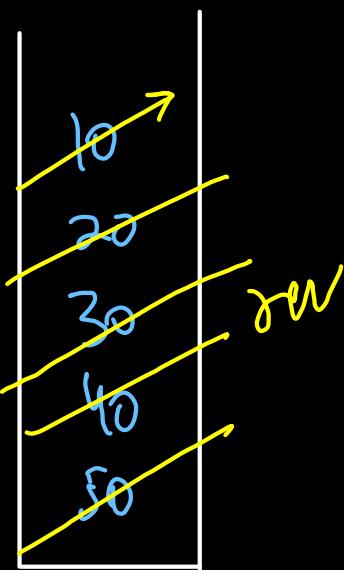
```

$O(n)$   
per  
remove



$\sigma.\text{peek}()$   
 $\hookrightarrow \text{front} : 10$

$\text{int res} = \text{rev}.\text{peek}();$

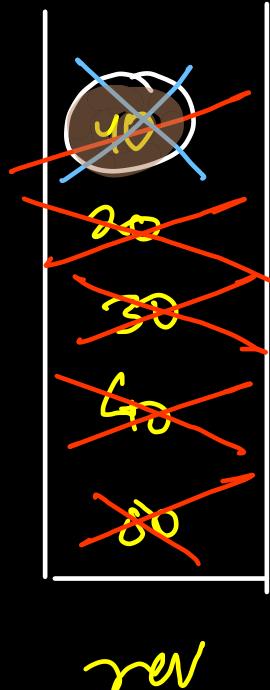
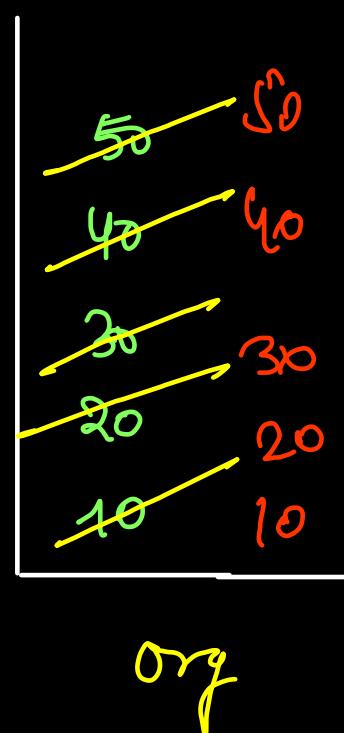


```
Stack<Integer> org = new Stack<>();  
  
public void push(int x) {  
    org.push(x); // Insert: O(1)  
}  
  
public int pop() {  
    Stack<Integer> rev = new Stack<>();  
    while(org.size() > 0){  
        rev.push(org.pop());  
    }  
  
    int res = rev.pop();  
  
    while(rev.size() > 0){  
        org.push(rev.pop());  
    }  
  
    return res;  
}
```

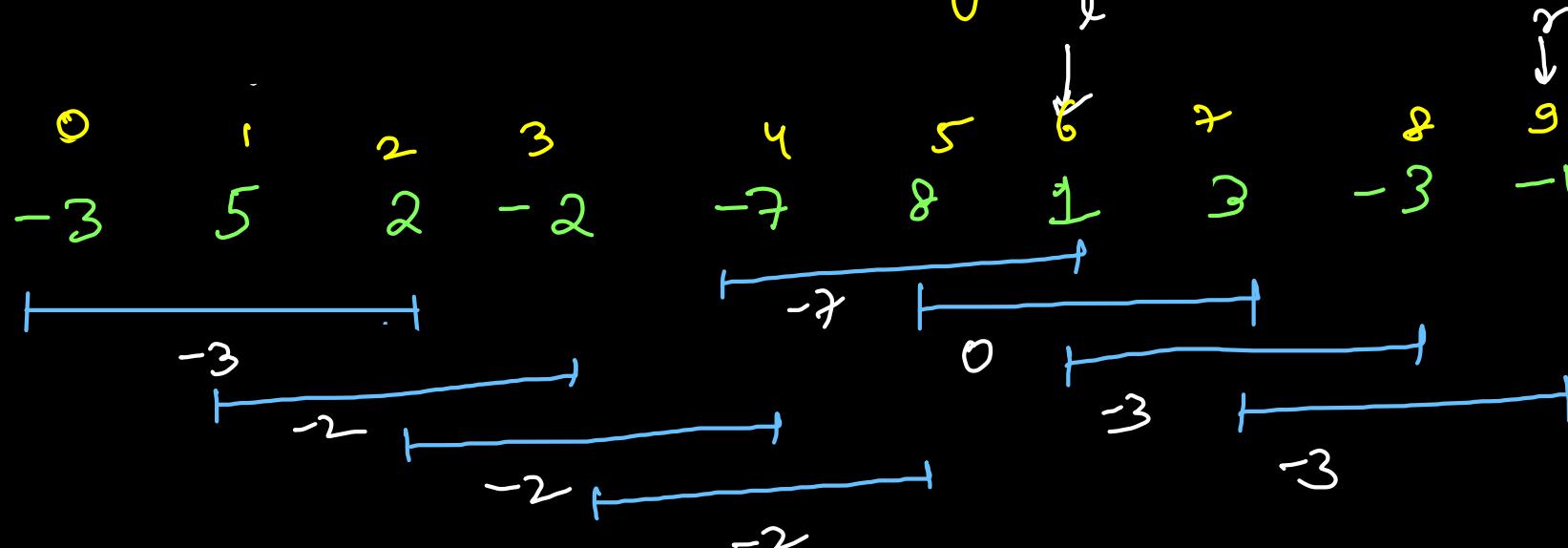
```
public int peek() {  
    Stack<Integer> rev = new Stack<>();  
    while(org.size() > 0){  
        rev.push(org.pop());  
    }  
  
    int res = rev.peek();  
  
    while(rev.size() > 0){  
        org.push(rev.pop());  
    }  
  
    return res;  
}  
  
public boolean empty() {  
    return org.empty();  
}
```

$q.add(10 // 20 / 30 / 40 / 50);$

```
Stack<Integer> org = new Stack<>();  
  
O(n) public void push(int x) {  
    org.push(x); // Insert: O(1)  
}  
  
public int pop() {  
    Stack<Integer> rev = new Stack<>();  
    while(org.size() > 0){  
        rev.push(org.pop());  
    }  
  
    O(n) int res = rev.pop(); reverse rev.pop()  
  
    while(rev.size() > 0){  
        org.push(rev.pop());  
    }  
  
    return res; // 10  
}
```



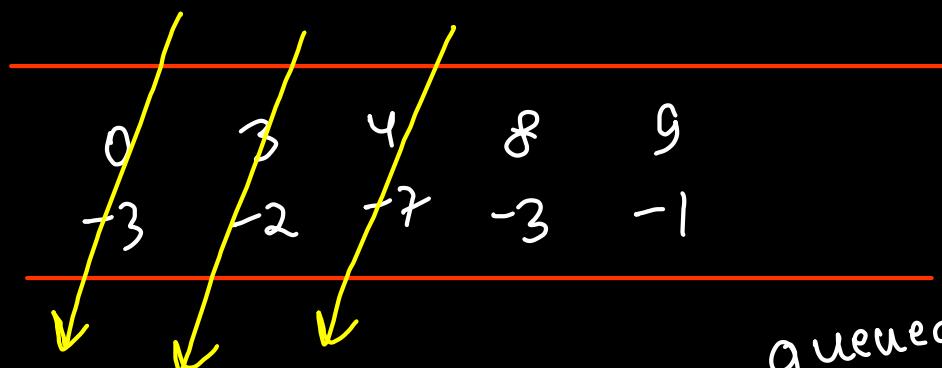
# First Negative Integer



window = ③  
static sliding window

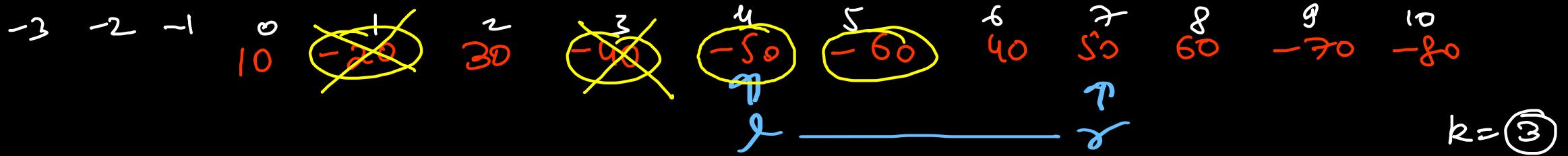
$$\begin{array}{ll} r=5 & l=2 \\ r=4 & l=1 \\ r=3 & l=0 \\ r=0 & l=-3 \end{array}$$

left → exclude  
right → include



queue of  
negative integer  
in sliding  
window

Queue  
First Neg In  
First Out



```

Queue<Integer> q = new ArrayDeque<>();

int right = 0, left = -k;
while(right < n){
    // include the right element } O(1)
    if(arr[right] < 0)
        q.add(arr[right]);

    // exclude the left element } O(1)
    if(left >= 0 && arr[left] < 0)
        q.remove();

    left++; right++;
    if(left < 0) continue;
    if(q.size() == 0) System.out.print(0 + " ");
    else System.out.print(q.peek() + " ");
}

```

$-20, -20, -40, -40, -50$   
 Time  $\approx n * (\text{include} + \text{exclude})$   
 $= n * 2 \Rightarrow \underline{\underline{O(n)}}$   
 (mean extra space queue  $\approx \underline{\underline{O(k)}}$ )

## Priority Queue

preference/importance

→ Ramleela

→ Ram Dusar (VUVP)

→ Laxman Dusar (VIP)

→ Common People

→ Stream BTech

CS

IT

ECE

Mech

BioTech

→ windows Task manager

→ Chrome ↑CPU

→ timer, disk cleanup, utility tasks

↳ running in background

JEE Rank List

Min-Heap by default

Rank 1  $\rightarrow$  min<sup>m</sup> no  $\Rightarrow$  highest priority

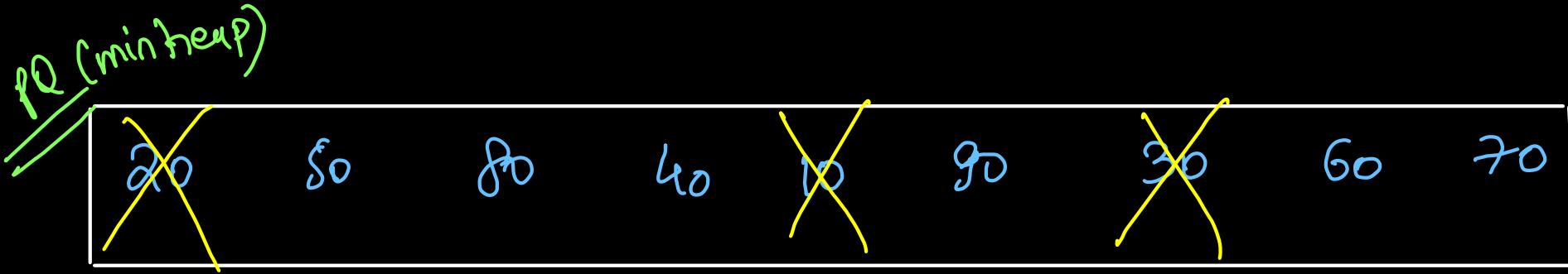
Rank  $\rightarrow$  ~~rank~~  
higher no  $\Rightarrow$  lower priority

Marks

Max-Heap

0/100  $\rightarrow$  lowest marks  $\rightarrow$  lowest priority

100/100  $\rightarrow$  higher marks  $\rightarrow$  highest priority



$q.\text{peek}() \Rightarrow 10$

$q.\text{remove}() \Rightarrow 10$

$q.\text{peek}() \Rightarrow 20$

$q.\text{remove}()$

$q.\text{peek}() \Rightarrow 30$

Single operations {

- insert  $\rightarrow O(\log n)$
- peek  $\rightarrow O(1)$
- remove  $\rightarrow O(\log n)$

Nodes {

- insert  $\rightarrow N \log N$
- peek  $\rightarrow N$
- remove  $\rightarrow N \log N$

```
PriorityQueue<Integer> q = new PriorityQueue<>();
// Min Heap: Highest Priority - Minimum Value

// Single Operation: O(log N)
q.add(20);
q.add(50);
q.add(80);
q.add(40);
q.add(10);
q.add(90);
q.add(30);
q.add(60);
q.add(70);

// Random Order : Unsorted Arrays
System.out.println(q);

// Heap Sort: O(N * log N) Time
while(q.size() > 0){
    System.out.print(q.remove() + " : ");
    System.out.println(q);
}
```

Finished in 90 ms

```
[10, 20, 30, 50, 40, 90, 80, 60, 70]
10 : [20, 40, 30, 50, 70, 90, 80, 60]
20 : [30, 40, 60, 50, 70, 90, 80]
30 : [40, 50, 60, 80, 70, 90]
40 : [50, 70, 60, 80, 90]
50 : [60, 70, 90, 80]
60 : [70, 80, 90]
70 : [80, 90]
80 : [90]
90 : []
```

M<sup>in</sup>H<sub>eap</sub>

```
PriorityQueue<Integer> q = new PriorityQueue<>(Collections.reverseOrder());
// Max Heap: Highest Priority - Maximum Value

// Single Operation: O(log N)
q.add(20);
q.add(50);
q.add(80);
q.add(40);
q.add(10);
q.add(90);
q.add(30);
q.add(60);
q.add(70);

// Random Order : Unsorted Arrays
System.out.println(q);

// Heap Sort in Decreasing Order: O(N * log N) Time
while(q.size() > 0){
    System.out.print(q.remove() + " : ");
    System.out.println(q);
}
```

Finished in 86 ms

```
[90, 70, 80, 60, 10, 50, 30, 20, 40]
90 : [80, 70, 50, 60, 10, 40, 30, 20]
80 : [70, 60, 50, 20, 10, 40, 30]
70 : [60, 30, 50, 20, 10, 40]
60 : [50, 30, 40, 20, 10]
50 : [40, 30, 10, 20]
40 : [30, 20, 10]
30 : [20, 10]
20 : [10]
10 : []
```

( $O(n \log n)$ )  
1. Declare an Empty *PriorityQueue*  $q$ .  $\rightarrow PQ < \text{Integer} \gg$   $pq = \text{new } PQ<>()$

2. Take Single Integer  $T$  as input.

3. For next  $T$  Lines format (*case, x(optional)*)

- case 1. *Print* the *size* of the *PriorityQueue* in a separate line.  $\rightarrow pq.size() \rightarrow O(1)$
- case 2. *Remove* the smallest element from the *PriorityQueue*. If the queue is empty then print  $-1$  in a separate line.  $\rightarrow pq.remove() \rightarrow O(\log n)$
- case 3. *Add* Integer  $x$  to the *PriorityQueue* s.  $\rightarrow pq.add(x) \rightarrow O(\log n)$
- case 4. *Print* the smallest element of the *PriorityQueue*. If queue is empty print  $-1$  in a seperate line.

$\rightarrow pq.peek() \rightarrow O(1)$

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int tests = scn.nextInt();

    PriorityQueue<Integer> pq = new PriorityQueue<>();

    while(tests-- > 0){
        int choice = scn.nextInt();
        if(choice == 1){
            // size: O(1)
            System.out.println(pq.size());
        } else if(choice == 2){
            // remove: Highest Priority: Minimum Value: O(log N)
            if(pq.size() == 0) System.out.println(-1);
            else pq.remove();
        } else if(choice == 3){
            // Insert: O(log N)
            int x = scn.nextInt();
            pq.add(x);
        } else {
            // Get the Smallest Element: Highest Priority: O(1)
            if(pq.size() == 0) System.out.println(-1);
            else System.out.println(pq.peek());
        }
    }
}
```

# Break Stone

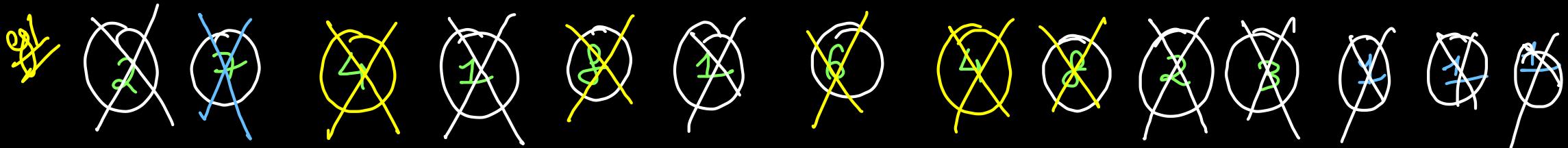
You are given an array of integers **stones** where **stones[i]** is the weight of the  $i_{th}$  stone.

We are playing a game with the stones. On each turn, we choose the **heaviest two stones** and smash them together. Suppose the heaviest two stones have weights  $x$  and  $y$  with  $x \leq y$ . The result of this smash is:

- If  $x == y$ , both stones are destroyed, and
- If  $x != y$ , the stone of weight  $x$  is destroyed, and the stone of weight  $y$  has new weight  $y - x$ .

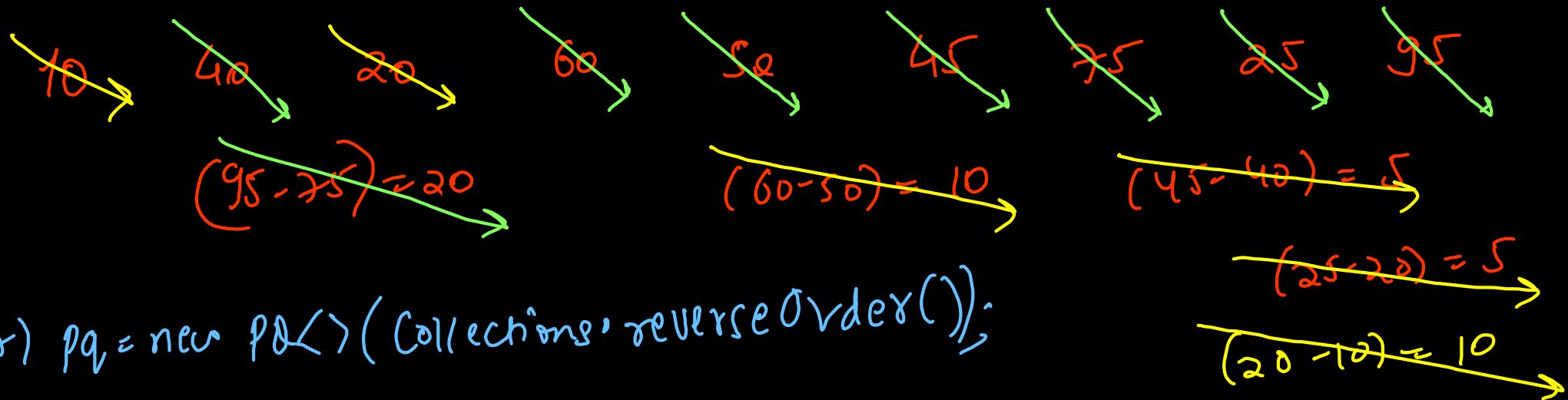
At the end of the game, there is **at most one** stone left.

*Return the weight of the last remaining stone. If there are no stones left, return 0.*



return 0;

A diagram illustrating a Max-Heap structure. At the top, the label "Max-Heap" is written in blue, with a blue arrow pointing downwards towards the heap. The heap itself consists of several yellow nodes arranged in a tree-like structure. The root node is labeled "egg". The left child of the root is labeled "apple", and its right child is labeled "orange". The left child of "apple" is labeled "banana", and its right child is labeled "cherry". The left child of "orange" is labeled "date", and its right child is labeled "grape". All labels are in yellow, matching the color of the heap nodes.



```

public static void main(String[] args) {
    PriorityQueue<Integer> pq =
        new PriorityQueue<>(Collections.reverseOrder());
    // max-heap

    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();

    for(int idx = 0; idx < n; idx++){
        int x = scn.nextInt();
        pq.add(x); → logn
    }

    while(pq.size() > 1){ → 2n logn
        int top1 = pq.remove(); → logn
        int top2 = pq.remove(); → logn
        if(top1 > top2) pq.add(top1 - top2); → logn
    }

    if(pq.size() == 0) System.out.println(0);
    else System.out.println(pq.peek());
}

```

$\downarrow n \log n$   
 $\Rightarrow \underline{\underline{O(n \log n)}}$

# Problem Statement

You are given an  $m \times n$  binary matrix  $\text{mat}$  of 1's (representing soldiers) and 0's (representing civilians). The soldiers are positioned in front of the civilians. That is, all the 1's will appear to the left of all the 0's in each row.

0	1	1	0	0	0
1	1	1	1	0	
2	1	0	0	0	0
3	1	1	0	0	0
4	1	1	1	1	1

A row  $i$  is **weaker** than a row  $j$  if one of the following is true:

- The number of soldiers in row  $i$  is less than the number of soldiers in row  $j$ .
- Both rows have the same number of soldiers and  $i < j$ .

*Return the indices of the  $k$  weakest rows in the matrix ordered from weakest to strongest.*

$$\gamma_0 < \gamma_1 \quad \gamma_2 < \gamma_3$$

$$\gamma_0 < \gamma_3 \quad \gamma_4 > \gamma_3$$

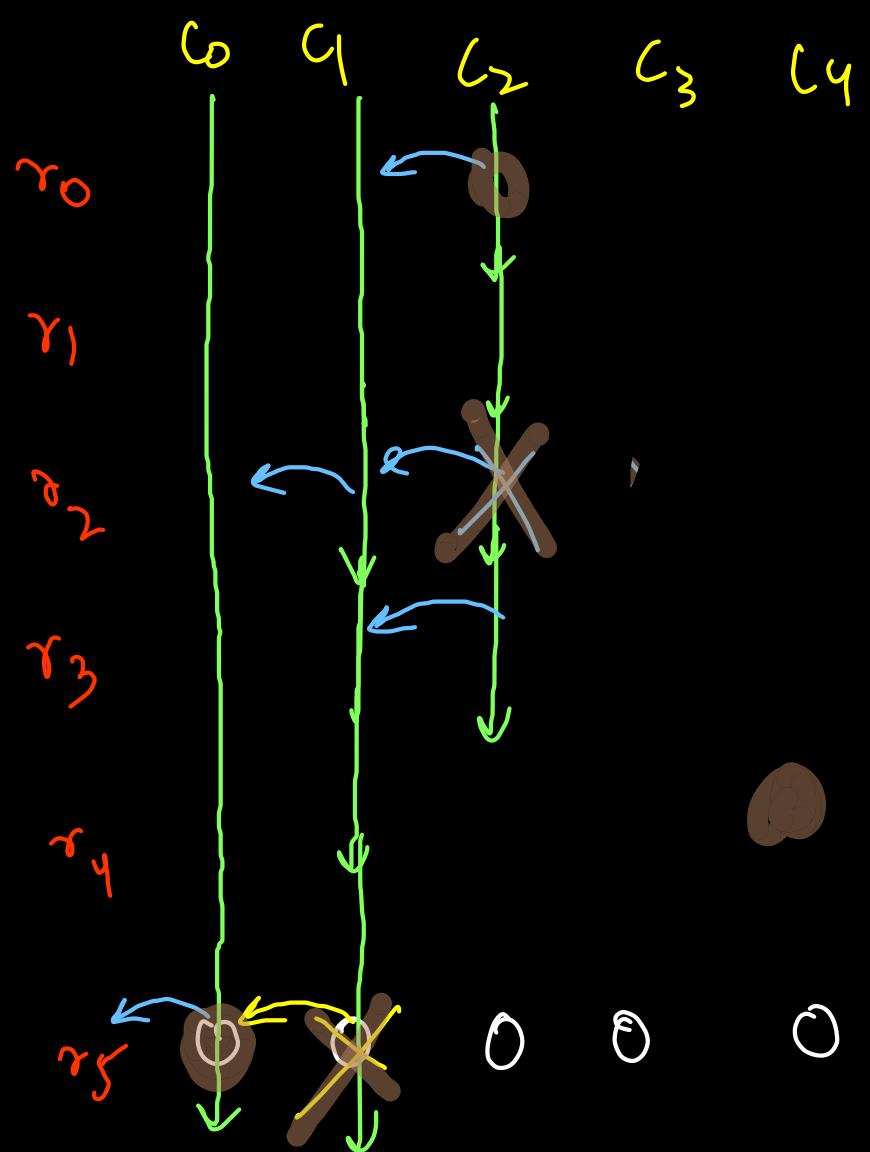
$$\gamma_1 < \gamma_4$$

$k=3$

$2^{\text{nd}} \rightarrow 1 \text{ soldier}$

$0^m \rightarrow 2 \text{ soldiers}$

$3^{\text{rd}} \rightarrow 2 \text{ soldiers}$



$R=4$

weakest  $\rightarrow$  Strongest  
min<sup>m</sup> soldiers      max<sup>m</sup> soldiers  
top to down

$r_5 \Rightarrow 0$  soldiers

$r_2 \Rightarrow 1$  soldier

$r_6 \Rightarrow 2$  soldiers

$r_3 \Rightarrow 2$  soldiers

Time till 8:34 PM

to code.

```
public static void weakestRows(int[][] mat, int rows, int cols, int k){  
    for(int c = 0; c < cols; c++){  
        for(int r = 0; r < rows; r++){  
            if(c - 1 >= 0 && mat[r][c - 1] == 0) continue;  
            if(mat[r][c] == 0){  
                System.out.print(r + " ");  
                k--;  
                if(k == 0) return;  
            }  
        }  
    }  
  
    for(int r = 0; r < rows; r++){  
        if(mat[r][cols - 1] == 1){  
            // entire row consists of 1s: No 0s  
            System.out.print(r + " ");  
            k--;  
            if(k == 0) return;  
        }  
    }  
}
```

*corner case*

If row is already printed  
ignore it to  
avoid printing  
it again.

Given an array of digits (values are from 0 to 9), find the **minimum possible sum** of two numbers formed from digits of the array. All digits of the given array must be used to form the two numbers.

Any combination of digits may be used to form the two numbers to be summed. Leading zeroes are permitted.

If forming two numbers is **impossible** (i.e. when n==0) then the "sum" is the value of the only number that can be formed.

0	1	2	3	4	5
1	2	3	4	5	6

$$(123 + 456) = 579$$

$$(145 + 236) = 381$$

$$(165 + 234) = 399$$

$$(135 + 246) = 381$$

$$12345 + 6 = 12351$$

$$1234 + 56 = 1290$$

$\text{eg}^2$

2 ✓ 0 ✓ 9 ✓ 4 ✓ 8 ✓ 5 ✓ 6 ✓ 2 ✓ 3 ✓ 7 ✓ 1 ✓

~~n1~~ ~~n2~~ ~~n1~~ ~~n2~~ ~~n1~~ ~~n2~~ ~~n1~~ ~~n2~~ ~~n1~~ ~~n2~~ ~~n1~~ ~~n2~~ ~~n1~~

" 0 2 3 5 7 9 "



2 3 5 7 9

+

+

" 1 2 4 6 8 "



1 2 4 6 8

```
Scanner scn = new Scanner(System.in);
PriorityQueue<Integer> pq = new PriorityQueue<>();

int n = scn.nextInt();
while(n-- > 0){
    int x = scn.nextInt();
    pq.add(x);
}

int v1 = 0, v2 = 0; if
int counter = 0;
while(pq.size() > 0){
    int digit = pq.remove();

    if(counter % 2 == 0)
        v1 = v1 * 10 + digit;
    else v2 = v2 * 10 + digit;

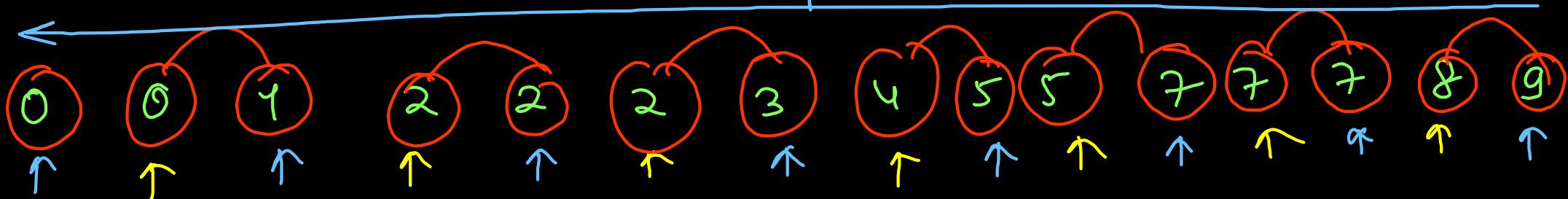
    counter++;
}

System.out.println(v1 + v2);
```

int vt=0; 4 bytes(stack)  
int v2 = 0; 4 bytes(stack)  
st

Answer  $\Rightarrow$   
overflowing  
integer  
range

# Mimic sum of 2 nos



String  $(d_1 + d_2 + c) \mod 10$   
 $c_{res} = "75306410"$  reverse  $\Rightarrow \underline{\underline{01460357}}$

$(d_1 + d_2 + c) \mod 10$   
 carry =  $\emptyset \neq \neq \neq \neq \neq \emptyset \emptyset \emptyset \emptyset$

"01235779" + "0224578"

```

PriorityQueue<Integer> pq = new PriorityQueue<>(Collections.reverseOrder());
int n = scn.nextInt();
while(n-- > 0){
    int x = scn.nextInt();
    pq.add(x);
}

StringBuilder res = new StringBuilder();

int carry = 0;

while(pq.size() > 0 || carry > 0){
    int d1 = (pq.size() == 0) ? 0 : pq.remove();
    int d2 = (pq.size() == 0) ? 0 : pq.remove();

    res.append((d1 + d2 + carry) % 10);
    carry = (d1 + d2 + carry) / 10;
}

// delete leading zeros
for(int idx = res.length() - 1; idx >= 0; idx--){
    if(res.charAt(idx) == '0') res.deleteCharAt(idx);
    else break;
}
res = res.reverse();
System.out.println(res);

```

max heap

arr:



$\log n$

carry = 0  
7 1 0  
0

" 86951 "

↓ reverse  
X 15965

```
String solve(int[] arr, int n) {
    PriorityQueue<Integer> pq = new PriorityQueue<>
        (Collections.reverseOrder());
    for(int val: arr)
        pq.add(val);

    StringBuilder res = new StringBuilder();
    int carry = 0;

    while(pq.size() > 0 || carry > 0) {
        int d1 = (pq.size() == 0) ? 0 : pq.remove();
        int d2 = (pq.size() == 0) ? 0 : pq.remove();

        if(d1 == 0 && carry == 0) break;
        // dont add leading zeros

        res.append((d1 + d2 + carry) % 10);
        carry = (d1 + d2 + carry) / 10;
    }

    return res.reverse().toString();
}
```

*Groceries  
Shop*

Fraction

Rice	wheat	Cereals	Salt	Sugar
10 kg	5 kg	2 kg	20 kg	20 kg
700 rs	1000 rs	2000 rs	500 rs	3000 rs

20 rs/kg

200rs/kg  
②

100rs/kg  
①

25 rs/kg

150rs/kg  
③

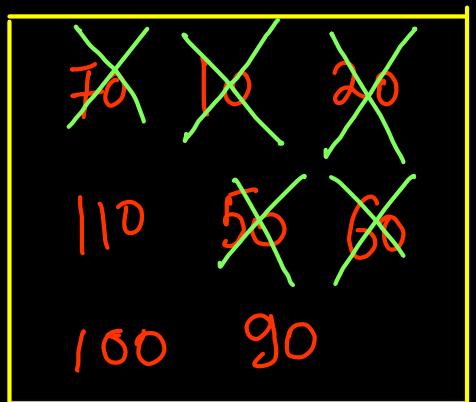
thief

~~15kg~~

Cereals + wheat + sugar  
2kg 5kg 8kg

# K Largest Elements

70    10    20    110    50    60    100    40    90    30    80  
↑    ↑    ↑    ↑    ↑    ↑    ↑    ↑    ↑    ↑    ↑



$$k=3 \\ \text{par} \cdot 3 \cdot 3 \\ \underbrace{\text{add}}_{\text{remove}} \left\{ \log k \right\}$$

min heap (3 largest elements)

$k = 3$   
Time  
 $\hookrightarrow \underline{N \log k}$

~~Approach 1:~~

Insert all elements + Pop k  $\Rightarrow$  max k elements  
in maxHeap  
 $\hookrightarrow O(N \log N)$

$O(k \log n)$

110, 100, 90

~~Approach 2:~~

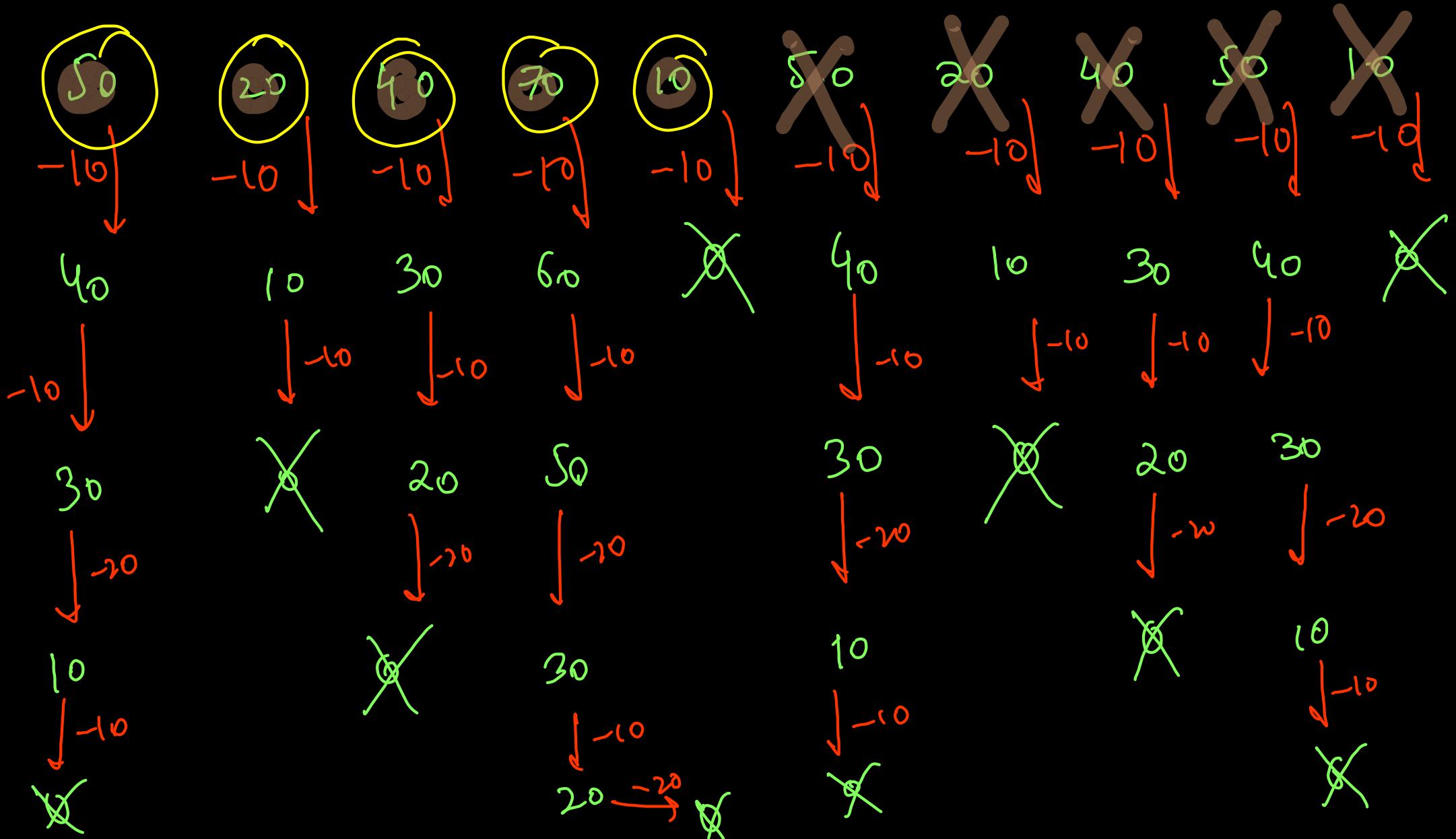
minHeap

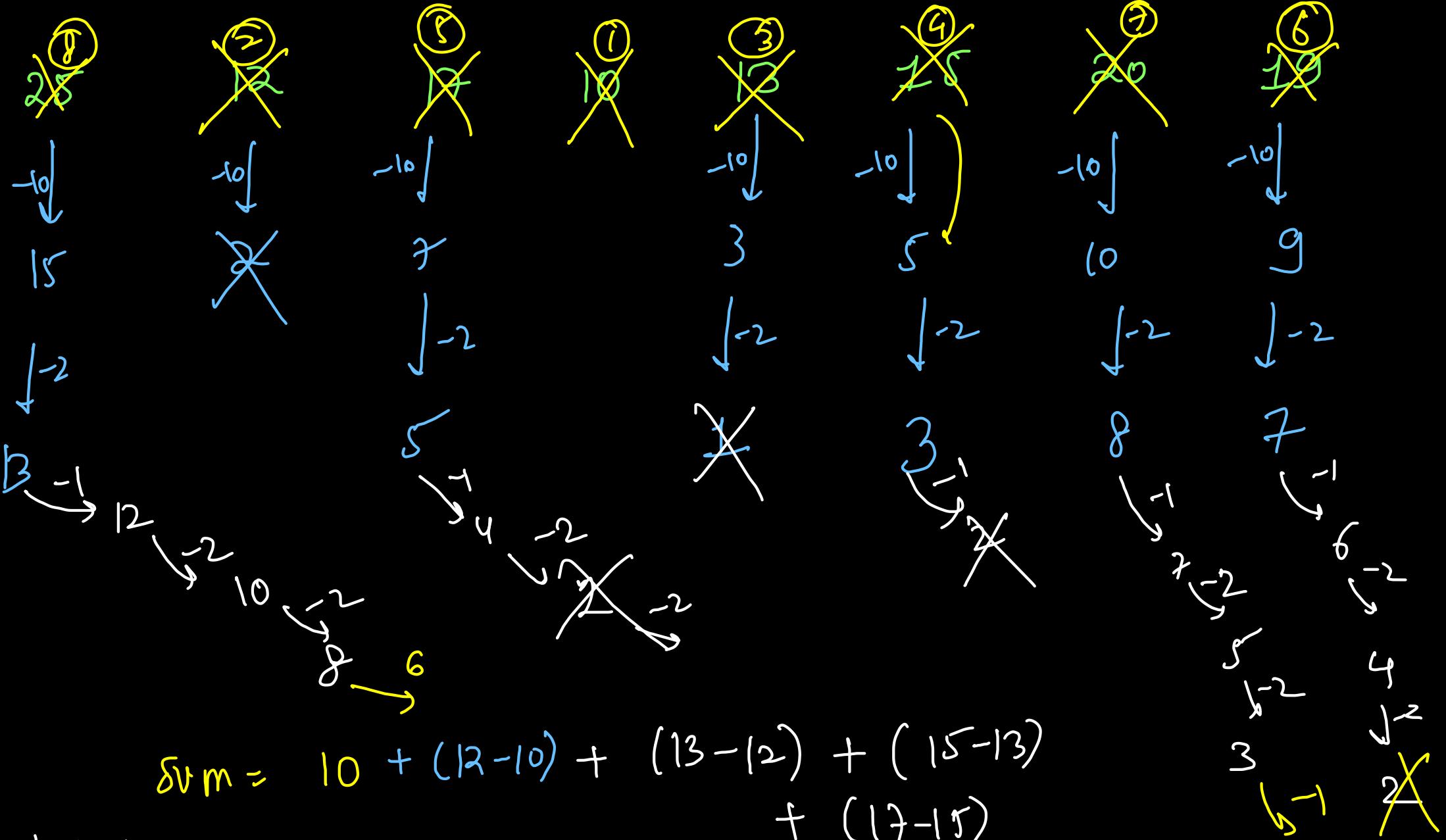
```
//Function to return k largest elements from an array.
public static ArrayList<Integer> kLargest(int arr[], int n, int k)
{
    PriorityQueue<Integer> pq = new PriorityQueue<>();

    for(int val: arr){
        if(pq.size() < k){
            pq.add(val); → there is vacancy
        } else if(pq.peek() < val){
            pq.remove(); → replace the worst performer
            pq.add(val);
        }
    }

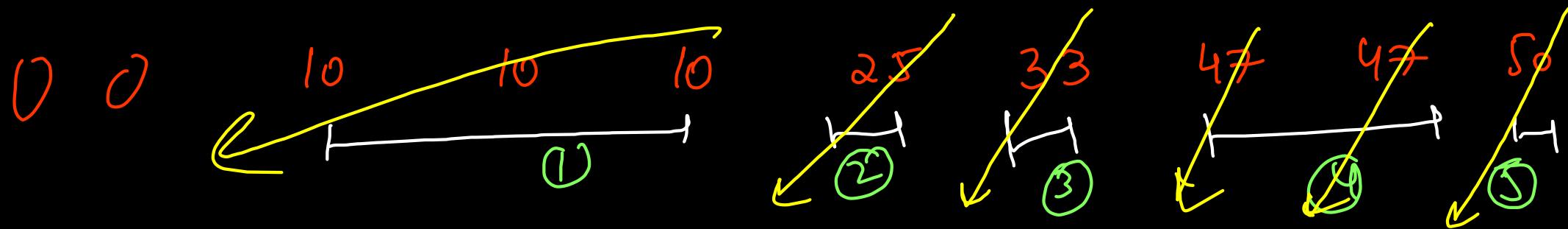
    ArrayList<Integer> res = new ArrayList<>();
    while(pq.size() > 0){
        res.add(pq.remove());
    }
    Collections.reverse(res);
    return res;
}
```

# Subtract Numbers





Step 1 + 1 + 1 + 1 + 1



```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();

    HashSet<Integer> set = new HashSet<>();

    for(int idx = 0; idx < n; idx++){
        int x = scn.nextInt();
        if(x > 0) set.add(x);
    }

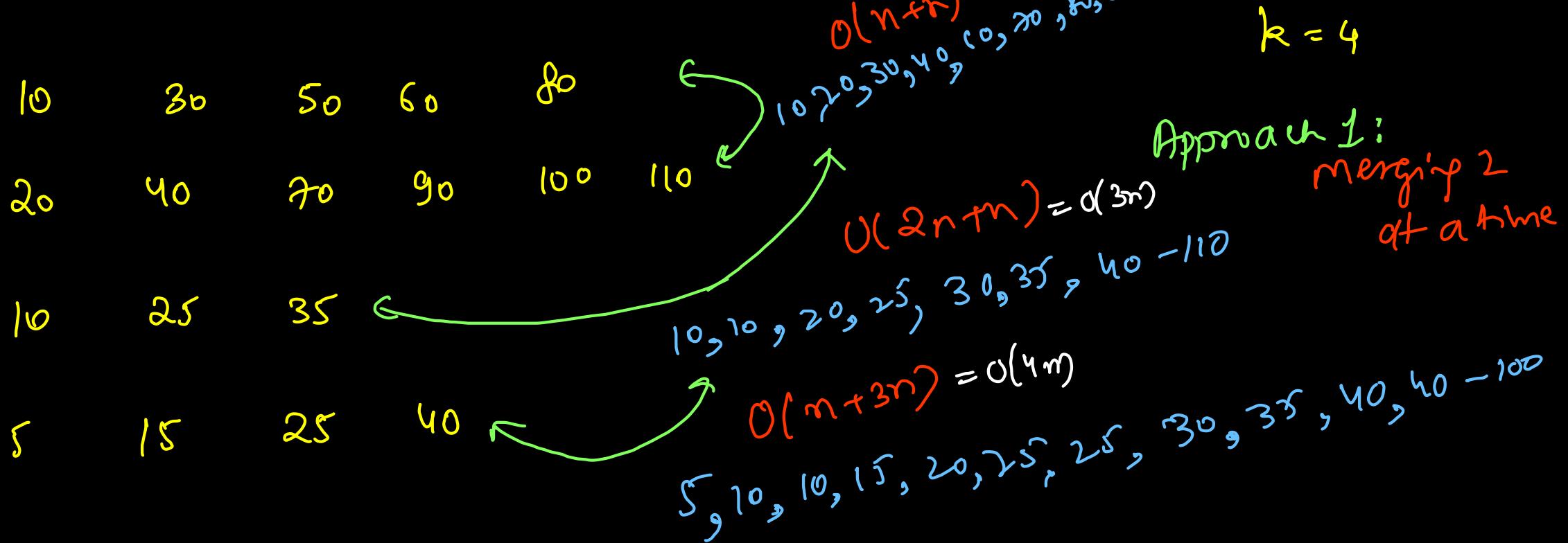
    System.out.println(set.size());
}

```

- ① 10      HashSet
- ②  $(25 - 10) = 15$  ↓  
HashMap
- ③  $(33 - 25) = 8$  w/o value
- ④  $(47 - 33) = 14$
- ⑤  $(50 - 47) = 3$

↳ number of unique non-zero elements

## Merge K Sorted Arrays

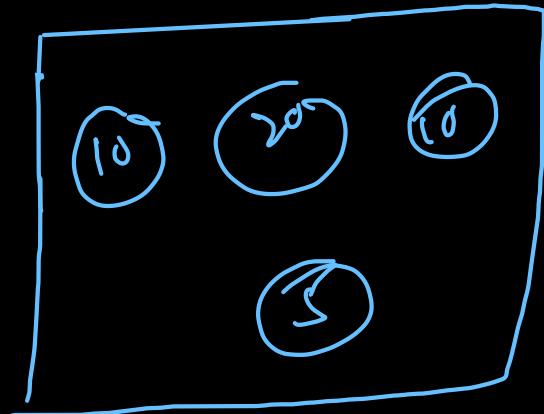


$$\begin{aligned}
 \text{Total Time Comp.} &\rightarrow O(2n + 3n + 4n + \dots + kn) \\
 &= O\left(\frac{n \times k \times (k+1)}{2}\right) = \underline{\underline{O(nk^2)}}
 \end{aligned}$$

Approach 2: Priority Queue

insert all  $(n \times k)$  elements in PQ  
and sort them  $\Rightarrow O(n \times k \log n \times k)$   
 $\Rightarrow O(n \times k \log n)$

$pQ[0]$	10	30	50	60	80
$pQ[1]$	20	40	70	90	100
$pQ[2]$	10	25	35		
$pQ[3]$	5	15	25	40	



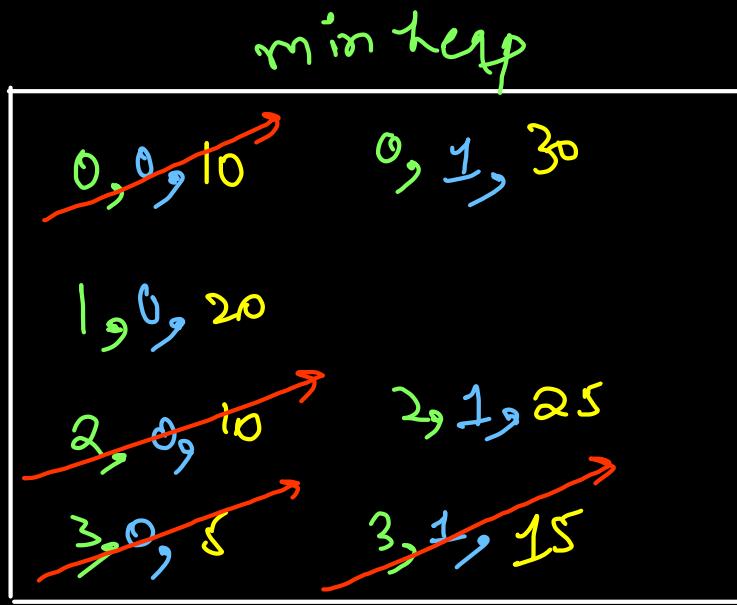
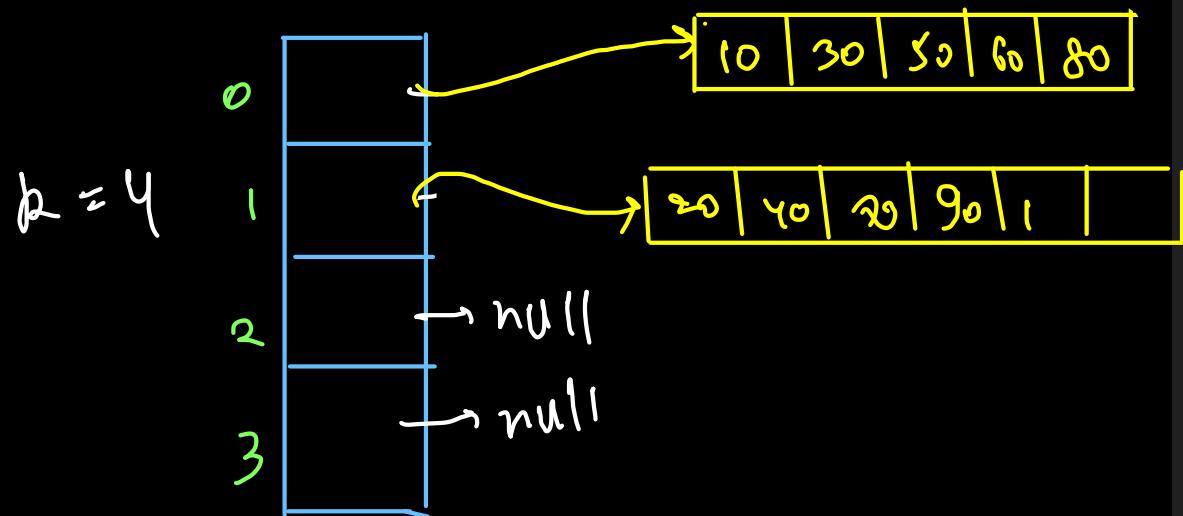
min-Heap

way1: push indices only  
↳ wrong

way2: push values only  
↳ wrong

$pQ < \text{int}[]?$   
↓  
{ row, col, value }

0	10		1	30	2	50	3	60	4	80	5	
1	20			40		70		90		100		110
2	10			25				35				
3	5			15								



```

public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);

    int rows = scn.nextInt();
    int[][] mat = new int[rows][];

    for(int row = 0; row < rows; row++) {

        int cols = scn.nextInt();
        mat[row] = new int[cols];

        for(int col = 0; col < cols; col++)
            mat[row][col] = scn.nextInt();
    }

    mergeKSorted(mat);
}

```

```

int rows = mat.length;

PriorityQueue<int[]> pq = new PriorityQueue<>();

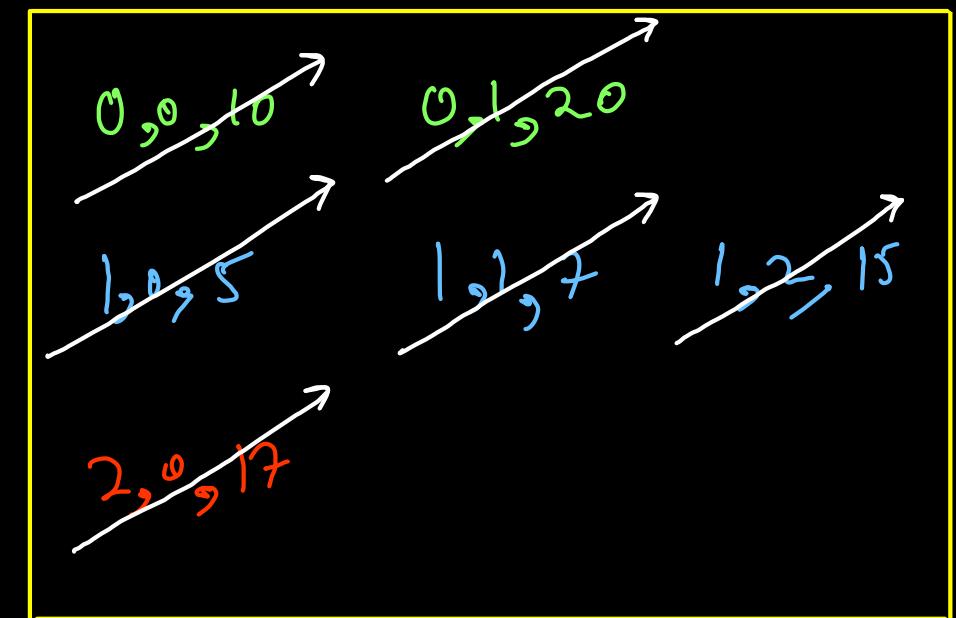
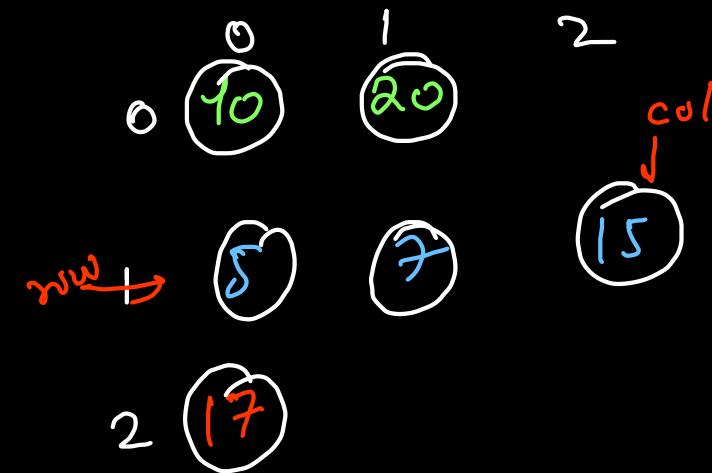
for(int row = 0; row < rows; row++){
    int[] ele = new int[3];
    ele[0] = row; ele[1] = 0; ele[2] = mat[row][0];
    pq.add(ele); // {row, col, value}
}

while(pq.size() > 0){
    int[] ele = pq.remove();
    int row = ele[0], col = ele[1], value = mat[row][col];

    System.out.println(value);

    if(col + 1 == mat[row].length) continue;
    col++; value = mat[row][col];
    pq.add(ele);
}

```



min heap

```
public static void mergeKSorted(int[][] mat){
    int rows = mat.length;
    PriorityQueue<int[]> pq = new PriorityQueue<>((e1, e2) -> e1[2] - e2[2]);
```

Comparator(after OOPS)

```
for(int row = 0; row < rows; row++){
    int[] ele = new int[3];
    ele[0] = row; ele[1] = 0; ele[2] = mat[row][0];
    pq.add(ele); // {row, col, value}
}

while(pq.size() > 0){
    int[] ele = pq.remove();
    int row = ele[0], col = ele[1], value = mat[row][col];
    System.out.print(value + " ");

    if(col + 1 == mat[row].length) continue;
    ele[1]++;
    ele[2] = mat[row][col + 1];
    pq.add(ele);
}
```

# Running Median

10    20     $\textcircled{30}$     40    50    odd length  
median

10    20     $\boxed{30 \quad 40}$     50    60    even length

$\underbrace{30 + 40}_{2} = \textcircled{35}$

## Running Median

$$[10] \rightarrow 10$$

$$[10, 50] \rightarrow (10 + 50)/2 = 30$$

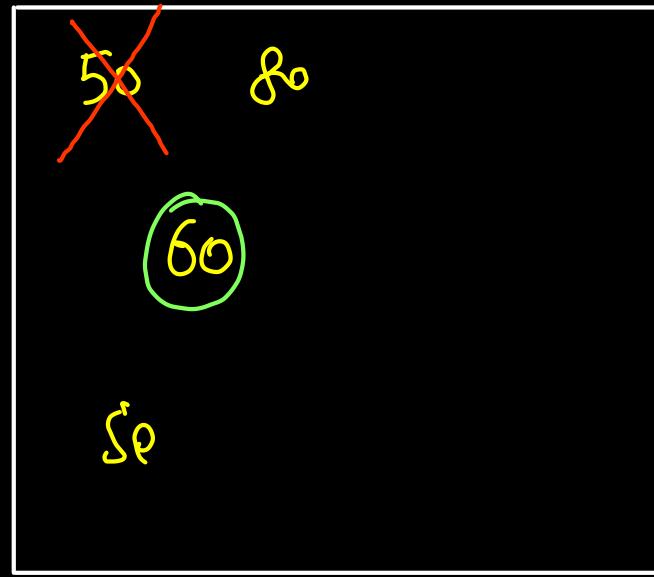
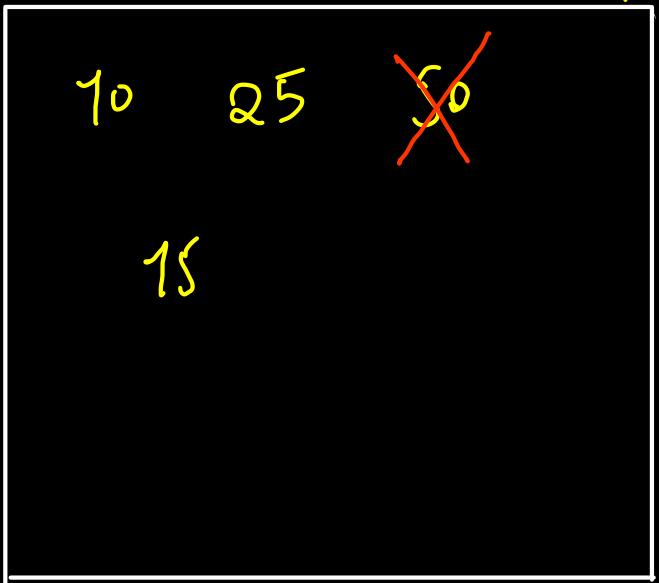
$$[10, 50, 80] \rightarrow 50$$

$$[10, [50, 60, 80]] \rightarrow (50 + 60)/2$$

$$[10, 30, 50, 60, 80] \rightarrow \{$$

$$[10, 30, 50, 60, 80, 100] \rightarrow \frac{50 + 60}{2} = 55$$

$$[10, 30, 50, 60, 80, 100, 110] \rightarrow 60$$



left  
 { Max Heap }

insert 10  $\Rightarrow$  (10)

insert 50  $\Rightarrow$   $(10+50)/2 = 30$

insert 80  $\Rightarrow$  (50)

right { min Heap }

insert 60  $\Rightarrow$   $(50+60)/2 = 55$

insert 25  $\Rightarrow$  50

insert 15  $\Rightarrow$

$n=5$   
odd

left  
3

right  
2

$$l = r + 1$$

or

$n=6$

left right  
3 3

$$l = r$$

balanced  
 $\equiv$

```

public void addNum(int num) {
    // Normal Insertion
    if(left.size() == 0 || num < left.peek()){
        left.add(num);
    } else right.add(num);

    // Balanced Condition
    if(left.size() == right.size() || left.size() == right.size() + 1)
        return;

    // Rebalancing
    if(right.size() > left.size()) left.add(right.remove());
    else right.add(left.remove());
}

```

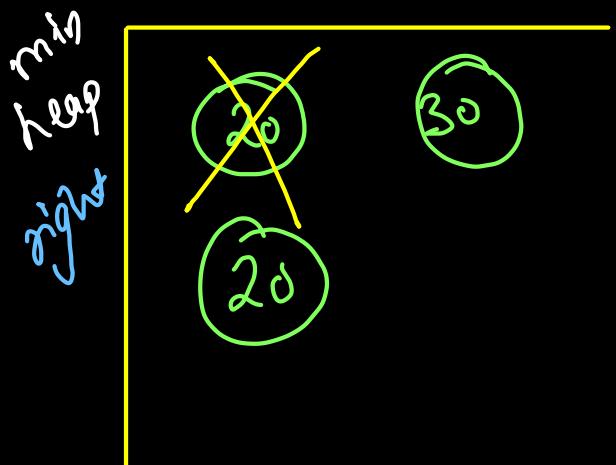
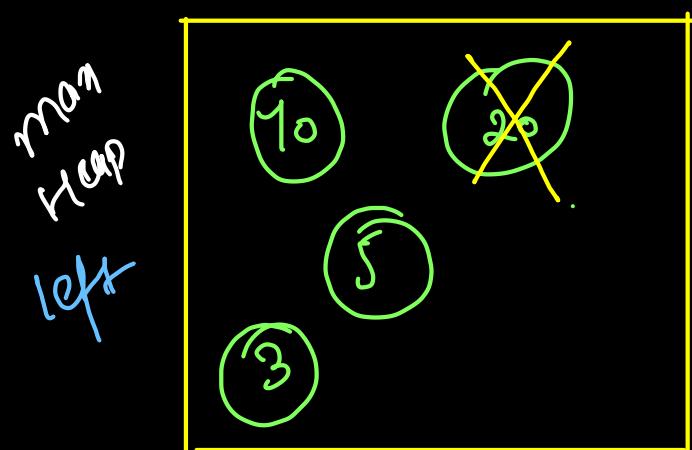
insert 10

insert 20

insert 30

insert 5

insert 3



```
class MedianFinder {
    PriorityQueue<Integer> left = new PriorityQueue<>(Collections.reverseOrder());
    PriorityQueue<Integer> right = new PriorityQueue<>();

    public void addNum(int num) {
        // Normal Insertion
        if(left.size() == 0 || num < left.peek()){
            left.add(num);
        } else right.add(num);

        // Balanced Condition
        if(left.size() == right.size() || left.size() == right.size() + 1)
            return;

        // Rebalancing
        if(right.size() > left.size()) left.add(right.remove());
        else right.add(left.remove());
    }

    public double findMedian() {
        if(left.size() == right.size() + 1) return left.peek();
        return (left.peek() + right.peek()) / 2.0;
    }
}
```