# **Model Recognition as Planning**

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#### Abstract

Given the partial observation of a plan execution and a set of possible planning models (that share the same state variables but define different action models), model recognition is the task of identifying which model in the set produced the given observation. The paper formalizes the model recognition task and proposes a novel method to assess the probability of a STRIPS model to produce a given partially observed plan execution. This method, that we called model recognition as planning, is robust to missing intermediate states and actions in the observed plan execution besides, it is built on top of offthe-shelf classical planning algorithms, that are used to elicit the likelihood of the observations given a candidate model. The effectiveness of model recognition as planning is shown in a set of STRIPS models encoding different kinds of finite state machines. We show that model recognition as planning succeeds to identify the executed automata despite internal machine states or actual applied transitions, are unobserved.

#### Introduction

*Plan recognition* is the task of predicting the future actions of an agent provided observations of its current behavior. Plan recognition is considered *automated planning* in reverse; while automated planning aims to compute a sequence of actions that accounts for a given goals, plan recognition aims to compute the goals that account for an observed sequence of actions (Geffner and Bonet 2013).

Diverse approaches has been proposed for plan recognition such as *rule-based systems*, *parsing*, *graph-covering*, *Bayesian nets*, etc (Carberry 2001; Sukthankar et al. 2014). *Plan recognition as planning* is the model-based approach for plan recognition (Ramírez 2012; Ramírez and Geffner 2009). This approach assumes that the action model of the observed agent is known and leverages it to compute the most likely goal of the agent, according to the observed plan execution.

This paper introduces the *model recognition* task where the target of the recognition process is not a goal but a *planning model*. Given a partial observation of a plan execution and a set of possible planning models (that share the same state variables but define different action models), *model recognition* is the task of identifying which model in the set

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Figure 1: Observation of the execution of a robot navigation plan in a  $5\times 5$  grid.

has the highest probability of producing the given observa-

To better illustrate *model recognition*, imagine a robot in a  $n \times n$  grid whose navigation is determined by the STRIPS model of Figure 2. According to this model the robot can increment its *x-coordinate* when it is at an *even* row while, at *odd rows*, can decrement the *x-coordinate*. Apart from this particular navigation model, different models can be defined within the same state variables (e.g. altering the way (state-q0) and (state-q1) are required and updated) and these models can determine different kinds of robot navigation. Given an observation of a plan execution, like the one illustrated at Figure , *model recognition* aims here to identify which navigation model produced that observation, despite key information is unobserved (e.g. the value of (state-q0) and (state-q1) or the particular applied actions).

Model recognition is of interest because once the planning model is recognized, then the model-based machinery for automated planning becomes applicable (Ghallab, Nau, and Traverso 2004). In addition, it enables identifying different kinds of automatae by observing their execution. It is well-known that diverse automatae representations, like finite state controllers, push-down automata, GOLOG programs or reactive policies, can be encoded as classical planning models (Baier, Fritz, and McIlraith 2007; Bonet, Palacios, and Geffner 2010; Segovia-Aguas, Jiménez, and Jonsson 2017; Ivankovic and Haslum 2015).

The paper introduces also *model recognition as planning*; a novel method to assess the probability of a given STRIPS model to produce an observed plan execution. The method is robust to missing intermediate states and actions in the observed plan execution besides, it is built on top of off-the-shelf classical planning algorithms, that are used to elicit

```
(:action inc-x
  :parameters (?v1 ?v2)
  :precondition (and (x-coord ?v1) (next ?v1 ?v2) (state-q0))
  :effect (and (not (x-coord ?v1)) (x-coord ?v2)))
(:action dec-x
 :parameters (?v1 ?v2)
  :precondition (and (x-coord ?v1) (next ?v2 ?v1) (state-q1))
 :effect (and (not (x-coord ?v1)) (x-coord ?v2)))
(:action inc-v-even
  :parameters (?v1 ?v2)
  :precondition (and (y-coord ?y1) (next ?y1 ?y2) (state-q0))
  :effect (and (not (y-coord ?y1)) (y-coord ?y2)
               (not (state-q0)) (state-q1))
(:action inc-y-odd
  :parameters (?v1 ?v2)
  :precondition (and (y-coord ?y1) (next ?y1 ?y2) (state-q1)))
 (:action dec-y-even
  :parameters (?y1 ?y2)
 :precondition (and (y-coord ?y1) (next ?y2 ?y1) (state-q0))
:effect (and (not (y-coord ?y1)) (y-coord ?y2)
               (not (state-q0)) (state-q1)))
(:action dec-v-odd
  :precondition (and (y-coord ?y1) (next ?y2 ?y1) (state-q1))
 :effect (and (not (y-coord ?y1)) (y-coord ?y2)
               (not (state-q1)) (state-q0)))
```

Figure 2: Action model for a robot navigation in a  $n \times n$  grid.

the likelihood of the observations given a candidate model. The paper evaluates the effectiveness of *model recognition* as planning with a set of STRIPS models that represent different finite state machines. All of these automatae are defined within the same alphabet and same machine states but different transition functions. We show that model recognition as planning succeeds to identify the executed automata despite internal machine states or actual applied transitions are unobserved.

# **Background**

This section formalizes the models for classical planning and for the observation of the execution of a classical plan.

#### Classical planning

We use F to denote the set of *fluents* (propositional variables) describing a state. A *literal* l is a valuation of a fluent  $f \in F$ , i.e. either l = f or  $l = \neg f$ . A set of literals L represents a partial assignment of values to fluents (without loss of generality, we will assume that L does not assign conflicting values to any fluent). We use  $\mathcal{L}(F)$  to denote the set of all literal sets on F, i.e. all partial assignments of values to fluents.

A state s is a full assignment of values to fluents and we explicitly include negative literals  $\neg f$  in states; i.e. |s| = |F|, so the size of the state space is  $2^{|F|}$ . Like in PDDL (Fox and Long 2003), we assume that fluents F are instantiated from a set of predicates  $\Psi$ . Each predicate  $p \in \Psi$  has an argument list of arity ar(p). Given a set of objects  $\Omega$ , the set of fluents F is induced by assigning objects in  $\Omega$  to the arguments of predicates in  $\Psi$ ; i.e.  $F = \{p(\omega) : p \in \Psi, \omega \in \Omega^{ar(p)}\}$  such that  $\Omega^k$  is the k-th Cartesian power of  $\Omega$ .

A classical planning frame is a pair  $\langle F,A\rangle$ , where F is a set of fluents and A is a set of actions whose semantics are specified with two functions:  $\rho(s,a)$  that denotes whether action a is applicable in a state s and  $\theta(s,a)$  that denotes the successor state that results of applying a in s. In this work we specify action semantics with the STRIPS model. With this regard, an action  $a \in A$  is defined with:

- $pre(a) \in \mathcal{L}(F)$ , the *preconditions* of a, is the set of literals that must hold for the action  $a \in A$  to be applicable.
- eff<sup>+</sup> $(a) \in \mathcal{L}(F)$ , the *positive effects* of a, is the set of literals that are true after the application of action  $a \in A$ .
- eff<sup>-</sup> $(a) \in \mathcal{L}(F)$ , the *negative effects* of a, is the set of literals that are false after the application of the action.

We say that an action  $a \in A$  is applicable in a state s iff  $pre(a) \subseteq s$ . The result of applying a in s is the successor state denoted by  $\theta(s,a) = \{s \setminus eff^{-}(a)\} \cup eff^{+}(a)\}$ .

A classical planning problem is a tuple  $P = \langle F, A, I, G \rangle$ , where I is an initial state and  $G \in \mathcal{L}(F)$  is a goal condition. A plan  $\pi$  for P is an action sequence  $\pi = \langle a_1, \ldots, a_n \rangle$  that induces the state trajectory  $s = \langle s_0, s_1, \ldots, s_n \rangle$  such that  $s_0 = I$  and, for each  $1 \leq i \leq n$ ,  $a_i$  is applicable in  $s_{i-1}$  and generates the successor state  $s_i = \theta(s_{i-1}, a_i)$ . The plan length is denoted with  $|\pi| = n$ . A plan  $\pi$  solves P iff  $G \subseteq s_n$ , i.e., if the goal condition is satisfied at the last state reached after following the application of the plan  $\pi$  in the initial state I. A solution plan for P is optimal if it has minimum length.

#### **Conditional effects**

Conditional effects are useful for compactly defining our *model recognition as planning* method because they allow planning actions to have different semantics according to the value of the current state.

An action  $a \in A$  with conditional effects is defined as a set of preconditions  $pre(a) \in \mathcal{L}(F)$  and a set of conditional effects cond(a). Each conditional effect  $C \rhd E \in cond(a)$  is composed of two sets of literals  $C \in \mathcal{L}(F)$ , the condition, and  $E \in \mathcal{L}(F)$ , the effect.

An action  $a \in A$  is *applicable* in a state s iff  $pre(a) \subseteq s$ , and the *triggered effects* resulting from the action application are the effects whose conditions hold in s:

$$triggered(s,a) = \bigcup_{C \triangleright E \in \mathsf{cond}(a), C \subseteq s} E,$$

The result of applying action a in state s is the successor state  $\theta(s,a) = \{s \setminus \mathsf{eff}^-_c(s,a)) \cup \mathsf{eff}^+_c(s,a)\}$  where  $\mathsf{eff}^-_c(s,a) \subseteq triggered(s,a)$  and  $\mathsf{eff}^+_c(s,a) \subseteq triggered(s,a)$  are, respectively, the triggered negative and positive effects.

# The observation model

Given a classical planning problem  $P=\langle F,A,I,G\rangle$  and a plan  $\pi$  that solves P, the observation of the execution of  $\pi$  on P is  $\tau=\langle a_1,\ldots,a_n,s_m\rangle$ , an interleaved combination of  $1\leq n\leq |\pi|$  observed actions and  $1\leq m\leq |\pi|+1$  observed states such that:

- Observed actions are *consistent* with  $\pi$  (Ramírez and Geffner 2009). This means that the sequence of observed actions  $\langle a_1, \ldots, a_n \rangle$  in  $\tau$  is a sub-sequence of the solution plan  $\pi$ .
- Observed states are *consistent* with  $\pi$ . On the one hand, the initial state  $s_0 = I$  is fully observed while the remaining observed states may be partial, the value of certain fluents in these states may be omitted, i.e.  $|s_i| \leq |F|$  for every  $1 \leq i \leq m$ . On the other hand, the sequence of observed states  $\langle s_1, \ldots, s_m \rangle$  in  $\tau$  is the same sequence of states traversed by  $\pi$  but certain states may also be omitted. Formally,  $0 \leq |s_i|$  for every  $1 \leq i \leq m$ . This means that the transitions between two consecutive observed states in  $\tau$  may require the execution of more than a single action  $(\theta(s_i, \langle a_1, \ldots, a_k \rangle) = s_{i+1}$ , where  $k \geq 1$  is unknown and unbound). Therefore having  $\tau$  does not implies knowing the actual length of  $\pi$ .

**Definition 1** ( $\Phi$ -observation). Given a subset of fluents  $\Phi \subseteq F$  we say that  $\tau$  is a  $\Phi$ -observation of the execution of  $\pi$  on P iff, for every  $0 \le i \le m$ , each observed state  $s_i$  only contains fluents in  $\Phi$ .

For instance, the observation of the plan execution shown in Figure illustrates a  $\Phi$ -observation because only the fluents of the kind (x-coord ?v) and (y-coord ?v) are observed while the value of the fluents (next ?v1 ?v2), (state-q0) and (state-q1) is unknown.

# **Model Recognition**

The *model recognition* task is a tuple  $\langle P, M, \tau \rangle$  where:

- $P = \langle F, A, I, G \rangle$  is a *classical planning problem* such that the semantics of the actions in A is unknown because functions  $\rho(s, a)$  and/or  $\theta(s, a)$  are undefined.
- $M = \{\mathcal{M}_1, \dots, \mathcal{M}_m\}$  is a finite and non-empty set of models for the actions in A. Each model in  $\mathcal{M} \in M$  defines a different function pair  $(\rho, \theta)$ .
- $\tau$  is an *observation* of the execution of a solution plan  $\pi$  for P.

A solution to the *model recognition* task is the discrete probability distribution  $P(\mathcal{M}|\tau)$  that expresses, for each model  $\mathcal{M} \in M$ , its probability of producing the observation  $\tau$ .

According to the *Bayes* rule, the probability of an hypothesis  $\mathcal{H}$ , provided the observation  $\mathcal{O}$ , is given by the expression  $P(\mathcal{H}|\mathcal{O}) = \frac{P(\mathcal{O}|\mathcal{H})P(\mathcal{H})}{P(\mathcal{O})}$ . In the *model recognition* task, hypotheses are about the possible action models  $\mathcal{M} \in M$  while the given observation is the partially observed plan execution  $\tau$ . The  $P(\mathcal{M}|\tau)$  probability distribution can then be estimated in three steps:

- 1. Computing the *a priori* probabilities.  $P(\tau)$ , that measures how surprising is the given observation and  $P(\mathcal{M})$ , that expresses if one model is a priori more likely than the others.
- 2. Computing the conditional probability  $P(\tau|\mathcal{M})$ . Our approach is to estimate this value according to the *amount* of edits required by the model  $\mathcal{M}$  to produce a plan  $\pi_{\tau}$  such that  $\tau$  is an observation of the execution of

```
;;; Propositional encoding for inc-x(?v1 ?v2)
(pre_x-coord_v1_inc-x) (pre_next_v1_v2_inc-x) (pre_even__inc-x)
(del_x-coord_v1_inc-x) (add_x-coord_v2_inc-x)
;;; Propositional encoding for dec-x(?v1 ?v2)
(pre_x-coord_v1_dec-x) (pre_next_v2_v1_dec-x) (pre_odd___dec-x)
(del_x-coord_v1_dec-x) (add_x-coord_v2_dec-x)
;;; Propositional encoding for inc-v-odd(?v1 ?v2)
(pre_y-coord_v1_inc-y-odd) (pre_next_v1_v2_inc-y-odd)
(pre_even__inc-y-odd)
(del_y-coord_vl_inc-y-odd) (del_odd__inc-y-odd)
(add_y-coord_v2_inc-y-odd) (add_even__inc-y-odd)
;;; Propositional encoding for inc-v-even(?v1 ?v2)
(pre_y-coord_v1_inc-y-even) (pre_next_v1_v2_inc-y-even)
(pre_even__inc-y-even)
(del_y-coord_v1_inc-y-even) (del_even__inc-y-even)
(add_y-coord_v2_inc-y-even) (add_odd_inc-y-even)
;;; Propositional encoding for dec-y-odd(?v1 ?v2)
(pre_y-coord_v1_dec-y-odd) (pre_next_v2_v1_dec-y-odd)
(pre odd dec-v-odd)
(del_y-coord_v1_dec-y-odd) (del_odd__dec-y-odd)
(add_y-coord_v2_dec-y-odd) (add_even__dec-y-odd)
;;; Propositional encoding for inc-y-even(v1 ?v2)
(pre_y-coord_v1_dec-y-even) (pre_next_v2_v1_dec-y-even)
(pre even dec-y-even)
(del_y-coord_v1_dec-y-even) (del_even__dec-y-even)
(add_y-coord_v2_dec-y-even) (add_odd__dec-y-even)
```

Figure 3: Propositional encoding for the six schema from Figure 2.

 $\pi_{\tau}$  on the classical planning problem P. The *edit distance* is a similarity metric that is traditionally computed over strings or graphs and that has been proved succesful for *pattern recognition* (Masek and Paterson 1980; Bunke 1997). In this work this similarity metric refers to the edition of classical planning models.

3. Applying the Bayes rule to obtain the normalized posterior probabilities. The  $P(\mathcal{M}|\tau)$  probabilities must sum 1 for all the  $\mathcal{M} \in M$ .

### **Recognition of STRIPS models**

Here we analyze the particular instantiation of the *model recognition* task where the semantics of the actions A (i.e. the  $\rho$  and  $\theta$  functions) are specified with STRIPS action schemas. We first formalize STRIPS schemas, then define the full space of possible STRIPS schema and eventually, we introduce an *edit distance* to estimate the  $P(\mathcal{M}|\tau)$  probabilities for STRIPS models.

# A propositional encoding for STRIPS action schema

STRIPS action schema provide a compact representation for specifying classical planning action models. In more detail, a STRIPS action schema  $\xi$  with name,  $name(\xi)$ , defines a list of  $parameters\ pars(\xi)$ , and the three list of predicates  $(pre(\xi),\ del(\xi))$  and  $add(\xi)$ ) that shape the kind of fluents that can appear in the preconditions,  $negative\ effects$  and  $positive\ effects$  of the actions induced from that schema. Figure 2 shows six action schema codded in PDDL for a robot navigation in a  $n \times n$  grid.

We say that two STRIPS schemes  $\xi$  and  $\xi'$  are *comparable* iff both share the same list of parameters. For instance, we claim that the six action schema of Figure 2 are *comparable* while stack (?v1, ?v2) and pickup (?v1) from a four opertator *blocksworld* (Slaney and Thiébaux 2001) are not.

```
(:action inc-x
  :parameters (?v1 ?v2)
  :precondition
    (and (or (not (pre_x-coord_v1_inc-x)) (x-coord ?v1))
                   (pre_x-coord_v2_inc-x)) (x-coord ?v2))
         (or (not (pre_y-coord_v1_inc-x)) (x-coord ?v1))
         (or (not (pre_y-coord_v2_inc-x)) (x-coord ?v2))
                   (pre_even__inc-x)) (even))
         (or (not (pre odd inc-x)) (odd)))
         (or (not (pre_next_v1_v1_inc-x)) (next ?v1 ?v1)))
         (or (not (pre_next_v1_v2_inc-x)) (next ?v1 ?v2)))
         (or (not (pre next v2 v1 inc-x)) (next ?v2 ?v1)))
         (or (not (pre_next_v2_v2_inc-x)) (next ?v2 ?v2)))
    :effect (and
       (when (del_x-coord_vl_inc-x) (not (x-coord ?v1)))
       (when (del_x-coord_v2\_inc-x) (not (x-coord ?v2)))
       (when (del y-coord v1 inc-x) (not (x-coord ?v1)))
       (when (del_y-coord_v2_inc-x) (not (x-coord ?v2)))
       (when (del_even__inc-x) (not (even)))
       (when (del_odd__inc-x) (not (odd)))
       (when (del_next_v1_v1_inc-x) (not (next ?v1 ?v1)))
       (when (del next v1 v2 inc-x) (not (next ?v1 ?v2)))
       (when (del_next_v2_v1_inc-x) (not (next ?v2 ?v1)))
       (when (del_next_v2_v2_inc-x) (not (next ?v2 ?v2)))
       (when (add_x-coord_v1_inc-x) (x-coord ?v1))
       (when (add_x-coord_v2_inc-x) (x-coord ?v2))
       (when (add y-coord v1 inc-x) (x-coord ?v1))
             (add_y-coord_v2_inc-x) (x-coord ?v2))
       (when (add_even__inc-x) (even))
(when (add_odd__inc-x) (odd))
             (add_next_v1_v1_inc-x) (next ?v1 ?v1))
       (when (add_next_v1_v2_inc-x) (next ?v1 ?v2))
(when (add_next_v2_v1_inc-x) (next ?v2 ?v1))
       (when (add_next_v2_v2_inc-x) (next ?v2 ?v2)))
```

Figure 4: Programmable version of the inc-x(?v1,?v2) schema for robot navigation in a  $n \times n$  grid.

Last but not least, two STRIPS models  $\mathcal{M}$  and  $\mathcal{M}'$  are *comparable* iff there exists a bijective function  $\mathcal{M} \mapsto \mathcal{M}^*$  that maps every action schema  $\xi \in \mathcal{M}$  to a comparable schema  $\xi' \in \mathcal{M}'$  and vice versa.

Given a STRIPS action schema  $\xi$ , a propositional encoding for the *preconditions*, *negative* and *positive* effects of that schema can be represented with fluents [pre|del|add]-p- $name(\xi)$ . Figure 3 shows the propositional encoding for the six action schema previously defined in Figure 2. The interest of having a propositional encoding for STRIPS action schema is that it allows to define *programmable actions* that is, actions whose semantics is given by the value of these particular fluents on the current state. For instance, Figure 4 shows the programmable version of the inc-x(?v1,?v2) schema for robot navigation in a  $n \times n$  grid. Note that this programmable schema, when the fluents of Figure 3 holds, behaves exactly as is defined in Figure 2. Further it allows to swap the semantics of two programable schemas provided that they are *comparable*.

# The space of STRIPS schema

The space of possible STRIPS schema is bound by these elements:

- 1.  $\Psi$ , is the set of *state predicates* that shape the propositional state variables of the model.
- 2.  $pars(\xi)$ , the list of *parameters* of the corresponding schema  $\xi$ .

For each action schema  $\xi$ , given the set of *state predicates*  $\Psi$ , then  $F_{\xi}$  defines the subset of elements that can appear

in the preconditions and effects of that action schema. For the actions schema defined in 2 this set is the same and contains the following ten elements,  $\{x-\operatorname{coord}(v_1), x-\operatorname{coord}(v_2), y-\operatorname{coord}(v_1), y-\operatorname{coord}(v_2), \operatorname{odd}(), \operatorname{even}(), \operatorname{next}(v_1,v_1), \operatorname{next}(v_1,v_2), \operatorname{next}(v_2,v_1), \operatorname{next}(v_2,v_2)\}$ 

3.  $\mathcal{C}$ , a set of *syntactic constraints*. These constraints include the STRIPS constraints (we assume that  $\operatorname{eff}^-(a) \subseteq \operatorname{pre}(a)$ ,  $\operatorname{eff}^-(a) \cap \operatorname{eff}^+(a) = \emptyset$  and  $\operatorname{pre}(a) \cap \operatorname{eff}^+(a) = \emptyset$ ) and domain-specific constraints. For instace, in a *robot navigation* domain like the modeled in Figure 2, predicates  $\operatorname{even}()$  and  $\operatorname{odd}()$  are exclusive so they cannot hold a the same time.

Predicates  $\Psi_a$  also shape an additional set of objects  $(\Omega \cap \Omega_v = \emptyset)$  that we denote as *variable names* and that is bound by the maximum arity of the given action headers, i.e.,  $\Omega_v = \{v_i\}_{i=1}^{\max_\xi ar(\xi)}$ . For instance, in a  $5 \times 5$  grid (like the one in Figure and modeled in Figure 2)  $\Omega = \{o_1, o_2, o_3, o_4, o_5\}$  while  $\Omega_v = \{v_1, v_2\}$  because all the actions schema have arity two.

**Definition 2** (The space of STRIPS models). Given the set of state predicates  $\Psi$ , the set of variable names  $\Omega_v$  and the set of syntactic constraints  $\mathcal{C}$ , the space of models for a given op predicate  $p_{\xi}$  that represents the header of an action schema  $\xi$  is given by three sets of inter

The size of the space of possible STRIPS models for an action schema  $\xi$  is given by the expression,  $2^{2\times|F_\xi|}$  (STRIPS constraints require negative effects appearing as preconditions, negative effects cannot be positive effects at the same time and also, positive effects cannot appear as preconditions). For the mentioned navigation model,  $2^{2\times|F_\xi|}=1,048,576$  for every action schema.

#### The STRIPS edit distance

We define two edit *operations* on a STRIPS model  $\mathcal{M} \in M$ :

- Deletion. A fluent  $pre_f(\xi)/del_f(\xi)/add_f(\xi)$  is removed from the operator schema  $\xi \in \mathcal{M}$ , such that  $f \in F_{\xi}$ .
- Insertion. A fluent  $pre_f(\xi)/del_f(\xi)/add_f(\xi)$  is added to the operator schema  $\xi \in \mathcal{M}$ , s.t.  $f \in F_{\xi}$ .

We can now formalize an *edit distance* that quantifies how similar two given STRIPS action models are. The distance is symmetric and meets the *metric axioms* provided that the two *edit operations*, deletion and insertion, have the same positive cost.

**Definition 3.** Let  $\mathcal{M}$  and  $\mathcal{M}'$  be two comparable STRIPS action models. The **edit distance**  $\delta(\mathcal{M}, \mathcal{M}')$  is the minimum number of edit operations that is required to transform  $\mathcal{M}$  into  $\mathcal{M}'$ .

Since  $F_v$  is a bound set, the maximum number of edits that can be introduced to a given action model defined within  $F_v$  is bound as well.

**Definition 4.** The maximum edit distance of an STRIPS model  $\mathcal{M}$  built from the set of possible elements  $F_v$  is  $\delta(\mathcal{M},*) = \sum_{\mathcal{E} \in \mathcal{M}} 3 \times |F_{\mathcal{E}}|$ .

We define now an edit distance to asses the matching of a STRIPS model with respect to an observation of a plan execution.

**Definition 5.** Given  $\tau$ , an observation of the execution of a plan for solving P and  $\mathcal{M}$ , a STRIPS action model built from  $F_v$ . The **observation edit distance**,  $\delta(\mathcal{M}, \tau)$ , is the minimal edit distance from  $\mathcal{M}$  to any comparable model  $\mathcal{M}'$  s.t.  $\mathcal{M}'$  produces a plan  $\pi_{\tau}^*$  optimal for P and compliant with  $\tau$ ;

$$\delta(\mathcal{M}, \tau) = \min_{\forall \mathcal{M}' \to \tau} \delta(\mathcal{M}, \mathcal{M}')$$

The  $\delta(\mathcal{M}, \tau)$  distance could also be defined assessing the edition required by the observed plan execution to match the given model. This implies defining *edit operations* that modify  $\tau$  instead of  $\mathcal{M}$  (Sohrabi, Riabov, and Udrea 2016). Our definition of the *observation edit distance* is more practical since normally  $F_v$  is smaller than F. In practice, the number of *variable objects* is smaller than the number of objects in a planning problem.

# The $P(\mathcal{M}|\tau)$ probability for STRIPS models

Since we assume that the dynamics of the actions  $A \in P$  is specified with STRIPS action schemas, we can estimate the  $P(\mathcal{M}|\tau)$  probability as follows:

- 1. Assuming a priori all action models  $\mathcal{M} \in M$  are equiprobable. There are no reasons to assume that one model is a priori more likely than the others so,  $P(\mathcal{M}) = \frac{1}{\prod_{E \in \mathcal{M}} 2^{2 \times |F_{\xi}|}}$ .
- 2. Assuming that all the observations of plan executions with maximum n observed actions and m observed states are equiprobable then,  $P(\tau) = \frac{1}{|A|^n \times 2^{m \times |F|}}$ .
- 3. Estimating the conditional probability  $P(\tau|\mathcal{M})$  by mapping the *observation edit distance* into a [0,1] likelihood,  $1-\frac{\delta(\mathcal{M},\tau)}{\delta(\mathcal{M},*)}$ .

# Model recognition as planning

This section shows that, when the dynamics of the actions  $A \in P$  are specified with STRIPS action schemas, then  $\delta(\mathcal{M},\tau)$  can be computed with a compilation of a classical planning with conditional effects. The intuition behind this compilation is that a solution to the resulting classical planning task is a sequence of actions that:

- 1. Edits the action model  $\mathcal{M}$  to build  $\mathcal{M}'$ . A solution plan starts with a *prefix* that modifies the preconditions and effects of the action schemes in  $\mathcal{M}$  using to the two *edit operations* defined above, *deletion* and *insertion*.
- 2. Validates the edited model  $\mathcal{M}'$ . The solution plan continues with a postfix that:
- (a) Induces an optimal solution plan  $\pi_{\tau}^*$  for the original classical planning problem P.
- (b) Validates that  $\tau$  is an observation of the execution of  $\pi_{\tau}^*$  on the classical planning problem P.

Figure 5 shows the plan with a prefix (steps [0,1]) for editing a given *blockswold* model where the positive effects (handempty) and (clear ?v1) of the stack schema are

```
00 : (insert_add_handempty_stack)
01 : (insert_add_clear_stack_varl)
02 : (apply_unstack blockB blockA i1 i2)
03 : (apply_putdown blockB i2 i3)
04 : (apply_pickup blockA i3 i4)
05 : (apply_stack blockA blockB i4 i5)
06 : (validate_1)
```

Figure 5: Plan for editing (steps [0-1]) and validating (steps [2-6]) a given STRIPS planning model for the *blockswold*.

missing. The postfix of the plan (steps [2,6]) validates the edited action model at the observation of a four action plan for inverting a two-block tower where intermediate states,  $s_1$ ,  $s_2$  and  $s_3$ , are unobserved.

Note that our interest is not in  $\mathcal{M}'$ , the edited model resulting from the compilation, but in the number of required edit operations (insertions and deletions) required by  $\mathcal{M}'$  to be validated. In the example of Figure 5  $\delta(\mathcal{M},\tau)=2$  and  $\delta(\mathcal{M},*)=3\times2\times(11+5)$  since there are 4 action schemes (pickup, putdown, stack and unstack) s.t.  $|F_v|=|F_{\rm stack}|=|F_{\rm unstack}|=11$  while  $|F_{\rm pickup}|=|F_{\rm putdown}|=5$ .

# The compilation formalization

Conditional effects allow us to compactly define our compilation. Given a STRIPS model  $\mathcal{M} \in M$  and the observation  $\tau$  of the execution of a plan for solving  $P = \langle F, A, I, G \rangle$ , our compilation outputs a classical planning task with conditionl effects  $P' = \langle F', A', I', G' \rangle$  such that:

- F' contains:
  - The original fluents F.
  - Fluents  $pre_f(\xi)/del_f(\xi)/add_f(\xi)$  modeling the space of STRIPS models.
  - The fluents  $F_{\pi} = \{plan(name(a_i), \Omega^{ar(a_i)}, i)\}_{1 \leq i \leq n}$  to code the  $i^{th}$  action in  $\tau$ . The static facts  $next_{i,i+1}$  and the fluents  $at_i$ ,  $1 \leq i < n$ , are also added to iterate through the n steps of  $\tau$ .
  - The fluents  $\{test_j\}_{1 \le j \le m}$ , indicating the state observation  $s_j \in \tau$  where the action model is validated.
  - The fluents  $mode_{edit}$  and  $mode_{val}$  to indicate whether the operator schemas are edited or validated.
- I' extends the original initial state I with the fluent  $mode_{edit}$  set to true as well as the fluents  $F_{\pi}$  plus fluents  $at_1$  and  $\{next_{i,i+1}\}$ ,  $1 \leq i < n$ , for tracking the plan step where the action model is validated. Our compilation assumes that initially  $\mathcal{M}'$  is defined as  $\mathcal{M}$ . Therefore fluents  $pre_f(\xi)/del_f(\xi)/add_f(\xi)$  hold as given by  $\mathcal{M}$ .
- $G' = G \bigcup \{at_n, test_m\}.$
- A' comprises three kinds of actions with conditional effects:
  - 1. Actions for *editing* operator schema  $\xi \in \mathcal{M}$ :
    - Actions for adding a precondition  $f \in F_v(\xi)$  from the action schema  $\xi \in \mathcal{M}$ .

```
\begin{split} & \operatorname{pre}(\operatorname{programPre}_{f,\xi}) = \{ \neg pre_f(\xi), \neg del_f(\xi), \neg add_f(\xi), mode_{edit} \}, \\ & \operatorname{cond}(\operatorname{programPre}_{f,\xi}) = \{\emptyset\} \rhd \{pre_f(\xi)\}. \end{split}
```

- Actions for adding a *negative* or *positive* effect  $f \in F_v(\xi)$  to the action schema  $\xi \in \mathcal{M}$ .

```
\begin{split} \operatorname{pre}(\operatorname{programEff_{f,\xi}}) = & \{ \neg del_f(\xi), \neg add_f(\xi), mode_{edit} \}, \\ \operatorname{cond}(\operatorname{programEff_{f,\xi}}) = & \{ pre_f(\xi) \} \rhd \{ del_f(\xi) \}, \\ & \{ \neg pre_f(\xi) \} \rhd \{ add_f(\xi) \}. \end{split}
```

Besides these actions, the A' set also contains the actions for *deleting* a precondition and a negative/positive effect.

2. Actions for *applying* an edited operator schema  $\xi \in \mathcal{M}$  bound with objects  $\omega \subseteq \Omega^{ar(\xi)}$ . Since operators headers are given as input, the variables  $pars(\xi)$  are bound to the objects in  $\omega$  that appear at the same position. Figure 6 shows the PDDL encoding of the action for applying a programmed operator stack from blocksworld.

```
\begin{split} \operatorname{pre}(\mathsf{apply}_{\xi,\omega}) = & \{pre_f(\xi) \implies p(\omega)\}_{\forall p \in \Psi, f = p(pars(\xi))} \\ & \cup \{\neg mode_{val}\}, \\ \operatorname{cond}(\mathsf{apply}_{\xi,\omega}) = & \{del_f(\xi)\} \rhd \{\neg p(\omega)\}_{\forall p \in \Psi, f = p(pars(\xi))}, \\ & \{add_f(\xi)\} \rhd \{p(\omega)\}_{\forall p \in \Psi, f = p(pars(\xi))}, \\ & \{mode_{edit}\} \rhd \{\neg mode_{edit}\}, \\ & \{\emptyset\} \rhd \{mode_{val}\}. \end{split}
```

When the observation  $\tau$  includes observed actions, then the extra conditional effects  $\{at_i, plan(name(a_i), \Omega^{ar(a_i)}, i)\} \rhd \{\neg at_i, at_{i+1}\}_{\forall i \in [1,n]}$  are included in the apply $\xi, \omega$  actions to validate that actions are applied, exclusively, in the same order as they appear in  $\tau$ .

3. Actions for *validating* the partially observed state  $s_j \in \tau$ ,  $1 \le j < m$ .

```
\begin{aligned} & \mathsf{pre}(\mathsf{validate_j}) = & s_j \cup \{test_{j-1}\}, \\ & \mathsf{cond}(\mathsf{validate_j}) = \{\emptyset\} \rhd \{\neg test_{j-1}, test_j, \neg mode_{val}\}. \end{aligned}
```

### **Evaluation**

To evaluate the empirical performance of *model recognition* as planning we defined a set of possible STRIPS models, each representing a different *Turing Machine*, but all sharing the same set of *machine states* and same *tape alphabet*.

# **Experimental setup**

We randomly generated a  $M = \{\mathcal{M}_1, \dots, \mathcal{M}_{100}\}$  set of one-hundred different *Turing Machines* where each  $\mathcal{M} \in M$  is a seven-symbol six-state *Turing Machine*. The classical planning frame  $\Phi = \langle F, A \rangle$  is the same for all the *Turing Machines*, there is an  $a \in A$  action for each pair of tape symbol and non-terminal state machine, while the  $\theta(s,a)$  function is defined differently for each the machines (using a different STRIPS action model).

```
(:action apply_stack
 :parameters (?o1 - object ?o2 - object)
  ·precondition
   (and (or (not (pre_on_stack_v1_v1)) (on ?o1 ?o1))
        (or (not (pre_on_stack_v1_v2)) (on ?o1 ?o2))
        (or (not (pre_on_stack_v2_v1)) (on ?o2 ?o1))
        (or (not (pre on stack v2 v2)) (on ?o2 ?o2))
       (or (not (pre_ontable_stack_v1)) (ontable ?o1))
       (or (not (pre_ontable_stack_v2)) (ontable ?o2))
        (or (not (pre_clear_stack_v1)) (clear ?o1))
        (or (not (pre_clear_stack_v2)) (clear ?o2))
        (or (not (pre_holding_stack_v1)) (holding ?o1))
        (or (not (pre_holding_stack_v2)) (holding ?o2))
        (or (not (pre_handempty_stack)) (handempty)))
  :effect
   (and (when (del_on_stack_v1_v1) (not (on ?o1 ?o1)))
        (when (del_on_stack_v1_v2) (not (on ?o1 ?o2)))
        (when (del_on_stack_v2_v1) (not (on ?o2 ?o1)))
        (when (del_on_stack_v2_v2) (not (on ?o2 ?o2)))
        (when (del_ontable_stack_v1) (not (ontable ?o1)))
        (when (del_ontable_stack_v2) (not (ontable ?o2)))
        (when (del_clear_stack_v1) (not (clear ?o1)))
        (when (del_clear_stack_v2) (not (clear ?o2)))
        (when (del_holding_stack_v1) (not (holding ?o1)))
       (when (del_holding_stack_v2) (not (holding ?o2)))
       (when (del_handempty_stack) (not (handempty)))
        (when (add_on_stack_v1_v1) (on ?o1 ?o1))
        (when (add_on_stack_v1_v2) (on ?o1 ?o2))
        (when (add_on_stack_v2_v1) (on ?o2 ?o1))
        (when (add on stack v2 v2) (on ?o2 ?o2))
        (when (add_ontable_stack_v1) (ontable ?o1))
        (when (add_ontable_stack_v2) (ontable ?o2))
        (when (add_clear_stack_v1) (clear ?o1))
        (when (add_clear_stack_v2) (clear ?o2))
        (when (add_holding_stack_v1) (holding ?o1))
        (when (add_holding_stack_v2) (holding ?o2))
        (when (add_handempty_stack) (handempty))
```

Figure 6: PDDL action for applying an already programmed schema *stack* (implications are coded as disjunctions).

(when (modeProg) (not (modeProg)))))

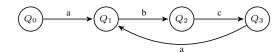


Figure 7: Four-symbol four-state regular automata for recognizing the  $\{(abc)^n:n\geq 1\}$  language.

We randomly choose a machine  $\mathcal{M} \in M$  and produce the observation  $\tau$  of a fifty-step execution plan. Finally, we follow our *model recognition as planning* method to identify the *Turing Machine* that produced  $\tau$ . This experiment is repeated for different amounts of missing information in the input trace  $\tau$ : unknown applied transitions, unknown internal machine state and unknown values of several tape cells.

**Reproducibility** MADAGASCAR is the classical planner we used to solve the instances that result from our compilations for its ability to deal with dead-ends (Rintanen 2014). Due to its SAT-based nature, MADAGASCAR can apply the actions for editing preconditions in a single planning step (in parallel) because there is no interaction between them. Actions for editing effects can also be applied in a single planning step, thus significantly reducing the planning horizon.

The compilation source code, evaluation scripts and benchmarks (including the used training and test sets) are fully available at this anonymous repository so any experimental data reported in the paper can be reproduced.

# Recognition of Regular Automatae

We analyze now *model recognition* when the input set of given set of models represent different *regular automatae*.

A Regular automatae is a tuple  $\mathcal{M} = \langle Q, q_o, Q_{\perp}, \Sigma \square, \delta \rangle$ :

- Q is a finite and non-empty set of machine states with the initial state  $q_0 \in Q$  and the terminal states  $Q_{\perp} \subseteq Q$ .
- $\Sigma$  is the *input alphabet*, a finite non-empty set of symbols and the *blank symbol*  $\square \in \Upsilon$  (the only symbol allowed to occur on the tape infinitely often).
- $\delta: \Sigma \times (Q \setminus Q_{\perp}) \to \Sigma \times Q \times \{left, right\}$  is the *transition* function. For each pair of tape symbol and non-terminal machine state  $\delta$  defines: (1), the tape symbol to print at the current position of the header (2), the new state of the machine and (3), whether the header is shifted *left* or *right* after the print operation. If  $\delta$  is not defined for the current pair of tape symbol and machine state, the machine halts.

Figure 7 illustrate a four-symbol four-state regular automata for recognizing the  $\{(abc)^n : n \ge 1\}$  language. The input alphabet is  $\Sigma = \{a, b, c, \square\}$ , and the machine states are  $Q = \{q_0, q_1, q_2, q_3\}$  (where  $q_3$  is the only acceptor state).

The STRIPS action schema of Figure ?? models the rule  $a, q_0 \rightarrow q_1$  of the *regular automatae* defined in Figure 7. The full encoding of the *automata* of Figure 7 produces a total of sixteen STRIPS action schema.

### **Recognition of Turing Machines**

We analyze now *model recognition* when the input set of given set of models represent different *Turing machines*.

```
\label{eq:continuous} \begin{tabular}{ll} (:action transition-1 & ;;; a, q_0 \rightarrow q_1 \\ :parameters & (?x ?xr) \\ :precondition & (and (head ?x) (symbol-a ?x) (state-q_0) \\ & & (next ?x ?xr)) \\ :effect & (and (not (head ?x)) \\ & & (not (symbol-a ?x)) (not (state-q_0)) \\ & & (head ?xr) & (state-q_1))) \end{tabular}
```

Figure 8: STRIPS action schema that models the transition  $q_0 \rightarrow q_1$  of the automata defined in Figure 7.

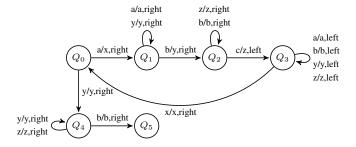


Figure 9: Seven-symbol six-state *Turing Machine* for recognizing the  $\{a^nb^nc^n: n \geq 1\}$  language.

A *Turing machine* is a tuple  $\mathcal{M} = \langle Q, q_o, Q_\perp, \Sigma, \Upsilon, \square, \delta \rangle$ :

- Q is a finite and non-empty set of machine states with the *initial state*  $q_0 \in Q$  and the *terminal states*  $Q_{\perp} \subseteq Q$ .
- $\Sigma$  is the *tape alphabet*, a finite non-empty set of symbols with the *input alphabet*  $\Upsilon \subseteq \Sigma$  (symbols allowed to initially appear in the tape) and the *blank symbol*  $\square \in \Upsilon$  (the only symbol allowed to occur on the tape infinitely often).
- δ: Σ×(Q\Q⊥) → Σ×Q×{left, right} is the transition function. For each pair of tape symbol and non-terminal machine state δ defines: (1), the tape symbol to print at the current position of the header (2), the new state of the machine and (3), whether the header is shifted left or right after the print operation. If δ is not defined for the current pair of tape symbol and machine state, the machine halts.

Figure 9 illustrate a seven-symbol six-state *Turing Machine* for recognizing the  $\{a^nb^nc^n:n\geq 1\}$  language. The *tape alphabet* is  $\Sigma=\{a,b,c,x,y,z,\Box\}$ , the *input alphabet*  $\Upsilon=\{a,b,c,\Box\}$  and the machine states are  $Q=\{q_0,q_1,q_2,q_3,q_4,\underline{q_5}\}$  (where  $\underline{q_5}$  is the only acceptor state).

The STRIPS action schema of Figure 10 models the rule  $a,q_0 \rightarrow x,r,q_1$  of the *Turing Machine* defined in Figure 9. The full encoding of the *Turing Machine* of Figure 9 produces a total of sixteen STRIPS action schema.

With regard to our STRIPS model for *Turing Machines*, executions of a *Turing Machine* are definable as an action sequence  $\langle a_1,\ldots,a_n\rangle$  that induces the *state trajectory*  $\langle s_0,s_1,\ldots,s_n\rangle$  such that  $a_i$   $(1\leq i\leq n)$  is applicable in  $s_{i-1}$  and generates the successor state  $s_i=\theta(s_{i-1},a_i)$ . For instance, the execution of the *Turing Machine* defined in Figure 9 with an initial tape  $abc\square\square\square\ldots$  produces the eight-action plan  $(a,q_0\to x,r,q_1),(b,q_1\to y,r,q_2),(c,q_2\to y,r,q_1)$ 

```
\label{eq:continuous} \begin{tabular}{ll} (:action transition-1 & ;;; a, q_0 \rightarrow x, r, q_1 \\ :parameters & (?x1 ?x ?xr) \\ :precondition & (and (head ?x) & (symbol-a ?x) & (state-q_0) \\ & & & & (next ?x1 ?x) & (next ?x ?xr)) \\ :effect & (and & (not & (head ?x)) & & & (not & (symbol-a ?x)) & (not & (state-q_0)) \\ & & & & & (head ?xr) & (symbol-x ?x) & (state-q_1))) \\ \end{tabular}
```

Figure 10: STRIPS action schema that models the transition  $a, q_0 \rightarrow x, r, q_1$  of the Turing Machine defined in Figure 9.

```
z, l, q_3), (y, q_3 \to y, l, q_3), (x, q_3 \to x, r, q_0), (y, q_0 \to y, r, q_4), (z, q_4 \to z, r, q_4), (\Box, q_4 \to \Box, r, q_5).
```

Assuming that the actual applied transitions is unknown means that the observation  $\tau$  of the execution of a Turing Machine contains no actions, it is simply a sequence of states  $\tau = \langle s_1, \dots, s_m \rangle$ . Further, assuming that the internal machine state is unknown means that  $\tau$  is a  $\Phi$ -observation and that the  $\Phi$  subset does not contain ( $\mathtt{state}-q$ ) fluents, with  $q \in Q$  and  $q \neq q_0$ . Finally, assuming that the values of several tape cells is unknown means that fluents of the kind ( $\mathtt{symbol}-\sigma$ ?x) are missing (i.e. unobserved) for some state  $s_i \in \tau$  s.t.  $1 \leq i \leq m$ . These facts affects to the a priori probability of the possible observations.

On the other hand, the model edition for Turing Machines is limited to these subset of possible positive effects: (head ?xr) or (head ?xl), (symbol- $\sigma$ ?x) with  $\sigma \in \Sigma$  and last but not least, (state-q) with  $q \in Q$ . No other positive effects, preconditions, or negative effects are required to be edited. This fact reduces the possible edition operations and affects to the a priori probability of the possible action models.

### **Results**

# **Conclusions**

The task of *model recognition* can also be understood as a classification task where each class is represented with a different planning model and the observed plan execution is the example to classify. With this regard, planning model that is associated to a class is acting as a class prototype that the summarizes all the plan executions that could be synthesized with that model or in other words, all the examples that belong to that class.

In this work we do not assume that the observed agent is acting rationally, like in *plan recognition as planning* (Ramírez 2012; Ramírez and Geffner 2009). A related approach is recently followed for *model reconciliation* (Chakraborti et al. 2017) where model edition is used to conform the PDDL models of two different agents.

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