

Semantics for coinductive structures in stochastic processes

Infinite product measures in quasi-Borel spaces

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CALCO 2025
Glasgow, UK - June 2025

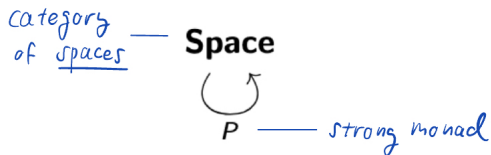
Overview

1. Random streams - a coalgebraic interpretation

2. Quasi-Borel spaces

3. Infinite product measures in QBS

Compositional Probability [Giry, 1982]



Monadic probabilistic programs define probabilistic models.

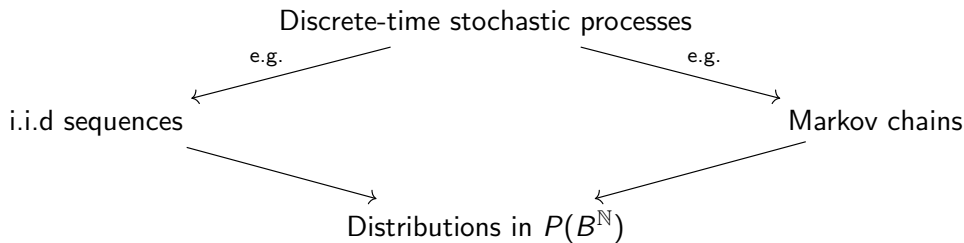
Space and P may be:

- Polish spaces and Borel/Radon measures
- Standard Borel spaces and probability measures

This talk:

- quasi-Borel spaces (QBS) and probability distributions [Heunen et al., 2017]

Stochastic Processes

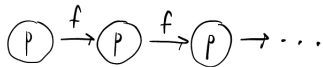


Many stochastic processes are generated step by step.

$$p \in \mathcal{PB}$$



$$b \in \mathcal{PB}, \quad f: \mathcal{B} \rightarrow \mathcal{PB}$$



Stochastic Processes as Random Streams

Deterministic streams

$B^{\mathbb{N}}$

Universal property [Rutten, 2000]:

final coalgebra of $X \mapsto B \times X$

step : $B^{\mathbb{N}} \xrightarrow{\langle \text{head}, \text{tail} \rangle} B \times B^{\mathbb{N}}$

Random streams

$P(B^{\mathbb{N}})$

$B^{\mathbb{N}}$ is final coalgebra in $\text{Kleisli}(P)$
in QBS for standard-Borel spaces,
but it is open in general

$B^{\mathbb{N}}$ is the final coalgebra of $X \mapsto B \times X$ in QBS:

$$\text{step} : B^{\mathbb{N}} \xrightarrow{\langle \text{head}, \text{tail} \rangle} B \times B^{\mathbb{N}}$$

$B^{\mathbb{N}}$ is the final coalgebra of $X \mapsto B \times X$ in QBS:

$$\text{step} : B^{\mathbb{N}} \xrightarrow{\langle \text{head}, \text{tail} \rangle} B \times B^{\mathbb{N}}$$

Question: Is $B^{\mathbb{N}}$ also the final coalgebra of $X \mapsto B \times X$ in $\text{Kleisli}(P)$?

For all Kleisli morphisms

$$S \xrightarrow{f} P(B \times S)$$

does there exist a unique Kleisli morphism

$$\text{chain}(f) : S \rightarrow P(B^{\mathbb{N}})$$

such that

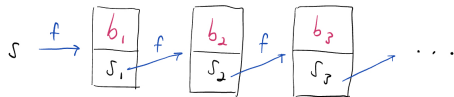
$$\begin{array}{ccc} S & \xrightarrow{f} & B \times S \\ \text{chain}(f) \downarrow & & \downarrow id \times \text{chain}(f) \\ B^{\mathbb{N}} & \xrightarrow{\text{step}} & B \times B^{\mathbb{N}} \end{array}$$

As a probabilistic program (LazyPPL - Haskell):

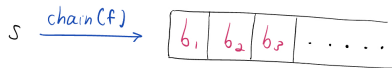
```
chain :: (S -> Prob (B,S))
      -> S -> Prob (Stream B)
chain f s = do {(b,s') <- f s;
                bs <- chain f s';
                return (b:bs)}
```

Iteratively applies f in an infinite chain

$$f : S \rightarrow P(B \times S)$$



$$\text{chain}(f) : S \rightarrow P(B^{\mathbb{N}})$$



Question: Is $\text{chain}(f)$ well-defined for any S and f ?

Infinite Product Measures in QBS

This talk: initial progress in proving that for all qbs B , $\text{chain}(f) : S \rightarrow P(B^{\mathbb{N}})$ is a well-defined qbs morphism for

$$\begin{aligned} S &= (PB)^{\mathbb{N}} \\ f : (PB)^{\mathbb{N}} &\rightarrow P(B \times (PB)^{\mathbb{N}}) \\ (p_n)_{n \geq 1} &\mapsto p_1 \otimes \delta(p_2, p_3, \dots) \end{aligned}$$

Then $\text{chain}(f)$ describes the stateless processes.

$$(p_n)_{n \geq 1} \in PB$$
A diagram illustrating an infinite sequence of elements. It consists of a series of circles, each containing a label. The first three circles are labeled p_1 , p_2 , and p_3 . This is followed by an ellipsis (\dots), then a circle labeled p_n , and another ellipsis (\dots). The circles are drawn with a hand-drawn style.

$$\begin{array}{ccc}
 (PB)^{\mathbb{N}} & \xrightarrow{(p_n)_{n \geq 1} \mapsto p_1 \otimes \delta_{(p_2, p_3, \dots)}} & B \times (PB)^{\mathbb{N}} \\
 \downarrow \text{chain}(f) & & \downarrow id \times \text{chain}(f) \\
 B^{\mathbb{N}} & \xrightarrow{\text{step}} & B \times B^{\mathbb{N}}
 \end{array}$$

$$\text{chain}(f)(p_1, p_2, p_3, \dots) = p_1 \otimes \text{chain}(f)(p_2, p_3, \dots)$$

Program interpretation:

```

prod :: Stream (Prob b) -> Prob (Stream b)
prod (p:ps) = do {x <- p; xs <- prod ps; return (x:xs)}

```

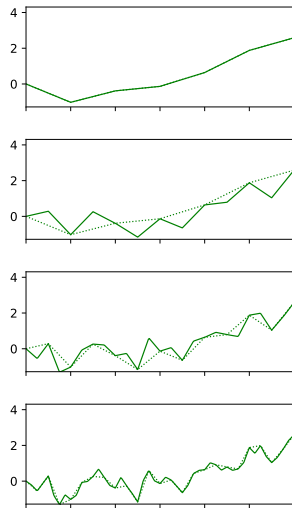
Intuition: the sequence of measures $(p_n)_{n \geq 1}$ are *independently stitched together*.

Application: Brownian motion in QBS

- [Karatzas and Shreve, 1998]: approximate Brownian motion by incrementally adding up a sequence of independent random functions.
- In QBS, can coinductively generate a random stream of functions

$$P((\mathbb{R}^{[0,T]})^{\mathbb{N}})$$

approximating Brownian motion.



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Higher order functions in probability

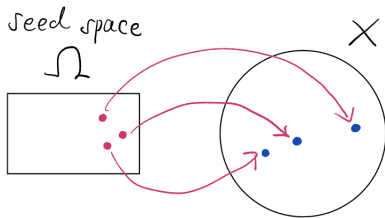
Theorem ([Aumann, 1961])

The category of measurable spaces is not Cartesian-closed.

Hence measure spaces cannot give semantics to higher-order PPLs.

Solution: **Quasi-Borel spaces (QBS)**

Random elements in X are
represented by $\Omega \rightarrow X$ where $\Omega = [0, 1]$.



- Measurable spaces use σ -algebras to implicitly specify which $\Omega \rightarrow X$ to admit as random elements.
- Quasi-Borel spaces axiomatise the admissible random elements $\Omega \rightarrow X$.

Quasi-Borel Spaces [Heunen, K, S, Yang, 2017]

Definition

A **quasi-Borel space** (X, M_X) is:

- X : set of points
- $M_X \subset [\Omega \rightarrow X]$: set of **admissible random elements**, such that:
 - it contains all the constant functions,
 - $\alpha \in M_X, f : \Omega \rightarrow \Omega$ measurable $\implies \alpha \circ f \in M_X$,
 - $\Omega = \bigcup_{i \in \mathbb{N}} S_i$ is Borel partition, $\alpha_i \in M_X, i \in \mathbb{N} \implies \lambda r. \text{case } r \text{ of } r \in S_i \rightarrow \alpha_i(r) \in M_X$.

where $\Omega = [0, 1]$.

Quasi-Borel Spaces [Heunen, K, S, Yang, 2017]

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where $\Omega = [0, 1]$.

Definition

A map $f : X \rightarrow Y$ is a **qbs morphism** if

$$(\Omega \xrightarrow{\alpha} X) \in M_X \implies (\Omega \xrightarrow{\alpha} X \xrightarrow{f} Y) \in M_Y$$

Quasi-Borel Spaces [Heunen, K, S, Yang, 2017]

Quasi-Borel spaces form a category **QBS**.

Properties of QBS:

- QBS is **Cartesian-closed**.
- QBS contains **standard-Borel spaces** as a full subcategory.
- QBS supports a **strong commutative probability monad** suitable for probabilistic programming.

QBS Probability Monad

$$\Sigma_{M_X} := \{U \subset X \mid \forall \alpha \in M_X. \alpha^{-1}U \in \Sigma_\Omega\}$$

Fix a probability measure μ on (Ω, Σ_Ω) . Then $\forall \alpha \in M_X$,

$$\begin{aligned} \alpha_*\mu &: \Sigma_{M_X} \rightarrow [0, 1] \\ U &\mapsto \mu(\alpha^{-1}U) \end{aligned}$$

is a probability measure on Σ_{M_X} , called the **pushforward measure** of α w.r.t μ .

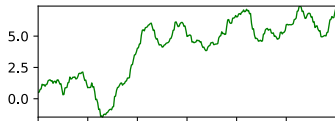
$$\int_X k \, d\alpha_*\mu = \int_\Omega (k \circ \alpha) \, d\mu$$

$\alpha, \beta \in M_X$ are **equal in law** if $\alpha_*\mu = \beta_*\mu$ as measures on Σ_{M_X} .

$$PX := M_X / \sim \text{ where } \sim = \text{law equality}$$

Random Functions

- **Random functions** are distributions on functions.
 - e.g. Brownian motion, stochastic differential equations (random functions $[0, \infty) \rightarrow \mathbb{R}$).



- Defined as parametrised family of random variables $(B_t)_{t \geq 0}$, $B_t : \Omega \rightarrow \mathbb{R} \ \forall t \geq 0$.
- Random functions $X \rightarrow Y$ are explicitly modelled by $P(X^Y)$ in QBS.
- Cartesian-closure of QBS allows us to venture beyond standard probability to directly use function spaces.

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Defining $\text{Ind} B^{\mathbb{N}}$

Characterising independence in QBS:

1. every finite prefix is independent of its suffix



2. satisfies a 0-1 law

Defining $\text{Ind} B^{\mathbb{N}}$

The strength and commutativity of P gives the operation

$$\otimes : PX \times PY \rightarrow P(X \times Y)$$

Hence for all finite n we have an operation

$$(PB)^n \rightarrow P(B^n) \\ (p_1, \dots, p_n) \mapsto p_1 \otimes \dots \otimes p_n$$

Independence for infinite n ?

QBS Products

The sigma algebra for the QBS product cannot be characterised by the product sigma algebra.

$$\bigotimes_n \Sigma_{M_{X_n}} \subsetneq \Sigma_{M_{\prod_n X_n}}$$

Defining $\text{Ind} B^{\mathbb{N}}$

Definition (0-1 law)

For all $n \geq 1$, define the sub sigma algebras of $\Sigma_{B^{\mathbb{N}}}$:

$$\Sigma_n := \left\{ \prod_{i=1}^n X_i \times S \mid S \in \Sigma_{\prod_{i \geq n+1} X_i} \right\}$$

$\bigcap_{n \geq 1} \Sigma_n$ is called the **tail sigma algebra**. Say E is a **tail event** if $E \in \bigcap_{n \geq 1} \Sigma_n$.
 $p \in P(B^{\mathbb{N}})$ satisfies the **0-1 law** if

$$\forall E \in \bigcap_{n \geq 1} \Sigma_n, \quad p(E) \in \{0, 1\}$$

$\bigcap_{n \geq 1} \Sigma_n$ is the sigma algebra of events invariant under finitely many changes.

Defining $\text{Ind}B^{\mathbb{N}}$

Definition

Define $\text{Ind}B^{\mathbb{N}} \subset P(B^{\mathbb{N}})$ as the subspace consisting of those measures $p \in P(B^{\mathbb{N}})$ such that:

1. There exist $\alpha_i \in M_{X_i}, i \geq 1$ such that for all n ,

$$p = \alpha_{1*}\mu \otimes \dots \otimes \alpha_{n*}\mu \otimes (p \triangleright \pi_{n+1,n+2,\dots})$$

2. p satisfies the 0-1 law.

Constructing $\text{prod} : (PB)^{\mathbb{N}} \rightarrow \text{Ind}B^{\mathbb{N}}$

Package a sequence of admissible random elements $(\alpha_n : \Omega \rightarrow B)_{n \geq 1}$ into a single admissible random element for $\text{Ind}B^{\mathbb{N}} \subset P(B^{\mathbb{N}})$:

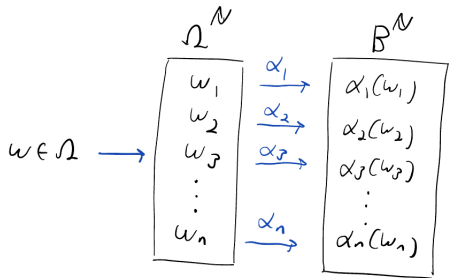
$$\boxed{\Omega \cong \Omega^{\mathbb{N}} \xrightarrow{\alpha_1 \times \alpha_2 \times \dots} B^{\mathbb{N}}}$$

$(\Omega \cong \Omega^{\mathbb{N}} \because \Omega = [0, 1] \text{ is standard-Borel})$

$$\text{ind} : (B^{\Omega})^{\mathbb{N}} \rightarrow (B^{\mathbb{N}})^{\Omega}$$

$$(\alpha_n)_{n \geq 1} \mapsto (\Omega \cong \Omega^{\mathbb{N}} \xrightarrow{\alpha_1 \times \alpha_2 \times \dots} B^{\mathbb{N}})$$

depends only on law $\alpha_{i*}\mu$.



Theorem (Infinite product measures in QBS)

There exists a unique morphism $\text{prod} : (PB)^{\mathbb{N}} \rightarrow \text{Ind}B^{\mathbb{N}}$ such that the diagram below commutes:

$$\begin{array}{ccc}
 (B^{\Omega})^{\mathbb{N}} & \xrightarrow{\text{law}^{\mathbb{N}}} & (PB)^{\mathbb{N}} \\
 \text{law} \circ \text{ind} \downarrow & \swarrow \text{prod} & \downarrow (p_n)_n \mapsto (p_1 \otimes \dots \otimes p_n)_n \\
 \text{Ind}B^{\mathbb{N}} & \xrightarrow{p \mapsto (p \triangleright \pi_1, \dots, \pi_n)_n} & \prod_{n \geq 1} P(B^n)
 \end{array}$$

where

$$\text{law} : B^{\Omega} \rightarrow PB$$

$$\alpha \mapsto \alpha_* \mu$$

Theorem (0-1 law for Ind)

For all $\vec{p} \in (PB)^{\mathbb{N}}$, $\text{prod}(\vec{p})(E) \in \{0, 1\}$ for all tail events E .

I.I.D Sequences in QBS

We also get an analogue of **i.i.d** sequences in QBS as a corollary.

$$\begin{aligned}\text{iid} : PB &\rightarrow \text{Ind}B^{\mathbb{N}} \\ p &\mapsto \text{prod}(p, p, \dots)\end{aligned}$$

$PB \xrightarrow{\text{iid}} \text{Ind}B^{\mathbb{N}} \hookrightarrow P(B^{\mathbb{N}})$ gives semantics to the program

```
iid :: Prob b -> Prob (Stream b)
iid p = do {x <- p; xs <- iid p; return (x:xs)}
```

Theorem (Hewitt-Savage 0-1 law)

For all $p \in PX$, $\text{iid}(p)$ satisfies the 0-1 law.

Future Work

- Full final coalgebra theorem for streams in $\text{Kleisli}(P)$ in QBS?
- Analogue of Ionescu-Tulcea theorem [Kallenberg, 2002] for $\text{Kleisli}(P)$ in QBS?

Summary

- Coalgebraic picture of discrete-time stochastic processes in QBS: $B^{\mathbb{N}}$ as the final $(B \times -)$ -coalgebra in $\text{Kleisli}(P)$.

$$\text{chain}(f) : S \rightarrow P(B^{\mathbb{N}})$$

- Infinite product measures in QBS:

$$\text{Ind}B^{\mathbb{N}} \subset P(B^{\mathbb{N}}), \text{ prod} : (PB)^{\mathbb{N}} \rightarrow \text{Ind}B^{\mathbb{N}}$$

```
prod :: Stream (Prob b) -> Prob (Stream b)
prod (p:ps) = do {x <- p; xs <- prod ps; return (x:xs)}
```

- Applications: random stream of functions $P((\mathbb{R}^{[0,T]})^{\mathbb{N}})$ approximating Brownian motion in QBS.

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